# **Constraint Satisfaction Problems**

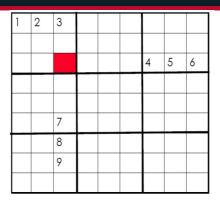
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**CSC384** 

Part 2

## Problems with Plain Backtracking



The backtracking search **won't** detect that the (3,3) cell has no possible value until all variables of the row/column/sub-square are assigned.

## **Constraint Propagation**

In CSPs, there might be variables that have no possible value, but BT doesn't detect
this until it tries to assign them a value.

This leads to the idea of Constraint Propagation (or Domain Filtering).

**Constraint Propagation**: "looking ahead" at the yet unassigned variables in the search, trying to detect obvious failures.

"Obvious" means things we can test/detect efficiently.

 Even if it doesn't detect an obvious failure, it might be possible to eliminate some parts of the future search.

## **Constraint Propagation**

- Propagation has to be applied during the search; potentially at every node of the search tree.
- Propagation itself is an inference step that needs some resources (in particular, time).
   If propagation is slow, this can slow the search down to the point where using propagation makes finding a solution take longer!
- Two main types of propagation: Forward Checking and Generalized Arc Consistency.

## Constraint Propagation: Forward Checking

**Forward Checking:** An extension of backtracking search. Employs a modest amount of propagation (look ahead).

Intuition: When instantiating a variable V, do the following for all constraints C that have only one uninstantiated variable X remaining:

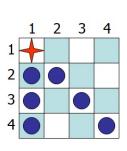
- Check all the values of X;
- Prune those values that violate C.

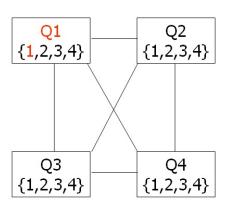
Undo the pruning when backtrack.

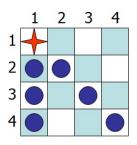
# Forward Checking – Example

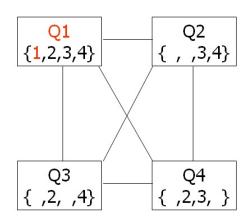
Each of  $Q_1, ..., Q_4$  denotes a queen per row.

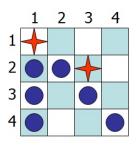
Forward checking prunes domains of  $Q_1,...,Q_4$  based on binary constraints over  $Q_1,...,Q_4$ .

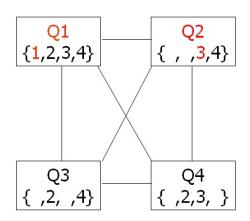


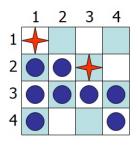


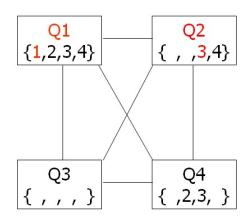


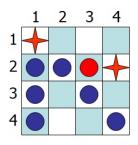


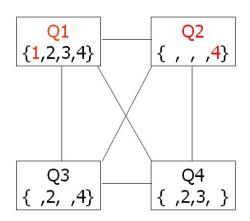


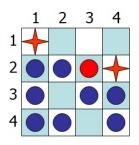


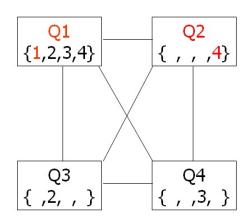


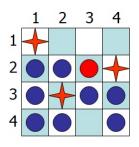


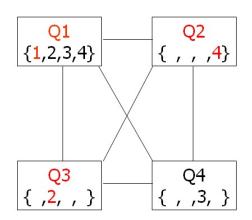


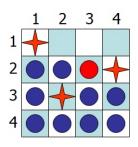


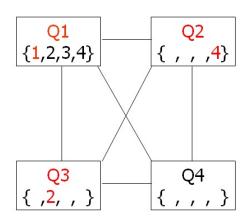


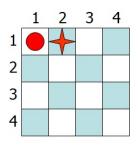


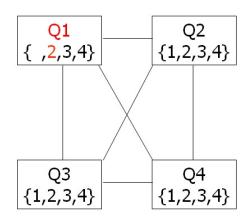


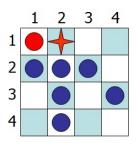


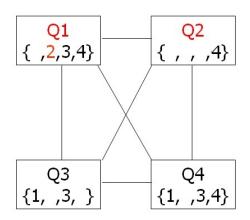


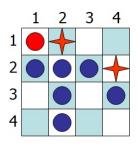


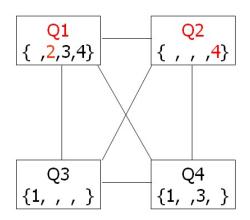


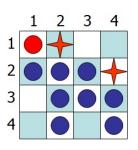


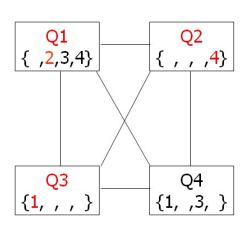


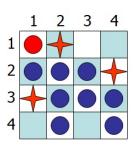


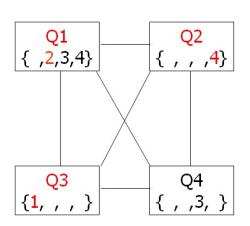


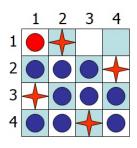


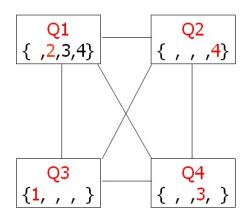


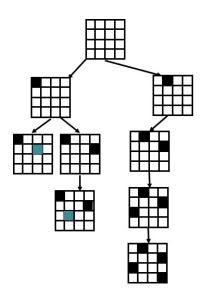












# Backtracking Search: The Algorithm

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```
def FC(Level):
1. if all Variables assigned
2.
      PRINT Value of each Variable
EXIT or RETURN
                                      # EXIT for only one solution
                                      # RETURN for more solutions
  V := PickUnassignedVariable()
   Assigned[V] := TRUE
6. for d := each member of CurDom(V)
7.
      Value[V] := d
  DWOoccured:= False
8.
9.
      for each constraint C over V such that C has only one
      unassigned variable X in its scope:
10.
            if FCCheck(C,X) == DWO: # X domain becomes empty
11.
               DWOoccurred:= True
12.
               BREAK
                                    # stop checking constraints
13.
           if NOT DWOoccured: # all constraints were ok
14.
               FC(Level+1)
15.
            RestoreAllValuesPrunedByFCCheck()
16.
    Assigned[V] := FALSE # UNDO as we have tried all of V's values

    RETURN
```

# Forward Checking: Restoring Values

- After we backtrack from the current assignment the values that were pruned (as a result of that assignment) must be restored.
- Some bookkeeping needs to be done to remember which values were pruned by which assignment.

# Forward Checking Heuristics: Human Analogy

What variables would you try first?

8	1	5	6					4
6				7	5		8	
				9				
9				4	1	7		
	4						2	
		6	2	3				8
				5				
	5		9	1				6
1					7	8	9	5

# Forward Checking: Variable and Value Ordering Heuristics

- · Heuristics can be used to determine
  - the order in which variables are assigned:
     PickUnassignedVariable()
  - the **order** of **values** tried for each variable.
- The choice of the next variable can vary from branch to branch. **Example:** Under the assignment  $V_1=a$  we might choose to assign  $V_4$  next, while under  $V_1=b$  we might choose to assign  $V_5$  next.
- This dynamically chosen variable ordering has a tremendous impact on performance.

## Forward Checking: Heuristics

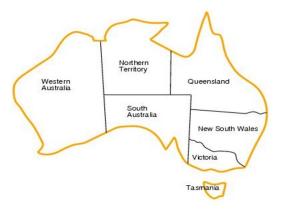
**Degree Heuristic:** Select the variable that is part of the most remaining unsatisfied constraints.

#### Minimum Remaining Values Heuris-cs (MRV):

- Always branch on a variable with the smallest remaining values (smallest CurDom).
   Intuition: If a variable has only one value left, that value is forced, so we should propagate its consequences immediately.
- This heuristic tends to produce skinny trees at the top.
   More variables can be instantiated with fewer nodes searched.
- More constraint propagation/DWO failures occur when the tree starts to branch out.
   Hence, inconsistencies can be found much faster.

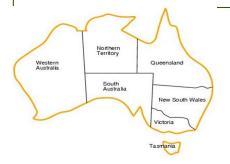
# Example: Map Colouring

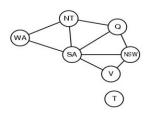
**Problem Statement:** Color the following map using **red**, **green**, and **blue** such that **adjacent regions have different colors**.

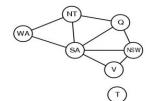


#### **Problem formulation:**

- · Variables:
- · Domains:
- Constraints:



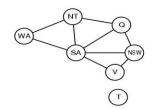


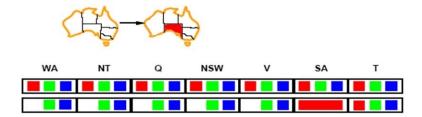




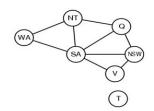


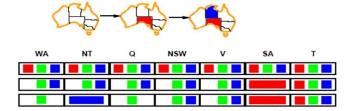
.  $\{SA = red\}$  (using Degree Heuristic)



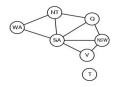


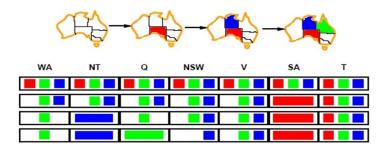
.  $\{SA = red, NT = blue\}$  (using Degree Heuristic)





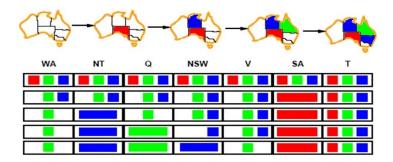
•  $\{SA = red, NT = blue, Q = green\}$  (using Degree Heuristic and MRV)



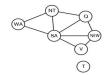


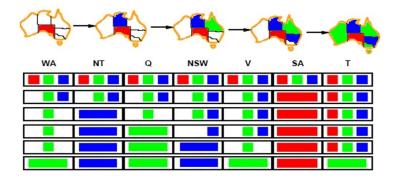
•  $\{SA = red, NT = blue, Q = green, NWS = blue\}$  (using Degree Heuristic)





 $. \; \{ \textit{SA} = \textit{red}, \textit{NT} = \textit{blue}, \textit{Q} = \textit{green}, \textit{NWS} = \textit{blue}, \textit{V} = \textit{green}, \textit{WA} = \textit{green}, \textit{T} = \textit{green} \}$ 





# Example: Map Colouring

Try the map coloring example without MRV and Degree heuristics.

## Forward Checking: Empirically

- · FC often is about 100 times faster than BT.
- FC with MRV (minimal remaining values) often 10000 times faster.
- On some problems the speed up can be much greater.
   Converts problems that are not solvable to problems that are solvable.
- Still FC is not that powerful.
   Other more powerful forms of constraint propagation are used in practice.