



Core C++ 2024.f

# Messing with Floating Point

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**“Never mess with floating point”**

# Agenda

**1.f** Floating Point Basics

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**2.f** Where it Breaks Down

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**3.f** IEEE 754 Standard

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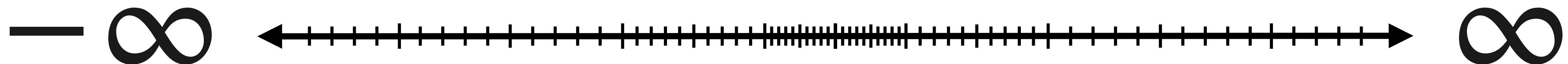
**4.f** Bonus Content

# 1.f Floating Point Basics

# Design Constraints

## Key Considerations

- Suitable for computations in  $\mathbb{R}$



- Practically implementable on hardware
- Consistent across platforms

# Scientific Notation

$$\pm m \times b^e \quad \left\{ \begin{array}{l} 1 \leq m < b \\ e \in \mathbb{Z} \end{array} \right.$$

$$m = x_0 \cdot x_1 x_2 x_3 \dots x_{p-1}$$

$\overbrace{\hspace{10em}}$   
*p digits*

$$b = 10 \quad p = 3$$

$$3.14 \times 10^0 \approx \pi$$

$$1.99 \times 10^{30} \approx \begin{array}{c} \text{orange sun icon} \\ \text{on a scale} \end{array}$$

$$8.40 \times 10^{-16} \approx \begin{array}{c} \text{red circle with white plus sign} \\ \text{with a horizontal line below it} \end{array}$$

$$3.14 \times 10^0$$
$$1.99 \times 10^{30}$$
$$8.40 \times 10^{-16}$$

-  Suitable for computations in  $\mathbb{R}$
-  Practical to store and manipulate
-  Consistent and reproducible

# Float Representation

## Binary Scientific Notation

S

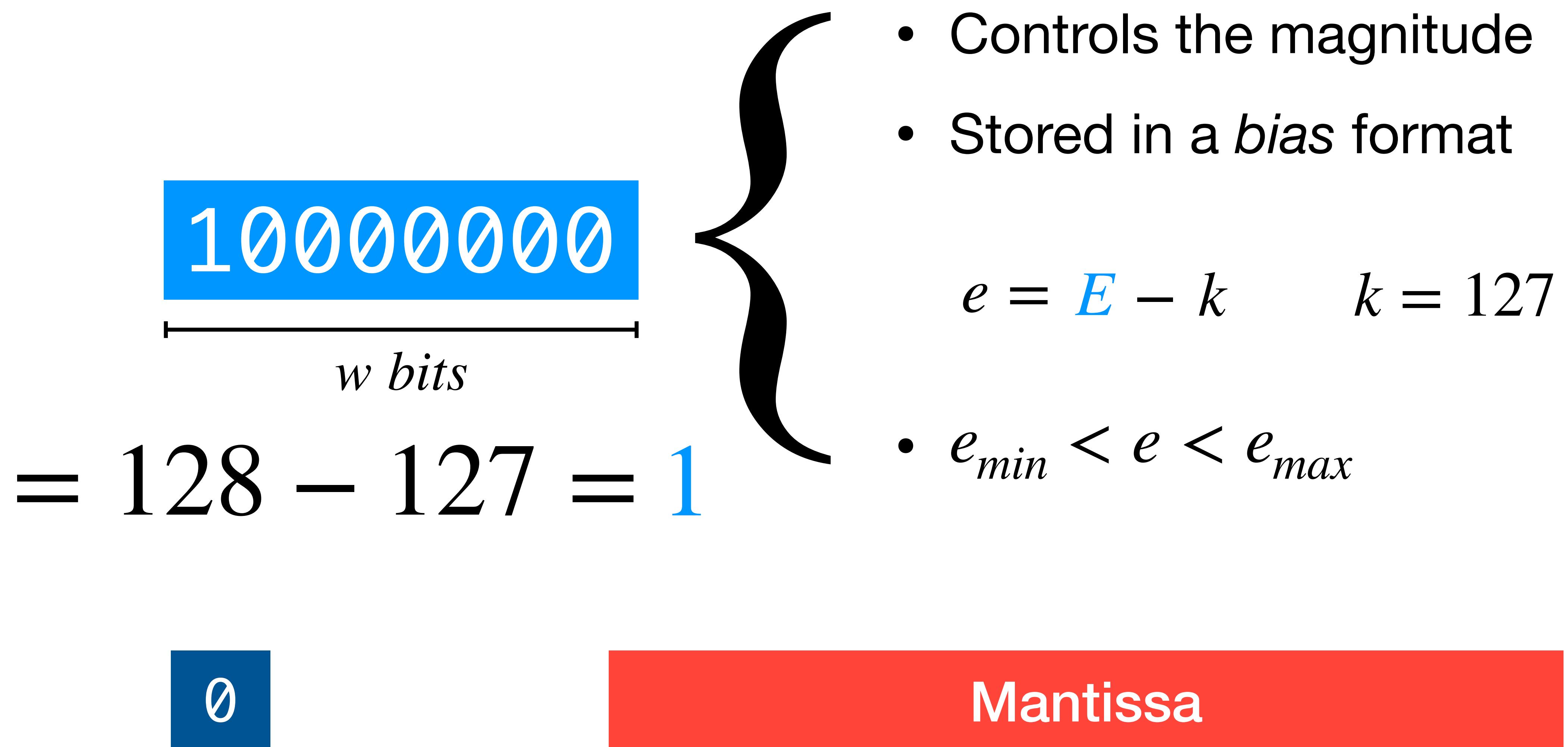
Exponent

Mantissa

- 
- + is 0, - is 1
  - Solely responsible for the sign
  - 0.0 and -0.0 are distinct

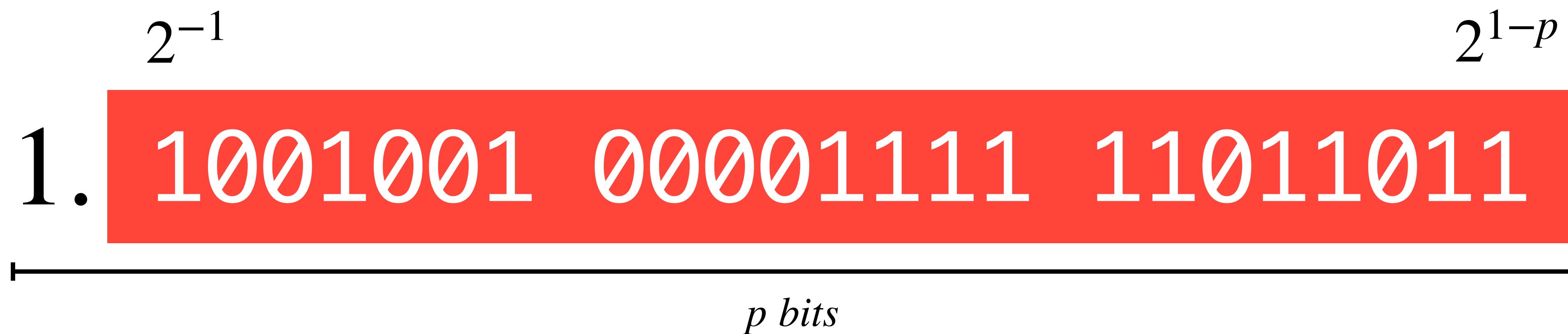
Exponent

Mantissa



$$1 \leq m < 2$$

$$\textcolor{red}{m} = 1 + 2^{-1} + 2^{-4} + 2^{-7} + \dots \approx 1.571$$



0	10000000
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$$+ 1.571 \times 2^1 \approx \pi$$

0 10000000 1001001 00001111 11011011

# Float Format Parameters

## IEEE 754 Specification

	16-bit	32-bit	64-bit	128-bit
$p$	11	24	53	113
$w$	5	8	11	15
$k$ (bias)	15	127	1023	16383



# 1.f Floating Point Basics



Design Constraints

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Scientific Notation

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Float Representation

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Float Format Parameters

# Agenda



## Floating Point Basics

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**2.f** Where it Breaks Down

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**3.f** IEEE 754 Standard

---

**4.f** Bonus Content

## 2.f Where it Breaks Down

```
float f = 0.1;  
assert(f == 0.1);
```

## 2.1f Representation Error

# Representation Error

$$\frac{1}{3} \rightarrow 3.33 \times 10^{-1}$$

333333...

$$\frac{1}{10} \rightarrow 1.100110... \times 2^{-4}$$

011001100...

**“Squeezing infinitely many real numbers  
into a finite number of bits requires an  
*approximate* representation.”**

**David Goldberg**

```
float f = 0.1;  
assert(f == 0.1);
```

32-bit

64-bit

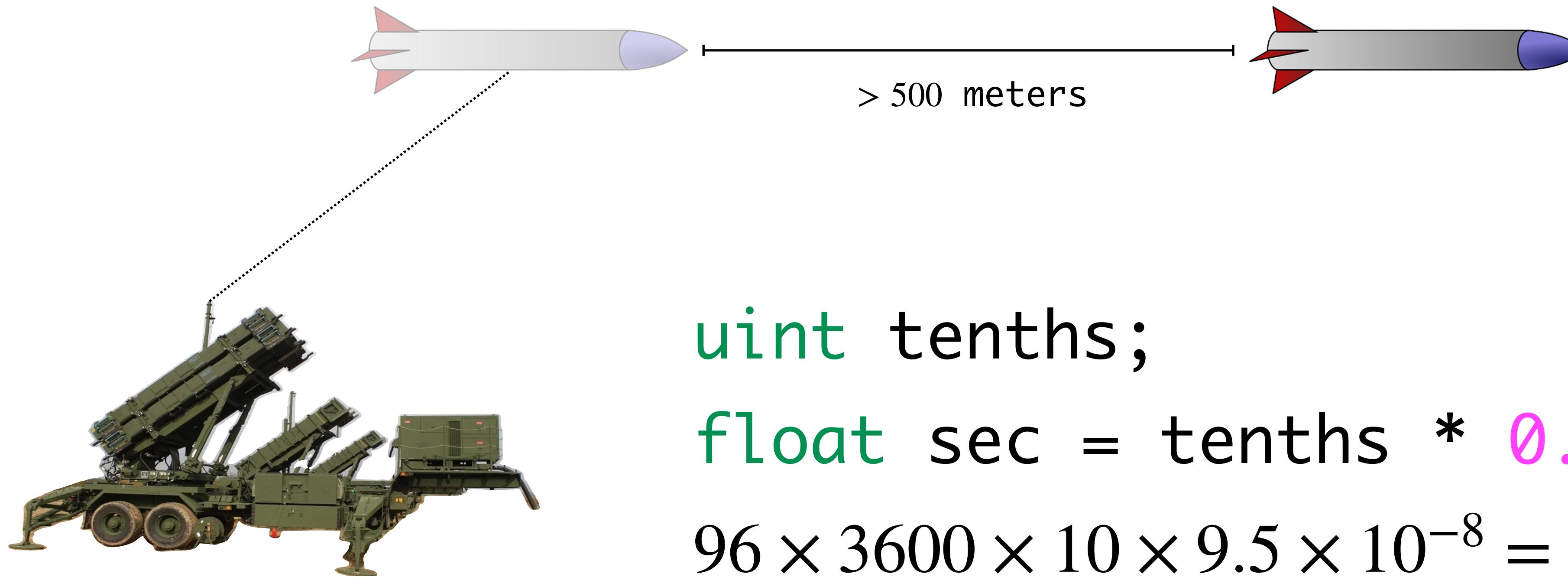
```
float f = 0.1;  
cout << setprecision(INT_MAX) << f << endl;
```

>> 0.10000001490116119384765625

`0.100000000000000000000000000000`

`0.10000001490116119384765625`

$\approx 1.5 \times 10^{-9}$



`uint tenths;`

`float sec = tenths * 0.1;`

$$96 \times 3600 \times 10 \times 9.5 \times 10^{-8} = 0.33$$

# Quantifying Error

## ULPs

- Units in the Last Place
- Variance in terms of least-significant digit
- Useful for representation error

## Relative Error

- Same as percent error

$$E = \frac{|f - r|}{r}$$

- Useful for post-calculation analysis

# Quantifying Error

$$\pi \rightarrow 3.14$$

$$\frac{\pi - 3.14}{\pi} = 0.0005$$

3.14159265... → 0.159 ULPs

# Mitigating Representation Error

- Use higher precision types
- Prefer exact values

$$\pm m \times 2^e \quad \longleftrightarrow \quad \text{A horizontal line with tick marks, representing the range of representable numbers.}$$

All integers up to  $2^p$  are exact

- Use stable algorithms

$$x \times 5 \iff \frac{x}{0.2}$$

$$x \times \frac{2}{3} \iff \frac{x}{1.5}$$

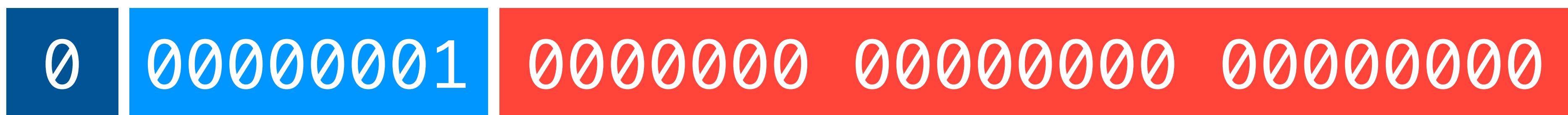
$$x \times 0.5 \iff \frac{x}{2}$$

```
float f = numeric_limits<float>::min();
// f = 1.17549e-38
cout << (f / 2.f) << endl;
```

>> 5.87747e-39

## 2.2f Denormal Numbers

$1.17549e-38$



$5.87747e-39$



**00000000** → Implicit bit = 0

# Denormal Numbers

$$\neg(1 \leq m < 2)$$

- $e = e_{min} = 00000000$
- Help prevent underflow



- Lower precision
- Poor performance

## 2.3f Float Comparison

```
float f = 0.1;  
assert(f == 0.1);
```

# Comparison

- **Bitwise:**  $a == b$  iff  $a$  and  $b$  have the same bit representation
- **IEEE 754:** Bitwise (  $\text{NaN} \neq x$      $-\infty < x < \infty$      $0 = -0$  )
- **Epsilon:**  $a == b$  iff  $|a - b| < \epsilon$
- **Relative:**  $a == b$  iff  $\frac{|a - b|}{a} < \epsilon$
- **ULPs:**  $a == b$  iff  $\text{ulp\_distance}(a, b) < \epsilon$

```
std::nextafter(float from, float to)  
std::numeric_limits<float>::epsilon()
```

```
int ulp_distance(float a, float b)  
{  
    uint32_t inta = *reinterpret_cast<uint32_t*>(&a);  
    uint32_t intb = *reinterpret_cast<uint32_t*>(&b);  
    return intb - inta;  
}
```

```
float sum = 0;  
  
for (int i = 0; i < 1'000'000'000; ++i)  
    sum += 1.f / 1'000'000'000.f;  
  
cout << sum << endl;
```

>> 0.03125

## 2.4f Resolution Error

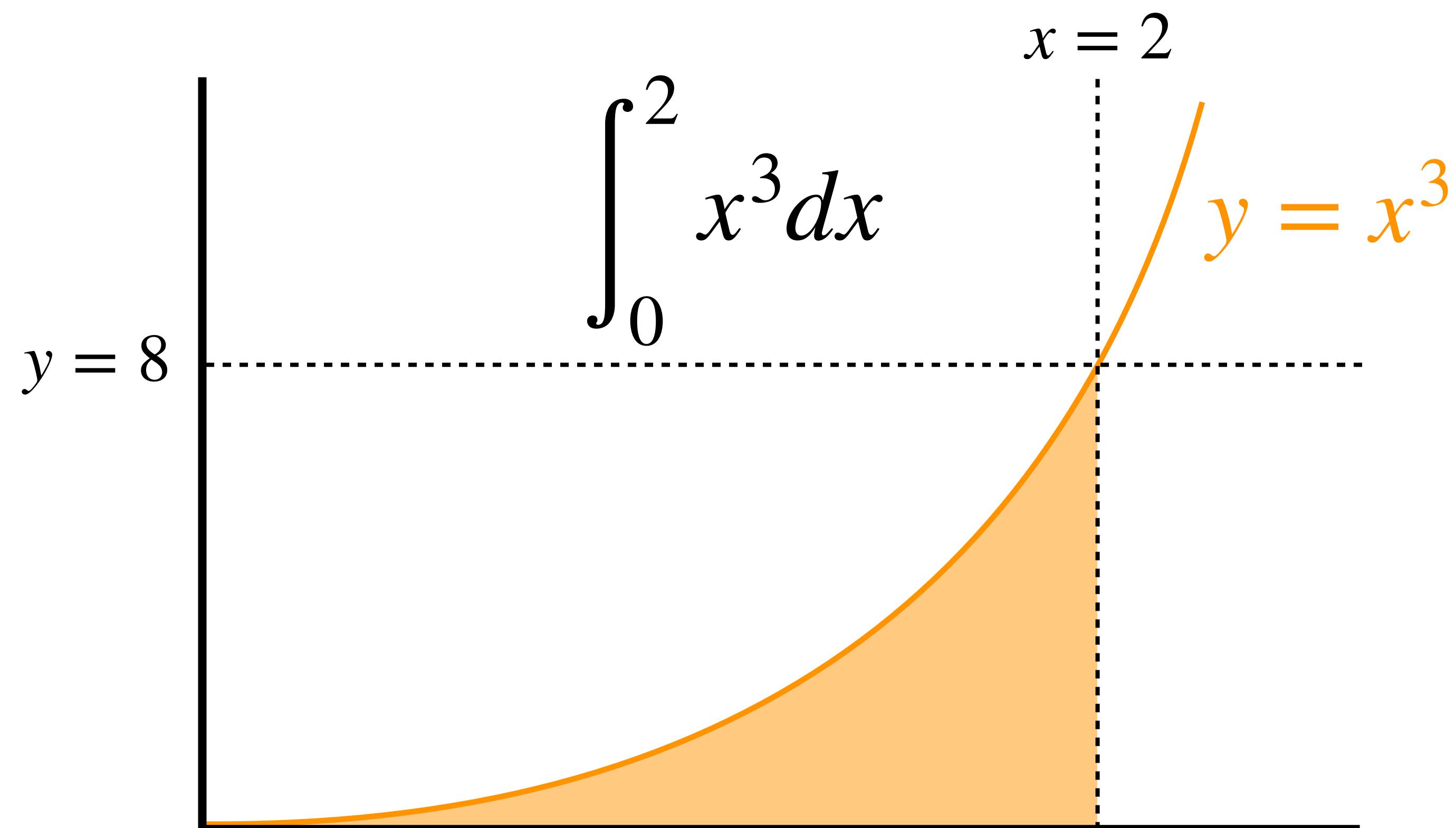
# Resolution Error

- Occurs when adding two very mismatched numbers

$$100 + 0.01 \rightarrow \begin{array}{r} 1.00 \times 10^2 \\ + 1.00 \times 10^{-2} \\ \hline \end{array} \rightarrow \frac{1.00 \times 10^2}{1.00 \times 10^2} + 0.00 \times 10^2 \rightarrow 100$$

- Group computations by magnitude
- Use more sophisticated algorithms

$$= 4$$



```
// integrate x**3 from 0 -> 2
float sum = 0;
const float dx = 0x1.0p-22;

for (float x = 0; x <= 2; x += dx)
    sum += x * x * x * dx;

cout << sum << endl;
```

>> 4.00028

```
// integrate x**3 from 0 -> 2
float sum = 0;
const float dx = 0x1.0p-22;

for (float x = 2; x >= 0; x -= dx)
    sum += x * x * x * dx;

cout << sum << endl;
```

>> 3.94732

```
float kahan_sum(const vector<float>& addends)
{
    float sum = 0; // accumulator
    float comp = 0; // running compensation for lost bits
    for (float f: addends)
    {
        float y = f - comp; // comp = 0 on first iter
        float t = sum + y; // sum > y... y gets truncated
        comp = (t - sum) - y; // comp recovers lost part of y
        sum = t; // error gets recovered next iter
    }
    return sum;
}
```

```
vector<float> nums(100'000'000);

random_device rd;
mt19937 gen(rd());
uniform_real_distribution<float> dist(0.f, 1.f);

generate(nums.begin(), nums.end(), [&] () { return dist(gen); });

cout << int(kahan_sum(nums)) << endl;
cout << int(naive_sum(nums)) << endl;
```

```
>> 49997528
>> 16777216
```



- 131072: diagonal walking is bumpy
- 262144: strings become invisible
- 2097152: blocks render in 2D
- 4194304: walking is impossible
- 8388608: fall through the floor
- 1677216: every 2<sup>nd</sup> block renders

## 2.f Where it Breaks Down



Representation Error

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Denormalization

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Float Comparison

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Resolution Error

# Agenda



Floating Point Basics

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Where it Breaks Down

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**3.f** IEEE 754 Standard

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**4.f** Bonus Content

## 3.f IEEE 754 Standard

# IEEE 754 Standard

## Introduction

- Developed in 1985, updated in 2019
- IEEE 754:2019 == ISO/IEC 60559:2020
- Representation, exceptional values, exceptions, and operations



- **Independent from C++**

# C++ Standard

- Does not mandate IEEE 754 compliance
- C++23 has very little to say about floating point
- C++26 / C++29 might have more

# C++26 / C++29

## P3375R0

- Compiler flags: -fno-fast-math
- Specification overhaul: define behavior for `float`
- New type: `std::IEEE_float32_t`
- Qualification:

`float func(float a, float b) reproducible`

# Floating Point Environment

- To check IEEE 754 compliance:

```
std::numeric_limits<float>::is_iec559
```

- To access floating point environment:

```
#include <cfenv>
```

```
#pragma STDC FENV_ACCESS ON
```

# 3.1f Exceptional Values

# Exceptional Values

## Zero

$0 =$	0	00000000	0000000 0000000 0000000
$-0 =$	1	00000000	0000000 0000000 0000000

“Comparisons shall ignore the sign of zero (so  $+0 = -0$ )”

$$x \times 0 = 0$$

$$x \times -0 = -0$$

$$0 \times -0 = -0$$

# Exceptional Values

## Infinity

$\infty =$	0	11111111	0000000 0000000 0000000
$-\infty =$	1	11111111	0000000 0000000 0000000

$$-\infty < x < \infty$$

$$\frac{x}{0} = \infty$$

$$\frac{x}{\infty} = 0$$

$$\infty + x = \infty$$

# Exceptional Values

**NaN**

$qNaN = \begin{array}{c} 0 \\ | \\ 11111111 \\ | \\ 1000000 \ 0000000 \ 0000000 \end{array}$

$sNaN = \begin{array}{c} 0 \\ | \\ 11111111 \\ | \\ 0????? \ ?????? \ ?????? \end{array}$

$$\frac{0}{0} = \text{NaN}$$

$$\text{NaN} + x = \text{NaN}$$

$$\infty - \infty = \text{NaN}$$

$$\text{NaN} \neq x \neq \text{NaN}$$

```
bool std::isnan(float f)
```

```
bool std::isinf(float f)
```

```
float std::numeric_limits<float>::quiet_NaN()
```

```
float std::numeric_limits<float>::signaling_NaN()
```

```
float std::numeric_limits<float>::infinity()
```

## 3.2f Exceptions

# Exceptions

## Invalid Operation

FE\_INVALID

“The *invalid operation* exception is signaled if and only if there is no usefully definable result.”

$$\frac{0}{0}$$

$$\frac{\infty}{\infty}$$

$$0 \times \infty$$

$$\sqrt{-x}$$

# Exceptions

## Division by Zero

FE\_DIVBYZERO

“The *divideByZero* exception shall be signaled if and only if an exact infinite result is defined for an operation on finite operands.”

$$\frac{x}{0} \quad \log(0)$$

# Exceptions

## Overflow

FE\_OVERFLOW

“The *overflow* exception shall be signaled if the destination format’s largest finite number is exceeded in magnitude...”

$$\text{FLT\_MAX} \times 2 = \begin{cases} \text{FLT\_MAX} \\ \infty \end{cases}$$

# Exceptions

## Underflow

FE\_UNDERFLOW

“The *underflow* exception shall be signaled when a tiny non-zero result is detected”

$$|x| < 2^{e_{min}}$$

# Exceptions

## Inexact

FE\_INEXACT

“Except as specified otherwise, an operation delivering a numerical result that signals no other exception shall signal inexact if its rounded result differs from what would have been computed were both the exponent range and precision unbounded.”

0.1

$\frac{1}{3}$

$\sqrt{2}$

```
int std::fetestexcept(int e)
```

```
int std::feclearexcept(int e)
```

```
int std::feraiseexcept(int e)
```

```
void show_fe_exceptions()
{
    cout << "exceptions raised:";

    if (fetestexcept(FE_INVALID))      cout << " invalid";
    if (fetestexcept(FE_DIVBYZERO))    cout << " divbyzero";
    if (fetestexcept(FE_OVERFLOW))     cout << " overflow";
    if (fetestexcept(FE_UNDERFLOW))    cout << " underflow";
    if (fetestexcept(FE_INEXACT))      cout << " inexact";

    feclearexcept(FE_ALL_EXCEPT);
    cout << endl;
}
```

```
show_fe_exceptions();

float f = 0.1; // 0.1 is inexact
show_fe_exceptions();

f /= 0; divide 0 = infinity
show_fe_exceptions();

f *= 0; // infinity * 0
show_fe_exceptions();
```

```
>> exceptions raised:
>> exceptions raised: inexact
>> exceptions raised: divbyzero
>> exceptions raised: invalid
```

## 3.3f Rounding Modes

# Rounding Modes

- Round to 0:  $x \rightarrow 0$  `FE_TOWARDZERO`
- Round up:  $x \rightarrow \infty$  `FE_UPWARD`
- Round down:  $x \rightarrow -\infty$  `FE_DOWNWARD`
- Round to nearest: `FE_TONEAREST`

$3.142 \rightarrow 3.14$

$2.718 \rightarrow 2.72$

$1.005 \rightarrow 1.00$

$1.015 \rightarrow 1.02$

```
int std::fegetround();
```

```
int std::fesetround(int round);
```

## 3.4f Arithmetic Guarantees

# Basic Operations

+ - \* / sqrt fma

- Deterministic and reproducible
- Exactly rounded (correct)
- $a + b = b + a$     $a \times b = b \times a$
- $(a + b) + c \neq a + (b + c)$     $(a \times b) \times c \neq a \times (b \times c)$
- Handling of special values

```
float std::fma(float x, float y, float z)
```

# Conversion Operations

float->int    int->float    float->float    float->string

- Exactly rounded
- For conversion to integers:
  - `int(float x)` always rounds to zero
  - `std::rint(float x)` obeys the current rounding mode
- Parse decimal strings
- Produce a unique string representation (*9 digits, 17 digits*)

## 3.f IEEE 754 Standard



Exceptional Values

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Exceptions

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Rounding Modes

---



Arithmetic Guarantees

# Agenda



Floating Point Basics

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Where it Breaks Down

---



IEEE 754 Standard

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**4.f** Bonus Content

## 4.f Bonus Content

# 4.1f Ryu Float-to-String

```
float f = 0.1;  
cout << setprecision(INT_MAX) << f << endl;
```

>> 0.10000001490116119384765625

0 01111011 1001100 11001100 11001101

$$e := \boxed{0111011} - 123 - 127 = -4$$

$$m := \boxed{1001100 \ 11001100 \ 11001101} = 1.6$$

$$e = -4$$

$$m = 1.6$$

$$1.6 \times 2^{-4} = 0.1$$

$$e = -4$$

$$m = 1.6 \ll 23$$

$$1.6 \times 2^{-4} = 0.1$$

$$e = -4$$

$$m = 1.6 \ll 2^3 = 13,421,773$$

$$1.6 \times 2^{-4} = 0.1$$

$$e = -4 - 23$$

$$m = 1.6 \ll 23 = 13,421,773$$

$$1.6 \times 2^{-4} = 0.1$$

$$e = -4 - 23 = -27$$

$$m = 1.6 \times 2^{23} = 13,421,773$$

$$1.6 \times 2^{-4} = 0.1$$

$$e = -4 - 23 = -27$$

$$m = 1.6 \ll 23 = 13,421,773$$

$$13,421,773 \times 2^{-27} = 0.1$$

e = -27

m = 13,421,773

$$e = -27$$

$$m = 13,421,773 * 5^{**27} * 5^{**-27}$$

$$e = -27$$

$$m = 13,421,773 * 5^{27} * 5^{-27}$$

$$13,421,773 * 2^{-27}$$

$$e = -27$$

$$m = 13,421,773 * 5^{27} * 5^{-27}$$

$$13,421,773 * 5^{27} * 5^{-27} * 2^{-27}$$

$$e = -27$$

$$m = 13,421,773 * 5^{27} * 5^{-27}$$

$$13,421,773 * 5^{27} * 10^{-27}$$

e = -27

m = 13,421,773 \* 5\*\*27 \* 5\*\*-27

10000001490116119384765625 \* 10\*\*-27

**0.10000001490116119384765625**

```
float f = 0.1;  
cout << setprecision(INT_MAX) << f << endl;
```

>> 0.10000001490116119384765625

# 4.2f Fast Inverse Square Root

```
float Q_rsqrt(float number)
{
    float x = number * 0.5f;
    float y = number;

    int i = * (int*) &y;          // evil bit hack
    i = 0x5f3759df - (i >> 1); // what the f?
    y = * (float*) &i;
    y *= (1.5f - (x * y * y)); // 1st iteration
    y *= (1.5f - (x * y * y)); // 2nd iteration

    return y;
}
```

## 4.3f Hexadecimal Float Literals

**0x<mantissa>p<exponent>**

# Thank you!

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