

Developing Best-In-Class Security Principles with Open Source Firmware

Vincent Zimmer Senior Principal Engineer Intel Corporation

STTS003



Agenda

• Problem Statement

• Ingredients

System Management Mode (SMM)

Open Platforms



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System Management Mode (SMM)

Open Platforms





Bootkits



Bootkits



BIOS Malware

Bootkits



BIOS Malware

Bootkits

SMM Rootkits



BIOS Malware UEFI Rootkits

Bootkits SMM Rootkits



BIOS Malware

UEFI Rootkits

Bootkits

SMM Rootkits

Evil Maid

HVM Rootkits (Blue Pill)



BIOS Malware UEFI Rootkits **Bootkits SMM Rootkits** Device FW Malware **ACPI Rootkits** Option ROM Malware | Evil Maid HVM Rootkits (Blue Pill) **HW Trojans**



Security Fundamentals

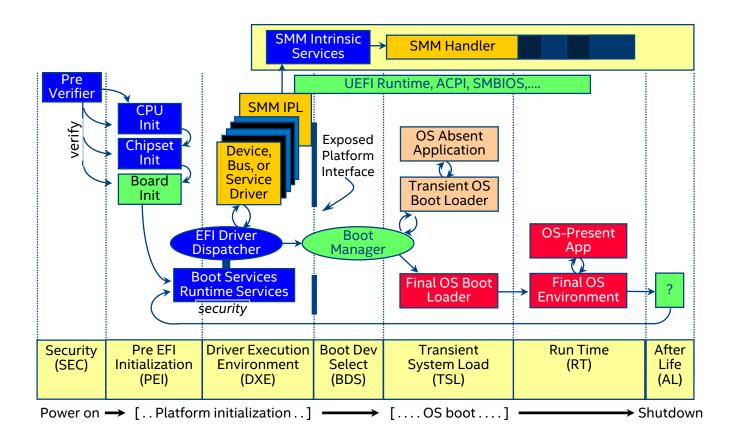




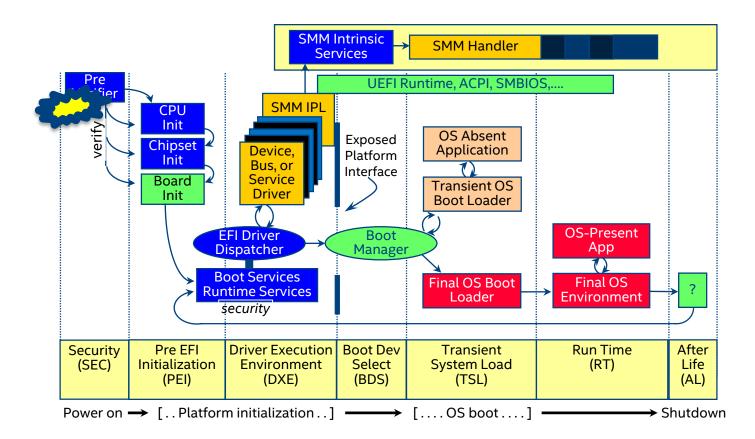
Security Fundamentals



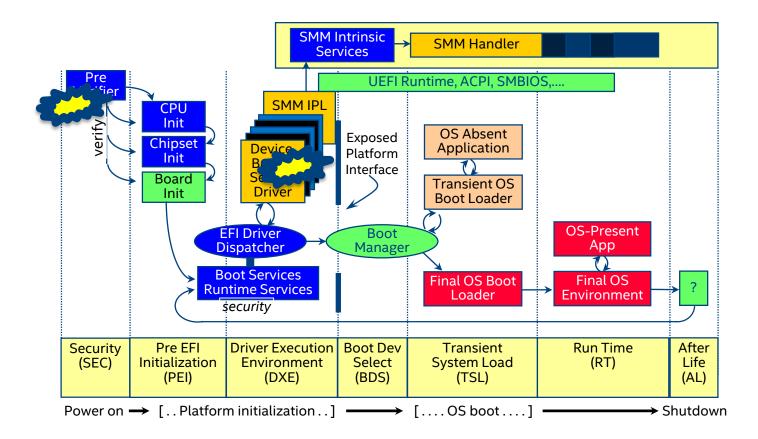




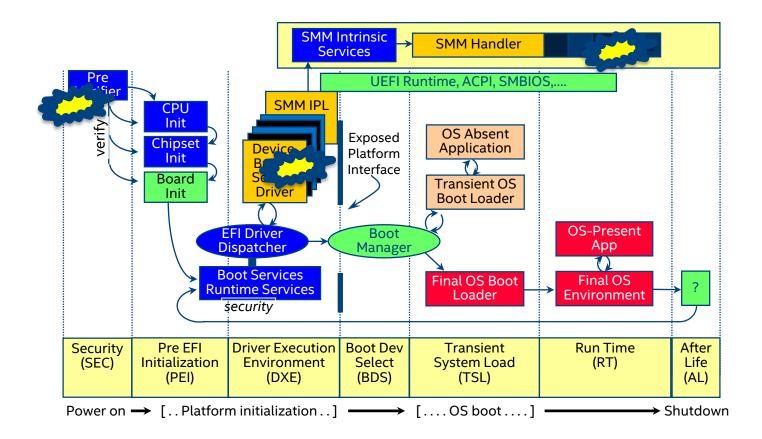




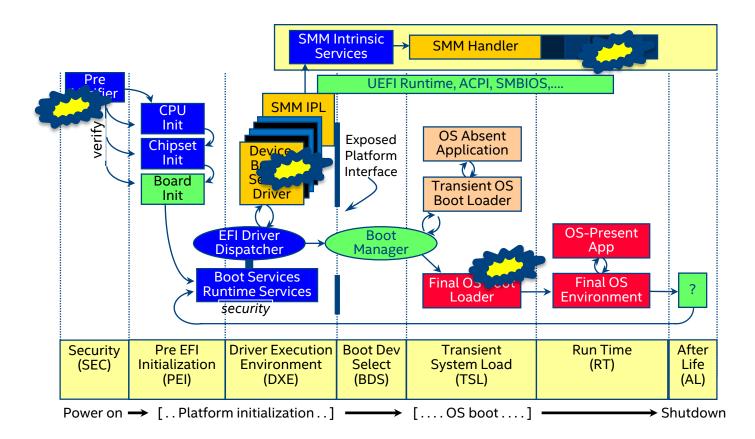














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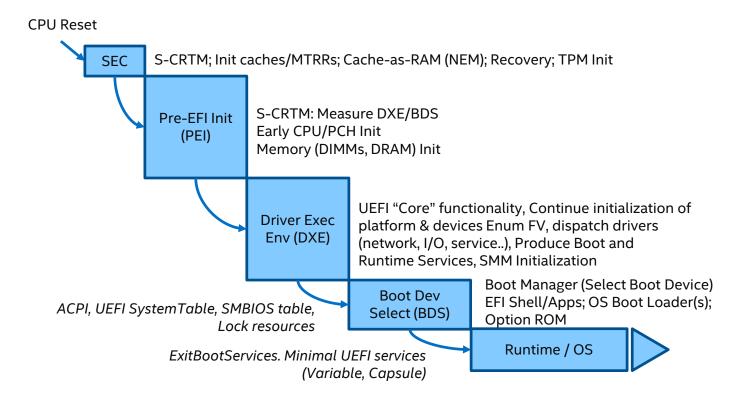
Ingredients

System Management Mode (SMM)

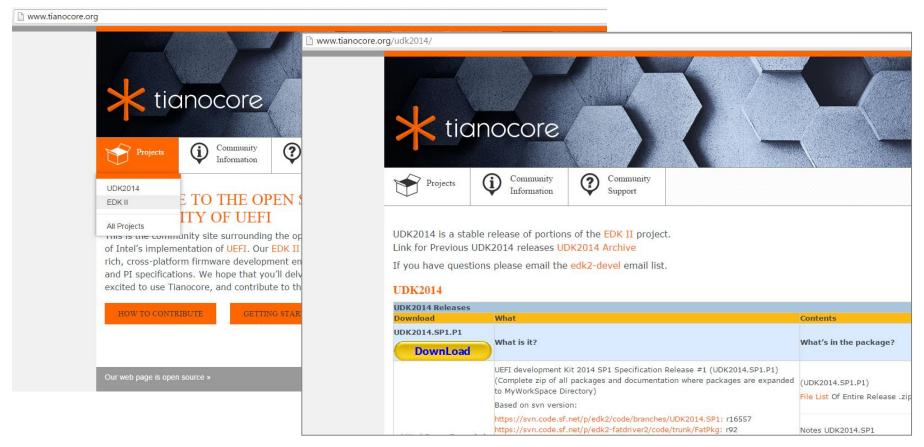
Open Platforms



UEFI PI [Compliant] Firmware



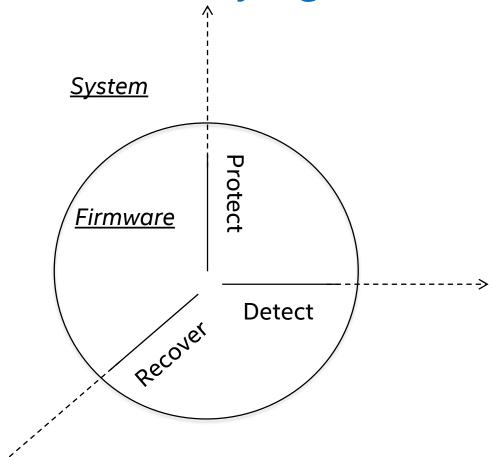




UDK2014 Available on Tianocore.org UDK2015 Coming Soon



Usage of the EDK II Security Ingredients





UEFI Secure Boot vs. TCG Trusted Boot

UEFI PI will measure OS loader & UEFI drivers into TPM (1.2 or 2.0) PCR **UEFI** authenticate OS **UEFI Firmware** (Platform Configuration 🦄 Register) (pub key and policy) UEFI OS Ldr, **Drivers**

Kernel

Drivers

Apps

Check signature of before loading

loader

- UEFI Secure boot will stop platform boot if signature not valid (OEM to provide remediation capability)
- UEFI will require remediation mechanisms if boot fails

 TCG Trusted boot will never fail

TPM

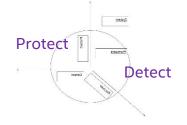
Protect

 Incumbent upon other software to make security decision using attestation



Detect

UEFI Development Kit 2014 SecurityPkg



RandomNumberGenerator

• UEFI driver implementing the EFI_RNG_PROTOCOL from the UEFI2.4 specification

Trusted Computing Group (TCG)

- PEI Modules & DXE drivers implementing Trusted Computing Group measured boot
- EFI_TCG_PROTOCOL and EFI_TREE_PROTOCOL from the TCG and Microsoft* MSDN websites, respectively

UserIdentification

- DXE drivers that support multi-factor user authentication
- Chapter 31 of the UEFI 2.4 specification

Library

• DxeVerificationLib for "UEFI Secure Boot", chapter 27.2 of the UEFI 2.4 specification + other support libs

VariableAuthenticated

• SMM and runtime DXE authenticated variable driver, chapter 7 of the UEFI2.4 specification

https://svn.code.sf.net/p/edk2/code/trunk/edk2/SecurityPkg



Additional Capabilities in Open Source



Variable Lock Protocol

Make variables read-only

https://github.com/tianocore/edk2/blob/master/MdeModulePkg/Include/Protocol/VariableLock.h

Lock Box

Protect content across re-starts

https://github.com/tianocore/edk2-MdeModulePkg/blob/master/Include/Protocol/LockBox.h

Capsule Update

Generic capsule update driver support

http://comments.gmane.org/gmane.comp.bios.tianocore.devel/8402

Recovery

Device support for recovery from PEI

https://svn.code.sf.net/p/edk2/code/trunk/edk2/MdeModulePkg/Include/Guid/RecoveryDevice.h

https://svn.code.sf.net/p/edk2/code/trunk/edk2/



Code Management

Analyze and Mark external Interfaces where input can be attacker controlled data, comment headers

```
/**
    Install child handles if the Handle supports GPT partition structure.
  Caution: This function may receive untrusted input.
   The GPT partition table is external input, so this routine will do basic validation
   for GPT partition table before install child handle for each GPT partition.
  @param[in] This Calling context.
  @param[in] Handle Parent Handle.
  @param[in] DevicePath Parent Device Path.
**/
EFI STATUS
PartitionInstallGptChildHandle
```

UEFI Development Kit 2010 example:

http://edk2.svn.sourceforge.net/svnroot/edk2/trunk/edk2/MdeModulePkg/Universal/Disk/PartitionDxe/Gpt.c



Code Management

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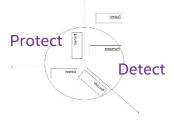
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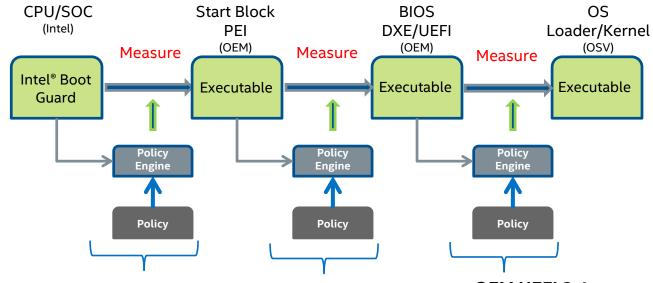
UEFI Development Kit 2010 example:

http://edk2.svn.sourceforge.net/svnroot/edk2/trunk/edk2/MdeModulePkg/Universal/Disk/PartitionDxe/Gpt.c



Full Verified Boot Sequence





Intel® Device Protection Technology with Boot Guard

http://www.intel.com/content/dam/www/public/us/en/documents/product-briefs/4th-gen-core-family-mobile-brief.pdf

OEM PI Verification Using PI Signed Firmware Volumes

Vol 3, section 3.2.1.1 of PI 1.3 Specification

OEM UEFI 2.4
Secure Boot

Chapter 27.2 of The UEFI 2.4 Specification



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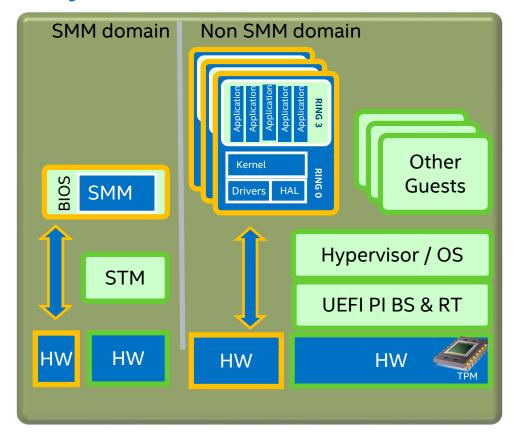
• Ingredients

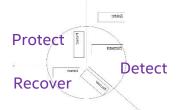
System Management Mode (SMM)

Open Platforms



Full System Picture – UEFI PI Boot and Runtime

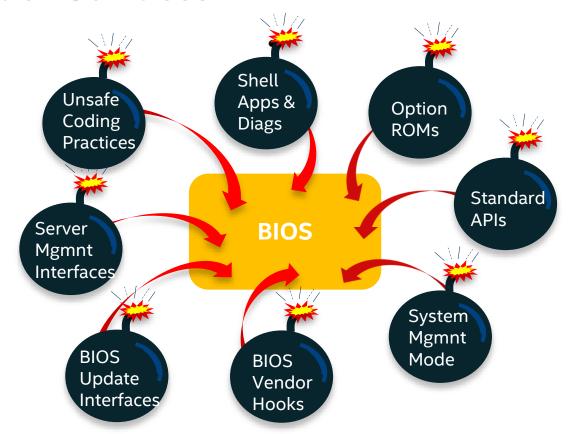




- Protect & Recover the UEFI PI implementation
 - UEFI Capsule Update
 - Hardware Secure Boot using Boot Guard on non-open platforms
- Detect if the Hypervisor and OS is expected one
 - UEFI Secure Boot (and TXT+LCP on non-open platforms)
 - EFI TCG Measured boot
- Protect at runtime
 - SMM Transfer Monitor (STM) to protect platform, hypervisor, and operating system (OS) from the BIOS SMM



BIOS Attack Surfaces



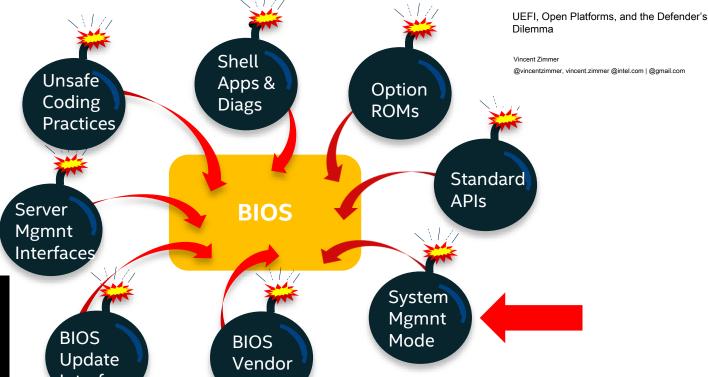


BIOS Attack Surfaces





Attacking Hypervisors Using Firmware and Hardware Bulygin, Matrosov, Gorobets, & Bazhaniuk





Interfaces

Hooks

System Management Mode (SMM)

- SMM
- SMM is the most privileged software in the system
- It has access to all host accessible resources
 - Memory, TPM, chipset registers, device registers
 - It can be used to protect flash which contains UEFI code and variables
- It is not affected by typical OS/VMM level software controls
 - Protection rings, paging, VMX...
 - System Management Interrupt (SMI) can't be masked
 - SMRAM can't be inspected & is transparent to typical system software
- It is commonly critical for proper system operation
- Mitigations
 - code review, validate internal/external input, no call outs



System Management Mode RAM (SMRAM)

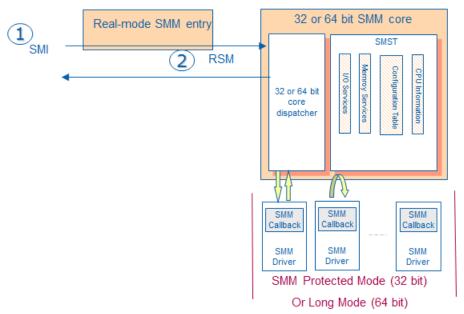
- Three SMRAM regions today: CSEG, HSEG (CSEG alias), and TSEG
 - CPU core's view of SMRAM based on internal register, SMBASE
 - SMM state save area
 - SMI entry point
- MSEG cleaved from top of TSEG on a 4K boundary
 - Related registers (programmed by BIOS):
 - IA32_SMM_MONITOR_CTL.MSEG_BASE

HSEG (CSEG alias) MSEG (inside TSEG) **TSEG** A000 .. BFFF **CSEG**

4 GB



System Management Mode with UEFI PI

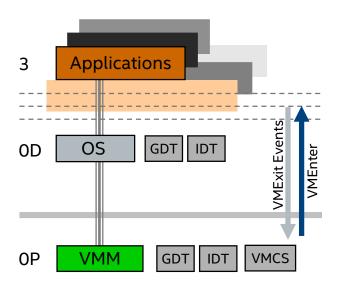


- Orange regions are SMRAM
- Software model defined in PI 1.4 specification, volume 4
- Implementation at edk2\MdeModulePkg\Core\PiSmmCore



Intel® Virtualization Technology (Intel® VT)

Virtual Machines



Virtual Machine Monitor

VMEXIT Conditions

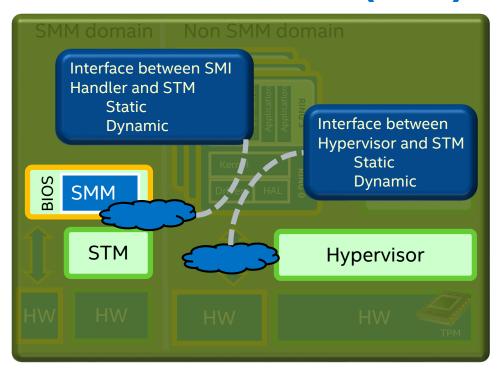
- CRO, CR4 accesses (basic CPU operations)
- CR3 writes (address space changes)
- CR3 reads, INVLPG (paging)
- MSR & debug register accesses
- I/O instructions (per-port bitmap)
- CPUID, INVD
- Exceptions

VM Control Structure (VMCS)

- Which operations cause VMEXITs
- Which states change on VMEXITs and VMENTER
- VMM state area (state loaded on VMEXITs)
- Guest state area (saved on VMEXIT, restored on VMENTER)



SMI Transfer Monitor (STM)



• STM user guide defines software interfaces to manage:

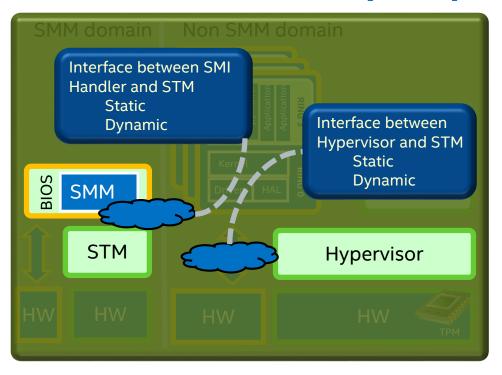
- Setup
- Teardown
- Steady state (runtime)
- ... it also defines some optional ACPI and SMI based interfaces

STM user guide does NOT define:

- Protection policy
- How protections are achieved This is done using normal Intel®
 Architecture mechanisms (Intel
 VT, paging, etc.)



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The STM provides isolation from the SMI handler

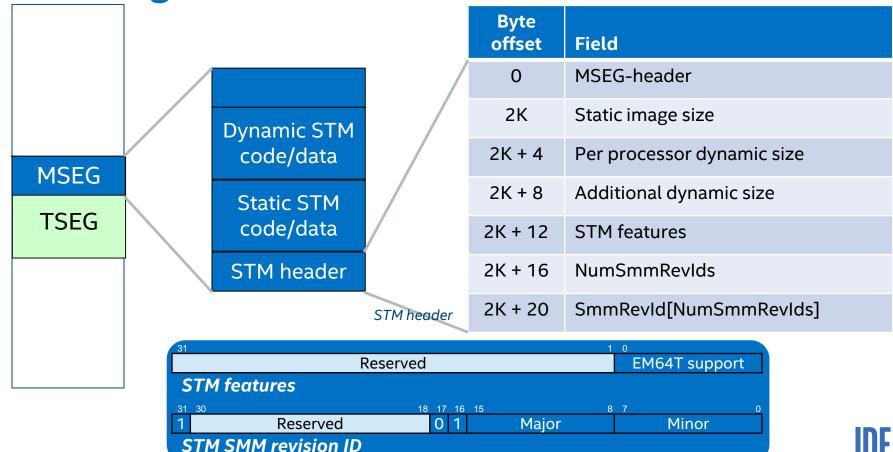


BIOS STM Opt-in

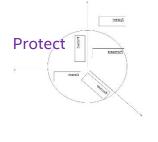
- BIOS should vigorously defend SMRAM
 - ...because of its power, and critical importance to platform function
 - Therefore, BIOS must not enable an arbitrary or unknown STM
- BIOS populates MSEG with an STM image
 - BIOS should enforce it's own policy regarding what is an "acceptable" STM
 - Likely policy:
 - ❖STM image supplied as part of BIOS flash image
 - ❖ BIOS flash image has controlled updates (e.g. signed)
 - ❖Therefore: "I found it in my flash, so it's acceptable"
 - Many other BIOS policy options are possible
- IA32_SMM_MONITOR_CTL.[0] /* Valid bit */
 - BIOS sets to 1 if STM is present, BIOS clears to 0 (default) if no STM is present
 - Must be programmed identically across all CPU threads, Register only writable from SMM
- STM is idle and quiescent until it is "configured"
 - SMI is handled via legacy SMI mechanism

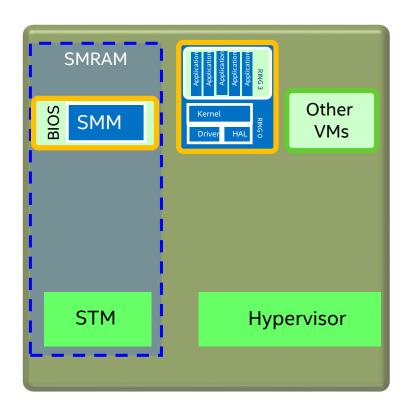


STM Image Format



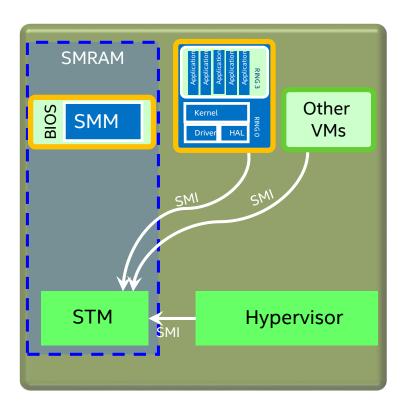








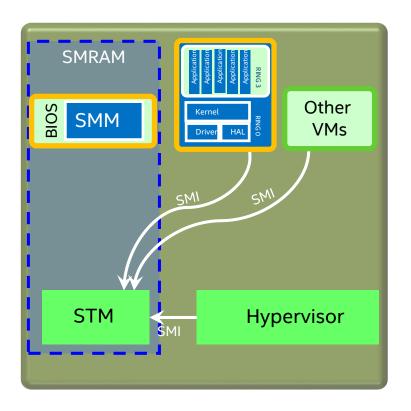




 SMI occurs - control is transferred to STM (VMEXIT)



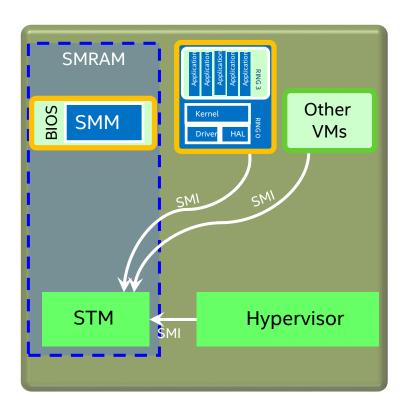




- SMI occurs control is transferred to STM (VMEXIT)
- STM creates SMM state save area for BIOS



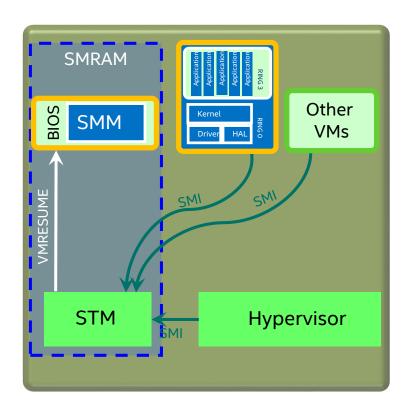




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 - Scrubs register state if protected code has been interrupted by SMI



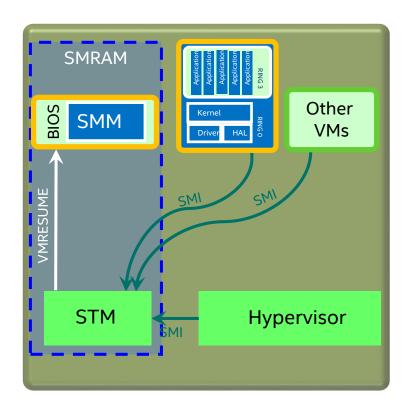




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- STM resumes BIOS SMI handler in guest VM



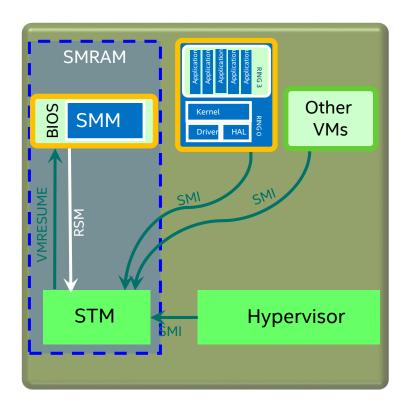




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- STM resumes BIOS SMI handler in guest VM
- BIOS SMM code handles SMI
- STM traps on protected hardware accesses
 - Based on negotiated protection profile



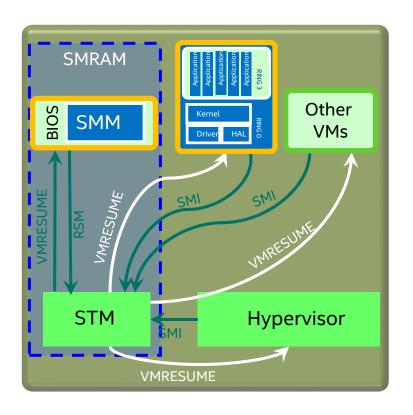




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- BIOS SMM executes RSM (VMEXIT) to STM







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- BIOS SMM code handles SMI
- STM traps on protected hardware accesses
 - Based on negotiated protection profile
- BIOS SMM executes RSM (VMEXIT) to STM
- STM restores interrupted VMs and resumes them

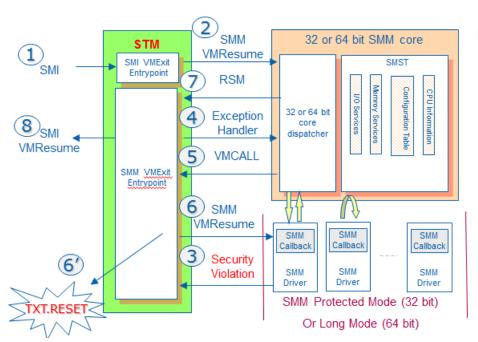


How to Declare Resources Allowed for SMM

- STM allocates hardware resources to BIOS SMI and MLE on a firstcome-first-served basis. BIOS always has first opportunity to make a request
 - This is done statically via the BiosHwResourceRequirementsPtr
 - 64 bit physical pointer to a STM_RESOURCE_LIST



SMM Flow with STM



SMI Transfer Monitor (Sim)

- firmware.intel.com to find STM user guide
- STM Reference implementation built on EDKII infrastructure
 - Build system, Mde Libraries, test driver and MinnowMax integration





Beyond STM Isolation, Moving to Testing

Usenix* WOOT 2015: KLEE → S2E →....

Symbolic execution for BIOS security¹

Oleksandr Bazhaniuk, John Loucaides, Lee Rosenbaum, Mark R. Tuttle, Vincent Zimmer² *Intel Corporation*May 25, 2015

Abstract

We are building a tool that uses symbolic execution to search for BIOS security vulnerabilities including dangerous memory references (call outs) by SMM interrupt handlers in UEFI-compliant implementations of BIOS. Our tool currently applies only to interrupt handlers for SMM variables. Given a snapshot of SMRAM, the base address of SMRAM, and the address of the variable interrupt handler in SMRAM, the tool uses S²E to run the KLEE symbolic execution engine to search for concrete

This point exploded into public view [1] at the Can-SecWest conference in March 2015. Among several interesting results was a paper provocatively titled "How many million BIOSes would you like to infect?" [2]. The authors made the following observation: Almost all machines are vulnerable because almost all machines are running unpatched BIOS (most consumers don't know patches exist, and even sophisticated consumers apply patches with their fingers crossed), and widespread software reuse in the BIOS community (normally considered

WOOT 2015 Paper

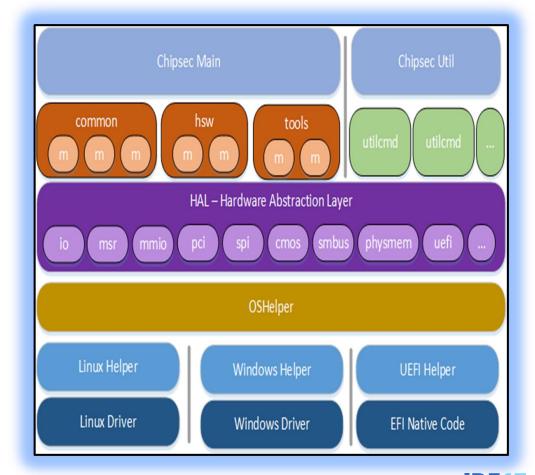


chipsec

 A platform security assessment framework for risk assessment

 Can be extended to meet specific platform security concerns

Open sourced
 https://github.com/chipsec/chipsec/chipsec





Agenda

Problem Statement

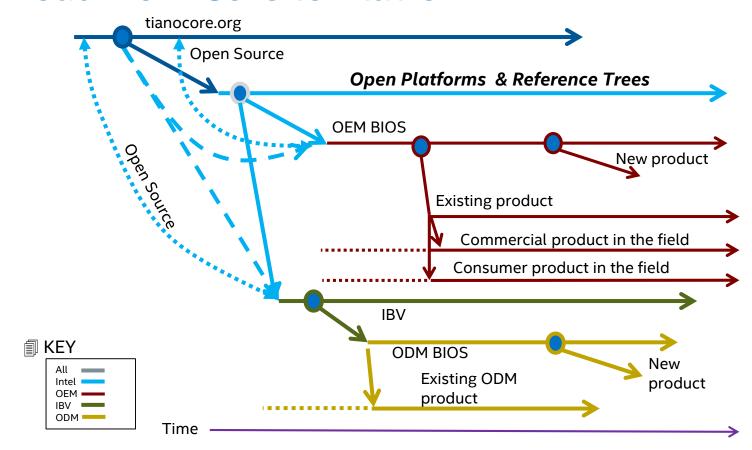
• Ingredients

System Management Mode (SMM)

Open Platforms

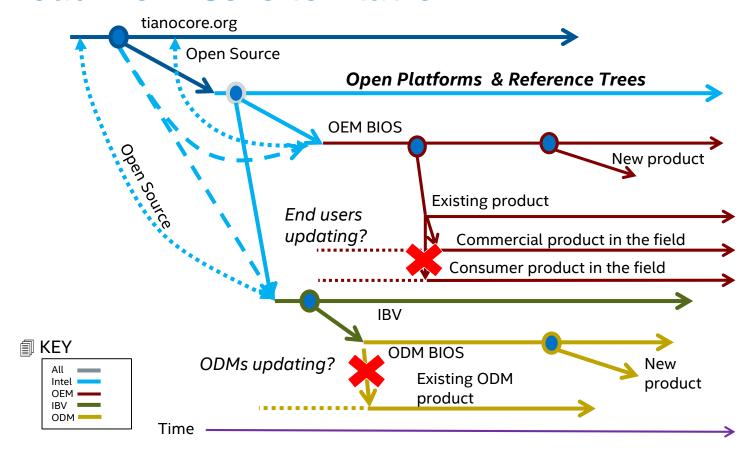


The Road from Core to Platform





The Road from Core to Platform





MinnowBoard Max

- Open hardware platform
- Intel® Atom™ SoC E38xx Series SoC single or dual core
- From http://firmware.intel.com/projects
- This project focuses in on the firmware source code (and binary modules) required to create the boot firmware image for the MinnowBoard MAX. The UEFI Open Source (EDKII project) packages for MinnowBoard MAX are available at http://tianocore.sourceforge.net/wiki/EDK2. To learn more about getting involved in the UEFI EDKII project visit the How to Contribute page.
- The source code builds using Microsoft Visual Studios* and GNU* C Compiler (for both 32 and 64 bit images) production and debug execution environments. The source code builds the same UEFI firmware image shipping on MinnowBoard MAX.
- See more at: http://firmware.intel.com/projects#sthash.100c8srY.dpuf





MinnowBoard Max

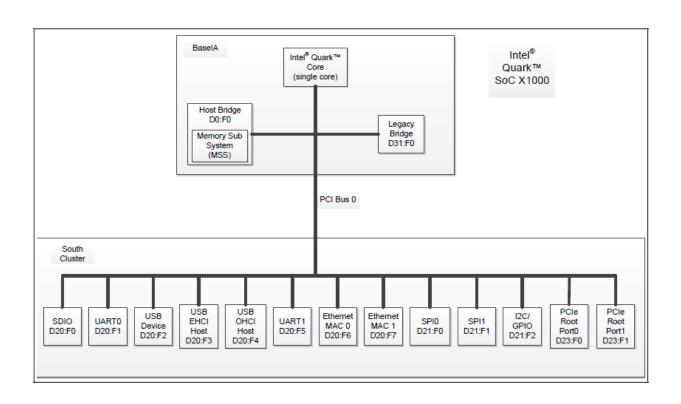
Focused on the maker community, but....

- 64-bit Intel[®] Atom[™] SoC E38xx Series
- Has UEFI Secure Boot
- Built off of live tree
- Supports the SMM Transfer Monitor (STM) without Intel® Trusted Execution Technology (Intel® TXT)
- Ability to update with latest capabilities on http://www.tianocore.org



Intel® Quark™ SoC – Hardware Overview

- ISA-class 32 bit Intel® Pentium® processor
- PCI
- USB
- 12C
- Single core





UEFI for Intel® Quark™ SoC

- First fully open source Intel® Galileo based platform
- Builds on Intel® UDK2014 packages like MdePkg, MdeModulePkg w/ a 32-bit build, adding
 - IA32FamilyCpuBasePkg
 - QuarkPlatformPkg
 - QuarkSocPkg
- Standard build is 1 Mbyte image w/full features
 - Capsule update, SMM, S3, PCI, recovery, full UEFI OS support, FAT OS support, UEFI variables



Intel® Quark™ SoC and Security

Protect Detect
Recover

- Support for I2C-attached TPM
- Hardware Secure Boot option
- UEFI Secure Boot implementation
- UEFI Capsule update support w/ hardware verification assist
- Demonstrates one way to build out UEFI Security Features with a full open source platform tree, with the following summary

	Capsule update	UEFI Secure Boot	TCG Measured Boot	STM	chipsec
MinnowBoard Max	Yes – with open source Capsule driver	Yes	Yes – Integrated TPM	Yes – VT w/o TXT	Yes
Intel® Quark™ Intel® Galileo	Yes – with BootROM support	Yes	Yes – I2C TPM	No	Yes



Summary and Next Steps

- Many security problems, including SMM escalation
- Open source ingredients
- New approach to handle SMM and Testing
- Use open platforms to demonstrate ingredients



Additional Sources of Information

- A PDF of this presentation is available from our Technical Session Catalog: <u>www.intel.com/idfsessionsSF</u>. This URL is also printed on the top of Session Agenda Pages in the Pocket Guide.
- Booth info: #511

More information on security

- UEFI and PI specification http://www.uefi.org
- EDK II Implementation http://www.tianocore.org
- Platform Security information: https://firmware.intel.com/blog/
- EDK II Security Fixes: http://www.tianocore.or/security
- SMM attacks from https://cansecwest.com/agenda.html "Corey Kallenberg & Xeno Kovah, LegbaCore How many million BIOSes would you like to infect?", "Rafal Wojtczuk & Corey Kallenberg Attacks on UEFI Security", "John Loucaides & Andrew Furtak, Intel A new class of vulnerability in SMI Handlers of BIOS/UEFI Firmware"
- STM Specification and code: https://firmware.intel.com/content/smi-transfer-monitor-stm
- CHIPSEC: https://github.com/chipsec/chipsec
- Intel® Quark™ Soc X1000 Version 1.1.0 BIOS https://downloadcenter.intel.com/download/23197/Intel-Quark-BSP
- MinnowMax http://www.minnowboard.org/meet-minnowboard-max/



Other Technical Sessions

	Session ID	Title	Day	Time	Room
✓	STTS001	Firmware in the Data Center: Building a Modern Deployment Framework Using UEFI and Redfish REST APIs	Tue	11:00	2002
√	STTS002	Building a Firmware Component Ecosystem with the Intel® Firmware Engine	Tue	1:15	2002
√	STTS003	Developing Best-in-Class Security Principles with Open Source Firmware	Tue	2:30	2002
	STTS004	Planning and Predicting Big Data Clusters for Spark*, NoSQL and SQL-on- Hadoop* Deployments	Tue	4:00	2002
	STTS005	Accelerating Real-time Analytics Insights with Open Source Software from Intel	Wed	9:30	2002





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Risk Factors

The above statements and any others in this document that refer to plans and expectations for the second quarter, the year and the future are forwardlooking statements that involve a number of risks and uncertainties. Words such as "anticipates," "expects," "intends," "plans," "believes," "seeks," "estimates," "may," "will," "should" and their variations identify forward-looking statements. Statements that refer to or are based on projections, uncertain events or assumptions also identify forward-looking statements. Many factors could affect Intel's actual results, and variances from Intel's current expectations regarding such factors could cause actual results to differ materially from those expressed in these forward-looking statements. Intel presently considers the following to be important factors that could cause actual results to differ materially from the company's expectations. Demand for Intel's products is highly variable and could differ from expectations due to factors including changes in business and economic conditions; consumer confidence or income levels; the introduction, availability and market acceptance of Intel's products, products used together with Intel products and competitors' products; competitive and pricing pressures, including actions taken by competitors; supply constraints and other disruptions affecting customers; changes in customer order patterns including order cancellations; and changes in the level of inventory at customers. Intel's gross margin percentage could vary significantly from expectations based on capacity utilization; variations in inventory valuation, including variations related to the timing of qualifying products for sale; changes in revenue levels; segment product mix; the timing and execution of the manufacturing ramp and associated costs; excess or obsolete inventory; changes in unit costs; defects or disruptions in the supply of materials or resources; and product manufacturing quality/yields. Variations in gross margin may also be caused by the timing of Intel product introductions and related expenses, including marketing expenses, and Intel's ability to respond quickly to technological developments and to introduce new products or incorporate new features into existing products, which may result in restructuring and asset impairment charges. Intel's results could be affected by adverse economic, social, political and physical/infrastructure conditions in countries where Intel, its customers or its suppliers operate, including military conflict and other security risks, natural disasters, infrastructure disruptions, health concerns and fluctuations in currency exchange rates. Results may also be affected by the formal or informal imposition by countries of new or revised export and/or import and doing-business regulations, which could be changed without prior notice. Intel operates in highly competitive industries and its operations have high costs that are either fixed or difficult to reduce in the short term. The amount, timing and execution of Intel's stock repurchase program could be affected by changes in Intel's priorities for the use of cash, such as operational spending, capital spending, acquisitions, and as a result of changes to Intel's cash flows or changes in tax laws. Product defects or errata (deviations from published specifications) may adversely impact our expenses, revenues and reputation. Intel's results could be affected by litigation or regulatory matters involving intellectual property, stockholder, consumer, antitrust, disclosure and other issues. An unfavorable ruling could include monetary damages or an injunction prohibiting Intel from manufacturing or selling one or more products, precluding particular business practices, impacting Intel's ability to design its products, or requiring other remedies such as compulsory licensing of intellectual property. Intel's results may be affected by the timing of closing of acquisitions, divestitures and other significant transactions. A detailed discussion of these and other factors that could affect Intel's results is included in Intel's SEC filings, including the company's most recent reports on Form 10-Q, Form 10-K and earnings release.

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