

Search Problems

Uniform Cost Search

- Type of uninformed search
- Explores in all directions, so it isn't optimal

Informed Search

- Heuristic := a function that estimates how close a state is to a goal
 - designed for a particular search problem
 - Takes in a state and outputs a guess for how close you are to the goal
- Manhattan distance := distance from a point making only one 90° turn
 - Possible heuristic

Greedy Search

- Uninformed? search alg
- Fringe based algorithm
- Expand the node that seems closest with respect to some heuristic
- Not always optimal
- Somewhat like DFS because it explores a full path at a time

A* Search

- Considers both cost and heuristic value
- Like a combination of greedy and UCS
- Heuristic:

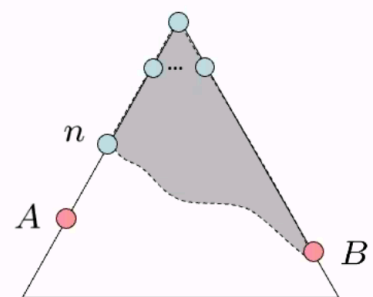
$$f(n) = g(n) + h(n) \quad \begin{cases} g(n) := \text{distance traveled so far} \\ h(n) := \text{heuristic value for a point} \end{cases}$$

- Note: you must evaluate the goal before you can declare success
- Not optimal b/c weightings of $g(n)$ and $h(n)$ can vary
 - Need optimal heuristic for optimality
- Expands mainly toward goal, but does hedge its bets to ensure optimality

Optimality

Proof:

- Imagine B is on the fringe
- Some ancestor n of A is on the fringe, too (maybe A!)
- Claim: n will be expanded before B
 1. $f(n)$ is less or equal to $f(A)$
 2. $f(A)$ is less than $f(B)$
 3. n expands before B
- All ancestors of A expand before B
- A expands before B
- A* search is optimal



Proof: If A is an optimal goal where $\{A, B\} \in \text{goals}$, A must be popped off the fringe before B - Claim:

$B \in \text{fringe} \implies \exists \text{ Ancestor of } A \in \text{fringe}$ - Argument: Ancestor of A will be popped off before B - Known:

$$\begin{aligned} n \in \text{ancestors}_A &\implies f(n) \leq f(A) \\ B \in \text{goals}_{\text{suboptimal}} &\implies f(A) < f(B) \end{aligned}$$

The heuristic is admissible, so it must be an underestimate $\implies h(A) = h(B) = 0 \therefore \{f(A) = g(A), f(B) = g(B)\}$

Given that B is suboptimal, we know the path cost to $A <$ the path cost to B : $g(A) < g(B)$

$\therefore f(n) \leq f(A) < f(B)$ so n must be expanded before $B \implies$ all ancestors of A are expanded before $B \therefore A^*$ is optimal.

Heuristics

- Two types:
 1. Optimistic
 2. Pessimistic
- Admissible (optimistic) iff $0 \leq h(n) \leq h^*(n) |_{h^*(n) := \text{true cost to the nearest goal}}$

Pessimistic (not admissible)

- Regardless of $h(n)$, consider $g(n)$ as more important
- Breaks optimality by trapping good plans on the fringe
- We thought the plan was bad, but it ended up being a good plan

Optimistic (admissible)

- Gives better estimates than the reality
- Ensures optimality because it forces excess computation (explore things that are wrong)

Creating An Admissible Heuristic

- Most of the hard work is for a search problem comes down to the heuristic
- Often we evaluate the cost of an easier problem as our heuristic
 - Manhattan distance, straight-line distance