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CTU.NegativeZero

Contents			5	o.5 Zeckendort's theorem	1.
1	Contest 1.1 C++ 1.2 Debug 1.3 Java 1.4 sublime-build 1.5 .bashrc	2 2 2 2 2 3	5 5 5	5.6 Bitwise operation	12 13 13
2	Data structures 2.1 Sparse table 2.2 Ordered set 2.3 Dsu 2.4 Segment tree 2.5 Efficient segment tree 2.6 Persistent lazy segment tree 2.7 Disjoint sparse table 2.8 Fenwick tree 2.9 Implicit treap	3 3 3 4 4 5 5 6 6	7 0	Linear algebra 5.1 Gauss elimination	1 4
	Mathematics 3.1 Trigonometry 3.2 Sums String 4.1 Prefix function 4.2 Counting occurrences of each prefix 4.3 Knuth–Morris–Pratt algorithm 4.4 Suffix array 4.5 Manacher's algorithm 4.6 Trie 4.7 Hashing	7 7 8 8 8 8 8 8 9 9 9	8 8 8 8	Graph 3.1 Bellman-Ford algorithm 3.2 Articulation point and Bridge 3.3 Strongly connected components 3.4 Topo sort 3.5 K-th smallest shortest path 3.6 Eulerian path	15 18 18 18
5	Number Theory 5.1 Euler's totient function 5.2 Mobius function 5.3 Primes 5.4 Wilson's theorem	10 10 11 11 11	9	Misc. O.1 Ternary search	20

Can Tho University Page 2 of 20

1 Contest

1.1 C++

```
1 #include <bits/stdc++.h>
using namespace std;
4 #ifdef LOCAL
5 #include "cp/debug.h"
6 #else
7 #define debug(...)
8 #endif
nt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
12 const int MOD = (int) 1e9 + 7;
13 const int INF = 0x3f3f3f3f3f:
14
15 int main() {
      ios::sync_with_stdio(false); cin.tie(nullptr);
      // freopen("input.txt", "r", stdin);
      // freopen("output.txt", "w", stdout);
18
      return 0:
20
21 }
1.2
     Debug
1 #define debug(...) { string _s = #__VA_ARGS__; replace(begin(_s), end(_s),
      ',', ''); stringstream _ss(_s); istream_iterator<string> _it(_ss);
      out_error(_it, __VA_ARGS__);}
3 void out_error(istream_iterator<string> it) { cerr << '\n'; }</pre>
5 template<typename T, typename ...Args>
6 void out_error(istream_iterator<string> it, T a, Args... args) {
      cerr << " [" << *it << " = " << a << "] ";
      out_error(++it, args...);
9 }
11 template<typename T, typename G> ostream& operator<<(ostream &os, const
      pair<T, G> &p) {
      return os << "(" << p.first << ", " << p.second << ")";</pre>
12
13 }
15 template < class Con, class = decltype(begin(declval < Con > ())) >
16 typename enable_if<!is_same<Con, string>::value, ostream&>::type
operator << (ostream& os, const Con& container) {</pre>
      os << "{";
18
      for (auto it = container.begin(); it != container.end(); ++it)
19
20
          os << (it == container.begin() ? "" : ", ") << *it;
      return os << "}";</pre>
22 }
1.3
      Java
```

```
import java.io.BufferedReader;
2 import java.util.StringTokenizer;
3 import java.io.IOException;
4 import java.io.InputStreamReader;
5 import java.io.PrintWriter;
6 import java.util.ArrayList;
7 import java.util.Arrays;
8 import java.util.Collections;
9 import java.util.Random;
 public class Main {
      public static void main(String[] args) {
          FastScanner fs = new FastScanner();
          PrintWriter out = new PrintWriter(System.out);
          int n = fs.nextInt();
          out.println(n);
          out.close(); // don't forget this line.
      static class FastScanner {
          BufferedReader br:
          StringTokenizer st;
          public FastScanner() {
              br = new BufferedReader(new InputStreamReader(System.in));
              st = null;
          public String next() {
              while (st == null || st.hasMoreTokens() == false) {
                  try {
                      st = new StringTokenizer(br.readLine());
                  catch (IOException e) {
                      throw new RuntimeException(e);
              return st.nextToken();
          }
          public int nextInt() {
              return Integer.parseInt(next());
          public long nextLong() {
              return Long.parseLong(next());
          public double nextDouble() {
              return Double.parseDouble(next());
          }
      sublime-build
1 {
```

Can Tho University Page 3 of 20

```
"cmd": ["g++", "-std=c++17", "-fmax-errors=5", "-DLOCAL", "-Wall",
      "-Wextra", "-o", "${file_path}/${file_base_name}.out", "${file}"],
      "file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)?:? (.*)$",
      "working_dir": "${file_path}",
      "selector": "source.cpp, source.c++"
6 }
1.5
     .bashrc
1 alias cpp='g++ -std=c++17 -fmax-errors=5 -DLOCAL -Wall -Wextra'
3 #Stress-testing
4 function test {
    SOL = 1
    CHECKER=$2
    for i in {1..100};
        ./gen.out > in && ./"$CHECKER.out" < in > ans && ./"$SOL.out" < in >
      out && diff -Z out ans && echo "Test $i passed!!" || break;
      done
11 }
```

2 Data structures

2.1 Sparse table

```
1 int st[MAXN][K + 1];
2 for (int i = 0; i < N; i++) {
      st[i][0] = f(array[i]);
4 }
5 for (int j = 1; j \le K; j++) {
      for (int i = 0; i + (1 << j) <= N; i++) {
          st[i][j] = f(st[i][j-1], st[i+(1 << (j-1))][j-1]);
8
9 }
10 // Range Minimum Queries.
int lg[MAXN + 1];
12 lg[1] = 0;
13 for (int i = 2; i \le MAXN; i++) {
      lg[i] = lg[i / 2] + 1;
15 }
16 int j = lg[R - L + 1];
int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);</pre>
18 // Range Sum Queries.
19 long long sum = 0;
20 for (int j = K; j >= 0; j--) {
      if ((1 << j) <= R - L + 1) {
21
22
          sum += st[L][j];
          L += 1 << j;
23
24
      }
25 }
```

2.2 Ordered set

```
1 #include <ext/pb_ds/assoc_container.hpp>
2 #include <ext/pb_ds/tree_policy.hpp>
```

```
3 using namespace __gnu_pbds;
4
5 template<typename key_type>
6 using set_t = tree<key_type, null_type, less<key_type>, rb_tree_tag,
       tree_order_statistics_node_update>;
9 void example() {
       vector < int > nums = \{1, 2, 3, 5, 10\};
       set_t<int> st(nums.begin(), nums.end());
       cout << *st.find_by_order(0) << '\n'; // 1</pre>
       assert(st.find_by_order(-INF) == st.end());
       assert(st.find_by_order(INF) == st.end());
       cout << st.order_of_key(2) << '\n'; // 1</pre>
       cout << st.order_of_key(4) << '\n'; // 3
       cout << st.order_of_key(9) << '\n'; // 4</pre>
       cout << st.order_of_key(-INF) << '\n'; // 0</pre>
       cout << st.order_of_key(INF) << '\n'; // 5</pre>
22 }
 2.3
       Dsu
1 struct Dsu {
      int n:
       vector<int> par, sz;
       Dsu(int _n) : n(_n) {
           sz.resize(n, 1);
           par.resize(n);
           iota(par.begin(), par.end(), 0);
       int find(int v) {
          // finding leader/parrent of set that contains the element v.
          // with {path compression optimization}.
           return (v == par[v] ? v : par[v] = find(par[v]));
       bool same(int u, int v) {
           return find(u) == find(v);
       bool unite(int u, int v) {
          u = find(u); v = find(v);
          if (u == v) return false;
          if (sz[u] < sz[v]) swap(u, v);
           par[v] = u;
           sz[u] += sz[v];
23
          return true;
       vector<vector<int>> groups() {
           // returns the list of the "list of the vertices in a connected
       component".
           vector<int> leader(n);
           for (int i = 0; i < n; ++i) {
               leader[i] = find(i);
```

Can Tho University Page 4 of 20

```
vector<int> id(n, -1);
31
          int count = 0;
32
33
          for (int i = 0; i < n; ++i) {
              if (id[leader[i]] == -1) {
34
                  id[leader[i]] = count++;
              }
          }
          vector<vector<int>>> result(count);
          for (int i = 0; i < n; ++i) {
39
              result[id[leader[i]]].push_back(i);
40
          return result;
44 };
      Segment tree
2 * Description: A segment tree with range updates and sum queries that
      supports three types of operations:
      + Increase each value in range [1, r] by x (i.e. a[i] += x).
      + Set each value in range [1, r] to x (i.e. a[i] = x).
      + Determine the sum of values in range [1, r].
6 */
7 struct SegmentTree {
      int n;
      vector<long long> tree, lazy_add, lazy_set;
      SegmentTree(int _n) : n(_n) {
10
          int p = 1;
11
          while (p < n) p *= 2;
12
          tree.resize(p * 2);
13
          lazy_add.resize(p * 2);
14
          lazy_set.resize(p * 2);
15
      }
16
      long long merge(const long long &left, const long long &right) {
          return left + right;
18
19
      void build(int id, int l, int r, const vector<int> &arr) {
20
          if (1 == r) {
21
              tree[id] += arr[1];
22
              return;
23
24
          int mid = (1 + r) >> 1;
25
          build(id * 2, 1, mid, arr);
          build(id * 2 + 1, mid + 1, r, arr);
          tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
28
      }
29
      void push(int id, int l, int r) {
30
          if (lazy_set[id] == 0 && lazy_add[id] == 0) return;
31
          int mid = (1 + r) >> 1;
32
          for (int child : {id * 2, id * 2 + 1}) {
33
              int range = (child == id * 2 ? mid - 1 + 1 : r - mid);
34
              if (lazy_set[id] != 0) {
                  lazy_add[child] = 0;
```

```
lazy_set[child] = lazy_set[id];
                  tree[child] = range * lazy_set[id];
              lazy_add[child] += lazy_add[id];
              tree[child] += range * lazy_add[id];
          lazy_add[id] = lazy_set[id] = 0;
      void update(int id, int l, int r, int u, int v, int amount, bool
      set_value = false) {
          if (r < u \mid \mid 1 > v) return;
          if (u <= 1 && r <= v) {
              if (set_value) {
                  tree[id] = 1LL * amount * (r - l + 1);
                  lazy_set[id] = amount;
                  lazy_add[id] = 0; // clear all previous updates.
              }
                  tree[id] += 1LL * amount * (r - 1 + 1);
                  lazy_add[id] += amount;
              }
              return;
          push(id, 1, r);
          int mid = (1 + r) >> 1;
          update(id * 2, 1, mid, u, v, amount, set_value);
          update(id * 2 + 1, mid + 1, r, u, v, amount, set_value);
          tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
      long long get(int id, int l, int r, int u, int v) {
          if (r < u \mid \mid 1 > v) return 0;
          if (u <= 1 && r <= v) {
              return tree[id];
          push(id, 1, r);
          int mid = (1 + r) >> 1;
          long long left = get(id * 2, 1, mid, u, v);
          long long right = get(id * 2 + 1, mid + 1, r, u, v);
          return merge(left, right);
77 };
2.5 Efficient segment tree
1 template < typename T> struct SegmentTree {
      int n;
      vector<T> tree:
      SegmentTree(int _n) : n(_n), tree(2 * n) {}
      T merge(const T &left, const T &right) {
          return left + right;
      template < typename G>
      void build(const vector<G> &initial) {
```

Can Tho University Page 5 of 20

```
assert((int) initial.size() == n);
                                                                                                    tree.back().has_changed = true;
10
          for (int i = 0; i < n; ++i) {</pre>
                                                                                                    return tree.size() - 1;
                                                                                     27
11
12
               tree[i + n] = initial[i];
                                                                                                }
                                                                                                int mid = (1 + r) >> 1;
13
          for (int i = n - 1; i > 0; --i) {
                                                                                                push(x, 1, mid, r);
14
              tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
                                                                                                int left = add(tree[x].1, 1, mid, u, v, amt);
15
                                                                                                int right = add(tree[x].r, mid + 1, r, u, v, amt);
      }
                                                                                                tree.emplace_back(left, right, tree[left].val + tree[right].val, 0);
      void modify(int i, int v) {
                                                                                                return tree.size() - 1;
18
          tree[i += n] = v;
19
          for (i /= 2; i > 0; i /= 2) {
                                                                                            long long get_sum(int x, int l, int r, int u, int v) {
                                                                                                if (r < u \mid | 1 > v) return 0;
              tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
21
          }
                                                                                                if (u <= 1 && r <= v) return tree[x].val;</pre>
22
      }
                                                                                                int mid = (1 + r) / 2;
23
      T get_sum(int 1, int r) {
                                                                                                push(x, 1, mid, r);
24
          // sum of elements from 1 to r - 1.
                                                                                                return get_sum(tree[x].1, 1, mid, u, v) + get_sum(tree[x].r, mid + 1,
25
          T ret{};
26
                                                                                            r, u, v);
          for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
                                                                                            }
27
              if (1 & 1) ret = merge(ret, tree[1++]);
                                                                                            void push(int x, int 1, int mid, int r) {
28
              if (r & 1) ret = merge(ret, tree[--r]);
                                                                                                if (!tree[x].has changed) return:
          }
                                                                                                Vertex left = tree[tree[x].1];
                                                                                                Vertex right = tree[tree[x].r];
31
          return ret;
                                                                                                tree.emplace_back(left);
32
                                                                                                tree[x].l = tree.size() - 1;
33 };
                                                                                                tree.emplace_back(right);
     Persistent lazy segment tree
                                                                                                tree[x].r = tree.size() - 1;
1 struct Vertex {
      int 1, r;
                                                                                                tree[tree[x].l].val += tree[x].lazy * (mid - 1 + 1);
      long long val, lazy;
                                                                                                tree[tree[x].1].lazy += tree[x].lazy;
      bool has_changed = false;
                                                                                                tree[tree[x].r].val += tree[x].lazy * (r - mid);
      Vertex() {}
      Vertex(int _1, int _r, long long _val, int _lazy = 0) : l(_1), r(_r),
                                                                                                tree[tree[x].r].lazy += tree[x].lazy;
      val(_val), lazy(_lazy) {}
                                                                                                tree[tree[x].1].has_changed = true;
7 };
                                                                                                tree[tree[x].r].has_changed = true;
8 struct PerSegmentTree {
      vector<Vertex> tree;
                                                                                                tree[x].lazy = 0;
                                                                                                tree[x].has_changed = false;
      vector<int> root;
10
      int build(const vector<int> &arr, int 1, int r) {
11
                                                                                      63 };
          if (1 == r) {
12
              tree.emplace_back(-1, -1, arr[1]);
13
                                                                                            Disjoint sparse table
              return tree.size() - 1;
14
15
          int mid = (1 + r) / 2;
                                                                                      2 * Description: range query on a static array.
          int left = build(arr, 1, mid);
                                                                                         * Time: O(1) per query.
                                                                                      4 * Tested: stress-test.
18
          int right = build(arr, mid + 1, r);
           tree.emplace_back(left, right, tree[left].val + tree[right].val);
19
          return tree.size() - 1;
                                                                                      6 const int MOD = (int) 1e9 + 7;
20
                                                                                      7 struct DisjointSparseTable { // product queries.
21
      int add(int x, int 1, int r, int u, int v, int amt) {
                                                                                            int n, h;
22
          if (1 > v \mid | r < u) return x;
                                                                                            vector<vector<int>> dst;
23
          if (u <= 1 && r <= v) {
                                                                                            vector<int> lg;
24
              tree.emplace_back(tree[x].l, tree[x].r, tree[x].val + 1LL * amt *
                                                                                            DisjointSparseTable(int _n) : n(_n) {
      (r - l + 1), tree[x].lazy + amt);
                                                                                     12
                                                                                                h = 1; // in case n = 1: h = 0 !!.
```

Can Tho University Page 6 of 20

```
int p = 1;
                                                                                                 for (int i = 1; i < n; ++i) {</pre>
13
                                                                                      19
          while (p < n) p *= 2, h++;
                                                                                                      fenw_coeff[i] = fenw_coeff[i - 1] + (long long) diff[i] * (n - i);
14
                                                                                      20
15
          lq.resize(p); lq[1] = 0;
                                                                                                     fenw[i] = fenw[i - 1] + diff[i];
          for (int i = 2; i < p; ++i) {
                                                                                                 for (int i = n - 1; i >= 0; --i) {
              lg[i] = 1 + lg[i / 2];
17
                                                                                                     int j = (i \& (i + 1)) - 1;
18
                                                                                                     if (i >= 0) {
          dst.resize(h, vector<int>(n));
      }
                                                                                                          fenw_coeff[i] -= fenw_coeff[j];
20
      void build(const vector<int> &A) {
                                                                                                          fenw[i] -= fenw[j];
21
                                                                                                     }
          for (int lv = 0; lv < h; ++lv) {
22
23
               int len = (1 << lv);</pre>
                                                                                                 }
               for (int k = 0; k < n; k += len * 2) {
24
                   int mid = min(k + len, n);
                                                                                             void add(vector<tree_type> &fenw, int i, tree_type val) {
                   dst[lv][mid - 1] = A[mid - 1] % MOD;
                                                                                                 while (i < n) {
                   for (int i = mid - 2; i >= k; --i) {
                                                                                                     fenw[i] += val;
27
                       dst[lv][i] = 1LL * A[i] * dst[lv][i + 1] % MOD;
                                                                                                     i |= (i + 1);
                   if (mid == n) break;
                                                                                             }
                   dst[lv][mid] = A[mid] % MOD;
                                                                                             tree_type __prefix_sum(vector<tree_type> &fenw, int i) {
                   for (int i = mid + 1; i < min(mid + len, n); ++i) {</pre>
                                                                                                 tree type res{}:
32
                       dst[lv][i] = 1LL * A[i] * dst[lv][i - 1] % MOD;
                                                                                                 while (i >= 0) {
                                                                                                     res += fenw[i];
34
              }
                                                                                                     i = (i \& (i + 1)) - 1;
35
          }
                                                                                      42
      }
                                                                                                 return res;
37
      int get(int 1, int r) {
          if (1 == r) {
                                                                                             tree_type prefix_sum(int i) {
39
                                                                                                 return __prefix_sum(fenw_coeff, i) - __prefix_sum(fenw, i) * (n - i -
40
               return dst[0][1];
                                                                                             1);
          int i = lg[l ^ r];
42
                                                                                             void range_add(int 1, int r, tree_type val) {
          return 1LL * dst[i][l] * dst[i][r] % MOD;
43
                                                                                                 add(fenw_coeff, 1, (n - 1) * val);
44
                                                                                                 add(fenw_coeff, r + 1, (n - r - 1) * (-val));
45 };
                                                                                                 add(fenw, 1, val);
      Fenwick tree
                                                                                      52
                                                                                                 add(fenw, r + 1, -val);
using tree_type = long long;
2 struct FenwickTree {
                                                                                             tree_type range_sum(int 1, int r) {
      int n;
                                                                                                 return prefix_sum(r) - prefix_sum(l - 1);
      vector<tree_type> fenw_coeff, fenw;
                                                                                      57 };
      FenwickTree() {}
      FenwickTree(int _n) : n(_n) {
                                                                                             Implicit treap
          fenw\_coeff.assign(n, 0); // fenwick tree with coefficient (n - i).
          fenw.assign(n, 0); // normal fenwick tree.
                                                                                       1 struct Node {
      }
                                                                                             int val, prior, cnt;
      void build(const vector<int> &A) {
                                                                                             bool rev;
10
           assert((int) A.size() == n);
                                                                                             Node *left, *right;
11
          vector<int> diff(n);
                                                                                             Node() {}
12
          diff[0] = A[0];
                                                                                             Node(int _val) : val(_val), prior(rng()), cnt(1), rev(false),
13
          for (int i = 1; i < n; ++i) {</pre>
                                                                                             left(nullptr), right(nullptr) {}
14
               diff[i] = A[i] - A[i - 1];
15
                                                                                       7 };
                                                                                       8 // Binary search tree + min-heap.
          fenw\_coeff[0] = (long long) diff[0] * n;
                                                                                       9 struct Treap {
          fenw[0] = diff[0];
                                                                                             Node *root;
18
```

Can Tho University Page 7 of 20

```
Treap() : root(nullptr) {}
11
      int get_cnt(Node *n) { return n ? n->cnt : 0; }
12
13
      void upd_cnt(Node *&n) {
          if (n) n->cnt = get_cnt(n->left) + get_cnt(n->right) + 1;
14
15
      void push_rev(Node *treap) {
16
17
           if (!treap || !treap->rev) return;
           treap->rev = false;
           swap(treap->left, treap->right);
          if (treap->left) treap->left->rev ^= true;
20
21
          if (treap->right) treap->right->rev ^= true;
22
      pair<Node*, Node*> split(Node *treap, int x, int smaller = 0) {
23
          if (!treap) return {};
24
           push_rev(treap);
25
           int idx = smaller + get_cnt(treap->left); // implicit val.
27
          if (idx <= x) {
               auto pr = split(treap->right, x, idx + 1);
28
               treap->right = pr.first;
               upd_cnt(treap);
               return {treap, pr.second};
32
           else {
33
               auto pl = split(treap->left, x, smaller);
34
               treap->left = pl.second;
               upd_cnt(treap);
               return {pl.first, treap};
      }
      Node* merge(Node *1, Node *r) {
40
           push_rev(l); push_rev(r);
41
          if (!l || !r) return (l ? l : r);
42
43
           if (l->prior < r->prior) {
               1->right = merge(1->right, r);
               upd_cnt(1);
               return 1;
           else {
               r->left = merge(l, r->left);
               upd_cnt(r);
51
               return r;
52
53
      void insert(int pos, int val) {
54
          if (!root) {
55
               root = new Node(val);
               return:
          Node *1, *m, *r;
          m = new Node(val);
          tie(l, r) = split(root, pos - 1);
          root = merge(l, merge(m, r));
63
```

```
void erase(int pos_l, int pos_r) {
          Node *1, *m, *r;
          tie(l, r) = split(root, pos_l - 1);
          tie(m, r) = split(r, pos_r - pos_l);
          root = merge(1, r);
      void reverse(int pos_l, int pos_r) {
          Node *1, *m, *r;
          tie(l, r) = split(root, pos_l - 1);
          tie(m, r) = split(r, pos_r - pos_l);
          m->rev ^= true;
          root = merge(l, merge(m, r));
      int query(int pos_l, int pos_r);
           // returns answer for corresponding types of query.
      void inorder(Node *n) {
          if (!n) return;
          push_rev(n);
          inorder(n->left);
           cout << n->val << ' ';
          inorder(n->right);
      void print() {
          inorder(root);
          cout << '\n';
90 };
```

3 Mathematics

3.1 Trigonometry

3.1.1 Sum - difference identities

$$\sin(u \pm v) = \sin(u)\cos(v) \pm \cos(u)\sin(v)$$

$$\cos(u \pm v) = \cos(u)\cos(v) \mp \sin(u)\sin(v)$$

$$\tan(u \pm v) = \frac{\tan(u) \pm \tan(v)}{1 \mp \tan(u)\tan(v)}$$

3.1.2 Sum to product identities

$$\cos(u) + \cos(v) = 2\cos(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\cos(u) - \cos(v) = -2\sin(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

$$\sin(u) + \sin(v) = 2\sin(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\sin(u) - \sin(v) = 2\cos(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

Can Tho University Page 8 of 20

3.1.3 Product identities

$$\cos(u)\cos(v) = \frac{1}{2}[\cos(u+v) + \cos(u-v)]$$

$$\sin(u)\sin(v) = -\frac{1}{2}[\cos(u+v) - \cos(u-v)]$$

$$\sin(u)\cos(v) = \frac{1}{2}[\sin(u+v) + \sin(u-v)]$$

3.1.4 Double - triple angle identities

$$\sin(2u) = 2\sin(u)\cos(u)$$

$$\cos(2u) = 2\cos^{2}(u) - 1 = 1 - 2\sin^{2}(u)$$

$$\tan(2u) = \frac{2\tan(u)}{1 - \tan^{2}(u)}$$

$$\sin(3u) = 3\sin(u) - 4\sin^{3}(u)$$

$$\cos(3u) = 4\cos^{3}(u) - 3\cos(u)$$

$$\tan(3u) = \frac{3\tan(u) - \tan^{3}(u)}{1 - 3\tan^{2}(u)}$$

3.2 Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$c + 2c^{2} + \dots + nc^{n} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c - 1)^{2}}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(n+1)(2n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \left(\frac{n(n+1)}{2}\right)^{2}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$1^{5} + 2^{5} + 3^{5} + \dots + n^{5} = \frac{n^{2}(n+1)^{2}(2n^{2} + 2n - 1)}{12}$$

$$1^{6} + 2^{6} + 3^{6} + \dots + n^{6} = \frac{n(n+1)(2n+1)(3n^{4} + 6n^{3} - 3n + 1)}{42}$$

$$1^{7} + 2^{7} + 3^{7} + \dots + n^{7} = \frac{n^{2}(n+1)^{2}(3n^{4} + 6n^{3} - n^{2} - 4n + 2)}{24}$$

4 String

4.1 Prefix function

```
2 * Description: The prefix function of a string 's' is defined as an array pi
      where pi[i] is the length of the longest proper prefix of the substring
      s[0..i] which is also a suffix of this substring.
   * Time complexity: O(|S|).
7 vector<int> prefix_function(const string &s) {
      int n = (int) s.length();
      vector<int> pi(n);
      pi[0] = 0;
      for (int i = 1; i < n; ++i) {
          int j = pi[i - 1]; // try length pi[i - 1] + 1.
          while (j > 0 \&\& s[j] != s[i]) {
              i = pi[i - 1];
          if (s[j] == s[i]) {
              pi[i] = i + 1;
      return pi;
      Counting occurrences of each prefix
vector<int> count_occurrences(const string &s) {
      vector<int> pi = prefix_function(s);
      int n = (int) s.size();
      vector<int> ans(n + 1);
      for (int i = 0; i < n; ++i) {
          ans[pi[i]]++;
      for (int i = n - 1; i > 0; --i) {
          ans[pi[i - 1]] += ans[i];
      for (int i = 0; i <= n; ++i) {
          ans[i]++;
      return ans;
      // Input: ABACABA
      // Output: 4 2 2 1 1 1 1
17 }
      Knuth-Morris-Pratt algorithm
     Searching for a substring in a string.
     Time complexity: O(N + M).
5 vector<int> KMP(const string &text, const string &pattern) {
      int n = (int) text.length();
      int m = (int) pattern.length();
      string s = pattern + '$' + text;
      vector<int> pi = prefix_function(s);
      vector<int> indices;
```

Can Tho University Page 9 of 20

```
for (int i = 0; i < (int) s.length(); ++i) {</pre>
11
           if (pi[i] == m) {
12
13
               indices.push_back(i - 2 * m);
14
      }
15
      return indices;
16
17 }
      Suffix array
struct SuffixArray {
       string s;
       int n, lim;
3
      vector<int> sa, lcp, rank;
       SuffixArray(const string &_s, int _{lim} = 256) : s(_s), n(s.length() + 1),
       \lim(_{\lim}, \sin(n), \log(n), rank(n))
          s += '$':
          build();
           kasai();
           sa.erase(sa.begin());
          lcp.erase(lcp.begin());
10
11
           s.pop_back();
      }
12
      void build() {
13
14
           vector<int> nrank(n), norder(n), cnt(max(n, lim));
           for (int i = 0; i < n; ++i) {
15
               sa[i] = i; rank[i] = s[i];
17
           for (int k = 0, rank_cnt = 0; rank_cnt < n - 1; k = max(1, k * 2),
18
       lim = rank_cnt + 1) {
               // counting sort.
               for (int i = 0; i < n; ++i) norder[i] = (sa[i] - k + n) % n;
20
               for (int i = 0; i < n; ++i) cnt[rank[i]]++;</pre>
21
               for (int i = 1; i < lim; ++i) cnt[i] += cnt[i - 1];</pre>
22
               for (int i = n - 1; i \ge 0; --i) sa[--cnt[rank[norder[i]]]] =
23
       norder[i];
               rank[sa[0]] = rank\_cnt = 0;
24
               for (int i = 1; i < n; ++i) {
25
                   int u = sa[i], v = sa[i - 1];
                   int nu = u + k, nv = v + k;
                   if (nu >= n) nu -= n;
                   if (nv >= n) nv -= n;
                   if (rank[u] != rank[v] || rank[nu] != rank[nv]) ++rank_cnt;
                   nrank[sa[i]] = rank_cnt;
31
33
               for (int i = 0; i < rank_cnt + 1; ++i) cnt[i] = 0;</pre>
               rank.swap(nrank);
34
          }
35
      }
       void kasai() {
37
           for (int i = 0; i < n; ++i) rank[sa[i]] = i;</pre>
38
           for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
39
               int j = sa[rank[i] - 1];
               while (s[i + k] == s[j + k]) k++;
```

```
lcp[rank[i]] = k;
42
43
44
          // Note: lcp[i] = longest common prefix(sa[i - 1], sa[i]).
46 };
      Manacher's algorithm
1 /**
2 * Description: for each position, computes d[0][i] = half length of
3 longest palindrome centered on i (rounded up), d[1][i] = half length of
   longest palindrome centered on i and i - 1.
   * Time complexity: O(N).
   * Tested: https://judge.yosupo.jp/problem/enumerate_palindromes,
       stress-tested.
8 array<vector<int>, 2> manacher(const string &s) {
      int n = (int) s.size();
       array<vector<int>, 2> d;
       for (int z = 0; z < 2; ++z) {
          d[z].resize(n);
           int 1 = 0, r = 0;
           for (int i = 0; i < n; ++i) {
              int mirror = l + r - i + z;
              d[z][i] = (i > r ? 0 : min(d[z][mirror], r - i));
              int L = i - d[z][i] - z, R = i + d[z][i];
              while (L >= 0 \&\& R < n \&\& s[L] == s[R]) {
                   d[z][i]++; L--; R++;
              if (R > r) {
                  l = L; r = R;
              }
          }
      }
       return d;
      Trie
1 struct Trie {
       const static int ALPHABET = 26;
       const static char minChar = 'a';
       struct Vertex {
           int next[ALPHABET];
          bool leaf;
          Vertex() {
              leaf = false;
              fill(next, next + ALPHABET, -1);
      };
       vector<Vertex> trie;
      Trie() { trie.emplace_back(); }
15
       void insert(const string &s) {
          int i = 0;
```

Can Tho University Page 10 of 20

```
for (const char &ch : s) {
17
               int j = ch - minChar;
18
               if (trie[i].next[j] == -1) {
                   trie[i].next[j] = trie.size();
20
                   trie.emplace_back();
21
22
23
               i = trie[i].next[j];
           trie[i].leaf = true;
25
26
27
      bool find(const string &s) {
          int i = 0:
28
           for (const char &ch : s) {
29
               int j = ch - minChar;
30
               if (trie[i].next[j] == -1) {
31
                   return false:
32
33
               i = trie[i].next[j];
34
          return (trie[i].leaf ? true : false);
37
38 };
      Hashing
1 struct Hash61 {
      static const uint64_t MOD = (1LL << 61) - 1;</pre>
      static uint64_t BASE;
      static vector<uint64_t> pw;
      uint64_t addmod(uint64_t a, uint64_t b) const {
          a += b;
          if (a >= MOD) a -= MOD;
          return a;
      }
9
      uint64_t submod(uint64_t a, uint64_t b) const {
          a += MOD - b;
11
          if (a >= MOD) a -= MOD;
12
13
          return a;
      }
14
      uint64_t mulmod(uint64_t a, uint64_t b) const {
15
          uint64_t low1 = (uint32_t) a, high1 = (a >> 32);
16
17
          uint64_t low2 = (uint32_t) b, high2 = (b >> 32);
18
           uint64_t low = low1 * low2;
          uint64_t mid = low1 * high2 + low2 * high1;
          uint64_t high = high1 * high2;
21
22
          uint64_t ret = (low & MOD) + (low >> 61) + (high << 3) + (mid >> 29)
23
      + (mid << 35 >> 3) + 1;
          // ret %= MOD:
24
           ret = (ret >> 61) + (ret & MOD);
25
          ret = (ret >> 61) + (ret & MOD);
           return ret - 1;
28
```

```
void ensure_pw(int m) {
29
           int n = (int) pw.size();
30
           if (n >= m) return;
          pw.resize(m);
           for (int i = n; i < m; ++i) {
               pw[i] = mulmod(pw[i - 1], BASE);
      }
       vector<uint64_t> pref;
       template < typename T > Hash61(const T &s) { // strings or arrays.
          n = (int) s.size();
           ensure_pw(n);
           pref.resize(n + 1);
          pref[0] = 0;
           for (int i = 0; i < n; ++i) {
               pref[i + 1] = addmod(mulmod(pref[i], BASE), s[i]);
      inline uint64_t operator()(const int from, const int to) const {
           assert(0 \le from \&\& from \le to \&\& to < n);
          // pref[to + 1] - pref[from] * pw[to - from + 1]
           return submod(pref[to + 1], mulmod(pref[from], pw[to - from + 1]));
54 };
55 mt19937 rng((unsigned int)
       chrono::steady_clock::now().time_since_epoch().count());
|_{56} uint64_t Hash61::BASE = (MOD >> 2) + rng() % (MOD >> 1);
vector<uint64_t> Hash61::pw = vector<uint64_t>(1, 1);
```

5 Number Theory

5.1 Euler's totient function

- Euler's totient function, also known as **phi-function** $\phi(n)$ counts the number of integers between 1 and n inclusive, that are **coprime to** n.
- Properties:
 - Divisor sum property: $\sum_{d|n} \phi(d) = n$.
 - $\phi(n)$ is a **prime number** when n = 3, 4, 6.
 - If p is a prime number, then $\phi(p) = p 1$.
 - If *p* is a prime number and $k \ge 1$, then $\phi(p^k) = p^k p^{k-1}$.
 - If *a* and *b* are **coprime**, then $\phi(ab) = \phi(a) \cdot \phi(b)$.
 - In general, for **not coprime** a and b, with d = gcd(a, b) this equation holds: $\phi(ab) = \phi(a) \cdot \phi(b) \cdot \frac{d}{\phi(d)}$.

Can Tho University Page 11 of 20

- With $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\phi(n) = \phi(p_1^{k_1}) \cdot \phi(p_2^{k_2}) \cdots \phi(p_m^{k_m})$$
$$= n \cdot \left(1 - \frac{1}{p_1}\right) \cdot \left(1 - \frac{1}{p_2}\right) \cdots \cdot \left(1 - \frac{1}{p_m}\right)$$

- Application in Euler's theorem:
 - If gcd(a, M) = 1, then:

$$a^{\phi(M)} \equiv 1 \pmod{M} \Rightarrow a^n \equiv a^{n \bmod{\phi(M)}} \pmod{M}$$

- In general, for arbitrary a, M and n ≥ $\log_2 M$:

$$a^n \equiv a^{\phi(M) + [n \bmod \phi(M)]} \pmod{M}$$

5.2 Mobius function

• For a positive integer $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\mu(n) = \begin{cases} 1, & \text{if } n = 1\\ 0, & \text{if } \exists k_i > 1\\ (-1)^m & \text{otherwise} \end{cases}$$

- Properties:
 - $-\sum_{d|n}\mu(d)=[n=1].$
 - If *a* and *b* are **coprime**, then $\mu(ab) = \mu(a) \cdot \mu(b)$.
 - Mobius inversion: let *f* and *g* be arithmetic functions:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right)g(d)$$

5.3 Primes

Approximating the number of primes up to *n*:

n	$\pi(n)$	$\frac{n}{\ln n - 1}$
$100 (1e^2)$	25	28
$500 (5e^2)$	95	96
$1000 (1e^3)$	168	169
$5000 (5e^3)$	669	665
$10000 (1e^4)$	1229	1218
$50000 (5e^4)$	5133	5092
$100000 (1e^5)$	9592	9512
$500000 (5e^5)$	41538	41246
$1000000 (1e^6)$	78498	78030
$5000000 (5e^6)$	348513	346622

 $(\pi(n))$ = the number of primes less than or equal to n, $\frac{n}{\ln n - 1}$ is used to approximate $\pi(n)$).

5.4 Wilson's theorem

A positive integer *n* is a prime if and only if:

$$(n-1)! \equiv n-1 \pmod{n}$$

5.5 Zeckendorf's theorem

The Zeckendorf's theorem states that every positive integer n can be represented uniquely as a sum of one or more distinct non-consecutive Fibonacci numbers. For example:

$$64 = 55 + 8 + 1$$
$$85 = 55 + 21 + 8 + 1$$

```
vector<int> zeckendoft_theorem(int n) {
       vector < int > fibs = \{1, 1\};
       int sz = 2;
       while (fibs.back() <= n) {</pre>
           fibs.push_back(fibs[sz - 1] + fibs[s - 2]);
           sz++:
       fibs.pop_back();
       vector<int> nums;
       int p = sz - 1;
       while (n > 0) {
           if (n >= fibs[p]) {
               nums.push_back(fibs[p]);
               n -= fibs[p];
           p--;
18
       return nums;
19 }
```

Can Tho University Page 12 of 20

5.6 Bitwise operation

```
• a \mid (a \& b) = a
• a + b = (a \oplus b) + 2(a \& b)
• a \mid b = (a \oplus b) + (a \& b)
                                                                  • a & (a | b) = a
• a \& (b \oplus c) = (a \& b) \oplus (a \& c)
                                                                 • n = 2^k \Leftrightarrow !(n \& (n-1)) = 1
• a \mid (b \& c) = (a \mid b) \& (a \mid c)
                                                                  • -a = \sim a + 1
• a \& (b \mid c) = (a \& b) \mid (a \& c)
                                                                  • 4i \oplus (4i + 1) \oplus (4i + 2) \oplus (4i + 3) = 0
```

• Iterating over all subsets of a set and iterating over all submasks of a mask:

```
1 for (int mask = 0; mask < (1 << n); ++mask) {
      for (int i = 0; i < n; ++i) {
          if (mask & (1 << i)) {</pre>
              // do something...
      }
      // Time complexity: 0(n * 2^n).
8 }
9 for (int mask = 0; mask < (1 << n); ++mask) {
      for (int submask = mask; ; submask = (submask - 1) & mask) {
          // do something...
11
          if (submask == 0) break;
12
      // Time complexity: 0(3^n).
14
15 }
```

Pollard's rho algorithm

```
1 using num_t = long long;
const int PRIME_MAX = (int) 4e4; // for handle numbers <= 1e9.</pre>
3 const int LIMIT = (int) 1e9;
4 vector<int> primes;
5 int small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 73, 113,
      193, 311, 313, 407521, 299210837};
6 void linear_sieve(int n);
7 num_t mulmod(num_t a, num_t b, num_t mod);
8 num_t powmod(num_t a, num_t n, num_t mod);
9 bool miller_rabin(num_t a, num_t d, int s, num_t mod) {
      num_t x = powmod(a, d, mod);
      if (x == mod - 1 || x == 1) {
11
12
          return true:
13
      for (int i = 0; i < s - 1; ++i) {
14
          x = mulmod(x, x, mod);
15
          if (x == mod - 1) return true;
16
17
      return false:
19 }
20 bool is_prime(num_t n, int tests = 10) {
      if (n < 4) return (n > 1);
21
      num_t d = n - 1;
22
      int s = 0;
      while (d % 2 == 0) { d >>= 1; s++; }
```

```
for (int i = 0; i < tests; ++i) {</pre>
                        int a = small_primes[i];
                        if (n == a) return true;
                        if (n % a == 0 || !miller_rabin(a, d, s, n)) return false;
               return true;
|x| = 1 |x| 
               x = mulmod(x, x, mod);
               x += c;
               if (x >= mod) x -= mod;
               return x;
37 }
38 num_t pollard_rho(num_t n, int c) {
               // algorithm to find a random divisor of 'n'.
               // using random function: f(x) = (x^2 + c) \% n.
               num_t x = 2, y = x, d;
              long long p = 1;
               int dist = 0;
               while (true) {
                        y = f(y, c, n);
                        dist++;
                        d = \_gcd(llabs(x - y), n);
                        if (d > 1) break;
                        if (dist == p) { dist = 0; p *= 2; x = y; }
               return d;
| void factorize(int n, vector<num_t> &factors);
54 void llfactorize(num_t n, vector<num_t> &factors) {
               if (n < 2) return;</pre>
               if (is_prime(n)) {
                         factors.emplace_back(n);
                        return;
               if (n < LIMIT) {</pre>
                         factorize(n, factors);
                        return;
               num_t d = n;
               for (int c = 2; d == n; c++) {
                        d = pollard_rho(n, c);
               llfactorize(d, factors);
               llfactorize(n / d, factors);
|71 vector<num_t> gen_divisors(vector<pair<num_t, int>> &factors) {
               vector<num_t> divisors = {1};
               for (auto &x : factors) {
                        int sz = (int) divisors.size();
                        for (int i = 0; i < sz; ++i) {
                                  num_t cur = divisors[i];
                                  for (int j = 0; j < x.second; ++j) {
```

Can Tho University Page 13 of 20

```
cur *= x.first;
                   divisors.push_back(cur);
      return divisors; // this array is NOT sorted yet.
       Bitset sieve
1 /**
     Description: sieve of eratosthenes for large n (up to 1e9).
     Time and space (tested on codeforces):
     + For n = 1e8: ~200 ms, 6 MB.
      + For n = 1e9: ~4000 \text{ ms}, 60 \text{ MB}.
7 const int N = (int) 1e8;
8 bitset<N / 2 + 1> isPrime;
9 void sieve(int n = N) {
      isPrime.flip();
      isPrime[0] = false;
11
      for (int i = 3; i <= (int) sqrt(n); i += 2) {
13
          if (isPrime[i >> 1]) {
               for (int j = i * i; j \le n; j += 2 * i) {
14
                   isPrime[i >> 1] = false:
19 }
20 void example(int n) {
      sieve(n);
21
      int primeCnt = (n >= 2);
22
      for (int i = 3; i \le n; i += 2) {
23
          if (isPrime[i >> 1]) {
24
               primeCnt++;
          }
      cout << primeCnt << '\n';</pre>
28
29 }
 5.9
       Block sieve
     Description: very fast sieve of eratosthenes for large n (up to 1e9).
     Source: kactl.
* Time and space (tested on codeforces):
     + For n = 1e8: ~160 ms, 60 MB.
     + For n = 1e9: ~1600 ms, 505 MB.
     Need to check memory limit.
9 const int N = (int) 1e8;
10 bitset<N + 1> is_prime;
vector<int> fast_sieve() {
      const int S = (int) sqrt(N), R = N / 2;
      vector<int> primes = {2};
```

```
vector<bool> sieve(S + 1, true);
vector<array<int, 2>> cp;
for (int i = 3; i <= S; i += 2) {
    if (sieve[i]) {
        cp.push_back({i, i * i / 2});
        for (int j = i * i; j \le S; j += 2 * i) {
            sieve[i] = false;
       }
   }
for (int L = 1; L <= R; L += S) {
    array<bool, S> block{};
    for (auto &[p, idx] : cp) {
        for (; idx < S + L; idx += p) block[idx - L] = true;</pre>
    for (int i = 0; i < min(S, R - L); ++i) {
        if (!block[i]) primes.push_back((L + i) * 2 + 1);
for (int p : primes) is_prime[p] = true;
return primes;
```

5.10 Combinatorics

5.10.1 Fibonacci numbers

$$F_n = \begin{cases} 0 & \text{if } n = 0\\ 1 & \text{if } n = 1\\ F_{n-1} + F_{n-2} & \text{otherwise} \end{cases}$$

• The first 20 Fibonacci numbers (n = 0, 1, 2, ..., 19):

 $F_n = 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, 2584, 4181$

• Properties:

$$F_{2n+1} = F_n^2 + F_{n+1}^2$$

$$F_{2n} = F_{n-1} \cdot F_n + F_n \cdot F_{n+1}$$

$$F_{n+1} \cdot F_{n-1} - F_n^2 = (-1)^n$$

$$n \mid m \iff F_n \mid F_m$$

$$(F_n, F_m) = F_{(n,m)}$$

5.10.2 Catalan numbers

$$C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{n!(n+1)!}$$

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}, C_0 = 1, C_n = \frac{4n-2}{n+1} C_{n-1}$$

Can Tho University Page 14 of 20

• The first 12 Catalan numbers $(n = 0, 1, 2, \dots, 11)$:

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$$

- Applications of Catalan numbers:
 - difference binary search trees with *n* vertices from 1 to *n*.
 - rooted binary trees with n + 1 leaves (vertices are not numbered).
 - correct bracket sequence of length 2 * n.
 - permutation [n] with no 3-term increasing subsequence (i.e. doesn't exist i < j < k for which a[i] < a[j] < a[k]).
 - ways a convex polygon of n + 2 sides can split into triangles by connecting vertices.

5.10.3 Stirling numbers of the second kind

Partitions of *n* distinct elements into exactly *k* non-empty groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

5.10.4 Derangements

Permutation of the elements of a set, such that no element appears in its original position (no fixied point). Recursive formulas:

$$D(n) = (n-1)[D(n-1) + D(n-2)] = nD(n-1) + (-1)^n$$

6 Linear algebra

6.1 Gauss elimination

```
for (int i=col; i<=m; ++i)</pre>
        swap (a[sel][i], a[row][i]);
    where[col] = row;
    for (int i=0; i<n; ++i)</pre>
        if (i != row) {
             double c = a[i][col] / a[row][col];
             for (int j=col; j<=m; ++j)
                 a[i][j] -= a[row][j] * c;
        }
    ++row;
ans.assign (m, 0);
for (int i=0; i<m; ++i)</pre>
    if (where[i] != -1)
        ans[i] = a[where[i]][m] / a[where[i]][i];
for (int i=0; i<n; ++i) {
    double sum = 0;
    for (int j=0; j<m; ++j)
        sum += ans[j] * a[i][j];
    if (abs (sum - a[i][m]) > EPS)
        return 0;
}
for (int i=0; i<m; ++i)
    if (where[i] == -1)
        return INF;
return 1:
```

7 Geometry

7.1 Fundamentals

7.1.1 **Point**

```
const double PI = acos(-1);
const double EPS = 1e-9;
3 typedef double ftype;
4 struct point {
      ftype x, y;
      point(ftype _x = 0, ftype _y = 0): x(_x), y(_y) {}
      point& operator+=(const point& other) {
         x += other.x; y += other.y; return *this;
     point& operator -=(const point& other) {
         x -= other.x; y -= other.y; return *this;
     point& operator*=(ftype t) {
         x *= t; y *= t; return *this;
     point& operator/=(ftype t) {
         x /= t; y /= t; return *this;
      point operator+(const point& other) const {
         return point(*this) += other;
```

Can Tho University Page 15 of 20

```
21
       point operator-(const point& other) const {
22
23
           return point(*this) -= other;
24
      point operator*(ftype t) const {
25
           return point(*this) *= t;
26
27
       point operator/(ftype t) const {
          return point(*this) /= t;
29
30
31
      point rotate(double angle) const {
           return point(x * cos(angle) - y * sin(angle), x * sin(angle) + y *
       cos(angle));
33
       friend istream& operator>>(istream &in, point &t);
34
       friend ostream& operator<<(ostream &out, const point& t);</pre>
35
       bool operator < (const point& other) const {</pre>
          if (fabs(x - other.x) < EPS)</pre>
37
               return y < other.y;</pre>
38
          return x < other.x:</pre>
39
      }
41 };
42
43 istream& operator>>(istream &in, point &t) {
       in >> t.x >> t.y;
       return in;
46 }
47 ostream& operator<<(ostream &out, const point& t) {
      out << t.x << ' ' << t.y;
      return out;
50 }
52 ftype dot(point a, point b) {return a.x * b.x + a.y * b.y;}
53 ftype norm(point a) {return dot(a, a);}
54 ftype abs(point a) {return sqrt(norm(a));}
55 ftype angle(point a, point b) {return acos(dot(a, b) / (abs(a) * abs(b)));}
56 ftype proj(point a, point b) {return dot(a, b) / abs(b);}
57 ftype cross(point a, point b) {return a.x * b.y - a.y * b.x;}
58 bool ccw(point a, point b, point c) {return cross(b - a, c - a) > EPS;}
59 bool collinear(point a, point b, point c) {return fabs(cross(b - a, c - a)) <</pre>
       EPS:}
60 point intersect(point a1, point d1, point a2, point d2) {
       double t = cross(a2 - a1, d2) / cross(d1, d2);
       return a1 + d1 * t;
63 }
7.1.2 Line
1 struct line {
       double a, b, c;
      line (double _a = 0, double _b = 0, double _c = 0): a(a), b(b), c(c) {}
      friend ostream & operator<<(ostream& out, const line& 1);</pre>
4
5 };
6 ostream & operator << (ostream& out, const line& 1) {
      out << 1.a << ' ' << 1.b << ' ' << 1.c;
```

```
8
      return out:
9 }
10 void pointsToLine(const point& p1, const point& p2, line& 1) {
      if (fabs(p1.x - p2.x) < EPS)
          1 = \{1.0, 0.0, -p1.x\};
      else {
          1.a = - (double)(p1.y - p2.y) / (p1.x - p2.x);
          1.b = 1.0;
          1.c = -1.a * p1.x - 1.b * p1.y;
      }
18 }
19 void pointsSlopeToLine(const point& p, double m, line& 1) {
      1.a = -m;
21
      1.b = 1:
      1.c = -1.a * p.x - 1.b * p.y;
22
23 }
24 bool areParallel(const line& 11, const line& 12) {
      return fabs(l1.a - l2.a) < EPS && fabs(l1.b - l2.b) < EPS;</pre>
bool areSame(const line& 11. const line& 12) {
      return areParallel(l1, l2) && fabs(l1.c - l2.c) < EPS;</pre>
30 bool areIntersect(line l1, line l2, point& p) {
      if (areParallel(l1, l2)) return false;
      p.x = -(11.c * 12.b - 11.b * 12.c) / (11.a * 12.b - 11.b * 12.a);
      if (fabs(11.b) > EPS) p.y = -(11.c + 11.a * p.x);
      else p.y = -(12.c + 12.a * p.x);
      return 1;
37 double distToLine(point p, point a, point b, point& c) {
      double t = dot(p - a, b - a) / norm(b - a);
      c = a + (b - a) * t;
      return abs(c - p);
42 double distToSegment(point p, point a, point b, point& c) {
      double t = dot(p - a, b - a) / norm(b - a);
      if (t > 1.0)
          c = point(b.x, b.y);
      else if (t < 0.0)
          c = point(a.x, a.y);
          c = a + (b - a) * t;
      return abs(c - p);
ftype ABxAC = cross(b - a, c - a);
      ftype ABxAD = cross(b - a, d - a);
      ftype CDxCA = cross(d - c, a - c);
      ftype CDxCB = cross(d - c, b - c);
      if (ABxAC == 0 | | ABxAD == 0 | | CDxCA == 0 | | CDxCB == 0) {
          if (ABxAC == 0 && dot(a - c, b - c) <= 0) return true;
          if (ABxAD == 0 && dot(a - d, b - d) <= 0) return true;
          if (CDxCA == 0 && dot(c - a, d - a) <= 0) return true;</pre>
```

Can Tho University Page 16 of 20

```
if (CDxCB == 0 && dot(c - b, d - b) <= 0) return true;
61
          return false:
62
      }
63
      return (ABxAC * ABxAD < 0 && CDxCA * CDxCB < 0);</pre>
64
65 }
66 void perpendicular(line 11, point p, line& 12) {
      if (fabs(l1.a) < EPS)
          12 = \{1.0, 0.0, -p.x\};
      else {
          12.a = -11.b / 11.a;
70
71
          12.b = 1.0;
          12.c = -12.a * p.x - 12.b * p.y;
72
      }
73
74 }
7.1.3 Circle
int insideCircle(const point& p, const point& center, ftype r) {
      ftype d = norm(p - center);
      ftype rSq = r * r;
      return fabs(d - rSq) < EPS ? 0 : (d - rSq >= EPS ? 1 : -1);
5 }
6 bool circle2PointsR(const point& p1, const point& p2, ftype r, point& c) {
      double h = r * r - norm(p1 - p2) / 4.0;
      if (fabs(h) < 0) return false;</pre>
      h = sqrt(h);
      point perp = (p2 - p1).rotate(PI / 2.0);
      point m = (p1 + p2) / 2.0;
11
      c = m + perp * (h / abs(perp));
      return true;
14 }
7.1.4 Triangle
double areaTriangle(double ab, double bc, double ca) {
      double p = (ab + bc + ca) / 2;
      return sqrt(p) * sqrt(p - ab) * sqrt(p - bc) * sqrt(p - ca);
3
4 }
5 double rInCircle(double ab, double bc, double ca) {
      double p = (ab + bc + ca) / 2;
      return areaTriangle(ab, bc, ca) / p;
8 }
9 double rInCircle(point a, point b, point c) {
      return rInCircle(abs(a - b), abs(b - c), abs(c - a));
11 }
12 bool inCircle(point p1, point p2, point p3, point &ctr, double &r) {
      r = rInCircle(p1, p2, p3);
      if (fabs(r) < EPS) return false;</pre>
14
      line 11, 12;
15
      double ratio = abs(p2 - p1) / abs(p3 - p1);
      point p = p2 + (p3 - p2) * (ratio / (1 + ratio));
      pointsToLine(p1, p, l1);
      ratio = abs(p1 - p2) / abs(p2 - p3);
      p = p1 + (p3 - p1) * (ratio / (1 + ratio));
20
21
      pointsToLine(p2, p, 12);
      areIntersect(l1, l2, ctr);
```

```
23
      return true:
24 }
25 double rCircumCircle(double ab, double bc, double ca) {
      return ab * bc * ca / (4.0 * areaTriangle(ab, bc, ca));
27 }
28 double rCircumCircle(point a, point b, point c) {
       return rCircumCircle(abs(b - a), abs(c - b), abs(a - c));
30 }
 7.1.5 Convex hull
vector<point> CH_Andrew(vector<point> &Pts) { // overall 0(n log n)
      int n = Pts.size(), k = 0;
       vector<point> H(2 * n);
       sort(Pts.begin(), Pts.end());
       for (int i = 0; i < n; ++i) {
          while ((k \ge 2) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
          H[k++] = Pts[i];
      }
      for (int i = n - 2, t = k + 1; i >= 0; --i) {
          while ((k >= t) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
          H[k++] = Pts[i];
      }
      H.resize(k);
      return H;
15 }
 7.1.6 Polygon
double perimeter(const vector<point> &P) {
       double ans = 0.0;
       for (int i = 0; i < (int)P.size() - 1; ++i)</pre>
           ans += abs(P[i] - P[i + 1]);
       return ans;
7 double area(const vector<point> &P) {
       double ans = 0.0;
       for (int i = 0; i < (int)P.size() - 1; ++i)
           ans += (P[i].x * P[i + 1].y - P[i + 1].x * P[i].y);
      return fabs(ans) / 2.0;
12 }
bool isConvex(const vector<point> &P) {
      int n = (int)P.size();
      if (n <= 3) return false;</pre>
      bool firstTurn = ccw(P[0], P[1], P[2]);
       for (int i = 1; i < n - 1; ++i)
          if (ccw(P[i], P[i + 1], P[(i + 2) == n ? 1 : i + 2]) != firstTurn)
              return false;
       return true:
21 }
22 int insidePolygon(point pt, const vector<point> &P) {
      int n = (int)P.size();
      if (n <= 3) return -1;
      bool on_polygon = false;
       for (int i = 0; i < n - 1; ++i)
          if (fabs(abs(P[i] - pt) + abs(pt - P[i + 1]) - abs(P[i] - P[i + 1]))
```

Can Tho University Page 17 of 20

```
< EPS)
               on_polygon = true;
28
      if (on_polygon) return 0;
      double sum = 0.0;
30
      for (int i = 0; i < n - 1; ++i) {
          if (ccw(pt, P[i], P[i + 1]))
32
33
               sum += angle(P[i] - pt, P[i + 1] - pt);
          else
               sum -= angle(P[i] - pt, P[i + 1] - pt);
35
      }
36
37
      return fabs(sum) > PI ? 1 : -1;
38 }
       Minimum enclosing circle
```

36 }

```
1 /**
     Description: computes the minimum circle that encloses all the given
3 */
4 double abs(point a) { return sqrt(a.X * a.X + a.Y * a.Y); }
6 point center_from(double bx, double by, double cx, double cy) {
      double B = bx * bx + by * by, C = cx * cx + cy * cy, D = bx * cy - by *
      return point((cy * B - by * C) / (2 * D), (bx * C - cx * B) / (2 * D));
9 }
11 circle circle_from(point A, point B, point C) {
      point I = center_from(B.X - A.X, B.Y - A.Y, C.X - A.X, C.Y - A.Y);
      return circle(I + A, abs(I));
13
14 }
16 const int N = 100005;
17 int n, x[N], y[N];
18 point a[N];
20 circle emo_welzl(int n, vector<point> T) {
      if (T.size() == 3 || n == 0) {
21
          if (T.size() == 0) return circle(point(0, 0), -1);
22
          if (T.size() == 1) return circle(T[0], 0);
          if (T.size() == 2) return circle((T[0] + T[1]) / 2, abs(T[0] - T[1])
24
      / 2);
          return circle_from(T[0], T[1], T[2]);
25
26
      random\_shuffle(a + 1, a + n + 1);
      circle Result = emo_welzl(0, T);
28
      for (int i = 1; i <= n; i++)</pre>
29
          if (abs(Result.X - a[i]) > Result.Y + 1e-9) {
31
               T.push_back(a[i]);
               Result = emo_welzl(i - 1, T);
32
33
              T.pop_back();
34
      return Result;
35
```

Graph

Bellman-Ford algorithm

```
1 /**
2 * Description: single source shortest path in a weighted (negative or
      positive) directed graph.
   * Time: O(N * M).
   * Tested: https://open.kattis.com/problems/shortestpath3
6 const int64_t INF = (int64_t) 2e18;
7 struct Edge {
      int u, v; // u -> v
      int64_t w;
      Edge() {}
      Edge(int u, int v, int64t w) : u(u), v(v), w(w) {}
vector<int64_t> bellmanFord(int s) {
      // dist[stating] = 0.
      // dist[u] = +INF, if u is unreachable.
      // dist[u] = -INF, if there is a negative cycle on the path from s to u.
      // -INF < dist[u] < +INF, otherwise.</pre>
      vector<int64_t> dist(n, INF);
      dist[s] = 0;
      for (int i = 0; i < n - 1; ++i) {
          bool any = false;
          for (auto [u, v, w] : edges) {
              if (dist[u] != INF && dist[v] > w + dist[u]) {
                  dist[v] = w + dist[u];
                  any = true;
              }
          if (!any) break;
      // handle negative cycles
      for (int i = 0; i < n - 1; ++i) {
          for (auto [u, v, w] : edges) {
              if (dist[u] != INF && dist[v] > w + dist[u]) {
                  dist[v] = -INF;
          }
      }
      return dist;
      Articulation point and Bridge
1 /**
2 * Description: finding articulation points and bridges in a simple
      undirected graph.
3 * Tested: https://oj.vnoi.info/problem/graph_
4 */
5 const int N = (int) 1e5;
6 vector<int> g[N];
| 7 int num[N], low[N], dfs_timer;
```

Can Tho University Page 18 of 20

```
8 bool joint[N];
9 vector<pair<int, int>> bridges;
void dfs(int u, int prev) {
      low[u] = num[u] = ++dfs_timer;
      int child = 0;
12
      for (int v : g[u]) {
13
          if (v == prev) continue;
14
          if (num[v]) low[u] = min(low[u], num[v]);
           else {
16
               dfs(v, u);
17
               low[u] = min(low[u], low[v]);
               child++:
               if (low[v] >= num[v]) {
                   bridges.emplace_back(u, v);
21
22
               if (u != prev && low[v] >= num[u]) joint[u] = true;
23
24
      }
25
      if (u == prev && child > 1) joint[u] = true;
26
27 }
28
29 int main() {
      int n, m;
      cin >> n >> m;
31
      for (int i = 0; i < m; ++i) {
32
33
          int u, v;
           cin >> u >> v;
34
35
          u--; v--;
          g[u].push_back(v);
          g[v].push_back(u);
37
      }
38
      for (int i = 0; i < n; ++i) {</pre>
          if (!num[i]) dfs(i, i);
      }
      return 0;
42
43 }
       Strongly connected components
 8.3
     Description: Tarjan's algorithm finds strongly connected components
      then scc_id[u] == scc_id[v].
     Tested: https://judge.yosupo.jp/problem/scc
```

```
1  /**
2  * Description: Tarjan's algorithm finds strongly connected components
3  * in a directed graph. If vertices u and v belong to the same component,
4  * then scc_id[u] == scc_id[v].
5  * Tested: https://judge.yosupo.jp/problem/scc
6  */
7  const int N = (int) 5e5;
8  vector<int> g[N], st;
9  int low[N], num[N], dfs_timer, scc_id[N], scc;
10  bool used[N];
11  void Tarjan(int u) {
12   low[u] = num[u] = ++dfs_timer;
13   st.push_back(u);
14   for (int v : g[u]) {
15    if (used[v]) continue;
```

```
if (num[v] == 0) {
              Tarjan(v);
17
              low[u] = min(low[u], low[v]);
          else {
              low[u] = min(low[u], num[v]);
      }
      if (low[u] == num[u]) {
          int v;
          do {
              v = st.back(); st.pop_back();
              debug(u, v)
              used[v] = true;
              scc_id[v] = scc;
          } while (v != u);
          scc++;
      }
34 }
      Topo sort
* Description: A topological sort of a directed acyclic graph
      is a linear ordering of its vertices such that for every directed edge
      from vertex u to vertex v, u comes before v in the ordering.
   * Note: If there are cycles, the returned list will have size smaller than n
       (i.e, topo.size() < n).
* Tested: https://judge.yosupo.jp/problem/scc
8 vector<int> topo_sort(const vector<vector<int>> &g) {
      int n = (int) g.size();
      vector<int> indeg(n);
      for (int u = 0; u < n; ++u) {
           for (int v : g[u]) indeg[v]++;
      queue<int> q; // Note: use min-heap to get the smallest lexicographical
      for (int u = 0; u < n; ++u) {
          if (indeg[u] == 0) q.emplace(u);
      vector<int> topo;
      while (!q.empty()) {
          int u = q.front(); q.pop();
          topo.emplace_back(u);
22
          for (int v : g[u]) {
              if (--indeg[v] == 0) q.emplace(v);
25
      return topo;
27 }
      K-th smallest shortest path
```

 $_{1}$ /** Finding the k-th smallest shortest path from vertex s to vertex t,

Can Tho University Page 19 of 20

```
each vertex can be visited more than once.
3 */
4 using adj_list = vector<vector<pair<int, int>>>;
5 vector<int> k_smallest(const adj_list &g, int k, int s, int t) {
      int n = (int) g.size();
      vector<long long> ans;
      vector<int> cnt(n);
      using pli = pair<long long, int>;
      priority_queue<pli, vector<pli>, greater<pli>>> pq;
10
      pq.emplace(0, s);
11
12
      while (!pq.empty() && cnt[t] < k) {</pre>
           int u = pq.top().second;
13
          long long d = pq.top().first;
14
15
          pq.pop();
          if (cnt[u] == k) continue;
           cnt[u]++;
          if (u == t) {
18
               ans.push_back(d);
19
20
21
           for (auto [v, cost] : g[u]) {
               pq.emplace(d + cost, v);
          }
23
24
25
      assert(ans.size() == k);
      return ans;
27 }
```

8.6 Eulerian path

8.6.1 Directed graph

```
1 /**
2 * Hierholzer's algorithm.
     Description: An Eulerian path in a directed graph is a path that visits
       all edges exactly once.
      An Eulerian cycle is a Eulerian path that is a cycle.
     Time complexity: O(|E|).
6 */
   vector<int> find_path_directed(const vector<vector<int>> &g, int s) {
      int n = (int) g.size();
      vector<int> stack, cur_edge(n), vertices;
      stack.push_back(s);
10
11
      while (!stack.empty()) {
           int u = stack.back();
12
13
           stack.pop_back();
          while (cur_edge[u] < (int) g[u].size()) {</pre>
14
15
               stack.push_back(u);
               u = g[u][cur\_edge[u]++];
17
           vertices.push_back(u);
18
19
      reverse(vertices.begin(), vertices.end());
20
      return vertices;
21
22 }
```

8.6.2 Undirected graph

```
1 /**
* Hierholzer's algorithm.
3 * Description: An Eulerian path in a undirected graph is a path that visits
      all edges exactly once.
      An Eulerian cycle is a Eulerian path that is a cycle.
   * Time complexity: O(|E|).
6 */
7 struct Edge {
      int to;
      list<Edge>::iterator reverse_edge;
      Edge(int _to) : to(_to) {}
12 vector<int> vertices;
13 void find_path(vector<list<Edge>> &g, int u) {
      while (!g[u].empty()) {
          int v = g[u].front().to;
          g[v].erase(g[u].front().reverse_edge);
          g[u].pop_front();
          find_path(g, v);
      vertices.emplace_back(u); // reversion list.
22 void add_edge(int u, int v) {
      g[u].emplace_front(v);
      g[v].emplace_front(u);
      g[u].front().reverse_edge = g[v].begin();
      g[v].front().reverse_edge = g[u].begin();
```

9 Misc.

9.1 Ternary search

```
const double eps = 1e-9;
2 double ternary_search_max(double 1, double r) {
       // find x0 such that: f(x0) > f(x), \all x: l \ll x \ll r.
       while (r - 1 > eps) {
           double mid1 = 1 + (r - 1) / 3;
          double mid2 = r - (r - 1) / 3;
          if (f(mid1) < f(mid2)) l = mid1;
          else r = mid2:
      }
      return 1;
12 double ternary_search_min(double 1, double r) {
       // find x0 such that: f(x0) < f(x), \all x: 1 <= x <= r.
       while (r - 1 > eps) {
           double mid1 = 1 + (r - 1) / 3;
          double mid2 = r - (r - 1) / 3;
17
          if (f(mid1) > f(mid2)) 1 = mid1;
           else r = mid2;
      }
      return 1;
```

Can Tho University Page 20 of 20

```
21 }
9.2
      Dutch flag national problem
void dutch_flag_national(vector<int> &arr) {
      // All elements that are LESS than pivot are moved to the LEFT.
      // All elements that are GREATER than pivot are moved to the RIGHT.
      // E.g. [1, 2, 0, 0, 2, 2, 1], pivot = 1 -> [0, 0, 1, 1, 2, 2, 2].
      int n = (int) arr.size();
      int i = 0, j = 0, k = n - 1;
      int pivot = 1;
      // 0....i....j....k....n
      while (j <= k) {
          if (arr[j] < pivot) {</pre>
               swap(arr[i], arr[j]);
11
12
              i++;
              j++;
13
           else if (arr[j] > pivot) {
15
               swap(arr[j], arr[k]);
              k--;
          }
           else {
              j++;
20
          }
21
22
      // 0 <= index <= i - 1: arr[index] < mid.
      // i <= index <= k: arr[index] = mid.</pre>
24
      // k + 1 \le index < sz: arr[index] > mid.
25
26 }
      Matrix
 9.3
1 struct Matrix {
      static const matrix_type INF = numeric_limits<matrix_type>::max();
      vector<vector<matrix_type>> mat;
      Matrix(int _N, int _M, matrix_type v = 0) : N(_N), M(_M) {
          mat.assign(N, vector<matrix_type>(M, v));
      static Matrix identity(int n) { // return identity matrix.
          Matrix I(n, n);
10
           for (int i = 0; i < n; ++i) {
11
               I[i][i] = 1;
          return I;
14
      }
15
      vector<matrix_type>& operator[](int r) { return mat[r]; }
17
      const vector<matrix_type>& operator[](int r) const { return mat[r]; }
18
      Matrix& operator*=(const Matrix &other) {
20
          assert(M == other.N); // [N x M] [other.N x other.M]
21
          Matrix res(N, other.M);
```

```
for (int r = 0; r < N; ++r) {
    for (int c = 0; c < other.M; ++c) {
        long long square_mod = (long long) MOD * MOD;
        long long sum = 0;
        for (int g = 0; g < M; ++g) {
            sum += (long long) mat[r][g] * other[g][c];
            if (sum >= square_mod) sum -= square_mod;
        }
        res[r][c] = sum % MOD;
}

mat.swap(res.mat); return *this;
}
```