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CTU.NegativeZero

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1 Contest

1.1 C++

```
#include <bits/stdc++.h>
using namespace std;
4 #ifdef LOCAL
5 #include "cp/debug.h"
6 #else
7 #define debug(...)
8 #endif
nt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
12 const int MOD = (int) 1e9 + 7;
13 const int INF = 0x3f3f3f3f3f;
14
15 int main() {
      ios::sync_with_stdio(false); cin.tie(nullptr);
      // freopen("input.txt", "r", stdin);
17
18
      // freopen("output.txt", "w", stdout);
20
      return 0;
21 }
1.2 Debug
#define debug(...) { string _s = #__VA_ARGS__; replace(begin(_s), end(_s),
       ',', ''); stringstream _ss(_s); istream_iterator<string> _it(_ss);
      out_error(_it, __VA_ARGS__);}
void out_error(istream_iterator<string> it) { cerr << '\n'; }</pre>
5 template < typename T, typename ...Args >
6 void out_error(istream_iterator<string> it, T a, Args... args) {
      cerr << " [" << *it << " = " << a << "] ";
      out_error(++it, args...);
9 }
11 template<typename T, typename G> ostream& operator<<(ostream &os, const
      pair<T, G> &p) {
      return os << "(" << p.first << ", " << p.second << ")";</pre>
13 }
15 template < class Con, class = decltype(begin(declval < Con > ())) >
16 typename enable_if<!is_same<Con, string>::value, ostream&>::type
operator << (ostream& os, const Con& container) {</pre>
      os << "{";
19
      for (auto it = container.begin(); it != container.end(); ++it)
          os << (it == container.begin() ? "" : ", ") << *it;
20
      return os << "}";
21
22 }
```

1.3 **Java**

```
import java.io.BufferedReader;
import java.util.StringTokenizer;
3 import java.io.IOException;
4 import java.io.InputStreamReader;
5 import java.io.PrintWriter;
6 import java.util.ArrayList;
7 import java.util.Arrays;
8 import java.util.Collections;
9 import java.util.Random;
 public class Main {
     public static void main(String[] args) {
          FastScanner fs = new FastScanner();
          PrintWriter out = new PrintWriter(System.out);
          int n = fs.nextInt();
          out.println(n):
          out.close(); // don't forget this line.
      static class FastScanner {
          BufferedReader br;
          StringTokenizer st:
          public FastScanner() {
              br = new BufferedReader(new InputStreamReader(System.in));
              st = null;
          public String next() {
              while (st == null || st.hasMoreTokens() == false) {
                  try {
                      st = new StringTokenizer(br.readLine());
                  catch (IOException e) {
                      throw new RuntimeException(e);
              return st.nextToken();
         }
          public int nextInt() {
              return Integer.parseInt(next());
         public long nextLong() {
              return Long.parseLong(next());
         }
          public double nextDouble() {
              return Double.parseDouble(next());
     }
```

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1.4 sublime-build

```
"cmd": ["g++", "-std=c++17", "-fmax-errors=5", "-DLOCAL", "-Wall", "-
Wextra", "-o", "${file_path}/${file_base_name}.out", "${file}"],
"file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)?:? (.*)$",
"working_dir": "${file_path}",
"selector": "source.cpp, source.c++"

1.5 .bashrc

alias cpp='g++ -std=c++17 -fmax-errors=5 -DLOCAL -Wall -Wextra'

#Stress-testing
```

out && diff -Z out ans && echo "Test \$i passed!!" || break;

./gen.out > in && ./"\$CHECKER.out" < in > ans && ./"\$SOL.out" < in >

2 Data structures

for i in {1..100};

2.1 Sparse table

4 function test {

CHECKER=\$2

SOL = 1

11 }

25 }

```
1 int st[MAXN][K + 1];
2 for (int i = 0; i < N; i++) {
      st[i][0] = f(array[i]);
4 }
5 for (int j = 1; j \le K; j++) {
      for (int i = 0; i + (1 << j) <= N; i++) {
          st[i][j] = f(st[i][j-1], st[i+(1 << (j-1))][j-1]);
9 }
10 // Range Minimum Queries.
int lg[MAXN + 1];
12 \log[1] = 0;
13 for (int i = 2; i \le MAXN; i++) {
      lg[i] = lg[i / 2] + 1;
14
15 }
int j = lg[R - L + 1];
int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);</pre>
18 // Range Sum Queries.
19 long long sum = 0;
20 for (int j = K; j >= 0; j--) {
      if ((1 << j) <= R - L + 1) {
          sum += st[L][i];
22
          L += 1 << j;
23
24
```

2.2 Ordered set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
 using namespace __gnu_pbds;
5 template<typename key_type>
using set_t = tree<key_type, null_type, less<key_type>, rb_tree_tag,
      tree_order_statistics_node_update>;
9 void example() {
      vector < int > nums = \{1, 2, 3, 5, 10\};
      set_t<int> st(nums.begin(), nums.end());
      cout << *st.find_by_order(0) << '\n'; // 1
      assert(st.find_by_order(-INF) == st.end());
      assert(st.find_by_order(INF) == st.end());
      cout << st.order_of_key(2) << '\n'; // 1</pre>
      cout << st.order_of_key(4) << '\n'; // 3</pre>
      cout << st.order_of_key(9) << '\n'; // 4</pre>
      cout << st.order_of_key(-INF) << '\n'; // 0</pre>
      cout << st.order_of_key(INF) << '\n'; // 5</pre>
22 }
2.3
     Dsu
1 struct Dsu {
      int n;
      vector<int> par, sz;
      Dsu(int _n) : n(_n) {
          sz.resize(n, 1);
          par.resize(n);
          iota(par.begin(), par.end(), 0);
      int find(int v) {
          // finding leader/parrent of set that contains the element v.
          // with {path compression optimization}.
          return (v == par[v] ? v : par[v] = find(par[v]));
      bool same(int u, int v) {
          return find(u) == find(v);
      bool unite(int u, int v) {
          u = find(u); v = find(v);
          if (u == v) return false;
          if (sz[u] < sz[v]) swap(u, v);
          par[v] = u;
          sz[u] += sz[v];
          return true:
      vector<vector<int>> groups() {
          // returns the list of the "list of the vertices in a connected
```

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```
component".
                                                                                            if (lazy_set[id] == 0 && lazy_add[id] == 0) return;
          vector<int> leader(n);
                                                                                  32
                                                                                            int mid = (1 + r) >> 1;
27
                                                                                            for (int child : {id * 2, id * 2 + 1}) {
          for (int i = 0; i < n; ++i) {
              leader[i] = find(i);
                                                                                                int range = (child == id * 2 ? mid - l + 1 : r - mid);
          }
                                                                                                if (lazy_set[id] != 0) {
          vector<int> id(n, -1);
                                                                                                     lazy_add[child] = 0;
31
          int count = 0;
                                                                                                     lazy_set[child] = lazy_set[id];
32
          for (int i = 0; i < n; ++i) {
                                                                                                     tree[child] = range * lazy_set[id];
33
              if (id[leader[i]] == -1) {
34
                                                                                                lazy_add[child] += lazy_add[id];
                  id[leader[i]] = count++;
              }
                                                                                                tree[child] += range * lazy_add[id];
          }
          vector<vector<int>> result(count);
                                                                                            lazy_add[id] = lazy_set[id] = 0;
          for (int i = 0; i < n; ++i) {
39
              result[id[leader[i]]].push_back(i);
                                                                                        }
40
          }
                                                                                        void update(int id, int l, int r, int u, int v, int amount, bool
41
42
          return result;
                                                                                         set_value = false) {
                                                                                            if (r < u \mid \mid 1 > v) return;
43
44 };
                                                                                            if (u <= 1 && r <= v) {
                                                                                                if (set_value) {
      Segment tree
 2.4
                                                                                                     tree[id] = 1LL * amount * (r - l + 1);
                                                                                                     lazy_set[id] = amount;
                                                                                                     lazy_add[id] = 0; // clear all previous updates.
2 * Description: A segment tree with range updates and sum queries that
      supports three types of operations:
                                                                                                else {
* + Increase each value in range [1, r] by x (i.e. a[i] += x).
                                                                                                     tree[id] += 1LL * amount * (r - 1 + 1);
  * + Set each value in range [1, r] to x (i.e. a[i] = x).
                                                                                                    lazy_add[id] += amount;
   * + Determine the sum of values in range [1, r].
                                                                                                }
6 */
                                                                                                return:
7 struct SegmentTree {
                                                                                            }
      int n;
                                                                                            push(id, 1, r);
      vector<long long> tree, lazy_add, lazy_set;
                                                                                            int mid = (1 + r) >> 1;
      SegmentTree(int _n) : n(_n) {
10
                                                                                            update(id * 2, 1, mid, u, v, amount, set_value);
11
          int p = 1:
                                                                                            update(id * 2 + 1, mid + 1, r, u, v, amount, set_value);
          while (p < n) p *= 2;
12
                                                                                            tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
13
          tree.resize(p * 2);
          lazy_add.resize(p * 2);
14
                                                                                        long long get(int id, int l, int r, int u, int v) {
          lazy_set.resize(p * 2);
15
                                                                                            if (r < u | | 1 > v) return 0;
16
                                                                                            if (u <= 1 && r <= v) {
      long long merge(const long long &left, const long long &right) {
17
                                                                                                return tree[id];
          return left + right;
18
19
                                                                                            push(id, 1, r);
20
      void build(int id, int 1, int r, const vector<int> &arr) {
                                                                                            int mid = (1 + r) >> 1;
          if (1 == r) {
21
                                                                                            long long left = get(id * 2, 1, mid, u, v);
              tree[id] += arr[l];
22
                                                                                            long long right = get(id * 2 + 1, mid + 1, r, u, v);
              return;
23
                                                                                            return merge(left, right);
24
          }
          int mid = (1 + r) >> 1;
25
                                                                                  77 };
          build(id * 2, 1, mid, arr);
          build(id * 2 + 1, mid + 1, r, arr);
27
                                                                                         Efficient segment tree
          tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
28
                                                                                  1 template < typename T> struct SegmentTree {
29
      void push(int id, int 1, int r) {
                                                                                         int n;
```

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```
vector<T> tree:
                                                                                            tree.emplace_back(left, right, tree[left].val + tree[right].val);
      SegmentTree(int _n) : n(_n), tree(2 * n) {}
                                                                                 20
                                                                                            return tree.size() - 1;
      T merge(const T &left, const T &right) {
          return left + right;
                                                                                        int add(int x, int 1, int r, int u, int v, int amt) {
                                                                                            if (1 > v \mid | r < u) return x;
      template<tvpename G>
                                                                                            if (u <= 1 && r <= v) {
      void build(const vector<G> &initial) {
                                                                                                 tree.emplace_back(tree[x].l, tree[x].r, tree[x].val + 1LL * amt
                                                                                         * (r - l + 1), tree[x].lazy + amt);
          assert((int) initial.size() == n);
10
                                                                                                tree.back().has_changed = true;
          for (int i = 0; i < n; ++i) {
11
              tree[i + n] = initial[i];
                                                                                                return tree.size() - 1;
12
13
          for (int i = n - 1; i > 0; --i) {
                                                                                            int mid = (1 + r) >> 1:
14
              tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
                                                                                            push(x, 1, mid, r);
15
                                                                                            int left = add(tree[x].1, 1, mid, u, v, amt);
          }
16
17
      }
                                                                                            int right = add(tree[x].r, mid + 1, r, u, v, amt);
      void modify(int i, int v) {
18
                                                                                            tree.emplace_back(left, right, tree[left].val + tree[right].val, 0)
19
          tree[i += n] = v;
          for (i /= 2; i > 0; i /= 2) {
                                                                                            return tree.size() - 1;
20
21
              tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
          }
                                                                                        long long get_sum(int x, int l, int r, int u, int v) {
22
23
                                                                                            if (r < u \mid | 1 > v) return 0;
      T get_sum(int 1, int r) {
                                                                                            if (u <= 1 && r <= v) return tree[x].val;
24
25
          // sum of elements from 1 to r - 1.
                                                                                            int mid = (1 + r) / 2;
          T ret{};
                                                                                            push(x, 1, mid, r);
26
                                                                                            return get_sum(tree[x].1, 1, mid, u, v) + get_sum(tree[x].r, mid +
          for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
              if (1 & 1) ret = merge(ret, tree[1++]);
                                                                                         1, r, u, v);
28
              if (r & 1) ret = merge(ret, tree[--r]);
29
          }
                                                                                        void push(int x, int 1, int mid, int r) {
30
                                                                                            if (!tree[x].has_changed) return;
          return ret;
31
                                                                                            Vertex left = tree[tree[x].1];
32
33 };
                                                                                            Vertex right = tree[tree[x].r];
                                                                                            tree.emplace_back(left);
2.6 Persistent lazy segment tree
                                                                                            tree[x].l = tree.size() - 1;
                                                                                            tree.emplace_back(right);
1 struct Vertex {
                                                                                            tree[x].r = tree.size() - 1;
      int 1, r;
      long long val, lazy;
                                                                                             tree[tree[x].1].val += tree[x].lazy * (mid - 1 + 1);
      bool has_changed = false;
                                                                                             tree[tree[x].1].lazy += tree[x].lazy;
      Vertex() {}
      Vertex(int _1, int _r, long long _val, int _lazy = 0) : l(_l), r(_r),
                                                                                            tree[tree[x].r].val += tree[x].lazy * (r - mid);
      val(_val), lazy(_lazy) {}
                                                                                            tree[tree[x].r].lazy += tree[x].lazy;
7 };
8 struct PerSegmentTree {
                                                                                            tree[tree[x].1].has_changed = true;
      vector<Vertex> tree;
                                                                                            tree[tree[x].r].has_changed = true;
      vector<int> root:
                                                                                            tree[x].lazy = 0;
      int build(const vector<int> &arr, int 1, int r) {
11
                                                                                             tree[x].has_changed = false;
          if (1 == r) {
12
              tree.emplace_back(-1, -1, arr[l]);
13
```

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63 };

Disjoint sparse table

2 * Description: range query on a static array.

14

15

16

17

}

return tree.size() - 1;

int left = build(arr, 1, mid);

int right = build(arr, mid + 1, r);

int mid = (1 + r) / 2;

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```
* Time: O(1) per query.
                                                                                            fenw.assign(n, 0); // normal fenwick tree.
4 * Tested: stress-test.
                                                                                        }
5 */
                                                                                        void build(const vector<int> &A) {
6 const int MOD = (int) 1e9 + 7;
                                                                                             assert((int) A.size() == n);
7 struct DisjointSparseTable { // product queries.
                                                                                            vector<int> diff(n);
      int n. h:
                                                                                            diff[0] = A[0];
      vector<vector<int>> dst;
                                                                                             for (int i = 1; i < n; ++i) {
      vector<int> lq;
                                                                                                 diff[i] = A[i] - A[i - 1];
10
      DisjointSparseTable(int _n) : n(_n) {
11
          h = 1; // in case n = 1: h = 0 !!.
                                                                                             fenw_coeff[0] = (long long) diff[0] * n;
12
          int p = 1;
                                                                                             fenw[0] = diff[0];
13
          while (p < n) p *= 2, h++;
                                                                                            for (int i = 1; i < n; ++i) {
14
                                                                                                 fenw_coeff[i] = fenw_coeff[i - 1] + (long long) diff[i] * (n -
          lg.resize(p); lg[1] = 0;
          for (int i = 2; i < p; ++i) {
                                                                                        i);
16
              lg[i] = 1 + lg[i / 2];
                                                                                                 fenw[i] = fenw[i - 1] + diff[i];
18
          dst.resize(h, vector<int>(n));
                                                                                            for (int i = n - 1; i >= 0; --i) {
19
                                                                                                 int j = (i & (i + 1)) - 1;
20
21
      void build(const vector<int> &A) {
                                                                                                 if (j >= 0) {
          for (int lv = 0; lv < h; ++lv) {
                                                                                                     fenw_coeff[i] -= fenw_coeff[j];
22
              int len = (1 << lv);
                                                                                                     fenw[i] -= fenw[j];
23
              for (int k = 0; k < n; k += len * 2) {
                                                                                                }
24
                  int mid = min(k + len, n);
                                                                                            }
25
                  dst[lv][mid - 1] = A[mid - 1] % MOD;
                                                                                        }
26
                  for (int i = mid - 2; i >= k; --i) {
                                                                                         void add(vector<tree_type> &fenw, int i, tree_type val) {
                       dst[lv][i] = 1LL * A[i] * dst[lv][i + 1] % MOD;
                                                                                            while (i < n) {
                                                                                                fenw[i] += val;
                  if (mid == n) break;
                                                                                                 i = (i + 1);
                  dst[lv][mid] = A[mid] % MOD;
                  for (int i = mid + 1; i < min(mid + len, n); ++i) {
                                                                                        }
                       dst[lv][i] = 1LL * A[i] * dst[lv][i - 1] % MOD;
                                                                                         tree_type __prefix_sum(vector<tree_type> &fenw, int i) {
                                                                                            tree_type res{};
34
              }
                                                                                            while (i >= 0) {
35
36
                                                                                                 res += fenw[i];
                                                                                                i = (i \& (i + 1)) - 1;
37
      int get(int 1, int r) {
38
          if (1 == r) {
                                                                                            return res;
39
              return dst[0][1];
                                                                                         tree_type prefix_sum(int i) {
42
          int i = lg[l ^ r];
                                                                                            return __prefix_sum(fenw_coeff, i) - __prefix_sum(fenw, i) * (n - i
          return 1LL * dst[i][l] * dst[i][r] % MOD;
                                                                                         - 1):
43
44
45 };
                                                                                         void range_add(int 1, int r, tree_type val) {
                                                                                             add(fenw_coeff, 1, (n - 1) * val);
      Fenwick tree
                                                                                             add(fenw\_coeff, r + 1, (n - r - 1) * (-val));
                                                                                             add(fenw, 1, val);
using tree_type = long long;
                                                                                             add(fenw, r + 1, -val);
2 struct FenwickTree {
      int n;
                                                                                         tree_type range_sum(int 1, int r) {
      vector<tree_type> fenw_coeff, fenw;
                                                                                            return prefix_sum(r) - prefix_sum(l - 1);
      FenwickTree() {}
      FenwickTree(int _n) : n(_n) {
          fenw_coeff.assign(n, 0); // fenwick tree with coefficient (n - i).
```

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2.9 Implicit treap

```
1 struct Node {
      int val, prior, cnt;
      bool rev;
      Node *left, *right;
      Node() {}
      Node(int _val) : val(_val), prior(rng()), cnt(1), rev(false), left(
      nullptr), right(nullptr) {}
7 };
8 // Binary search tree + min-heap.
9 struct Treap {
      Node *root;
      Treap() : root(nullptr) {}
11
      int get_cnt(Node *n) { return n ? n->cnt : 0; }
      void upd_cnt(Node *&n) {
13
          if (n) n->cnt = get_cnt(n->left) + get_cnt(n->right) + 1;
14
15
      void push_rev(Node *treap) {
16
17
          if (!treap || !treap->rev) return;
          treap->rev = false;
18
          swap(treap->left, treap->right);
19
20
          if (treap->left) treap->left->rev ^= true;
          if (treap->right) treap->right->rev ^= true;
21
22
      pair<Node*, Node*> split(Node *treap, int x, int smaller = 0) {
23
24
          if (!treap) return {};
          push_rev(treap);
25
          int idx = smaller + get_cnt(treap->left); // implicit val.
26
          if (idx <= x) {
27
               auto pr = split(treap->right, x, idx + 1);
28
               treap->right = pr.first;
29
              upd_cnt(treap);
              return {treap, pr.second};
31
          }
32
          else {
33
               auto pl = split(treap->left, x, smaller);
               treap->left = pl.second;
              upd_cnt(treap);
37
              return {pl.first, treap};
          }
38
39
      Node* merge(Node *1, Node *r) {
40
          push_rev(1); push_rev(r);
41
          if (!l || !r) return (l ? l : r);
42
          if (l->prior < r->prior) {
              1->right = merge(1->right, r);
44
              upd_cnt(1);
              return 1;
          }
          else {
48
              r->left = merge(1, r->left);
49
```

```
upd_cnt(r);
              return r;
      }
      void insert(int pos, int val) {
          if (!root) {
              root = new Node(val);
              return:
          Node *1, *m, *r;
          m = new Node(val);
          tie(1, r) = split(root, pos - 1);
          root = merge(1, merge(m, r));
      void erase(int pos_l, int pos_r) {
          Node *1, *m, *r;
          tie(l, r) = split(root, pos_l - 1);
          tie(m, r) = split(r, pos_r - pos_l);
          root = merge(1, r);
      void reverse(int pos_l, int pos_r) {
          Node *1, *m, *r;
          tie(l, r) = split(root, pos_l - 1);
          tie(m, r) = split(r, pos_r - pos_l);
          m->rev ^= true;
          root = merge(1, merge(m, r));
      int query(int pos_l, int pos_r);
           // returns answer for corresponding types of query.
      void inorder(Node *n) {
          if (!n) return;
          push_rev(n);
          inorder(n->left);
          cout << n->val << ' ';
          inorder(n->right);
      void print() {
          inorder(root);
          cout << '\n';</pre>
90 };
```

3 Mathematics

3.1 Trigonometry

3.1.1 Sum - difference identities

```
\sin(u \pm v) = \sin(u)\cos(v) \pm \cos(u)\sin(v)
\cos(u \pm v) = \cos(u)\cos(v) \mp \sin(u)\sin(v)
\tan(u \pm v) = \frac{\tan(u) \pm \tan(v)}{1 \mp \tan(u)\tan(v)}
```

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3.1.2 Sum to product identities

$$\cos(u) + \cos(v) = 2\cos(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\cos(u) - \cos(v) = -2\sin(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

$$\sin(u) + \sin(v) = 2\sin(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\sin(u) - \sin(v) = 2\cos(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

3.1.3 Product identities

$$\cos(u)\cos(v) = \frac{1}{2}[\cos(u+v) + \cos(u-v)]$$

$$\sin(u)\sin(v) = -\frac{1}{2}[\cos(u+v) - \cos(u-v)]$$

$$\sin(u)\cos(v) = \frac{1}{2}[\sin(u+v) + \sin(u-v)]$$

3.1.4 Double - triple angle identities

$$\sin(2u) = 2\sin(u)\cos(u)$$

$$\cos(2u) = 2\cos^{2}(u) - 1 = 1 - 2\sin^{2}(u)$$

$$\tan(2u) = \frac{2\tan(u)}{1 - \tan^{2}(u)}$$

$$\sin(3u) = 3\sin(u) - 4\sin^{3}(u)$$

$$\cos(3u) = 4\cos^{3}(u) - 3\cos(u)$$

$$\tan(3u) = \frac{3\tan(u) - \tan^{3}(u)}{1 - 3\tan^{2}(u)}$$

3.2 Sums

$$n^{a} + n^{a+1} + \dots + n^{b} = \frac{n^{b+1} - n^{a}}{n-1}, \ n \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(n+1)(2n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \left(\frac{n(n+1)}{2}\right)^{2}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$1^{5} + 2^{5} + 3^{5} + \dots + n^{5} = \frac{n^{2}(n+1)^{2}(2n^{2} + 2n - 1)}{12}$$

$$1^{6} + 2^{6} + 3^{6} + \dots + n^{6} = \frac{n(n+1)(2n+1)(3n^{4} + 6n^{3} - 3n + 1)}{42}$$

$$1^{7} + 2^{7} + 3^{7} + \dots + n^{7} = \frac{n^{2}(n+1)^{2}(3n^{4} + 6n^{3} - n^{2} - 4n + 2)}{24}$$

4 String

4.1 Prefix function

```
2 * Description: The prefix function of a string 's' is defined as an array
      pi of length n,
     where pi[i] is the length of the longest proper prefix of the substring
     s[0..i] which is also a suffix of this substring.
  * Time complexity: O(|S|).
vector<int> prefix_function(const string &s) {
     int n = (int) s.length();
      vector<int> pi(n);
     pi[0] = 0;
      for (int i = 1; i < n; ++i) {
          int j = pi[i - 1]; // try length pi[i - 1] + 1.
         while (j > 0 && s[j] != s[i]) {
             j = pi[j - 1];
         if (s[j] == s[i]) {
             pi[i] = j + 1;
     return pi;
```

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4.2 Counting occurrences of each prefix

```
vector<int> count_occurrences(const string &s) {
      vector<int> pi = prefix_function(s);
      int n = (int) s.size();
      vector<int> ans(n + 1);
      for (int i = 0; i < n; ++i) {
           ans[pi[i]]++;
      for (int i = n - 1; i > 0; --i) {
           ans[pi[i - 1]] += ans[i];
10
      for (int i = 0; i <= n; ++i) {</pre>
11
           ans[i]++;
12
      }
13
      return ans;
14
      // Input: ABACABA
15
16
      // Output: 4 2 2 1 1 1 1
17 }
```

4.3 Knuth-Morris-Pratt algorithm

```
1 /**
  * Searching for a substring in a string.
  * Time complexity: O(N + M).
5 vector<int> KMP(const string &text, const string &pattern) {
      int n = (int) text.length();
      int m = (int) pattern.length();
      string s = pattern + '$' + text;
      vector<int> pi = prefix_function(s);
      vector<int> indices;
10
      for (int i = 0; i < (int) s.length(); ++i) {</pre>
11
           if (pi[i] == m) {
12
               indices.push_back(i - 2 * m);
13
14
          }
15
      return indices;
16
17 }
```

4.4 Suffix array

```
struct SuffixArray {
    string s;
    int n, lim;
    vector<int> sa, lcp, rank;
    SuffixArray(const string &_s, int _lim = 256) : s(_s), n(s.length() +
    1), lim(_lim), sa(n), lcp(n), rank(n) {
        s += '$';
        build();
        kasai();
        sa.erase(sa.begin());
    lcp.erase(lcp.begin());
    s.pop_back();
```

```
void build() {
          vector<int> nrank(n), norder(n), cnt(max(n, lim));
          for (int i = 0; i < n; ++i) {
              sa[i] = i; rank[i] = s[i];
          for (int k = 0, rank_cnt = 0; rank_cnt < n - 1; k = max(1, k * 2),
      lim = rank_cnt + 1) {
              // counting sort.
              for (int i = 0; i < n; ++i) norder[i] = (sa[i] - k + n) % n;
              for (int i = 0; i < n; ++i) cnt[rank[i]]++;</pre>
              for (int i = 1; i < lim; ++i) cnt[i] += cnt[i - 1];</pre>
              for (int i = n - 1; i >= 0; --i) sa[--cnt[rank[norder[i]]]] =
      norder[i];
              rank[sa[0]] = rank\_cnt = 0;
              for (int i = 1; i < n; ++i) {
                  int u = sa[i], v = sa[i - 1];
                  int nu = u + k, nv = v + k;
                  if (nu >= n) nu -= n;
                  if (nv >= n) nv -= n;
                  if (rank[u] != rank[v] || rank[nu] != rank[nv]) ++rank_cnt;
                  nrank[sa[i]] = rank_cnt;
              for (int i = 0; i < rank_cnt + 1; ++i) cnt[i] = 0;</pre>
              rank.swap(nrank);
          }
      }
      void kasai() {
          for (int i = 0; i < n; ++i) rank[sa[i]] = i;
          for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
              int j = sa[rank[i] - 1];
              while (s[i + k] == s[j + k]) k++;
              lcp[rank[i]] = k;
          // Note: lcp[i] = longest common prefix(sa[i - 1], sa[i]).
46 };
      Manacher's algorithm
2 * Description: for each position, computes d[0][i] = half length of
longest palindrome centered on i (rounded up), d[1][i] = half length of
  longest palindrome centered on i and i - 1.
  * Time complexity: O(N).
   * Tested: https://judge.yosupo.jp/problem/enumerate_palindromes, stress-
      tested.
8 array<vector<int>, 2> manacher(const string &s) {
      int n = (int) s.size();
      array<vector<int>, 2> d;
      for (int z = 0; z < 2; ++z) {
```

d[z].resize(n);

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```
int 1 = 0, r = 0;
           for (int i = 0; i < n; ++i) {
14
               int mirror = 1 + r - i + z;
              d[z][i] = (i > r ? 0 : min(d[z][mirror], r - i));
               int L = i - d[z][i] - z, R = i + d[z][i];
               while (L >= 0 \&\& R < n \&\& s[L] == s[R]) {
18
                   d[z][i]++; L--; R++;
19
              }
20
              if (R > r) {
                   1 = L; r = R;
22
              }
23
          }
24
      }
25
26
      return d;
27 }
      Trie
1 struct Trie {
      const static int ALPHABET = 26;
      const static char minChar = 'a';
      struct Vertex {
          int next[ALPHABET];
          bool leaf;
          Vertex() {
              leaf = false:
               fill(next, next + ALPHABET, -1);
          }
10
      };
11
12
      vector<Vertex> trie;
13
      Trie() { trie.emplace_back(); }
14
      void insert(const string &s) {
15
16
          int i = 0:
           for (const char &ch : s) {
17
18
               int j = ch - minChar;
              if (trie[i].next[j] == -1) {
19
                   trie[i].next[j] = trie.size();
20
                   trie.emplace_back();
21
              i = trie[i].next[j];
          }
24
25
          trie[i].leaf = true;
26
      bool find(const string &s) {
27
          int i = 0;
28
          for (const char &ch : s) {
29
               int j = ch - minChar;
30
               if (trie[i].next[i] == -1) {
31
                   return false;
32
33
              i = trie[i].next[j];
34
35
          }
```

```
return (trie[i].leaf ? true : false);
      }
38 };
      Hashing
 struct Hash61 {
      static const uint64_t MOD = (1LL << 61) - 1;</pre>
      static uint64_t BASE;
      static vector<uint64_t> pw;
      uint64_t addmod(uint64_t a, uint64_t b) const {
          a += b;
          if (a >= MOD) a -= MOD;
          return a;
      uint64_t submod(uint64_t a, uint64_t b) const {
          a += MOD - b;
          if (a >= MOD) a -= MOD;
          return a;
      uint64_t mulmod(uint64_t a, uint64_t b) const {
          uint64_t low1 = (uint32_t) a, high1 = (a >> 32);
          uint64_t low2 = (uint32_t) b, high2 = (b >> 32);
          uint64_t low = low1 * low2;
          uint64_t mid = low1 * high2 + low2 * high1;
          uint64_t high = high1 * high2;
          uint64_t ret = (low & MOD) + (low >> 61) + (high << 3) + (mid >>
       29) + (mid << 35 >> 3) + 1;
          // ret %= MOD:
          ret = (ret >> 61) + (ret & MOD);
          ret = (ret >> 61) + (ret & MOD);
          return ret - 1;
      void ensure_pw(int m) {
          int n = (int) pw.size();
          if (n >= m) return;
          pw.resize(m);
          for (int i = n; i < m; ++i) {</pre>
              pw[i] = mulmod(pw[i - 1], BASE);
          }
      }
      vector<uint64_t> pref;
       template < typename T > Hash61(const T &s) { // strings or arrays.
          n = (int) s.size();
          ensure_pw(n);
          pref.resize(n + 1);
44
          pref[0] = 0;
45
          for (int i = 0; i < n; ++i) {
              pref[i + 1] = addmod(mulmod(pref[i], BASE), s[i]);
```

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5 Number Theory

5.1 Euler's totient function

- Euler's totient function, also known as **phi-function** $\phi(n)$ counts the number of integers between 1 and n inclusive, that are **coprime to** n.
- Properties:
 - Divisor sum property: $\sum_{d|n} \phi(d) = n$.
 - $\phi(n)$ is a **prime number** when n = 3, 4, 6.
 - If *p* is a prime number, then $\phi(p) = p 1$.
 - If *p* is a prime number and *k* ≥ 1, then $\phi(p^k) = p^k p^{k-1}$.
 - If *a* and *b* are **coprime**, then $\phi(ab) = \phi(a) \cdot \phi(b)$.
 - In general, for **not coprime** a and b, with d = gcd(a, b) this equation holds: $\phi(ab) = \phi(a) \cdot \phi(b) \cdot \frac{d}{\phi(d)}$.
 - With $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\phi(n) = \phi(p_1^{k_1}) \cdot \phi(p_2^{k_2}) \cdots \phi(p_m^{k_m})$$
$$= n \cdot \left(1 - \frac{1}{p_1}\right) \cdot \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_m}\right)$$

- Application in Euler's theorem:
 - If gcd(a, M) = 1, then:

$$a^{\phi(M)} \equiv 1 \pmod{M} \Rightarrow a^n \equiv a^{n \bmod{\phi(M)}} \pmod{M}$$

- In general, for arbitrary a, M and n ≥ $log_2 M$:

$$a^n \equiv a^{\phi(M) + [n \bmod \phi(M)]} \pmod{M}$$

5.2 Mobius function

• For a positive integer $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\mu(n) = \begin{cases} 1, & \text{if } n = 1\\ 0, & \text{if } \exists k_i > 1\\ (-1)^m & \text{otherwise} \end{cases}$$

- Properties:
 - $-\sum_{d|n}\mu(d)=[n=1].$
 - If *a* and *b* are **coprime**, then $\mu(ab) = \mu(a) \cdot \mu(b)$.
 - Mobius inversion: let *f* and *g* be arithmetic functions:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right)g(d)$$

5.3 Primes

Approximating the number of primes up to *n*:

n	$\pi(n)$	$\frac{n}{\ln n - 1}$		
$100 (1e^2)$	25	28		
$500 (5e^2)$	95	96		
$1000 (1e^3)$	168	169		
$5000 (5e^3)$	669	665		
$10000 (1e^4)$	1229	1218		
$50000 (5e^4)$	5133	5092		
$100000 (1e^5)$	9592	9512		
$500000 (5e^5)$	41538	41246		
$1000000 (1e^6)$	78498	78030		
$5000000 (5e^6)$	348513	346622		

 $(\pi(n))$ = the number of primes less than or equal to n, $\frac{n}{\ln n - 1}$ is used to approximate $\pi(n)$).

5.4 Wilson's theorem

A positive integer *n* is a prime if and only if:

$$(n-1)! \equiv n-1 \pmod{n}$$

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5.5 Zeckendorf's theorem

The Zeckendorf's theorem states that every positive integer *n* can be represented uniquely as a sum of one or more distinct non-consecutive Fibonacci numbers. For example:

64 = 55 + 8 + 185 = 55 + 21 + 8 + 1

```
vector<int> zeckendoft_theorem(int n) {
       vector<int> fibs = {1, 1};
       int sz = 2;
       while (fibs.back() <= n) {</pre>
           fibs.push_back(fibs[sz - 1] + fibs[s - 2]);
           SZ++:
       fibs.pop_back();
      vector<int> nums;
      int p = sz - 1;
       while (n > 0) {
11
          if (n >= fibs[p]) {
12
               nums.push_back(fibs[p]);
               n -= fibs[p];
14
          }
15
          p--;
16
17
18
       return nums;
19 }
```

5.6 Bitwise operation

- $a + b = (a \oplus b) + 2(a \& b)$ • $a \mid b = (a \oplus b) + (a \& b)$ • $a \& (b \oplus c) = (a \& b) \oplus (a \& c)$ • $a \mid (b \& c) = (a \mid b) \& (a \mid c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$ • $a \& (b \mid c) = (a \& b) \mid (a \& c)$
- Iterating over all subsets of a set and iterating over all submasks of a mask:

```
// Time complexity: 0(3^n).
      Pollard's rho algorithm
const int PRIME_MAX = (int) 4e4; // for handle numbers <= 1e9.</pre>
const int LIMIT = (int) 1e9;
3 vector<int> primes;
5 void linear_sieve(int n);
6 num_type mulmod(num_type a, num_type b, num_type mod);
 num_type powmod(num_type a, num_type n, num_type mod);
9 bool miller_rabin(num_type a, num_type d, int s, num_type mod) {
      // \mod - 1 = a (d * 2^s).
      num_type x = powmod(a, d, mod);
      if (x == 1 || x == mod - 1) return true;
      for (int i = 1; i \le s - 1; ++i) {
          x = mulmod(x, x, mod);
          if (x == mod - 1) return true;
      return false:
19 bool is_prime(num_type n, int ITERATION = 10) {
      if (n < 4) return (n == 2 || n == 3);
      if (n % 2 == 0 || n % 3 == 0) return false;
      num\_type d = n - 1;
      int s = 0;
      while (d % 2 == 0) {
          d /= 2;
          s++;
      for (int i = 0; i < ITERATION; ++i) {
          num_type a = (num_type) (rand() % (n - 2)) + 2;
          if (miller_rabin(a, d, s, n) == false) {
              return false:
          }
      }
      return true;
36 num_type f(num_type x, int c, num_type mod) { // f(x) = (x^2 + c) % mod.
      x = mulmod(x, x, mod);
      x += c;
      if (x >= mod) x -= mod;
      return x;
42 num_type pollard_rho(num_type n, int c) {
      // algorithm to find a random divisor of 'n'.
      // using random function: f(x) = (x^2 + c) \% n.
      num_type x = 2, y = x, d;
      long long p = 1;
      int dist = 0;
```

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```
while (true) {
48
          y = f(y, c, n);
49
50
          dist++;
          d = \_gcd(llabs(x - y), n);
          if (d > 1) break;
52
          if (dist == p) { dist = 0; p *= 2; x = y; }
53
54
      return d;
55
57 void factorize(int n, vector<num_type> &factors);
void llfactorize(num_type n, vector<num_type> &factors) {
      if (n < 2) return:
      if (n < LIMIT) {</pre>
          factorize(n, factors);
61
62
          return:
63
64
      if (is_prime(n)) {
          factors.emplace_back(n);
65
          return:
66
67
      }
      num_type d = n;
68
      for (int c = 2; d == n; c++) {
69
          d = pollard_rho(n, c);
70
71
72
      llfactorize(d, factors);
      llfactorize(n / d, factors);
73
74 }
75 vector<num_type> gen_divisors(vector<pair<num_type, int>> &factors) {
      vector<num_type> divisors = {1};
76
      for (auto &x : factors) {
          int sz = (int) divisors.size();
78
          for (int i = 0; i < sz; ++i) {
79
              num_type cur = divisors[i];
80
81
              for (int j = 0; j < x.second; ++j) {
                   cur *= x.first;
                   divisors.push_back(cur);
              }
84
85
          }
86
87
      return divisors; // this array is NOT sorted yet.
88 }
      Bitset sieve
* Description: sieve of eratosthenes for large n (up to 1e9).
* Time and space (tested on codeforces):
* + For n = 1e8: ~200 ms, 6 MB.
* + For n = 1e9: ^{\circ}4000 ms, 60 MB.
6 */
7 const int N = (int) 1e8;
8 bitset<N / 2 + 1> isPrime;
9 void sieve(int n = N) {
```

```
isPrime.flip();
      isPrime[0] = false;
      for (int i = 3; i <= (int) sqrt(n); i += 2) {
          if (isPrime[i >> 1]) {
              for (int j = i * i; j <= n; j += 2 * i) {
                  isPrime[j >> 1] = false;
      }
19 }
20 void example(int n) {
      sieve(n):
      int primeCnt = (n >= 2);
      for (int i = 3; i \le n; i += 2) {
          if (isPrime[i >> 1]) {
              primeCnt++;
          }
      cout << primeCnt << '\n';</pre>
29 }
      Block sieve
1 /**
* Description: very fast sieve of eratosthenes for large n (up to 1e9).
3 * Source: kactl.
* Time and space (tested on codeforces):
   * + For n = 1e8: ~160 ms, 60 MB.
   * + For n = 1e9: ~1600 ms, 505 MB.
7 * Need to check memory limit.
9 const int N = (int) 1e8;
bitset<N + 1> is_prime;
vector<int> fast_sieve() {
      const int S = (int) sqrt(N), R = N / 2;
      vector<int> primes = {2};
      vector<bool> sieve(S + 1, true);
      vector<array<int, 2>> cp;
      for (int i = 3; i \le S; i += 2) {
          if (sieve[i]) {
              cp.push_back({i, i * i / 2});
              for (int j = i * i; j <= S; j += 2 * i) {
                  sieve[j] = false;
              }
          }
      }
      for (int L = 1; L <= R; L += S) {
          array<bool, S> block{};
          for (auto &[p, idx] : cp) {
              for (; idx < S + L; idx += p) block[idx - L] = true;</pre>
          for (int i = 0; i < min(S, R - L); ++i) {
              if (!block[i]) primes.push_back((L + i) * 2 + 1);
```

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```
31     }
32     }
33     for (int p : primes) is_prime[p] = true;
34     return primes;
35 }
```

5.10 Combinatorics

5.10.1 Catalan numbers

$$C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{n!(n+1)!}$$

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}, \ C_0 = 1, \ C_n = \frac{4n-2}{n+1} C_{n-1}$$

• The first 12 Catalan numbers (n = 0, 1, 2, ..., 11):

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$$

- Applications of Catalan numbers:
 - difference binary search trees with *n* vertices from 1 to *n*.
 - rooted binary trees with n + 1 leaves (vertices are not numbered).
 - correct bracket sequence of length 2 * n.
 - permutation [n] with no 3-term increasing subsequence (i.e. doesn't exist i < j < k for which a[i] < a[j] < a[k]).
 - ways a convex polygon of n + 2 sides can split into triangles by connecting vertices.

5.10.2 Stirling numbers of the second kind

Partitions of *n* distinct elements into exactly *k* non-empty groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} \binom{k}{i} i^n$$

5.10.3 Derangements

Permutation of the elements of a set, such that no element appears in its original position (no fixied point). Recursive formulas:

$$D(n) = (n-1)[D(n-1) + D(n-2)] = nD(n-1) + (-1)^n$$

6 Linear algebra

6.1 Gauss elimination

```
const double EPS = 1e-9;
 const int INF = 2; // it doesn't actually have to be infinity or a big
int gauss (vector < vector<double> > a, vector<double> & ans) {
      int n = (int) a.size();
      int m = (int) a[0].size() - 1;
      vector<int> where (m, -1);
      for (int col=0, row=0; col<m && row<n; ++col) {</pre>
          int sel = row;
          for (int i=row; i<n; ++i)</pre>
              if (abs (a[i][col]) > abs (a[sel][col]))
                  sel = i;
          if (abs (a[sel][col]) < EPS)</pre>
              continue:
          for (int i=col; i<=m; ++i)
              swap (a[sel][i], a[row][i]);
          where[col] = row;
          for (int i=0; i<n; ++i)</pre>
              if (i != row) {
                  double c = a[i][col] / a[row][col];
                  for (int j=col; j<=m; ++j)</pre>
                       a[i][j] -= a[row][j] * c;
          ++row;
     }
      ans.assign (m, 0);
      for (int i=0; i<m; ++i)
          if (where[i] != -1)
              ans[i] = a[where[i]][m] / a[where[i]][i];
      for (int i=0; i<n; ++i) {
          double sum = 0;
          for (int j=0; j<m; ++j)
              sum += ans[j] * a[i][j];
          if (abs (sum - a[i][m]) > EPS)
              return 0;
      for (int i=0; i<m; ++i)</pre>
          if (where[i] == -1)
              return INF;
      return 1;
```

7 Geometry

7.1 Fundamentals

7.1.1 Point

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```
const double PI = acos(-1);
const double EPS = 1e-9;
3 typedef double ftype;
4 struct point {
      ftype x, y;
       point(ftype _x = 0, ftype _y = 0): x(_x), y(_y) {}
       point& operator+=(const point& other) {
           x += other.x; y += other.y; return *this;
       point& operator -= (const point& other) {
10
          x -= other.x; y -= other.y; return *this;
11
12
      point& operator*=(ftype t) {
13
           x *= t; y *= t; return *this;
14
15
       point& operator/=(ftype t) {
16
17
          x /= t; y /= t; return *this;
18
19
      point operator+(const point& other) const {
           return point(*this) += other;
20
21
       point operator-(const point& other) const {
22
           return point(*this) -= other;
23
24
       point operator*(ftype t) const {
25
           return point(*this) *= t;
26
27
      point operator/(ftype t) const {
28
           return point(*this) /= t;
29
30
       point rotate(double angle) const {
31
32
           return point(x * cos(angle) - y * sin(angle), x * sin(angle) + y *
       cos(angle));
33
      friend istream& operator>>(istream &in, point &t);
34
       friend ostream& operator<<(ostream &out, const point& t);</pre>
35
       bool operator<(const point& other) const {</pre>
36
37
          if (fabs(x - other.x) < EPS)</pre>
               return y < other.y;</pre>
39
          return x < other.x;</pre>
      }
40
41 };
42
43 istream& operator>>(istream &in, point &t) {
       in >> t.x >> t.y;
       return in;
45
47 ostream& operator<<(ostream &out, const point& t) {
       out << t.x << ' ' << t.y;
49
      return out;
50 }
51
```

```
symple ftype dot(point a, point b) {return a.x * b.x + a.y * b.y;}
53 ftype norm(point a) {return dot(a, a);}
54 ftype abs(point a) {return sqrt(norm(a));}
ftype angle(point a, point b) {return acos(dot(a, b) / (abs(a) * abs(b)));}
ftype proj(point a, point b) {return dot(a, b) / abs(b);}
57 ftype cross(point a, point b) {return a.x * b.y - a.y * b.x;}
58 bool ccw(point a, point b, point c) {return cross(b - a, c - a) > EPS;}
59 bool collinear(point a, point b, point c) {return fabs(cross(b - a, c - a))
        < EPS:}
60 point intersect(point a1, point d1, point a2, point d2) {
       double t = cross(a2 - a1, d2) / cross(d1, d2);
       return a1 + d1 * t:
63 }
 7.1.2 Line
struct line {
       double a, b, c;
      line (double _a = \emptyset, double _b = \emptyset, double _c = \emptyset): a(_a), b(_b), c(_c)
      friend ostream & operator<<(ostream& out, const line& 1);</pre>
5 };
6 ostream & operator << (ostream & out, const line & 1) {
       out << 1.a << ' ' << 1.b << ' ' << 1.c;
      return out:
9 }
void pointsToLine(const point& p1, const point& p2, line& l) {
      if (fabs(p1.x - p2.x) < EPS)
          1 = \{1.0, 0.0, -p1.x\};
      else {
          1.a = - (double)(p1.y - p2.y) / (p1.x - p2.x);
          1.b = 1.0;
          1.c = -1.a * p1.x - 1.b * p1.y;
      }
void pointsSlopeToLine(const point& p, double m, line& 1) {
      1.a = -m:
      1.b = 1:
      1.c = -1.a * p.x - 1.b * p.y;
24 bool areParallel(const line& l1, const line& l2) {
      return fabs(11.a - 12.a) < EPS && fabs(11.b - 12.b) < EPS;
27 bool areSame(const line& 11, const line& 12) {
      return areParallel(l1, l2) && fabs(l1.c - l2.c) < EPS;</pre>
29 }
30 bool areIntersect(line 11, line 12, point& p) {
      if (areParallel(l1, l2)) return false;
      p.x = -(11.c * 12.b - 11.b * 12.c) / (11.a * 12.b - 11.b * 12.a);
      if (fabs(11.b) > EPS) p.y = -(11.c + 11.a * p.x);
      else p.y = -(12.c + 12.a * p.x);
35
      return 1;
36 }
```

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```
37 double distToLine(point p, point a, point b, point& c) {
      double t = dot(p - a, b - a) / norm(b - a);
      c = a + (b - a) * t;
      return abs(c - p);
41 }
42 double distToSegment(point p, point a, point b, point& c) {
      double t = dot(p - a, b - a) / norm(b - a);
43
      if (t > 1.0)
44
          c = point(b.x, b.y);
45
      else if (t < 0.0)
46
          c = point(a.x, a.y);
47
48
          c = a + (b - a) * t;
49
      return abs(c - p);
50
51 }
52 bool intersectTwoSegment(point a, point b, point c, point d) {
      ftype ABxAC = cross(b - a, c - a);
      ftype ABxAD = cross(b - a, d - a);
54
      ftype CDxCA = cross(d - c, a - c);
55
      ftype CDxCB = cross(d - c, b - c);
56
      if (ABxAC == 0 \mid | ABxAD == 0 \mid | CDxCA == 0 \mid | CDxCB == 0) {
57
          if (ABxAC == 0 && dot(a - c, b - c) <= 0) return true;
58
          if (ABxAD == 0 && dot(a - d, b - d) <= 0) return true;</pre>
59
          if (CDxCA == 0 && dot(c - a, d - a) <= 0) return true;</pre>
60
          if (CDxCB == 0 \&\& dot(c - b, d - b) <= 0) return true;
61
          return false;
62
63
      return (ABxAC * ABxAD < 0 && CDxCA * CDxCB < 0);</pre>
64
65 }
66 void perpendicular(line l1, point p, line& l2) {
      if (fabs(l1.a) < EPS)
67
          12 = \{1.0, 0.0, -p.x\};
68
69
      else {
70
          12.a = -11.b / 11.a;
          12.b = 1.0;
71
72
          12.c = -12.a * p.x - 12.b * p.y;
      }
73
74 }
7.1.3 Circle
int insideCircle(const point& p, const point& center, ftype r) {
      ftype d = norm(p - center);
      ftype rSq = r * r;
      return fabs(d - rSq) < EPS ? 0 : (d - rSq >= EPS ? 1 : -1);
5 }
6 bool circle2PointsR(const point& p1, const point& p2, ftype r, point& c) {
      double h = r * r - norm(p1 - p2) / 4.0;
      if (fabs(h) < 0) return false;</pre>
      h = sqrt(h);
10
      point perp = (p2 - p1).rotate(PI / 2.0);
11
      point m = (p1 + p2) / 2.0;
      c = m + perp * (h / abs(perp));
12
```

```
13
      return true:
14 }
 7.1.4 Triangle
double areaTriangle(double ab, double bc, double ca) {
       double p = (ab + bc + ca) / 2;
       return sqrt(p) * sqrt(p - ab) * sqrt(p - bc) * sqrt(p - ca);
4 }
5 double rInCircle(double ab, double bc, double ca) {
       double p = (ab + bc + ca) / 2;
      return areaTriangle(ab, bc, ca) / p;
8 }
9 double rInCircle(point a, point b, point c) {
      return rInCircle(abs(a - b), abs(b - c), abs(c - a));
11 }
12 bool inCircle(point p1, point p2, point p3, point &ctr, double &r) {
      r = rInCircle(p1, p2, p3);
      if (fabs(r) < EPS) return false;</pre>
      line 11, 12;
      double ratio = abs(p2 - p1) / abs(p3 - p1);
      point p = p2 + (p3 - p2) * (ratio / (1 + ratio));
      pointsToLine(p1, p, l1);
      ratio = abs(p1 - p2) / abs(p2 - p3);
      p = p1 + (p3 - p1) * (ratio / (1 + ratio));
      pointsToLine(p2, p, 12);
      areIntersect(l1, l2, ctr);
      return true;
25 double rCircumCircle(double ab, double bc, double ca) {
       return ab * bc * ca / (4.0 * areaTriangle(ab, bc, ca));
28 double rCircumCircle(point a, point b, point c) {
      return rCircumCircle(abs(b - a), abs(c - b), abs(a - c));
30 }
 7.1.5 Convex hull
vector<point> CH_Andrew(vector<point> &Pts) { // overall O(n log n)
      int n = Pts.size(), k = 0;
      vector<point> H(2 * n);
      sort(Pts.begin(), Pts.end());
      for (int i = 0; i < n; ++i) {
          while ((k \ge 2) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
          H[k++] = Pts[i];
      for (int i = n - 2, t = k + 1; i >= 0; --i) {
          while ((k \ge t) \& ! ccw(H[k - 2], H[k - 1], Pts[i])) --k;
          H[k++] = Pts[i];
12
      }
      H.resize(k);
      return H:
15 }
```

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7.1.6 Polygon

cx;

```
double perimeter(const vector<point> &P) {
      double ans = 0.0;
      for (int i = 0; i < (int)P.size() - 1; ++i)
          ans += abs(P[i] - P[i + 1]);
      return ans;
6 }
7 double area(const vector<point> &P) {
      double ans = 0.0;
      for (int i = 0; i < (int)P.size() - 1; ++i)</pre>
          ans += (P[i].x * P[i + 1].y - P[i + 1].x * P[i].y);
      return fabs(ans) / 2.0;
11
12 }
bool isConvex(const vector<point> &P) {
      int n = (int)P.size();
      if (n <= 3) return false;</pre>
15
      bool firstTurn = ccw(P[0], P[1], P[2]);
16
      for (int i = 1; i < n - 1; ++i)
          if (ccw(P[i], P[i + 1], P[(i + 2) == n ? 1 : i + 2]) != firstTurn)
18
              return false;
19
20
      return true;
21 }
22 int insidePolygon(point pt, const vector<point> &P) {
      int n = (int)P.size();
24
      if (n <= 3) return -1;
      bool on_polygon = false;
25
26
      for (int i = 0; i < n - 1; ++i)
          if (fabs(abs(P[i] - pt) + abs(pt - P[i + 1]) - abs(P[i] - P[i + 1])
27
      ) < EPS
              on_polygon = true;
28
      if (on_polygon) return 0;
29
      double sum = 0.0;
30
      for (int i = 0; i < n - 1; ++i) {
31
          if (ccw(pt, P[i], P[i + 1]))
32
              sum += angle(P[i] - pt, P[i + 1] - pt);
33
          else
34
              sum -= angle(P[i] - pt, P[i + 1] - pt);
35
36
      return fabs(sum) > PI ? 1 : -1;
37
      Minimum enclosing circle
* Description: computes the minimum circle that encloses all the given
      points.
4 double abs(point a) { return sqrt(a.X * a.X + a.Y * a.Y); }
6 point center_from(double bx, double by, double cx, double cy) {
```

double B = bx * bx + by * by, C = cx * cx + cy * cy, D = bx * cy - by *

```
return point((cy * B - by * C) / (2 * D), (bx * C - cx * B) / (2 * D));
9 }
11 circle circle_from(point A, point B, point C) {
      point I = center_from(B.X - A.X, B.Y - A.Y, C.X - A.X, C.Y - A.Y);
      return circle(I + A, abs(I));
14 }
16 const int N = 100005;
17 int n, x[N], y[N];
18 point a[N];
20 circle emo_welzl(int n, vector<point> T) {
      if (T.size() == 3 || n == 0) {
          if (T.size() == 0) return circle(point(0, 0), -1);
          if (T.size() == 1) return circle(T[0], 0);
          if (T.size() == 2) return circle((T[0] + T[1]) / 2, abs(T[0] - T[1])
      [1]) / 2);
          return circle_from(T[0], T[1], T[2]);
      random\_shuffle(a + 1, a + n + 1);
      circle Result = emo_welzl(0, T);
      for (int i = 1; i <= n; i++)
          if (abs(Result.X - a[i]) > Result.Y + 1e-9) {
              T.push_back(a[i]);
              Result = emo_welzl(i - 1, T);
              T.pop_back();
      return Result;
```

8 Graph

8.1 Bellman-Ford algorithm

```
1 /**
2 * Description: single source shortest path in a weighted (negative or positive) directed graph.
3 * Time: O(N * M).
4 * Tested: https://open.kattis.com/problems/shortestpath3
5 */
6 const int64_t INF = (int64_t) 2e18;
7 struct Edge {
8    int u, v; // u -> v
9    int64_t w;
10    Edge() {}
11    Edge(int _u, int _v, int64_t _w) : u(_u), v(_v), w(_w) {}
12 };
13 vector<int64_t> bellmanFord(int s) {
14    // dist[stating] = 0.
15    // dist[u] = +INF, if u is unreachable.
16    // dist[u] = -INF, if there is a negative cycle on the path from s to u
```

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```
// -INF < dist[u] < +INF, otherwise.</pre>
18
       vector<int64_t> dist(n, INF);
       dist[s] = 0;
       for (int i = 0; i < n - 1; ++i) {</pre>
20
           bool any = false;
21
           for (auto [u, v, w] : edges) {
22
               if (dist[u] != INF && dist[v] > w + dist[u]) {
23
24
                    dist[v] = w + dist[u];
                    any = true;
               }
26
27
28
           if (!any) break;
29
30
       // handle negative cycles
31
       for (int i = 0; i < n - 1; ++i) {
32
           for (auto [u, v, w] : edges) {
               if (dist[u] != INF && dist[v] > w + dist[u]) {
33
34
                    dist[v] = -INF;
35
               }
           }
36
37
38
       return dist;
39 }
```

Articulation point and Bridge

```
* Description: finding articulation points and bridges in a simple
      undirected graph.
  * Tested: https://oj.vnoi.info/problem/graph_
4 */
5 const int N = (int) 1e5;
6 vector<int> g[N];
7 int num[N], low[N], dfs_timer;
8 bool joint[N];
9 vector<pair<int, int>> bridges;
void dfs(int u, int prev) {
      low[u] = num[u] = ++dfs_timer;
      int child = 0;
12
      for (int v : g[u]) {
13
          if (v == prev) continue;
14
          if (num[v]) low[u] = min(low[u], num[v]);
15
          else {
16
              dfs(v, u);
              low[u] = min(low[u], low[v]);
18
              child++;
19
              if (low[v] >= num[v]) {
20
21
                  bridges.emplace_back(u, v);
              if (u != prev && low[v] >= num[u]) joint[u] = true;
23
          }
24
25
      }
```

```
if (u == prev && child > 1) joint[u] = true;
27 }
29 int main() {
      int n, m;
      cin >> n >> m;
      for (int i = 0; i < m; ++i) {
          int u, v;
          cin >> u >> v;
          u--; v--;
          g[u].push_back(v);
          g[v].push_back(u);
      for (int i = 0; i < n; ++i) {
          if (!num[i]) dfs(i, i);
      return 0;
43 }
      Strongly connected components
```

```
* Description: Tarjan's algorithm finds strongly connected components
      in a directed graph. If vertices u and v belong to the same component,
      then scc_id[u] == scc_id[v].
  * Tested: https://judge.yosupo.jp/problem/scc
7 const int N = (int) 5e5;
8 vector<int> g[N], st;
9 int low[N], num[N], dfs_timer, scc_id[N], scc;
10 bool used[N];
void Tarjan(int u) {
      low[u] = num[u] = ++dfs_timer;
      st.push_back(u);
      for (int v : g[u]) {
          if (used[v]) continue;
          if (num[v] == 0) {
              Tarjan(v);
              low[u] = min(low[u], low[v]);
          }
          else {
              low[u] = min(low[u], num[v]);
      if (low[u] == num[u]) {
          int v;
          do {
              v = st.back(); st.pop_back();
              debug(u, v)
              used[v] = true;
              scc_id[v] = scc;
          } while (v != u);
          scc++;
```

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```
34 }
 8.4 Topo sort
2 * Description: A topological sort of a directed acyclic graph
  * is a linear ordering of its vertices such that for every directed edge
  * from vertex u to vertex v, u comes before v in the ordering.
5 * Note: If there are cycles, the returned list will have size smaller than
       n (i.e, topo.size() < n).
* Tested: https://judge.yosupo.jp/problem/scc
8 vector<int> topo_sort(const vector<vector<int>> &q) {
      int n = (int) g.size();
      vector<int> indea(n):
      for (int u = 0; u < n; ++u) {
          for (int v : g[u]) indeg[v]++;
13
      queue < int > q; // Note: use min-heap to get the smallest lexicographical
14
      for (int u = 0; u < n; ++u) {
15
          if (indeg[u] == 0) q.emplace(u);
17
      vector<int> topo;
18
      while (!q.empty()) {
19
          int u = q.front(); q.pop();
20
21
          topo.emplace_back(u);
          for (int v : g[u]) {
23
              if (--indeg[v] == 0) q.emplace(v);
          }
24
25
      return topo;
26
27 }
      K-th smallest shortest path
1 /** Finding the k-th smallest shortest path from vertex s to vertex t,
      each vertex can be visited more than once.
3 */
4 using adj_list = vector<vector<pair<int, int>>>;
5 vector<int> k_smallest(const adj_list &q, int k, int s, int t) {
      int n = (int) g.size();
      vector<long long> ans;
      vector<int> cnt(n);
      using pli = pair<long long, int>;
      priority_queue<pli, vector<pli>, greater<pli>> pq;
      pq.emplace(0, s);
11
12
      while (!pq.empty() && cnt[t] < k) {</pre>
          int u = pq.top().second;
13
```

long long d = pq.top().first;

if (cnt[u] == k) continue;

pq.pop();

14

15

```
cnt[u]++:
          if (u == t) {
              ans.push_back(d);
          for (auto [v, cost] : q[u]) {
              pq.emplace(d + cost, v);
      assert(ans.size() == k);
      return ans;
27 }
      Eulerian path
 8.6.1 Directed graph
* Hierholzer's algorithm.
   * Description: An Eulerian path in a directed graph is a path that visits
      all edges exactly once.
      An Eulerian cycle is a Eulerian path that is a cycle.
   * Time complexity: O(|E|).
6 */
   vector<int> find_path_directed(const vector<vector<int>> &g, int s) {
      int n = (int) q.size();
      vector<int> stack, cur_edge(n), vertices;
      stack.push_back(s);
      while (!stack.empty()) {
          int u = stack.back();
          stack.pop_back();
          while (cur_edge[u] < (int) g[u].size()) {</pre>
              stack.push_back(u);
              u = g[u][cur_edge[u]++];
          vertices.push_back(u);
      reverse(vertices.begin(), vertices.end());
      return vertices;
 8.6.2 Undirected graph
* Hierholzer's algorithm.
* Description: An Eulerian path in a undirected graph is a path that
      visits all edges exactly once.
   * An Eulerian cycle is a Eulerian path that is a cycle.
5 * Time complexity: O(|E|).
6 */
7 struct Edge {
      int to;
      list<Edge>::iterator reverse_edge;
      Edge(int _to) : to(_to) {}
11 };
```

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```
vector<int> vertices;
void find_path(vector<list<Edge>> &g, int u) {
      while (!g[u].empty()) {
          int v = g[u].front().to;
          g[v].erase(g[u].front().reverse_edge);
16
17
          g[u].pop_front();
          find_path(g, v);
18
19
      vertices.emplace_back(u); // reversion list.
20
21 }
void add_edge(int u, int v) {
      g[u].emplace_front(v);
      g[v].emplace_front(u);
24
      g[u].front().reverse_edge = g[v].begin();
25
26
      g[v].front().reverse_edge = g[u].begin();
27 }
```

9 Misc.

9.1 Ternary search

```
const double eps = 1e-9;
2 double ternary_search_max(double 1, double r) {
      // find x0 such that: f(x0) > f(x), \all x: 1 <= x <= r.
      while (r - 1 > eps) {
          double mid1 = 1 + (r - 1) / 3;
          double mid2 = r - (r - 1) / 3;
          if (f(mid1) < f(mid2)) l = mid1;
          else r = mid2;
      return 1;
10
11 }
double ternary_search_min(double 1, double r) {
      // find x0 such that: f(x0) < f(x), \all x: 1 <= x <= r.
      while (r - 1 > eps) {
14
          double mid1 = 1 + (r - 1) / 3;
15
16
          double mid2 = r - (r - 1) / 3;
          if (f(mid1) > f(mid2)) 1 = mid1;
          else r = mid2;
18
      }
19
20
      return 1:
21 }
```

9.2 Dutch flag national problem

```
while (j <= k) {
          if (arr[j] < pivot) {</pre>
              swap(arr[i], arr[j]);
              i++:
              j++;
          else if (arr[j] > pivot) {
              swap(arr[j], arr[k]);
              k--;
          }
          else {
              j++;
          }
      // 0 <= index <= i - 1: arr[index] < mid.
      // i <= index <= k: arr[index] = mid.</pre>
      // k + 1 <= index < sz: arr[index] > mid.
26 }
9.3 Matrix
1 struct Matrix {
      static const matrix_type INF = numeric_limits<matrix_type>::max();
      int N, M;
      vector<vector<matrix_type>> mat;
      Matrix(int _N, int _M, matrix_type v = 0) : N(_N), M(_M) {
          mat.assign(N, vector<matrix_type>(M, v));
      static Matrix identity(int n) { // return identity matrix.
          Matrix I(n, n);
          for (int i = 0; i < n; ++i) {
              I[i][i] = 1;
          return I;
      }
      vector<matrix_type>& operator[](int r) { return mat[r]; }
      const vector<matrix_type>& operator[](int r) const { return mat[r]; }
      Matrix& operator*=(const Matrix &other) {
          assert(M == other.N); // [N x M] [other.N x other.M]
          Matrix res(N, other.M);
          for (int r = 0; r < N; ++r) {
              for (int c = 0; c < other.M; ++c) {</pre>
                  long long square_mod = (long long) MOD * MOD;
                  long long sum = 0;
                  for (int g = 0; g < M; ++g) {
                       sum += (long long) mat[r][g] * other[g][c];
                      if (sum >= square_mod) sum -= square_mod;
```

res[r][c] = sum % MOD;

}

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