

ACM - ICPC 2022

TEAM NOTEBOOK

Can Tho University

Contents

1 Mathematics	2	5 Linear algebra	10
1.1 Trigonometry	2	5.1 Gauss elimination	10
1.2 Sums	2		
2 Data structures	2	6 Geometry	11
2.1 Sparse table	2	6.1 Fundamentals	11
2.2 Ordered set	2	6.2 Minimum enclosing circle	13
2.3 Dsu	3		
2.4 Segment tree	3	7 Graph	14
2.5 Efficient segment tree	4	7.1 K-th smallest shortest path	14
2.6 Persistent lazy segment tree	4	7.2 Eulerian path	14
2.7 Fenwick tree	5		
3 String	6	8 Misc.	14
3.1 Prefix function	6	8.1 Ternary search	14
3.2 Counting occurrences of each prefix	6	8.2 Dutch flag national problem	15
3.3 Knuth–Morris–Pratt algorithm	6	8.3 Matrix	15
3.4 Manacher’s algorithm	6	8.4 Debugging	15
3.5 Trie	6		
3.6 Hashing	7		
4 Number Theory	7		
4.1 Euler’s totient function	7		
4.2 Mobius function	8		
4.3 Primes	8		
4.4 Wilson’s theorem	8		
4.5 Zeckendorf’s theorem	8		
4.6 Bitwise operation	9		
4.7 Combinatorics	9		
4.8 Pollard’s rho algorithm	9		

1 Mathematics

1.1 Trigonometry

1.1.1 Sum - difference identities

$$\sin(u \pm v) = \sin(u) \cos(v) \pm \cos(u) \sin(v)$$

$$\cos(u \pm v) = \cos(u) \cos(v) \mp \sin(u) \sin(v)$$

$$\tan(u \pm v) = \frac{\tan(u) \pm \tan(v)}{1 \mp \tan(u) \tan(v)}$$

1.1.2 Sum to product identities

$$\cos(u) + \cos(v) = 2 \cos\left(\frac{u+v}{2}\right) \cos\left(\frac{u-v}{2}\right)$$

$$\cos(u) - \cos(v) = -2 \sin\left(\frac{u+v}{2}\right) \sin\left(\frac{u-v}{2}\right)$$

$$\sin(u) + \sin(v) = 2 \sin\left(\frac{u+v}{2}\right) \cos\left(\frac{u-v}{2}\right)$$

$$\sin(u) - \sin(v) = 2 \cos\left(\frac{u+v}{2}\right) \sin\left(\frac{u-v}{2}\right)$$

1.1.3 Product identities

$$\cos(u) \cos(v) = \frac{1}{2} [\cos(u+v) + \cos(u-v)]$$

$$\sin(u) \sin(v) = -\frac{1}{2} [\cos(u+v) - \cos(u-v)]$$

$$\sin(u) \cos(v) = \frac{1}{2} [\sin(u+v) + \sin(u-v)]$$

1.1.4 Double - triple angle identities

$$\sin(2u) = 2 \sin(u) \cos(u)$$

$$\cos(2u) = 2 \cos^2(u) - 1 = 1 - 2 \sin^2(u)$$

$$\tan(2u) = \frac{2 \tan(u)}{1 - \tan^2(u)}$$

$$\sin(3u) = 3 \sin(u) - 4 \sin^3(u)$$

$$\cos(3u) = 4 \cos^3(u) - 3 \cos(u)$$

$$\tan(3u) = \frac{3 \tan(u) - \tan^3(u)}{1 - 3 \tan^2(u)}$$

1.2 Sums

$$n^a + n^{a+1} + \dots + n^b = \frac{n^{b+1} - n^a}{n - 1}, \quad n \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$$

$$1^3 + 2^3 + 3^3 + \dots + n^3 = \left(\frac{n(n+1)}{2}\right)^2$$

$$1^4 + 2^4 + 3^4 + \dots + n^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

2 Data structures

2.1 Sparse table

```

1 int st[MAXN][K + 1];
2 for (int i = 0; i < N; i++) {
3     st[i][0] = f(array[i]);
4 }
5 for (int j = 1; j <= K; j++) {
6     for (int i = 0; i + (1 << j) <= N; i++) {
7         st[i][j] = f(st[i][j - 1], st[i + (1 << (j - 1))][j - 1]);
8     }
9 }
10 // Range Minimum Queries.
11 int lg[MAXN + 1];
12 lg[1] = 0;
13 for (int i = 2; i <= MAXN; i++) {
14     lg[i] = lg[i / 2] + 1;
15 }
16 int j = lg[R - L + 1];
17 int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);
18 // Range Sum Queries.
19 long long sum = 0;
20 for (int j = K; j >= 0; j--) {
21     if ((1 << j) <= R - L + 1) {
22         sum += st[L][j];
23         L += 1 << j;
24     }
25 }

```

2.2 Ordered set

```

1 #include <ext/pb_ds/assoc_container.hpp>
2 #include <ext/pb_ds/tree_policy.hpp>
3 using namespace __gnu_pbds;
4

```

```

5 template<typename key_type>
6 using set_t = tree<key_type, null_type, less<key_type>, rb_tree_tag,
7     tree_order_statistics_node_update>;
8
9 void example() {
10     vector<int> nums = {1, 2, 3, 5, 10};
11     set_t<int> st(nums.begin(), nums.end());
12
13     cout << *st.find_by_order(0) << '\n'; // 1
14     assert(st.find_by_order(-INF) == st.end());
15     assert(st.find_by_order(INF) == st.end());
16
17     cout << st.order_of_key(2) << '\n'; // 1
18     cout << st.order_of_key(4) << '\n'; // 3
19     cout << st.order_of_key(9) << '\n'; // 4
20     cout << st.order_of_key(-INF) << '\n'; // 0
21     cout << st.order_of_key(INF) << '\n'; // 5
22 }

```

2.3 Dsu

```

1 struct Dsu {
2     int n;
3     vector<int> par, sz;
4     Dsu(int _n) : n(_n) {
5         sz.resize(n, 1);
6         par.resize(n);
7         iota(par.begin(), par.end(), 0);
8     }
9     int find(int v) {
10         // finding leader/parent of set that contains the element v.
11         // with {path compression optimization}.
12         return (v == par[v] ? v : par[v] = find(par[v]));
13     }
14     bool same(int u, int v) {
15         return find(u) == find(v);
16     }
17     bool unite(int u, int v) {
18         u = find(u); v = find(v);
19         if (u == v) return false;
20         if (sz[u] < sz[v]) swap(u, v);
21         par[v] = u;
22         sz[u] += sz[v];
23         return true;
24     }
25     vector<vector<int>> groups() {
26         // returns the list of the "list of the vertices in a connected
27         // component".
28         vector<int> leader(n);
29         for (int i = 0; i < n; ++i) {
30             leader[i] = find(i);
31         }
32         vector<int> id(n, -1);

```

```

32     int count = 0;
33     for (int i = 0; i < n; ++i) {
34         if (id[leader[i]] == -1) {
35             id[leader[i]] = count++;
36         }
37     }
38     vector<vector<int>> result(count);
39     for (int i = 0; i < n; ++i) {
40         result[id[leader[i]]].push_back(i);
41     }
42     return result;
43 }
44 };

```

2.4 Segment tree

```

1 /**
2  * Description: A segment tree with range updates and sum queries that
3  * supports three types of operations:
4  * + Increase each value in range [l, r] by x (i.e. a[i] += x).
5  * + Set each value in range [l, r] to x (i.e. a[i] = x).
6  * + Determine the sum of values in range [l, r].
7  */
8 struct SegmentTree {
9     int n;
10    vector<long long> tree, lazy_add, lazy_set;
11    SegmentTree(int _n) : n(_n) {
12        int p = 1;
13        while (p < n) p *= 2;
14        tree.resize(p * 2);
15        lazy_add.resize(p * 2);
16        lazy_set.resize(p * 2);
17    }
18    long long merge(const long long &left, const long long &right) {
19        return left + right;
20    }
21    void build(int id, int l, int r, const vector<int> &arr) {
22        if (l == r) {
23            tree[id] += arr[l];
24            return;
25        }
26        int mid = (l + r) >> 1;
27        build(id * 2, l, mid, arr);
28        build(id * 2 + 1, mid + 1, r, arr);
29        tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
30    }
31    void push(int id, int l, int r) {
32        if (lazy_set[id] == 0 && lazy_add[id] == 0) return;
33        int mid = (l + r) >> 1;
34        for (int child : {id * 2, id * 2 + 1}) {
35            int range = (child == id * 2 ? mid - l + 1 : r - mid);
36            if (lazy_set[id] != 0) {

```

```

37         lazy_set[child] = lazy_set[id];
38         tree[child] = range * lazy_set[id];
39     }
40     lazy_add[child] += lazy_add[id];
41     tree[child] += range * lazy_add[id];
42 }
43 lazy_add[id] = lazy_set[id] = 0;
44
45 }
46 void update(int id, int l, int r, int u, int v, int amount, bool
set_value = false) {
47     if (r < u || l > v) return;
48     if (u <= l && r <= v) {
49         if (set_value) {
50             tree[id] = 1LL * amount * (r - l + 1);
51             lazy_set[id] = amount;
52             lazy_add[id] = 0; // clear all previous updates.
53         }
54         else {
55             tree[id] += 1LL * amount * (r - l + 1);
56             lazy_add[id] += amount;
57         }
58         return;
59     }
60     push(id, l, r);
61     int mid = (l + r) >> 1;
62     update(id * 2, l, mid, u, v, amount, set_value);
63     update(id * 2 + 1, mid + 1, r, u, v, amount, set_value);
64     tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
65 }
66 long long get(int id, int l, int r, int u, int v) {
67     if (r < u || l > v) return 0;
68     if (u <= l && r <= v) {
69         return tree[id];
70     }
71     push(id, l, r);
72     int mid = (l + r) >> 1;
73     long long left = get(id * 2, l, mid, u, v);
74     long long right = get(id * 2 + 1, mid + 1, r, u, v);
75     return merge(left, right);
76 }
77 };

```

2.5 Efficient segment tree

```

1 template<typename T> struct SegmentTree {
2     int n;
3     vector<T> tree;
4     SegmentTree(int _n) : n(_n), tree(2 * n) {}
5     T merge(const T &left, const T &right) {
6         return left + right;
7     }
8     template<typename G>

```

```

9     void build(const vector<G> &initial) {
10         assert((int) initial.size() == n);
11         for (int i = 0; i < n; ++i) {
12             tree[i + n] = initial[i];
13         }
14         for (int i = n - 1; i > 0; --i) {
15             tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
16         }
17     }
18     void modify(int i, int v) {
19         tree[i + n] = v;
20         for (i /= 2; i > 0; i /= 2) {
21             tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
22         }
23     }
24     T get_sum(int l, int r) {
25         // sum of elements from l to r - 1.
26         T ret{};
27         for (l += n, r += n; l < r; l /= 2, r /= 2) {
28             if (l & 1) ret = merge(ret, tree[l++]);
29             if (r & 1) ret = merge(ret, tree[--r]);
30         }
31         return ret;
32     }
33 };

```

2.6 Persistent lazy segment tree

```

1 struct Vertex {
2     int l, r;
3     long long val, lazy;
4     bool has_changed = false;
5     Vertex() {}
6     Vertex(int _l, int _r, long long _val, int _lazy = 0) : l(_l), r(_r),
val(_val), lazy(_lazy) {}
7 };
8 struct PerSegmentTree {
9     vector<Vertex> tree;
10    vector<int> root;
11    int build(const vector<int> &arr, int l, int r) {
12        if (l == r) {
13            tree.emplace_back(-1, -1, arr[l]);
14            return tree.size() - 1;
15        }
16        int mid = (l + r) / 2;
17        int left = build(arr, l, mid);
18        int right = build(arr, mid + 1, r);
19        tree.emplace_back(left, right, tree[left].val + tree[right].val);
20        return tree.size() - 1;
21    }
22    int add(int x, int l, int r, int u, int v, int amt) {
23        if (l > v || r < u) return x;
24        if (u <= l && r <= v) {

```

```

25         tree.emplace_back(tree[x].l, tree[x].r, tree[x].val + 1LL * amt
26         * (r - l + 1), tree[x].lazy + amt);
27         tree.back().has_changed = true;
28         return tree.size() - 1;
29     }
30     int mid = (l + r) >> 1;
31     push(x, l, mid, r);
32     int left = add(tree[x].l, l, mid, u, v, amt);
33     int right = add(tree[x].r, mid + 1, r, u, v, amt);
34     tree.emplace_back(left, right, tree[left].val + tree[right].val, 0)
35     ;
36     return tree.size() - 1;
37 }
38 long long get_sum(int x, int l, int r, int u, int v) {
39     if (r < u || l > v) return 0;
40     if (u <= l && r <= v) return tree[x].val;
41     int mid = (l + r) / 2;
42     push(x, l, mid, r);
43     return get_sum(tree[x].l, l, mid, u, v) + get_sum(tree[x].r, mid +
44     1, r, u, v);
45 }
46 void push(int x, int l, int mid, int r) {
47     if (!tree[x].has_changed) return;
48     Vertex left = tree[tree[x].l];
49     Vertex right = tree[tree[x].r];
50     tree.emplace_back(left);
51     tree[x].l = tree.size() - 1;
52     tree.emplace_back(right);
53     tree[x].r = tree.size() - 1;
54
55     tree[tree[x].l].val += tree[x].lazy * (mid - l + 1);
56     tree[tree[x].l].lazy += tree[x].lazy;
57
58     tree[tree[x].r].val += tree[x].lazy * (r - mid);
59     tree[tree[x].r].lazy += tree[x].lazy;
60
61     tree[tree[x].l].has_changed = true;
62     tree[tree[x].r].has_changed = true;
63     tree[x].lazy = 0;
64     tree[x].has_changed = false;
65 }
66 };

```

2.7 Fenwick tree

```

1 using tree_type = long long;
2 struct FenwickTree {
3     int n;
4     vector<tree_type> fenw_coeff, fenw;
5     FenwickTree() {}
6     FenwickTree(int _n) : n(_n) {
7         fenw_coeff.assign(n, 0); // fenwick tree with coefficient (n - i).
8         fenw.assign(n, 0); // normal fenwick tree.

```

```

9     }
10 void build(const vector<int> &A) {
11     assert((int) A.size() == n);
12     vector<int> diff(n);
13     diff[0] = A[0];
14     for (int i = 1; i < n; ++i) {
15         diff[i] = A[i] - A[i - 1];
16     }
17     fenw_coeff[0] = (long long) diff[0] * n;
18     fenw[0] = diff[0];
19     for (int i = 1; i < n; ++i) {
20         fenw_coeff[i] = fenw_coeff[i - 1] + (long long) diff[i] * (n -
21         i);
22         fenw[i] = fenw[i - 1] + diff[i];
23     }
24     for (int i = n - 1; i >= 0; --i) {
25         int j = (i & (i + 1)) - 1;
26         if (j >= 0) {
27             fenw_coeff[i] -= fenw_coeff[j];
28             fenw[i] -= fenw[j];
29         }
30     }
31 void add(vector<tree_type> &fenw, int i, tree_type val) {
32     while (i < n) {
33         fenw[i] += val;
34         i |= (i + 1);
35     }
36 }
37 tree_type __prefix_sum(vector<tree_type> &fenw, int i) {
38     tree_type res{};
39     while (i >= 0) {
40         res += fenw[i];
41         i = (i & (i + 1)) - 1;
42     }
43     return res;
44 }
45 tree_type prefix_sum(int i) {
46     return __prefix_sum(fenw_coeff, i) - __prefix_sum(fenw, i) * (n - i
47     - 1);
48 }
49 void range_add(int l, int r, tree_type val) {
50     add(fenw_coeff, l, (n - 1) * val);
51     add(fenw_coeff, r + 1, (n - r - 1) * (-val));
52     add(fenw, l, val);
53     add(fenw, r + 1, -val);
54 }
55 tree_type range_sum(int l, int r) {
56     return prefix_sum(r) - prefix_sum(l - 1);
57 };

```

3 String

3.1 Prefix function

```

1 /**
2  * Description: The prefix function of a string 's' is defined as an array
3  * pi of length n,
4  * where pi[i] is the length of the longest proper prefix of the substring
5  * s[0..i] which is also a suffix of this substring.
6  * Time complexity: O(|S|).
7  */
8 vector<int> prefix_function(const string &s) {
9     int n = (int) s.length();
10    vector<int> pi(n);
11    pi[0] = 0;
12    for (int i = 1; i < n; ++i) {
13        int j = pi[i - 1]; // try length pi[i - 1] + 1.
14        while (j > 0 && s[j] != s[i]) {
15            j = pi[j - 1];
16        }
17        if (s[j] == s[i]) {
18            pi[i] = j + 1;
19        }
20    }
21    return pi;
22 }
```

3.2 Counting occurrences of each prefix

```

1 vector<int> count_occurrences(const string &s) {
2     vector<int> pi = prefix_function(s);
3     int n = (int) s.size();
4     vector<int> ans(n + 1);
5     for (int i = 0; i < n; ++i) {
6         ans[pi[i]]++;
7     }
8     for (int i = n - 1; i > 0; --i) {
9         ans[pi[i - 1]] += ans[i];
10    }
11    for (int i = 0; i <= n; ++i) {
12        ans[i]++;
13    }
14    return ans;
15    // Input: ABACABA
16    // Output: 4 2 2 1 1 1 1
17 }
```

3.3 Knuth–Morris–Pratt algorithm

```

1 /**
2  * Searching for a substring in a string.
3  * Time complexity: O(N + M).
4  */
5 vector<int> KMP(const string &text, const string &pattern) {
```

```

6     int n = (int) text.length();
7     int m = (int) pattern.length();
8     string s = pattern + '$' + text;
9     vector<int> pi = prefix_function(s);
10    vector<int> indices;
11    for (int i = 0; i < (int) s.length(); ++i) {
12        if (pi[i] == m) {
13            indices.push_back(i - 2 * m);
14        }
15    }
16    return indices;
17 }
```

3.4 Manacher's algorithm

```

1 /**
2  * Description: for each position, computes d[0][i] = half length of
3  * longest palindrome centered on i (rounded up), d[1][i] = half length of
4  * longest palindrome centered on i and i - 1.
5  * Time complexity: O(N).
6  * Tested: https://judge.yosupo.jp/problem/enumerate\_palindromes, stress-
7  * tested.
8  */
9 array<vector<int>, 2> manacher(const string &s) {
10    int n = (int) s.size();
11    array<vector<int>, 2> d;
12    for (int z = 0; z < 2; ++z) {
13        d[z].resize(n);
14        int l = 0, r = 0;
15        for (int i = 0; i < n; ++i) {
16            int mirror = l + r - i + z;
17            d[z][i] = (i > r ? 0 : min(d[z][mirror], r - i));
18            int L = i - d[z][i] - z, R = i + d[z][i];
19            while (L >= 0 && R < n && s[L] == s[R]) {
20                d[z][i]++; L--; R++;
21            }
22            if (R > r) {
23                l = L; r = R;
24            }
25        }
26    }
27    return d;
28 }
```

3.5 Trie

```

1 struct Trie {
2     const static int ALPHABET = 26;
3     const static char minChar = 'a';
4     struct Vertex {
5         int next[ALPHABET];
6         bool leaf;
7         Vertex() {
```

```

8         leaf = false;
9         fill(next, next + ALPHABET, -1);
10    }
11};
12vector<Vertex> trie;
13Trie() { trie.emplace_back(); }
14
15void insert(const string &s) {
16    int i = 0;
17    for (const char &ch : s) {
18        int j = ch - minChar;
19        if (trie[i].next[j] == -1) {
20            trie[i].next[j] = trie.size();
21            trie.emplace_back();
22        }
23        i = trie[i].next[j];
24    }
25    trie[i].leaf = true;
26}
27bool find(const string &s) {
28    int i = 0;
29    for (const char &ch : s) {
30        int j = ch - minChar;
31        if (trie[i].next[j] == -1) {
32            return false;
33        }
34        i = trie[i].next[j];
35    }
36    return (trie[i].leaf ? true : false);
37}
38};

```

3.6 Hashing

```

1 struct Hash61 {
2     static const uint64_t MOD = (1LL << 61) - 1;
3     static uint64_t BASE;
4     static vector<uint64_t> pw;
5     uint64_t addmod(uint64_t a, uint64_t b) const {
6         a += b;
7         if (a >= MOD) a -= MOD;
8         return a;
9     }
10    uint64_t submod(uint64_t a, uint64_t b) const {
11        a += MOD - b;
12        if (a >= MOD) a -= MOD;
13        return a;
14    }
15    uint64_t mulmod(uint64_t a, uint64_t b) const {
16        uint64_t low1 = (uint32_t) a, high1 = (a >> 32);
17        uint64_t low2 = (uint32_t) b, high2 = (b >> 32);
18
19        uint64_t low = low1 * low2;

```

```

20        uint64_t mid = low1 * high2 + low2 * high1;
21        uint64_t high = high1 * high2;
22
23        uint64_t ret = (low & MOD) + (low >> 61) + (high << 3) + (mid >>
24        29) + (mid << 35 >> 3) + 1;
25        // ret %= MOD;
26        ret = (ret >> 61) + (ret & MOD);
27        ret = (ret >> 61) + (ret & MOD);
28        return ret - 1;
29    }
30    void ensure_pw(int m) {
31        int n = (int) pw.size();
32        if (n >= m) return;
33        pw.resize(m);
34        for (int i = n; i < m; ++i) {
35            pw[i] = mulmod(pw[i - 1], BASE);
36        }
37    }
38    vector<uint64_t> pref;
39    int n;
40    template<typename T> Hash61(const T &s) { // strings or arrays.
41        n = (int) s.size();
42        ensure_pw(n);
43        pref.resize(n + 1);
44        pref[0] = 0;
45        for (int i = 0; i < n; ++i) {
46            pref[i + 1] = addmod(mulmod(pref[i], BASE), s[i]);
47        }
48    }
49    inline uint64_t operator()(const int from, const int to) const {
50        assert(0 <= from && from <= to && to < n);
51        // pref[to + 1] - pref[from] * pw[to - from + 1]
52        return submod(pref[to + 1], mulmod(pref[from], pw[to - from + 1]));
53    }
54};
55mt19937 rng((unsigned int) chrono::steady_clock::now().time_since_epoch().
56    count());
57uint64_t Hash61::BASE = (MOD >> 2) + rng() % (MOD >> 1);
58vector<uint64_t> Hash61::pw = vector<uint64_t>(1, 1);

```

4 Number Theory

4.1 Euler's totient function

- Euler's totient function, also known as **phi-function** $\phi(n)$ counts the number of integers between 1 and n inclusive, that are **coprime to** n .

- Properties:

– Divisor sum property: $\sum_{d|n} \phi(d) = n$.

- $\phi(n)$ is a **prime number** when $n = 3, 4, 6$.
- If p is a prime number, then $\phi(p) = p - 1$.
- If p is a prime number and $k \geq 1$, then $\phi(p^k) = p^k - p^{k-1}$.
- If a and b are **coprime**, then $\phi(ab) = \phi(a) \cdot \phi(b)$.
- In general, for **not coprime** a and b , with $d = \gcd(a, b)$ this equation holds: $\phi(ab) = \phi(a) \cdot \phi(b) \cdot \frac{d}{\phi(d)}$.
- With $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\begin{aligned}\phi(n) &= \phi(p_1^{k_1}) \cdot \phi(p_2^{k_2}) \cdots \phi(p_m^{k_m}) \\ &= n \cdot \left(1 - \frac{1}{p_1}\right) \cdot \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_m}\right)\end{aligned}$$

- Application in Euler's theorem:

- If $\gcd(a, M) = 1$, then:

$$a^{\phi(M)} \equiv 1 \pmod{M} \Rightarrow a^n \equiv a^{n \bmod M} \pmod{M}$$

- In general, for arbitrary a, M and $n \geq \log_2 M$:

$$a^n \equiv a^{\phi(M) + [n \bmod \phi(M)]} \pmod{M}$$

4.2 Mobius function

- For a positive integer $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\mu(n) = \begin{cases} 1, & \text{if } n = 1 \\ 0, & \text{if } \exists k_i > 1 \\ (-1)^m & \text{otherwise} \end{cases}$$

- Properties:

- $\sum_{d|n} \mu(d) = [n = 1]$.
- If a and b are **coprime**, then $\mu(ab) = \mu(a) \cdot \mu(b)$.
- Mobius inversion: let f and g be arithmetic functions:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right) g(d)$$

4.3 Primes

Approximating the number of primes up to n :

n	$\pi(n)$	$\frac{n}{\ln n - 1}$
100 ($1e^2$)	25	28
500 ($5e^2$)	95	96
1000 ($1e^3$)	168	169
5000 ($5e^3$)	669	665
10000 ($1e^4$)	1229	1218
50000 ($5e^4$)	5133	5092
100000 ($1e^5$)	9592	9512
500000 ($5e^5$)	41538	41246
1000000 ($1e^6$)	78498	78030
5000000 ($5e^6$)	348513	346622

($\pi(n)$ = the number of primes less than or equal to n , $\frac{n}{\ln n - 1}$ is used to approximate $\pi(n)$).

4.4 Wilson's theorem

A positive integer n is a prime if and only if:

$$(n - 1)! \equiv n - 1 \pmod{n}$$

4.5 Zeckendorf's theorem

The Zeckendorf's theorem states that every positive integer n can be represented uniquely as a sum of one or more distinct non-consecutive Fibonacci numbers. For example:

$$64 = 55 + 8 + 1$$

$$85 = 55 + 21 + 8 + 1$$

```

1 vector<int> zeckendoft_theorem(int n) {
2     vector<int> fibs = {1, 1};
3     int sz = 2;
4     while (fibs.back() <= n) {
5         fibs.push_back(fibs[sz - 1] + fibs[sz - 2]);
6         sz++;
7     }
8     fibs.pop_back();
9     vector<int> nums;
10    int p = sz - 1;
11    while (n > 0) {
12        if (n >= fibs[p]) {
13            nums.push_back(fibs[p]);
14            n -= fibs[p];
15        }
16    }
17 }
```



```

16     p--;
17 }
18 return nums;
19 }

```

4.6 Bitwise operation

- $a + b = (a \oplus b) + 2(a \& b)$
- $a | b = (a \oplus b) + (a \& b)$
- $a \& (b \oplus c) = (a \& b) \oplus (a \& c)$
- $a | (b \& c) = (a | b) \& (a | c)$
- $a \& (b | c) = (a \& b) | (a \& c)$
- $a | (a \& b) = a$
- $a \& (a | b) = a$
- $n = 2^k \Leftrightarrow !(n \& (n - 1)) = 1$
- $-a = \sim a + 1$
- $(4i) \oplus (4i+1) \oplus (4i+2) \oplus (4i+3) = 0$

- Iterating over all subsets of a set and iterating over all submasks of a mask:

```

1 for (int mask = 0; mask < (1 << n); ++mask) {
2     for (int i = 0; i < n; ++i) {
3         if (mask & (1 << i)) {
4             // do something...
5         }
6     }
7     // Time complexity: O(n * 2^n).
8 }
9 for (int mask = 0; mask < (1 << n); ++mask) {
10    for (int submask = mask; ; submask = (submask - 1) & mask) {
11        // do something...
12        if (submask == 0) break;
13    }
14    // Time complexity: O(3^n).
15 }

```

4.7 Combinatorics

4.7.1 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{n!(n+1)!}$$

$$C_{n+1} = \sum_{i=0}^n C_i C_{n-i}, \quad C_0 = 1, \quad C_n = \frac{4n-2}{n+1} C_{n-1}$$

- The first 12 Catalan numbers ($n = 0, 1, 2, \dots, 12$):

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$$

- Applications of Catalan numbers:

- difference binary search trees with n vertices from 1 to n .

- rooted binary trees with $n + 1$ leaves (vertices are not numbered).
- correct bracket sequence of length $2 * n$.
- permutation $[n]$ with no 3-term increasing subsequence (i.e. doesn't exist $i < j < k$ for which $a[i] < a[j] < a[k]$).
- ways a convex polygon of $n + 2$ sides can split into triangles by connecting vertices.

4.7.2 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k non-empty groups.

$$S(n, k) = S(n-1, k-1) + kS(n-1, k)$$

$$S(n, 1) = S(n, n) = 1$$

$$S(n, k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} \binom{k}{i} i^n$$

4.7.3 Derangements

Permutation of the elements of a set, such that no element appears in its original position (no fixed point). Recursive formulas:

$$D(n) = (n-1)[D(n-1) + D(n-2)] = nD(n-1) + (-1)^n$$

4.8 Pollard's rho algorithm

```

1 const int PRIME_MAX = (int) 4e4; // for handle numbers <= 1e9.
2 const int LIMIT = (int) 1e9;
3 vector<int> primes;
4
5 void linear_sieve(int n);
6 num_type mulmod(num_type a, num_type b, num_type mod);
7 num_type powmod(num_type a, num_type n, num_type mod);
8
9 bool miller_rabin(num_type a, num_type d, int s, num_type mod) {
10     // mod - 1 = a ^ (d * 2^s).
11     num_type x = powmod(a, d, mod);
12     if (x == 1 || x == mod - 1) return true;
13     for (int i = 1; i <= s - 1; ++i) {
14         x = mulmod(x, x, mod);
15         if (x == mod - 1) return true;
16     }
17     return false;
18 }
19 bool is_prime(num_type n, int ITERATION = 10) {
20     if (n < 4) return (n == 2 || n == 3);
21     if (n % 2 == 0 || n % 3 == 0) return false;
22     num_type d = n - 1;
23     int s = 0;
24     while (d % 2 == 0) {
25         d /= 2;

```

```

26     s++;
27 }
28 for (int i = 0; i < ITERATION; ++i) {
29     num_type a = (num_type) (rand() % (n - 2)) + 2;
30     if (miller_rabin(a, d, s, n) == false) {
31         return false;
32     }
33 }
34 return true;
35 }
36 num_type f(num_type x, int c, num_type mod) { // f(x) = (x^2 + c) % mod.
37     x = mulmod(x, x, mod);
38     x += c;
39     if (x >= mod) x -= mod;
40     return x;
41 }
42 num_type pollard_rho(num_type n, int c) {
43     // algorithm to find a random divisor of 'n'.
44     // using random function: f(x) = (x^2 + c) % n.
45
46     // ***** Floyd's cycle detection algorithm *****
47     // move 1 step and 2 steps.
48     // num_type x = 2, y = 2, d;
49     // while (true) {
50     //     x = f(x, c, n);
51     //     y = f(y, c, n);
52     //     y = f(y, c, n);
53     //     d = __gcd(llabs(x - y), n);
54     //     if (d > 1) break;
55     // }
56     // return d;
57
58     // ***** Brent's cycle detection algorithm *****
59     // move power of two steps.
60     num_type x = 2, y = x, d;
61     long long p = 1;
62     int dist = 0;
63     while (true) {
64         y = f(y, c, n);
65         dist++;
66         d = __gcd(llabs(x - y), n);
67         if (d > 1) break;
68         if (dist == p) { dist = 0; p *= 2; x = y; }
69     }
70     return d;
71 }
72 void factorize(int n, vector<num_type> &factors);
73 void llfactorize(num_type n, vector<num_type> &factors) {
74     if (n < 2) return;
75     if (n < LIMIT) {
76         factorize(n, factors);
77         return;

```

```

78     }
79     if (is_prime(n)) {
80         factors.emplace_back(n);
81         return;
82     }
83     num_type d = n;
84     for (int c = 2; d == n; c++) {
85         d = pollard_rho(n, c);
86     }
87     llfactorize(d, factors);
88     llfactorize(n / d, factors);
89 }
90 vector<num_type> gen_divisors(vector<pair<num_type, int>> &factors) {
91     vector<num_type> divisors = {1};
92     for (auto &x : factors) {
93         int sz = (int) divisors.size();
94         for (int i = 0; i < sz; ++i) {
95             num_type cur = divisors[i];
96             for (int j = 0; j < x.second; ++j) {
97                 cur *= x.first;
98                 divisors.push_back(cur);
99             }
100         }
101     }
102     return divisors; // this array is NOT sorted yet.
103 }

```

5 Linear algebra

5.1 Gauss elimination

```

1  const double EPS = 1e-9;
2  const int INF = 2; // it doesn't actually have to be infinity or a big
   number
3  int gauss (vector < vector<double> > a, vector<double> & ans) {
4      int n = (int) a.size();
5      int m = (int) a[0].size() - 1;
6      vector<int> where (m, -1);
7      for (int col=0, row=0; col<m && row<n; ++col) {
8          int sel = row;
9          for (int i=row; i<n; ++i)
10             if (abs (a[i][col]) > abs (a[sel][col]))
11                 sel = i;
12             if (abs (a[sel][col]) < EPS)
13                 continue;
14             for (int i=col; i<=m; ++i)
15                 swap (a[sel][i], a[row][i]);
16             where[col] = row;
17
18             for (int i=0; i<n; ++i)
19                 if (i != row) {
20                     double c = a[i][col] / a[row][col];

```

```

21         for (int j=col; j<=m; ++j)
22             a[i][j] -= a[row][j] * c;
23     }
24     ++row;
25 }
26 ans.assign (m, 0);
27 for (int i=0; i<m; ++i)
28     if (where[i] != -1)
29         ans[i] = a[where[i]][m] / a[where[i]][i];
30 for (int i=0; i<n; ++i) {
31     double sum = 0;
32     for (int j=0; j<m; ++j)
33         sum += ans[j] * a[i][j];
34     if (abs (sum - a[i][m]) > EPS)
35         return 0;
36 }
37 for (int i=0; i<m; ++i)
38     if (where[i] == -1)
39         return INF;
40 return 1;
41 }

```

6 Geometry

6.1 Fundamentals

6.1.1 Point

```

1 const double PI = acos(-1);
2 const double EPS = 1e-9;
3 typedef double ftype;
4 struct point {
5     ftype x, y;
6     point(ftype _x = 0, ftype _y = 0): x(_x), y(_y) {}
7     point& operator+=(const point& other) {
8         x += other.x; y += other.y; return *this;
9     }
10    point& operator-=(const point& other) {
11        x -= other.x; y -= other.y; return *this;
12    }
13    point& operator*=(ftype t) {
14        x *= t; y *= t; return *this;
15    }
16    point& operator/=(ftype t) {
17        x /= t; y /= t; return *this;
18    }
19    point operator+(const point& other) const {
20        return point(*this) += other;
21    }
22    point operator-(const point& other) const {
23        return point(*this) -= other;
24    }
25    point operator*(ftype t) const {

```

```

26        return point(*this) *= t;
27    }
28    point operator/(ftype t) const {
29        return point(*this) /= t;
30    }
31    point rotate(double angle) const {
32        return point(x * cos(angle) - y * sin(angle), x * sin(angle) + y *
33        cos(angle));
34    }
35    friend istream& operator>>(istream &in, point &t);
36    friend ostream& operator<<(ostream &out, const point& t);
37    bool operator<(const point& other) const {
38        if (fabs(x - other.x) < EPS)
39            return y < other.y;
40        return x < other.x;
41    };
42
43    istream& operator>>(istream &in, point &t) {
44        in >> t.x >> t.y;
45        return in;
46    }
47    ostream& operator<<(ostream &out, const point& t) {
48        out << t.x << ' ' << t.y;
49        return out;
50    }
51
52    ftype dot(point a, point b) {return a.x * b.x + a.y * b.y;}
53    ftype norm(point a) {return dot(a, a);}
54    ftype abs(point a) {return sqrt(norm(a));}
55    ftype angle(point a, point b) {return acos(dot(a, b) / (abs(a) * abs(b)));}
56    ftype proj(point a, point b) {return dot(a, b) / abs(b);}
57    ftype cross(point a, point b) {return a.x * b.y - a.y * b.x;}
58    bool ccw(point a, point b, point c) {return cross(b - a, c - a) > EPS;}
59    bool collinear(point a, point b, point c) {return fabs(cross(b - a, c - a))
60    < EPS;}
61    point intersect(point a1, point d1, point a2, point d2) {
62        double t = cross(a2 - a1, d2) / cross(d1, d2);
63        return a1 + d1 * t;
64    }

```

6.1.2 Line

```

1 struct line {
2     double a, b, c;
3     line (double _a = 0, double _b = 0, double _c = 0): a(_a), b(_b), c(_c)
4     {}
5     friend ostream & operator<<(ostream& out, const line& l);
6 };
7 ostream & operator<<(ostream& out, const line& l) {
8     out << l.a << ' ' << l.b << ' ' << l.c;
9     return out;
10 }

```

```

10 void pointsToLine(const point& p1, const point& p2, line& l) {
11     if (fabs(p1.x - p2.x) < EPS)
12         l = {1.0, 0.0, -p1.x};
13     else {
14         l.a = - (double)(p1.y - p2.y) / (p1.x - p2.x);
15         l.b = 1.0;
16         l.c = - l.a * p1.x - l.b * p1.y;
17     }
18 }
19 void pointsSlopeToLine(const point& p, double m, line& l) {
20     l.a = -m;
21     l.b = 1;
22     l.c = -l.a * p.x - l.b * p.y;
23 }
24 bool areParallel(const line& l1, const line& l2) {
25     return fabs(l1.a - l2.a) < EPS && fabs(l1.b - l2.b) < EPS;
26 }
27 bool areSame(const line& l1, const line& l2) {
28     return areParallel(l1, l2) && fabs(l1.c - l2.c) < EPS;
29 }
30 bool areIntersect(line l1, line l2, point& p) {
31     if (areParallel(l1, l2)) return false;
32     p.x = - (l1.c * l2.b - l1.b * l2.c) / (l1.a * l2.b - l1.b * l2.a);
33     if (fabs(l1.b) > EPS) p.y = - (l1.c + l1.a * p.x);
34     else p.y = - (l2.c + l2.a * p.x);
35     return 1;
36 }
37 double distToLine(point p, point a, point b, point& c) {
38     double t = dot(p - a, b - a) / norm(b - a);
39     c = a + (b - a) * t;
40     return abs(c - p);
41 }
42 double distToSegment(point p, point a, point b, point& c) {
43     double t = dot(p - a, b - a) / norm(b - a);
44     if (t > 1.0)
45         c = point(b.x, b.y);
46     else if (t < 0.0)
47         c = point(a.x, a.y);
48     else
49         c = a + (b - a) * t;
50     return abs(c - p);
51 }
52 bool intersectTwoSegment(point a, point b, point c, point d) {
53     ftype ABxAC = cross(b - a, c - a);
54     ftype ABxAD = cross(b - a, d - a);
55     ftype CDxCA = cross(d - c, a - c);
56     ftype CDxCB = cross(d - c, b - c);
57     if (ABxAC == 0 || ABxAD == 0 || CDxCA == 0 || CDxCB == 0) {
58         if (ABxAC == 0 && dot(a - c, b - c) <= 0) return true;
59         if (ABxAD == 0 && dot(a - d, b - d) <= 0) return true;
60         if (CDxCA == 0 && dot(c - a, d - a) <= 0) return true;
61         if (CDxCB == 0 && dot(c - b, d - b) <= 0) return true;

```

```

62         return false;
63     }
64     return (ABxAC * ABxAD < 0 && CDxCA * CDxCB < 0);
65 }
66 void perpendicular(line l1, point p, line& l2) {
67     if (fabs(l1.a) < EPS)
68         l2 = {1.0, 0.0, -p.x};
69     else {
70         l2.a = -l1.b / l1.a;
71         l2.b = 1.0;
72         l2.c = -l2.a * p.x - l2.b * p.y;
73     }
74 }

```

6.1.3 Circle

```

1 int insideCircle(const point& p, const point& center, ftype r) {
2     ftype d = norm(p - center);
3     ftype rSq = r * r;
4     return fabs(d - rSq) < EPS ? 0 : (d - rSq >= EPS ? 1 : -1);
5 }
6 bool circle2PointsR(const point& p1, const point& p2, ftype r, point& c) {
7     double h = r * r - norm(p1 - p2) / 4.0;
8     if (fabs(h) < 0) return false;
9     h = sqrt(h);
10    point perp = (p2 - p1).rotate(PI / 2.0);
11    point m = (p1 + p2) / 2.0;
12    c = m + perp * (h / abs(perp));
13    return true;
14 }

```

6.1.4 Triangle

```

1 double areaTriangle(double ab, double bc, double ca) {
2     double p = (ab + bc + ca) / 2;
3     return sqrt(p) * sqrt(p - ab) * sqrt(p - bc) * sqrt(p - ca);
4 }
5 double rInCircle(double ab, double bc, double ca) {
6     double p = (ab + bc + ca) / 2;
7     return areaTriangle(ab, bc, ca) / p;
8 }
9 double rInCircle(point a, point b, point c) {
10    return rInCircle(abs(a - b), abs(b - c), abs(c - a));
11 }
12 bool inCircle(point p1, point p2, point p3, point &ctr, double &r) {
13    r = rInCircle(p1, p2, p3);
14    if (fabs(r) < EPS) return false;
15    line l1, l2;
16    double ratio = abs(p2 - p1) / abs(p3 - p1);
17    point p = p2 + (p3 - p2) * (ratio / (1 + ratio));
18    pointsToLine(p1, p, l1);
19    ratio = abs(p1 - p2) / abs(p2 - p3);
20    p = p1 + (p3 - p1) * (ratio / (1 + ratio));
21    pointsToLine(p2, p, l2);

```

```

22     areIntersect(l1, l2, ctr);
23     return true;
24 }
25 double rCircumCircle(double ab, double bc, double ca) {
26     return ab * bc * ca / (4.0 * areaTriangle(ab, bc, ca));
27 }
28 double rCircumCircle(point a, point b, point c) {
29     return rCircumCircle(abs(b - a), abs(c - b), abs(a - c));
30 }

```

6.1.5 Convex hull

```

1 vector<point> CH_Andrew(vector<point> &Pts) { // overall O(n log n)
2     int n = Pts.size(), k = 0;
3     vector<point> H(2 * n);
4     sort(Pts.begin(), Pts.end());
5     for (int i = 0; i < n; ++i) {
6         while ((k >= 2) && !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
7         H[k++] = Pts[i];
8     }
9     for (int i = n - 2, t = k + 1; i >= 0; --i) {
10        while ((k >= t) && !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
11        H[k++] = Pts[i];
12    }
13    H.resize(k);
14    return H;
15 }

```

6.1.6 Polygon

```

1 double perimeter(const vector<point> &P) {
2     double ans = 0.0;
3     for (int i = 0; i < (int)P.size() - 1; ++i)
4         ans += abs(P[i] - P[i + 1]);
5     return ans;
6 }
7 double area(const vector<point> &P) {
8     double ans = 0.0;
9     for (int i = 0; i < (int)P.size() - 1; ++i)
10        ans += (P[i].x * P[i + 1].y - P[i + 1].x * P[i].y);
11    return fabs(ans) / 2.0;
12 }
13 bool isConvex(const vector<point> &P) {
14     int n = (int)P.size();
15     if (n <= 3) return false;
16     bool firstTurn = ccw(P[0], P[1], P[2]);
17     for (int i = 1; i < n - 1; ++i)
18         if (ccw(P[i], P[i + 1], P[(i + 2) == n ? 1 : i + 2]) != firstTurn)
19             return false;
20     return true;
21 }
22 int insidePolygon(point pt, const vector<point> &P) {
23     int n = (int)P.size();
24     if (n <= 3) return -1;

```

```

25     bool on_polygon = false;
26     for (int i = 0; i < n - 1; ++i)
27         if (fabs(abs(P[i] - pt) + abs(pt - P[i + 1]) - abs(P[i] - P[i + 1]))
28             < EPS)
29             on_polygon = true;
30     if (on_polygon) return 0;
31     double sum = 0.0;
32     for (int i = 0; i < n - 1; ++i) {
33         if (ccw(pt, P[i], P[i + 1]))
34             sum += angle(P[i] - pt, P[i + 1] - pt);
35         else
36             sum -= angle(P[i] - pt, P[i + 1] - pt);
37     }
38     return fabs(sum) > PI ? 1 : -1;

```

6.2 Minimum enclosing circle

```

1 /**
2  * Description: computes the minimum circle that encloses all the given
3  * points.
4  */
5 double abs(point a) { return sqrt(a.X * a.X + a.Y * a.Y); }
6 point center_from(double bx, double by, double cx, double cy) {
7     double B = bx * bx + by * by, C = cx * cx + cy * cy, D = bx * cy - by *
8     cx;
9     return point((cy * B - by * C) / (2 * D), (bx * C - cx * B) / (2 * D));
10 }
11 circle circle_from(point A, point B, point C) {
12     point I = center_from(B.X - A.X, B.Y - A.Y, C.X - A.X, C.Y - A.Y);
13     return circle(I + A, abs(I));
14 }
15
16 const int N = 100005;
17 int n, x[N], y[N];
18 point a[N];
19
20 circle emo_welzl(int n, vector<point> T) {
21     if (T.size() == 3 || n == 0) {
22         if (T.size() == 0) return circle(point(0, 0), -1);
23         if (T.size() == 1) return circle(T[0], 0);
24         if (T.size() == 2) return circle((T[0] + T[1]) / 2, abs(T[0] - T
25         [1]) / 2);
26         return circle_from(T[0], T[1], T[2]);
27     }
28     random_shuffle(a + 1, a + n + 1);
29     circle Result = emo_welzl(0, T);
30     for (int i = 1; i <= n; i++)
31         if (abs(Result.X - a[i]) > Result.Y + 1e-9) {
32             T.push_back(a[i]);
33             Result = emo_welzl(i - 1, T);

```

```

33         T.pop_back();
34     }
35     return Result;
36 }

```

7 Graph

7.1 K-th smallest shortest path

```

1  /** Finding the k-th smallest shortest path from vertex s to vertex t,
2   * each vertex can be visited more than once.
3   */
4  using adj_list = vector<vector<pair<int, int>>>;
5  vector<int> k_smallest(const adj_list &g, int k, int s, int t) {
6      int n = (int) g.size();
7      vector<long long> ans;
8      vector<int> cnt(n);
9      using pli = pair<long long, int>;
10     priority_queue<pli, vector<pli>, greater<pli>> pq;
11     pq.emplace(0, s);
12     while (!pq.empty() && cnt[t] < k) {
13         int u = pq.top().second;
14         long long d = pq.top().first;
15         pq.pop();
16         if (cnt[u] == k) continue;
17         cnt[u]++;
18         if (u == t) {
19             ans.push_back(d);
20         }
21         for (auto [v, cost] : g[u]) {
22             pq.emplace(d + cost, v);
23         }
24     }
25     assert(ans.size() == k);
26     return ans;
27 }

```

7.2 Eulerian path

7.2.1 Directed graph

```

1  /**
2   * Hierholzer's algorithm.
3   * Description: An Eulerian path in a directed graph is a path that visits
4   * all edges exactly once.
5   * An Eulerian cycle is a Eulerian path that is a cycle.
6   * Time complexity: O(|E|).
7   */
8  vector<int> find_path_directed(const vector<vector<int>> &g, int s) {
9      int n = (int) g.size();
10     vector<int> stack, cur_edge(n), vertices;
11     stack.push_back(s);
12     while (!stack.empty()) {
13         int u = stack.back();

```

```

13         stack.pop_back();
14         while (cur_edge[u] < (int) g[u].size()) {
15             stack.push_back(u);
16             u = g[u][cur_edge[u]++];
17         }
18         vertices.push_back(u);
19     }
20     reverse(vertices.begin(), vertices.end());
21     return vertices;
22 }

```

7.2.2 Undirected graph

```

1  /**
2   * Hierholzer's algorithm.
3   * Description: An Eulerian path in a undirected graph is a path that
4   * visits all edges exactly once.
5   * An Eulerian cycle is a Eulerian path that is a cycle.
6   * Time complexity: O(|E|).
7   */
8  struct Edge {
9      int to;
10     list<Edge>::iterator reverse_edge;
11     Edge(int _to) : to(_to) {}
12 };
13 vector<int> vertices;
14 void find_path(vector<list<Edge>> &g, int u) {
15     while (!g[u].empty()) {
16         int v = g[u].front().to;
17         g[v].erase(g[u].front().reverse_edge);
18         g[u].pop_front();
19         find_path(g, v);
20     }
21     vertices.emplace_back(u); // reversion list.
22 }
23 void add_edge(int u, int v) {
24     g[u].emplace_front(v);
25     g[v].emplace_front(u);
26     g[u].front().reverse_edge = g[v].begin();
27     g[v].front().reverse_edge = g[u].begin();
28 }

```

8 Misc.

8.1 Ternary search

```

1  const double eps = 1e-9;
2  double ternary_search_max(double l, double r) {
3      // find x0 such that: f(x0) > f(x), \all x: l <= x <= r.
4      while (r - l > eps) {
5          double mid1 = l + (r - l) / 3;
6          double mid2 = r - (r - l) / 3;
7          if (f(mid1) < f(mid2)) l = mid1;
8          else r = mid2;

```

```

9     }
10    return l;
11 }
12 double ternary_search_min(double l, double r) {
13     // find x0 such that: f(x0) < f(x), \all x: l <= x <= r.
14     while (r - l > eps) {
15         double mid1 = l + (r - l) / 3;
16         double mid2 = r - (r - l) / 3;
17         if (f(mid1) > f(mid2)) l = mid1;
18         else r = mid2;
19     }
20     return l;
21 }

```

8.2 Dutch flag national problem

```

1 void dutch_flag_national(vector<int> &arr) {
2     // All elements that are LESS than pivot are moved to the LEFT.
3     // All elements that are GREATER than pivot are moved to the RIGHT.
4     // E.g. [1, 2, 0, 0, 2, 2, 1], pivot = 1 -> [0, 0, 1, 1, 2, 2, 2].
5     int n = (int) arr.size();
6     int i = 0, j = 0, k = n - 1;
7     int pivot = 1;
8     // 0....i....j....k....n
9     while (j <= k) {
10        if (arr[j] < pivot) {
11            swap(arr[i], arr[j]);
12            i++;
13            j++;
14        }
15        else if (arr[j] > pivot) {
16            swap(arr[j], arr[k]);
17            k--;
18        }
19        else {
20            j++;
21        }
22    }
23    // 0 <= index <= i - 1: arr[index] < mid.
24    // i <= index <= k: arr[index] = mid.
25    // k + 1 <= index < sz: arr[index] > mid.
26 }

```

8.3 Matrix

```

1 struct Matrix {
2     static const matrix_type INF = numeric_limits<matrix_type>::max();
3     int N, M;
4     vector<vector<matrix_type>> mat;
5
6     Matrix(int _N, int _M, matrix_type v = 0) : N(_N), M(_M) {
7         mat.assign(N, vector<matrix_type>(M, v));
8     }

```

```

9     static Matrix identity(int n) { // return identity matrix.
10        Matrix I(n, n);
11        for (int i = 0; i < n; ++i) {
12            I[i][i] = 1;
13        }
14        return I;
15    }
16
17    vector<matrix_type>& operator[](int r) { return mat[r]; }
18    const vector<matrix_type>& operator[](int r) const { return mat[r]; }
19
20    Matrix& operator*=(const Matrix &other) {
21        assert(M == other.N); // [N x M] [other.N x other.M]
22        Matrix res(N, other.M);
23        for (int r = 0; r < N; ++r) {
24            for (int c = 0; c < other.M; ++c) {
25                long long square_mod = (long long) MOD * MOD;
26                long long sum = 0;
27                for (int g = 0; g < M; ++g) {
28                    sum += (long long) mat[r][g] * other[g][c];
29                    if (sum >= square_mod) sum -= square_mod;
30                }
31                res[r][c] = sum % MOD;
32            }
33        }
34        mat.swap(res.mat); return *this;
35    }
36 };

```

8.4 Debugging

```

1 #define debug(...) { string _s = #__VA_ARGS__; replace(begin(_s), end(_s),
2     ',', ' '); stringstream _ss(_s); istream_iterator<string> _it(_ss);
3     out_error(_it, __VA_ARGS__); }
4
5 void out_error(istream_iterator<string> it) { cerr << '\n'; }
6
7 template<typename T, typename ...Args>
8 void out_error(istream_iterator<string> it, T a, Args... args) {
9     cerr << " [" << *it << " = " << a << "]" ";
10    out_error(++it, args...);
11 }
12
13 template<typename T, typename G> ostream& operator<<(ostream &os, const
14    pair<T, G> &p) {
15    return os << "(" << p.first << ", " << p.second << ")";
16 }
17
18 template<class Con, class = decltype(begin(declval<Con>()))>
19 typename enable_if<!is_same<Con, string>::value, ostream&>::type
20 operator<<(ostream& os, const Con& container) {
21     os << "{";
22     for (auto it = container.begin(); it != container.end(); ++it)

```

```
20         os << (it == container.begin() ? "" : ", ") << *it;
21     return os << "}";
22 }
```