

Can Tho University

CTU.NegativeZero

Trong Huu Nguyen
Ngoc Thi Pham
Minh Thien Nguyen

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Contents

1 Contest	2		
1.1 C++	2		
1.2 Debug	2		
1.3 Java	2		
1.4 sublime-build	2		
1.5 .bashrc	3		
2 Data structures	3		
2.1 Sparse table	3		
2.2 Ordered set	3		
2.3 Dsu	3		
2.4 Segment tree	4		
2.5 Efficient segment tree	4		
2.6 Persistent lazy segment tree	5		
2.7 Disjoint sparse table	5		
2.8 Fenwick tree	6		
2.9 Fenwick tree 2D	6		
2.10 Implicit treap	7		
3 Mathematics	8		
3.1 Trigonometry	8		
3.2 Sums	8		
4 String	8		
4.1 Prefix function	8		
4.2 Counting occurrences of each prefix	9		
4.3 Knuth–Morris–Pratt algorithm	9		
4.4 Suffix array	9		
4.5 Suffix array slow	10		
4.6 Manacher’s algorithm	10		
4.7 Trie	10		
4.8 Hashing	10		
4.9 Minimum rotation	11		
5 Number Theory	11		
5.1 Euler’s totient function	11		
5.2 Mobius function	12		
5.3 Primes	12		
5.4 Wilson’s theorem	12		
5.5 Zeckendorf’s theorem	12		
5.6 Bitwise operation	13		
		5.7 Pollard’s rho algorithm	13
		5.8 Segment divisor sieve	14
		5.9 Bitset sieve	14
		5.10 Block sieve	14
		5.11 Combinatorics	14
		6 Geometry	15
		6.1 Fundamentals	15
		6.2 Minimum enclosing circle	17
		7 Linear algebra	18
		7.1 Gauss elimination	18
		8 Graph	18
		8.1 Bellman-Ford algorithm	18
		8.2 Articulation point and Bridge	19
		8.3 Strongly connected components	19
		8.4 Topo sort	19
		8.5 K-th smallest shortest path	20
		8.6 Eulerian path	20
		9 Misc.	20
		9.1 Ternary search	20
		9.2 Ternary search 2	21
		9.3 Dutch flag national problem	21
		9.4 Matrix	21

1 Contest

1.1 C++

```

1 #include <bits/stdc++.h>
2 using namespace std;
3
4 #ifdef LOCAL
5 #include "cp/debug.h"
6 #else
7 #define debug(...)
8 #endif
9
10 mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
11
12 const int MOD = (int) 1e9 + 7;
13 const int INF = 0x3f3f3f3f;
14
15 int main() {
16     ios::sync_with_stdio(false); cin.tie(nullptr);
17     // freopen("input.txt", "r", stdin);
18     // freopen("output.txt", "w", stdout);
19
20     return 0;
21 }

```

1.2 Debug

```

1 #define debug(...) { string _s = #__VA_ARGS__; replace(begin(_s), end(_s),
    ', ', ' '); stringstream _ss(_s); istream_iterator<string> _it(_ss);
    out_error(_it, __VA_ARGS__); }
2
3 void out_error(istream_iterator<string> it) { cerr << '\n'; }
4
5 template<typename T, typename ...Args>
6 void out_error(istream_iterator<string> it, T a, Args... args) {
7     cerr << " [" << *it << " = " << a << "]" ";
8     out_error(++it, args...);
9 }
10
11 template<typename T, typename G> ostream& operator<<(ostream &os, const
    pair<T, G> &p) {
12     return os << "(" << p.first << ", " << p.second << ")";
13 }
14
15 template<class Con, class = decltype(begin(declval<Con>()))>
16 typename enable_if<!is_same<Con, string>::value, ostream&>::type
17 operator<<(ostream& os, const Con& container) {
18     os << "{";
19     for (auto it = container.begin(); it != container.end(); ++it)
20         os << (it == container.begin() ? "" : ", ") << *it;
21     return os << "}";
22 }

```

1.3 Java

```

1 import java.io.BufferedReader;
2 import java.util.StringTokenizer;
3 import java.io.IOException;
4 import java.io.InputStreamReader;
5 import java.io.PrintWriter;
6 import java.util.ArrayList;
7 import java.util.Arrays;
8 import java.util.Collections;
9 import java.util.Random;
10
11 public class Main {
12     public static void main(String[] args) {
13         FastScanner fs = new FastScanner();
14         PrintWriter out = new PrintWriter(System.out);
15         int n = fs.nextInt();
16         out.println(n);
17         out.close(); // don't forget this line.
18     }
19     static class FastScanner {
20         BufferedReader br;
21         StringTokenizer st;
22         public FastScanner() {
23             br = new BufferedReader(new InputStreamReader(System.in));
24             st = null;
25         }
26         public String next() {
27             while (st == null || st.hasMoreTokens() == false) {
28                 try {
29                     st = new StringTokenizer(br.readLine());
30                 }
31                 catch (IOException e) {
32                     throw new RuntimeException(e);
33                 }
34             }
35             return st.nextToken();
36         }
37
38         public int nextInt() {
39             return Integer.parseInt(next());
40         }
41
42         public long nextLong() {
43             return Long.parseLong(next());
44         }
45
46         public double nextDouble() {
47             return Double.parseDouble(next());
48         }
49     }
50 }

```

1.4 sublime-build

```

1 {

```

```

2   "cmd": ["g++", "-std=c++17", "-fmax-errors=5", "-DLOCAL", "-Wall",
   "-Wextra", "-o", "${file_path}/${file_base_name}.out", "${file}"],
3   "file_regex": "^(..[:]*):([0-9]+):?([0-9]+)?(?:\\.*)$",
4   "working_dir": "${file_path}",
5   "selector": "source.cpp, source.c++"
6 }

```

1.5 .bashrc

```

1 alias c++='g++ -std=c++2a -fmax-errors=5 -DLOCAL -Wall -Wextra -O2 -s'
2
3 #Stress-testing
4 function test {
5     SOL=$1
6     CHECKER=$2
7     for i in {1..100};
8     do
9         ./gen.out > in && $CHECKER < in > ans && $SOL < in > out && diff -Zb
10        out ans && echo "Test $i passed!!" || break;
11    done
12 }

```

2 Data structures

2.1 Sparse table

```

1 int st[MAXN][K + 1];
2 for (int i = 0; i < N; i++) {
3     st[i][0] = f(array[i]);
4 }
5 for (int j = 1; j <= K; j++) {
6     for (int i = 0; i + (1 << j) <= N; i++) {
7         st[i][j] = f(st[i][j - 1], st[i + (1 << (j - 1))][j - 1]);
8     }
9 }
10 // Range Minimum Queries.
11 int lg[MAXN + 1];
12 lg[1] = 0;
13 for (int i = 2; i <= MAXN; i++) {
14     lg[i] = lg[i / 2] + 1;
15 }
16 int j = lg[R - L + 1];
17 int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);
18 // Range Sum Queries.
19 long long sum = 0;
20 for (int j = K; j >= 0; j--) {
21     if ((1 << j) <= R - L + 1) {
22         sum += st[L][j];
23         L += 1 << j;
24     }
25 }

```

2.2 Ordered set

```

1 #include <ext/pb_ds/assoc_container.hpp>
2 #include <ext/pb_ds/tree_policy.hpp>

```

```

3 using namespace __gnu_pbds;
4
5 template<typename key_type>
6 using set_t = tree<key_type, null_type, less<key_type>, rb_tree_tag,
7     tree_order_statistics_node_update>;
8
9 void example() {
10     vector<int> nums = {1, 2, 3, 5, 10};
11     set_t<int> st(nums.begin(), nums.end());
12
13     cout << *st.find_by_order(0) << '\n'; // 1
14     assert(st.find_by_order(-INF) == st.end());
15     assert(st.find_by_order(INF) == st.end());
16
17     cout << st.order_of_key(2) << '\n'; // 1
18     cout << st.order_of_key(4) << '\n'; // 3
19     cout << st.order_of_key(9) << '\n'; // 4
20     cout << st.order_of_key(-INF) << '\n'; // 0
21     cout << st.order_of_key(INF) << '\n'; // 5
22 }

```

2.3 Dsu

```

1 struct Dsu {
2     int n;
3     vector<int> par, sz;
4     Dsu(int _n) : n(_n) {
5         sz.resize(n, 1);
6         par.resize(n);
7         iota(par.begin(), par.end(), 0);
8     }
9     int find(int v) {
10         // finding leader/parent of set that contains the element v.
11         // with {path compression optimization}.
12         return (v == par[v] ? v : par[v] = find(par[v]));
13     }
14     bool same(int u, int v) {
15         return find(u) == find(v);
16     }
17     bool unite(int u, int v) {
18         u = find(u); v = find(v);
19         if (u == v) return false;
20         if (sz[u] < sz[v]) swap(u, v);
21         par[v] = u;
22         sz[u] += sz[v];
23         return true;
24     }
25     vector<vector<int>> groups() {
26         // returns the list of the "list of the vertices in a connected
27         // component".
28         vector<int> leader(n);
29         for (int i = 0; i < n; ++i) {
30             leader[i] = find(i);
31         }
32     }
33 }

```

```

31     vector<int> id(n, -1);
32     int count = 0;
33     for (int i = 0; i < n; ++i) {
34         if (id[leader[i]] == -1) {
35             id[leader[i]] = count++;
36         }
37     }
38     vector<vector<int>> result(count);
39     for (int i = 0; i < n; ++i) {
40         result[id[leader[i]]].push_back(i);
41     }
42     return result;
43 }
44 };

```

2.4 Segment tree

```

1  /**
2   * Description: A segment tree with range updates and sum queries that
3   * supports three types of operations:
4   * + Increase each value in range [l, r] by x (i.e. a[i] += x).
5   * + Set each value in range [l, r] to x (i.e. a[i] = x).
6   * + Determine the sum of values in range [l, r].
7   */
8  struct SegmentTree {
9      int n;
10     vector<long long> tree, lazy_add, lazy_set;
11     SegmentTree(int _n) : n(_n) {
12         int p = 1;
13         while (p < n) p *= 2;
14         tree.resize(p * 2);
15         lazy_add.resize(p * 2);
16         lazy_set.resize(p * 2);
17     }
18     long long merge(const long long &left, const long long &right) {
19         return left + right;
20     }
21     void build(int id, int l, int r, const vector<int> &arr) {
22         if (l == r) {
23             tree[id] += arr[l];
24             return;
25         }
26         int mid = (l + r) >> 1;
27         build(id * 2, l, mid, arr);
28         build(id * 2 + 1, mid + 1, r, arr);
29         tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
30     }
31     void push(int id, int l, int r) {
32         if (lazy_set[id] == 0 && lazy_add[id] == 0) return;
33         int mid = (l + r) >> 1;
34         for (int child : {id * 2, id * 2 + 1}) {
35             int range = (child == id * 2 ? mid - l + 1 : r - mid);
36             if (lazy_set[id] != 0) {

```

```

37                 lazy_set[child] = lazy_set[id];
38                 tree[child] = range * lazy_set[id];
39             }
40             lazy_add[child] += lazy_add[id];
41             tree[child] += range * lazy_add[id];
42         }
43         lazy_add[id] = lazy_set[id] = 0;
44     }
45 }
46 void update(int id, int l, int r, int u, int v, int amount, bool
47 set_value = false) {
48     if (r < u || l > v) return;
49     if (u <= l && r <= v) {
50         if (set_value) {
51             tree[id] = 1LL * amount * (r - l + 1);
52             lazy_set[id] = amount;
53             lazy_add[id] = 0; // clear all previous updates.
54         }
55         else {
56             tree[id] += 1LL * amount * (r - l + 1);
57             lazy_add[id] += amount;
58         }
59         return;
60     }
61     push(id, l, r);
62     int mid = (l + r) >> 1;
63     update(id * 2, l, mid, u, v, amount, set_value);
64     update(id * 2 + 1, mid + 1, r, u, v, amount, set_value);
65     tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
66 }
67 long long get(int id, int l, int r, int u, int v) {
68     if (r < u || l > v) return 0;
69     if (u <= l && r <= v) {
70         return tree[id];
71     }
72     push(id, l, r);
73     int mid = (l + r) >> 1;
74     long long left = get(id * 2, l, mid, u, v);
75     long long right = get(id * 2 + 1, mid + 1, r, u, v);
76     return merge(left, right);
77 };

```

2.5 Efficient segment tree

```

1  template<typename T> struct SegmentTree {
2      int n;
3      vector<T> tree;
4      SegmentTree(int _n) : n(_n), tree(2 * n) {}
5      T merge(const T &left, const T &right) {
6          return left + right;
7      }
8      template<typename G>
9      void build(const vector<G> &initial) {

```

```

10     assert((int) initial.size() == n);
11     for (int i = 0; i < n; ++i) {
12         tree[i + n] = initial[i];
13     }
14     for (int i = n - 1; i > 0; --i) {
15         tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
16     }
17 }
18 void modify(int i, int v) {
19     tree[i += n] = v;
20     for (i /= 2; i > 0; i /= 2) {
21         tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
22     }
23 }
24 T get_sum(int l, int r) {
25     // sum of elements from l to r - 1.
26     T ret{};
27     for (l += n, r += n; l < r; l /= 2, r /= 2) {
28         if (l & 1) ret = merge(ret, tree[l++]);
29         if (r & 1) ret = merge(ret, tree[--r]);
30     }
31     return ret;
32 }
33 };

```

2.6 Persistent lazy segment tree

```

1 struct Vertex {
2     int l, r;
3     long long val, lazy;
4     bool has_changed = false;
5     Vertex() {}
6     Vertex(int _l, int _r, long long _val, int _lazy = 0) : l(_l), r(_r),
7         val(_val), lazy(_lazy) {}
8 }
9 struct PerSegmentTree {
10     vector<Vertex> tree;
11     vector<int> root;
12     int build(const vector<int> &arr, int l, int r) {
13         if (l == r) {
14             tree.emplace_back(-1, -1, arr[l]);
15             return tree.size() - 1;
16         }
17         int mid = (l + r) / 2;
18         int left = build(arr, l, mid);
19         int right = build(arr, mid + 1, r);
20         tree.emplace_back(left, right, tree[left].val + tree[right].val);
21         return tree.size() - 1;
22     }
23     int add(int x, int l, int r, int u, int v, int amt) {
24         if (l > v || r < u) return x;
25         if (u <= l && r <= v) {
26             tree.emplace_back(tree[x].l, tree[x].r, tree[x].val + 1LL * amt *
27                 (r - l + 1), tree[x].lazy + amt);

```

```

26         tree.back().has_changed = true;
27         return tree.size() - 1;
28     }
29     int mid = (l + r) >> 1;
30     push(x, l, mid, r);
31     int left = add(tree[x].l, l, mid, u, v, amt);
32     int right = add(tree[x].r, mid + 1, r, u, v, amt);
33     tree.emplace_back(left, right, tree[left].val + tree[right].val, 0);
34     return tree.size() - 1;
35 }
36 long long get_sum(int x, int l, int r, int u, int v) {
37     if (r < u || l > v) return 0;
38     if (u <= l && r <= v) return tree[x].val;
39     int mid = (l + r) / 2;
40     push(x, l, mid, r);
41     return get_sum(tree[x].l, l, mid, u, v) + get_sum(tree[x].r, mid + 1,
42         r, u, v);
43 }
44 void push(int x, int l, int mid, int r) {
45     if (!tree[x].has_changed) return;
46     Vertex left = tree[tree[x].l];
47     Vertex right = tree[tree[x].r];
48     tree.emplace_back(left);
49     tree[x].l = tree.size() - 1;
50     tree.emplace_back(right);
51     tree[x].r = tree.size() - 1;
52
53     tree[tree[x].l].val += tree[x].lazy * (mid - l + 1);
54     tree[tree[x].l].lazy += tree[x].lazy;
55
56     tree[tree[x].r].val += tree[x].lazy * (r - mid);
57     tree[tree[x].r].lazy += tree[x].lazy;
58
59     tree[tree[x].l].has_changed = true;
60     tree[tree[x].r].has_changed = true;
61     tree[x].lazy = 0;
62     tree[x].has_changed = false;
63 }

```

2.7 Disjoint sparse table

```

1 /**
2  * Description: range query on a static array.
3  * Time: O(1) per query.
4  * Tested: stress-test.
5  */
6 const int MOD = (int) 1e9 + 7;
7 struct DisjointSparseTable { // product queries.
8     int n, h;
9     vector<vector<int>> dst;
10    vector<int> lg;
11    DisjointSparseTable(int _n) : n(_n) {
12        h = 1; // in case n = 1: h = 0 !!

```

```

13     int p = 1;
14     while (p < n) p *= 2, h++;
15     lg.resize(p); lg[1] = 0;
16     for (int i = 2; i < p; ++i) {
17         lg[i] = 1 + lg[i / 2];
18     }
19     dst.resize(h, vector<int>(n));
20 }
21 void build(const vector<int> &A) {
22     for (int lv = 0; lv < h; ++lv) {
23         int len = (1 << lv);
24         for (int k = 0; k < n; k += len * 2) {
25             int mid = min(k + len, n);
26             dst[lv][mid - 1] = A[mid - 1] % MOD;
27             for (int i = mid - 2; i >= k; --i) {
28                 dst[lv][i] = 1LL * A[i] * dst[lv][i + 1] % MOD;
29             }
30             if (mid == n) break;
31             dst[lv][mid] = A[mid] % MOD;
32             for (int i = mid + 1; i < min(mid + len, n); ++i) {
33                 dst[lv][i] = 1LL * A[i] * dst[lv][i - 1] % MOD;
34             }
35         }
36     }
37 }
38 int get(int l, int r) {
39     if (l == r) {
40         return dst[0][l];
41     }
42     int i = lg[l ^ r];
43     return 1LL * dst[i][l] * dst[i][r] % MOD;
44 }
45 };

```

2.8 Fenwick tree

```

1 /**
2  * Description: range update and range sum query.
3  */
4
5 using tree_type = long long;
6 struct FenwickTree {
7     int n;
8     vector<tree_type> fenw_coeff, fenw;
9     FenwickTree() {}
10    FenwickTree(int _n) : n(_n) {
11        fenw_coeff.assign(n, 0); // fenwick tree with coefficient (n - i).
12        fenw.assign(n, 0); // normal fenwick tree.
13    }
14    template<typename G>
15    void build(const vector<G> &A) {
16        assert((int) A.size() == n);
17        vector<int> diff(n);
18        diff[0] = A[0];

```

```

19        for (int i = 1; i < n; ++i) {
20            diff[i] = A[i] - A[i - 1];
21        }
22        fenw_coeff[0] = (long long) diff[0] * n;
23        fenw[0] = diff[0];
24        for (int i = 1; i < n; ++i) {
25            fenw_coeff[i] = fenw_coeff[i - 1] + (long long) diff[i] * (n - i);
26            fenw[i] = fenw[i - 1] + diff[i];
27        }
28        for (int i = n - 1; i >= 0; --i) {
29            int j = (i & (i + 1)) - 1;
30            if (j >= 0) {
31                fenw_coeff[i] -= fenw_coeff[j];
32                fenw[i] -= fenw[j];
33            }
34        }
35    }
36    void add(vector<tree_type> &fenw, int i, tree_type val) {
37        while (i < n) {
38            fenw[i] += val;
39            i |= (i + 1);
40        }
41    }
42    tree_type __prefix_sum(vector<tree_type> &fenw, int i) {
43        tree_type res{};
44        while (i >= 0) {
45            res += fenw[i];
46            i = (i & (i + 1)) - 1;
47        }
48        return res;
49    }
50    tree_type prefix_sum(int i) {
51        return __prefix_sum(fenw_coeff, i) - __prefix_sum(fenw, i) * (n - i - 1);
52    }
53    void range_add(int l, int r, tree_type val) {
54        add(fenw_coeff, l, (n - l) * val);
55        add(fenw_coeff, r + 1, (n - r - 1) * (-val));
56        add(fenw, l, val);
57        add(fenw, r + 1, -val);
58    }
59    tree_type range_sum(int l, int r) {
60        return prefix_sum(r) - prefix_sum(l - 1);
61    }
62 };

```

2.9 Fenwick tree 2D

```

1 /**
2  * Description: range update and range sum query on a 2D array.
3  */
4 using tree_type = long long;
5 struct FenwickTree2D {
6     int n, m;

```

```

7   vector<vector<tree_type>> > fenw[4];
8   FenwickTree2D(int _n, int _m) : n(_n), m(_m) {
9       for (int i = 0; i < 4; i++) {
10          fenw[i].resize(n, vector<tree_type>(m));
11      }
12  }
13  void add(int u, int v, tree_type val) {
14      for (int i = u; i < n; i |= (i + 1)) {
15          for (int j = v; j < m; j |= (j + 1)) {
16              fenw[0][i][j] += val;
17              fenw[1][i][j] += (u + 1) * val;
18              fenw[2][i][j] += (v + 1) * val;
19              fenw[3][i][j] += (u + 1) * (v + 1) * val;
20          }
21      }
22  }
23  void range_add(int r, int c, int rr, int cc, tree_type val) { // [r, rr]
24      x [c, cc].
25      add(r, c, val);
26      add(r, cc + 1, -val);
27      add(rr + 1, c, -val);
28      add(rr + 1, cc + 1, val);
29  }
30  tree_type prefix_sum(int u, int v) {
31      tree_type res{};
32      for (int i = u; i >= 0; i = (i & (i + 1)) - 1) {
33          for (int j = v; j >= 0; j = (j & (j + 1)) - 1) {
34              res += (u + 2) * (v + 2) * fenw[0][i][j];
35              res -= (v + 2) * fenw[1][i][j];
36              res -= (u + 2) * fenw[2][i][j];
37              res += fenw[3][i][j];
38          }
39      }
40      return res;
41  }
42  tree_type range_sum(int r, int c, int rr, int cc) { // [r, rr] x [c, cc].
43      return prefix_sum(rr, cc) - prefix_sum(r - 1, cc) - prefix_sum(rr, c
44      - 1) + prefix_sum(r - 1, c - 1);
45  }
46  };

```

2.10 Implicit treap

```

1  struct Node {
2      int val, prior, cnt;
3      bool rev;
4      Node *left, *right;
5      Node() {}
6      Node(int _val) : val(_val), prior(rng()), cnt(1), rev(false),
7      left(nullptr), right(nullptr) {}
8  };
9  // Binary search tree + min-heap.
10 struct Treap {
11     Node *root;

```

```

11 Treap() : root(nullptr) {}
12 int get_cnt(Node *n) { return n ? n->cnt : 0; }
13 void upd_cnt(Node *n) {
14     if (n) n->cnt = get_cnt(n->left) + get_cnt(n->right) + 1;
15 }
16 void push_rev(Node *treap) {
17     if (!treap || !treap->rev) return;
18     treap->rev = false;
19     swap(treap->left, treap->right);
20     if (treap->left) treap->left->rev ^= true;
21     if (treap->right) treap->right->rev ^= true;
22 }
23 pair<Node*, Node*> split(Node *treap, int x, int smaller = 0) {
24     if (!treap) return {};
25     push_rev(treap);
26     int idx = smaller + get_cnt(treap->left); // implicit val.
27     if (idx <= x) {
28         auto pr = split(treap->right, x, idx + 1);
29         treap->right = pr.first;
30         upd_cnt(treap);
31         return {treap, pr.second};
32     }
33     else {
34         auto pl = split(treap->left, x, smaller);
35         treap->left = pl.second;
36         upd_cnt(treap);
37         return {pl.first, treap};
38     }
39 }
40 Node* merge(Node *l, Node *r) {
41     push_rev(l); push_rev(r);
42     if (!l || !r) return (l ? l : r);
43     if (l->prior < r->prior) {
44         l->right = merge(l->right, r);
45         upd_cnt(l);
46         return l;
47     }
48     else {
49         r->left = merge(l, r->left);
50         upd_cnt(r);
51         return r;
52     }
53 }
54 void insert(int pos, int val) {
55     if (!root) {
56         root = new Node(val);
57         return;
58     }
59     Node *l, *m, *r;
60     m = new Node(val);
61     tie(l, r) = split(root, pos - 1);
62     root = merge(l, merge(m, r));
63 }

```



```

64 void erase(int pos_l, int pos_r) {
65     Node *l, *m, *r;
66     tie(l, r) = split(root, pos_l - 1);
67     tie(m, r) = split(r, pos_r - pos_l);
68     root = merge(l, r);
69 }
70 void reverse(int pos_l, int pos_r) {
71     Node *l, *m, *r;
72     tie(l, r) = split(root, pos_l - 1);
73     tie(m, r) = split(r, pos_r - pos_l);
74     m->rev ^= true;
75     root = merge(l, merge(m, r));
76 }
77 int query(int pos_l, int pos_r);
78 // returns answer for corresponding types of query.
79 void inorder(Node *n) {
80     if (!n) return;
81     push_rev(n);
82     inorder(n->left);
83     cout << n->val << ' ';
84     inorder(n->right);
85 }
86 void print() {
87     inorder(root);
88     cout << '\n';
89 }
90 };

```

3 Mathematics

3.1 Trigonometry

3.1.1 Sum - difference identities

$$\sin(u \pm v) = \sin(u) \cos(v) \pm \cos(u) \sin(v)$$

$$\cos(u \pm v) = \cos(u) \cos(v) \mp \sin(u) \sin(v)$$

$$\tan(u \pm v) = \frac{\tan(u) \pm \tan(v)}{1 \mp \tan(u) \tan(v)}$$

3.1.2 Sum to product identities

$$\cos(u) + \cos(v) = 2 \cos\left(\frac{u+v}{2}\right) \cos\left(\frac{u-v}{2}\right)$$

$$\cos(u) - \cos(v) = -2 \sin\left(\frac{u+v}{2}\right) \sin\left(\frac{u-v}{2}\right)$$

$$\sin(u) + \sin(v) = 2 \sin\left(\frac{u+v}{2}\right) \cos\left(\frac{u-v}{2}\right)$$

$$\sin(u) - \sin(v) = 2 \cos\left(\frac{u+v}{2}\right) \sin\left(\frac{u-v}{2}\right)$$

3.1.3 Product identities

$$\cos(u) \cos(v) = \frac{1}{2} [\cos(u+v) + \cos(u-v)]$$

$$\sin(u) \sin(v) = -\frac{1}{2} [\cos(u+v) - \cos(u-v)]$$

$$\sin(u) \cos(v) = \frac{1}{2} [\sin(u+v) + \sin(u-v)]$$

3.1.4 Double - triple angle identities

$$\sin(2u) = 2 \sin(u) \cos(u)$$

$$\cos(2u) = 2 \cos^2(u) - 1 = 1 - 2 \sin^2(u)$$

$$\tan(2u) = \frac{2 \tan(u)}{1 - \tan^2(u)}$$

$$\sin(3u) = 3 \sin(u) - 4 \sin^3(u)$$

$$\cos(3u) = 4 \cos^3(u) - 3 \cos(u)$$

$$\tan(3u) = \frac{3 \tan(u) - \tan^3(u)}{1 - 3 \tan^2(u)}$$

3.2 Sums

$$c^a + c^{a+1} + \dots + c^b = \frac{c^{b+1} - c^a}{c - 1}, \quad c \neq 1$$

$$c + 2c^2 + \dots + nc^n = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^2}, \quad c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$$

$$1^3 + 2^3 + 3^3 + \dots + n^3 = \left(\frac{n(n+1)}{2}\right)^2$$

$$1^4 + 2^4 + 3^4 + \dots + n^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

$$1^5 + 2^5 + 3^5 + \dots + n^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12}$$

$$1^6 + 2^6 + 3^6 + \dots + n^6 = \frac{n(n+1)(2n+1)(3n^4+6n^3-3n+1)}{42}$$

$$1^7 + 2^7 + 3^7 + \dots + n^7 = \frac{n^2(n+1)^2(3n^4+6n^3-n^2-4n+2)}{24}$$

4 String

4.1 Prefix function

```

1 /**
2  * Description: The prefix function of a string 's' is defined as an array pi
3  * of length n,
4  * where pi[i] is the length of the longest proper prefix of the substring
5  * s[0..i] which is also a suffix of this substring.
6  * Time complexity: O(|S|).
7  */
8 vector<int> prefix_function(const string &s) {
9     int n = (int) s.length();
10    vector<int> pi(n);
11    pi[0] = 0;
12    for (int i = 1; i < n; ++i) {
13        int j = pi[i - 1]; // try length pi[i - 1] + 1.
14        while (j > 0 && s[j] != s[i]) {
15            j = pi[j - 1];
16        }
17        if (s[j] == s[i]) {
18            pi[i] = j + 1;
19        }
20    }
21    return pi;
22 }

```

4.2 Counting occurrences of each prefix

```

1 vector<int> count_occurrences(const string &s) {
2     vector<int> pi = prefix_function(s);
3     int n = (int) s.size();
4     vector<int> ans(n + 1);
5     for (int i = 0; i < n; ++i) {
6         ans[pi[i]]++;
7     }
8     for (int i = n - 1; i > 0; --i) {
9         ans[pi[i - 1]] += ans[i];
10    }
11    for (int i = 0; i <= n; ++i) {
12        ans[i]++;
13    }
14    return ans;
15    // Input: ABACABA
16    // Output: 4 2 2 1 1 1 1
17 }

```

4.3 Knuth–Morris–Pratt algorithm

```

1 /**
2  * Searching for a substring in a string.
3  * Time complexity: O(N + M).
4  */
5 vector<int> KMP(const string &text, const string &pattern) {
6     int n = (int) text.length();
7     int m = (int) pattern.length();
8     string s = pattern + '$' + text;
9     vector<int> pi = prefix_function(s);
10    vector<int> indices;

```

```

11    for (int i = 0; i < (int) s.length(); ++i) {
12        if (pi[i] == m) {
13            indices.push_back(i - 2 * m);
14        }
15    }
16    return indices;
17 }

```

4.4 Suffix array

```

1 struct SuffixArray {
2     string s;
3     int n, lim;
4     vector<int> sa, lcp, rank;
5     SuffixArray(const string &s, int _lim = 256) : s(s), n(s.length() + 1),
6         lim(_lim), sa(n), lcp(n), rank(n) {
7         s += '$';
8         build();
9         kasai();
10        sa.erase(sa.begin());
11        lcp.erase(lcp.begin());
12        s.pop_back();
13    }
14    void build() {
15        vector<int> nrank(n), norder(n), cnt(max(n, lim));
16        for (int i = 0; i < n; ++i) {
17            sa[i] = i; rank[i] = s[i];
18        }
19        for (int k = 0, rank_cnt = 0; rank_cnt < n - 1; k = max(1, k * 2),
20            lim = rank_cnt + 1) {
21            // counting sort.
22            for (int i = 0; i < n; ++i) norder[i] = (sa[i] - k + n) % n;
23            for (int i = 0; i < n; ++i) cnt[rank[i]]++;
24            for (int i = 1; i < lim; ++i) cnt[i] += cnt[i - 1];
25            for (int i = n - 1; i >= 0; --i) sa[--cnt[rank[norder[i]]]] =
26                norder[i];
27            rank[sa[0]] = rank_cnt = 0;
28            for (int i = 1; i < n; ++i) {
29                int u = sa[i], v = sa[i - 1];
30                int nu = u + k, nv = v + k;
31                if (nu >= n) nu -= n;
32                if (nv >= n) nv -= n;
33                if (rank[u] != rank[v] || rank[nu] != rank[nv]) ++rank_cnt;
34                nrank[sa[i]] = rank_cnt;
35            }
36            for (int i = 0; i < rank_cnt + 1; ++i) cnt[i] = 0;
37            rank.swap(nrank);
38        }
39    }
40    void kasai() {
41        for (int i = 0; i < n; ++i) rank[sa[i]] = i;
42        for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
43            int j = sa[rank[i] - 1];
44            while (s[i + k] == s[j + k]) k++;

```

```

42         lcp[rank[i]] = k;
43     }
44     // Note: lcp[i] = longest common prefix(sa[i - 1], sa[i]).
45 }
46 };

```

4.5 Suffix array slow

```

1 vector<int> suffix_array_slow(string &s) {
2     s += '$';
3     int n = (int) s.size();
4     vector<int> order(n), rank(n);
5     for (int i = 0; i < n; ++i) {
6         order[i] = i; rank[i] = s[i];
7     }
8     for (int k = 0; k < n; k = max(1, k * 2)) {
9         stable_sort(sa.begin(), sa.end(), [&](int i, int j) {
10             return make_pair(rank[i], rank[(i + k) % n]) < make_pair(rank[j],
11                 rank[(j + k) % n]);
12         });
13         vector<int> nrank(n);
14         for (int i = 0, cnt = 0; i < n; ++i) {
15             if (i > 0 && rank[order[i]] != rank[order[i - 1]]) ++cnt;
16             else if (i > 0 && rank[(order[i] + k) % n] != rank[(order[i - 1]
17                 + k) % n]) ++cnt;
18             nrank[order[i]] = cnt;
19         }
20         rank.swap(nrank);
21     }
22     s.pop_back(); order.erase(order.begin());
23     return order;
24 }
// Time complexity: O(N * log(N)^2).

```

4.6 Manacher's algorithm

```

1 /**
2  * Description: for each position, computes d[0][i] = half length of
3  * longest palindrome centered on i (rounded up), d[1][i] = half length of
4  * longest palindrome centered on i and i - 1.
5  * Time complexity: O(N).
6  * Tested: https://judge.yosupo.jp/problem/enumerate\_palindromes,
7  * stress-tested.
8  */
9 array<vector<int>, 2> manacher(const string &s) {
10     int n = (int) s.size();
11     array<vector<int>, 2> d;
12     for (int z = 0; z < 2; ++z) {
13         d[z].resize(n);
14         int l = 0, r = 0;
15         for (int i = 0; i < n; ++i) {
16             int mirror = l + r - i + z;
17             d[z][i] = (i > r ? 0 : min(d[z][mirror], r - i));
18             int L = i - d[z][i] - z, R = i + d[z][i];
19             while (L >= 0 && R < n && s[L] == s[R]) {

```

```

19         d[z][i]++; L--; R++;
20     }
21     if (R > r) {
22         l = L; r = R;
23     }
24 }
25 }
26 return d;
27 }

```

4.7 Trie

```

1 struct Trie {
2     const static int ALPHABET = 26;
3     const static char minChar = 'a';
4     struct Vertex {
5         int next[ALPHABET];
6         bool leaf;
7         Vertex() {
8             leaf = false;
9             fill(next, next + ALPHABET, -1);
10        }
11    };
12    vector<Vertex> trie;
13    Trie() { trie.emplace_back(); }
14
15    void insert(const string &s) {
16        int i = 0;
17        for (const char &ch : s) {
18            int j = ch - minChar;
19            if (trie[i].next[j] == -1) {
20                trie[i].next[j] = trie.size();
21                trie.emplace_back();
22            }
23            i = trie[i].next[j];
24        }
25        trie[i].leaf = true;
26    }
27    bool find(const string &s) {
28        int i = 0;
29        for (const char &ch : s) {
30            int j = ch - minChar;
31            if (trie[i].next[j] == -1) {
32                return false;
33            }
34            i = trie[i].next[j];
35        }
36        return (trie[i].leaf ? true : false);
37    }
38 };

```

4.8 Hashing

```

1 struct Hash61 {
2     static const uint64_t MOD = (1LL << 61) - 1;

```

```

3  static uint64_t BASE;
4  static vector<uint64_t> pw;
5  uint64_t addmod(uint64_t a, uint64_t b) const {
6      a += b;
7      if (a >= MOD) a -= MOD;
8      return a;
9  }
10 uint64_t submod(uint64_t a, uint64_t b) const {
11     a += MOD - b;
12     if (a >= MOD) a -= MOD;
13     return a;
14 }
15 uint64_t mulmod(uint64_t a, uint64_t b) const {
16     uint64_t low1 = (uint32_t) a, high1 = (a >> 32);
17     uint64_t low2 = (uint32_t) b, high2 = (b >> 32);
18
19     uint64_t low = low1 * low2;
20     uint64_t mid = low1 * high2 + low2 * high1;
21     uint64_t high = high1 * high2;
22
23     uint64_t ret = (low & MOD) + (low >> 61) + (high << 3) + (mid >> 29)
24 + (mid << 35 >> 3) + 1;
25     // ret %= MOD;
26     ret = (ret >> 61) + (ret & MOD);
27     ret = (ret >> 61) + (ret & MOD);
28     return ret - 1;
29 }
30 void ensure_pw(int m) {
31     int n = (int) pw.size();
32     if (n >= m) return;
33     pw.resize(m);
34     for (int i = n; i < m; ++i) {
35         pw[i] = mulmod(pw[i - 1], BASE);
36     }
37 }
38 vector<uint64_t> pref;
39 int n;
40 template<typename T> Hash61(const T &s) { // strings or arrays.
41     n = (int) s.size();
42     ensure_pw(n);
43     pref.resize(n + 1);
44     pref[0] = 0;
45     for (int i = 0; i < n; ++i) {
46         pref[i + 1] = addmod(mulmod(pref[i], BASE), s[i]);
47     }
48 }
49 inline uint64_t operator()(const int from, const int to) const {
50     assert(0 <= from && from <= to && to < n);
51     // pref[to + 1] - pref[from] * pw[to - from + 1]
52     return submod(pref[to + 1], mulmod(pref[from], pw[to - from + 1]));
53 }
54 };

```

```

55 mt19937 rng((unsigned int)
56     chrono::steady_clock::now().time_since_epoch().count());
57 uint64_t Hash61::BASE = (MOD >> 2) + rng() % (MOD >> 1);
58 vector<uint64_t> Hash61::pw = vector<uint64_t>(1, 1);

```

4.9 Minimum rotation

```

1  /**
2   * Author: Stjepan Glavina
3   * License: Unlicense
4   * Source: https://github.com/stjepang/snippets/blob/master/min_rotation.cpp
5   * Description: Finds the lexicographically smallest rotation of a string.
6   * Time: O(N)
7   * Usage:
8   * rotate(v.begin(), v.begin()+minRotation(v), v.end());
9   * Status: Stress-tested
10  */
11 #pragma once
12
13 int minRotation(string s) {
14     int a = 0, N = (int) s.size(); s += s;
15     rep(b, 0, N) rep(k, 0, N) {
16         if (a + k == b || s[a + k] < s[b + k]) {b += max(0, k - 1); break;}
17         if (s[a + k] > s[b + k]) {a = b; break;}
18     }
19     return a;
20 }

```

5 Number Theory

5.1 Euler's totient function

- Euler's totient function, also known as **phi-function** $\phi(n)$ counts the number of integers between 1 and n inclusive, that are **coprime to** n .

- Properties:

– Divisor sum property: $\sum_{d|n} \phi(d) = n$.

– $\phi(n)$ is a **prime number** when $n = 3, 4, 6$.

– If p is a prime number, then $\phi(p) = p - 1$.

– If p is a prime number and $k \geq 1$, then $\phi(p^k) = p^k - p^{k-1}$.

– If a and b are **coprime**, then $\phi(ab) = \phi(a) \cdot \phi(b)$.

– In general, for **not coprime** a and b , with $d = \gcd(a, b)$ this equation holds:

$$\phi(ab) = \phi(a) \cdot \phi(b) \cdot \frac{d}{\phi(d)}.$$

– With $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\begin{aligned}\phi(n) &= \phi(p_1^{k_1}) \cdot \phi(p_2^{k_2}) \cdots \phi(p_m^{k_m}) \\ &= n \cdot \left(1 - \frac{1}{p_1}\right) \cdot \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_m}\right)\end{aligned}$$

• Application in Euler's theorem:

– If $\gcd(a, M) = 1$, then:

$$a^{\phi(M)} \equiv 1 \pmod{M} \Rightarrow a^n \equiv a^{n \bmod \phi(M)} \pmod{M}$$

– In general, for arbitrary a, M and $n \geq \log_2 M$:

$$a^n \equiv a^{\phi(M) + [n \bmod \phi(M)]} \pmod{M}$$

5.2 Mobius function

• For a positive integer $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\mu(n) = \begin{cases} 1, & \text{if } n = 1 \\ 0, & \text{if } \exists k_i > 1 \\ (-1)^m & \text{otherwise} \end{cases}$$

• Properties:

$$\sum_{d|n} \mu(d) = [n = 1].$$

– If a and b are **coprime**, then $\mu(ab) = \mu(a) \cdot \mu(b)$.

– Mobius inversion: let f and g be arithmetic functions:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right) g(d)$$

5.3 Primes

Approximating the number of primes up to n :

n	$\pi(n)$	$\frac{n}{\ln n - 1}$
100 ($1e^2$)	25	28
500 ($5e^2$)	95	96
1000 ($1e^3$)	168	169
5000 ($5e^3$)	669	665
10000 ($1e^4$)	1229	1218
50000 ($5e^4$)	5133	5092
100000 ($1e^5$)	9592	9512
500000 ($5e^5$)	41538	41246
1000000 ($1e^6$)	78498	78030
5000000 ($5e^6$)	348513	346622

($\pi(n)$ = the number of primes less than or equal to n , $\frac{n}{\ln n - 1}$ is used to approximate $\pi(n)$).

5.4 Wilson's theorem

A positive integer n is a prime if and only if:

$$(n - 1)! \equiv n - 1 \pmod{n}$$

5.5 Zeckendorf's theorem

The Zeckendorf's theorem states that every positive integer n can be represented uniquely as a sum of one or more distinct non-consecutive Fibonacci numbers. For example:

$$64 = 55 + 8 + 1$$

$$85 = 55 + 21 + 8 + 1$$

```
1 vector<int> zeckendofth_theorem(int n) {
2     vector<int> fibs = {1, 1};
3     int sz = 2;
4     while (fibs.back() <= n) {
5         fibs.push_back(fibs[sz - 1] + fibs[sz - 2]);
6         sz++;
7     }
8     fibs.pop_back();
9     vector<int> nums;
10    int p = sz - 1;
11    while (n > 0) {
12        if (n >= fibs[p]) {
13            nums.push_back(fibs[p]);
14            n -= fibs[p];
15        }
16        p--;
17    }
18    return nums;
19 }
```

5.6 Bitwise operation

- $a + b = (a \oplus b) + 2(a \& b)$
- $a | b = (a \oplus b) + (a \& b)$
- $a \& (b \oplus c) = (a \& b) \oplus (a \& c)$
- $a | (b \& c) = (a | b) \& (a | c)$
- $a \& (b | c) = (a \& b) | (a \& c)$
- $a | (a \& b) = a$
- $a \& (a | b) = a$
- $n = 2^k \Leftrightarrow (n \& (n - 1)) = 1$
- $-a = \sim a + 1$
- $4i \oplus (4i + 1) \oplus (4i + 2) \oplus (4i + 3) = 0$
- Iterating over all subsets of a set and iterating over all submasks of a mask:

```

1 for (int mask = 0; mask < (1 << n); ++mask) {
2     for (int i = 0; i < n; ++i) {
3         if (mask & (1 << i)) {
4             // do something...
5         }
6     }
7     // Time complexity: O(n * 2^n).
8 }
9 for (int mask = 0; mask < (1 << n); ++mask) {
10    for (int submask = mask; ; submask = (submask - 1) & mask) {
11        // do something...
12        if (submask == 0) break;
13    }
14    // Time complexity: O(3^n).
15 }

```

5.7 Pollard's rho algorithm

```

1 using num_t = long long;
2 const int PRIME_MAX = (int) 4e4; // for handle numbers <= 1e9.
3 const int LIMIT = (int) 1e9;
4 vector<int> primes;
5 int small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 73, 113,
6     193, 311, 313, 407521, 299210837};
7 void linear_sieve(int n);
8 num_t mulmod(num_t a, num_t b, num_t mod);
9 num_t powmod(num_t a, num_t n, num_t mod);
10 bool miller_rabin(num_t a, num_t d, int s, num_t mod) {
11     num_t x = powmod(a, d, mod);
12     if (x == mod - 1 || x == 1) {
13         return true;
14     }
15     for (int i = 0; i < s - 1; ++i) {
16         x = mulmod(x, x, mod);
17         if (x == mod - 1) return true;
18     }
19     return false;
20 }
21 bool is_prime(num_t n, int tests = 10) {
22     if (n < 4) return (n > 1);
23     num_t d = n - 1;
24     int s = 0;
25     while (d % 2 == 0) { d >>= 1; s++; }

```

```

25     for (int i = 0; i < tests; ++i) {
26         int a = small_primes[i];
27         if (n == a) return true;
28         if (n % a == 0 || !miller_rabin(a, d, s, n)) return false;
29     }
30     return true;
31 }
32 num_t f(num_t x, int c, num_t mod) { // f(x) = (x^2 + c) % mod.
33     x = mulmod(x, x, mod);
34     x += c;
35     if (x >= mod) x -= mod;
36     return x;
37 }
38 num_t pollard_rho(num_t n, int c) {
39     // algorithm to find a random divisor of 'n'.
40     // using random function: f(x) = (x^2 + c) % n.
41     num_t x = 2, y = x, d;
42     long long p = 1;
43     int dist = 0;
44     while (true) {
45         y = f(y, c, n);
46         dist++;
47         d = __gcd(llabs(x - y), n);
48         if (d > 1) break;
49         if (dist == p) { dist = 0; p *= 2; x = y; }
50     }
51     return d;
52 }
53 void factorize(int n, vector<num_t> &factors);
54 void llfactorize(num_t n, vector<num_t> &factors) {
55     if (n < 2) return;
56     if (is_prime(n)) {
57         factors.emplace_back(n);
58         return;
59     }
60     if (n < LIMIT) {
61         factorize(n, factors);
62         return;
63     }
64     num_t d = n;
65     for (int c = 2; d == n; c++) {
66         d = pollard_rho(n, c);
67     }
68     llfactorize(d, factors);
69     llfactorize(n / d, factors);
70 }
71 vector<num_t> gen_divisors(vector<pair<num_t, int>> &factors) {
72     vector<num_t> divisors = {1};
73     for (auto &x : factors) {
74         int sz = (int) divisors.size();
75         for (int i = 0; i < sz; ++i) {
76             num_t cur = divisors[i];
77             for (int j = 0; j < x.second; ++j) {

```

```

78         cur *= x.first;
79         divisors.push_back(cur);
80     }
81 }
82 }
83 return divisors; // this array is NOT sorted yet.
84 }

```

5.8 Segment divisor sieve

```

1  const int MAXN = (int) 1e6; // R - L + 1 <= N.
2  int divisor_count[MAXN + 3];
3  void segment_divisor_sieve(long long L, long long R) {
4      for (long long i = 1; i <= (long long) sqrt(R); ++i) {
5          long long start1 = ((L + i - 1) / i) * i;
6          long long start2 = i * i;
7          long long j = max(start1, start2);
8          if (j == start2) {
9              divisor_count[j - L] += 1;
10             j += i;
11         }
12         for (; j <= R; j += i) {
13             divisor_count[j - L] += 2;
14         }
15     }
16 }

```

5.9 Bitset sieve

```

1  /**
2   * Description: sieve of eratosthenes for large n (up to 1e9).
3   * Time and space (tested on codeforces):
4   * + For n = 1e8: ~200 ms, 6 MB.
5   * + For n = 1e9: ~4000 ms, 60 MB.
6   */
7  const int N = (int) 1e8;
8  bitset<N / 2 + 1> isPrime;
9  void sieve(int n = N) {
10     isPrime.flip();
11     isPrime[0] = false;
12     for (int i = 3; i <= (int) sqrt(n); i += 2) {
13         if (isPrime[i >> 1]) {
14             for (int j = i * i; j <= n; j += 2 * i) {
15                 isPrime[j >> 1] = false;
16             }
17         }
18     }
19 }
20 void example(int n) {
21     sieve(n);
22     int primeCnt = (n >= 2);
23     for (int i = 3; i <= n; i += 2) {
24         if (isPrime[i >> 1]) {
25             primeCnt++;
26         }
27     }
28 }

```

```

27     }
28     cout << primeCnt << '\n';
29 }

```

5.10 Block sieve

```

1  /**
2   * Description: very fast sieve of eratosthenes for large n (up to 1e9).
3   * Source: kactl.
4   * Time and space (tested on codeforces):
5   * + For n = 1e8: ~160 ms, 60 MB.
6   * + For n = 1e9: ~1600 ms, 505 MB.
7   * Need to check memory limit.
8   */
9  const int N = (int) 1e8;
10 bitset<N + 1> is_prime;
11 vector<int> fast_sieve() {
12     const int S = (int) sqrt(N), R = N / 2;
13     vector<int> primes = {2};
14     vector<bool> sieve(S + 1, true);
15     vector<array<int, 2>> cp;
16     for (int i = 3; i <= S; i += 2) {
17         if (sieve[i]) {
18             cp.push_back({i, i * i / 2});
19             for (int j = i * i; j <= S; j += 2 * i) {
20                 sieve[j] = false;
21             }
22         }
23     }
24     for (int L = 1; L <= R; L += S) {
25         array<bool, S> block{};
26         for (auto &[p, idx] : cp) {
27             for (; idx < S + L; idx += p) block[idx - L] = true;
28         }
29         for (int i = 0; i < min(S, R - L); ++i) {
30             if (!block[i]) primes.push_back((L + i) * 2 + 1);
31         }
32     }
33     for (int p : primes) is_prime[p] = true;
34     return primes;
35 }

```

5.11 Combinatorics

5.11.1 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{n!(n+1)!}$$

$$C_{n+1} = \sum_{i=0}^n C_i C_{n-i}, \quad C_0 = 1, \quad C_n = \frac{4n-2}{n+1} C_{n-1}$$

- The first 12 Catalan numbers ($n = 0, 1, 2, \dots, 11$):

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$$

- Applications of Catalan numbers:
 - difference binary search trees with n vertices from 1 to n .
 - rooted binary trees with $n + 1$ leaves (vertices are not numbered).
 - correct bracket sequence of length $2 * n$.
 - permutation $[n]$ with no 3-term increasing subsequence (i.e. doesn't exist $i < j < k$ for which $a[i] < a[j] < a[k]$).
 - ways a convex polygon of $n + 2$ sides can split into triangles by connecting vertices.

5.11.2 Fibonacci numbers

$$F_n = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F_{n-1} + F_{n-2}, & \text{otherwise} \end{cases}$$

- The first 20 Fibonacci numbers ($n = 0, 1, 2, \dots, 19$):

$$F_n = 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, 2584, 4181$$

- Properties:

$$\left. \begin{aligned} F_{2n+1} &= F_n^2 + F_{n+1}^2 \\ F_{2n} &= F_{n-1} \cdot F_n + F_n \cdot F_{n+1} \\ F_{n+1} \cdot F_{n-1} - F_n^2 &= (-1)^n \end{aligned} \right| \begin{aligned} n \mid m &\iff F_n \mid F_m \\ (F_n, F_m) &= F_{(n,m)} \end{aligned}$$

5.11.3 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k non-empty groups.

$$S(n, k) = S(n-1, k-1) + kS(n-1, k)$$

$$S(n, 1) = S(n, n) = 1$$

$$S(n, k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} \binom{k}{i} i^n$$

5.11.4 Derangements

Permutation of the elements of a set, such that no element appears in its original position (no fixed point). Recursive formulas:

$$D(n) = (n-1)[D(n-1) + D(n-2)] = nD(n-1) + (-1)^n$$

6 Geometry

6.1 Fundamentals

6.1.1 Point

```

1  const double PI = acos(-1);
2  const double EPS = 1e-9;
3  typedef double ftype;
4  struct point {
5      ftype x, y;
6      point(ftype _x = 0, ftype _y = 0): x(_x), y(_y) {}
7      point& operator+=(const point& other) {
8          x += other.x; y += other.y; return *this;
9      }
10     point& operator-=(const point& other) {
11         x -= other.x; y -= other.y; return *this;
12     }
13     point& operator*=(ftype t) {
14         x *= t; y *= t; return *this;
15     }
16     point& operator/=(ftype t) {
17         x /= t; y /= t; return *this;
18     }
19     point operator+(const point& other) const {
20         return point(*this) += other;
21     }
22     point operator-(const point& other) const {
23         return point(*this) -= other;
24     }
25     point operator*(ftype t) const {
26         return point(*this) *= t;
27     }
28     point operator/(ftype t) const {
29         return point(*this) /= t;
30     }
31     point rotate(double angle) const {
32         return point(x * cos(angle) - y * sin(angle), x * sin(angle) + y *
33             cos(angle));
34     }
35     friend istream& operator>>(istream& in, point& t);
36     friend ostream& operator<<(ostream& out, const point& t);
37     bool operator<(const point& other) const {
38         if (fabs(x - other.x) < EPS)
39             return y < other.y;
40         return x < other.x;
41     };
42
43     istream& operator>>(istream& in, point& t) {
44         in >> t.x >> t.y;
45         return in;
46     }
47     ostream& operator<<(ostream& out, const point& t) {
48         out << t.x << ' ' << t.y;

```



```

49     return out;
50 }
51
52 ftype dot(point a, point b) {return a.x * b.x + a.y * b.y;}
53 ftype norm(point a) {return dot(a, a);}
54 ftype abs(point a) {return sqrt(norm(a));}
55 ftype angle(point a, point b) {return acos(dot(a, b) / (abs(a) * abs(b)));}
56 ftype proj(point a, point b) {return dot(a, b) / abs(b);}
57 ftype cross(point a, point b) {return a.x * b.y - a.y * b.x;}
58 bool ccw(point a, point b, point c) {return cross(b - a, c - a) > EPS;}
59 bool collinear(point a, point b, point c) {return fabs(cross(b - a, c - a)) <
    EPS;}
60 point intersect(point a1, point d1, point a2, point d2) {
61     double t = cross(a2 - a1, d2) / cross(d1, d2);
62     return a1 + d1 * t;
63 }

```

6.1.2 Line

```

1 struct line {
2     double a, b, c;
3     line (double _a = 0, double _b = 0, double _c = 0): a(_a), b(_b), c(_c) {}
4     friend ostream & operator<<(ostream& out, const line& l);
5 };
6 ostream & operator<<(ostream& out, const line& l) {
7     out << l.a << ' ' << l.b << ' ' << l.c;
8     return out;
9 }
10 void pointsToLine(const point& p1, const point& p2, line& l) {
11     if (fabs(p1.x - p2.x) < EPS)
12         l = {1.0, 0.0, -p1.x};
13     else {
14         l.a = - (double)(p1.y - p2.y) / (p1.x - p2.x);
15         l.b = 1.0;
16         l.c = - l.a * p1.x - l.b * p1.y;
17     }
18 }
19 void pointsSlopeToLine(const point& p, double m, line& l) {
20     l.a = -m;
21     l.b = 1;
22     l.c = -l.a * p.x - l.b * p.y;
23 }
24 bool areParallel(const line& l1, const line& l2) {
25     return fabs(l1.a - l2.a) < EPS && fabs(l1.b - l2.b) < EPS;
26 }
27 bool areSame(const line& l1, const line& l2) {
28     return areParallel(l1, l2) && fabs(l1.c - l2.c) < EPS;
29 }
30 bool areIntersect(line l1, line l2, point& p) {
31     if (areParallel(l1, l2)) return false;
32     p.x = - (l1.c * l2.b - l1.b * l2.c) / (l1.a * l2.b - l1.b * l2.a);
33     if (fabs(l1.b) > EPS) p.y = - (l1.c + l1.a * p.x);
34     else p.y = - (l2.c + l2.a * p.x);
35     return 1;
36 }

```

```

37 double distToLine(point p, point a, point b, point& c) {
38     double t = dot(p - a, b - a) / norm(b - a);
39     c = a + (b - a) * t;
40     return abs(c - p);
41 }
42 double distToSegment(point p, point a, point b, point& c) {
43     double t = dot(p - a, b - a) / norm(b - a);
44     if (t > 1.0)
45         c = point(b.x, b.y);
46     else if (t < 0.0)
47         c = point(a.x, a.y);
48     else
49         c = a + (b - a) * t;
50     return abs(c - p);
51 }
52 bool intersectTwoSegment(point a, point b, point c, point d) {
53     ftype ABxAC = cross(b - a, c - a);
54     ftype ABxAD = cross(b - a, d - a);
55     ftype CDxCA = cross(d - c, a - c);
56     ftype CDxCB = cross(d - c, b - c);
57     if (ABxAC == 0 || ABxAD == 0 || CDxCA == 0 || CDxCB == 0) {
58         if (ABxAC == 0 && dot(a - c, b - c) <= 0) return true;
59         if (ABxAD == 0 && dot(a - d, b - d) <= 0) return true;
60         if (CDxCA == 0 && dot(c - a, d - a) <= 0) return true;
61         if (CDxCB == 0 && dot(c - b, d - b) <= 0) return true;
62         return false;
63     }
64     return (ABxAC * ABxAD < 0 && CDxCA * CDxCB < 0);
65 }
66 void perpendicular(line l1, point p, line& l2) {
67     if (fabs(l1.a) < EPS)
68         l2 = {1.0, 0.0, -p.x};
69     else {
70         l2.a = -l1.b / l1.a;
71         l2.b = 1.0;
72         l2.c = -l2.a * p.x - l2.b * p.y;
73     }
74 }

```

6.1.3 Circle

```

1 int insideCircle(const point& p, const point& center, ftype r) {
2     ftype d = norm(p - center);
3     ftype rSq = r * r;
4     return fabs(d - rSq) < EPS ? 0 : (d - rSq >= EPS ? 1 : -1);
5 }
6 bool circle2PointsR(const point& p1, const point& p2, ftype r, point& c) {
7     double h = r * r - norm(p1 - p2) / 4.0;
8     if (fabs(h) < 0) return false;
9     h = sqrt(h);
10    point perp = (p2 - p1).rotate(PI / 2.0);
11    point m = (p1 + p2) / 2.0;
12    c = m + perp * (h / abs(perp));
13    return true;
14 }

```

6.1.4 Triangle

```

1 double areaTriangle(double ab, double bc, double ca) {
2     double p = (ab + bc + ca) / 2;
3     return sqrt(p) * sqrt(p - ab) * sqrt(p - bc) * sqrt(p - ca);
4 }
5 double rInCircle(double ab, double bc, double ca) {
6     double p = (ab + bc + ca) / 2;
7     return areaTriangle(ab, bc, ca) / p;
8 }
9 double rInCircle(point a, point b, point c) {
10    return rInCircle(abs(a - b), abs(b - c), abs(c - a));
11 }
12 bool inCircle(point p1, point p2, point p3, point &ctr, double &r) {
13    r = rInCircle(p1, p2, p3);
14    if (fabs(r) < EPS) return false;
15    line l1, l2;
16    double ratio = abs(p2 - p1) / abs(p3 - p1);
17    point p = p2 + (p3 - p2) * (ratio / (1 + ratio));
18    pointsToLine(p1, p, l1);
19    ratio = abs(p1 - p2) / abs(p2 - p3);
20    p = p1 + (p3 - p1) * (ratio / (1 + ratio));
21    pointsToLine(p2, p, l2);
22    areIntersect(l1, l2, ctr);
23    return true;
24 }
25 double rCircumCircle(double ab, double bc, double ca) {
26    return ab * bc * ca / (4.0 * areaTriangle(ab, bc, ca));
27 }
28 double rCircumCircle(point a, point b, point c) {
29    return rCircumCircle(abs(b - a), abs(c - b), abs(a - c));
30 }

```

6.1.5 Convex hull

```

1 vector<point> CH_Andrew(vector<point> &Pts) { // overall O(n log n)
2     int n = Pts.size(), k = 0;
3     vector<point> H(2 * n);
4     sort(Pts.begin(), Pts.end());
5     for (int i = 0; i < n; ++i) {
6         while ((k >= 2) && !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
7         H[k++] = Pts[i];
8     }
9     for (int i = n - 2, t = k + 1; i >= 0; --i) {
10        while ((k >= t) && !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
11        H[k++] = Pts[i];
12    }
13    H.resize(k);
14    return H;
15 }

```

6.1.6 Polygon

```

1 double perimeter(const vector<point> &P) {
2     double ans = 0.0;
3     for (int i = 0; i < (int)P.size() - 1; ++i)

```

```

4         ans += abs(P[i] - P[i + 1]);
5     return ans;
6 }
7 double area(const vector<point> &P) {
8     double ans = 0.0;
9     for (int i = 0; i < (int)P.size() - 1; ++i)
10        ans += (P[i].x * P[i + 1].y - P[i + 1].x * P[i].y);
11    return fabs(ans) / 2.0;
12 }
13 bool isConvex(const vector<point> &P) {
14    int n = (int)P.size();
15    if (n <= 3) return false;
16    bool firstTurn = ccw(P[0], P[1], P[2]);
17    for (int i = 1; i < n - 1; ++i)
18        if (ccw(P[i], P[i + 1], P[(i + 2) == n ? 1 : i + 2]) != firstTurn)
19            return false;
20    return true;
21 }
22 int insidePolygon(point pt, const vector<point> &P) {
23    int n = (int)P.size();
24    if (n <= 3) return -1;
25    bool on_polygon = false;
26    for (int i = 0; i < n - 1; ++i)
27        if (fabs(abs(P[i] - pt) + abs(pt - P[i + 1]) - abs(P[i] - P[i + 1]))
28            < EPS)
29            on_polygon = true;
30    if (on_polygon) return 0;
31    double sum = 0.0;
32    for (int i = 0; i < n - 1; ++i) {
33        if (ccw(pt, P[i], P[i + 1]))
34            sum += angle(P[i] - pt, P[i + 1] - pt);
35        else
36            sum -= angle(P[i] - pt, P[i + 1] - pt);
37    }
38    return fabs(sum) > PI ? 1 : -1;

```

6.2 Minimum enclosing circle

```

1 /**
2  * Description: computes the minimum circle that encloses all the given
3  * points.
4  */
5 double abs(point a) { return sqrt(a.X * a.X + a.Y * a.Y); }
6 point center_from(double bx, double by, double cx, double cy) {
7     double B = bx * bx + by * by, C = cx * cx + cy * cy, D = bx * cy - by *
8     cx;
9     return point((cy * B - by * C) / (2 * D), (bx * C - cx * B) / (2 * D));
10 }
11 circle circle_from(point A, point B, point C) {
12     point I = center_from(B.X - A.X, B.Y - A.Y, C.X - A.X, C.Y - A.Y);
13     return circle(I + A, abs(I));

```

```

14 }
15
16 const int N = 100005;
17 int n, x[N], y[N];
18 point a[N];
19
20 circle emo_welzl(int n, vector<point> T) {
21     if (T.size() == 3 || n == 0) {
22         if (T.size() == 0) return circle(point(0, 0), -1);
23         if (T.size() == 1) return circle(T[0], 0);
24         if (T.size() == 2) return circle((T[0] + T[1]) / 2, abs(T[0] - T[1])
25 / 2);
26         return circle_from(T[0], T[1], T[2]);
27     }
28     random_shuffle(a + 1, a + n + 1);
29     circle Result = emo_welzl(0, T);
30     for (int i = 1; i <= n; i++)
31         if (abs(Result.X - a[i]) > Result.Y + 1e-9) {
32             T.push_back(a[i]);
33             Result = emo_welzl(i - 1, T);
34             T.pop_back();
35         }
36     return Result;
37 }

```

7 Linear algebra

7.1 Gauss elimination

```

1 const double EPS = 1e-9;
2 const int INF = 2; // it doesn't actually have to be infinity or a big number
3 int gauss (vector < vector<double> > a, vector<double> & ans) {
4     int n = (int) a.size();
5     int m = (int) a[0].size() - 1;
6     vector<int> where (m, -1);
7     for (int col=0, row=0; col<m && row<n; ++col) {
8         int sel = row;
9         for (int i=row; i<n; ++i)
10             if (abs (a[i][col]) > abs (a[sel][col]))
11                 sel = i;
12         if (abs (a[sel][col]) < EPS)
13             continue;
14         for (int i=col; i<=m; ++i)
15             swap (a[sel][i], a[row][i]);
16         where[col] = row;
17
18         for (int i=0; i<n; ++i)
19             if (i != row) {
20                 double c = a[i][col] / a[row][col];
21                 for (int j=col; j<=m; ++j)
22                     a[i][j] -= a[row][j] * c;
23             }
24         ++row;
25     }

```

```

26     ans.assign (m, 0);
27     for (int i=0; i<m; ++i)
28         if (where[i] != -1)
29             ans[i] = a[where[i]][m] / a[where[i]][i];
30     for (int i=0; i<n; ++i) {
31         double sum = 0;
32         for (int j=0; j<m; ++j)
33             sum += ans[j] * a[i][j];
34         if (abs (sum - a[i][m]) > EPS)
35             return 0;
36     }
37     for (int i=0; i<m; ++i)
38         if (where[i] == -1)
39             return INF;
40     return 1;
41 }

```

8 Graph

8.1 Bellman-Ford algorithm

```

1 /**
2  * Description: single source shortest path in a weighted (negative or
3  * positive) directed graph.
4  * Time: O(N * M).
5  * Tested: https://open.kattis.com/problems/shortestpath3
6  */
7 const int64_t INF = (int64_t) 2e18;
8 struct Edge {
9     int u, v; // u -> v
10     int64_t w;
11     Edge() {}
12     Edge(int _u, int _v, int64_t _w) : u(_u), v(_v), w(_w) {}
13 };
14 vector<int64_t> bellmanFord(int s) {
15     // dist[stating] = 0.
16     // dist[u] = +INF, if u is unreachable.
17     // dist[u] = -INF, if there is a negative cycle on the path from s to u.
18     // -INF < dist[u] < +INF, otherwise.
19     vector<int64_t> dist(n, INF);
20     dist[s] = 0;
21     for (int i = 0; i < n - 1; ++i) {
22         bool any = false;
23         for (auto [u, v, w] : edges) {
24             if (dist[u] != INF && dist[v] > w + dist[u]) {
25                 dist[v] = w + dist[u];
26                 any = true;
27             }
28         }
29         if (!any) break;
30     }
31     // handle negative cycles
32     for (int i = 0; i < n - 1; ++i) {
33         for (auto [u, v, w] : edges) {

```

```

33         if (dist[u] != INF && dist[v] > w + dist[u]) {
34             dist[v] = -INF;
35         }
36     }
37 }
38 return dist;
39 }

```

8.2 Articulation point and Bridge

```

1 /**
2  * Description: finding articulation points and bridges in a simple
3  * undirected graph.
4  * Tested: https://oj.vnoi.info/problem/graph_
5  */
6 const int N = (int) 1e5;
7 vector<int> g[N];
8 int num[N], low[N], dfs_timer;
9 bool joint[N];
10 vector<pair<int, int>> bridges;
11 void dfs(int u, int prev) {
12     low[u] = num[u] = ++dfs_timer;
13     int child = 0;
14     for (int v : g[u]) {
15         if (v == prev) continue;
16         if (num[v]) low[u] = min(low[u], num[v]);
17         else {
18             dfs(v, u);
19             low[u] = min(low[u], low[v]);
20             child++;
21             if (low[v] >= num[u]) {
22                 bridges.emplace_back(u, v);
23             }
24             if (u != prev && low[v] >= num[u]) joint[u] = true;
25         }
26     }
27     if (u == prev && child > 1) joint[u] = true;
28 }
29 int main() {
30     int n, m;
31     cin >> n >> m;
32     for (int i = 0; i < m; ++i) {
33         int u, v;
34         cin >> u >> v;
35         u--; v--;
36         g[u].push_back(v);
37         g[v].push_back(u);
38     }
39     for (int i = 0; i < n; ++i) {
40         if (!num[i]) dfs(i, i);
41     }
42     return 0;
43 }

```

8.3 Strongly connected components

```

1 /**
2  * Description: Tarjan's algorithm finds strongly connected components
3  * in a directed graph. If vertices u and v belong to the same component,
4  * then scc_id[u] == scc_id[v].
5  * Tested: https://judge.yosupo.jp/problem/scc
6  */
7 const int N = (int) 5e5;
8 vector<int> g[N], st;
9 int low[N], num[N], dfs_timer, scc_id[N], scc;
10 bool used[N];
11 void Tarjan(int u) {
12     low[u] = num[u] = ++dfs_timer;
13     st.push_back(u);
14     for (int v : g[u]) {
15         if (used[v]) continue;
16         if (num[v] == 0) {
17             Tarjan(v);
18             low[u] = min(low[u], low[v]);
19         }
20         else {
21             low[u] = min(low[u], num[v]);
22         }
23     }
24     if (low[u] == num[u]) {
25         int v;
26         do {
27             v = st.back(); st.pop_back();
28             debug(u, v)
29             used[v] = true;
30             scc_id[v] = scc;
31         } while (v != u);
32         scc++;
33     }
34 }

```

8.4 Topo sort

```

1 /**
2  * Description: A topological sort of a directed acyclic graph
3  * is a linear ordering of its vertices such that for every directed edge
4  * from vertex u to vertex v, u comes before v in the ordering.
5  * Note: If there are cycles, the returned list will have size smaller than n
6  * (i.e, topo.size() < n).
7  * Tested: https://judge.yosupo.jp/problem/scc
8  */
9 vector<int> topo_sort(const vector<vector<int>> &g) {
10     int n = (int) g.size();
11     vector<int> indeg(n);
12     for (int u = 0; u < n; ++u) {
13         for (int v : g[u]) indeg[v]++;
14     }
15     queue<int> q; // Note: use min-heap to get the smallest lexicographical
16     // order.

```

```

15     for (int u = 0; u < n; ++u) {
16         if (indeg[u] == 0) q.emplace(u);
17     }
18     vector<int> topo;
19     while (!q.empty()) {
20         int u = q.front(); q.pop();
21         topo.emplace_back(u);
22         for (int v : g[u]) {
23             if (--indeg[v] == 0) q.emplace(v);
24         }
25     }
26     return topo;
27 }

```

8.5 K-th smallest shortest path

```

1  /** Finding the k-th smallest shortest path from vertex s to vertex t,
2  *   each vertex can be visited more than once.
3  */
4  using adj_list = vector<vector<pair<int, int>>>;
5  vector<int> k_smallest(const adj_list &g, int k, int s, int t) {
6      int n = (int) g.size();
7      vector<long long> ans;
8      vector<int> cnt(n);
9      using pli = pair<long long, int>;
10     priority_queue<pli, vector<pli>, greater<pli>> pq;
11     pq.emplace(0, s);
12     while (!pq.empty() && cnt[t] < k) {
13         int u = pq.top().second;
14         long long d = pq.top().first;
15         pq.pop();
16         if (cnt[u] == k) continue;
17         cnt[u]++;
18         if (u == t) {
19             ans.push_back(d);
20         }
21         for (auto [v, cost] : g[u]) {
22             pq.emplace(d + cost, v);
23         }
24     }
25     assert(ans.size() == k);
26     return ans;
27 }

```

8.6 Eulerian path

8.6.1 Directed graph

```

1  /**
2  * Hierholzer's algorithm.
3  * Description: An Eulerian path in a directed graph is a path that visits
4  *   all edges exactly once.
5  * An Eulerian cycle is a Eulerian path that is a cycle.
6  * Time complexity: O(|E|).
7  */
8  vector<int> find_path_directed(const vector<vector<int>> &g, int s) {

```

```

8     int n = (int) g.size();
9     vector<int> stack, cur_edge(n), vertices;
10    stack.push_back(s);
11    while (!stack.empty()) {
12        int u = stack.back();
13        stack.pop_back();
14        while (cur_edge[u] < (int) g[u].size()) {
15            stack.push_back(u);
16            u = g[u][cur_edge[u]++];
17        }
18        vertices.push_back(u);
19    }
20    reverse(vertices.begin(), vertices.end());
21    return vertices;
22 }

```

8.6.2 Undirected graph

```

1  /**
2  * Hierholzer's algorithm.
3  * Description: An Eulerian path in a undirected graph is a path that visits
4  *   all edges exactly once.
5  * An Eulerian cycle is a Eulerian path that is a cycle.
6  * Time complexity: O(|E|).
7  */
8  struct Edge {
9      int to;
10     list<Edge>::iterator reverse_edge;
11     Edge(int _to) : to(_to) {}
12 };
13 vector<int> vertices;
14 void find_path(vector<list<Edge>> &g, int u) {
15     while (!g[u].empty()) {
16         int v = g[u].front().to;
17         g[v].erase(g[u].front().reverse_edge);
18         g[u].pop_front();
19         find_path(g, v);
20     }
21     vertices.emplace_back(u); // reversion list.
22 }
23 void add_edge(int u, int v) {
24     g[u].emplace_front(v);
25     g[v].emplace_front(u);
26     g[u].front().reverse_edge = g[v].begin();
27     g[v].front().reverse_edge = g[u].begin();
28 }

```

9 Misc.

9.1 Ternary search

```

1  const double eps = 1e-9;
2  double ternary_search_max(double l, double r) {
3      // find x0 such that: f(x0) > f(x), \forall x: l <= x <= r.
4      while (r - l > eps) {

```

```

5         double mid1 = l + (r - l) / 3;
6         double mid2 = r - (r - l) / 3;
7         if (f(mid1) < f(mid2)) l = mid1;
8         else r = mid2;
9     }
10    return l;
11 }
12 double ternary_search_min(double l, double r) {
13     // find x0 such that: f(x0) < f(x), \all x: l <= x <= r.
14     while (r - l > eps) {
15         double mid1 = l + (r - l) / 3;
16         double mid2 = r - (r - l) / 3;
17         if (f(mid1) > f(mid2)) l = mid1;
18         else r = mid2;
19     }
20    return l;
21 }

```

9.2 Ternary search 2

```

1 /**
2  * Source: https://github.com/icpcvn/icpcvn.github.io/blob/master/2017/
3  * national/solutions/E/RR_ternary.cpp
4  */
5
6 Point A, B, C;
7 double f(Point P) {
8     return (P - A).len() + (P - B).len() + (P - C).len();
9 }
10 double f(double x, double& y) {
11     double l = -2000, r = 2000;
12     for (int i = 0; i < 500; ++i) {
13         double y1 = (l * 2 + r) / 3.0;
14         double y2 = (l + r * 2) / 3.0;
15
16         double f1 = f(Point(x, y1));
17         double f2 = f(Point(x, y2));
18
19         if (f1 < f2) r = y2;
20         else l = y1;
21     }
22     y = (l + r) / 2.0;
23     return f(Point(x, y));
24 }
25 int32_t main() {
26     ios::sync_with_stdio(0);
27     cin.tie(0);
28     cout << fixed << setprecision(9) << boolalpha;
29     cin >> A >> B >> C;
30
31     double l = -2000, r = 2000;
32     for (int i = 0; i < 500; ++i) {
33         double x1 = (l * 2 + r) / 3.0;
34         double x2 = (l + r * 2) / 3.0;

```

```

35
36         double y;
37         double f1 = f(x1, y);
38         double f2 = f(x2, y);
39
40         if (f1 < f2) r = x2;
41         else l = x1;
42     }
43     double y;
44     f(l, y);
45     cout << l << ' ' << y << '\n';
46     return 0;
47 }

```

9.3 Dutch flag national problem

```

1 void dutch_flag_national(vector<int> &arr) {
2     // All elements that are LESS than pivot are moved to the LEFT.
3     // All elements that are GREATER than pivot are moved to the RIGHT.
4     // E.g. [1, 2, 0, 0, 2, 2, 1], pivot = 1 -> [0, 0, 1, 1, 2, 2, 2].
5     int n = (int) arr.size();
6     int i = 0, j = 0, k = n - 1;
7     int pivot = 1;
8     // 0....i....j....k....n
9     while (j <= k) {
10         if (arr[j] < pivot) {
11             swap(arr[i], arr[j]);
12             i++;
13             j++;
14         }
15         else if (arr[j] > pivot) {
16             swap(arr[j], arr[k]);
17             k--;
18         }
19         else {
20             j++;
21         }
22     }
23     // 0 <= index <= i - 1: arr[index] < mid.
24     // i <= index <= k: arr[index] = mid.
25     // k + 1 <= index < sz: arr[index] > mid.
26 }

```

9.4 Matrix

```

1 struct Matrix {
2     static const matrix_type INF = numeric_limits<matrix_type>::max();
3     int N, M;
4     vector<vector<matrix_type>> mat;
5
6     Matrix(int _N, int _M, matrix_type v = 0) : N(_N), M(_M) {
7         mat.assign(N, vector<matrix_type>(M, v));
8     }
9     static Matrix identity(int n) { // return identity matrix.
10        Matrix I(n, n);

```

```
11     for (int i = 0; i < n; ++i) {
12         I[i][i] = 1;
13     }
14     return I;
15 }
16
17 vector<matrix_type>& operator[](int r) { return mat[r]; }
18 const vector<matrix_type>& operator[](int r) const { return mat[r]; }
19
20 Matrix& operator*=(const Matrix &other) {
21     assert(M == other.N); // [N x M] [other.N x other.M]
22     Matrix res(N, other.M);
23     for (int r = 0; r < N; ++r) {
24         for (int c = 0; c < other.M; ++c) {
25             long long square_mod = (long long) MOD * MOD;
26             long long sum = 0;
27             for (int g = 0; g < M; ++g) {
28                 sum += (long long) mat[r][g] * other[g][c];
29                 if (sum >= square_mod) sum -= square_mod;
30             }
31             res[r][c] = sum % MOD;
32         }
33     }
34     mat.swap(res.mat); return *this;
35 }
36 };
```