B38DF

Computer Architecture and Embedded Systems

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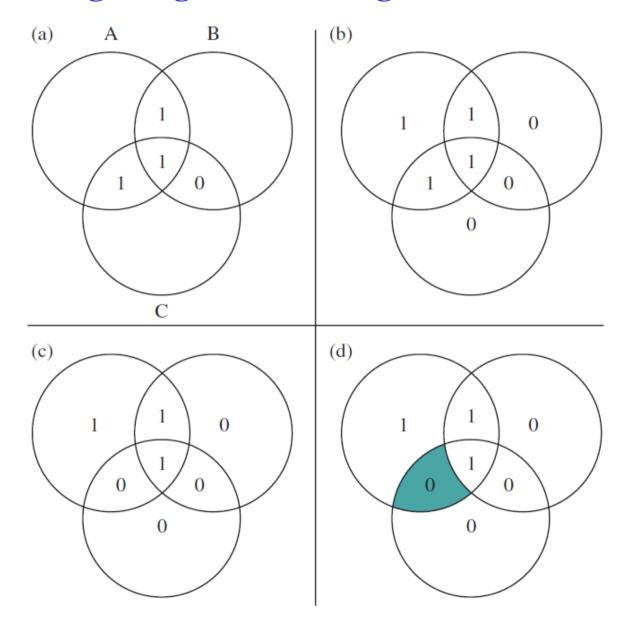
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Based on the slides prepared by Dr. Mustafa Suphi Erden



Hamming's Algorithm – Single error correction

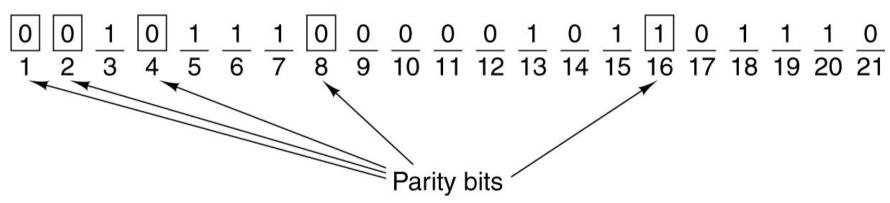


Hamming's Algorithm - Single error correction for any size word

Hamming Code

r parity bits added to an m-bit word, forming a new word of length m+r

Memory word 1111000010101110

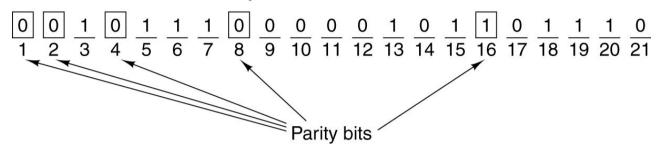


Construction of the Hamming code for the memory word 1111000001011110 by adding 5 check bits to the 16 data bits.

- Number the bits from left-to-right starting from 1.
- All bits whose bit number is a power of 2 are parity bits
- For example with a 16-bit word, 5 parity bits are added: bits 1, 2, 4, 8, 16
 are parity bits, all the rest are data bits. In all the memory word has 21 bits
 (16 data, 5 parity).

Hamming's Algorithm – Single error correction for any size word

Memory word 1111000010101110



- Each parity bit checks specific bit positions:
 - Bit 1 checks bits
- 1, 3, 5, 7, 9, 11, 13, 15, 17, 19, 21
- * * * * 1

- Bit 2 checks bits
- 2, 3, 6, 7, 10, 11, 14, 15, 18, 19
- * * * 1 *

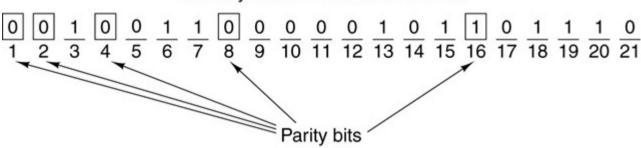
- Bit 4 checks bits
- 4, 5, 6, 7, 12, 13, 14, 15, 20, 21
- * * 1 * *

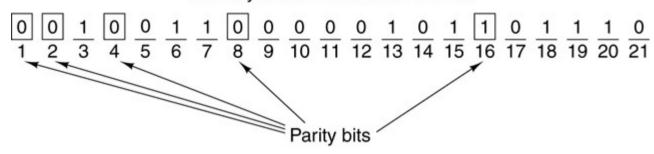
- Bit 8 checks bits
- 8, 9, 10, 11, 12, 13, 14, 15

* 1 * * *

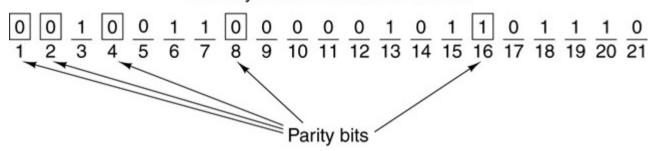
- > Bit **16** checks bits
- 16, 17, 18, 19, 20, 21

- 1 * * * *
- Rule: Bit b is checked by those bits b1, b2, ..., bj such that b1+b2+...+bj = b
- e.g. Bit 5 is checked by bits 1 and 4 because 1+4 = 5; bit 6 is checked by bits 2 and 4 because 2 + 4 = 6.
- The parity bit is set so that the total number of 1s in the checked positions is even.

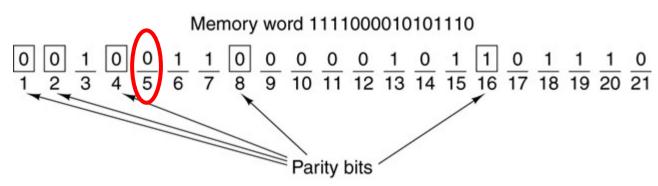




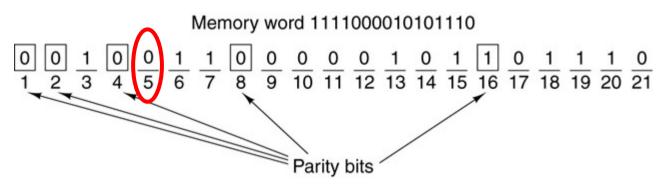
- Where is the error?
 - Parity bit 1 incorrect (1, 3, 5, 7, 9, 11, 13, 15, 17, 19, 21 contain five 1s)
 - Parity bit 2 correct (2, 3, 6, 7, 10, 11, 14, 15, 18, 19 contain six 1s)
 - > Parity bit 4 incorrect (4, 5, 6, 7, 12, 13, 14, 15, 20, 21 contains five 1s)
 - Parity bit 8 correct (8, 9, 10, 11, 12, 13, 14, 15 contain two 1s)
 - > Parity bit **16** correct (16, 17, 18, 19, 20, 21 contain four 1s)



- Where is the error?
 - Parity bit 1 incorrect (1, 3, 5) (7, 9, 11, 13, 15, 17, 19, 21) contain five 1s)
 - Parity bit 2 correct (2, 3, 6, 7, 10, 11, 14, 15, 18, 19 contain six 1s)
 - Parity bit 4 incorrect (4, 5, 6, 7, 12, 13, 14, 15, 20, 21) contains five 1s)
 - Parity bit 8 correct (8, 9, 10, 11, 12, 13, 14, 15 contain two 1s)
 - > Parity bit **16** correct (16, 17, 18, 19, 20, 21 contain four 1s)

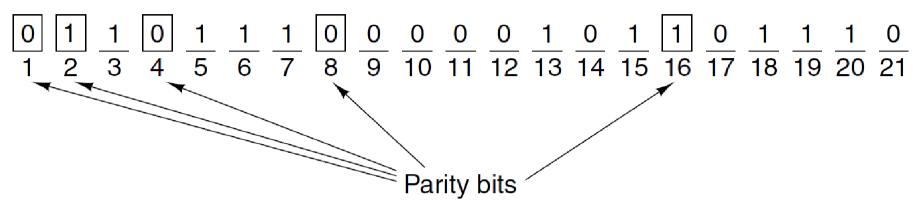


- Where is the error?
 - Parity bit 1 incorrect (1, 3, 5) 7, 9, 11, 13, 15, 17, 19, 21 contain five 1s)
 - > Parity bit **2** correct (2, 3, 6, (7), 10, 11, 14, (15) 18, 19 contain six 1s)
 - Parity bit 4 incorrect (4, 5, 6, 7), 12, 13, 14, 15, 20, 21 contains five 1s)
 - Parity bit 8 correct (8, 9, 10, 11, 12, 13, 14, 15 contain two 1s)
 - Parity bit 16 correct (16, 17, 18, 19, 20, 21) contain four 1s)

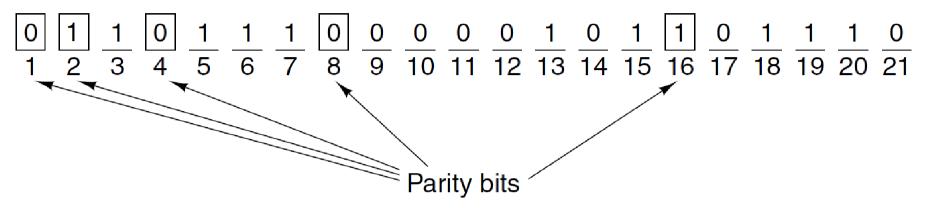


- Where is the error?
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 - Parity bit 8 correct (8, 9, 10, 11, 12, 13, 14, 15 contain two 1s)
 - Parity bit **16** correct (16, 17, 18, 19, 20, 21) contain four 1s)
- Solution: Add up all the incorrect parity bits. The resulting sum is the position of the incorrect bit! (1 + 4 = 5)
- Answer: Bit 5 → Bit 5 should be corrected by the computer from 0 to 1.

Hamming's Algorithm – Another Exercise

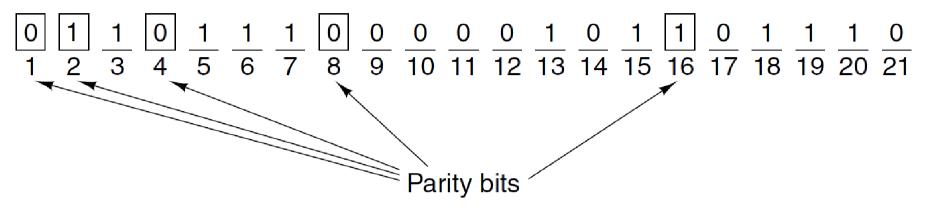


Hamming's Algorithm – Another Exercise



- Where is the error?
 - Parity bit 1 correct (1, 3, 5, 7, 9, 11, 13, 15, 17, 19, 21 contain six 1s)
 - Parity bit 2 incorrect (2, 3, 6, 7, 10, 11, 14, 15, 18, 19 contain seven 1s)
 - > Parity bit 4 correct (4, 5, 6, 7, 12, 13, 14, 15, 20, 21 contains five 1s)
 - Parity bit 8 correct (8, 9, 10, 11, 12, 13, 14, 15 contain two 1s)
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Hamming's Algorithm – Another Exercise



- Where is the error?
 - Parity bit 1 correct (1, 3, 5, 7, 9, 11, 13, 15, 17, 19, 21 contain six 1s)
 - > Parity bit **2** incorrect (2, **3**, **6**, **7**, **10**, **11**, **14**, **15**, **18**, **19** contain seven 1s)
 - > Parity bit 4 correct (4, 5, 6, 7, 12, 13, 14, 15, 20, 21 contains five 1s)
 - Parity bit 8 correct (8, 9, 10, 11, 12, 13, 14, 15 contain two 1s)
 - > Parity bit **16** correct (16, 17, 18, 19, 20, 21 contain four 1s)
- Solution: Add up all the incorrect parity bits. The resulting sum is the position of the incorrect bit! (2 = 2)

Hamming's Algorithm: Single Error Correction (SEC) + Double Error Detection (DED) by adding an extra bit

With Even Parity

With Two Errors

Word

SEC Attempt

(a) $_{\rm L}(b)$ $_{\rm I}(c)$ 0 (d) (e) (f) 0 0 0

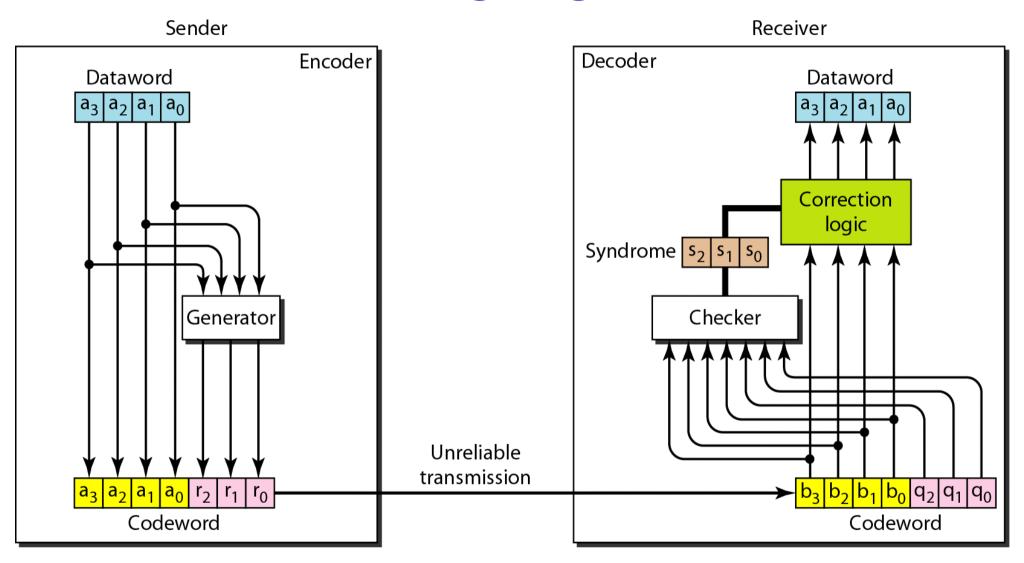
IS SEC Correct? Extra Bit Confirms DE

Hamming's Algorithm: SEC-DED

A single-error-correcting, double-error-detecting (SEC-DED) code is typically used for semiconductor memory.

	Single-Erro	r Correction	Single-Error Double-Erro	Correction/ or Detection
Data Bits	Check Bits	% Increase	Check Bits	% Increase
8	4	50	5	62.5
16	5	31.25	6	37.5
32	6	18.75	7	21.875
64	7	10.94	8	12.5
128	8	6.25	9	7.03
256	9	3.52	10	3.91

Hamming's Algorithm



1. Devise a 7-bit even-parity Hamming code for the digits 0 to 9.

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- $0 \rightarrow 0000$
- $1 \rightarrow 0001$
- $2 \rightarrow 0010$
- $3 \rightarrow 0011$
- $4 \rightarrow 0100$
- $5 \rightarrow 0101$
- $6 \rightarrow 0110$
- $7 \rightarrow 0111$
- $8 \rightarrow 1000$
- $9 \rightarrow 1001$

1. Devise a 7-bit even-parity Hamming code for the digits 0 to 9.

```
\rightarrow 0000
                        → **0*000
    \rightarrow 0001
                         → **0*001
    \rightarrow 0010
                              **()*()1()
    \rightarrow 0011
                              **0*011
    \rightarrow 0100
                              **0*100
    \rightarrow 0101
                    5 \rightarrow **0*101
    \rightarrow 0110
                              **0*110
    \rightarrow 0111
                    7 \rightarrow **0*111
   \rightarrow 1000
                    8 \rightarrow **1*000
9
    \rightarrow 1001
                    9 \rightarrow **1*001
```

Bit 1 checks bits 1, 3, 5, 7, 9

Bit 2 checks bits 2, 3, 6, 7

1. Devise a 7-bit even-parity Hamming code for the digits 0 to 9.

```
0
    \rightarrow 0000
                             **()*()()
                                            0 \rightarrow 0*0*000
    \rightarrow 0001
                             **0*001
                                                \rightarrow 1*0*001
         0010
                             **0*010
                                                \rightarrow 0*0*010
    \rightarrow 0011
                             **0*011
                                                     1*0*011
    \rightarrow 0100
                             **0*100
                                                \rightarrow 1*0*100
    \rightarrow 0101
                             **()*1()1
                                                \rightarrow 0*0*101
    \rightarrow 0110
                             **0*110
                                                \rightarrow 1*0*110
    \rightarrow 0111
                             **0*111
                                                     0*0*111
8
   \rightarrow 1000
                             **1*000
                                                     1*1*000
9
        1001
                             **1*001
                                                    0*1*001
```

Bit 1 checks bits 1, 3, 5, 7, 9

Bit 2 checks bits 2, 3, 6, 7

1. Devise a 7-bit even-parity Hamming code for the digits 0 to 9.

```
0
       0000
                        **0*000
                                        \rightarrow 0*0*000
                                                                 000*000
       0001
                        **0*001
                                                                 110 * 001
                                            1*0*001
       0010
                        **()*()1()
                                            0*0*010
                                                                 010*010
3
       0011
                        **()*()11
                                                                 100 * 011
                3
                                    3
                                            1*0*011
   \rightarrow
   \rightarrow 0100
                        **0*100
                                            1*0*100
                                                                 100 * 100
      0101
                        **()*1()1
                                            0*0*101
                                                                 010*101
6
   \rightarrow 0110
                        **0*110
                                            1*0*110
                                                                 110*110
7
       0111
                        **0*111
                                            0*0*111
                                                                 000*111
8
       1000
                        **1*000
                                    8
                                            1*1*000
                                                                 111*000
   \rightarrow
                                        \rightarrow
9
       1001
                9
                        **1*001
                                    9
                                                                 001*001
                                            0*1*001
```

Bit 1 checks bits 1, 3, 5, 7, 9

Bit 2 checks bits 2, 3, 6, 7

1. Devise a 7-bit even-parity Hamming code for the digits 0 to 9.

```
0
       0000
                       **0*000
                                                               000*000
                                      \rightarrow 0*0*000
                                                                                  0000000
       0001
                       **0*001
                                                               110 * 001
                                                                                  1101001
                                          1*0*001
       0010
                                                               010*010
                       **0*010
                                          0*0*010
                                                                                  0101010
3
       0011
                                                               100 * 011
                                                                                  1000011
                3
                       **0*011
                                   3
                                           1*()*()11
                                                       3
                                                                           3
   \rightarrow
       0100
                                                               100 * 100
                       **0*100
                                          1*0*100
                                                                                  1001100
5
       0101
                       **()*1()1
                                          0*0*101
                                                               010*101
                                                                           5
                                                                                  0100101
6
       0110
                       **0*110
                                           1*0*110
                                                               110*110
                                                                                  1100110
7
       0111
                                                               000*111
                       **0*111
                                           0*0*111
                                                                                  0001111
8
       1000
                       **1*000
                                   8
                                           1*1*000
                                                       8
                                                               111*000
                                                                           8
                                                                                  1110000
   \rightarrow
                                      \rightarrow
9
       1001
                                   9
                                                               001*001
                                                                                  0011001
                       **1*001
                                           0*1*001
                                                       9
                                                                           9
```

Bit 1 checks bits 1, 3, 5, 7, 9

Bit 2 checks bits 2, 3, 6, 7

2. Devise a code for the digits 0 to 9 whose Hamming distance is 2.

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Solution: Just add a parity bit:

Problems

2. Devise a code for the digits 0 to 9 whose Hamming distance is 2.

Solution: Just add a parity bit:

```
\begin{array}{cccc}
0 & \to & 0000 \\
1 & \to & 0001 \\
2 & \to & 0010 \\
3 & \to & 0011 \\
4 & \to & 0100 \\
4 & \to & 0100 \\
5 & \to & 0101 \\
6 & \to & 0110 \\
7 & \to & 0111 \\
8 & \to & 1000 \\
9 & \to & 1001 \\
0
\end{array}
```

3. In a Hamming code, some bits are "wasted" in the sense that they are used for checking and not information. What is the percentage of wasted bits for messages whose total length (data + check bits) is $2^n - 1$? Evaluate this expression numerically for values of n from 3 to 10.

3. In a Hamming code, some bits are "wasted" in the sense that they are used for checking and not information. What is the percentage of wasted bits for messages whose total length (data + check bits) is $2^n - 1$? Evaluate this expression numerically for values of n from 3 to 10.

Solution: If the total length is $2^n - 1$ bits, there are n check bits. Consequently, the percentage of "wasted" bits is $n/(2^n - 1) \times 100\%$. Numerically for n from 3 to 10 we get: 42.9%, 26.7%, 16.1%, 9.5%, 5.5%, 3.1%, 1.8%, and 1.0%.

4. An extended ASCII character is represented by an 8-bit quantity. The associated Hamming encoding of each character can then be represented by a string of three hex digits. Encode the following extended five-character ASCII text using an even-parity Hamming code: **Earth**. Show your answer as a string of hex digits.

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Hex	Char	Hex	Char	Hex	Char	Hex	Char	Hex	Char	Hex	Char
20	(Space)	30	0	40	@	50	Р	60	6	70	р
21	!	31	1	41	Α	51	Q	61	а	71	q
22	п	32	2	42	В	52	R	62	b	72	r
23	#	33	3	43	С	53	S	63	С	73	s
24	\$	34	4	44	D	54	Т	64	d	74	t
25	%	35	5	45	Е	55	U	65	е	75	u
26	&	36	6	46	F	56	V	66	f	76	V
27	,	37	7	47	G	57	W	67	g	77	w
28	(38	8	48	Н	58	Χ	68	h	78	X
29)	39	9	49	1	59	Υ	69	i	79	у
2A	*	ЗА	:	4A	J	5A	Z	6A	j	7A	Z
2B	+	3B	;	4B	K	5B	[6B	k	7B	{
2C	,	3C	<	4C	L	5C	\	6C	1	7C	1
2D	-	3D	=	4D	M	5D]	6D	m	7D	}
2E		3E	>	4E	Ν	5E	^	6E	n	7E	~
2F	/	3F	?	4F	Ο	5F	_	6F	0	7F	DEL

ASCII (American Standard Code for Information Interchange)

$$E \rightarrow 45$$

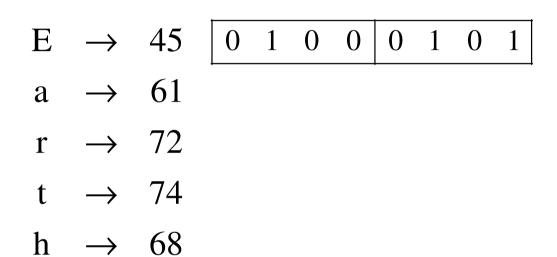
$$a \rightarrow 61$$

$$r \rightarrow 72$$

$$t \rightarrow 74$$

$$h \rightarrow 68$$

4. An extended ASCII character is represented by an 8-bit quantity. The associated Hamming encoding of each character can then be represented by a string of three hex digits. Encode the following extended five-character ASCII text using an even-parity Hamming code: **Earth**. Show your answer as a string of hex digits.



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E	\rightarrow	45	0	1	0	0	0	1	0	1									
a	\rightarrow	61					1	1	2	3	4	5	6	7	8	9	10	11	12
r	\rightarrow	72					*	<	*	0	*	1	0	0	*	() 1	0	1
t	_	71																	

 $h \rightarrow 68$

Bit 1 checks bits 1, 3, 5, 7, 9, 11

Bit 2 checks bits 2, 3, 6, 7, 10, 11

Bit 4 checks bits 4, 5, 6, 7, 12

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a	\rightarrow	61	•													10		
r	\rightarrow	72					*	*	0	*	1	0	0	*	0	1	0	1
t	\rightarrow	74					1	*	0	*	1	0	0	*	0	1	0	1
h	\rightarrow	68																

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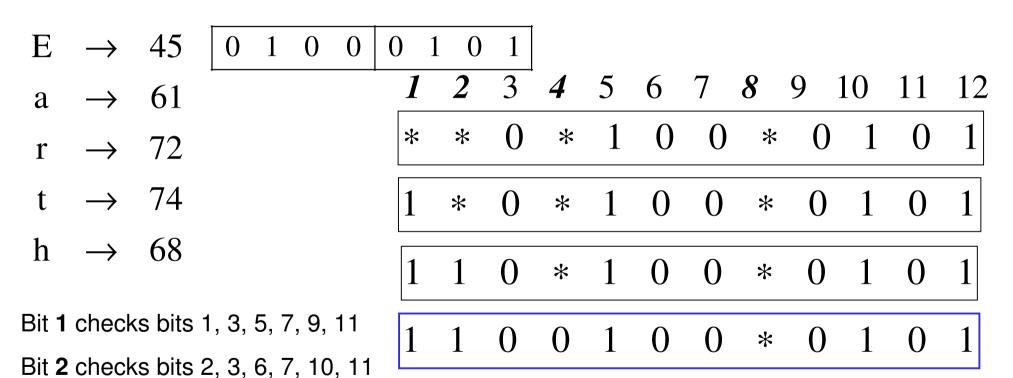
E	\rightarrow	45	0	1	0	0	0	1 0	1									
a	\rightarrow	61	•			•	1	2	3	4	5	6	7	8	9	10	11	12
r	\rightarrow	72					*	*	0	*	1	0	0	*	0	1	0	1
t	\rightarrow	74					1	*	0	*	1	0	0	*	0	1	0	1
h	\rightarrow	68																
							1	1	U	*	1	U	U	*	U	1	U	1

Bit 1 checks bits 1, 3, 5, 7, 9, 11

Bit 2 checks bits 2, 3, 6, 7, 10, 11

Bit 4 checks bits 4, 5, 6, 7, 12

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Bit 4 checks bits 4, 5, 6, 7, 12

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a	\rightarrow	61	-			•	1	2	3	4	5	6	7	8	9	10	11	12
r	\rightarrow	72					*	*	0	*	1	0	0	*	0	1	0	1
t	\rightarrow	74					1	*	0	*	1	0	0	*	0	1	0	1
h	\rightarrow	68					1	1	0	*	1	0	0	*	0	1	0	1
Bit 1	check	s bits	1, 3,	5, 7	, 9, ·	11	1	1	0	0	1	0	0	*	0	1	0	1
Bit 2	check	s bits	2, 3,	6, 7	, 10 _.	, 11	1		<u> </u>	<u> </u>	1			•	0			
Bit 4	check	s bits	4, 5,	6, 7	', 12		1	1	0	0	1	0	0	0	0	1	0	1

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a	\rightarrow	61	•				1	2	3	4	5	6	7	8	9	10	11	12
r	\rightarrow	72					*	*	0	*	1	0	0	*	0	1	0	1
t	\rightarrow	74					1	*	0	*	1	0	0	*	0	1	0	1
h	\rightarrow	68					1	1	0	*	1	0	0	*	0	1	0	1
		s bits					1	1	0	0	1	0	0	*	0	1	0	1
Bit 2	check	s bits	2, 3,	6, 7	, 10	, 11												
Bit 4	check	s bits	4, 5,	6, 7	, 12		1	1	0	0	1	0	0	0	0	1	0	1
			_															

C85

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5. The following string of hex digits encodes extended ASCII characters in an even-parity Hamming code: 0D3 DD3 0F2 5C1 1C5 CE3. Decode this string and write down the characters that are encoded.

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Solution: Each 8-bit ASCII character is encoded into three hex digits.

The first set of hex digits is 0D3:

0 0 0 0	1 1 0 1	0 0 1 1
---------	---------	---------

5. The following string of hex digits encodes extended ASCII characters in an even-parity Hamming code: 0D3 DD3 0F2 5C1 1C5 CE3. Decode this string and write down the characters that are encoded.

Solution: Each 8-bit ASCII character is encoded into three hex digits.

Bit 1 checks bits 1, 3, 5, 7, 9, 11

Bit 2 checks bits 2, 3, 6, 7, 10, 11

Bit 4 checks bits 4, 5, 6, 7, 12

Bit 8 checks bits 8, 9, 10, 11, 12

The first set of hex digits is 0D3:

1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	1	1	0	1	0	0	1	1

5. The following string of hex digits encodes extended ASCII characters in an even-parity Hamming code: 0D3 DD3 0F2 5C1 1C5 CE3. Decode this string and write down the characters that are encoded.

Solution: Each 8-bit ASCII character is encoded into three hex digits.

Bit 1 checks bits 1, 3, 5, 7, 9, 11

Bit 2 checks bits 2, 3, 6, 7, 10, 11

Bit 4 checks bits 4, 5, 6, 7, 12

Bit 8 checks bits 8, 9, 10, 11, 12

The first set of hex digits is 0D3:

0 0 0 0	1 1 0 1	0 0 1 1
---------	---------	---------

1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	1	1	0	1	0	0	1	1

bit 1 is correct

bit 2 is correct

bit **4** is incorrect (has the wrong parity)

bit 8 is incorrect (has the wrong parity)

5. The following string of hex digits encodes extended ASCII characters in an even-parity Hamming code: 0D3 DD3 0F2 5C1 1C5 CE3. Decode this string and write down the characters that are encoded.

Solution: Each 8-bit ASCII character is encoded into three hex digits.

Bit 1 checks bits 1, 3, 5, 7, 9, 11

Bit 2 checks bits 2, 3, 6, 7, 10, 11

Bit 4 checks bits 4, 5, 6, 7, 12

Bit 8 checks bits 8, 9, 10, 11, 12

Add up all the incorrect parity bits. The resulting sum is the position of the incorrect bit: $\mathbf{4} + \mathbf{8} = 12$

The first set of hex digits is 0D3:

0 0 0 0	1 1 0 1	0 0 1 1
---------	---------	---------

1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	1	1	0	1	0	0	1	1

bit 1 is correct

bit 2 is correct

bit **4** is incorrect (has the wrong parity)

bit **8** is incorrect (has the wrong parity)

corrected code:

0	0	0	0	1	1	0	1	0	0	1	0

5. The following string of hex digits encodes extended ASCII characters in an even-parity Hamming code: 0D3 DD3 0F2 5C1 1C5 CE3. Decode this string and write down the characters that are encoded.

Solution: Each 8-bit ASCII character is encoded into three hex digits.

Bit 1 checks bits 1, 3, 5, 7, 9, 11

Bit 2 checks bits 2, 3, 6, 7, 10, 11

Bit 4 checks bits 4, 5, 6, 7, 12

Bit 8 checks bits 8, 9, 10, 11, 12

Add up all the incorrect parity bits. The resulting sum is the position of the incorrect bit: 4 + 8 = 12

The first set of hex digits is 0D3:

0 0 0 0 1 1 0 1 0 0 1 1	0 (0	0 0	0)	0	0	0		0	(0)	0	(
-------------------------	-----	---	-----	---	---	---	---	---	--	---	---	--	---	---	---	---

1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	1	1	0	1	0	0	1	1

bit 1 is correct

bit 2 is correct

bit **4** is incorrect (has the wrong parity)

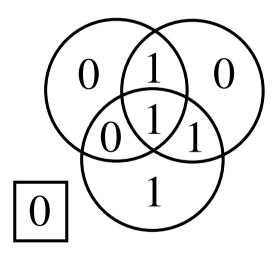
bit **8** is incorrect (has the wrong parity)

corrected code:

5. The following string of hex digits encodes extended ASCII characters in an even-parity Hamming code: 0D3 DD3 0F2 5C1 1C5 CE3. Decode this string and write down the characters that are encoded.

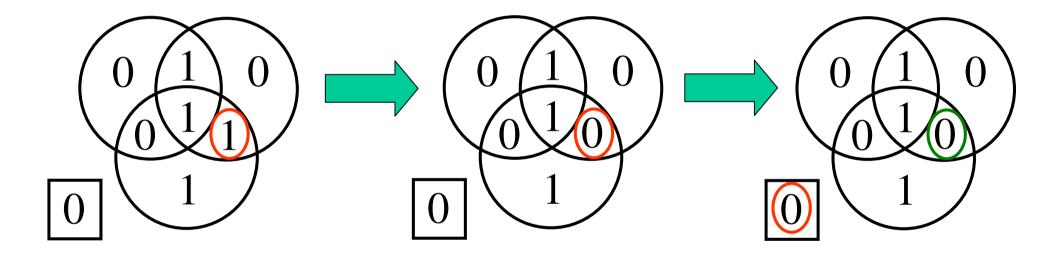
Solution: Each 8-bit ASCII character is encoded into three hex digits. The first set of hex digits: 0D3, has an error in bit 12 (as indicated by the fact that bit 4 and bit 8 have the wrong parity). The next set, DD3 has bit 11 wrong; the set 0F2 has bit 7 wrong; the set 5C1 has bit 9 wrong; the set 1C5 has bit 1 wrong; the last set CE3 does not contain any errors. After the bit positions are corrected and the data extracted from the code words and looked up in the ASCII table, the encoded characters are: **babies**.

6. Check if there is a single/double error. In the single error case, correct the error.



6. Check if there is a single/double error. In the single error case, correct the error.

Solution:



We have a double error.