

*Definition of Recursion ... see Recursion*

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# Recursion

Self help

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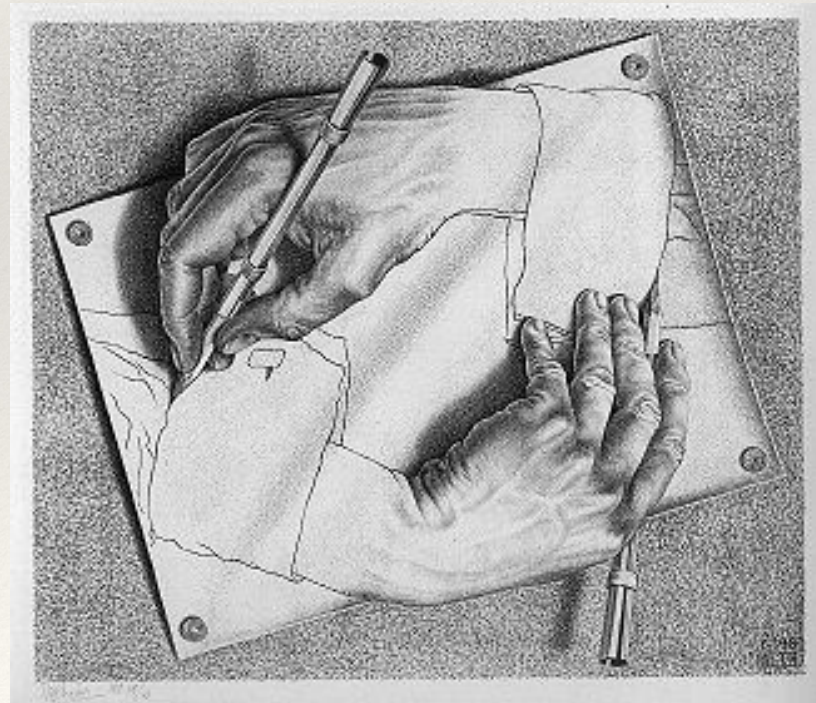
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# Simply Recursive ...

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- ❖ These notes on recursion are the basic concepts you will need to answer A4 question 2
- ❖ More detailed notes on recursion including the concept of “tail recursion as well as animated slides will be coming out next week ... stay tuned

**Drawing Hands**  
by M. C. Escher  
lithograph  
January 1948





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# Recursion

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- ❖ *Definition:* When a subroutine **invokes** *itself*.
- ❖ Or when a series of subroutines eventually invoke the first subroutine again.
- ❖ The intent is to break a large problem into smaller and simpler problems. These smaller solutions are then combined to solve the larger problem.




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# Example of Recursion

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- ❖ Factorials can be calculated recursively.

$$n! = n \times (n - 1) \times (n - 2) \times 1$$

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \times (n - 1)! & \text{if } n > 0 \end{cases}$$


- ❖ This part suggests recursion. The calculation is done with reference to itself.

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# Factorial Example

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$$4! = 4 \times 3!$$

$$= 4 \times (3 \times 2!)$$

$$= 4 \times (3 \times (2 \times 1!))$$

$$= 4 \times (3 \times (2 \times (1 \times 0!)))$$

$$= 4 \times (3 \times (2 \times (1 \times 1)))$$

❖ Three levels of recursion



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# Recursion

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- ❖ Every recursive process requires two things:
  1. A **base case** that is processed *without* recursion. This requires an **ending condition** that knows when to apply the base case.
  2. A method that reduces a particular case to one or more smaller cases. This requires a **recursive call**.



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# Recursion

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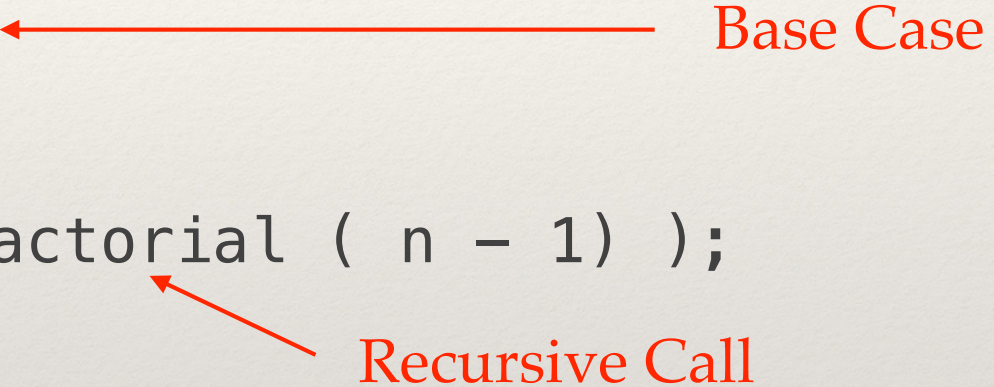
- ❖ In the factorial example the base case is
  - ❖  $n! = 1$  if  $n = 0$  *no further recursion is needed*
- ❖ and the method to reduce a case to a smaller one is
  - ❖  $n! = n \times (n - 1)$  if  $n > 0$
- ❖ Now, let's translate this to C code...

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# factorial Subroutine

---

```
int factorial ( int n ) {  
    if ( n == 0 ) {  
        return ( 1 );  
    else {  
        return ( n * factorial ( n - 1 ) );  
    }  
}
```




The diagram illustrates the execution of the factorial subroutine. A red arrow points from the text "Base Case" to the condition `if ( n == 0 )`. Another red arrow points from the text "Recursive Call" to the recursive call `factorial ( n - 1 )`.

The call to `factorial` from `factorial` is *recursive*.

When executed this will continue until the `if` statement is true (*i.e.* `n == 0` ).



```
int factorial ( 4 ) {  
    if ...  
    ...  
    else  
        return (4 * factorial(4-1));  
        if ...  
        else  
            return (3 * factorial(3-1));  
            if ...  
            else  
                return (2 * factorial(2-1));  
                if ...  
                else  
                    return (1 * factorial(1-1));  
                    if ( n == 0 )  
                        return (1);
```



4 \* 3 \* 2 \* 1 \* 1

3 \* 2 \* 1 \* 1

2 \* 1 \* 1

1 \* 1



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# Why Use Recursion?

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- ❖ Recursive solutions can be **concise** and *elegant*.
  - ❖ Small amount of code required.
- ❖ **But**, it requires that the programmer understand the problem and the recursive solution *very well*.
- ❖ What problems are candidates for recursion?
  - ❖ Problems that are readily subdivided can be solved recursively in a small amount of code.
  - ❖ Problems that have a long chain of partial results can benefit from recursion.



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# Other Recursive Algorithm Examples

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- ❖ **Quicksort** is a divide and conquer algorithm.
- ❖ A large array is divided into two smaller sub-arrays called low and high.
- ❖ The sub-arrays are then recursively sorted.



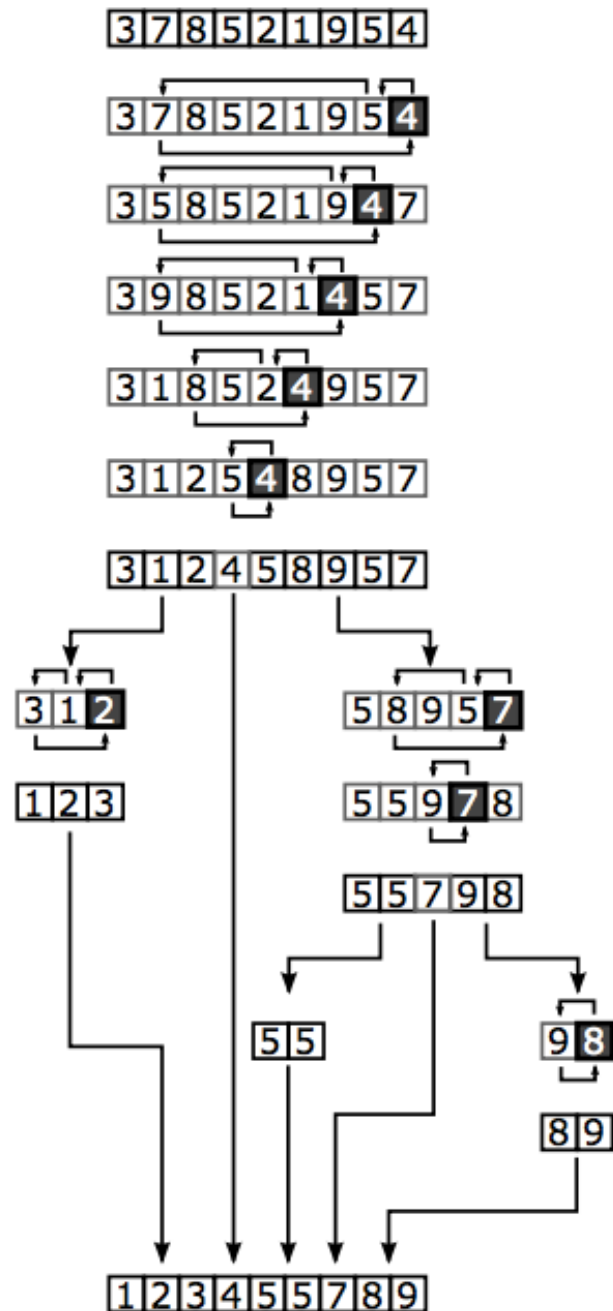
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# Quicksort Algorithm

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- ❖ Pick an element, called a **pivot**, from the array.
- ❖ Partition the array by reordering it so that all elements with values **less** than the **pivot** come before the pivot and all elements **greater** than the **pivot** come after it..
- ❖ Recursively apply the above steps to the two sub-array created by the pivot: the sub-array of **lesser** values and the sub-array of elements of **greater** values.





Select the last element to be the **pivot**.

Compare the **pivot** to other elements and put **greater** element after it.

After the **partition** is finished - *recursively partition* each sub-array on either side of the position of the **pivot**.



```
#include <stdio.h>
#include <stdlib.h>

void quickSort ( int *arr, int low, int high ) {
    int pivot, i, temp;

    /*
     * Select a pivot element
     * - the last element
     */
    pivot = high;

    if ( low < high ) {
        i = low;
        while ( i < pivot ) {
            /*
             * Go from the lower boundary until you
             * get a number greater than the pivot index
             */
            while ( arr[i] <= arr[pivot] && i < pivot ) {
                i++;
            }
        }
    }
}
```



```
/*  
 * If you find an element that is higher than the pivot  
 * swap with the element in front of the pivot  
 */
```

```
temp = arr[i];  
arr[i] = arr[pivot-1];  
arr[pivot-1] = temp;
```

```
/*  
 * Swap the pivot with the element in front of it  
 */
```

```
temp = arr[pivot];  
arr[pivot] = arr[pivot-1];  
arr[pivot-1] = temp;  
pivot = pivot - 1;
```

```
}
```

```
/*  
 * Recursion: perform quickSort for the two sub-arrays,  
 * one to the left of pivot and one to the right of the pivot  
 */
```

```
quickSort(arr, low, pivot-1);  
quickSort(arr, pivot+1, high);
```

```
}
```

```
}
```



Generating the numbers to be sorted:

1 45 89 53 **33**

quicksort ( 0, 4 )

1 53 89 **33** 45

1 89 **33** 53 45

1 **33** 89 53 45

quicksort ( 0, 1 )

1 33

quicksort ( 0, 0 )

quicksort ( 1, 1 )

quicksort ( 2, 4 )

53 **45** 89

**45** 53 89

quicksort ( 2, 2 )

quicksort ( 3, 4 )

53 89

quicksort ( 3, 3 )

quicksort ( 4, 4 )

Sorted array:

1 33 45 53 89

( 1 45 89 53 33 )



Generating the numbers to be sorted:

4 87 32 21 5 25 11 59 4 **18**

quicksort ( 0, 9 )

4 4 32 21 5 25 11 59 **18** 87

4 4 59 21 5 25 11 **18** 32 87

4 4 11 21 5 25 **18** 59 32 87

4 4 11 25 5 **18** 21 59 32 87

4 4 11 5 **18** 25 21 59 32 87

4 4 11 5 **18** 25 21 59 32 87

quicksort ( 0, 3 )

4 4 **5** 11

quicksort ( 0, 2 )

4 4 **5**

quicksort ( 0, 1 )

4 **4**

quicksort ( 4, 9 )

18 25 21 59 32 **87**

quicksort ( 4, 8 )

18 25 21 **32** 59

quicksort ( 4, 7 )

18 25 **21** 32

quicksort ( 4, 6 )

18 **21** 25

quicksort ( 4, 5 )

18 **21**

Sorted array:

4 4 5 11 18 21 25 32 59 87  
( 4 87 32 21 5 25 11 59 4 18 )



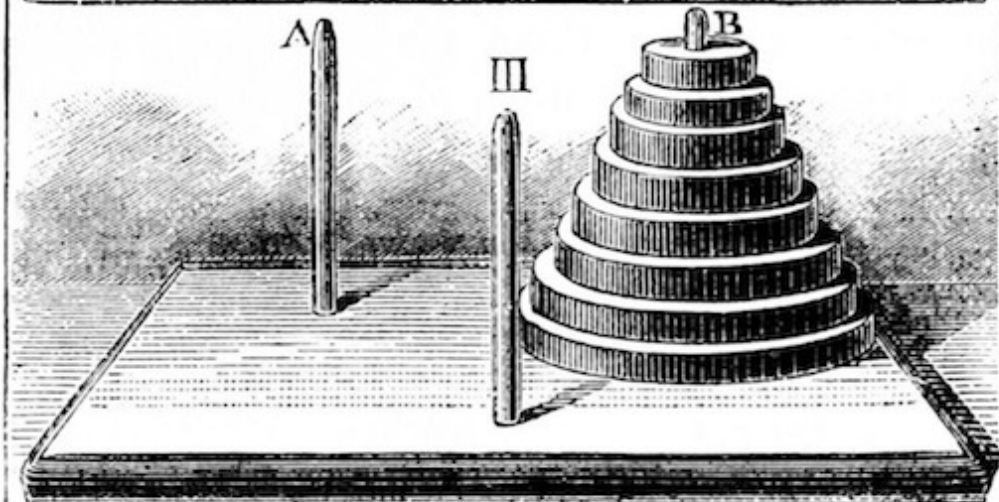
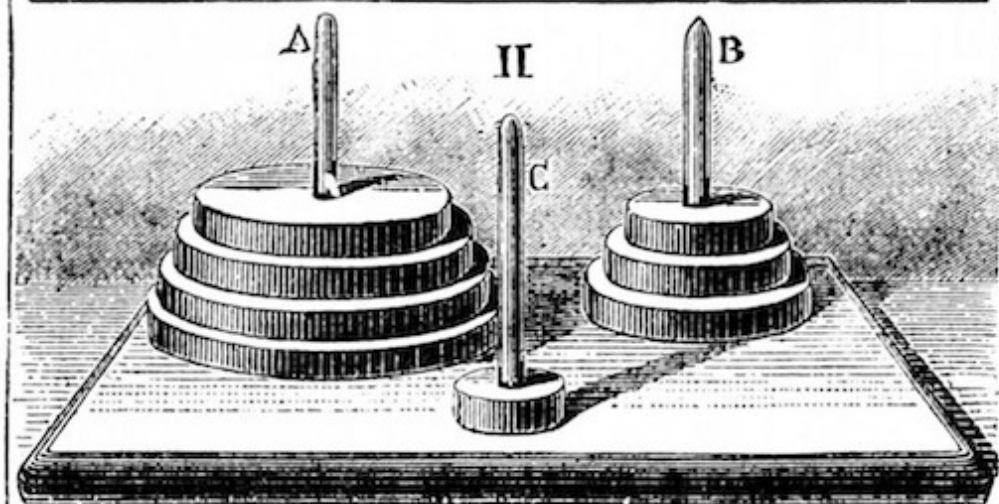
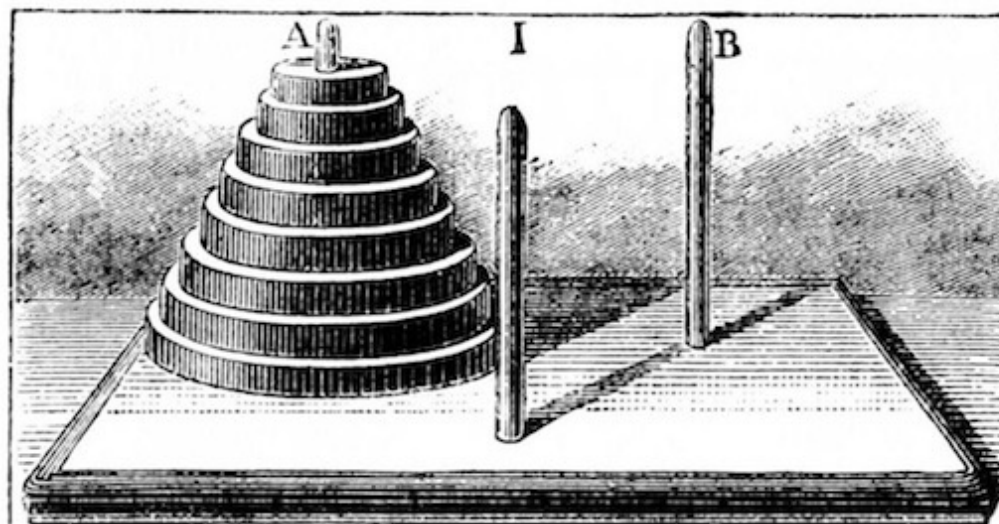
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# Towers of Hanoi

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- ❖ Towers of Hanoi is a puzzle that consists of three posts, and a set of disks of different sizes that can be stacked on the posts.
- ❖ At the start, all the disks are stacked on one post by size (largest on bottom).
- ❖ The challenge is to transfer the stack from the first post to the third, using the second post for temporary storage.
- ❖ Only one disk can be moved at a time, and a larger disk can never be put on top of a smaller disk.







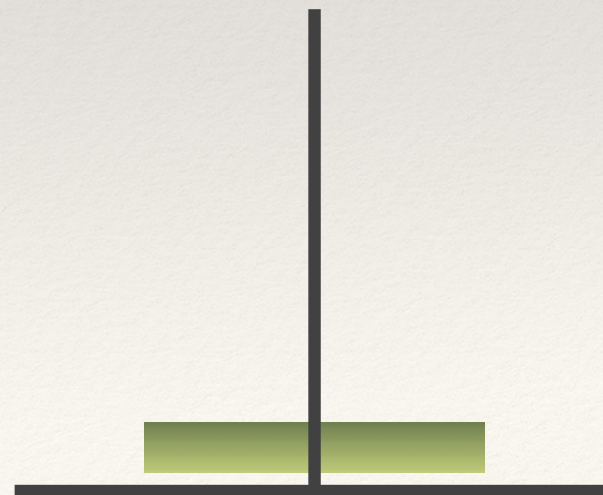
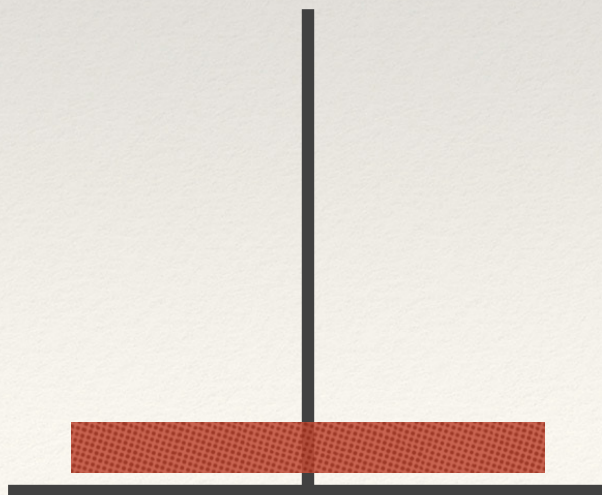
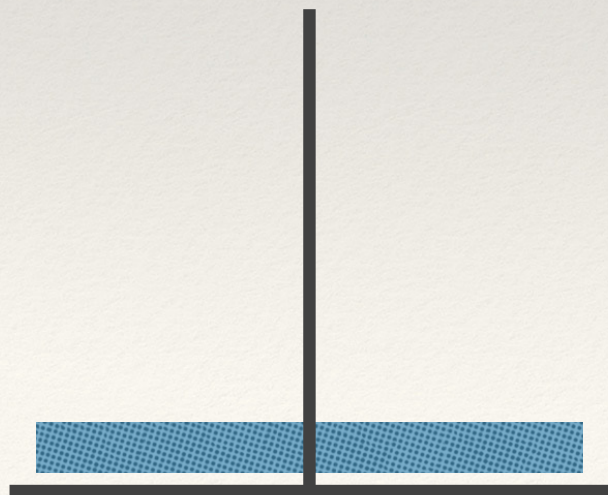
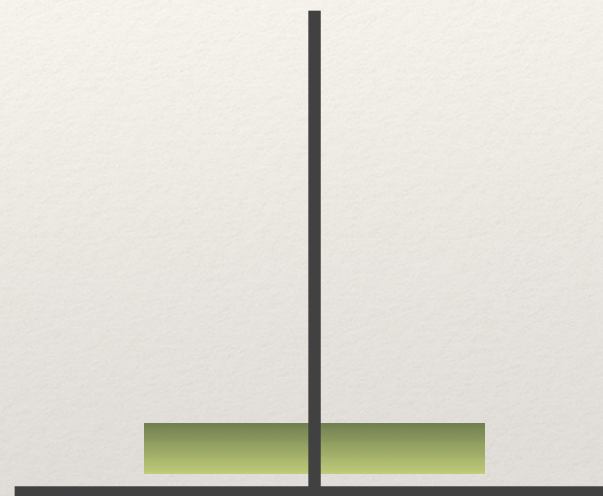
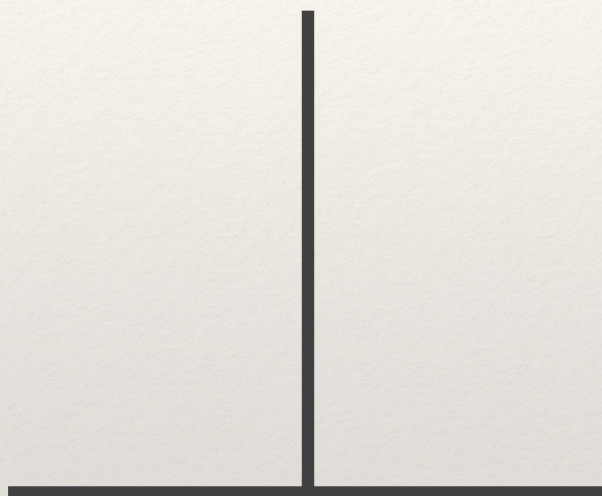
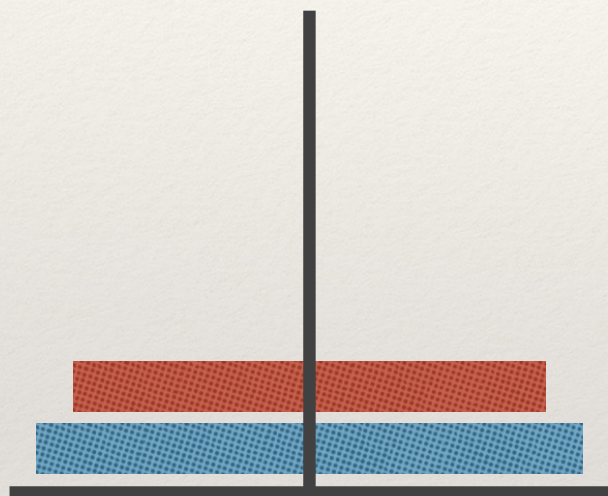
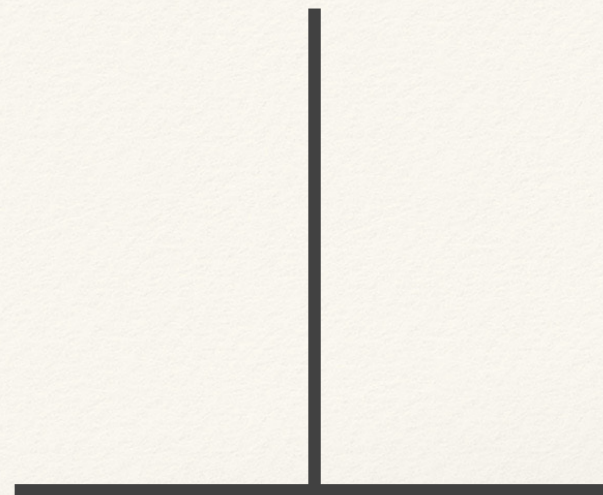
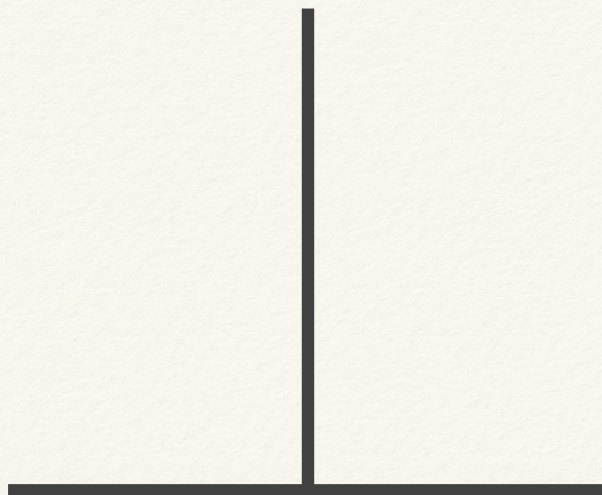
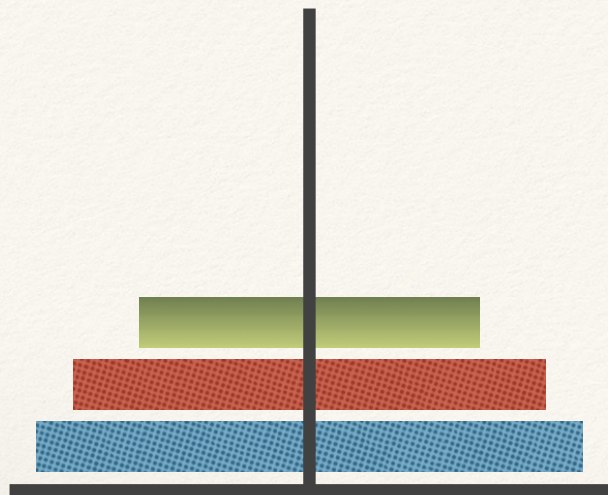
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# The Recursive Algorithm

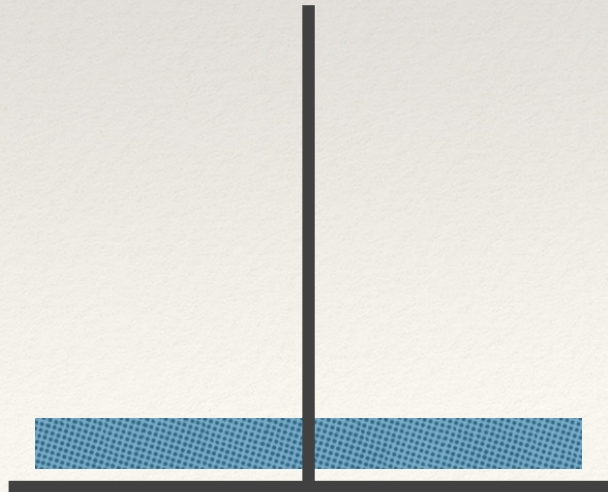
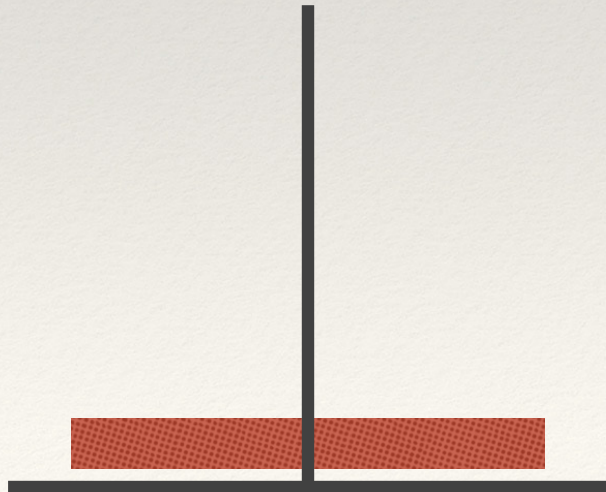
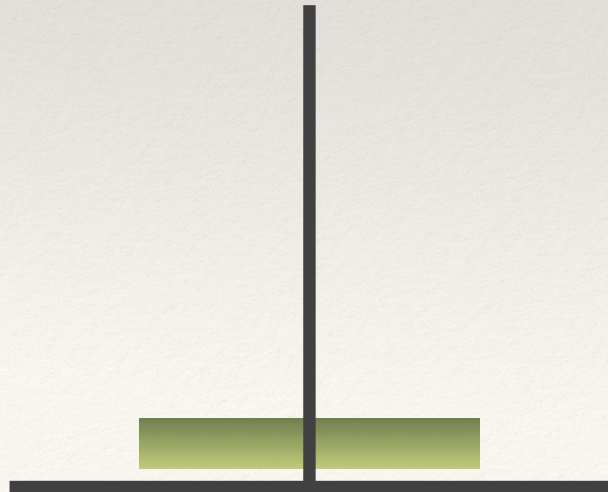
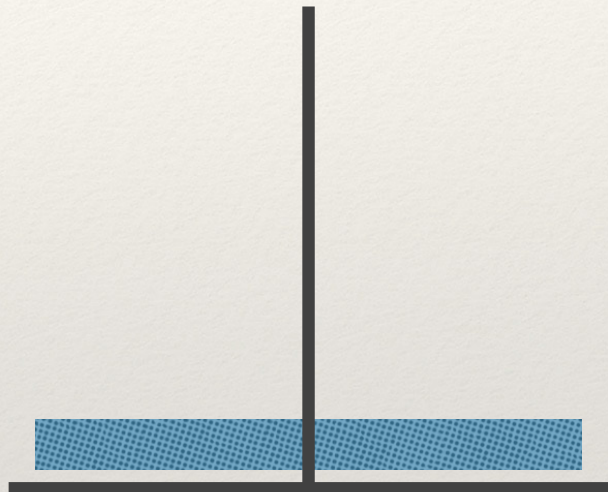
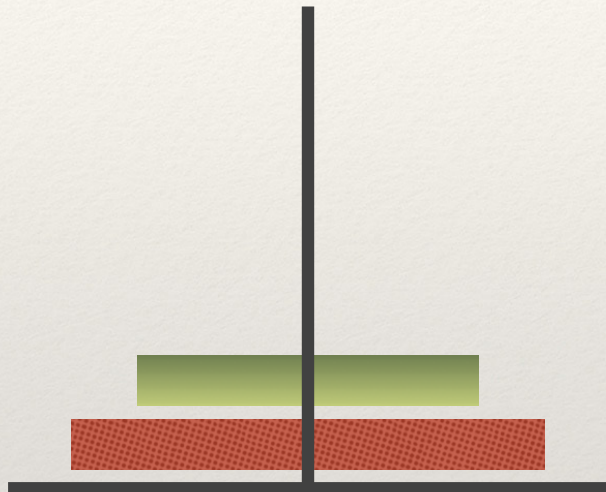
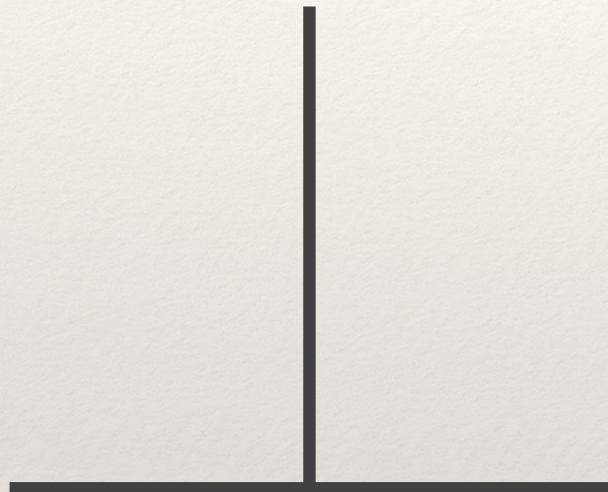
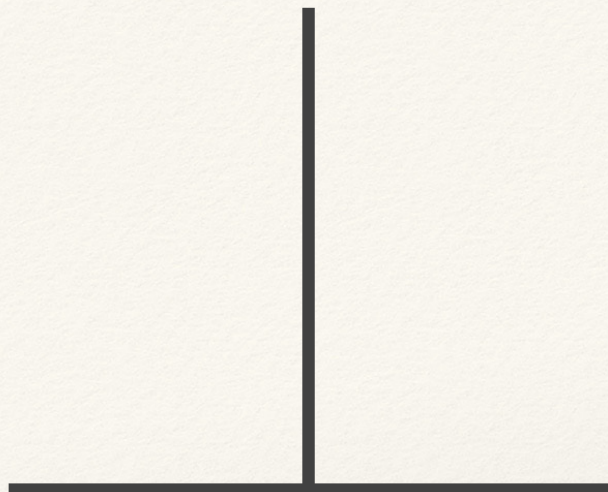
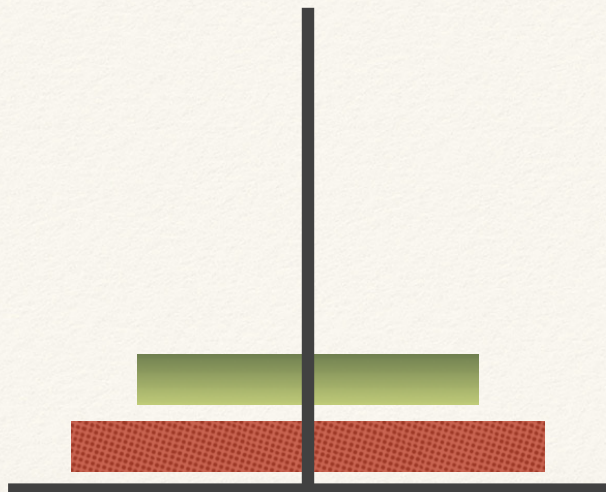
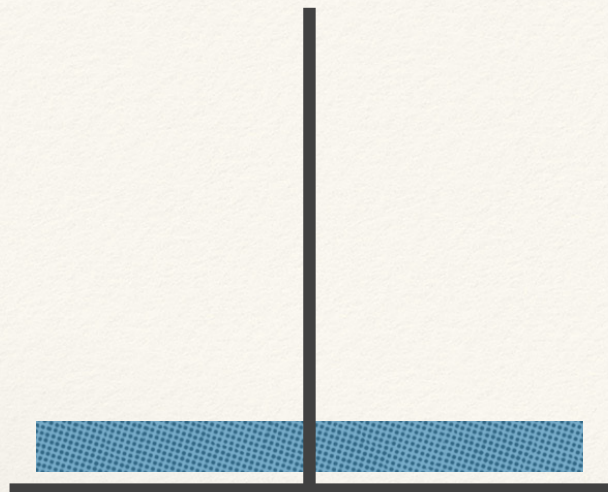
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- ❖ Step 1. Move the stack of all but the largest disk from the first to the second post.
- ❖ Step 2. Then move the largest disk from the first post to the third post.
- ❖ Step 3. Then move the stack from the second to the third post.
- ❖ Steps 1 and 3 are *recursive*.

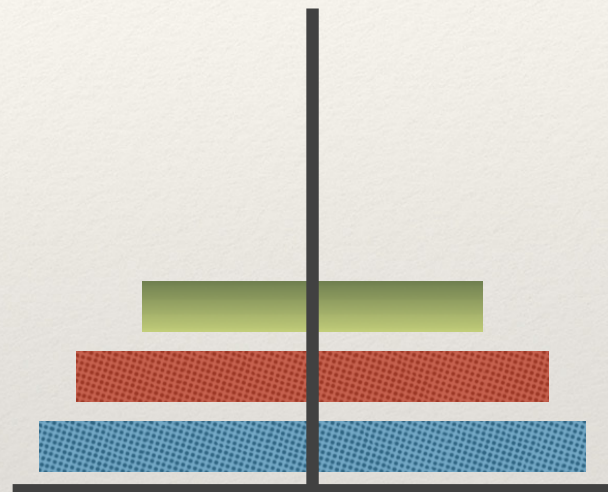
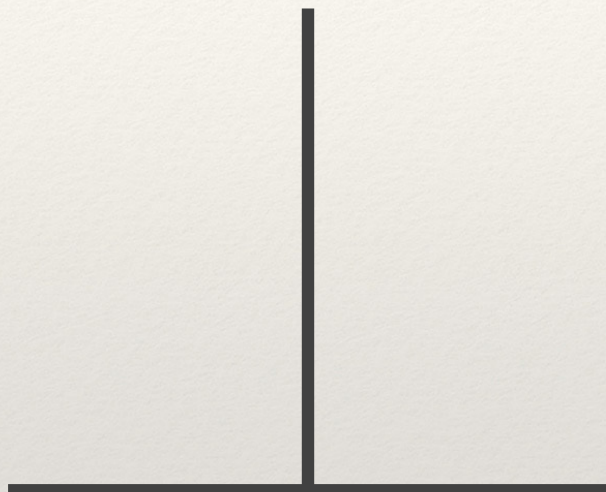
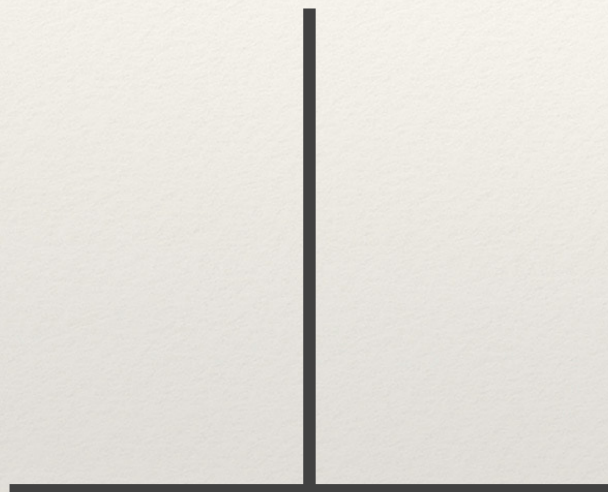
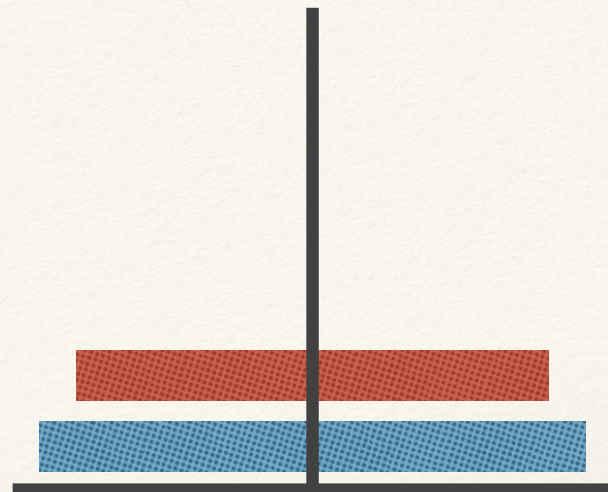
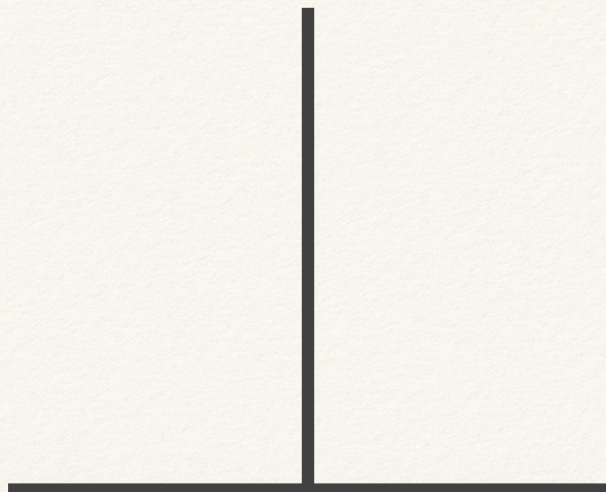
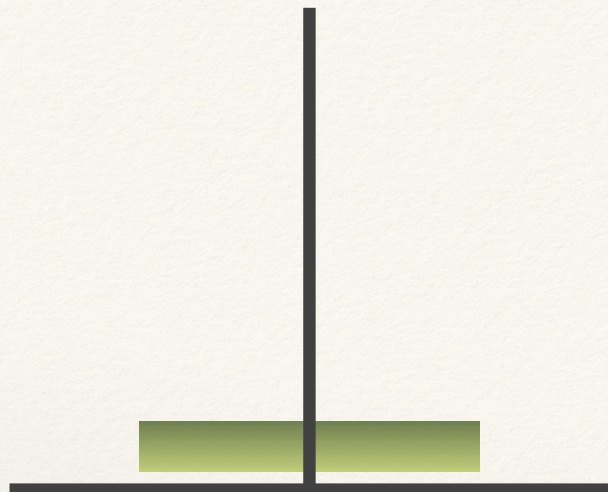














```
void hanoi (int height, int one, int two, int three)
{
    if (height <= 0) {
        return;
    }

    hanoi ( height-1, one, three, two );

    printf ("Move disk %d from %d to %d\n", height, one, three);

    hanoi ( height-1, two, one, three);
}
```

Solution to Towers of Hanoi with 3 disks:     **hanoi ( 3, 1, 2, 3 )**  
Move disk 1 from 1 to 3  
Move disk 2 from 1 to 2  
Move disk 1 from 3 to 2  
Move disk 3 from 1 to 3  
Move disk 1 from 2 to 1  
Move disk 2 from 2 to 3  
Move disk 1 from 1 to 3



---

# But when is recursion not the answer?

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- ❖ The Fibonacci sequence can be defined as follows:

$$F_0 = 0$$

$$F_1 = 1$$

$$F_n = F_{n-1} + F_{n-2} \text{ for } n \geq 2$$

- ❖ The sequence looks like:

0, 1, 1, 2, 3, 5, 8, 13, 21, ...

$F_0$   $F_1$   $F_2$   $F_3$   $F_4$   $F_5$   $F_6$   $F_7$   $F_8$



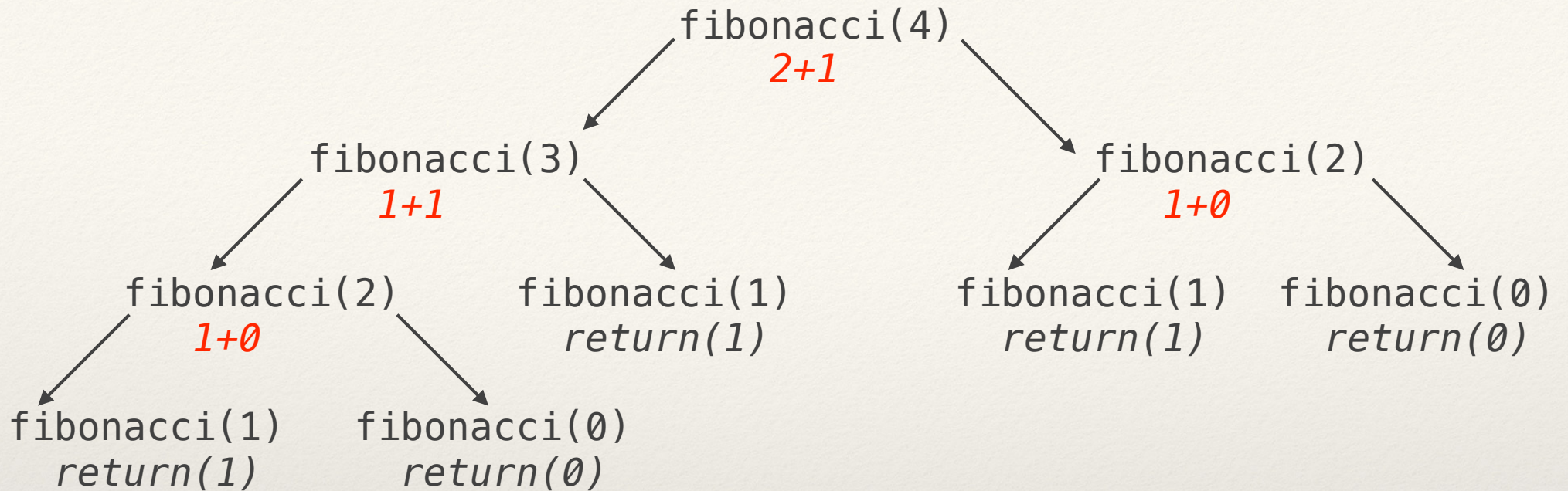
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# Recursive Fibonacci

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```
int fibonacci ( int n ) {  
    if ( n <= 0 ) {  
        return ( 0 );  
    } else if ( n == 1 ) {  
        return ( 1 );  
    } else {  
        return ( fibonacci(n-1) + fibonacci(n-2) );  
    }  
}
```





```
int main ( int argc, char *argv[] )
{
    int i, n;
    n = atoi ( argv[1] );
    for ( i=0; i<=n; i++ ) {
        printf ( "%d ", fibonacci(i) );
    }
    printf ( "\n" );
    return (0);
}
```

```
$ ./fib 8
0 1 1 2 3 5 8 13 21
```



---

# But...

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- ❖ The recursive solution for the Fibonacci sequence is very inefficient.
  - ❖ It requires a large number of function calls.
- ❖ The problem is that we have stated the problem incorrectly. We are not looking to find the  $n$ th term but are instead looking at a way to construct a sequence.
- ❖ This is better done as an iterative (loop-based) function.



```
int fibonacci ( int n ) {  
    int i;  
    int oneBack, twoBack, current;  
  
    if ( n <= 0 ) {  
        return (0);  
    } else if ( n == 1 ) {  
        return (1);  
    } else {  
        twoBack = 0;  
        oneBack = 1;  
        for ( i=2; i<=n; i++ ) {  
            current = twoBack + oneBack;  
            twoBack = oneBack;  
            oneBack = current;  
        }  
        return (current);  
    }  
}
```