# PROJECT 2 Priority Scheduling



Real time systems Course 2022/2023

# Participants:

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# PART 1: 1-bit samples

# Software module

#### Scheduler

The tasks are not independent because there is a shared variable. Therefore, we will use the priority inheritance method. All times are in milliseconds:

|                         | T=D  | С | Р   | LOCK |
|-------------------------|------|---|-----|------|
| T1: PLAYBACK TASK       | 512  | 7 | 3 + | 1    |
| T2: RECEIVE COMMANDS    | 1000 | 1 | 2   | 1    |
| T3: SHOW PLAYBACK STATE | 2500 | 4 | 1 - | 1    |

CPU usage: 
$$U = 7/512 + 1/1000 + 4/2500 = 0.0163 < 1$$

$$B(T1) = 1 + 1 = 2$$
  $B(T2) = 1$   $B(T3) = 0$ 

$$R(T1) = 7 + 2 + 0 = 9 \le 512 = D(T1)$$

$$R(T2) = 1 + 1 + 7 * \lceil R(T2)/512 \rceil$$
  
 $W0 = 1 + 1 + 7 * \lceil 1/512 \rceil = 9$   
 $W1 = 1 + 1 + 7 * \lceil 9/512 \rceil = 9$ 

$$R(T2) = 9 \le 1000 = D(T2)$$

$$R(T3) = 4 + 0 + 7 * [R(T3)/512] + 1 * [R(T3)/1000]$$
  
 $W0 = 4 + 7 * [1/512] + 1 * [1/1000] = 12$   
 $W1 = 4 + 7 * [12/512] + 1 * [12/1000] = 12$   
 $R(T3) = 12 \le 2500 = D(T3)$ 

# Hardware module

### Scheduler

The tasks are not independent because there is a shared variable. Therefore, we will use the priority inheritance method. All times are in microseconds:

|                             | T=D  | С  | Р   | LOCK |
|-----------------------------|------|----|-----|------|
| T1: PLAYBACK                | 250  | 28 | 2 + | 28   |
| T23 (2 & 3): BUTTON AND LED | 1000 | 24 | 1 - | 8    |

CPU usage: 
$$U = 28/250 + 24/1000 = 0.136 < 1$$

$$B(T1) = 8$$
  $B(T23) = 0$ 

$$R(T1) = 28 + 8 + 0 = 36 \le 250 = D(T1)$$

$$R(T23) = 24 + 0 + 28 * [R(T23)/250]$$
  
 $W0 = 24 + 0 + 28 * [1/250] = 52$   
 $W1 = 24 + 0 + 28 * [52/250] = 52$   
 $R(T23) = 52 \le 1000 = D(T23)$ 

# PART 2: 8-bit samples

# Software module

#### Scheduler

The tasks are not independent because there is a shared variable. Therefore, we will use the priority inheritance method. All times are in milliseconds:

|                         | T=D  | С | Р   | LOCK |
|-------------------------|------|---|-----|------|
| T1: PLAYBACK TASK       | 64   | 7 | 3 + | 1    |
| T2: RECEIVE COMMANDS    | 1000 | 1 | 2   | 1    |
| T3: SHOW PLAYBACK STATE | 2500 | 4 | 1 - | 1    |

CPU usage: 
$$U = 7/64 + 1/1000 + 4/2500 = 0.112 < 1$$

$$B(T1) = 1 + 1 = 2$$
  $B(T2) = 1$   $B(T3) = 0$ 

$$R(T1) = 7 + 2 + 0 = 9 \le 64 = D(T1)$$

$$R(T2) = 1 + 1 + 7 * \lceil R(T2)/64 \rceil$$
  
 $W0 = 1 + 1 + 7 * \lceil 1/64 \rceil = 9$   
 $W1 = 1 + 1 + 7 * \lceil 9/64 \rceil = 9$ 

$$R(T2) = 9 \le 1000 = D(T2)$$

$$R(T3) = 4 + 0 + 7 * \lceil R(T3)/64 \rceil + 1 * \lceil R(T3)/1000 \rceil$$
  
 $W0 = 4 + 7 * \lceil 1/64 \rceil + 1 * \lceil 1/1000 \rceil = 12$   
 $W1 = 4 + 7 * \lceil 12/64 \rceil + 1 * \lceil 12/1000 \rceil = 12$   
 $R(T3) = 12 \le 2500 = D(T3)$ 

## Hardware module

#### Scheduler

The tasks are not independent because there is a shared variable. Therefore, we will use the priority inheritance method. All times are in microseconds:

|                             | T=D  | С  | Р   | LOCK |
|-----------------------------|------|----|-----|------|
| T1: PLAYBACK                | 250  | 28 | 2 + | 28   |
| T23 (2 & 3): BUTTON AND LED | 1000 | 24 | 1 - | 8    |

CPU usage: 
$$U = 28/250 + 24/1000 = 0.136 < 1$$
 $B(T1) = 8$ 
 $B(T23) = 0$ 
 $R(T1) = 28 + 8 + 0 = 36 \le 250 = D(T1)$ 
 $R(T23) = 24 + 0 + 28 * \lceil R(T23)/250 \rceil$ 
 $W0 = 24 + 0 + 28 * \lceil 1/250 \rceil = 52$ 
 $W1 = 24 + 0 + 28 * \lceil 52/250 \rceil = 52$ 
 $R(T23) = 52 \le 1000 = D(T23)$ 

## Remarks

All compute times and lock times were estimated running the actual programs. Real compute times and lock times may be smaller than those used for priority scheduler design. Since these times depend on the hardware used, we should note that the hardware module was run on an Arduino Uno microcontroller, and the software module was run on QEMU on top of a Debian linux bare metal install.

We considered that the lock time of the playback task (task 1) in the hardware module for both parts was the same as the total compute time, since this task is essentially fully enclosed in a critical section. This is because this task is executed upon an interrupt, which can never be interrupted by another task, as we do not use other interrupts.

 $OCR1A = Arduino\_frequency/(1/t\_interrupt) - 1 = 16000000/(1/0.00025) - 1 = 3999$