ACM - ICPC 2022

TEAM NOTEBOOK Can Tho University

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1 Contest

1.1 C++

```
#include <bits/stdc++.h>
using namespace std;
4 #ifdef LOCAL
5 #include "cp/debug.h"
6 #else
7 #define debug(...)
8 #endif
nt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
12 const int MOD = (int) 1e9 + 7;
13 const int INF = 0x3f3f3f3f3f;
14
15 int main() {
      ios::sync_with_stdio(false); cin.tie(nullptr);
      // freopen("input.txt", "r", stdin);
17
18
      // freopen("output.txt", "w", stdout);
20
      return 0;
21 }
1.2 Debug
#define debug(...) { string _s = #__VA_ARGS__; replace(begin(_s), end(_s),
       ',', ''); stringstream _ss(_s); istream_iterator<string> _it(_ss);
      out_error(_it, __VA_ARGS__);}
void out_error(istream_iterator<string> it) { cerr << '\n'; }</pre>
5 template < typename T, typename ...Args >
6 void out_error(istream_iterator<string> it, T a, Args... args) {
      cerr << " [" << *it << " = " << a << "] ";
      out_error(++it, args...);
9 }
11 template<typename T, typename G> ostream& operator<<(ostream &os, const
      pair<T, G> &p) {
      return os << "(" << p.first << ", " << p.second << ")";</pre>
13 }
15 template < class Con, class = decltype(begin(declval < Con > ())) >
16 typename enable_if<!is_same<Con, string>::value, ostream&>::type
operator << (ostream& os, const Con& container) {</pre>
      os << "{";
19
      for (auto it = container.begin(); it != container.end(); ++it)
          os << (it == container.begin() ? "" : ", ") << *it;
20
      return os << "}";</pre>
21
22 }
```

1.3 **Java**

```
import java.io.BufferedReader;
import java.util.StringTokenizer;
3 import java.io.IOException;
4 import java.io.InputStreamReader;
5 import java.io.PrintWriter;
6 import java.util.ArrayList;
7 import java.util.Arrays;
8 import java.util.Collections;
9 import java.util.Random;
 public class Main {
     public static void main(String[] args) {
          FastScanner fs = new FastScanner();
          PrintWriter out = new PrintWriter(System.out);
          int n = fs.nextInt();
          out.println(n):
          out.close(); // don't forget this line.
      static class FastScanner {
          BufferedReader br;
          StringTokenizer st:
          public FastScanner() {
              br = new BufferedReader(new InputStreamReader(System.in));
              st = null;
          public String next() {
              while (st == null || st.hasMoreTokens() == false) {
                  try {
                      st = new StringTokenizer(br.readLine());
                  catch (IOException e) {
                      throw new RuntimeException(e);
              return st.nextToken();
         }
          public int nextInt() {
              return Integer.parseInt(next());
         }
         public long nextLong() {
              return Long.parseLong(next());
         }
          public double nextDouble() {
              return Double.parseDouble(next());
     }
```

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1.4 sublime-build

1 {

11 }

25 }

```
"cmd": ["g++", "-std=c++17", "-fmax-errors=5", "-DLOCAL", "-Wall", "-
Wextra", "-o", "${file_path}/${file_base_name}.out", "${file}"],

"file_regex": "^(...[^:]*):([0-9]+):?([0-9]+)?:? (.*)$",

"working_dir": "${file_path}",

"selector": "source.cpp, source.c++"

1.5 .bashrc

alias cpp='g++ -std=c++17 -fmax-errors=5 -DLOCAL -Wall -Wextra'

#Stress-testing
```

out && diff -Z out ans && echo "Test \$i passed!!" || break;

./gen.out > in && ./"\$CHECKER.out" < in > ans && ./"\$SOL.out" < in >

2 Data structures

for i in {1..100};

2.1 Sparse table

4 function test {

CHECKER=\$2

SOL = 1

```
1 int st[MAXN][K + 1];
2 for (int i = 0; i < N; i++) {
      st[i][0] = f(array[i]);
4 }
5 for (int j = 1; j \le K; j++) {
      for (int i = 0; i + (1 << j) <= N; i++) {
          st[i][j] = f(st[i][j-1], st[i+(1 << (j-1))][j-1]);
9 }
10 // Range Minimum Queries.
int lg[MAXN + 1];
12 \log[1] = 0;
13 for (int i = 2; i \le MAXN; i++) {
      lg[i] = lg[i / 2] + 1;
14
15 }
int j = lg[R - L + 1];
int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);</pre>
18 // Range Sum Queries.
19 long long sum = 0;
20 for (int j = K; j >= 0; j--) {
      if ((1 << j) <= R - L + 1) {
          sum += st[L][i];
22
          L += 1 << j;
23
24
```

2.2 Ordered set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
5 template<typename key_type>
using set_t = tree<key_type, null_type, less<key_type>, rb_tree_tag,
      tree_order_statistics_node_update>;
9 void example() {
      vector < int > nums = \{1, 2, 3, 5, 10\};
      set_t<int> st(nums.begin(), nums.end());
      cout << *st.find_by_order(0) << '\n'; // 1</pre>
      assert(st.find_by_order(-INF) == st.end());
      assert(st.find_by_order(INF) == st.end());
      cout << st.order_of_key(2) << '\n'; // 1</pre>
      cout << st.order_of_key(4) << '\n'; // 3</pre>
      cout << st.order_of_key(9) << '\n'; // 4</pre>
      cout << st.order_of_key(-INF) << '\n'; // 0</pre>
      cout << st.order_of_key(INF) << '\n'; // 5</pre>
22 }
2.3
     Dsu
1 struct Dsu {
      int n;
      vector<int> par, sz;
      Dsu(int _n) : n(_n) {
          sz.resize(n, 1);
          par.resize(n);
          iota(par.begin(), par.end(), 0);
      int find(int v) {
          // finding leader/parrent of set that contains the element v.
          // with {path compression optimization}.
          return (v == par[v] ? v : par[v] = find(par[v]));
      bool same(int u, int v) {
          return find(u) == find(v);
      bool unite(int u, int v) {
          u = find(u); v = find(v);
          if (u == v) return false;
          if (sz[u] < sz[v]) swap(u, v);
          par[v] = u;
          sz[u] += sz[v];
          return true:
      vector<vector<int>> groups() {
          // returns the list of the "list of the vertices in a connected
```

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```
component".
                                                                                             if (lazy_set[id] == 0 && lazy_add[id] == 0) return;
          vector<int> leader(n);
                                                                                  32
                                                                                             int mid = (1 + r) >> 1;
27
                                                                                             for (int child : {id * 2, id * 2 + 1}) {
          for (int i = 0; i < n; ++i) {
              leader[i] = find(i);
                                                                                                 int range = (child == id * 2 ? mid - l + 1 : r - mid);
          }
                                                                                                 if (lazy_set[id] != 0) {
          vector<int> id(n, -1);
                                                                                                     lazy_add[child] = 0;
31
          int count = 0;
                                                                                                     lazy_set[child] = lazy_set[id];
32
          for (int i = 0; i < n; ++i) {
                                                                                                     tree[child] = range * lazy_set[id];
33
              if (id[leader[i]] == -1) {
34
                                                                                                 lazy_add[child] += lazy_add[id];
                  id[leader[i]] = count++;
              }
                                                                                                 tree[child] += range * lazy_add[id];
          }
          vector<vector<int>> result(count);
                                                                                             lazy_add[id] = lazy_set[id] = 0;
          for (int i = 0; i < n; ++i) {
39
              result[id[leader[i]]].push_back(i);
                                                                                        }
40
          }
                                                                                         void update(int id, int l, int r, int u, int v, int amount, bool
41
42
          return result;
                                                                                         set_value = false) {
                                                                                             if (r < u \mid \mid 1 > v) return;
43
44 };
                                                                                             if (u <= 1 && r <= v) {
                                                                                                 if (set_value) {
      Segment tree
 2.4
                                                                                                     tree[id] = 1LL * amount * (r - l + 1);
                                                                                                     lazy_set[id] = amount;
                                                                                                     lazy_add[id] = 0; // clear all previous updates.
2 * Description: A segment tree with range updates and sum queries that
      supports three types of operations:
                                                                                                 else {
  * + Increase each value in range [1, r] by x (i.e. a[i] += x).
                                                                                                     tree[id] += 1LL * amount * (r - l + 1);
  * + Set each value in range [1, r] to x (i.e. a[i] = x).
                                                                                                     lazy_add[id] += amount;
   * + Determine the sum of values in range [1, r].
                                                                                                 }
6 */
                                                                                                 return:
7 struct SegmentTree {
                                                                                             }
      int n;
                                                                                             push(id, 1, r);
      vector<long long> tree, lazy_add, lazy_set;
                                                                                             int mid = (1 + r) >> 1;
      SegmentTree(int _n) : n(_n) {
10
                                                                                             update(id * 2, 1, mid, u, v, amount, set_value);
11
          int p = 1:
                                                                                             update(id * 2 + 1, mid + 1, r, u, v, amount, set_value);
          while (p < n) p *= 2;
12
                                                                                             tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
13
          tree.resize(p * 2);
          lazy_add.resize(p * 2);
14
                                                                                        long long get(int id, int l, int r, int u, int v) {
          lazy_set.resize(p * 2);
15
                                                                                             if (r < u \mid \mid 1 > v) return 0;
16
                                                                                             if (u <= 1 && r <= v) {
      long long merge(const long long &left, const long long &right) {
17
                                                                                                 return tree[id];
          return left + right;
18
19
                                                                                             push(id, 1, r);
20
      void build(int id, int 1, int r, const vector<int> &arr) {
                                                                                             int mid = (1 + r) >> 1;
          if (1 == r) {
21
                                                                                             long long left = get(id * 2, 1, mid, u, v);
              tree[id] += arr[l];
22
                                                                                             long long right = get(id * 2 + 1, mid + 1, r, u, v);
              return;
23
                                                                                             return merge(left, right);
24
          }
          int mid = (1 + r) >> 1;
25
                                                                                  77 };
          build(id * 2, 1, mid, arr);
          build(id * 2 + 1, mid + 1, r, arr);
27
                                                                                         Efficient segment tree
          tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
28
                                                                                   1 template < typename T> struct SegmentTree {
29
      void push(int id, int 1, int r) {
                                                                                         int n;
```

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```
vector<T> tree:
                                                                                            tree.emplace_back(left, right, tree[left].val + tree[right].val);
      SegmentTree(int _n) : n(_n), tree(2 * n) {}
                                                                                 20
                                                                                            return tree.size() - 1;
      T merge(const T &left, const T &right) {
          return left + right;
                                                                                        int add(int x, int 1, int r, int u, int v, int amt) {
                                                                                            if (1 > v \mid | r < u) return x;
      template<tvpename G>
                                                                                            if (u <= 1 && r <= v) {
      void build(const vector<G> &initial) {
                                                                                                 tree.emplace_back(tree[x].l, tree[x].r, tree[x].val + 1LL * amt
                                                                                          * (r - l + 1), tree[x].lazy + amt);
          assert((int) initial.size() == n);
10
                                                                                                tree.back().has_changed = true;
          for (int i = 0; i < n; ++i) {
11
              tree[i + n] = initial[i];
                                                                                                return tree.size() - 1;
12
13
          for (int i = n - 1; i > 0; --i) {
                                                                                            int mid = (1 + r) >> 1:
14
              tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
                                                                                            push(x, 1, mid, r);
15
                                                                                            int left = add(tree[x].1, 1, mid, u, v, amt);
          }
16
17
      }
                                                                                            int right = add(tree[x].r, mid + 1, r, u, v, amt);
      void modify(int i, int v) {
18
                                                                                            tree.emplace_back(left, right, tree[left].val + tree[right].val, 0)
19
          tree[i += n] = v;
          for (i /= 2; i > 0; i /= 2) {
                                                                                            return tree.size() - 1;
20
21
              tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
          }
                                                                                        long long get_sum(int x, int l, int r, int u, int v) {
22
23
                                                                                            if (r < u \mid | 1 > v) return 0;
      T get_sum(int 1, int r) {
                                                                                            if (u <= 1 && r <= v) return tree[x].val;
24
25
          // sum of elements from 1 to r - 1.
                                                                                            int mid = (1 + r) / 2;
                                                                                            push(x, 1, mid, r);
          T ret{};
26
                                                                                            return get_sum(tree[x].1, 1, mid, u, v) + get_sum(tree[x].r, mid +
          for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
              if (1 & 1) ret = merge(ret, tree[1++]);
                                                                                         1, r, u, v);
28
              if (r & 1) ret = merge(ret, tree[--r]);
29
          }
                                                                                        void push(int x, int 1, int mid, int r) {
30
                                                                                            if (!tree[x].has_changed) return;
          return ret;
31
                                                                                            Vertex left = tree[tree[x].1];
32
33 };
                                                                                            Vertex right = tree[tree[x].r];
                                                                                            tree.emplace_back(left);
2.6 Persistent lazy segment tree
                                                                                            tree[x].l = tree.size() - 1;
                                                                                            tree.emplace_back(right);
1 struct Vertex {
                                                                                            tree[x].r = tree.size() - 1;
      int 1, r;
      long long val, lazy;
                                                                                             tree[tree[x].1].val += tree[x].lazy * (mid - 1 + 1);
      bool has_changed = false;
                                                                                             tree[tree[x].1].lazy += tree[x].lazy;
      Vertex() {}
      Vertex(int _1, int _r, long long _val, int _lazy = 0) : l(_l), r(_r),
                                                                                            tree[tree[x].r].val += tree[x].lazy * (r - mid);
      val(_val), lazy(_lazy) {}
                                                                                            tree[tree[x].r].lazy += tree[x].lazy;
7 };
8 struct PerSegmentTree {
                                                                                            tree[tree[x].1].has_changed = true;
      vector<Vertex> tree;
                                                                                            tree[tree[x].r].has_changed = true;
      vector<int> root:
                                                                                            tree[x].lazy = 0;
      int build(const vector<int> &arr, int 1, int r) {
11
                                                                                             tree[x].has_changed = false;
          if (1 == r) {
12
              tree.emplace_back(-1, -1, arr[l]);
13
                                                                                  63 };
```

CTU.NegativeZero

Disjoint sparse table

² * Description: range query on a static array.

14

15

16

17

}

return tree.size() - 1;

int left = build(arr, 1, mid);

int right = build(arr, mid + 1, r);

int mid = (1 + r) / 2;

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```
* Time: 0(1) per query.
                                                                                            fenw.assign(n, 0); // normal fenwick tree.
4 * Tested: stress-test.
                                                                                        }
5 */
                                                                                        void build(const vector<int> &A) {
6 const int MOD = (int) 1e9 + 7;
                                                                                             assert((int) A.size() == n);
7 struct DisjointSparseTable { // product queries.
                                                                                            vector<int> diff(n);
      int n. h:
                                                                                            diff[0] = A[0];
      vector<vector<int>> dst;
                                                                                             for (int i = 1; i < n; ++i) {
      vector<int> lq;
                                                                                                 diff[i] = A[i] - A[i - 1];
10
      DisjointSparseTable(int _n) : n(_n) {
11
          h = 1; // in case n = 1: h = 0 !!.
                                                                                             fenw_coeff[0] = (long long) diff[0] * n;
12
          int p = 1;
                                                                                             fenw[0] = diff[0];
13
          while (p < n) p *= 2, h++;
                                                                                            for (int i = 1; i < n; ++i) {
14
                                                                                                 fenw_coeff[i] = fenw_coeff[i - 1] + (long long) diff[i] * (n -
          lg.resize(p); lg[1] = 0;
          for (int i = 2; i < p; ++i) {
                                                                                        i);
16
              lg[i] = 1 + lg[i / 2];
                                                                                                fenw[i] = fenw[i - 1] + diff[i];
18
          dst.resize(h, vector<int>(n));
                                                                                            for (int i = n - 1; i >= 0; --i) {
19
                                                                                                int j = (i & (i + 1)) - 1;
20
21
      void build(const vector<int> &A) {
                                                                                                if (j >= 0) {
          for (int lv = 0; lv < h; ++lv) {
                                                                                                     fenw_coeff[i] -= fenw_coeff[j];
22
              int len = (1 << lv);
                                                                                                     fenw[i] -= fenw[j];
23
              for (int k = 0; k < n; k += len * 2) {
                                                                                                }
24
                  int mid = min(k + len, n);
                                                                                            }
25
                  dst[lv][mid - 1] = A[mid - 1] % MOD;
                                                                                        }
26
                  for (int i = mid - 2; i >= k; --i) {
                                                                                         void add(vector<tree_type> &fenw, int i, tree_type val) {
                       dst[lv][i] = 1LL * A[i] * dst[lv][i + 1] % MOD;
                                                                                            while (i < n) {
                                                                                                fenw[i] += val;
                  if (mid == n) break;
                                                                                                i = (i + 1);
                  dst[lv][mid] = A[mid] % MOD;
                  for (int i = mid + 1; i < min(mid + len, n); ++i) {
                                                                                        }
                       dst[lv][i] = 1LL * A[i] * dst[lv][i - 1] % MOD;
                                                                                         tree_type __prefix_sum(vector<tree_type> &fenw, int i) {
                                                                                            tree_type res{};
34
              }
                                                                                            while (i >= 0) {
35
36
                                                                                                res += fenw[i];
                                                                                                i = (i \& (i + 1)) - 1;
37
      int get(int 1, int r) {
38
          if (1 == r) {
                                                                                            return res;
39
              return dst[0][1];
                                                                                         tree_type prefix_sum(int i) {
42
          int i = lg[l ^ r];
                                                                                            return __prefix_sum(fenw_coeff, i) - __prefix_sum(fenw, i) * (n - i
          return 1LL * dst[i][l] * dst[i][r] % MOD;
                                                                                         - 1):
43
44
45 };
                                                                                         void range_add(int 1, int r, tree_type val) {
                                                                                             add(fenw_coeff, 1, (n - 1) * val);
      Fenwick tree
                                                                                             add(fenw\_coeff, r + 1, (n - r - 1) * (-val));
                                                                                             add(fenw, 1, val);
using tree_type = long long;
                                                                                             add(fenw, r + 1, -val);
2 struct FenwickTree {
      int n;
                                                                                         tree_type range_sum(int 1, int r) {
      vector<tree_type> fenw_coeff, fenw;
                                                                                            return prefix_sum(r) - prefix_sum(l - 1);
      FenwickTree() {}
      FenwickTree(int _n) : n(_n) {
          fenw_coeff.assign(n, 0); // fenwick tree with coefficient (n - i).
```

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3 Mathematics

3.1 Trigonometry

3.1.1 Sum - difference identities

$$\sin(u \pm v) = \sin(u)\cos(v) \pm \cos(u)\sin(v)$$

$$\cos(u \pm v) = \cos(u)\cos(v) \mp \sin(u)\sin(v)$$

$$\tan(u \pm v) = \frac{\tan(u) \pm \tan(v)}{1 \mp \tan(u)\tan(v)}$$

3.1.2 Sum to product identities

$$\cos(u) + \cos(v) = 2\cos(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\cos(u) - \cos(v) = -2\sin(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

$$\sin(u) + \sin(v) = 2\sin(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\sin(u) - \sin(v) = 2\cos(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

3.1.3 Product identities

$$\cos(u)\cos(v) = \frac{1}{2}[\cos(u+v) + \cos(u-v)]$$

$$\sin(u)\sin(v) = -\frac{1}{2}[\cos(u+v) - \cos(u-v)]$$

$$\sin(u)\cos(v) = \frac{1}{2}[\sin(u+v) + \sin(u-v)]$$

3.1.4 Double - triple angle identities

$$\sin(2u) = 2\sin(u)\cos(u)$$

$$\cos(2u) = 2\cos^{2}(u) - 1 = 1 - 2\sin^{2}(u)$$

$$\tan(2u) = \frac{2\tan(u)}{1 - \tan^{2}(u)}$$

$$\sin(3u) = 3\sin(u) - 4\sin^{3}(u)$$

$$\cos(3u) = 4\cos^{3}(u) - 3\cos(u)$$

$$\tan(3u) = \frac{3\tan(u) - \tan^{3}(u)}{1 - 3\tan^{2}(u)}$$

3.2 Sums

$$n^{a} + n^{a+1} + \dots + n^{b} = \frac{n^{b+1} - n^{a}}{n-1}, \ n \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(n+1)(2n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \left(\frac{n(n+1)}{2}\right)^{2}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

4 String

4.1 Prefix function

for (int i = 0; i < n; ++i) {

ans[pi[i]]++;

```
Description: The prefix function of a string 's' is defined as an array
     pi of length n,
     where pi[i] is the length of the longest proper prefix of the substring
     s[0..i] which is also a suffix of this substring.
    Time complexity: O(|S|).
  */
vector<int> prefix_function(const string &s) {
     int n = (int) s.length();
     vector<int> pi(n);
     pi[0] = 0;
     for (int i = 1; i < n; ++i) {
         int j = pi[i - 1]; // try length pi[i - 1] + 1.
         while (j > 0 && s[j] != s[i]) {
             j = pi[j - 1];
         if (s[j] == s[i]) {
             pi[i] = j + 1;
         }
     return pi;
     Counting occurrences of each prefix
vector<int> count_occurrences(const string &s) {
     vector<int> pi = prefix_function(s);
     int n = (int) s.size();
     vector<int> ans(n + 1);
```

}

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```
for (int i = n - 1; i > 0; --i) {
                                                                                                // counting sort.
          ans[pi[i - 1]] += ans[i];
                                                                                                 for (int i = 0; i < n; ++i) norder[i] = (sa[i] - k + n) % n;
                                                                                                for (int i = 0; i < n; ++i) cnt[rank[i]]++;</pre>
10
      for (int i = 0; i <= n; ++i) {
                                                                                                for (int i = 1; i < lim; ++i) cnt[i] += cnt[i - 1];</pre>
11
                                                                                                for (int i = n - 1; i >= 0; --i) sa[--cnt[rank[norder[i]]]] =
          ans[i]++;
12
                                                                                        norder[i]:
13
      return ans;
                                                                                                rank[sa[0]] = rank\_cnt = 0;
14
      // Input: ABACABA
                                                                                                for (int i = 1; i < n; ++i) {
15
      // Output: 4 2 2 1 1 1 1
                                                                                                     int u = sa[i], v = sa[i - 1];
16
17 }
                                                                                                     int nu = u + k, nv = v + k;
                                                                                                     if (nu >= n) nu -= n;
      Knuth-Morris-Pratt algorithm
                                                                                                    if (nv >= n) nv -= n:
                                                                                                    if (rank[u] != rank[v] || rank[nu] != rank[nv]) ++rank_cnt;
                                                                                                     nrank[sa[i]] = rank_cnt;
  * Searching for a substring in a string.
                                                                                                }
3 * Time complexity: O(N + M).
                                                                                                for (int i = 0; i < rank_cnt + 1; ++i) cnt[i] = 0;</pre>
                                                                                                rank.swap(nrank);
5 vector<int> KMP(const string &text, const string &pattern) {
                                                                                            }
      int n = (int) text.length();
                                                                                        }
      int m = (int) pattern.length();
                                                                                        void kasai() {
      string s = pattern + '$' + text;
                                                                                             for (int i = 0; i < n; ++i) rank[sa[i]] = i;
      vector<int> pi = prefix_function(s);
                                                                                            for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
10
      vector<int> indices;
                                                                                                int j = sa[rank[i] - 1];
      for (int i = 0; i < (int) s.length(); ++i) {
11
                                                                                                while (s[i + k] == s[j + k]) k++;
          if (pi[i] == m) {
12
                                                                                                lcp[rank[i]] = k;
              indices.push_back(i - 2 * m);
13
          }
14
                                                                                            // Note: lcp[i] = longest common prefix(sa[i - 1], sa[i]).
15
                                                                                        }
      return indices;
16
17 }
                                                                                        Manacher's algorithm
 4.4 Suffix array
struct SuffixArray {
                                                                                  2 * Description: for each position, computes d[0][i] = half length of
      string s;
                                                                                  3 longest palindrome centered on i (rounded up), d[1][i] = half length of
      int n, lim;
                                                                                  4 longest palindrome centered on i and i - 1.
      vector<int> sa, lcp, rank;
                                                                                     * Time complexity: O(N).
      SuffixArray(const string &_s, int _lim = 256) : s(_s), n(s.length() +
                                                                                  * Tested: https://judge.yosupo.jp/problem/enumerate_palindromes, stress-
      1), lim(_lim), sa(n), lcp(n), rank(n) {
                                                                                        tested.
          s += '$';
          build();
                                                                                  8 array<vector<int>, 2> manacher(const string &s) {
          kasai();
                                                                                        int n = (int) s.size();
          sa.erase(sa.begin());
                                                                                        array<vector<int>, 2> d;
          lcp.erase(lcp.begin());
10
                                                                                        for (int z = 0; z < 2; ++z) {
          s.pop_back();
11
                                                                                            d[z].resize(n);
12
                                                                                            int 1 = 0, r = 0;
13
      void build() {
                                                                                            for (int i = 0; i < n; ++i) {
          vector<int> nrank(n), norder(n), cnt(max(n, lim));
14
                                                                                                int mirror = l + r - i + z;
15
          for (int i = 0; i < n; ++i) {
                                                                                                d[z][i] = (i > r ? 0 : min(d[z][mirror], r - i));
              sa[i] = i; rank[i] = s[i];
16
                                                                                                int L = i - d[z][i] - z, R = i + d[z][i];
17
                                                                                                while (L >= 0 \&\& R < n \&\& s[L] == s[R])  {
          for (int k = 0, rank_cnt = 0; rank_cnt < n - 1; k = max(1, k * 2),
18
```

lim = rank_cnt + 1) {

d[z][i]++; L--; R++;

}

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```
if (R > r) {
22
                   1 = L; r = R;
          }
25
26
      return d;
27 }
      Trie
 4.6
struct Trie {
      const static int ALPHABET = 26;
      const static char minChar = 'a';
      struct Vertex {
          int next[ALPHABET];
          bool leaf;
          Vertex() {
              leaf = false;
               fill(next, next + ALPHABET, -1);
          }
10
      };
11
      vector<Vertex> trie;
12
13
      Trie() { trie.emplace_back(); }
14
      void insert(const string &s) {
15
          int i = 0:
16
           for (const char &ch : s) {
               int j = ch - minChar;
18
              if (trie[i].next[j] == -1) {
19
20
                   trie[i].next[j] = trie.size();
21
                   trie.emplace_back();
              i = trie[i].next[j];
23
24
          }
          trie[i].leaf = true;
25
26
      bool find(const string &s) {
27
          int i = 0;
28
          for (const char &ch : s) {
29
               int j = ch - minChar;
              if (trie[i].next[j] == -1) {
31
                   return false;
32
33
              i = trie[i].next[j];
          return (trie[i].leaf ? true : false);
36
37
38 };
4.7 Hashing
1 struct Hash61 {
      static const uint64_t MOD = (1LL << 61) - 1;</pre>
```

```
static uint64_t BASE;
static vector<uint64_t> pw;
uint64_t addmod(uint64_t a, uint64_t b) const {
    a += b:
    if (a >= MOD) a -= MOD;
    return a;
uint64_t submod(uint64_t a, uint64_t b) const {
    a += MOD - b;
    if (a >= MOD) a -= MOD;
    return a;
uint64_t mulmod(uint64_t a, uint64_t b) const {
    uint64_t low1 = (uint32_t) a, high1 = (a >> 32);
    uint64_t low2 = (uint32_t) b, high2 = (b >> 32);
    uint64_t low = low1 * low2;
    uint64_t mid = low1 * high2 + low2 * high1;
    uint64_t high = high1 * high2;
    uint64_t ret = (low & MOD) + (low >> 61) + (high << 3) + (mid >>
29) + (mid << 35 >> 3) + 1;
    // ret %= MOD:
    ret = (ret >> 61) + (ret & MOD);
    ret = (ret >> 61) + (ret & MOD);
    return ret - 1;
void ensure_pw(int m) {
    int n = (int) pw.size();
    if (n >= m) return;
    pw.resize(m);
    for (int i = n; i < m; ++i) {
        pw[i] = mulmod(pw[i - 1], BASE);
}
vector<uint64_t> pref;
template < typename T > Hash61(const T &s) { // strings or arrays.
    n = (int) s.size();
    ensure_pw(n);
    pref.resize(n + 1);
    pref[0] = 0;
    for (int i = 0; i < n; ++i) {</pre>
        pref[i + 1] = addmod(mulmod(pref[i], BASE), s[i]);
    }
inline uint64_t operator()(const int from, const int to) const {
    assert(0 \le from \&\& from \le to \&\& to < n);
    // pref[to + 1] - pref[from] * pw[to - from + 1]
    return submod(pref[to + 1], mulmod(pref[from], pw[to - from + 1]));
```

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5 Number Theory

5.1 Euler's totient function

- Euler's totient function, also known as **phi-function** $\phi(n)$ counts the number of integers between 1 and n inclusive, that are **coprime to** n.
- Properties:
 - Divisor sum property: $\sum_{d|n} \phi(d) = n$.
 - $\phi(n)$ is a **prime number** when n = 3, 4, 6.
 - If *p* is a prime number, then $\phi(p) = p 1$.
 - If *p* is a prime number and *k* ≥ 1, then $\phi(p^k) = p^k p^{k-1}$.
 - If *a* and *b* are **coprime**, then $\phi(ab) = \phi(a) \cdot \phi(b)$.
 - In general, for **not coprime** a and b, with d = gcd(a, b) this equation holds: $\phi(ab) = \phi(a) \cdot \phi(b) \cdot \frac{d}{\phi(d)}$.
 - With $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\phi(n) = \phi(p_1^{k_1}) \cdot \phi(p_2^{k_2}) \cdots \phi(p_m^{k_m})$$
$$= n \cdot \left(1 - \frac{1}{p_1}\right) \cdot \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_m}\right)$$

- Application in Euler's theorem:
 - If gcd(a, M) = 1, then:

$$a^{\phi(M)} \equiv 1 \pmod{M} \Rightarrow a^n \equiv a^{n \bmod M} \pmod{M}$$

- In general, for arbitrary a, M and n ≥ $\log_2 M$:

$$a^n \equiv a^{\phi(M) + [n \mod \phi(M)]} \pmod{M}$$

5.2 Mobius function

• For a positive integer $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\mu(n) = \begin{cases} 1, & \text{if } n = 1 \\ 0, & \text{if } \exists k_i > 1 \\ (-1)^m & \text{otherwise} \end{cases}$$

• Properties:

$$-\sum_{d|n}\mu(d)=[n=1].$$

- If *a* and *b* are **coprime**, then $\mu(ab) = \mu(a) \cdot \mu(b)$.
- Mobius inversion: let f and g be arithmetic functions:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right)g(d)$$

5.3 Primes

Approximating the number of primes up to *n*:

n	$\pi(n)$	$\frac{n}{\ln n - 1}$	
$100 (1e^2)$	25	28	
$500 (5e^2)$	95	96	
$1000 (1e^3)$	168	169	
$5000 (5e^3)$	669	665	
$10000 (1e^4)$	1229	1218	
$50000 (5e^4)$	5133	5092	
$100000 (1e^5)$	9592	9512	
$500000 (5e^5)$	41538	41246	
$1000000 (1e^6)$	78498	78030	
$5000000 (5e^6)$	348513	346622	

 $(\pi(n))$ = the number of primes less than or equal to n, $\frac{n}{\ln n - 1}$ is used to approximate $\pi(n)$).

5.4 Wilson's theorem

A positive integer *n* is a prime if and only if:

$$(n-1)! \equiv n-1 \pmod{n}$$

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5.5 Zeckendorf's theorem

The Zeckendorf's theorem states that every positive integer n can be represented uniquely as a sum of one or more distinct non-consecutive Fibonacci numbers. For example: 64 = 55 + 8 + 1

```
85 = 55 + 21 + 8 + 1

vector<int> zeckendoft_theorem(int n) {
 vector<int> fibs = {1, 1};
 int sz = 2;
```

```
while (fibs.back() <= n) {</pre>
           fibs.push_back(fibs[sz - 1] + fibs[s - 2]);
           SZ++:
       fibs.pop_back();
       vector<int> nums;
       int p = sz - 1;
       while (n > 0) {
11
           if (n >= fibs[p]) {
12
               nums.push_back(fibs[p]);
               n -= fibs[p];
14
          }
15
          p--;
16
17
18
       return nums;
19 }
```

5.6 Bitwise operation

- $a + b = (a \oplus b) + 2(a \& b)$ • $a \mid b = (a \oplus b) + (a \& b)$ • $a \& (b \oplus c) = (a \& b) \oplus (a \& c)$ • $a \mid (b \& c) = (a \mid b) \& (a \mid c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$
- Iterating over all subsets of a set and iterating over all submasks of a mask:

```
// Time complexity: 0(3^n).
      Pollard's rho algorithm
const int PRIME_MAX = (int) 4e4; // for handle numbers <= 1e9.</pre>
const int LIMIT = (int) 1e9;
3 vector<int> primes;
5 void linear_sieve(int n);
6 num_type mulmod(num_type a, num_type b, num_type mod);
 num_type powmod(num_type a, num_type n, num_type mod);
9 bool miller_rabin(num_type a, num_type d, int s, num_type mod) {
      // \mod - 1 = a (d * 2^s).
      num_type x = powmod(a, d, mod);
      if (x == 1 || x == mod - 1) return true;
      for (int i = 1; i \le s - 1; ++i) {
          x = mulmod(x, x, mod);
          if (x == mod - 1) return true;
      return false:
19 bool is_prime(num_type n, int ITERATION = 10) {
      if (n < 4) return (n == 2 || n == 3);
      if (n % 2 == 0 || n % 3 == 0) return false;
      num\_type d = n - 1;
      int s = 0;
      while (d % 2 == 0) {
          d /= 2;
          s++;
      for (int i = 0; i < ITERATION; ++i) {
          num_type a = (num_type) (rand() % (n - 2)) + 2;
          if (miller_rabin(a, d, s, n) == false) {
               return false:
          }
      }
      return true;
36 num_type f(\text{num_type } x, \text{ int } c, \text{ num_type mod}) \{ // f(x) = (x^2 + c) \% \text{ mod.} \}
      x = mulmod(x, x, mod);
      x += c;
      if (x >= mod) x -= mod;
      return x;
42 num_type pollard_rho(num_type n, int c) {
      // algorithm to find a random divisor of 'n'.
      // using random function: f(x) = (x^2 + c) \% n.
      num_type x = 2, y = x, d;
      long long p = 1;
      int dist = 0;
```

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```
while (true) {
48
          y = f(y, c, n);
49
50
          dist++;
          d = \_gcd(llabs(x - y), n);
          if (d > 1) break;
53
          if (dist == p) { dist = 0; p *= 2; x = y; }
54
      return d;
55
57 void factorize(int n, vector<num_type> &factors);
  void llfactorize(num_type n, vector<num_type> &factors) {
       if (n < 2) return:
      if (n < LIMIT) {</pre>
           factorize(n, factors);
61
62
           return:
63
      if (is_prime(n)) {
           factors.emplace_back(n);
65
           return:
66
67
       num_type d = n;
68
      for (int c = 2; d == n; c++) {
           d = pollard_rho(n, c);
70
71
72
      llfactorize(d, factors);
      llfactorize(n / d, factors);
73
74 }
75 vector<num_type> gen_divisors(vector<pair<num_type, int>> &factors) {
       vector<num_type> divisors = {1};
76
       for (auto &x : factors) {
          int sz = (int) divisors.size();
78
79
           for (int i = 0; i < sz; ++i) {
               num_type cur = divisors[i];
80
81
               for (int j = 0; j < x.second; ++j) {
                   cur *= x.first;
                   divisors.push_back(cur);
               }
84
85
          }
       return divisors; // this array is NOT sorted yet.
88 }
```

5.8 Combinatorics

5.8.1 Catalan numbers

$$C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{n!(n+1)!}$$

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}, \ C_0 = 1, \ C_n = \frac{4n-2}{n+1} C_{n-1}$$

• The first 12 Catalan numbers (n = 0, 1, 2, ..., 12):

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$$

- Applications of Catalan numbers:
 - difference binary search trees with *n* vertices from 1 to *n*.
 - rooted binary trees with n + 1 leaves (vertices are not numbered).
 - correct bracket sequence of length 2 * n.
 - permutation [n] with no 3-term increasing subsequence (i.e. doesn't exist i < j < k for which a[i] < a[j] < a[k]).
 - ways a convex polygon of n + 2 sides can split into triangles by connecting vertices.

5.8.2 Stirling numbers of the second kind

Partitions of *n* distinct elements into exactly *k* non-empty groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

5.8.3 Derangements

Permutation of the elements of a set, such that no element appears in its original position (no fixied point). Recursive formulas:

$$D(n) = (n-1)[D(n-1) + D(n-2)] = nD(n-1) + (-1)^n$$

6 Linear algebra

6.1 Gauss elimination

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```
continue:
           for (int i=col; i<=m; ++i)</pre>
14
                swap (a[sel][i], a[row][i]);
           where[col] = row;
           for (int i=0; i<n; ++i)</pre>
18
                if (i != row) {
19
                    double c = a[i][col] / a[row][col];
20
                    for (int j=col; j<=m; ++j)
                         a[i][j] -= a[row][j] * c;
22
               }
23
           ++row:
24
       }
25
26
       ans.assign (m, 0);
27
       for (int i=0; i<m; ++i)</pre>
28
           if (where[i] != -1)
29
                ans[i] = a[where[i]][m] / a[where[i]][i];
       for (int i=0; i<n; ++i) {</pre>
30
31
           double sum = 0:
           for (int j=0; j<m; ++j)
32
                sum += ans[j] * a[i][j];
33
           if (abs (sum - a[i][m]) > EPS)
34
                return 0;
35
36
       for (int i=0; i<m; ++i)</pre>
37
           if (where[i] == -1)
38
               return INF;
39
       return 1;
40
41 }
```

7 Geometry

7.1 Fundamentals

7.1.1 Point

```
const double PI = acos(-1);
const double EPS = 1e-9;
3 typedef double ftype;
4 struct point {
      ftype x, y;
      point(ftype _x = 0, ftype _y = 0): x(_x), y(_y) {}
      point& operator+=(const point& other) {
          x += other.x; y += other.y; return *this;
10
      point& operator -= (const point& other) {
          x -= other.x; y -= other.y; return *this;
11
12
      point& operator*=(ftype t) {
13
14
          x *= t; y *= t; return *this;
15
      point& operator/=(ftype t) {
16
          x /= t; y /= t; return *this;
17
```

```
point operator+(const point& other) const {
          return point(*this) += other;
      point operator-(const point& other) const {
          return point(*this) -= other;
      point operator*(ftype t) const {
          return point(*this) *= t;
      point operator/(ftype t) const {
          return point(*this) /= t;
      point rotate(double angle) const {
          return point(x * cos(angle) - y * sin(angle), x * sin(angle) + y *
       cos(angle));
      friend istream& operator>>(istream &in, point &t);
      friend ostream& operator<<(ostream &out, const point& t);</pre>
      bool operator<(const point& other) const {</pre>
          if (fabs(x - other.x) < EPS)</pre>
              return y < other.y;</pre>
          return x < other.x;</pre>
      }
41 };
43 istream& operator>>(istream &in, point &t) {
      in >> t.x >> t.y;
      return in;
46 }
47 ostream& operator<<(ostream &out, const point& t) {
      out << t.x << ' ' << t.y;
      return out;
50 }
symple ftype dot(point a, point b) {return a.x * b.x + a.y * b.y;}
ftype norm(point a) {return dot(a, a);}
ftype abs(point a) {return sqrt(norm(a));}
ftype angle(point a, point b) {return acos(dot(a, b) / (abs(a) * abs(b)));}
ftype proj(point a, point b) {return dot(a, b) / abs(b);}
ftype cross(point a, point b) {return a.x * b.y - a.y * b.x;}
58 bool ccw(point a, point b, point c) {return cross(b - a, c - a) > EPS;}
bool collinear(point a, point b, point c) {return fabs(cross(b - a, c - a))
60 point intersect(point a1, point d1, point a2, point d2) {
       double t = cross(a2 - a1, d2) / cross(d1, d2);
       return a1 + d1 * t;
63 }
 7.1.2 Line
1 struct line {
       double a, b, c;
```

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```
line (double _a = 0, double _b = 0, double _c = 0): a(_a), b(_b), c(_c)
       {}
      friend ostream & operator<<(ostream& out, const line& 1);</pre>
5 };
6 ostream & operator << (ostream& out, const line& 1) {
      out << 1.a << ' ' << 1.b << ' ' << 1.c;
      return out;
9 }
void pointsToLine(const point& p1, const point& p2, line& 1) {
      if (fabs(p1.x - p2.x) < EPS)
          1 = \{1.0, 0.0, -p1.x\};
12
13
14
          1.a = - (double)(p1.y - p2.y) / (p1.x - p2.x);
15
          1.b = 1.0;
16
          1.c = -1.a * p1.x - 1.b * p1.y;
17
      }
18 }
void pointsSlopeToLine(const point& p, double m, line& 1) {
      1.a = -m:
      1.b = 1:
21
      1.c = -1.a * p.x - 1.b * p.y;
22
23 }
24 bool areParallel(const line& 11, const line& 12) {
      return fabs(l1.a - l2.a) < EPS && fabs(l1.b - l2.b) < EPS;</pre>
25
26 }
27 bool areSame(const line& 11, const line& 12) {
      return areParallel(11, 12) && fabs(11.c - 12.c) < EPS;</pre>
29 }
30 bool areIntersect(line l1, line l2, point& p) {
      if (areParallel(l1, l2)) return false;
      p.x = -(11.c * 12.b - 11.b * 12.c) / (11.a * 12.b - 11.b * 12.a);
33
      if (fabs(11.b) > EPS) p.y = -(11.c + 11.a * p.x);
      else p.y = -(12.c + 12.a * p.x);
34
35
      return 1;
36 }
37 double distToLine(point p, point a, point b, point& c) {
      double t = dot(p - a, b - a) / norm(b - a);
      c = a + (b - a) * t;
      return abs(c - p);
40
41 }
42 double distToSegment(point p, point a, point b, point& c) {
      double t = dot(p - a, b - a) / norm(b - a);
43
      if (t > 1.0)
44
          c = point(b.x, b.y);
45
      else if (t < 0.0)
46
47
          c = point(a.x, a.y);
      else
48
          c = a + (b - a) * t;
49
50
      return abs(c - p);
51 }
52 bool intersectTwoSegment(point a, point b, point c, point d) {
      ftype ABxAC = cross(b - a, c - a);
```

```
ftype ABxAD = cross(b - a, d - a);
      ftype CDxCA = cross(d - c, a - c);
      ftype CDxCB = cross(d - c, b - c);
      if (ABxAC == 0 | | ABxAD == 0 | | CDxCA == 0 | | CDxCB == 0) {
          if (ABxAC == 0 &\& dot(a - c, b - c) <= 0) return true;
          if (ABxAD == 0 && dot(a - d, b - d) <= 0) return true;
          if (CDxCA == 0 &\& dot(c - a, d - a) <= 0) return true;
          if (CDxCB == 0 &\& dot(c - b, d - b) <= 0) return true;
          return false;
      }
      return (ABxAC * ABxAD < 0 && CDxCA * CDxCB < 0);</pre>
66 void perpendicular(line l1, point p, line& l2) {
      if (fabs(l1.a) < EPS)</pre>
          12 = \{1.0, 0.0, -p.x\};
       else {
          12.a = -11.b / 11.a;
          12.b = 1.0;
          12.c = -12.a * p.x - 12.b * p.y;
      }
74 }
 7.1.3 Circle
int insideCircle(const point& p, const point& center, ftype r) {
      ftype d = norm(p - center);
      ftype rSq = r * r;
      return fabs(d - rSq) < EPS ? 0 : (d - rSq >= EPS ? 1 : -1);
5 }
6 bool circle2PointsR(const point& p1, const point& p2, ftype r, point& c) {
      double h = r * r - norm(p1 - p2) / 4.0;
      if (fabs(h) < 0) return false;</pre>
      h = sqrt(h);
      point perp = (p2 - p1).rotate(PI / 2.0);
      point m = (p1 + p2) / 2.0;
      c = m + perp * (h / abs(perp));
      return true:
14 }
 7.1.4 Triangle
 double areaTriangle(double ab, double bc, double ca) {
       double p = (ab + bc + ca) / 2;
      return sqrt(p) * sqrt(p - ab) * sqrt(p - bc) * sqrt(p - ca);
4 }
5 double rInCircle(double ab, double bc, double ca) {
       double p = (ab + bc + ca) / 2;
      return areaTriangle(ab, bc, ca) / p;
8 }
9 double rInCircle(point a, point b, point c) {
      return rInCircle(abs(a - b), abs(b - c), abs(c - a));
11 }
12 bool inCircle(point p1, point p2, point p3, point &ctr, double &r) {
      r = rInCircle(p1, p2, p3);
```

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```
if (fabs(r) < EPS) return false;</pre>
14
15
      line 11, 12;
      double ratio = abs(p2 - p1) / abs(p3 - p1);
16
      point p = p2 + (p3 - p2) * (ratio / (1 + ratio));
      pointsToLine(p1, p, l1);
18
      ratio = abs(p1 - p2) / abs(p2 - p3);
19
      p = p1 + (p3 - p1) * (ratio / (1 + ratio));
20
      pointsToLine(p2, p, 12);
21
      areIntersect(l1, l2, ctr);
22
23
      return true:
24 }
25 double rCircumCircle(double ab, double bc, double ca) {
      return ab * bc * ca / (4.0 * areaTriangle(ab, bc, ca));
27 }
28 double rCircumCircle(point a, point b, point c) {
29
      return rCircumCircle(abs(b - a), abs(c - b), abs(a - c));
30 }
7.1.5 Convex hull
vector<point> CH_Andrew(vector<point> &Pts) { // overall 0(n log n)
      int n = Pts.size(), k = 0;
      vector<point> H(2 * n);
      sort(Pts.begin(), Pts.end());
      for (int i = 0; i < n; ++i) {
          while ((k \ge 2) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
          H[k++] = Pts[i];
      for (int i = n - 2, t = k + 1; i >= 0; --i) {
          while ((k \ge t) \& (H[k - 2], H[k - 1], Pts[i])) --k;
10
          H[k++] = Pts[i];
11
12
      H.resize(k);
13
      return H;
14
15 }
 7.1.6 Polygon
double perimeter(const vector<point> &P) {
      double ans = 0.0;
      for (int i = 0; i < (int)P.size() - 1; ++i)
          ans += abs(P[i] - P[i + 1]);
      return ans;
6 }
7 double area(const vector<point> &P) {
      double ans = 0.0;
      for (int i = 0; i < (int)P.size() - 1; ++i)</pre>
          ans += (P[i].x * P[i + 1].y - P[i + 1].x * P[i].y);
      return fabs(ans) / 2.0;
11
12 }
bool isConvex(const vector<point> &P) {
      int n = (int)P.size():
      if (n <= 3) return false;</pre>
      bool firstTurn = ccw(P[0], P[1], P[2]);
```

```
for (int i = 1; i < n - 1; ++i)
          if (ccw(P[i], P[i + 1], P[(i + 2) == n ? 1 : i + 2]) != firstTurn)
              return false;
      return true;
21 }
int insidePolygon(point pt, const vector<point> &P) {
      int n = (int)P.size();
      if (n <= 3) return -1;
      bool on_polygon = false;
      for (int i = 0; i < n - 1; ++i)
          if (fabs(abs(P[i] - pt) + abs(pt - P[i + 1]) - abs(P[i] - P[i + 1])
              on_polygon = true;
      if (on_polygon) return 0;
      double sum = 0.0;
      for (int i = 0; i < n - 1; ++i) {
          if (ccw(pt, P[i], P[i + 1]))
              sum += angle(P[i] - pt, P[i + 1] - pt);
          else
              sum -= angle(P[i] - pt, P[i + 1] - pt);
      return fabs(sum) > PI ? 1 : -1;
38 }
      Minimum enclosing circle
2 * Description: computes the minimum circle that encloses all the given
      points.
4 double abs(point a) { return sqrt(a.X * a.X + a.Y * a.Y); }
6 point center_from(double bx, double by, double cx, double cy) {
      double B = bx * bx + by * by, C = cx * cx + cy * cy, D = bx * cy - by *
      return point((cy * B - by * C) / (2 * D), (bx * C - cx * B) / (2 * D));
9 }
11 circle circle_from(point A, point B, point C) {
      point I = center_from(B.X - A.X, B.Y - A.Y, C.X - A.X, C.Y - A.Y);
      return circle(I + A, abs(I));
14 }
16 const int N = 100005;
17 int n, x[N], y[N];
18 point a[N];
20 circle emo_welzl(int n, vector<point> T) {
      if (T.size() == 3 || n == 0) {
          if (T.size() == 0) return circle(point(0, 0), -1);
          if (T.size() == 1) return circle(T[0], 0);
          if (T.size() == 2) return circle((T[0] + T[1]) / 2, abs(T[0] - T[1])
24
```

[1]) / 2);

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```
return circle_from(T[0], T[1], T[2]);
25
26
       random\_shuffle(a + 1, a + n + 1);
27
       circle Result = emo_welzl(0, T);
       for (int i = 1; i <= n; i++)</pre>
29
           if (abs(Result.X - a[i]) > Result.Y + 1e-9) {
30
               T.push_back(a[i]);
31
               Result = emo_welzl(i - 1, T);
32
33
               T.pop_back();
          }
34
       return Result;
35
36 }
```

8 Graph

8.1 Strongly connected components

```
* Description: Tarjan's algorithm finds strongly connected components
3 * in a directed graph. If vertices u and v belong to the same component,
  * then scc_id[u] == scc_id[v].
5 * Tested: https://judge.yosupo.jp/problem/scc
7 const int N = (int) 5e5;
8 vector<int> g[N], st;
9 int low[N], num[N], dfs_timer, scc_id[N], scc;
10 bool used[N];
void Tarjan(int u) {
      low[u] = num[u] = ++dfs_timer;
      st.push_back(u);
13
      for (int v : g[u]) {
14
          if (used[v]) continue;
15
          if (num[v] == 0) {
16
              Tarjan(v);
              low[u] = min(low[u], low[v]);
18
          }
19
          else {
20
              low[u] = min(low[u], num[v]);
21
          }
22
23
24
      if (low[u] == num[u]) {
          int v;
25
26
          do {
              v = st.back(); st.pop_back();
27
              debug(u, v)
28
              used[v] = true;
29
              scc_id[v] = scc;
          } while (v != u);
          scc++;
32
33
34 }
```

8.2 Topo sort

```
1 /**
* Description: A topological sort of a directed acyclic graph
3 * is a linear ordering of its vertices such that for every directed edge
* from vertex u to vertex v, u comes before v in the ordering.
5 * Note: If there are cycles, the returned list will have size smaller than
       n (i.e. topo.size() < n).
  * Tested: https://judge.yosupo.jp/problem/scc
8 vector<int> topo_sort(const vector<vector<int>> &g) {
      int n = (int) g.size();
      vector<int> indeg(n);
      for (int u = 0; u < n; ++u) {
          for (int v : g[u]) indeg[v]++;
      queue < int > q; // Note: use min-heap to get the smallest lexicographical
       order.
      for (int u = 0; u < n; ++u) {
          if (indeg[u] == 0) q.emplace(u);
      }
      vector<int> topo;
      while (!q.empty()) {
          int u = q.front(); q.pop();
          topo.emplace_back(u);
          for (int v : g[u]) {
              if (--indeg[v] == 0) q.emplace(v);
          }
      }
      return topo;
      K-th smallest shortest path
1 /** Finding the k-th smallest shortest path from vertex s to vertex t,
      each vertex can be visited more than once.
3 */
4 using adj_list = vector<vector<pair<int, int>>>;
5 vector<int> k_smallest(const adj_list &g, int k, int s, int t) {
      int n = (int) g.size();
      vector<long long> ans;
      vector<int> cnt(n);
      using pli = pair<long long, int>;
      priority_queue<pli, vector<pli>, greater<pli>> pq;
      pg.emplace(0, s);
      while (!pq.empty() && cnt[t] < k) {</pre>
          int u = pq.top().second;
          long long d = pq.top().first;
          pq.pop();
          if (cnt[u] == k) continue;
          cnt[u]++;
          if (u == t) {
```

ans.push_back(d);

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```
for (auto [v, cost] : g[u]) {
21
              pq.emplace(d + cost, v);
          }
24
25
      assert(ans.size() == k);
      return ans;
26
27 }
      Eulerian path
 8.4.1 Directed graph
1 /**
* Hierholzer's algorithm.
3 * Description: An Eulerian path in a directed graph is a path that visits
      all edges exactly once.
  * An Eulerian cycle is a Eulerian path that is a cycle.
   * Time complexity: O(|E|).
   vector<int> find_path_directed(const vector<vector<int>> &g, int s) {
      int n = (int) q.size();
      vector<int> stack, cur_edge(n), vertices;
      stack.push_back(s);
      while (!stack.empty()) {
11
          int u = stack.back();
12
          stack.pop_back();
13
          while (cur_edge[u] < (int) g[u].size()) {</pre>
14
              stack.push_back(u);
15
              u = g[u][cur\_edge[u]++];
16
          }
          vertices.push_back(u);
18
19
      reverse(vertices.begin(), vertices.end());
20
21
      return vertices;
22 }
 8.4.2 Undirected graph
1 /**
* Hierholzer's algorithm.
3 * Description: An Eulerian path in a undirected graph is a path that
      visits all edges exactly once.
* An Eulerian cycle is a Eulerian path that is a cycle.
* Time complexity: O(|E|).
6 */
7 struct Edge {
      list<Edge>::iterator reverse_edge;
      Edge(int _to) : to(_to) {}
11 };
vector<int> vertices;
void find_path(vector<list<Edge>> &g, int u) {
      while (!g[u].empty()) {
```

```
int v = g[u].front().to;
          g[v].erase(g[u].front().reverse_edge);
          g[u].pop_front();
          find_path(g, v);
      vertices.emplace_back(u); // reversion list.
void add_edge(int u, int v) {
      g[u].emplace_front(v);
      g[v].emplace_front(u);
      g[u].front().reverse_edge = g[v].begin();
      g[v].front().reverse_edge = g[u].begin();
27 }
     Misc.
9.1 Ternary search
const double eps = 1e-9;
2 double ternary_search_max(double 1, double r) {
      // find x0 such that: f(x0) > f(x), \all x: 1 <= x <= r.
      while (r - l > eps) {
          double mid1 = 1 + (r - 1) / 3;
          double mid2 = r - (r - 1) / 3;
          if (f(mid1) < f(mid2)) l = mid1;
          else r = mid2;
      }
      return 1;
12 double ternary_search_min(double 1, double r) {
      // find x0 such that: f(x0) < f(x), \all x: l \ll x \ll r.
      while (r - 1 > eps) {
          double mid1 = 1 + (r - 1) / 3;
          double mid2 = r - (r - 1) / 3;
          if (f(mid1) > f(mid2)) 1 = mid1;
          else r = mid2;
      }
      return 1;
21 }
      Dutch flag national problem
void dutch_flag_national(vector<int> &arr) {
      // All elements that are LESS than pivot are moved to the LEFT.
      // All elements that are GREATER than pivot are moved to the RIGHT.
      // E.g. [1, 2, 0, 0, 2, 2, 1], pivot = 1 -> [0, 0, 1, 1, 2, 2, 2].
      int n = (int) arr.size();
      int i = 0, j = 0, k = n - 1;
      int pivot = 1;
      // 0....i....j....k....n
      while (j <= k) {
          if (arr[j] < pivot) {</pre>
              swap(arr[i], arr[j]);
```

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```
36 };
              i++;
12
13
              j++;
14
          }
          else if (arr[j] > pivot) {
               swap(arr[j], arr[k]);
16
17
              k--;
          }
18
          else {
19
20
              j++;
          }
21
22
23
      // 0 <= index <= i - 1: arr[index] < mid.
      // i <= index <= k: arr[index] = mid.</pre>
24
      // k + 1 <= index < sz: arr[index] > mid.
25
26 }
 9.3
      Matrix
1 struct Matrix {
      static const matrix_type INF = numeric_limits<matrix_type>::max();
      int N, M;
      vector<vector<matrix_type>> mat;
      Matrix(int _N, int _M, matrix_type v = 0) : N(_N), M(_M) {
           mat.assign(N, vector<matrix_type>(M, v));
      static Matrix identity(int n) { // return identity matrix.
           Matrix I(n, n);
          for (int i = 0; i < n; ++i) {</pre>
11
12
              I[i][i] = 1;
13
14
          return I;
      }
15
16
      vector<matrix_type>& operator[](int r) { return mat[r]; }
17
18
      const vector<matrix_type>& operator[](int r) const { return mat[r]; }
19
      Matrix& operator*=(const Matrix &other) {
20
           assert(M == other.N); // [N x M] [other.N x other.M]
21
          Matrix res(N, other.M);
22
          for (int r = 0; r < N; ++r) {
23
               for (int c = 0; c < other.M; ++c) {</pre>
24
                   long long square_mod = (long long) MOD * MOD;
25
                   long long sum = 0;
                   for (int g = 0; g < M; ++g) {
                       sum += (long long) mat[r][g] * other[g][c];
                       if (sum >= square_mod) sum -= square_mod;
29
                   res[r][c] = sum % MOD;
              }
          }
33
          mat.swap(res.mat); return *this;
34
35
```