

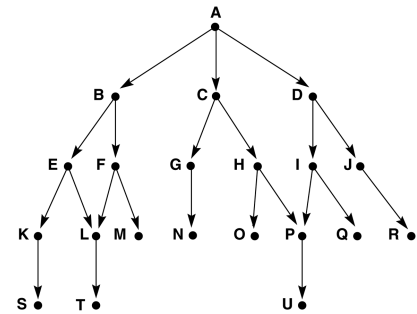
COMP 472 Artificial Intelligence (Winter 2024)

Worksheet #1: Solving Problems by Searching

These are active learning exercises; we'll work on them during the lecture in teams of two!

Breadth-First Search. Let's apply the BFS algorithm discussed in the lecture on an example:

```
begin
  open := [Start];           % initialize
  closed := [];              % states remain
  while open ≠ [] do
    begin
      remove leftmost state from open, call it X; % goal found
      if X is a goal then return SUCCESS
      else begin
        generate children of X;
        put X on closed;
        discard children of X if already on open or closed; % loop check
        put remaining children on right end of open % queue
      end
    end
  end
  return FAIL
end.
```



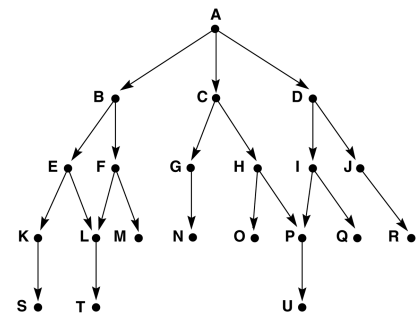
Assume **U** is the **goal state**. The subscript indicates the parent node, e.g., B_A means we reached node B from node A. Note that **open** is a **queue**:

1. $open = [A_{\text{null}}]$, $closed = []$
2. $open = [B_A \ C_A \ D_A]$, $closed = [A]$
3. $open = [C_A \ D_A \ E_B \ F_B]$, $closed = [B \ A]$
4. $open = [\dots]$, $closed = [\dots]$
5. $open = [\dots]$, $closed = [\dots]$
6. $open = [\dots]$, $closed = [\dots]$
7. $open = [\dots]$, $closed = [\dots]$

Depth-First Search. Now we do the same for the DFS algorithm:

Function depth_first_search algorithm

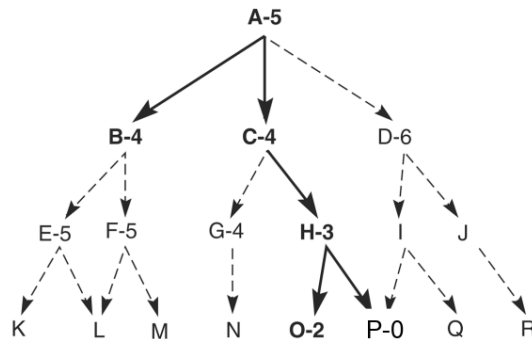
```
begin
  open := [Start];           % initialize
  closed := [];              % states remain
  while open ≠ [] do
    begin
      remove leftmost state from open, call it X; % goal found
      if X is a goal then return SUCCESS
      else begin
        generate children of X;
        put X on closed;
        discard children of X if already on open or closed; % loop check
        put remaining children on left end of open % stack
      end
    end
  end
  return FAIL
end.
```



Again, assume **U** is the **goal state**. Note that **open** is a **stack**:

1. $open = [A_{\text{null}}]$, $closed = []$
2. $open = [B_A \ C_A \ D_A]$, $closed = [A]$
3. $open = [E_B \ F_B \ C_A \ D_A]$, $closed = [B \ A]$
4. $open = [K_E \ L_E \ F_B \ C_A \ D_A]$, $closed = [E \ B \ A]$
5. $open = [S_K \ L_E \ F_B \ C_A \ D_A]$, $closed = [K \ E \ B \ A]$
6. $open = [\dots]$, $closed = [\dots]$
7. $open = [\dots]$, $closed = [\dots]$
8. $open = [\dots]$, $closed = [\dots]$
9. $open = [\dots]$, $closed = [\dots]$
10. $open = [\dots]$, $closed = [\dots]$

Best-First Search. Next, we try a best-first (greedy) search. We have a heuristic $h(n)$ that estimates the cost for each path. The goal is **P**. At each step, expand the node with the *lowest* cost (as predicted by the heuristic), i.e., sort the open list by $h(n)$ (smallest first):



1. open = [A_{null}^5], closed = []
2. open = [B_A^4 C_A^4 D_A^6] (random choice), closed = [A]
3. open = [C_A^4 E_B^5 F_B^5 D_A^6], closed = [B A]
4. open = [.....], closed = [.....]
5. open = [.....], closed = [.....]
6. ???

Finally, *extract the path to the solution* from the search result using the recorded parent nodes:
 Note the difference between *search path* and *solution path*!

Algorithm A. Compute the next step of the Algorithm A on the 8-puzzle:

1	2	3
8		4
7	6	5

Goal

where:

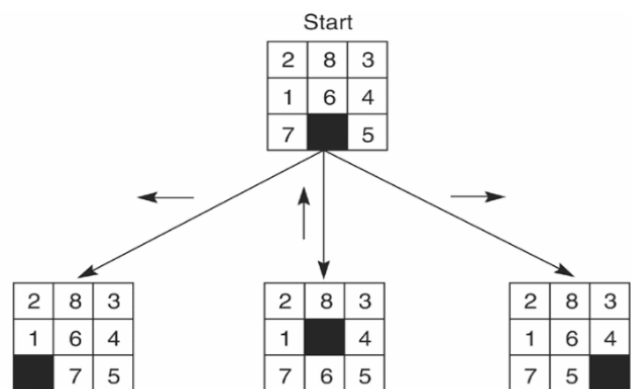
$$f(n) = g(n) + h(n),$$

$g(n)$ = actual distance from n
to the start state, and

$h(n)$ = number of tiles out of place.

$$g(n) = 0$$

$$g(n) = 1$$



Values of $f(n)$ for each state,

6

4

6

1. Pick the state with the *lowest* total cost $f(n)$
2. and compute the next possible search states, including the new values of $f(n)$, $g(n)$ and $h(n)$.