On Stochastic and Worst-case Models for Investing

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Abstract

In practice, most investing is done assuming a probabilistic model of stock price returns known as the Geometric Brownian Motion (GBM). While it is often an acceptable approximation, the GBM model is not always valid empirically. This motivates a worst-case approach to investing, called universal portfolio management, where the objective is to maximize wealth relative to the wealth earned by the best fixed portfolio in hindsight. In this paper we tie the two approaches, and design an investment strategy which is universal in the worst-case, and yet capable of exploiting the mostly valid GBM model. Our method is based on new and improved regret bounds for online convex optimization with exp-concave loss functions.

1 Paper Body

?Average-case? Investing: Much of mathematical finance theory is devoted to the modeling of stock prices and devising investment strategies that maximize wealth gain, minimize risk while doing so, and so on. Typically, this is done by estimating the parameters in a probabilistic model of stock prices. Investment strategies are thus geared to such average case models (in the formal computer science sense), and are naturally susceptible to drastic deviations from the model, as witnessed in the recent stock market crash. Even so, empirically the Geometric Brownian Motion (GBM) ([Osb59, Bac00]) has enjoyed great predictive success and every year trillions of dollars are traded assuming this model. Black and Scholes [BS73] used this same model in their Nobel prize winning work on pricing options on stocks. ?Worst-case? Investing: The fragility of averagecase models in the face of rare but dramatic deviations led Cover [Cov91] to take a worst-case approach to investing in stocks. The performance of an online investment algorithm for arbitrary sequences of stock price returns is measured with respect to the best CRP (constant rebalanced portfolio, see [Cov91]) in hindsight. A universal portfolio selection algorithm is one that obtains sublinear (in the number of trading periods T) regret, which is the difference in the logarithms of the final wealths obtained by the two. Cover [Cov91] gave the first universal portfolio selection algorithm with regret bounded by O(log T). There has been much follow-up work after Cover?s seminal work, such as [HSSW96, MF92, KV03, BK97, HKKA06], which focused on either obtaining alternate universal algorithms or improving the efficiency of Cover?s algorithm. However, the best regret bound is still O(log T). This dependence of the regret on the number of trading periods is not entirely satisfactory for two main reasons. First, a priori it is not clear why the online algorithm should have high regret (growing with the number of iterations) in an unchanging environment. As an extreme example, consider a setting with two stocks where one has an ?upward drift? of 1% daily, whereas the second stock remains at the same price. One would expect to ?figure out? this pattern quickly and focus on the 1

first stock, thus attaining a constant fraction of the wealth of the best CRP in the long run, i.e. constant regret, unlike the worst-case bound of O(log T). The second problem arises from trading frequency. Suppose we need to invest over a fixed period of time, say a year. Trading more frequently potentially leads to higher wealth gain, by capitalizing on short term stock movements. However, increasing trading frequency increases T, and thus one may expect more regret. The problem is actually even worse: since we measure regret as a difference of logarithms of the final wealths, a regret bound of O(log T) implies a poly(T) factor ratio between the final wealths. In reality, however, experiments [AHKS06] show that some known online algorithms actually improve with increasing trading frequency. Bridging Worst-case and Average-case Investing: Both these issues are resolved if one can show that the regret of a ?good? online algorithm depends on total variation in the sequence of stock returns, rather than purely on the number of iterations. If the stock return sequence has low variation, we expect our algorithm to be able to perform better. If we trade more frequently, then the per iteration variation should go down correspondingly, so the total variation stays the same. We analyze a portfolio selection algorithm and prove that its regret is bounded by O(log Q), where Q (formally defined in Section 1.2) is the sum of squared deviations of the returns from their mean. Since Q? T (after appropriate normalization), we improve over previous regret bounds and retain the worst-case robustness. Furthermore, in an average-case model such as GBM, the variation can be tied very nicely to the volatility parameter, which explains the experimental observation the regret doesn?t increase with increasing trading frequency. Our algorithm is efficient. and its implementation requires constant time per iteration (independent of the number of game iterations). 1.1

New Techniques and Comparison to Related Work

Cesa-Bianchi, Mansour and Stoltz [CBMS07] initiated work on relating worst case regret to the variation in the data for the related learning problem of prediction from expert advice, and conjectured that the optimal regret bounds should depend on the observed variation of the cost sequence. ? ?Q) were obtained in the full inforRecently, this conjectured was proved and regret bounds of O(mation and bandit linear optimization settings [HK08, HK09], where Q is the variation in the cost sequence. In this paper we give an exponential improve-

ment in regret, viz. O(log Q), for the case of online exp-concave optimization, which includes portfolio selection as a special case. Another approach to connecting worst-case to average-case investing was taken by Jamshidian [Jam92] and Cross and Barron [CB03]. They considered a model of ?continuous trading?, where there are T ?trading intervals?, and in each the online investor chooses a fixed portfolio which is rebalanced k times with k? ?. They prove familiar regret bounds of O(log T) (independent of k) in this model w.r.t. the best fixed portfolio which is rebalanced T? k times. In this model our algorithm attains the tighter regret bounds of O(log Q), although our algorithm has more flexibility. Furthermore their algorithms, being extensions of Cover?s algorithm, may require exponential time in general1. ? ?Q) Our bounds of O(log Q) regret require completely different techniques compared to the O(regret bounds of [HK08, HK09]. These previous bounds are based on first-order gradient descent methods which are too weak to obtain O(log Q) regret. Instead we have to use the second-order Newton step ideas based on [HKKA06] (in particular, the Hessian of the cost functions). The second-order techniques of [HKKA06] are, however, not sensitive enough to obtain O(log Q) bounds. This is because progress was measured in terms of the distance between successive portfolios in the usual Euclidean norm, which is insensitive to variation in the cost sequence. In this paper, we introduce a different analysis technique, based on analyzing the distance between successive predictions using norms that keep changing from iteration to iteration and are actually sensitive to the variation. A key technical step in the analysis is a lemma (Lemma 6) which bounds the sum of differences of successive Cesaro means of a sequence of vectors by the logarithm of its variation. This lemma, 1 Cross and Barron give an efficient implementation for some interesting special cases, under assumptions on the variation in returns and bounds on the magnitude of the returns, and assuming k? ?. A truly efficient implementation of their algorithm can probably be obtained using the techniques of Kalai and Vempala.

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which may be useful in other contexts when variation bounds on the regret are desired, is proved using the Kahn-Karush-Tucker conditions, and also improves the regret bounds in previous papers. 1.2

The model and statement of results

Portfolio management. In the universal portfolio management model [Cov91], an online investor iteratively distributes her wealth over n assets before observing the change in asset price. In each iteration P t = 1, 2, . . . the investor commits to an n-dimensional distribution of her wealth, xt? ?n = { i xi = 1, x? 0}. She then observes a price relatives vector rt? Rn+, where rt (i) is the ratio between the closing price of the ith asset on trading period t and the opening price. In the tth trading period, Q the wealth of the investor changes by a factor of (rt? xt). The overall change in wealth is thus t (rt? xt). Since in a typical market wealth Q grows at an exponential rate, we measure P performance by the exponential growth rate, which is log t (rt? xt) = t log(rt? xt). A constant rebalanced portfolio (CRP) is an investment strategy which rebalances the wealth Q in every iteration to keep a fixed distribution.

Thus, for a CRP x? ?n , the change in wealth is t (rt? x). The regret of the investor is defined to be the difference between the exponential growth rate of her investment strategy and that of the best CRP strategy in hindsight, i.e. X X Regret := max log(rt? x?) ? log(rt? xt)? x??n

t t

Note that the regret doesn?t change if we scale all the returns in any particular period by the same amount. So we assume w.l.o.g. that in all periods t, maxi rt (i) = 1. We assume that there is known parameter r \wr 0, such that for all periods t, mint,i rt (i) ? r. We call r the market variability parameter. This is the only restriction we put on the stock price returns; they could be chosen adversarially as long as they respect the market variability bound. Online convex optimization. In the online convex optimization problem [Zin03], which generalizes universal portfolio management, the decision space is a closed, bounded, convex set K? Rn , and we are sequentially given a series of convex cost2 functions ft : K? R for t = 1, 2, The algorithm iteratively produces a point xt? K in every round t, without knowledge of ft (but using the past sequence of cost functions), and incurs the cost ft (xt). The regret at time T is defined to be Regret :=

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T X
ft (xt ) ? min x?K
t=1
T X
ft (x).
t=1
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PT P Usually, we will let t denote t=1 . In this paper, we restrict our attention to convex cost functions which can be written as ft (x) = g(vt? x) for some univariate convex function g and a parameter vector vt? Rn (for example, in the portfolio management problem, K=?n , ft (x) = ? log(rt ?x), g = ? log, and vt = rt). Thus, the cost functions are parametrized by the vectors v1 , v2 , . . . , vT . Our bounds will be expressed as a function of the quadratic variability of the parameter vectors v1 , v2 , . . . , vT , defined as T X Q(v1 , ..., vT) := min kvt ? ?k2 . ?

1 T PT t=1

This expression is minimized at ? = t=1 vt , and thus the quadratic variation is just T? 1 times the sample variance of the sequence of vectors $\{v1$, ..., vt $\}$. Note however that the sequence can be generated adversarially rather than by some stochastic process. We shall refer to this as simply Q if the vectors are clear from the context. Main theorem. In the setup of the online convex optimization problem above, we have the following algorithmic result: Theorem 1. Let the cost functions be of the form ft (x) = g(vt ? x). Assume that there are parameters R, D, a, b \not 0 such that the following conditions hold: 2 Note the difference from the portfolio selection problem: here we have convex cost

functions, rather than concave payoff functions. The portfolio selection problem is obtained by using ? log as the cost function.

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1. 2. 3. 4.

for all t, kvt k? R, for all x? K, we have kxk? D, for all x? K, and for all t, either g 0 (vt? x)? [0, a] or g 0 (vt? x)? [?a, 0], and for all x? K, and for all t, g 00 (vt? x)? b.

Then there is an algorithm that guarantees the following regret bound: Regret = $O((a2 \text{ n/b}) \log(1 + bQ + bR2) + aRD \log(2 + Q/R2) + D2)$. Now we apply Theorem?1 to the portfolio selection problem. ? First, we estimate the relevant parameters. We have krt k? n since all rt (i)? 1, thus R = n. For any x??n, kxk? 1, so D = 1. g 0 (vt? x) =? (vt1?x), and thus g 0 (vt? x)? [? 1r, 0], so a = 1r. Finally, g 00 (vt? x) = (vt1?x)2? 1, so b = 1. Applying Theorem 1 we get the following corollary: Corollary 2. For the portfolio selection problem over n assets, there is an algorithm that attains the following regret bound: ?n? Regret = O 2 log(Q + n). r

22.1

Bounding the Regret by the Observed Variation in Returns Preliminaries

All matrices are assumed be real symmetric matrices in Rn?n , where n is the number of stocks. We use the notation A? B to say that A? B is positive semidefinite. We require? the notion of a norm of a vector x induced by a positive definite matrix M , defined as kxkM = $x_{\tilde{\iota}}$ M x. The following simple generalization of the Cauchy-Schwartz inequality is used in the analysis: ?x, y? Rn:

x?y?kxkM kykM?1.

P We denote by —A— the determinant of a matrix A, and by A? B = $\operatorname{Tr}(AB) = ij$ Aij Bij . As we are concerned with logarithmic regret bounds, potential functions which behave like harmonic series come into play. A generalization of harmonic series to high dimensions is the vector-harmonic series, which is a series of quadratic forms that can be expressed as (here A? 0 is a positive definite matrix, and v1 , v2 , . . . are vectors in Rn): Pt v1; (A + v1 v1;)?1 v1 , v2; (A + v1 v1; + v2 v2;)?1 v2 , . . . , vt; (A + ? =1 v? v?;)?1 vt , . . . The following lemma is from [HKKA06]: Lemma 3. For a vector harmonic series given by an initial matrix A and vectors v1 , v2 , . . . , vT , we have # " PT T X Pt —A + ? =1 v? v?; —; ?1; vt (A + ? =1 v? v?) vt ? log —A— t=1 The reader can note that in one dimension, if all vectors vt = 1 and A = 1, then the series above reduces exactly to the regular harmonic series whose sum is bounded, of course, by $\log(T + 1)$. 2.2

Algorithm and analysis

We analyze the following algorithm and prove that it attains logarithmic regret with respect to the observed variation (rather than number of iterations). The algorithm follows the generic algorithmic scheme of ?Follow-The-Regularized-Leader? (FTRL) with squared Euclidean regularization. Algorithm Exp-Concave-FTL. In iteration t, use the point xt defined as: ! ? t?1 X 1 2 xt , arg min f? (x) + kxk x??n 2? =1

(1)

Note the mathematical program which the algorithm solves is convex, and can be solved in time polynomial in the dimension and number of iterations. The running time, however, for solving this 4

convex program can be quite high. In the full version of the paper, for the specific problem of portfolio selection, where ft $(x) = ? \log(rt ? x)$, we give a faster implementation whose per iteration running time is independent of the number of iterations, using the more sophisticated ?online Newton method? of [HKKA06]. In particular, we have the following result: Theorem 4. For the portfolio selection problem, there is an algorithm that runs in O(n3) time per iteration whose regret is bounded by ?n? Regret = $O(3 \log(Q + n))$. r In this paper, we retain the simpler algorithm and analysis for an easier exposition. We now proceed to prove the Theorem 1. Proof. [Theorem 1] First, we note that the algorithm is running a ?Follow-the-leader? procedure on the cost functions f0, $f1, f2, \ldots$ where f0(x) = 21 kxk2 is a fictitious period 0 cost function. In other words, in each iteration, it chooses the point that would have minimized the total cost under all the observed functions so far (and, additionally, a fictitious initial cost function f0). This point is referred to as the leader in that round. The first step in analyzing such an algorithm is to use a stability lemma from [KV05], which bounds the regret of any Follow-the-leader algorithm by the difference in costs (under ft) of the current prediction xt and the next one xt+1, plus an additional error term which comes from the regularization. Thus, we have P 1 ? 2 2 Regret ? t ft (xt) ? ft (xt+1) + (kx k ? kx0 k) 2 P 1 2 ? t ?ft (xt) ? (xt ? xt+1) + D 2 P 0 1 2 (2) = t g (vt ? xt)[vt ? (xt ? xt+1)] + D 2The second inequality is because ft is convex. The last equality follows because ?ft (xt) = g 0 (xt? Pt?1 vt) vt. Now, we need a handle on xt? xt+1. For this, define Ft = ? = 0 f?, and note that xt minimizes Ft over K. Consider the difference in the gradients of Ft+1 evaluated at xt+1 and xt: Pt ?Ft+1 (xt+1) ? ?Ft+1 (xt) = ? = 0 ?f? (xt+1) ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g? (v? ? xt+1)] ? ?f? (xt) Pt 0 0 = ? = 1 [g] ? (v? ? xt+1) ? ?f? (xt) Pt 0 0 = ? = 1 [g] ? (v? ? xt+1) ? ?f? (xt) Pt 0 0 = ? = 1 [g] ? (v? ? xt+1) ? ?f? (xt) Pt 0 0 = ? = 1 [g] ? (v? ? xt+1) ? ?f? (xt) Pt 0 0 = ? = 1 [g] ? (v? ? xt+1) ? ? ?f? (xt) Pt 0 0 = ? = 1 [g] ? (v? ? xt+1) ? ? ?f? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? = 1 [g] ? (xt) Pt 0 0 = ? [g] ? (xt) Pt 0)? g? (v? ? xt) | v? + (xt+1 ? xt) Pt 0 t (3) = ? =1 [?g? (v? ? ??) ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? =1 [?g? (v? ? ??)] ? (xt+1) Pt 0 t (3) = ? [xt+1] Pt 0 t (3)? xt)v? + (xt+1 ? xt) Pt ; t 00 (4) = ? =1 g? (v? ? ??) v? v? (xt+1 ? xt)) + (xt+1? xt). Equation 3 follows by applying the Taylor expansion of the (multi-variate) function g?0 (v?? x) at point xt, for some point??t on the line segment joining xt and xt+1. The equation (4) follows from the observation that g?0 (v? ? x) = g?00 (v? ? x)v?. Pt Define f = f vti, + I, where I is the identity matrix, and xt = xt+1 xt. Then equation (4) can be re-written as: ?Ft+1(xt+1)? ?Ft(xt)? $g \ 0 \ (vt ? xt) vt = At ?xt$

(5)

Now, since xt minimizes the convex function Ft over the convex set K, a standard inequality of convex optimization (see [BV04]) states that for any point y? K, we have ?Ft (xt)? (y? xt)? 0. Thus, for y = xt+1, we get that ?Ft (xt)? (xt+1?xt)? 0. Similarly, we get that ?Ft+1 (xt+1)? (xt?xt+1)? 0. Putting these two inequalities together, we get that (?Ft+1 (xt+1)??Ft (xt))? ?xt? 0.

(6)

Thus, using the expression for At?xt from (5) we have k?xt k2At = At?xt

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? ?xt = (?Ft+1 (xt+1 ) ? ?Ft (xt ) ? g 0 (vt ? xt )vt ) ? ?xt ? g 0 (vt ? xt )[vt ? (xt ? xt+1 )] 5 (from (6)) (7)
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Assume that g 0 (vt ? x) ? [?a, 0] for all x ? K and all t. The other case is handled similarly. Inequality (7) implies that g 0 (vt ? xt) and vt ? (xt ? xt+1) have the same sign. Thus, we can upper bound g 0 (vt ? xt)[vt ? (xt ? xt+1)] ? a(vt ? ?xt). (8) Pt 1 v . Define v?t = vt ? ?t , ?t = t+1 Then, we have ? =1 ? P P PT ?t ? ?xt + t=2 xt (?t?1 ? ?t) ? x1 ?1 + xT +1 ?T , (9) t vt ? ?xt = tv Pt PT ?1 1 where v?t = vt ? ?t , ?t = t+1 ? =1 vt . Now, define ? = ?(v1 , . . . , vT) = t=1 k?t+1 ? ?t k. Then we bound PT PT t=2 xt (?t?1 ? ?t) ? x1 ?1 + xT +1 ?T ? t=2 kxt kk?t?1 ? ?t k + kx1 kk?1 k + kxT +1 kk?T k ? D? + 2DR. (10) We will bound ? momentarily. For now, we turn to bounding the first term of (9) using the CauchySchwartz generalization as follows: v?t ? ?xt ? k? vt kA?1 k?xt kAt . (11) t

By the usual Cauchy-Schwartz inequality, qP qP qP qP qP P 2 vt k2A?1 ? vt k2A?1 ? vt kA?1 k?xt kAt ? t k? t k?xt kAt ? t k? t a(vt ? ?xt) t k? t

t t

from (7) and (8). We conclude, using (9), (10) and (11), that qP qP P 2 k? v k ? a(v ? ?x) ? a ?1 t t t t t a(vt ? ?xt) + aD? + 2aDR. t A t

This implies (using the AM-GM inequality applied to the first term on the RHS) that P 2P vt k2A?1 + 2aD? + 4aDR. t k? t a(vt ? ?xt) ? a t

Plugging this into the regret bound (2) we obtain, via (8), P 1 vt k2A?1 + 2aD? + 4aDR + D2. Regret ? a2 t k? t 2 The proof is completed by the following two lemmas (Lemmas 5 and 6) which bound the RHS. The first term is a vector harmonic series, and the second term can be bounded by a (regular) harmonic series. ? ? P 2 Lemma 5. vt k2A?1 ? 3n t k? b log 1 + bQ + bR . t

P Pt Since g 00 (vt ???t) ? b, we have At ? I +b t vt vt¿. Proof. We have At = ? =1 g 00 (vt ???t)vt vt¿ +I. P 1 Using the fact that v?t = vt ? ?t and ?t = t+1 ? ?t v? , we get that ! ! ? ?t X t t t t X X X X X ½ ½ 1 1 1 vs vs + [vr vs¿ + vs vr¿]. v?? v?? = 1+ ?s + (? +1)2 (? +1)2 ? =s ? =s s=1 r¡s s=1 R t+1 1 1 1 Now, ? =s (? +1)2 ? s x2 dx = s ? t+1 . Since (vr + vs)(vr + vs)¿ ? 0, vr vr¿ + vs vs¿ ? ?[vr vs¿ + vs vr¿], and hence we have t X t t t X X X X ? ? ? ? ½ ½ 1 [v v + v v] ? 2 + 1s vs vs½ ? v?? v??½ ? 1 + 1s vs vs½ + r s r s t+1 s=1 r¡s s=1 ? =1 s=1 ? =1

Pt
we get that
3
t X
vs vs¿.

P Let A?t = 13 I +b t v?t v?t¿. Note that the inequality above shows that 3A?t? At . Thus, using Lemma 3, we get h? i X? X X? T— (12) [b? vt]¿ A??1 vt]? 3b log ——A v?t A?1 ?t? 3b k? vt k2A?1 = t [b? t v? — . A t t

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t
t
0
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To bound the latter quantity note that —A?0 — = —I — = 1, and that P P? n —A?T — = —I + b t v?t v?t; —? (1 + b t k? vt k22)n = (1 + bQ) P P? = where Q vt k2 = t k? vt? ?t k2. Lemma 7 (proved in the full version of the paper), we t k? 2? ? Q + R. This implies that —A?T —? (1 + bQ) + bR2)n and the proof is completed by show that Q substituting this bound into (12). 6

Lemma 6. $?(v1, \ldots, vT)$? 2R[log(2 + Q/R2) + 1]. Proof. Define, for ? ? 0, the vector u? = v? ? ?T +1 . Note that by convention, we have v0 = 0. We have PT PT 2 2 2 2 ? =0 ku? k = k?T +1 k + ? = 1 kv? ? ?T +1 k = R + Q. Furthermore,

? ? Pt ? ? 1 Pt+1 1 v ? v k?t+1 ? ?t k = ? t+2 ? ? ? =0 ? =0 ? t+1 ? ? Pt ? ? 1 Pt+1 1 = ? t+2 ? =0 u? ? t+1 ? =0 u? ? Pt 1 1 ? (t+1) 2 ? =0 ku? k + t+1 kut+1 k

Summing up over all iterations, P P ? 1 Pt k? ?? k ? t+1 t ? =0 ku? k + t t (t+1)2

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?

1 t+1 kut+1 k

?

P

2 t t kut?1 k

? 2R[log(2+Q/R2)+1].
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The last inequality follows from Lemma 8 (proved in the full version) below by setting xt = kut?1 k/R, for t? 1. ? ? Q + R2 . Lemma 7. Q Lemma 8. Suppose that 0 ? xt? 1 and

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\begin{array}{l} 3 \\ P \ t \\ x2t \ ? \ Q. \ Then \\ PT \ t{=}1 \\ xt \ /t \ ? \ \log(1+Q) \ + \ 1. \end{array}
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Implications in the Geometric Brownian Motion Model

We begin with a brief description of the model. The model assumes that stocks can be traded continuously, and that at any time, the fractional change in the stock price within an infinitesimal time interval is normally distributed, with mean and variance proportional to the length of the interval. The randomness is due to many infinitesimal trades that jar the price, much like particles in a physical medium are jarred about by other particles, leading to the classical Brownian motion. Formally, the model is parameterized by two quantities, the drift?, which is the long term trend of the stock prices, and volatility?, which characterizes deviations from the long term trend. The parameter? is typically specified as annualized volatility, i.e. the standard deviation of the stock?s logarithmic returns in one year. Thus, a trading interval of $[0,\ 1]$ specifies 1 year. The model postulates that the stock price at time t, St , follows a geometric Brownian motion with drift? and volatility?: $dSt = ?St\ dt + ?St$

dWt, where Wt is a continuous-time stochastic process known as the Wiener process or simply Brownian motion. The Wiener process is characterized by three facts: 1. W0 = 0, 2. Wt is almost surely continuous, and 3. for any two disjoint time intervals [s1, t1] and [s2, t2], the random variables Wt1 ? Ws1 and Wt2 ? Ws2 are independent zero mean Gaussian random variables with variance t1? s1 and t2? s2 respectively. Using It?o?s lemma (see, for example, [KS04]), it can be shown that the stock price at time t is given by St $= S0 \exp((????2/2)t + ?Wt)$. (13) Now, we consider a situation where we have n stocks in the GBM model. Let $? = (?1, ?2, \ldots, ?n)$ be the vector of drifts, and $? = (?1, ?2, \ldots, ?n)$ be the vector of (annualized) volatilities. Suppose we trade for one year. We now study the effect of trading frequency on the quadratic variation of the stock price returns. For this, assume that the year-long trading interval is sub-divided into T equally sized intervals of length 1/T, and we trade at the end of each such interval. Let rt = (rt (1), rt (2), , rt (n)) be the vector of stock returns in the tth trading period. We assume that T is ?(i) 2) for any i. ?large enough?, which is taken to mean that it is larger than ?(i), ?(i), (?(i) 7

Then using the facts of the Wiener process stated above, we can prove the following lemma, which shows that the expected quadratic variation, and its variance, is the essentially the same regardless of trading frequency. The proof is a straightforward calculation and deferred to the full version of this paper. Lemma 9. In the setup of trading n stocks in the GBM model over one year with T trading periods, there is a vector v such that h i PT 2 E? k?k2 (1 + O(T1)) t=1 krt? vk and VAR

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hP T t=1
2
i
krt ? vk
? 6k?k2 (1 + O( T1 )),
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regardless of how the stocks are correlated. Applying this bound in our algorithm, we obtain the following regret bound from Corollary 2. Theorem 10. In the setup of Lemma 9, for any? ; 0, with probability at least 1? 2e??, we have Regret ? $O(n(\log(k?k2 + n) + ?))$. Theorem 10 shows that one expects to achieve constant regret independent of the trading frequency, as long as the total trading period is fixed. This result is only useful if increasing trading frequency improves the performance of the best constant rebalanced portfolio. Indeed, this has been observed empirically (see e.g. [AHKS06], and more empirical evidence is given in the full version of this paper.). To obtain a theoretical justification for increasing trading frequency, we consider an example where we have two stocks that follow independent Black-Scholes models with the same drifts, but different volatilities ?1, ?2. The same drift assumption is necessary because in the long run, the best CRP is the one that puts all its wealth on the stock with the greater drift. We normalize the drifts to be equal to 0, this doesn?t change the performance in any qualitative manner. Since the drift is 0, the expected return of either stock in any trading period is 1; and since the returns in each period are independent, the expected final change in wealth, which is the product of the returns, is also 1. Thus, in expectation, any CRP (indeed, any portfolio selection strategy) has overall return 1. We therefore turn to a different criterion for selecting a CRP. The risk of an investment strategy is measured by the variance of its payoff; thus, if different investment strategies have the same expected payoff, then the one to choose is the one with minimum variance. We therefore choose the CRP with the least variance. We prove the following lemma in the full version of the paper: Lemma 11. In the setup where we trade two stocks with zero drift and volatilities ?1, ?2, the variance of the minimum variance CRP decreases as the trading frequency increases. Thus, increasing the trading frequency decreases the variance of the minimum variance CRP, which implies that it gets less risky to trade more frequently; in other words, the more frequently we trade, the more likely the payoff will be close to the expected value. On the other hand, as we show in Theorem 10, the regret does not change even if we trade more often; thus, one expects to see improving performance of our algorithm as the trading frequency increases.

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Conclusions and Future Work

We have presented an efficient algorithm for regret minimization with expconcave loss functions whose regret strictly improves upon the state of the art. For the problem of portfolio selection, the regret is bounded in terms of the observed variation in stock returns rather than the number of iterations. Recently, DeMarzo, Kremer and Mansour [DKM06] presented a novel game-theoretic framework for option pricing. Their method prices options using low regret algorithms, and it is possible that our analysis can be applied to options pricing via their method (although that would require a much tighter optimization of the constants involved). Increasing trading frequency in practice means increasing transaction costs. We have assumed no transaction costs in this paper. It would be very interesting to extend our portfolio selection algorithm to take into account transaction costs as in the work of Blum and Kalai [BK97]. 8

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