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ESOF 322 - Homework 5

11/5/2019

**ESOF 322 Software Engineering I** 

Homework #5 (32 points total)

Due: Tuesday November 5th

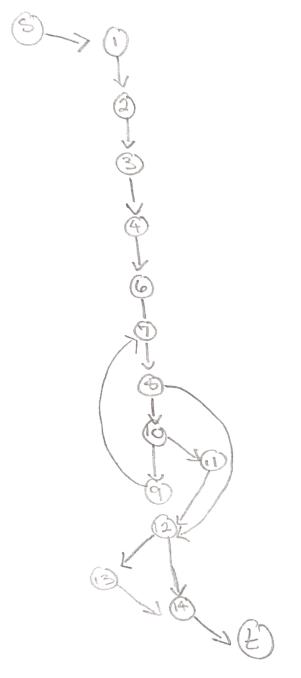
Hand in your printed copy in class, and submit a PDF file into D2L.

#### Question 1 (10 pts)

a) Create a control flowgraph for the *sieve* algorithm. To the left of the line numbers in the source code clearly identify the nodes that will be used in your graph. Once you have identified the nodes, draw the control graph. (4 pts)

```
    /* Find all primes from 2-upper bound using Sieve of Eratosthanes */

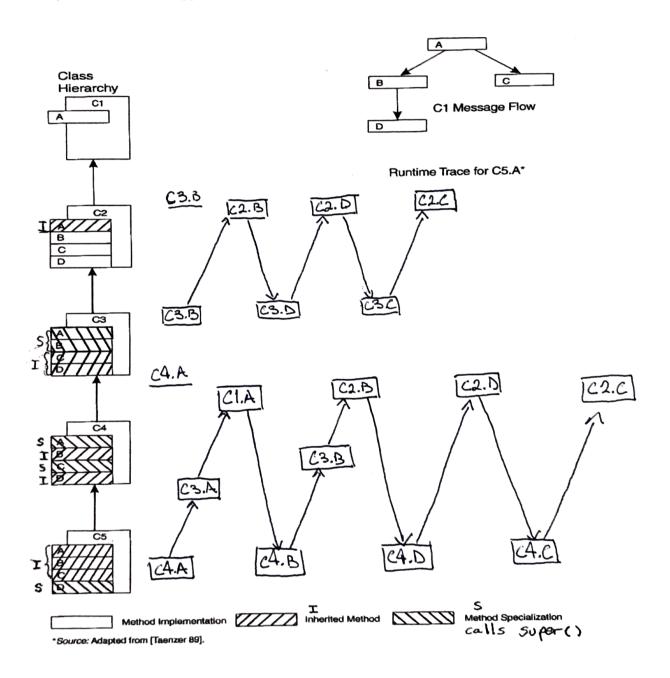
  2.
  3. #include
(\cline{f ()}4. typedef struct IntList {
                int value;
   6.
                struct IntList *next;
  7.
                } *INTLIST, INTCELL;
   8. INTLIST sieve ( int upper bound ) {
   9.
                                           /* list of primes found */
  10.
             INTLIST prime list = NULL;
                                            /* cursor into prime list */
  11.
             INTLIST cursor;
                                            /* a candidate prime number */
  12.
             int candidate;
                                            /* flag: 1=prime, 0=not prime */
  13.
             int is prime;
   14.
              /* try all numbers up to upper bound */
  15.
3 16.
              for (candidate=2;
  17.
(A) 18.
                    candidate <= upper bound;</pre>
19.
                    candidate++) {
   20.
               is_prime = 1; /* assume candidate is prime */
( ) 21.
               for(cursor = prime list;
  22.
   23.
24.
                    cursor;
                    cursor = cursor->next) {
  25.
   26.
                   if (candidate % cursor->value == 0) {
10)27.
   28.
   29.
                     /* candidate divisible by prime */
   30.
                    /* in list, can't be prime */
  31.
                    is prime = 0;
                     break; /* "for cursor" loop */
   32.
   33.
                   }
   34.
               if(is prime) {
  35.
   36.
   37.
                  /* add candidate to front of list */
   38.
                  cursor = (INTLIST) malloc(sizeof(INTCELL));
  39.
                  cursor->value = candidate;
                  cursor->next = prime list;
  40.
                  prime list = cursor;
  41.
  42.
               }
  43.
              }
   44.
              return prime list;
   45.
```



- \*b) Provide a set of test cases that would give 100% Node Coverage (NC). (2 pts)
- \*c) Provide a set of test cases that would give 100% Edge Coverage (EC). (2 pts)
- ∠d) Is 100% NC or 100% EC possible in general? Why, or why not? (2 pts)

# Question 2 (10 pts)

- a) Draw the execution of the calls that exhibit the YoYo problem for a runtime trace of C3.B, and for C4.A (Draw both on the diagram below). (8 pts)
- \*b) Describe what happens when we call C1.D (2 pts)



## **Question 1**

1b)  $T = \{\{1, 2, 3, 4, 6, 7, 8, 10, 9, 7, 8, 12, 13, 14, t\}, \{1, 2, 3, 4, 6, 7, 8, 10, 11, 12, 13, 14, t\}\}$ 

1c) 
$$T = \{\{1, 2, 3, 4, 6, 7, 8, 10, 11, 12, 13, 14, t\}, \{1, 2, 3, 4, 6, 7, 8, 10, 9, 7, 8, 12, 14, t\}\}$$

1d) While 100% Node coverage is possible in general, 100% Edge coverage is generally not. This is because to guarantee 100% Edge coverage you would need to have a solution to the Halting problem. Things like infinite loops make determining if every edge has been covered impossible.

#### **Question 2**

2b) When we call C1.D during the runtime trace of a call to a lower class, we will call the D method on the lower class (i.e. C5.D). However, if we just directly call D from C1 (C1.D) we will get an error as C1 does not have any visibility of a D method.

## **Question 3**

 $3a) T = \{1\}$ 

3b)  $T = \{0\}$ 

 $3c) T = {4}$ 

The test case that would kill all three mutants simultaneously would be T = {{1, 0, 4}}