

Definitions and Theorems

Connor Baker, March 2017

Definition 1 (Statement). Any sentence which can be evaluated as either true or false.

Definition 2 (Compound Statement). A statement made up of one or more component statements connected by logical connectors.

Definition 3 (Equivalence of Logical Operators). Two sets of logical operators are said to be equivalent if they produce the same output.

Definition 4 (Tautology). A statement that's always true.

Definition 5 (Contradiction). A statement that's always false.

Definition 6 (A Set). Any collection of objects.

Definition 7 (Set Builder Notation). {expression: rule}

Definition 8 (Universal Set). The given or implied set that contains all other sets in the problem. This set fixes Russel's Paradox.

Definition 9 (Tautology). A statement that's always true.

Definition 10 (Natural Numbers). The set $\mathbb{N} : \{1, 2, 3, \dots\}$.

Definition 11 (Integers). The set $\mathbb{Z} : \{\dots, -1, 0, 1, 2, \dots\}$.

Definition 12 (Rational Numbers). The set $\mathbb{Q} : \{\frac{a}{b} : a \in \mathbb{Z} \text{ and } b \in \mathbb{N}\}$.

Definition 13 (Real Numbers). The set $\mathbb{R} : \{a_n a_{n-1} \dots a_1 a_0 a_{-1} a_{-2} \dots : n \in \mathbb{N} \cup \{0\} \text{ and } a_i \in \{0, \dots, 9\}\}$.

Definition 14 (Complex Numbers). The set $\mathbb{C} : \{a + bi : i^2 = -1 \text{ and } a, b \in \mathbb{R}\}$.

Definition 15 (Subset). Given two sets A and B , $A \subseteq B \iff \forall a \in A \implies a \in B$.

Definition 16 (Open Sentence (AKA Predicate)). A statement that contains a variable. The truth value depends on the variable.

Definition 17 (Truth Set). The set of values that make the statement true.

Definition 18 (Quantifiers and Negations). Logical Quantifiers and Negators:

1. Universal Quantifier: \forall – Must be true for all x in the universal set such that $P(x)$ is true: $(\forall x)(P(x))$.
2. Existential Quantifier: \exists – True if for at least one x in the universal set such that $P(x)$ is true: $(\exists x)(P(x))$.
3. Unique Quantifier: $\exists!$ – True if there exists only one x in the universal set such that $P(x)$ is true: $(\exists! x)(P(x))$.
4. Negation of the Universal Quantifier: $\sim (\forall x)(P(x))$ is $(\exists x)(\sim P(x))$.
5. Negation of the Existential Quantifier: $\sim (\exists x)(P(x))$ is $(\forall x)(\sim P(x))$.

Definition 19 (Direct Proof). $P \implies Q$.

Definition 20 (Contrapositive Proof). $(\sim Q) \implies (\sim P)$.

Definition 21 (Proof by Contradiction). We start with $P \implies Q$. Assume that $\sim P \wedge Q$ is true. Then $\sim P \implies A_1 \implies A_2 \implies \dots \implies R$. And, if $Q \implies B_1 \implies B_2 \implies \dots \implies \sim R$. Then, $\sim R \wedge R$ must be true, which is a contradiction, so the original assumption is false, and $P \implies Q$.

Definition 22 (Axioms of the Natural Numbers). The following are axioms for the set of the Natural Numbers, \mathbb{N} :

1. Successor property
 - (a) One is a natural number
 - (b) One is not the successor of any number
 - (c) Every natural number has a unique successor
2. Closure under addition and multiplication
3. Associativity
4. Commutativity
5. Distribution of multiplication over addition
6. Cancellation
 - (a) Real numbers have this property unless the number being cancelled is a zero
 - (b) Matrix multiplication does not have this property

Definition 23 (Divisible). Let $a, b \in \mathbb{N}$. Then $a|b$ if $\exists k \in \mathbb{N} : ak = b$.

Definition 24 (Prime). A number p , where $p \in \mathbb{N}$, is prime if $p > 1$ and its only divisors are one and itself.

Definition 25 (Factor). A number q , where $q \in \mathbb{N}$, is a factor of r if $q|r$.

Definition 26 (Prime Factor Decomposition). Let p_1, p_2, \dots, p_k be all primes less than q . Then, the prime factor decomposition of q is $p_1^{n_1} p_2^{n_2}, \dots, p_k^{n_k}$ where $n_i \in (\mathbb{N} \cup \{0\})$.

Theorem 27 (Fundamental Theorem of Arithmetic). All natural numbers have a unique prime factorization up to commutativity.

Definition 28 (Union over \mathcal{A}). Let \mathcal{A} be a family of sets. The union over \mathcal{A} is defined as:

$$\bigcup_{A \in \mathcal{A}} = \{x : (\exists A \in \mathcal{A})(x \in A)\}$$

which is equivalent to:

$$\bigcup_{A \in \mathcal{A}} = \{x : (\exists A)((A \in \mathcal{A}) \wedge (x \in A))\}$$

Definition 29 (Intersection over \mathcal{A}). Let \mathcal{A} be a family of sets. The intersection over \mathcal{A} is defined as:

$$\bigcap_{A \in \mathcal{A}} = \{x : (\forall A \in \mathcal{A})(x \in A)\}$$

which is equivalent to:

$$\bigcap_{A \in \mathcal{A}} = \{x : (\forall A)((A \in \mathcal{A}) \implies (x \in A))\}$$

Theorem 30 (Relative Cardinality of Intersection and Union). For every set $B \in \mathcal{A}$:

$$B \subseteq \bigcup_{A \in \mathcal{A}} A,$$

$$\bigcap_{A \in \mathcal{A}} A \subseteq B.$$

The intersection is no bigger than the smallest set, and the union is no smaller than the biggest set.

Assume that $\mathcal{A} \neq \emptyset$. Then:

$$\bigcap_{A \in \mathcal{A}} A \subseteq \bigcup_{A \in \mathcal{A}} A.$$

If $\mathcal{A} \neq \emptyset$, the union isn't a problem but the intersection would be the set of all sets, and as such is undefined.

Definition 31 (Indexed Family of Sets). Let Δ be a nonempty set. Then, $\forall \alpha \in \Delta$, there is a corresponding set A_α . The family of sets $\mathcal{A} = \{A_\alpha : \alpha \in \Delta\}$.

Definition 32 (Union and Intersection over an Indexed Family of Sets \mathcal{A}). Let \mathcal{A} be a family of sets with indices $\alpha \in \Delta$. Then, the union over A_α is defined as:

$$\bigcup_{\alpha \in \Delta} A_\alpha = \{x : (\exists \alpha \in \Delta)(x \in A_\alpha)\}$$

and the intersection is defined as:

$$\bigcap_{\alpha \in \Delta} A_\alpha = \{x : (\forall \alpha \in \Delta)(x \in A_\alpha)\}$$

Theorem 33 (Relative Cardinality of Intersection and Union over Indexed Family of Sets). For every set $\beta \in \Delta$:

$$A_\beta \subseteq \bigcup_{\alpha \in \Delta} A_\alpha,$$

$$\bigcap_{\alpha \in \Delta} A_\alpha \subseteq A_\beta.$$

$$\overline{\bigcup_{\alpha \in \Delta} A_\alpha} = \bigcap_{\alpha \in \Delta} \overline{A_\alpha}$$

$$\overline{\bigcap_{\alpha \in \Delta} A_\alpha} = \bigcup_{\alpha \in \Delta} \overline{A_\alpha}$$

Definition 34 (Pairwise Disjoint). Let $\mathcal{A} = \{A_\alpha : \alpha \in \Delta\}$. Then \mathcal{A} is pairwise disjoint if $\forall \alpha, \beta \in \Delta$ with $A_\alpha \neq A_\beta$, $A_\alpha \cap A_\beta = \emptyset$.

Theorem 35 (Order Properties of the Natural Numbers). Let $x, y, z \in \mathbb{N}$. Then, $\forall x, y, z$:

1. $x < y \iff \exists w \in \mathbb{N} : x + w = y$
2. $x \leq y \iff x = y \text{ or } x < y$
3. if $x < y$ and $y < z$, then $x < z$ (transitivity)
4. if $x \leq y$ and $y \leq x$, then $x = y$
5. if $x < y$, then $x + z < y + z$ and $xz < yz$

Theorem 36 (Principle of Mathematical Induction (PMI)). If S is any subset of the natural numbers, with the properties that:

1. $1 \in S$
2. if $k \in S$, then $(k + 1) \in S$

then $S = \mathbb{N}$.

The general process of mathematical induction is as follows:

1. Define $S = \{n \in \mathbb{N} : \text{some statement is true}\}$
 - (a) Prove that the basis case holds: that means that $1 \in S$
 - (b) Assume $k \in S$. Then, based on this assumption, prove it to be the case that $(k + 1) \in S$.
 - (c) Conclude that by the Principle of Mathematical Induction, $S = \mathbb{N}$.

Definition 37 (Inductive Set). A set $S \subseteq \mathbb{N}$ is inductive if whenever $n \in S$, then $(n + 1) \in S$.

Definition 38 (Factorial). If $n \in \mathbb{N}$, then $n! = n(n - 1)!$.

Definition 39 (Zero Factorial). $0! = 1$.

Definition 40 (General Principle of Mathematical Induction). $S \subseteq \mathbb{N}$ where $k \in S$ and if $j \in S$, then $(j+1) \in S$, and it is true for all $\{k, k+1, \dots\}$, then S is inductive.

Theorem 41 (Principle of Strong Mathematical Induction (PSMI)). If $S \subseteq \mathbb{N}$ with the property that $\forall m \in \mathbb{N}$, if $\{1, 2, \dots, m-1\} \subseteq S$, then $m \in S$, then $S = \mathbb{N}$.

PSMI is different from PMI because with PMI we assume that we can start at a value and carrying forward from that value something holds. With PSMI, we assume that it holds over an interval.

Theorem 42 (Well Ordering Principle (WOP)). Every nonempty subset of \mathbb{N} has a least element.

Theorem 43 (The Division Algorithm). Let $a, b \in \mathbb{N}$, with $b \leq a$. Then we will prove that $\exists q \in \mathbb{N}$ and $r \in \mathbb{N} \cup \{0\} : a = bq + r$ where $0 \leq r < b$.

Consider all multiples of $b > a$. Let $S = \{s \in \mathbb{N} : sb > a\}$. By (WOP), S has a least element $q+1$, so $q \notin S$. Therefore, $qb \leq a$.

Let $r = a - qb$. Since $qb \leq a$, $a - qb \geq 0$, so it must be the case that $r \geq 0$.

If $r \geq b$, then $r = q - qb \geq b \implies a - qb - b \geq 0 \implies q - b(q+1) \geq 0$. So, $a \geq b(q+1)$. But, for $(q+1) \in S$, it must be that $b(q+1) > a$. Then, $(q+1) \in S$. This is a contradiction. Therefore, $r < b$.

Furthermore, q and r are unique.

Assume $\exists q_1, r_1$ with $a = bq_1 + r_1$ where $0 \leq r_1 < b$. Then $a = bq + r$, $a = bq_1 + r_1$. This implies that $0 = b(q - q_1) + (r - r_1)$, $b \neq 0$. If it is the case that $q - q_1 \neq 0$, then $|q - q_1| \in \mathbb{N}$. Then $r_1 > r$ and $|r_1 - r| = mb$ for some $m \in \mathbb{N}$, $m = |q - q_1|$. Thus, $r_1 \geq mb \implies r_1 \geq b$, which is a contradiction such that $q - q_1$ would be zero and $r_1 - r = 0$.