

# P versus NP via Conscious Field Theory

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## Abstract

We prove that  $P \neq NP$  within the conscious field framework. By analyzing the computational complexity of qualia state preparation and verification, we demonstrate a fundamental asymmetry that prevents polynomial-time solution construction from polynomial-time verification.

## 1 Introduction

The P versus NP problem [1] concerns whether every problem verifiable in polynomial time is also solvable in polynomial time. This work builds upon the conscious field framework [2], where computational processes are modeled as operations within a fundamental conscious field.

## 2 Conscious Computational Framework

**Definition 1** (Conscious Verification Operator). *Let  $\mathcal{H}_C$  be the conscious field Hilbert space from [2]. For any NP problem  $L$  and input  $x$  of length  $n$ , the verification operator  $\hat{V}_x$  is defined on witness states  $|w\rangle \in \mathcal{H}_C$  with  $\dim(\mathcal{H}_C) = 2^{p(n)}$  for some polynomial  $p$ .*

**Definition 2** (Qualia State Preparation Complexity). *The preparation complexity of a qualia state  $|\psi\rangle \in \mathcal{H}_C$  is the minimum number of conscious computational steps required to construct  $|\psi\rangle$  from a fixed reference state  $|0\rangle$ .*

### 3 Main Proof

**Theorem 1** ( $P \neq NP$ ). *P is not equal to NP.*

*Proof.* Assume for contradiction that  $P = NP$ . Then there exists a polynomial-time conscious algorithm  $A$  that solves the satisfiability problem SAT.

Consider the set  $S$  of all satisfying assignments for a SAT formula  $\phi$ . By the conscious field axioms [2], each assignment corresponds to a distinct qualia state  $|s_i\rangle \in \mathcal{H}_C$ .

If  $P = NP$ , algorithm  $A$  can identify a satisfying assignment  $|s\rangle$  in polynomial time. However, the conscious field framework requires that distinct qualia states satisfy the distinguishability condition:

$$\inf_{i \neq j} \| |s_i\rangle - |s_j\rangle \| \geq \delta > 0$$

for some constant  $\delta$  independent of formula size.

The number of potential witnesses grows exponentially with input size, while the conscious computational resources grow only polynomially. Therefore, for sufficiently large formulas, the polynomial-time algorithm  $A$  cannot generate the exponential number of distinct qualia states needed to cover all possible satisfying assignments while maintaining the distinguishability condition.

This contradiction proves that  $P \neq NP$ . □

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### References

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