CS 310: Stack and Queue

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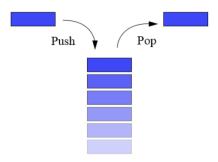
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Review

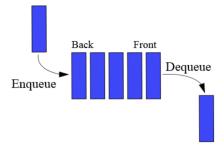
- Iterators
 - Motivation: why do we need iterators?
 - Implementation: how do we support efficient iterations?
 - Nested class / inner (anonymous) class
- Take-home
 - When you use a data structure, use an Iterator to improve efficiency and uniformity
 - When you design or implement a data structure, consider providing an Iterator for the above reason

New Topic

- Stack
 - A data structure that works like a stack (what a twist!)

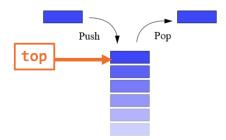


- Queue
 - A data structure that works like people waiting in a line (or queue if you're British)



Stack

- Features
 - LIFO
 - Always operates at the top of the stack
- Basic operations
 - $\mathtt{push}(\mathtt{x}) \colon \mathrm{add} \ \mathtt{x}$ to the top of the stack (grows the stack)
 - pop(): remove the top of the stack (shrinks the stack)
 - top(): return the top of the stack (size is not changed)
 - isEmpty(): true when nothing is in it, false otherwise
- Implementation
 - Based on array / linked list



Stack Example

• You need to be able to draw the stack contents

```
s = new Stack();
s.push(4);
s.push(10);
s.push(5);
s.pop();
s.push(11);
```

Stacks based on Arrays

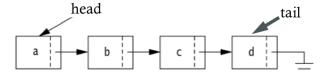
```
class AStack<T>{
    private ArrayList<T> stuff;
    public AStack(); // Constructor
    public void push(T x); // like add(x) or append(x)
    public void pop(); // like remove(size()-1)
    public T top(); // like get(size()-1)
    public boolean isEmpty(); // like size()==0
}
```

- \bullet Use an ${\tt ArrayList}$ as the underlying storage
- The top of the stack is the end of the array
 - Operations are performed only at the end which makes it faster with an array based implementation
- What's the Big-O?
 - -O(1)

Stack Based on Linked List

```
Class LStack<T>{
    private LinkedList<T> stuff;
    public LStack(); // Assume head as stack top
    public void push(T x); // like insert(0,x)
    public void pop(); // like remove(0)
    public T top(); // like get(0)
    public boolean isEmpty(); // like size()==0
}
```

- Use a Linked List as the underlying storage
 - Operate only at one end
- Big-O?
 - -O(1)



Stack Applications

• Check the symbolic balancing of an equation

$$- \{(<>[\{<>\}])\{\}\} \text{ vs. } \{(<[\{<>>\}])\{\}\}$$

• Postfix calculation

$$-6523 + 8 \times + 3 + \times =$$

• Infix to Postfix conversion

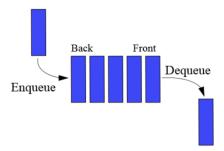
$$-a + b \times c + (d \times e + f) \times g \rightarrow abc \times + de \times f + g \times + de \times f + g \times f$$

- Call stack
 - fib(4)
- Tree traversal preorder traversal
- Graph search depth first search
- And a bunch of over applications

Queue

- Features
 - FIFO
 - Only remove from front
 - Only add to back
- Basic operations
 - enqueue(x) or add(x): x enters at the back

- dequeue() or poll(): front leaves
- getFront() or peek(): returns the item at the front
- isEmpty(): true when nothing is in it, false otherwise

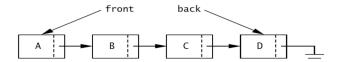


Queue Example

- You need to be able to draw the queue contents
- What is the value of v?

```
q = new Queue();
q.enqueue(4);
q.enqueue(10);
q.enqueue(5);
q.dequeue();
v = getFront();
q.dequeue(11);
q.enqueue(25);
```

Queue Based on Linked Lists

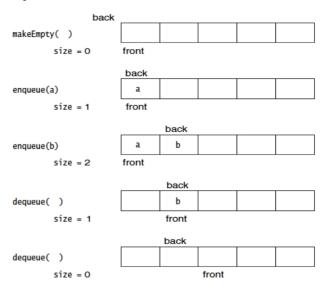


- Append to one end, and remove from the other end
 - For example, head \rightarrow front, tail \rightarrow back
 - enqueue(x): insert at the tail
 - dequeue(): remove from head
 - getFront(): return head contents
 - isEmpty(): size() == 0

Queue Based on Arrays

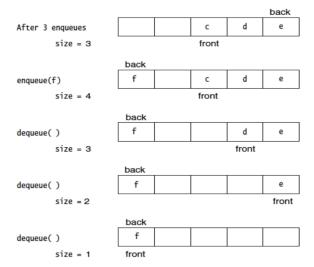
- Naive implementation:
 - enqueue(x): insert at the end
 - dequeue(): remove from start and shifting internally
 - * In fact, a lot of shifting! Shifting is done for every single dequeue()!
 - Alternatively, we could just mark the front and the back in the array and update them with enqueue and dequeue

Queue Based on Arrays



- Between the front and the back, we have a valid queue
 - There's no shifting: O(1) for dequeue()!
 - But it does use a sizeable amount of space

Queue: Array with Wraparound



• Exercise: what needs to be changed to implement the wraparound functionality?

Big-O Comparison

• Stack

Implementation	push()	pop()	top()	<pre>isEmpty()</pre>	size
Array	1*	1	1	1	1
Linked List	1	1	1	1	1

^{*}Amortized analysis

Implementation	enqueue()	dequeue()	<pre>getFront()</pre>	<pre>isEmpty()</pre>	size
Array	1*	1	1	1	1
Linked List	1	1	1	1	1

*Amortized analysis

Why use a Stack or Queue

- Restricted operations give us good worst cases
 - O(1) for all supported operations
 - -O(n) for space
- Simple data structures
 - Focus on limited operations
 - Can be made out of primitive data structures (arrays and linked lists)
- Good for representing time-related data
 - Call stack
 - Packet queues

Review: Queues

- FIFO
- Supported operations:
 - enqueue(x): insert at the tail
 - dequeue(): remove from head
 - getFront(): return head contents
 - size(): returns the size of the queue
 - isEmpty()
- Applications:
 - Simulate a process with a FIFO order
 - Scheduling queue of a CPU or disk or printer
 - Serve as a buffer for file I/O, network communications, etc.

Priority Queues

- Much of the time tasks that we use a queue for have different priorities
 - It is convention that the lower the priority, the better
 - Symmetric code if higher is better
 - Dequeue the ones with the "best" priority first
- Common priority queue operations
 - void insert(T x, int p): insert x with priority p
 - T findMin(): return the object with the "best" priority
 - deleteMin(): remove the object with the "best" priority

Priority Queue Design

Data Structure	<pre>insert()</pre>	<pre>findMin()</pre>	<pre>deleteMin()</pre>	Notes()
Sorted Array Sorted Linked List	O(n) $O(n)$	O(1) $O(1)$	$O(1) \\ O(1)$	min at high index min at head or tail

- Other data structures exist as good candidates of priority queues
 - Binary search trees
 - Heaps
 - We'll cover these later

Summary

- Stacks and queues
 - Try implementing them
 - Project 2
- Next lecture: Trees, recursion
 - Reading: Chapter 18, Chapter 7