On Polynomial Methods in Combinatorics

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Chapter 1

Introduction

1.1 Preface to this Draft

This draft is released to hopefully give the reader a clearer picture of the report before the presentations next week. The reader ought to be aware that the report is not fully complete, and is lacking polish in certain areas that we be remedied before the final submission. Polynomial methods are surprisingly powerful tools in combinatorics and are beautiful in their non-constructive elegance. To get a summary of what this manuscript covers, I recommend reading Section 2.3, the short Chapter 4, Sections 5.4 and 5.5, and finally Chapter 6 upon which I will be basing my presentation.

The following is a short exposition of the polynomial method in combinatorics. It is perhaps already misleading to call this the polynomial method as there are in fact a collection of polynomial methods that have applications in combinatorics, the first example of this being known since the 1990s where Alon's proved a combinatorial version of Hibert's Nullstellensatz.[1] We examine this result in detail in Chapter 3. In 2008 Dvir produced a remarkably short resolution to the finite field analogue of the Kakeya Conjecture which provided a new framework and enthusiasm for the polynomial method in combinatorial problems.[2] We will examine this proof in the next chapter.

The most striking feature of the following proofs is that they leverage polynomials in problems which on the surface appear not to have anything to do with polynomials. Generally, we will try to capture the combinatorial structures of these problems using polynomials and then use methods from algebraic geometry to study these polynomials and provide bounds for the combinatorial problems. We will sketch here the general strategies of the chapters of this manuscript.

In Chapter 2 we will use a low degree polynomial to interpolate a certain set of objects. We will then show that the zero set of the polynomial can not contain too many objects of this type due to the inherent structure of these objects. In Chapter 4 we will . . . . Chapter 5 introduces the idea of using a polynomial to partition the plane into a degree-controlled number of cells each containing at most some uniform amount of points. Chapter 6 will discuss a proof that first uses a similar strategy to Chapter 2, however in this proof we

do some-thing indeed

weak

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will be interpolating objects of one type and then proving that the polynomial's zero set cannot contain too many objects of some second related type. We then provide a new proof of the same theorem which uses an extension of the polynomial partitioning seen in 5, the key difference being that instead of controlling the number of points in each cell we will control the number of varieties that intersect any given cell. Finally, in Chapter 3 we will provide an example of the polynomial method being used outside the context of incidence geometry.

1.2 Why polynomials?

Points to possibly mention here:

> this section

1. choosing a polynomial has D^n degrees of freedom $[\mathrm{Poly}_D(\mathbb{F}^n) \sim D^n]$

2. polynomials behave rigidly on lines, having only D degrees of freedom.

3. "parameter counting - vanishing lemma" method.

4. Non constructive method of finding a polynomial is reminiscent of probabilistic method.[3] [4]

5. ill-formed thoughts that polynomials extract information about finite series analogous to generating functions and infinite series.

Notation

We introduce some convenient notation here. We write that $A \lesssim_n B$ to mean that there exists some constant $C(n)$ which depends on n such that $A \leq C(n)B$. Further, we write that $A \sim_n B$ if $A \lesssim_n B$ and $B \lesssim_n A$.

We write $\mathrm{Poly}_D(\mathbb{K}^n)$ to represent the space of polynomials in n variables with coefficients in $\mathbb{K}$ and degree at most D.

The indicator function $\mathbb{1}$ is defined on logical statements X as follows:

$$\mathbb{1}[X] = \begin{cases} 1 & \text{if X is true,} \\ 0 & \text{if X is false.} \end{cases}$$

For any function $f : \mathbb{R}^n \to \mathbb{R}$ let us denote the zero set of f by $Z(f) = \{x \in \mathbb{R}^n \mid f(x) = 0\}$. We borrow from Computer Science the big O notation. For functions $f, g : \mathbb{N}^+ \to \mathbb{R}$ we write:

$$f(N) = O(g(N)) \iff \exists N_0, M \in \mathbb{N} \text{ such that } f(n) \leq Mg(n) \forall n > N_0$$
$$f(N) = \Omega(g(n)) \iff g(N) = O(f(N)).$$

These can be thought of as asymtopic upper and lower bounds respectively.

Chapter 2

The Kakeya Problem in Finite Fields

Before we can discuss the Kakeya problem in finite fields, and its rather surprising resolution, we ought to first discuss the origin and history of the problem. Work on the Kakeya problem can be traced back to the Russian mathematician Abram Besicovitch in 1917. While working on a problem in Riemann integration, Besicovitch reduced it to the question of the existence of planar sets of measure zero which contain a line segment in every direction. In 1920, Besicovitch constructed such a set and published in a Russian Journal.

However, 1917 was a turbulent year as it marked the end of the Russian Empire and the start of the Russian civil war. Due to this and the ensuing blockade of Russian ports there was scarce communication with the outside world. Thus Besicovitch could not have known of a Japanese mathematician Kakeya who asked also in 1917 a related question: What is the smallest area of a convex set within which one can rotate a needle by 180 degrees in the plane? Julius Pal answered this question in 1921 with the equilateral triangle. The more interesting problem without the convexity condition remained open. In 1924, after leaving the newly formed Soviet Union for Copenhagen, Besicovitch discovered this problem and by modifying his previous construction produced a solution in 1925. This lead to the more general questions being asked about Kakeya sets in higher dimensions.

Definition 2.0.1 (Kakeya Set in $\mathbb{R}^n$). A Kakeya set is a set $A \subset \mathbb{R}^n$ that contains a unit segment in every direction.

Besicovitch's construction showed that these sets can have arbitrarily small measures, even attaining zero, in $\mathbb{R}^2$. Further, a straightforward construction produces these measure zero sets in dimensions > 2.

The natural question then arises, what is the dimension of such sets? There are many notions of dimensions that can be investigated, but we restrict ourselves to the Minkowski and Hausdorff dimensions.

Definition 2.0.2 (Minkowski Dimension). Given a set $S \subset \mathbb{R}^n$, define $N(\varepsilon)$ to be the

number of boxes of side length ε required to cover the set. The Minkowski Dimension of the set S is then defined as:

$$\dim_M(S) = \lim_{\varepsilon \to 0} \frac{\log(N(\varepsilon))}{\log(1/\varepsilon)}.$$

If this limit does not exist, one can still define the upper and lower Minkowski dimensions, $\dim_{M_{\mathrm{upper}}}$ and $\dim_{M_{\mathrm{lower}}}$, by taking the limit superior and limit inferior respectively.

Definition 2.0.3 (Hausdorff Dimension)**.** We define the d-dimensional Hausdorff measure of a set $S \subset \mathbb{R}^n$ as:

$$\mathcal{H}^d(S) = \lim_{r \to 0} \inf \left\{ \sum_i r_i^d : \text{ there is a countable cover of } S \text{ by balls with radii } 0 < r_i < r \right\}$$

Then we can define the Hausdorff dimension of the set S to be:

$$\dim_H(S) = \inf\{d \geq 0 : \mathcal{H}^d(S) = 0\}.$$

These dimensions are related by the following inequality when they are all defined:

$$\dim_H \leq \dim_{M_{\mathrm{lower}}} \leq \dim_{M_{\mathrm{upper}}}.$$

In 1971, Davies produced a solution for the 2 dimensional case, proving that although the measure of a Kakeya set can be arbitrarily small, it must have Hausdorff (and hence Minkowski) dimension of 2.[5] This resulted in the following conjectures:

Conjecture 1 (Kakeya Conjecture for the Minkowski Dimension)**.** Let A be a Kakeya set in $\mathbb{R}^n$. Then $\dim_M(A) = n$.

Conjecture 2 (Kakeya Conjecture for the Hausdorff Dimension)**.** Let A be a Kakeya set in $\mathbb{R}^n$. Then $\dim_H(A) = n$.

In a survey on the problem, Wolff proposed a finite analogue to the Kakeya Conjecture.[6]

A Kakeya set in $\mathbb{F}_p^n$ is a set that contains a line in every direction. The following section formalises the notion of a line in a finite set. The finite analogue to Conjectures 1 and 2 is:

Theorem 2.0.1 (Kakeya Conjecture in Finite Fields)**.** *If $A \subset \mathbb{F}_p^n$ contains a translate of every line, then $|A| \gtrsim_n p^n$.*

This conjecture had a significant influence on the subject, inspiring work on the sum-product phenomenon in finite fields. From its postulation in 1999, little progress was made in the following years and it was assumed that the problem was roughly as difficult as the Euclidean case.

In 2008, Dvir published a remarkably simple proof of Theorem 2.0.1 using basic facts about polynomials. This proof revitalised interest in the polynomial method, and we shall explore its details later in this chapter.
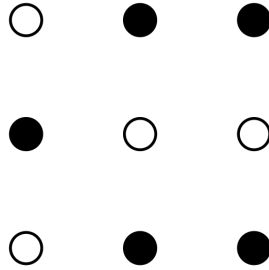
big example

Figure 2.1: An example of a Kakeya set (shaded) in $\mathbb{F}_3^2$.

2.1 Introduction to Finite Fields

Definition 2.1.1 (Finite Field). A finite field, $\mathbb{F}$, is a finite set that forms a field. That is, it is closed under addition, subtraction, multiplication, and non-zero division. The number of elements of a finite field, $|\mathbb{F}|$, is called the order of the finite field.

A finite field of order q exists if and only if $q = p^k$ for some prime p and integer k.

Lemma 2.1.1. *Each element X in a finite field $\mathbb{F}$ satisfies the identity:*

$$X^{|\mathbb{F}|} - X = 0$$

identically in $\mathbb{F}$.

This lemma follows immediately from Fermat's Little Theorem.

Analogous to the Euclidean case, we define lines in $\mathbb{F}_p^n$ as the set:

$$\mathcal{L} = \{x + ty : x, y \in \mathbb{F}_p^n, t \in \mathbb{F}_p\}.$$

It should be noted that a line in a finite field $\mathbb{F}_p^n$ contains exactly $|\mathbb{F}_p|$ points.

2.2 Combinatorial attempts at the proof

To truly show the power of the polynomial method, we shall first explore purely combinational attempts at estimating the sizes of Kakeya sets. All of the following estimates are much weaker than the bound in Theorem 2.0.1. We fix a finite field $\mathbb{F} = \mathbb{F}_{p^k}$ where p is a prime.

Our first estimate gives a good bound for the $n = 2$ case for the finite field Kakeya problem.

Lemma 2.2.1. *Suppose $s < |\mathbb{F}|$. If $l_1, \ldots, l_s$ are distinct lines in $\mathbb{F}^n$, then their union has cardinality at least $(1/2)|\mathbb{F}|s$.*

In particular, if $A \subset \mathbb{F}^n$ is a Kakeya set, then we have:

$$|A| \gtrsim |\mathbb{F}|^2$$

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Proof. We add the lines one at a time and track the cardinality of their union. The first line contains $|\mathbb{F}|$ points. The second line must contain at least $|\mathbb{F}| - 1$ points not in the first line. Similarly the third line must contain at least $|\mathbb{F}| - 2$ points not in the first two lines, and so on. Thus the number of distinct points in the union of all s lines is given by:

$$\sum_{i=1}^{s} |\mathbb{F}| - i + 1 > (1/2)|\mathbb{F}|s.$$

A Kakeya set always contains at least $|\mathbb{F}|$ lines, so setting $s = |\mathbb{F}|$ above yields:

$$|A| > (1/2)|\mathbb{F}|^2$$

$\square$

2.2.1 Bush Argument

Bourgain produced one of the first non-trivial estimates of the dimension in 1991.[7] We present the finite field analogue of his argument here.[4]

Theorem 2.2.2 (Bush Argument)**.** *If A is a Kakeya set, then we have:*

$$|A| \gtrsim |\mathbb{F}|^{\frac{n+1}{2}}$$

Proof. Let μ be a fixed multiplicity parameter to be chosen later. Then either there exists a point $p \in A$ such that there are μ lines passing through p, or else every point in A has less than μ lines passing through it.

In the first case, since the lines have distinct directions they must become disjoint when p is removed. Hence,

$$|A| \gtrsim \mu|\mathbb{F}|.$$

In the latter case, by considering the double sum and double counting:

$$\sum_{\ell \in K} \sum_{p \in A} \mathbb{1}[p \in \ell] = \sum_{\ell \in K} |\mathbb{F}| = |\mathbb{F}|^{n-1}|\mathbb{F}|$$

$$\sum_{p \in A} \sum_{\ell \in K} \mathbb{1}[p \in \ell] \leq \sum_{p \in A} \mu = |A|\mu$$

$$\implies \frac{|\mathbb{F}|^n}{\mu} \leq |A|.$$

Now optimising μ:

$$\frac{|\mathbb{F}|^n}{\mu} \sim \mu|\mathbb{F}|$$

$$\mu \sim |\mathbb{F}|^{\frac{n-1}{2}}.$$

This yields the bound:

$$|A| \gtrsim |\mathbb{F}||\mathbb{F}|^{\frac{n-1}{2}} \sim |\mathbb{F}|^{\frac{n+1}{2}}.$$

$\square$

2.2.2 Hairbrush Argument

Due to Wolff. [8]

Theorem 2.2.3 (Hair Brush Argument). *Suppose $l_1, \ldots, l_M$ are lines in $\mathbb{F}^n$, and that at most $|\mathbb{F}| + 1$ of the lines lie in any plane. Then their union has cardinality at least*

$$\frac{1}{3}|\mathbb{F}|^{3/2}M^{1/2}.$$

In particular, if A is a Kakeya set, then we have:

$$|A| \gtrsim |\mathbb{F}|^{\frac{n+2}{2}}$$

Proof. By Lemma 2.2.1, if $A_m \subset \mathbb{F}^2$ is a set containing a line in m different directions then $|A_m| \gtrsim m|\mathbb{F}|$.

Let μ be a fixed multiplicity parameter to be chosen later. We say a line ℓ is μ-rich if for at least $|\mathbb{F}|/2$ points $p \in \ell$ there are μ lines distinct from p in A passing through p. Then either there exists a μ-rich line or there does not.

Suppose a μ-rich line exists, and denote this line by ℓ_μ. Consider the family of 2-dimensional planes Π passing through ℓ_μ. If a line ℓ intersects ℓ_μ then there exists a unique plane $\pi \in \Pi$ such that $\ell, \ell_\mu \in \pi$. Let $\mathcal{L}_\pi := \{\ell \subset \pi \mid \ell \cap \ell_\mu \neq \emptyset\}$. The set $A \cap \pi$ is a set in $\mathbb{F}^2$ that contains at least $|\mathcal{L}_\pi|$ lines, and hence by Lemma 2.2.1 we have:

$$|A \cap \pi| \gtrsim |\mathcal{L}_\pi||\mathbb{F}|.$$

Thus,

$$|A| \geq \sum_{\pi \in \Pi} |(A \cap pi) \backslash \ell_\mu| \gtrsim |\mathbb{F}| \sum_{\pi \in \Pi} |\mathcal{L}_\pi| \gtrsim \mu|\mathbb{F}|^2,$$

where the last inequality comes from the fact ℓ_μ intersects at least $\mu|\mathbb{F}|/2$ lines, which produces $\mu|\mathbb{F}|/2$ planes as all these lines are distinct.

Suppose there does not exist a μ-rich line. Let

$$A' = \{p \in A \mid p \text{ belongs to at most } \mu \text{ lines in } A\}.$$

Since we assume there does not exist a μ-rich line in A, for any line $\ell \subset A$ we have $|A' \cap \ell| > |F|/2$. Proceeding by double counting,

$$|A| \geq |A'| \geq \frac{1}{\mu} \sum_{p \in A'} \sum_{\ell \in A} \mathbb{1}[p \in \ell]$$

$$= \frac{1}{\mu} \sum_{\ell \in A} |A' \cap \ell| \gtrsim \frac{1}{\mu} |\mathbb{F}|^{n-1} |\mathbb{F}|$$

$$\implies |A| \gtrsim \frac{|\mathbb{F}|^n}{\mu}.$$

We now optimise μ:

$$\frac{|\mathbb{F}|^n}{\mu} \sim \mu |\mathbb{F}|^2$$

$$\mu \sim |\mathbb{F}|^{\frac{n-2}{2}}$$

Thus we achieve the bound:

$$|A| \gtrsim |\mathbb{F}|^2 |\mathbb{F}|^{\frac{n-2}{2}} \sim |\mathbb{F}|^{\frac{n+2}{2}}.$$

$\square$

2.3 Dvir's Proof

We shall prove this theorem via 3 surprisingly simple lemmas. The general strategy of the proof is to assume a Kakeya set A is of a limited size and find a low-degree non-zero polynomial that interpolates (contains in its zero set) the points in A. We then show that due to the structure of the Kakeya set the polynomial's zero set must contain many more points. This contradicts the fact our polynomial is non-zero and establishes the Kakeya Conjecture over finite fields.

Lemma 2.3.1 (Parameter Counting). *Let $\mathbb{K}$ be a (not necessarily finite) field. If $A \subset \mathbb{K}^n$ and $|A| < \binom{n+D}{n}$, there exists a non-zero polynomial $P(x_1, \ldots, x_n)$ of degree D that vanishes on A.*

Proof. We first show the dimension of $\mathrm{Poly}_D(\mathbb{K}^n)$ is $\binom{D+n}{n}$. A basis for $\mathrm{Poly}_D(\mathbb{K}^n)$ is given by monomials of the form $x_1^{D_1} \ldots x_n^{D_n}$, where $\sum_i D_i \leq D$, hence we just need to count the number of monomials.

We can map a monomial $x_1^{D_1} \ldots x_n^{D_n}$ to a string of D $\star$'s and n |'s as follows. Begin with D_1 $\star$'s, then place one |. We put now D_2 $\star$'s, and place a second |. We continue until we have placed D_n $\star$'s followed by an nth |. Finally we place $D - \sum_i D_i \star$'s. This is a bijective map between the monomials in $\mathrm{Poly}_D(\mathbb{K}^n)$ and all the strings of D $\star$'s and n |'s. To count the strings, fix the n |'s. Now we have $n+1$ bins to distribute our D $\star$'s. Therefore we have by the stars and bars theorem:

$$\mathrm{Poly}_D(\mathbb{K}^n) = \binom{n+1+D-1}{n+1-1} = \binom{n+D}{n}.$$

Let now $p_1, \ldots, p_{|A|}$ be the points of A. We consider the evaluation map $E : \mathrm{Poly}_D(\mathbb{K}^n) \to \mathbb{K}^{|A|}$ defined by:

$$E(Q) = \big(Q(p_1), \ldots, Q(p_{|A|})\big).$$

This map is clearly linear. Its kernel $\ker E$ is exactly the set of polynomials in $\mathrm{Poly}_D(\mathbb{K}^n)$ that vanish on A. By assumption, the dimension of $\mathrm{Poly}_D(\mathbb{K}^n)$ is greater than A, so the dimension of the domain of E is greater than the codomain of E. By the rank-nullity theorem, we conclude E must have a non-trivial kernel. Thus there exists a non-zero polynomial $P \in \mathrm{Poly}_D(\mathbb{K}^n)$ that vanishes on A. $\square$

Note that if $D = |\mathbb{F}| - 1$, and $|A| \leq \binom{|\mathbb{F}|+n-1}{|\mathbb{F}|-1} = \binom{|\mathbb{F}|+n-1}{n}$ we have a polynomial of degree $|\mathbb{F}| - 1$ that vanishes on A. Since $\frac{|\mathbb{F}|^n}{n!} < \binom{|\mathbb{F}|+n-1}{n}$, we can definitely find such a polynomial when $|A| \leq \frac{|\mathbb{F}|^n}{n!}$.

Lemma 2.3.2. *Suppose $A \subset \mathbb{F}^n$ contains a line in every direction, and suppose that there exists a non-zero polynomial P with degree $D < |\mathbb{F}|$ that vanishes on A. Then there exists a non-zero degree D polynomial $\bar{P}$ that vanishes everywhere on $\mathbb{F}^n$.*

Proof. Choose a line in A, say $\ell = \{x + tz : t \in \mathbb{F}\}$ with $x \in \mathbb{F}^n$ and $z \in \mathbb{F}^n/\mathbb{F}^\times$. Now we consider the restriction of our polynomial P to the line ℓ, $P_{|\ell}$. Recall P is a sum of monomials, and we use multi-index notation here with $\alpha = (\alpha_1, \alpha_2, \ldots, \alpha_n)$, $\alpha_i \in \mathbb{N} \cup \{0\}$ and $|\alpha| = \sum \alpha_i$. P can be written as:

$$P(x_1, x_2, \ldots, x_n) = \sum_{|\alpha| \leq D} c_\alpha x_1^{\alpha_1} x_2^{\alpha_2} \ldots x_n^{\alpha_n}.$$

Now $P_{|\ell}$ can be written:

$$P_{|\ell} = P(x + tz) = Q_{x,z}(t) = \sum_{|\alpha| \leq D} c_\alpha \prod_i (x_i + tz_i)^{\alpha_i}.$$

We now wish to examine the degree D term of Q, which is achieved by picking the tz_i terms from each bracket in the product above. This gives the degree D component of Q, $Q_{x,z,D}$, which has the form:

$$Q_{x,z,D} = t^D Q_D(z) = t^D \sum_{|\alpha|=D} c_\alpha \prod_i z_i.$$

Now if $P_{|\ell}$ vanishes everywhere on ℓ, since its dependence on t is given by a polynomial of degree less than $|\mathbb{F}|$, all its coefficients must be zero. This is clear from the factor theorem, as we could write the roots of $P_{|\ell}$ as $(t - k_1)(t - k_2) \ldots (t - k_{|\mathbb{F}|})$, but this contradicts the fact P is of degree $D < |\mathbb{F}|$.

Notice that $Q_{x,z,D}$ no longer depends on x, but on z alone. In particular $Q_D(z) = 0$, but z was an arbitrary element of $\mathbb{F}^n/\mathbb{F}^\times$, and $Q_D(z)$ also vanishes at zero, so it vanishes everywhere. Thus we can pick $\bar{P} = Q_D$, and we are done. $\square$

Lemma 2.3.3. *Let P be a non-zero polynomial on $\mathbb{F}^n$ with degree less than $|\mathbb{F}|$. Then P does not vanish everywhere.*

Proof. We proceed by induction on n. For $n = 1$, a non-zero polynomial that vanishes everywhere has $|\mathbb{F}|$ roots, so must be at least of degree $|\mathbb{F}|$. Let us assume that the statement holds in $\mathbb{F}^{n-1}$, we now prove it must also hold for $\mathbb{F}^N$.

We let $x_1, \ldots, x_n$ be coordinates on $\mathbb{F}^n$, and we write P in the form:

$$P(x_1, \ldots, x_n) = \sum_{j=n}^{|\mathbb{F}|-1} P_j(x_1, \ldots x_{n-1}) x_n^j.$$

Each P_j are polynomials in $x_1, \ldots x_{n-1}$ of degree less than $|\mathbb{F}|$. Fix $x_1, \ldots x_{n-1}$, and let x_n vary. Now we have a polynomial in x_n of degree less than $|\mathbb{F}|$ that vanishes for all $x_n \in \mathbb{F}$. By the base case this must be the zero polynomial. So each $P_j(x_1, \ldots, x_{n-1}) = 0$ for all j and for all $(x_1, \ldots x_{n-1}) \in \mathbb{F}^{n-1}$. Now by induction on n, each P_j is the zero polynomial. Then P is the zero polynomial as well. $\qquad\square$

Proof of Theorem 2.0.1. Assume $A \subset \mathbb{F}^n$ is a Kakeya set, and that $|A| \leq \frac{|\mathbb{F}|^n}{n!}$. Then by Lemma 2.3.1 we can find a non-zero polynomial, say P, that vanishes on A. Now by Lemma 2.3.2 there exists a non-zero polynomial $\bar{P}$ that vanishes everywhere on $\mathbb{F}^n$, and has degree less than $|\mathbb{F}|$. Finally, Lemma 2.3.3 says that such a $\bar{P}$ is necessarily the zero polynomial, a contradiction. We conclude that $|A| > \frac{|\mathbb{F}|^n}{n!}$, or in other words:

$$|A| \gtrsim_n |\mathbb{F}|^n.$$

$\square$

Chapter 3

The Polynomial Method in Additive Combinatorics

So far our discussion has only centred around the usefulness of polynomial methods in combinatorial problems with a geometric flavour. In this chapter we shall discuss applications to the exciting field of additive combinatorics, where in 1990 Alon provided perhaps the first example of the polynomial method in action.[1] We present Michalek's short, elementary, and direct proof of the combinatorial nullstellensatz.[9]

3.1   Combinatorial Nullstellensatz

Theorem 3.1.1 (Combinatorial Nullstellensatz)**.** *Let $\mathbb{K}$ be a (not necessarily finite) field, and let $P(X_1, \ldots, X_n) \in \mathbb{K}[X_1, \ldots, X_n]$ be a polynomial in n variables with coefficients in $\mathbb{K}$. Suppose the coefficient of $X_1^{k_1} X_2^{k_2} \ldots X_n^{k_n}$ is non-zero and further suppose $\deg P = \sum_{i=1}^{n} k_i$, where each k_i is a non-negative integer.*

Then for any subsets $A_1, \ldots A_n$ of $\mathbb{K}$ satisfying $|A_i| > k_i$ for each $1 \le i \le n$ there exist $a_1 \in A_1, \ldots, a_n \in A_n$ such that $P(a_1, \ldots, a_n) \ne 0$.

Proof. We proceed by induction on $\deg P = D$. When $D = 1$, P is simply a linear combination of n variables say $P(X_1, \ldots X_m) = c_1 X_1 + \ldots c_n X_n$. Without loss of generality assume x_1 has a non-zero coefficient and consider the sets $A_i = \{a_{i,1}, a_{i,2}\}$. Suppose P at the point $(a_{1,1}, a_{2,1}, \ldots, a_{n,1})$ is zero. We can then determine $c_1 = -\frac{c_2 a_{2,1} + c_3 a_{3,1} + \cdots + c_n a_{n,1}}{a_{1,1}}$. Now evaluating at $(a_{1,2}, a_{2,1}, \ldots, a_{n,1})$ we see that this is zero only when $a_{1,1} = a_{1,2}$, so our theorem holds.

Now let us assume the theorem holds for $\deg P = D - 1$, and prove for $\deg P = D$. Suppose that P satisfies the assumptions of the theorem but $P(x) = 0$ for every $x \in A_1 \times \cdots \times A_n$. Without loss of generality $k_1 > 0$. Fixing $a \in A_1$ we can write

$$P = (x_1 - a)Q + R \tag{$\dagger$}$$

by the usual long division of polynomials. The degree of R in x_1 must be strictly less

than $\deg(X_1 - a)$, so R is independent of X_1 terms. Thus it follows that Q must have a monomial with non-zero coefficient of the form $X_1^{k_1-1}X_2^{k_2}\ldots X_n^{k_n}$ and $\deg(Q) = D - 1$.

Take any $x \in \{a\} \times A_2 \times \cdots \times A_n$ and evaluate (†). Since $P(x) = 0$ it follows that $R(x) = 0$. Hence R vanishes identically on the slice $\{a\} \times A_2 \times \cdots \times A_n$. Since R is independent of X_1 it must also vanish identically on $A_1 \times A_2 \times \cdots \times A_n$. Now take any $x \in A_1\backslash\{a\} \times A_2 \times \cdots \times A_n$ and evaluate (†). Since the $(X_1 - a)$ term is non-zero, $Q(x) = 0$. So Q vanishes on all $x \in A_1\backslash\{a\} \times A_2 \times \cdots \times A_n$, which contradicts the inductive hypothesis. $\qquad\square$

3.2 Cauchy-Davenport Theorem

Theorem 3.2.1 (Cauchy-Davenport Theorem)**.** *Let A, B be non-empty subsets of $\mathbb{Z}_p$ for some p prime. Define their sumset $A + B$ as follows:*

$$A + B = \{x \in \mathbb{Z}_p \mid x = a + b \text{ for some } a \in A, \ b \in B\}.$$

Then we have:

$$|A + B| \geq \min\{p, |A| + |B| - 1\}.$$

Generalise this!

Proof. Let us tackle the two cases separately. First, assume that $\min\{p, |A| + |B| - 1\} = p$. Then if $|A| + |B| > p$, A and B must intersect. For some $g \in \mathbb{Z}_p$ denote the set $\{g - x \mid x \in B, \} \subset \mathbb{Z}_p$ as $g - B$. Since $|g - B| = |B|$, we have that $g - B$ and A must intersect as well. Thus there exists some $a \in A$, $b \in B$ such that:

expand PHP

$$g - b = a$$
$$g = a + b.$$

Our choice of g was arbitrary, so it follows that $A + B = \mathbb{Z}_p$ and hence $|A + B| = p$.

Now assume that $\min\{p, |A| + |B| - 1\} = |A| + |B| - 1$. Then if the theorem is false we have $|A + B| \leq |A| + |B| - 2$, so there exists some $C \subset \mathbb{Z}_p$ such that $A + B \subset C$ and $|C| = |A| + |B| - 2$. Now let us define a polynomial $f(x, y) \in \mathbb{Z}_p[X, Y]$ as:

$$f(X, Y) = \prod_{c \in C}(X + Y - c).$$

Since $A + B \subset C$, $f(a, b) = 0$ for all $(a, b) \in A \times B$. Further, the degree of f is $\deg f = |C| = |A| + |B| - 2$. We can now appeal to the Combinatorial Nullstellensatz to yield a contradiction. Let $k_1 = |A| - 1$, and $k_2 = |B| - 1$. Now $\deg f = k_1 + k_2$, and the coefficient of $x^{k_1}y^{k_2}$ is $\binom{|A|+|B|-2}{|A|-1}$ which is non-zero in $\mathbb{Z}_p$ as the binomial coefficient is not divisible by p if all its factors are less than p. Applying Theorem 3.1.1 we see that there must exist some $(a, b) \in A \times B$ such that $f(a, b) \neq 0$, a contradiction. $\qquad\square$

Chapter 4

The Joints Problem

The joints problem was first posed in 1990 by Chazelle et al.[10] They focused on the 3-dimensional case of the problem, establishing a lower bound of $\Omega(N^{3/2})$ and an upper bound estimate of $O(N^{7/4})$. This upper bound exponent has fallen gradually throughout the years, the best result being due to Sharir and Feldman who established an upper bound of $O(N^{112/69})$. Their proof uses an array of tools from combinatorial geometry, including forbidden subgraphs in extremal graph theory, space decomposition techniques, and some results from the geometry of lines in space.

4.1   Background

Let $\mathcal{L}$ be a set of distinct lines in $\mathbb{R}^n$. A joint of $\mathcal{L}$ is a point which lies in three non-coplanar lines of $\mathcal{L}$. The joints problem consists of setting a sharp lower bound on the maximal number of joints that can be formed from a configuration of L distinct lines. We denote this quantity $J(L)$. In other words $J(L)$ is the supremeum over all configurations of lines in $\mathbb{R}^n$ in the number of joints.

4.2   Examples

We shall begin by examining an example based on a grid, with the hopes of gaining better intuition about the problem and formulating a conjecture.

Example 4.1. *Consider an $N \times N \times N$ regular grid of integer coordinates. We shall give a collection of lines such that each point of this grid is a joint for the collection. Let $\mathcal{L}$ be the collection of all lines parallel to any of the Cartesian axes that intersect this a point in this grid. For each horizontal $N \times N$ layer, there are $N + N = 2N$ lines that intersect our grid. There are N layers, so we obtain $2N^2$ distinct lines in this manner. Finally we need to account for the N^2 lines perpendicular to the $N \times N$ layers. This leaves us with $|\mathcal{L}| = 3N^2$ lines forming N^3 joints. The number of joints is thus $\sim |\mathcal{L}|^{3/2}$.*

We can extend this example to higher dimensional grids easily.
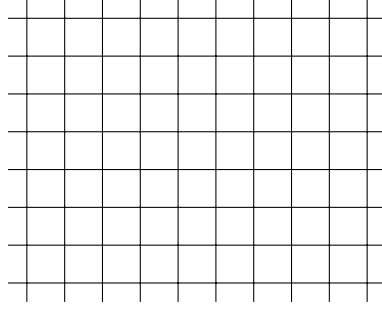
Figure 4.1: A $N \times N$ layer of our grid.

Example 4.2. *If we have an* $\underbrace{N \times \cdots \times N}_{n \text{ Dimensions}}$ *regular grid of integer coordinates in* $\mathbb{R}^n$, *we can construct an example by a straightforward extension of the above example. Each additional dimension increases the number of lines by a factor of* N, *this can be seen by considering each new dimension as a layering of the previous set along the new axis. Thus we can see that* $\sim N^{n-1}$ *lines form* N^n *joints in this manner. So the number of joints is* $\sim |\mathcal{L}|^{\frac{n}{n-1}}$.

It turns out that the examples illustrated above provide asymptotically maximal configurations, that is, disregarding the best constant C such that $J(L) \leq CL^{\frac{n}{n-1}}$.

4.3 Solution of the Joints Problem

This solution was first produced by Guth-Katz for the three-dimensional case,[11] and later extended to the general case by Quilodrán,[12] and independently at the same time by Kaplan-Sharir-Shustin.[13]

Theorem 4.3.1. *Any* L *lines in* $\mathbb{R}^n$ *determine* $\lesssim_n L^{\frac{n}{n-1}}$ *joints.*

We begin with the fundamental lemma to this proof.

Lemma 4.3.2. *If* $\mathcal{L}$ *is a set of lines in* $\mathbb{R}^n$ *that determines* J *joints, then one of the lines contains at most* $nJ^{\frac{1}{n}}$ *joints.*

Proof. Let P denote the lowest degree non-zero polynomial that vanishes at every joint of $\mathcal{L}$. By parameter counting, Lemma 2.3.1, the degree of P is at most $nJ^{\frac{1}{n}}$. (To see this, set $D = \lfloor nJ^{\frac{1}{n}} \rfloor$ and then $J < \binom{D+n}{n}$.)

We proceed by contradiction. Assume every line contains more than $nJ^{\frac{1}{n}}$ joints. Now P must vanish on every line in $\mathcal{L}$ as the degree of P is less than the number of joints it must interpolate.

We now examine the gradient of P at each joint in $\mathcal{L}$. We will need a fact about gradients for this, which we will encapsulate in the following lemma for clarity.

Lemma 4.3.3. *If* x *is a joint of* $\mathcal{L}$, *and if a smooth function* $F : \mathbb{R}^n \to \mathbb{R}$ *vanishes on the lines of* $\mathcal{L}$, *then* ∇F *vanishes at* x.

Proof. The joint x is contained in n non-coplanar lines $l_1, \ldots, l_n$, in directions $v_1, v_2, \ldots, v_n$ respectively. Now consider the directional derivative for a particular v_i:

$$\frac{\partial F}{\partial v_i} = \lim_{t \to 0} \frac{\overbrace{F(x + tv_i)}^{F \equiv 0 \text{ on a line in } \mathcal{L}} - \overbrace{F(x)}^{F \equiv 0 \text{ on joints}}}{t} = \frac{0}{t} = 0.$$

Notice that $\frac{\partial F}{\partial v_i} = \langle \nabla F, v_i \rangle$, so since we have this for each v_i, and the set of v_i's form a basis of $\mathbb{R}^n$, we have that $\nabla F(x) = 0$. $\qquad\square$

So we see that the partial derivatives of P vanish at each joint. The derivatives are polynomials of smaller degree than P and since P was assumed to be the minimal degree non-zero polynomial that vanishes at each joint, each derivative of P is identically zero. This implies P must be constant, which implies that there does not exist such a minimal degree polynomial, a contradiction. $\qquad\square$

Finally we can prove the main result.

Proof. Lemma 4.3.2 tells us that if we remove a line from our collection, we are removing at most $nJ(L)^{\frac{1}{n}}$ joints. By repeating this process, we get the chain of inequalities:

$$\begin{aligned}
J(L) &\leq J(L-1) + n(J(L))^{\frac{1}{n}} \\
&\leq J(L-2) + 2\left[n(J(L))^{\frac{1}{n}}\right] \\
&\leq J(L-3) + 3\left[n(J(L))^{\frac{1}{n}}\right] \\
&\vdots \\
&\leq L\left[n(J(L))^{\frac{1}{n}}\right].
\end{aligned}$$

Now we have:

$$\begin{aligned}
J(L) &\leq L\left[n(J(L))^{\frac{1}{n}}\right] \\
J(L)^{\frac{n-1}{n}} &\lesssim_n L \\
J(L) &\lesssim_n L^{\frac{n}{n-1}}
\end{aligned}$$

$\qquad\square$

Chapter 5

Szemerédi-Trotter Theorem

In this chapter we will study the application of the polynomial method to incidence geometry by proving a fundamental theorem in the field. Incidence geometry is the study of possible intersection patterns of simple geometric objects, such as lines or low degree curves. We have already seen an incidence problem in the previous chapter on the Joints problem. By developing the powerful tool of polynomial partitioning we shall see the key role that the topology of $\mathbb{R}^n$ can play in such problems, in contrast to the trivial topology of finite fields.

5.1 Background

The Szemerédi-Trotter theorem is a fundamental theorem to the field of incidence geometry, originally proved by an involved cell decomposition argument of Szemerédi and Trotter and later given a shorter proof using crossing numbers by Székely.

Theorem 5.1.1 (Szemerédi-Trotter). *Let $\mathcal{S} \subset \mathbb{R}^2$ be a finite set of points and let $\mathcal{L} \subset \mathbb{R}^2$ be a finite set of lines. We define*

$$I(\mathcal{S}, \mathcal{L}) = \{(p, \ell) \in \mathcal{S} \times \mathcal{L} \mid p \in \ell\}$$

to be the set of incidences between $\mathcal{S}$ and $\mathcal{L}$.
   Then:
$$|I(\mathcal{S}, \mathcal{L})| \lesssim (|\mathcal{S}||\mathcal{L}|)^{2/3} + |\mathcal{S}| + |\mathcal{L}|$$

5.2 The Trivial Bound

In planar geometry, we have the following dual statements: two points determine a line and every pair of lines intersect in at most one point. Using this we can prove the following bounds on $I(\mathcal{S}, \mathcal{L})$:

Theorem 5.2.1 (Trivial Bounds)**.** *For a set of points $\mathcal{S}$ and lines $\mathcal{L}$ we have*

$$I(\mathcal{S}, \mathcal{L})| \leq |\mathcal{S}|^2 + |\mathcal{L}|.$$

$$I(\mathcal{S}, \mathcal{L})| \leq |\mathcal{L}|^2 + |\mathcal{S}|.$$

Proof. To see this, count the lines that have at most one point in P on them. These contribute at most $|\mathcal{L}|$ incidences. Every other line has at least two points in $\mathcal{S}$. The total number of incidences on these lines is at most $|\mathcal{S}|^2$ as otherwise there would exist a $p \in \mathcal{S}$ that lies on over $|\mathcal{S}|$ lines, and each of these lines would have an additional point on it. This would imply there are more that $|\mathcal{S}|$ points, a contradiction.

Interchanging the roles of $\mathcal{L}$ and $\mathcal{S}$ achieves the other bound as two lines intersect in at most one point. $\qquad\square$

Theorem 5.2.2 (Second Trivial Incidence Bound)**.**

$$I(\mathcal{S}, \mathcal{L}) \lesssim |\mathcal{S}| \cdot |\mathcal{L}|^{\frac{1}{2}} + |\mathcal{L}|$$

and

$$I(\mathcal{S}, \mathcal{L}) \lesssim |\mathcal{L}| \cdot |\mathcal{S}|^{\frac{1}{2}} + |\mathcal{S}|.$$

Proof. We now bound the number of incidences.

$$
\begin{aligned}
|I(\mathcal{S}, \mathcal{L})|^2 &= \left(\sum_{\ell \in \mathcal{L}} \sum_{p \in \mathcal{S}} \mathbb{1}[p \in \ell] \right)^2 \\
&\leq |\mathcal{L}| \cdot \sum_{\ell \in \mathcal{L}} \left(\sum_{p \in \mathcal{S}} \mathbb{1}[p \in \ell] \right)^2 \qquad \text{(Cauchy-Schwarz on } \ell) \\
&= |\mathcal{L}| \cdot \sum_{p_1, p_2 \in \mathcal{S}} \sum_{\ell \in \mathcal{L}} \mathbb{1}[p_1 \in \ell] \mathbb{1}[p_2 \in \ell] \\
&\leq |\mathcal{L}| \cdot (|I(\mathcal{S}, \mathcal{L})| + |\mathcal{S}|^2) \\
&\leq |\mathcal{L}|^2 + 2|\mathcal{L}| \cdot |\mathcal{S}|^2 \qquad \text{(Using Theorem 5.2.1)}
\end{aligned}
$$

This implies

$$I(\mathcal{S}, \mathcal{L}) \lesssim |\mathcal{S}| \cdot |\mathcal{L}|^{\frac{1}{2}} + |L|.$$

extra step needed here - split up sum and do triple thing

Repeating the above proof interchanging the roles $\mathcal{S}$ and $\mathcal{L}$ achieves the other bound. $\quad\square$

5.3 Examples

We can not improve beyond our second trivial bounds in a finite field $\mathbb{F}^2$.

Example 5.1 (Finite Fields)**.** *Consider the set of points $\mathcal{S} = \mathbb{F}^2$ and let $\mathcal{L}$ be the set of all lines in $\mathbb{F}^2$. Every line contains exactly $|\mathbb{F}|$ many points of $\mathcal{S}$, so we have $|\mathbb{F}|^3$ incidences. So both sides of the second trivial bound are comparable:*

$$I(\mathcal{S}, \mathcal{L}) = |\mathbb{F}|^3 \sim (|\mathbb{F}|^2)(|\mathbb{F}|^2)^{1/2} + |\mathbb{F}|^2.$$

In contrast, the following examples seem to be the best possible over $\mathbb{R}$. We will later prove that these are the asymptotically tight case of the Szemerédi–Trotter Theorem. We define a line in $\mathbb{R}^2$ as follows:

$$\ell_{m,c} = \{(x,y) \in \mathbb{R}^2 \mid y = mx + c\}.$$

Example 5.2. *Consider the following collections in $\mathbb{R}^2$:*

$$\mathcal{S} = \{(a,b) \in \mathbb{Z}^2 \mid a \in [1,N], b \in [1, 2N^2]\}$$
$$\mathcal{L} = \{\ell_{m,c} \in \mathbb{R}^2 \mid m, c \in \mathbb{Z}, m \in [1,N], c \in [1, N^2]\}$$

The collection $\mathcal{S}$ contains $2N^3$ points and $\mathcal{L}$ contains N^3 lines. Every line in $\mathcal{L}$ contains N points in $\mathcal{S}$ as for each $x \in [1,N]$ the y-coordinate of $\ell_{m,c}$, $mx + c$, gives a different integer in $[1, 2N^2]$. Hence there are N^4 incidences. Both sides of the Szemerédi-Trotter inequality are comparable as

$$I(\mathcal{S}, \mathcal{L}) = N^4 \sim (N^3)^{\frac{2}{3}}(N^3)^{\frac{2}{3}} \sim |\mathcal{S}|^{2/3}|\mathcal{L}|^{2/3}$$

diagram?

Example 5.3. *Let $N \gg 1$ be a large even integer and let $1 < R \ll N$ be another integer. Consider the collections in $\mathbb{R}^2$:*

$$\mathcal{S} = \{(a,b) \in \mathbb{Z}^2 \mid (a,b) \in [-N/2, N/2] \times [-N/2, N/2]\}$$
$$\mathcal{L} = \{\ell \mid \ell \text{ contains between } R \text{ and } 2R \text{ points of } \mathcal{S}\}$$

We begin by estimating how many lines pass through a given point of the regular grid $\mathcal{S}$. Let $\ell \in \mathcal{L}$ and $p \in \mathcal{S}$. The closest point $p' \in \mathcal{S}$ such that $p \neq p'$ and $p' \in \ell$ must lie in a square centred at p of side length $\sim N/R$. This follows from the fact that there are at least $\sim R$ points of $\mathcal{S}$ in ℓ and hence the projections of these points to the axes can be separated by at most $\sim N/R$. Taking each possible combination of these we can conclude that there are $\lesssim N^2/R^2$ in $\mathcal{L}$ through a given point p.

We now claim that there are $\gtrsim N^2/R^2$ distinct such lines. We need only consider the points in the upper right quadrant of $\mathcal{S}$ as the problem is symmetrical. Further, we restrict ourselves to considering lines with slopes m satisfying $\frac{1}{2} < m < 2$. For such a line

to contain R points of $\mathcal{S}$ we require $m = \frac{l}{k} \in \mathbb{Q}$ with $\gcd(l, k) = 1$ and $l, k \in \left[\frac{N}{2R}, \frac{N}{R}\right]$. There are $\gtrsim N^2/R^2$ pairs, as the proportion of pairs that share a factor of 2 is $\frac{1}{2}^2$ and the proportion of pairs that share a factor of 3 is $\frac{1}{3}^2$, and in general the proportion that shares a factor of k is $\frac{1}{k}^2$. We have that $\sum_{k>1} \frac{1}{k}^2 < \frac{2}{3} < 1$ and hence there are $\gtrsim N^2/R^2$ distinct lines in $\mathcal{L}$ through each point. Taking account of what we have shown:

$$|\mathcal{S}| \sim N^2$$

$$|\mathcal{L}| \sim |\mathcal{S}|\frac{N^2}{R^2}\frac{1}{R} \sim \frac{N^4}{R^3}$$

$$|I(\mathcal{S}, \mathcal{L})| \sim |\mathcal{S}|\frac{N^2}{R^2} \sim \frac{N^4}{R^2}$$

we can see that both sides of the Szemerédi-Trotter inequality are comparable:

$$|I(\mathcal{S}, \mathcal{L})| \sim \frac{N^4}{R^2} \sim (N^2)^{\frac{2}{3}}\left(\frac{N^4}{R^3}\right)^{\frac{2}{3}} \sim |\mathcal{S}|^{\frac{2}{3}}|\mathcal{L}|^{\frac{2}{3}}$$

add diagram (!)

5.4 Ham Sandwich Theorems

The above examples suggest that the topology of $\mathbb{R}^3$ plays a key role in this incidence problem. We shall now introduce the method of polynomial partitioning, which can be seen as the topological analogy to the vanishing lemma we used in the previous chapters.

Let $\mathbb{S}^n$ denote the unit n-sphere in $\mathbb{R}^{n+1}$.

Theorem 5.4.1 (Borsuk-Ulam). *A map ϕ is said to be antipodal if it obeys $\phi(-x) = -\phi(x)$ for all x in its domain. Suppose $\phi : \mathbb{S}^N \to \mathbb{R}^N$ is a continuous antipodal mapping. Then the image of ϕ contains 0.*

The proof of this result is long and beyond the scope of this manuscript. We refer interested readers to a beautiful geometric proof in chapter 2 of Matousek's book *Using the Borsuk-Ulam theorem.*[14]

Let us now define some useful notation going forward.

Definition 5.4.1 (Bisection of a Set). A function $f : \mathbb{R}^n \to \mathbb{R}$ is said to bisect an open set U with volume $\mathrm{Vol}(U) < \infty$ if:

$$\mathrm{Vol}\{x \in U \mid f(x) > 0\} = \mathrm{Vol}\{x \in U \mid f(x) < 0\} = \frac{1}{2}\mathrm{Vol}(U).$$

Analogously, a function f is said to bisect a finite set S if both:

$$|\{x \in S \mid f(x) > 0\}| \leq \frac{|S|}{2}$$

and

$$|\{x \in S \mid f(x) < 0\}| \leq \frac{|S|}{2}.$$

Theorem 5.4.2 (General Ham Sandwich Theorem). *Let V be a finite dimensional vector space of continuous functions $f : \mathbb{R}^n \to \mathbb{R}$ such that for any non-zero function f, $Z(f)$ has zero Lebesgue measure. Let $U_1, U_2, \ldots, U_N \subset \mathbb{R}^n$ be finite volume open sets with $N < \dim V$.*

Then there exists a non-zero function $f \in V$ that bisects each U_i.

Proof. Define the functions $\{\phi_i\}_{i=1}^N$, $\phi_i : V \backslash \{0\} \to \mathbb{R}$ by:

$$\phi_i(f) = \mathrm{Vol}(\{x \in U_i \mid f(x) > 0\}) - \mathrm{Vol}(\{x \in U_i \mid f(x) > 0\}).$$

Since $Z(f)$ has measure zero, it is easy to see that $\phi_i(f) = 0$ if and only if f bisects U_i. Notice also that $\phi_i(-f) = -\phi_i(f)$, hence ϕ_i is antipodal.

We now show each $\phi_i(F)$ is continuous. It is enough to show that if U is a finite volume open set, then the measure of $\{x \in U \mid f(x) > 0\}$ depends continuously on $f \in V \backslash \{0\}$.

Suppose $f_n \to f$ in V for some $f, f_n \in V \backslash \{0\}$. f_n converges to f in the topology of V, so it follows it must converge pointwise. Pick any $\varepsilon > 0$. By Egorov's theorem, we can find a subset $E \subset U$ so that $f_n \to f$ uniformly pointwise on $U \backslash E$ with $m(E) < \varepsilon$. By hypothesis, $m(Z(f)) = 0$ and $m(U) < \infty$. Since the Lebesgue measure is continuous we can choose δ such that $m\left(\{x \in U \mid |f(x)| < \delta\}\right) < \varepsilon$.

Now we choose n sufficiently large that $|f_n(x) - f(x)| < \delta$ on $U \backslash E$. Then we have by the triangle inequality:

$$\begin{aligned} |m\left(\{x \in U \mid f_n(x) > 0\}\right) &- m\left(\{x \in U \mid f(x) > 0\}\right)| \\ &\leq |m\left(\{x \in U \mid f_n(x) > 0\}\right)| + |m\left(\{x \in U \mid f(x) > 0\}\right)| \\ &< 2\varepsilon. \end{aligned}$$

Since ε was arbitrary each ϕ_i is continuous.

We now combine each ϕ_i into the map $\phi : V \backslash \{0\} \to \mathbb{R}^N$. Since $\dim V > N$, select a subspace $U < V$ such that $\dim U = N + 1$. Now choose an isomorphism of U with $\mathbb{R}^{N+1}$, and think of $\mathbb{S}^N$ as a subset of U. Now the map $\phi : \mathbb{S}^N \to \mathbb{R}^N$ is antipodal and continuous. By the Borsuk-Ulam theorem, there exists an $f \in \mathbb{S}^N \subset V \backslash \{0\}$ such that $\phi(f) = 0$. $\square$

Corollary 5.4.2.1 (Finite Ham Sandwich Theorem). *Let $S_1, \ldots, S_N$ be finite sets in $\mathbb{R}^n$ and let D be such that $N < \binom{D+n}{n}$. Then there exists a non-zero $P \in \mathrm{Poly}_D(\mathbb{R}^n)$ that bisects each S_i.*

Proof. Let us equip the space with the L^1 norm. For each $\delta > 0$, define $U_{i,\delta}$ to be the union of $\delta - balls$ centred at the points of S_i. By Theorem 5.4.2, we can find a non-zero P_δ with degree less than D that bisects each $U_{i,\delta}$. By rescaling we can assume $P_\delta \in \mathbb{S}^N \subset \mathrm{Poly}_D(\mathbb{R}^n) \backslash \{0\}$. Since $\mathbb{S}^N$ compact, we can find a sequence $\delta_m \to 0$ so that

P_{δ_m} converges to P in $\mathbb{S}^N$. Since the coefficients of P_{δ_m} converge to P, P_{δ_m} converges to P uniformly on compact sets.

We claim P bisects each S_i. By contradiction, suppose $P > 0$ on more than half the points of S_i, say on the points of S_i^+. Choosing ε sufficiently small, we can assume $P > 0$ on the ε-ball around each point of S_i^+. Further, we can choose ε such that each ε-ball is disjoint. Since P_{δ_m} converges uniformly, we can find m sufficiently large such that $P_{\delta_m} > 0$ on the ε-ball around each point of S_i^+. By making m large, we can also arrange that $\delta_m < \varepsilon$. Thus $P_{\delta_m} > 0$ on more than half the points of U_{i,δ_m}. $\qquad\square$

Theorem 5.4.3 (Polynomial Partitioning). *For any n there exists a constant $c(n)$ such that if S is a finite subset of $\mathbb{R}^n$ and D is any degree, then there exists a polynomial P of degree D such that $\mathbb{R}\backslash Z(P)$ is a disjoint union of $\lesssim D^n$ open sets O_i each containing $\lesssim_n |S|D^{-n}$ points.*

Proof. The main idea of this proof is the repeated application of the Finite Ham Sandwich Theorem. We begin by finding a polynomial P_1 of degree 1 that bisects S. This partitions $\mathbb{R}\backslash Z(P_1)$ into two disjoint open sets according to the sign of P_1, P_1^+ and P_1^-, each containing at most $|S|/2$ points. We then bisect both of these sets using another polynomial P_2. There are four sign conditions on P_1 and P_2, these being the four possible intersections of the sets $P_1^{\pm}$ and $P_2^{\pm}$, and the subset for each sign condition contains at most $|S|/4$ points of S. Continuing this process to define polynomials $P_3, P_4, \ldots$, where the polynomial P_j simultaneously bisects 2^{j-1} finite sets. By the Finite Ham Sandwich Theorem, each P_j can have a degree $\lesssim 2^{j/n}$. Repeating this procedure J times, and defining $P = \prod_{i=1}^J P_i$, $\mathbb{R}^n\backslash Z(P)$ is the disjoint union of 2^J open sets each containing $\leq |S|2^{-J}$ points of S. Now we choose D such that $\deg(P) < D$ which is equivalent to $\sum_{j=0}^J c(n)2^{j/n} \leq D$. But $\sum_{j=0}^J 2^{j/n}$ is a geometric series so we can find $\deg(P) < D$ for $D \leq c(n)2^{J/n}$. The number of points in each O_i is $\leq |S|2^{-J} \leq c(n)|S|D^{-n}$ $\qquad\square$

There is a crucial point to note about polynomial partitioning. The above theorem does not guarantee anything about the distribution of points between $Z(P)$ and its compliment. This is made most clear looking at the extremal examples. If all points line in the compliment of $Z(P)$ then we have an optimal eqidistribution of points, and can often use trivial bounds in a divide-and-conquer style argument. On the otherhand, in the case all points are contained in $Z(P)$ we have many points in an algebraic surface of controlled degree, so we can try and use tools from algebraic geometry. Generally, there will be some points in both $Z(P)$ and its compliment, which we need to deal with separately.

5.5 Proof of the Szemerédi-Trotter Theorem

We now can prove the Szemerédi-Trotter theorem using polynomial partitioning.

Proof of the Szemerédi-Trotter Theorem. Let $|\mathcal{S}| = S$ and $|\mathcal{L}| = L$. We need only consider the case $S^{\frac{1}{2}} \leq L \leq S^2$, as otherwise the proof follows immediately from the lemma above.

Let D be a degree to be chosen later. By Theorem 5.4.3, there exists a polynomial P of degree D such that $\mathbb{R}^2 \backslash Z(P)$ splits into D^2 components each having $\lesssim SD^{-2}$ points. Let $O_{i \in \Pi}$ denote these components and let $\mathcal{S}_i = \mathcal{S} \cap O_i$, $\mathcal{L}_i$ denote the lines that intersect the interior of each O_i respectively. We define the following pairs of complimentary sets:

$$\mathcal{S}_c = \{x \in \mathcal{S} \mid x \notin Z(p)\}$$
$$\mathcal{S}_z = \{x \in \mathcal{S} \mid x \in Z(p)\}$$
$$\mathcal{L}_c = \{\ell \in \mathcal{L} \mid \ell \not\subset Z(p)\}$$
$$\mathcal{L}_z = \{\ell \in \mathcal{L} \mid \ell \subset Z(p)\}$$

Note that $\mathcal{S} = \mathcal{S}_c \cup \mathcal{S}_z$, $\mathcal{L} = \mathcal{L}_c \cup \mathcal{L}_z$. We can now write our total line-point incidences as the following sum

$$I(\mathcal{S}, \mathcal{L}) = I(\mathcal{S}_c, \mathcal{L}) + I(\mathcal{S}_z, \mathcal{L}_z) + I(\mathcal{S}_z, \mathcal{L}_c).$$

If a line ℓ is not contained entirely in $Z(P)$ then it can intersect P at most D times, so each line intersects at most $D + 1$ cells. Hence $\sum_{i \in \Pi} L_i \leq (D + 1)L$. We begin by examining the $I(\mathcal{S}_c, \mathcal{L})$ term:

$$I(\mathcal{S}_c, \mathcal{L}) = \sum_{i \in \Pi} I(\mathcal{S}_i, \mathcal{L}_i)$$

Using our trivial bound in each cell:

$$\leq \sum_{i \in \Pi} \mathcal{S}_i^2 + \sum_{i \in \Pi} \mathcal{L}_i$$
$$\lesssim LD + SD^{-2} \sum_{i \in \Pi} S_i$$
$$\leq LD + S^2 D^{-2}$$

The number of lines in $\mathcal{L}_z$ is at most D. So we have by our trivial bounds:

$$I(\mathcal{S}_z, \mathcal{L}_z) \leq S + D^2.$$

Each line in $\mathcal{L}_c$ has at most D intersection points with $Z(P)$ so it has at most D incidences with $\mathcal{S}_z$. Hence:

$$I(\mathcal{S}_z, \mathcal{L}_c) \leq LD.$$

Together we have now

$$I(\mathcal{S}, \mathcal{L}) \lesssim LD + S^2 D^{-2} + S + D^2.$$

We optimise $LD + S^2 D^{-2}$ by choosing D such that both terms comparable and hence $D \sim S^{\frac{2}{3}} L^{-\frac{1}{3}}$. From our restriction $S^{\frac{1}{2}} \leq L \leq S^2$ we have $S^{\frac{2}{3}} L^{-\frac{1}{3}} \geq 1$ and $D^2 \sim S^{\frac{4}{3}} L^{-\frac{2}{3}} \leq S$, so we achieve

$$I(\mathcal{S}, \mathcal{L}) \lesssim (SL)^{2/3} + S.$$

Considering the regime where $L > S^2$ and applying the trivial bound yields the full Szemerédi-Trotter inequality:

$$I(\mathcal{S}, \mathcal{L}) \lesssim (SL)^{2/3} + S + L.$$

$\square$

There are two key things to note about the above proof. First, the key role that the topology of $\mathbb{R}^3$ plays. Topology is used in the proof of polynomial partitioning as it relies on the Borsuk Ulam theorem. It is a worthwhile heuristic to develop that polynomial partitioning may be useful for incidence problems where the best examples in a finite field (which is only equipped with the trivial topology) do not coincide with the best known examples over the reals. Secondly, the above proof illustrates the surprising power of polynomial partitioning. We are able to use very trivial bounds in each cell to achieve a asymptotically tight overall bound.

Chapter 6

Counting Circle Tangencies

Here we shall discuss the problem of counting the number of tangencies in a suitably non-degenerate collection of circles. We say two circles are tangent if their intersection contains a single point. The set of unordered pairs of circles in a collection $\mathcal{C}$ which are mutually tangent are called the tangencies of the collection and is denoted $\tau(\mathcal{C})$.

6.1 Trivial Bounds

Theorem 6.1.1 (Trivial Bound). *Let $\mathcal{C}$ be an arbitrary finite collection of circles. Then the number of tangencies $|\tau(\mathcal{C})|$ is bounded as follows:*

$$|\tau(\mathcal{C})| \leq |\mathcal{C}|^2.$$

Stated for arbitrary collections of circles the problem is not particularly interesting as the above bound turns out to be asymptotically tight.

Example 6.1. *Denote a circle centred at (x, y) with radius r as $\gamma(x, y, r)$. Consider the following collection of $N + 1$ circles:*

$$C_0 = \gamma\left(0, 0, 1\right)$$

$$C_1 = \gamma\left(\frac{1}{2}, 0, \frac{1}{2}\right)$$

$$C_2 = \gamma\left(\frac{3}{4}, 0, \frac{1}{4}\right)$$

$$\vdots$$

$$C_N = \gamma\left(1 - \frac{1}{2^N}, 0, \frac{1}{2^N}\right)$$

Each circle in our collection $\mathcal{C} = \{C_i \mid 0 \leq i \leq N\}$ is tangent to N other circles at the point (1,0). Hence $|\tau(\mathcal{C})| \sim N^2 \sim |\mathcal{C}|^2$. diagram

Instead we shall look at collections of circles that satisfy a single non-degeneracy condition. We consider collections of circles such that no three are mutually tangent at a

common point. The best known examples leverage our asymptotically tight Szemerédi-Trotter Theorem examples:

Example 6.2. *Let $\mathcal{S}$, $\mathcal{L}$ be collections of N points and N lines respectively such that the Szemerédi-Trotter Theorem's bound is asymptotically tight. In particular, they determine $\sim N^{4/3}$ point line incidences. Let C_1 denote the collection of unit circles centred at the points of $\mathcal{S}$. Let C_2 denote the collection of lines obtained by translating each line $\ell in \mathcal{L}$ one unit in the $\ell^\perp$ direction. If $(p, \ell) \in \mathcal{S} \times \mathcal{L}$ is a point-line incidence from our original collection, then the corresponding circle-line pair will be tangent. Performing an inversion transform about a point that does not lie in any of the circles or lines, our collection $\mathcal{C} = C_1 \cup C_2$ becomes a collection of $2N$ circles that determine $N^{4/3}$ tangencies.*

We shall now present a bound from a recent paper of Ellenberg-Solymosi-Zahl which uses the polynomial method.[15]

Theorem 6.1.2. *Given a finite collection $\mathcal{C}$ of N circles in the plane such that no three are tangent at a common point, then the number of tangencies $|\tau(\mathcal{C})|$ obeys:*

$$|\tau(\mathcal{C})| \lesssim N^{3/2}$$

It is not obvious how to apply the polynomial method in the current formulation. We will perform a lifting of our circles into algebraic curves $\mathbb{R}^3$, preserving tangencies between circles in the plane as intersections between their lifted curves in $\mathbb{R}^3$. This transforms the problem from a tangency problem to an incidence problem, and reduces the required degree of polynomial needed to interpolate. In general Lemma 2.3.1 tells us that if we are to interpolate a set of M points in $\mathbb{R}^d$ by a polynomial P, then the degree of P is $O(M^{1/d})$, so increasing the dimension of the problem yields a more controlled degree of the polynomial required.

6.2 Lifting of circles to $\mathbb{R}^3$

We shall now discuss this lifting in detail. Let $\gamma \subset \mathbb{R}^2$ be a circle in the plane of radius r_γ centred at (x_γ, y_γ). We define the lifting transform of the circle as:

$$\beta(\gamma) = \left\{ (x, y, z) \in \mathbb{R}^3 \mid (x - x_\gamma)^2 + (y - y_\gamma)^2 = r_\gamma^2, \ z = -\frac{x - x_\gamma}{y - y_\gamma} \right\}.$$

Clearly $\beta(\gamma)$ is an algebraic curve, so now we shall examine the correspondence between mutually tangent circles. Notice that z is defined as the reciprocal of the slope of the tangent line at the point $(x, y) \in \gamma$.

Lemma 6.2.1. *Let β be the transform defined as above. Then two circles $\gamma, \gamma' \subset \mathbb{R}^2$ are tangent if and only if $\beta(\gamma) \cap \beta(\gamma') \neq \emptyset$.*

Proof. If γ and γ' are tangent then there exists a point $(x, y) \in \gamma \cap \gamma'$ and at a point of tangency we have that the slopes of the tangent lines coincide for γ and γ'. Explicitly $\frac{x - x'_\gamma}{y - y'_\gamma} = \frac{x - x_\gamma}{y - y_\gamma} = z$. Hence $(x, y, z) \in \beta(\gamma) \cap \beta(\gamma')$ and thus $\beta(\gamma) \cap \beta(\gamma') \neq \emptyset$.

In the other direction, assume there exists some $(x, y, z) \in \beta(\gamma) \cap \beta(\gamma')$. Clearly $(x, y) \in \gamma \cap \gamma'$ and the slopes of the tangent lines at this point are equal, hence we conclude γ is tangent to γ'. $\qquad\square$

This one-to-one correspondence between tangencies in $\mathbb{R}^2$ and incidences in $\mathbb{R}^3$ is the key idea behind the proof. For those more algebraically inclined, what we are doing here is equivalent to looking at the tangent bundles of our circle in the projective space, hence we are preserving intersections.

We will need a Bezout-type result which bounds the number of intersections between curves with no common components.

Lemma 6.2.2 (Bezout's Theorem for Plane Algebraic Curves). *Let Z_1 and Z_2 be the zero sets of two polynomials over $\mathbb{R}[X, Y]$ and suppose $\deg Z_1 > \deg Z_2$. Then one of the following holds:*

$$|Z_1 \cap Z_2| \leq \deg Z_1 \deg Z_2$$
$$\text{or } Z_2 \subset Z_1$$

6.3   Ellenberg-Solymosi-Zahl's Proof of Theorem 6.1.2

In this proof, we will need to ensure our tangencies are sufficiently uniformly distributed among our circles. We can refine our collection such that this is the case by the following lemma.

Lemma 6.3.1 (Uniform Refinement). *Let $\mathcal{C}$ be a collection of N circles and suppose that $|\tau(\mathcal{C})| \gtrsim N^\alpha$. Then we can refine our collection to a collection $\mathcal{C}' \subset \mathcal{C}$ such that every circle in $\mathcal{C}'$ is mutually tangent to $\gtrsim N^{\alpha-1}$ other circles in $\mathcal{C}'$, $|\tau(\mathcal{C}')| \gtrsim N^\alpha$, and $|\mathcal{C}'| \gtrsim N^{\alpha/2}$.*

Proof. We proceed by a stopping time argument. Let $\tau(\mathcal{C})$ be the set of tangencies of the circles in $\mathcal{C}$. Let $\mathcal{C}_0 = \mathcal{C}$ and let c_1 be a fixed constant to be chosen later. If there exists a circle $\gamma \in \mathcal{C}_0$ such that $|\{\gamma' \mid (\gamma, \gamma') \in \tau(\mathcal{C}_0)\}| < c_1 N^{1/2}$ remove it from the collection and label the new refined collection as $\mathcal{C}_1$. From this refined collection, if there exists a circle γ such that $|\{\gamma' \mid (\gamma, \gamma') \in \tau(\mathcal{C}_1)\}| < c_1 N^{1/2}$ we remove it and label the remaining collection as $\mathcal{C}_2$.

After repeating this process M times until there are no more circles γ that satisfy $|\{\gamma' \mid (\gamma, \gamma') \in \tau(\mathcal{C}_M)\}| < c_1 N^{1/2}$, at each step removing a circle that contributes only a small number of tangencies, we attain a collection $\mathcal{C}_M$. We claim that $|\tau(\mathcal{C}_M)| \gtrsim N^\alpha$, and that $|\mathcal{C}_M| \gtrsim N^{\alpha/2}$.

For the first claim, observe that at each step i we are reducing $\tau(\mathcal{C}_i)$ by at most $c_1 N^{\alpha-1}$. Thus,

$$
\begin{aligned}
|\tau(\mathcal{C}_M)| &\geq |\tau(\mathcal{C}_0)| - M c_1 N^{\alpha-1} \\
&> c_0 N^\alpha - M c_1 N^{\alpha-1} \\
&> c_0 N^\alpha - c_1 N^\alpha \\
|\tau(\mathcal{C}_M)| &> \frac{c_0}{2} N^\alpha. \qquad \text{(by choosing } c_1 = c_0/2)
\end{aligned}
$$

We now provide a lower bound on the cardinality of our refined set $\mathcal{C}_M$. We have the trivial inequality $|\tau(\mathcal{C}_M)| \leq |\mathcal{C}_M|^2$. Combining this with the result above, we attain $|\mathcal{C}_M| \gtrsim |N|^{\alpha/2}$. $\qquad \square$

In a similar fashion to our arguments in the Kakeya problem we will be arguing by contradiction, however here we will be arguing that if the zero set of a polynomial contains too many of a certain first type of object, it must contain many more of a second kind of object.

We can now prove the main theorem.

Proof of Theorem 6.1.2. Given an arbitrary collection of circles $\mathcal{C}$ with $\gtrsim N^{3/2}$ tangencies, Lemma 6.3.1 with $\alpha = \frac{3}{2}$ allows us to reduce to a collection Γ where each circle is tangent to $\gtrsim N^{1/2}$ other circles using the previous lemma. After applying a small rotation, we can assume that the tangent line at each point of tangency does not point vertically in the y-direction. Let $\beta(\Gamma) = \{\beta(\gamma) : \gamma \in \Gamma\}$, where β is the lifting transform defined earlier, Recall from Lemma 6.2.1 that two circles γ_1 and γ_2 are tangent if and only if $\beta(\gamma_1) \cap \beta(\gamma_2) \neq \emptyset$.

Suppose $(x, y, z) \in \beta(\gamma_1) \cap \beta(\gamma_2)$ for some $\gamma_1 \neq \gamma_2$. Then

$$
(0, 0, 1) \in \text{span}\left(T_{(x,y,z)}\beta(\gamma_1), T_{(x,y,z)}\beta(\gamma_2)\right).
$$

In other words, at the intersection of $\beta(\gamma_1)$ and $\beta(\gamma_2)$ their tangent vectors span a vertical subspace of $\mathbb{R}^3$. We can establish this by examining a parameterisation of γ_1 and γ_2 in the neighbourhood of (x, y). Define $f_i(t)$, $i \in \{1, 2\}$ such that $(t+x, f_i(t))$ is a parameterisation of γ_i in the neighbourhood of (x, y) for all t in a small neighbourhood of 0. In particular $f_i(0) = y$. Taking the first derivative and evaluating at $t = 0$:

$$
\frac{df_i}{dt}(0) = -\frac{x - x_{\gamma_i}}{y - y_{\gamma_i}}
$$

Since γ_1 is tangent to γ_2 at (x, y) the slopes of their tangent lines coincide at that point so $\frac{df_1}{dt}(0) = \frac{df_2}{dt}(0)$. Now taking the second derivative and evaluating at $t = 0$:

$$
\frac{d^2 f_i}{dt^2}(0) = -\frac{(y - y_{\gamma_i})^2 + (x - x_{\gamma_i})^2}{(y - y_{\gamma_i})^3} = -\frac{r_{\gamma_i}^2}{(y - y_{\gamma_i})^3}
$$

Since γ_1 and γ_2 are distinct quadratic curves, $\frac{d^2 f_1}{dt^2}(0) \neq \frac{d^2 f_2}{dt^2}(0)$.

In the neighbourhood of (x, y, z), $\beta(\gamma_i)$ is parametrised by $\left(t, f_i(t), \frac{df_1}{dt}(t)\right)$ as the slope of the tangent to the circle is given by $\frac{df_1}{dt}(t)$. It follows that the tangent vector $\left(1, \frac{df_i}{dt}(0), \frac{d^2 f_i}{dt^2}(0)\right)$ generates the vertical space $T_{(x,y,z)}\beta(\gamma_i)$. Thus

$$(0, 0, 1) \in \text{span}\left(\left(1, \frac{df_1}{dt}(0), \frac{d^2 f_1}{dt^2}(0)\right) - \left(1, \frac{df_2}{dt}(0), \frac{d^2 f_2}{dt^2}(0)\right)\right)$$
$$\subset \text{span}\left(T_{(x,y,z)}\beta(\gamma_1), T_{(x,y,z)}\beta(\gamma_2)\right).$$

We will now interpolate all points of intersection with a minimal polynomial of suitably low degree, and show that if this contains too many intersections it must also contain the curves. Then due to the tangent vector spanning the z-axis we will be able to achieve a contradiction.

Let $P \in \mathbb{R}[x, y, z]$ be a non-zero polynomial of minimal degree that vanishes on all intersections between the curves in $\beta(\Gamma)$. This polynomial interpolates $\sim N^{3/2}$ points in $\mathbb{R}^3$, so by Lemma 2.3.1 the degree of P is $\lesssim \left(N^{3/2}\right)^{\frac{1}{3}} = N^{1/2}$.

Due to our refinement each $\gamma \in \Gamma$ is tangent to $\gtrsim N^{1/2}$ circles and each of these tangencies occur at a distinct point by our non-degeneracy condition. Hence we have that P vanishes at $\gtrsim N^{1/2}$ points on each curve in $\beta(\Gamma)$. By Bézout's theorem we have that P vanishes on all curves in $\beta(\Gamma)$ as:

$$\deg(P) \deg(\beta(\gamma)) \lesssim \#\{P \cap \beta(\gamma)\}$$

for suitable choice of constants.

By our result above, if (x, y, z) is a point where two curves from $\beta(\Gamma)$ intersect, then $\partial_z P(x, y, z) = 0$. Thus by the same Bezout argument $\partial_z P$ is also a polynomial which vanishes on all curves in $\beta(\Gamma)$.

Since P was a non-zero polynomial of minimal degree that vanishes on all the curves in $\beta(\Gamma)$, we must conclude $\partial_z P \equiv 0$. By the minimality if P we must have that $P(x, y, z) = Q(x, y)$ for some $Q \in \mathbb{R}[x, y]$ with degree $\lesssim N^{1/2}$. But this implies that there must be at least of the $N^{3/4}$ circles in our original collection must also be contained in $Z(Q)$. This is a contradiction, as Q has degree $\sim N^{1/2}$ whereas $\cup\gamma$ has degree $2N^{3/4}$. Thus $\beta(\Gamma)$ has $\lesssim N^{3/2}$ curve-curve incidences, so we can conclude that Γ has $\lesssim N^{3/2}$ tangencies. $\square$

6.4 A New Proof via Polynomial Partitioning for Varieties

In this section we shall present a new proof of Theorem 6.1.2. There are a few key differences between both methods of proof. In the following proof, it is no longer required to ensure the uniform distribution of tangent points among each circle in our collection, so we will not have to use Lemma 6.3.1. Secondly, we will partition $\mathbb{R}^3$ using a polynomial of controlled degree and leverage our trivial bound in each cell. Slight care will be needed to

deal with intersections with curves in the zero set. We begin by introducing the main tool for this proof, an extension of the polynomial partitioning theorem to algebraic varieties instead of just points (which are in some sense 0-dimensional varieties).

Lemma 6.4.1 (Polynomial Partitioning for Algebraic Varieties). *Suppose Γ is a set of k-dimensional varieties in $\mathbb{R}^n$. For any positive integer D there exists a non-zero polynomial P of degree at most D such that each connected component of $\mathbb{R}^n \backslash Z(P)$ intersects $\lesssim D^{k-n}|\Gamma|$ varieties $\gamma \in \Gamma$.*

We shall not prove this here. A proof of this can be found as Theorem 0.3 in the original paper of Guth presenting the result.[16]

add sketch of proof?

Proof of Theorem 6.1.2 by Polynomial Partitioning. Relabel $\mathcal{C}$ as Γ for convenience. Again, we perform the lifting transform β on each $\gamma \in \Gamma$. We have a collection of N 1-dimensional varieties in $\mathbb{R}^3$ upon which we shall use our polynomial partitioning lemma to find a polynomial P such that each cell of $\mathbb{R}^n \backslash Z(P)$ intersects $\lesssim ND^{-2}$ varieties. $\mathbb{R}^3 \backslash Z(P)$ partitions the space into D^3 cells. Let us label the interior of each of these cells as $\{\Omega_i \mid 0 \leq i \lesssim D^3\}$, and further label the set of varieties in Γ that pass through a given cell Ω_i as Γ_i.

We can now define the following complimentary sets based on whether the variety is contained entirely in $Z(P)$:

$$C_1 = \{\beta(\gamma) \mid \beta(\gamma \not\subset Z(P)\}$$
$$C_2 = \{\beta(\gamma) \mid \beta(\gamma) \subset Z(P)\}$$

Notice here that $\beta(\Gamma) = C_1 \cup C_2$. Recalling the correspondence between incidences between our curves and tangencies of the circles, we define the following incidence sets and hence have the following expression for $|\tau(\Gamma)|$:

$$I(C_1, C_1) = \{(\beta(\gamma), \beta(\gamma')) \mid \beta(\gamma), \beta(\gamma') \in C_1, \ \beta(\gamma) \cap \beta(\gamma') \neq \emptyset\}$$
$$I(C_1, C_2) = \{(\beta(\gamma), \beta(\gamma')) \mid \beta(\gamma) \in C_1, \beta(\gamma') \in C_2, \ \beta(\gamma) \cap \beta(\gamma') \neq \emptyset\}$$
$$I(C_2, C_2) = \{(\beta(\gamma), \beta(\gamma')) \mid \beta(\gamma), \beta(\gamma') \in C_2, \ \beta(\gamma) \cap \beta(\gamma') \neq \emptyset\}$$
$$|\tau(\Gamma)| = |I(C_1, C_1)| + |I(C_1, C_2)| + |I(C_2, C_2)|$$

We now proceed by calculating the cardinality each of these sets. The case for $I(C_1, C_1)$ will use our trivial bound and bezout's theorem. Calculating for $I(C_1, C_2)$ is again straightforward by Bezout. The interesting case here is $I(C_2, C_2)$, where we will be forced to argue via a recursive style of argument.

We begin with the intersections that occur between varieties not entirely contained in the zero set:

$$I(C_1, C_1) = \{(\beta(\gamma), \beta(\gamma')) \mid \beta(\gamma), \beta(\gamma') \in C_1, \ \beta(\gamma) \cap \beta(\gamma') \in \mathbb{R}^3 \backslash Z(P)\}$$
$$\cup \{(\beta(\gamma), \beta(\gamma')) \mid \beta(\gamma), \beta(\gamma') \in C_1, \ \beta(\gamma) \cap \beta(\gamma') \in Z(P)\}$$

Hence we have:

$$|I(C_1, C_1)| = \sum_{\beta(\gamma), \beta(\gamma') \in C_1} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in \mathbb{R}^3 \backslash Z(P)] + \sum_{\beta(\gamma), \beta(\gamma') \in C_1} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in Z(P)]$$

$$= \sum_{\beta(\gamma), \beta(\gamma') \in C_1} \sum_i \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in \Omega_i] + \sum_{\beta(\gamma), \beta(\gamma') \in C_1} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in Z(P)]$$

$$= \sum_i \sum_{\beta(\gamma), \beta(\gamma') \in \Gamma_i} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in \Omega_i] + \sum_{\beta(\gamma), \beta(\gamma') \in C_1} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in Z(P)]$$

Using our trivial bound and the fact that there are $\lesssim ND^{-2}$ varieties intesecting a given cell we attain:

$$\lesssim \sum_i \left(\frac{N}{D^2}\right)^2 + \sum_{\beta(\gamma), \beta(\gamma') \in C_1} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in Z(P)]$$

$$= D^3 N^2 D^{-4} + \sum_{\beta(\gamma) \in C_1} \sum_{\beta(\gamma') \in C_1} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in Z(P)]$$

Fixing $\beta(\gamma)$ we see that the latter sum must be $\lesssim D$ by Bezout's lemma in conjunction with our non-degeneracy condition.

$$\lesssim N^2 D^{-1} + \sum_{\beta(\gamma) \in C_1} D$$

$$= N^2 D^{-1} + ND.$$

The argument for $I(C_1, C_2)$ is the similar to the above calculation, however there are now no intersections happening inside the cells by the definition of C_2. Fix $\beta(\gamma)$ in:

$$|I(C_1, C_2)| = \sum_{\beta(\gamma) \in C_1} \sum_{\beta(\gamma') \in C_2} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \in Z(P)]$$

and again the second sum is $\lesssim D$ due to Bezout and our non-degeneracy condition. Hence,

$$\lesssim \sum_{\beta(\gamma) \in C_1} D = ND.$$

Finally we need to handle the incidences between curves contained entirely in the zero-set, $I(C_2, C_2)$.

$$|I(C_2, C_2)| = \sum_{\beta(\gamma), \beta(\gamma') \in C_2} \mathbb{1}[\beta(\gamma) \cap \beta(\gamma') \neq \emptyset]$$

Suppose P was the minimal degree polynomial that partitioned the space as per Lemma 6.4.1. Then we can write C_2 as the union of disjoint sets:

$$C_2 = \{\beta(\gamma) \subset Z(\partial_z P)\} \cup \{\beta(\gamma) \subset \mathbb{R}^3 \backslash Z(\partial_z P)\}$$

If all the curves are in $\{\beta(\gamma) \subset Z(\partial_z P)\}$, then we use the exact same technique as in Ellenberg-Solymosi-Zahl's Proof to control the size of this contribution by $\lesssim ND$.

In the other case, there are at most a constant number of curves in $\{\beta(\gamma) \subset \mathbb{R}^3 \backslash Z(\partial_z P)\}$, each of which can intersect $Z(\partial_z P)$ at most D times. We recursively iterate this process now, examining now the set:

$$\{\beta(\gamma) \subset Z(\partial_z P)\} = \{\beta(\gamma) \subset Z(\partial_z^2 P)\} \cup \{\beta(\gamma) \subset \mathbb{R}^3 \backslash Z(\partial_z^2 P)\}$$

Again, if all the curves are in $Z(\partial_z^2 P)$ the process stops and we proceed by the argument in Ellenberg-Solymosi-Zahl's Proof. In the other case, $\mathbb{R}^3 \backslash Z(\partial_z^2 P)$ contains a constant number of curves which intersect $Z(\partial_z^2 P)$ at most D times each. We can iterate this process at most D times, so our sum becomes:

$$|I(C_2, C_2)| = ND + \sum_{i=1}^{D} D = ND + D^2$$

Adding together our bounds we achieve:

$$|I(C_1, C_1)| + |I(C_1, C_2)| + |I(C_2, C_2)| \lesssim N^2 D^{-1} + ND + ND + ND + D^2$$
$$\sim N^2 D^{-1} + ND + D^2.$$

We now optimise D by setting $N^2 D^{-1} \sim ND$ which gives $D \sim N^{1/2}$ and hence our bounds become:

$$|I(C_1, C_1)| + |I(C_1, C_2)| + |I(C_2, C_2)| \lesssim N^2 N^{-1/2} + NN^{1/2} + (N^{1/2})^2 \sim N^{3/2}.$$

$\square$

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