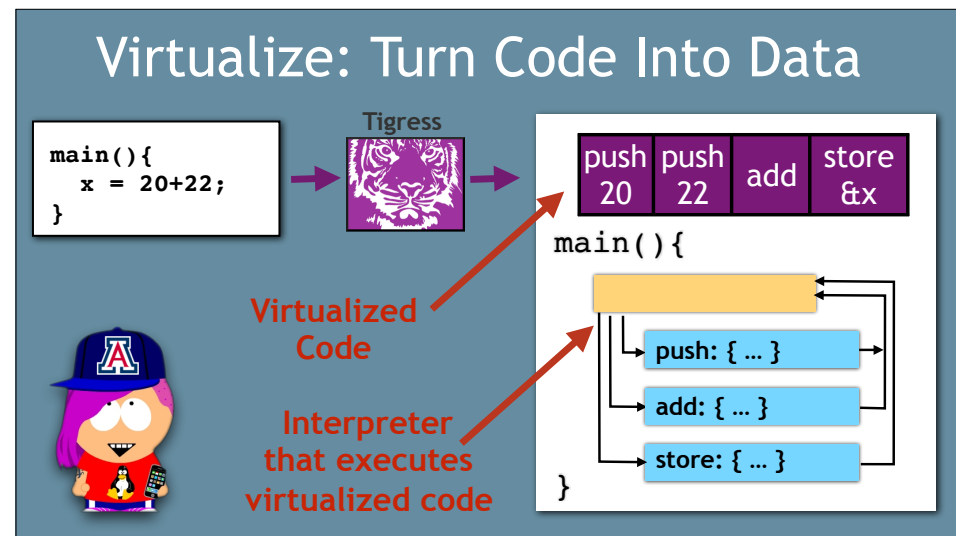
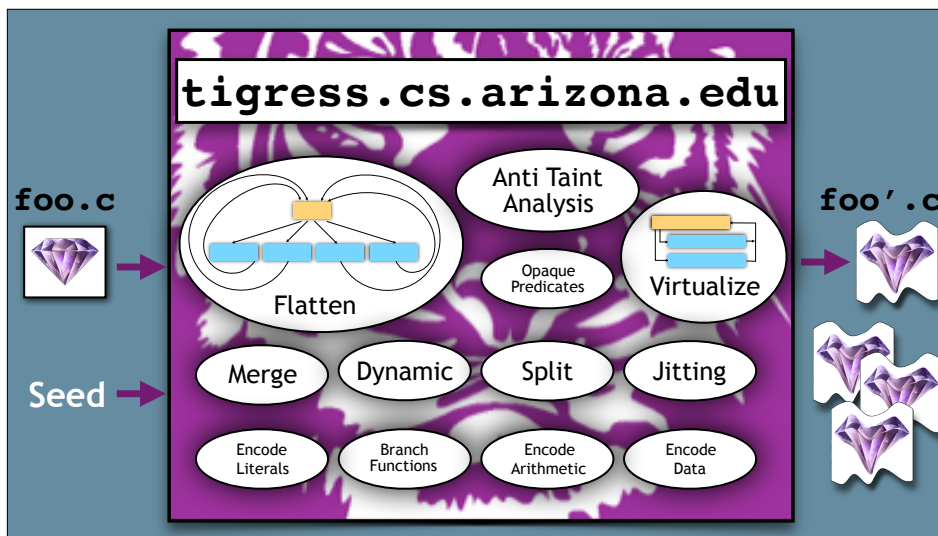


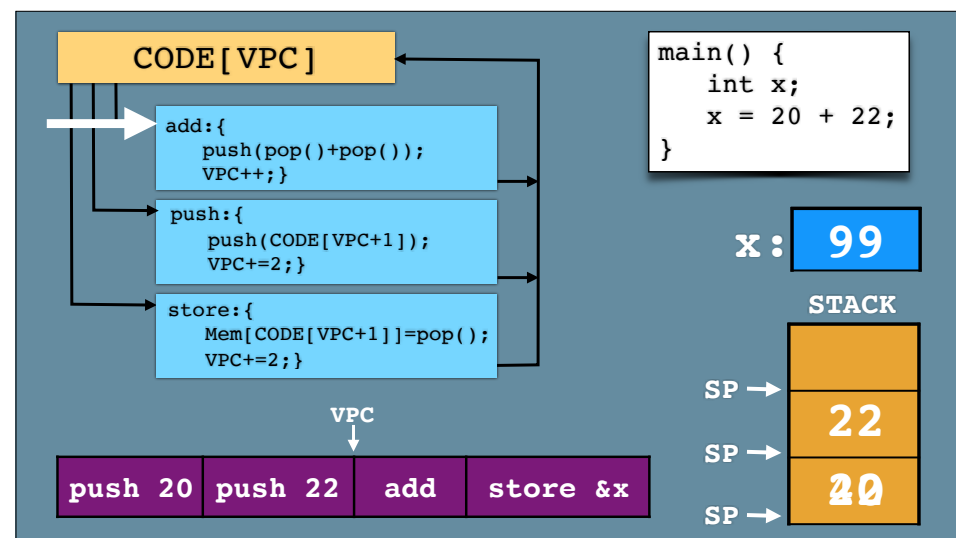
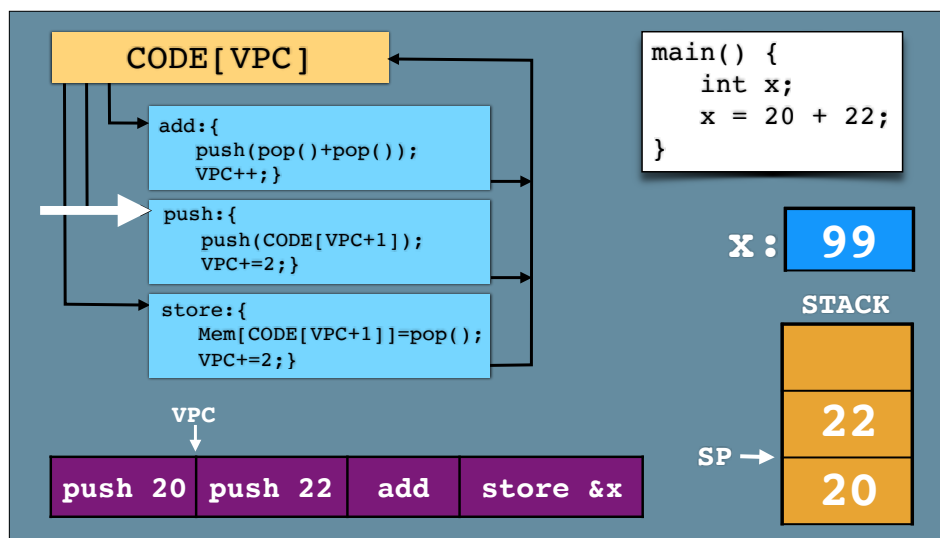
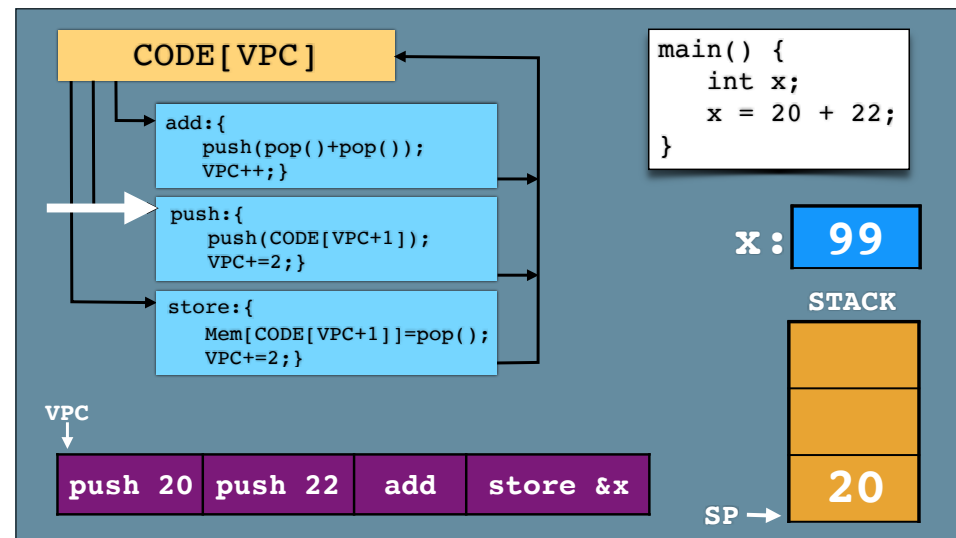
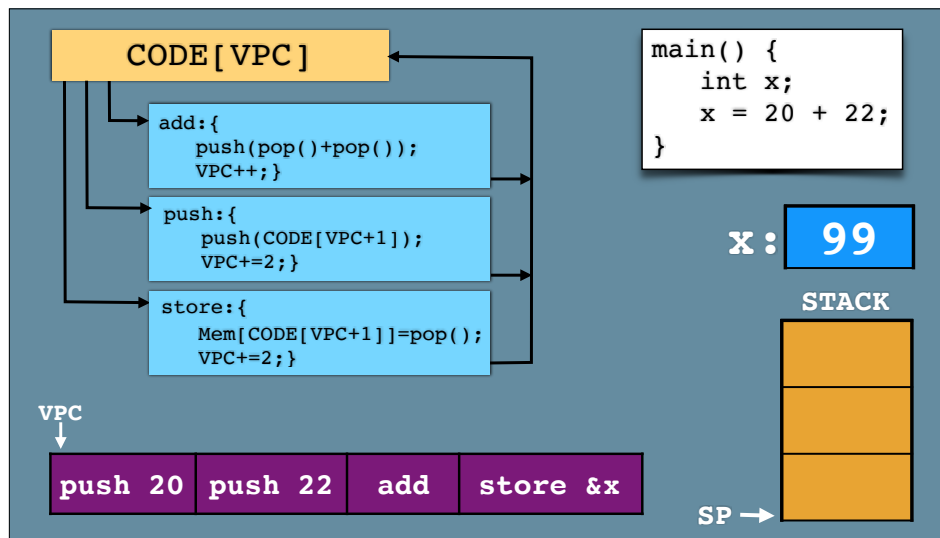
Obfuscation and Tamperproofing

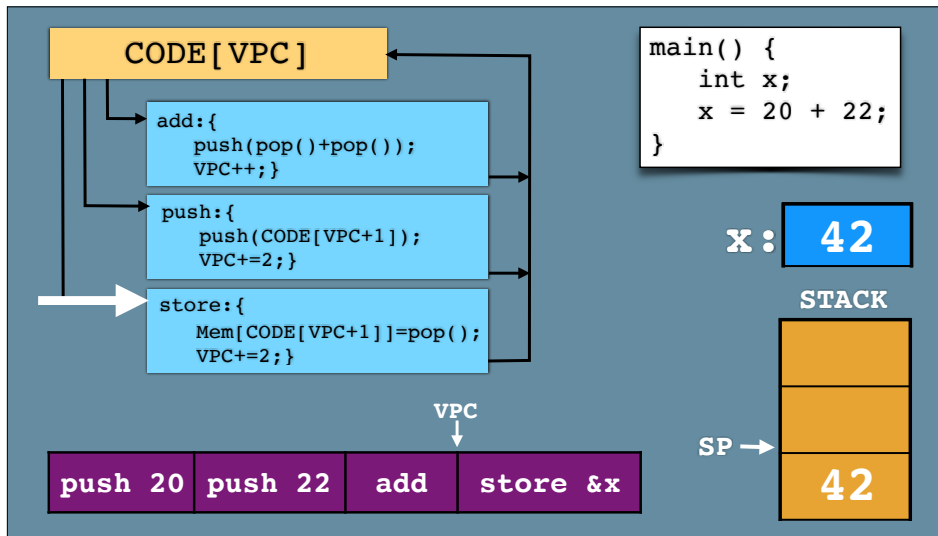
Lecture #4

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Obfuscation by Virtualization







Obfuscating Arithmetic

Encoding Integer Arithmetic

$$x+y = x - \neg y - 1$$

$$x+y = (x \oplus y) + 2 \cdot (x \wedge y)$$

$$x+y = (x \vee y) + (x \wedge y)$$

$$x+y = 2 \cdot (x \vee y) - (x \oplus y)$$

Example

One possible encoding of

$$z = x + y + w$$

is

$$z = (((x \wedge y) + ((x \& y) \ll 1)) \ll w) + (((x \wedge y) + ((x \& y) \ll 1)) \& w);$$

Many others are possible, which is good for diversity.

Exercise!

- The virtualizer's **add** instruction handler could still be identified by the fact that it uses a + operator!
- Try adding an arithmetic transformer:

```
tigress --Environment=x86_64:Linux:Gcc:4.6\
--Transform=Virtualize \
--Functions=fib \
--VirtualizeDispatch=switch\
--Transform=EncodeArithmetic \
--Functions=fib \
--out=fib5.c fib.c
```

- What differences do you notice between before and after arithmetic encoding?

```
int fib(int n ) {
    ...
    while (1) {
        switch (*( _1_fib_$pc[0])) {
            case PlusA: {
                ( _1_fib_$sp[0] + -1)->_int =
                    ( _1_fib_$sp[0] + -1)->_int
                    ( _1_fib_$sp[0] + 0)->_int;
                break;
            }
        }
    }
}
```

+

```
int fib(int n ) {
    ...
    while (1) {
        switch (*( _1_fib_$pc[0])) {
            case PlusA: {
                ( _1_fib_$sp[0] + -1)->_int =
                    (( _1_fib_$sp[0] + -1)->_int
                     ( _1_fib_$sp[0] + 0)->_int)
                    ((( _1_fib_$sp[0] + -1)->_int
                     ( _1_fib_$sp[0] + 0)->_int) << 1);
                break;
            }
        }
    }
}
```

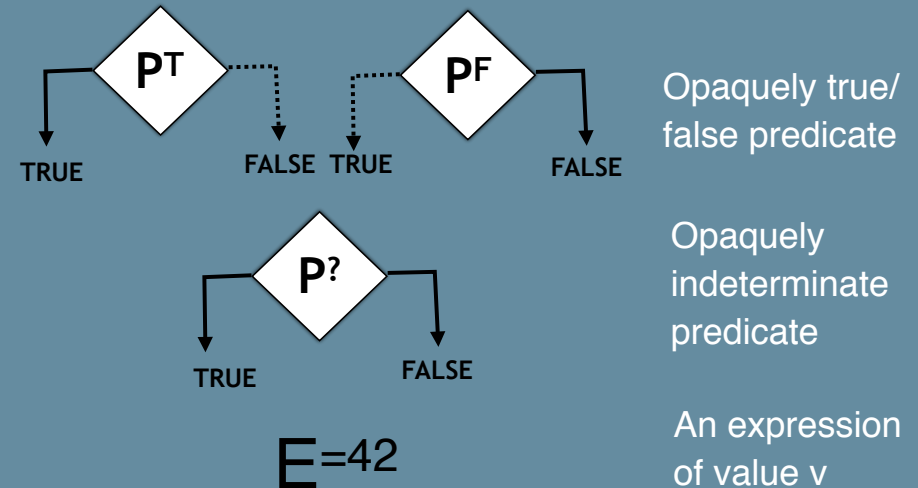
^
+
&

$$x+y = (x \oplus y) + 2 \cdot (x \wedge y)$$

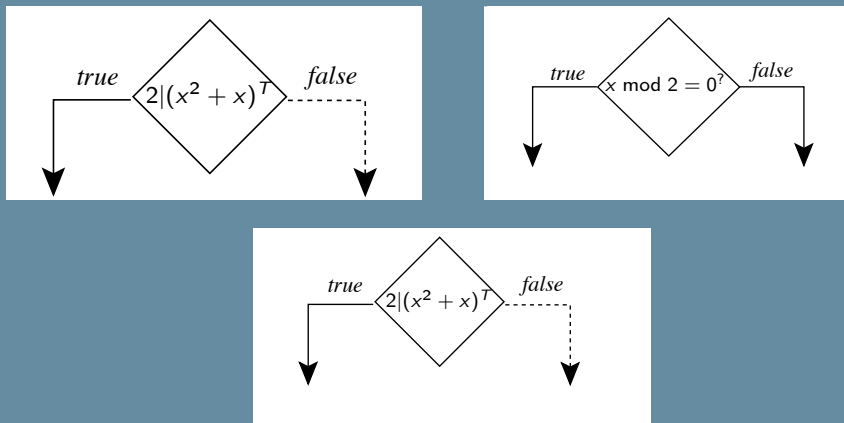
**Opaque
Expressions**

Opaque Expressions

An expression whose value is known to you as the defender (at obfuscation time) but which is difficult for an attacker to figure out



Examples



Inserting Bogus Control Flow

```
if (x[k] == 1)
    R = (s*y) % n
else
    R = s;
    s = R*R % n;
    L = R;
```



```
if (x[k] == E=1)
    R = (s*y) % n
else
    R = s;
    s = R*R % n;
    L = R;
```

Inserting Bogus Control Flow

```
if (x[k] == 1)
    R = (s*y) % n
else
    R = s;
s = R*R % n;
L = R;
```



```
if (x[k] == 1)
    R = (s*y) % n
else
    R = s;
if (expr=T)
    s = R*R % n;
else
    s = R*R * n;
L = R;
```

Inserting Bogus Control Flow

```
if (x[k] == 1)
    R = (s*y) % n
else
    R = s;
s = R*R % n;
L = R;
```

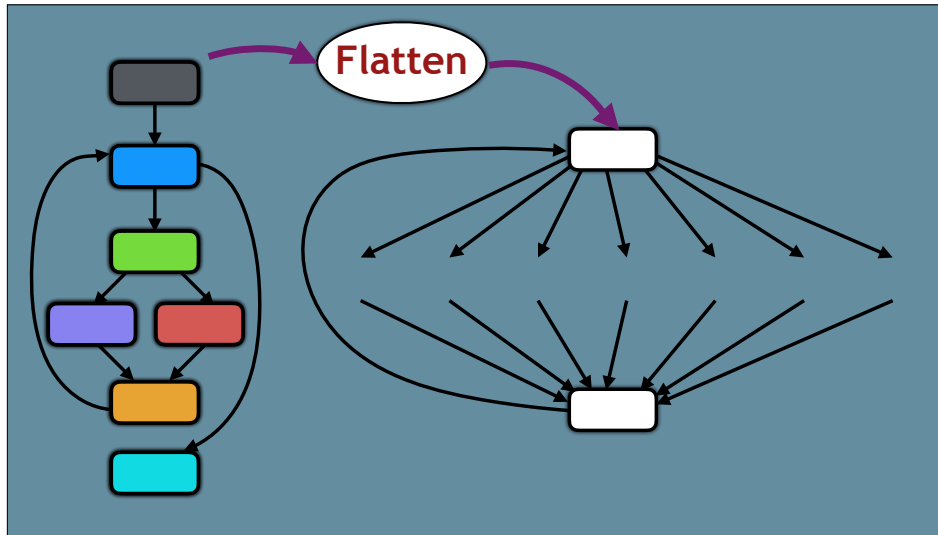


```
if (x[k] == 1)
    R = (s*y) % n
else
    R = s;
if (expr=?)
    s = R*R % n;
else
    s = (R%n)*(R%n)%n;
L = R;
```

Exercise!

```
*****
* 1) Opaque Predicates
*****
tigress --Environment=x86_64:Linux:Gcc:4.6 --Seed=0 \
--Transform=InitEntropy \
--InitEntropyKinds=vars \
--Transform=InitOpaque \
--Functions=main\
--InitOpaqueCount=2\
--InitOpaqueStructs=list,array \
--Transform=AddOpaque\
--Functions=fib\
--AddOpaqueKinds=question \
--AddOpaqueCount=10 \
--out=fib1.c fib.c
```

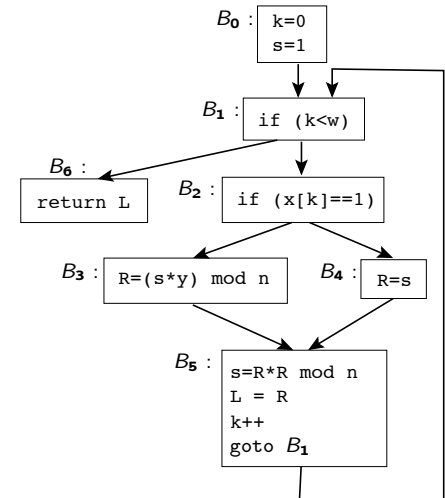
**Control Flow
Flattening**



```

int modexp(
  int y, int x[],
  int w, int n){
  int R, L;
  int k=0; int s=0;
  while (k < w) {
    if (x[k] == 1)
      R = (s*y) % n
    else
      R = s;
    s = R*R % n;
    L = R;
    k++;
  }
  return L;
}

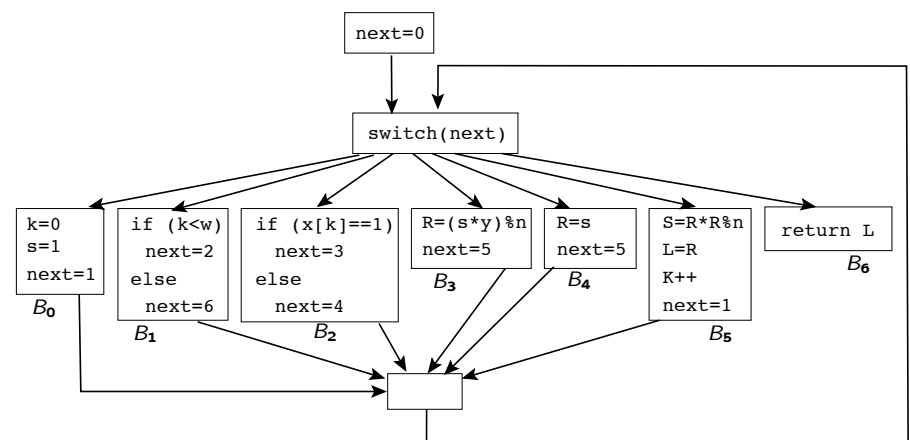
```



```

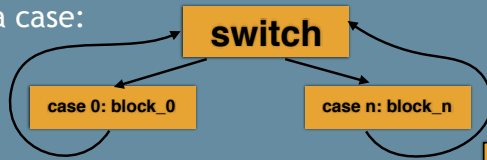
int modexp(int y, int x[], int w, int n) {
  int R, L, k, s;
  int next=0;
  for(;;)
    switch(next) {
      case 0 : k=0; s=1; next=1; break;
      case 1 : if (k<w) next=2; else next=6; break;
      case 2 : if (x[k]==1) next=3; else next=4; break;
      case 3 : R=(s*y)%n; next=5; break;
      case 4 : R=s; next=5; break;
      case 5 : s=R*R%n; L=R; k++; next=1; break;
      case 6 : return L;
    }
}

```



Flattening Algorithm

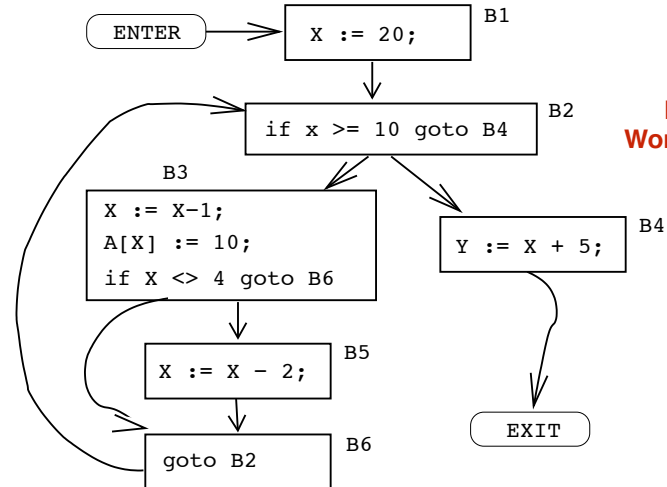
1. Construct the CFG
2. Add a new variable **int next=0;**
3. Create a switch inside an infinite loop, where every basic block is a case:



4. Add code to update the **next** variable:

```

case n: {
  if (expression)
    next = ...
  else
    next = ...
}
  
```

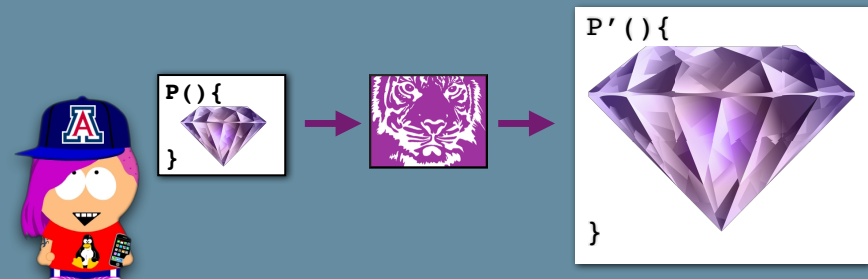


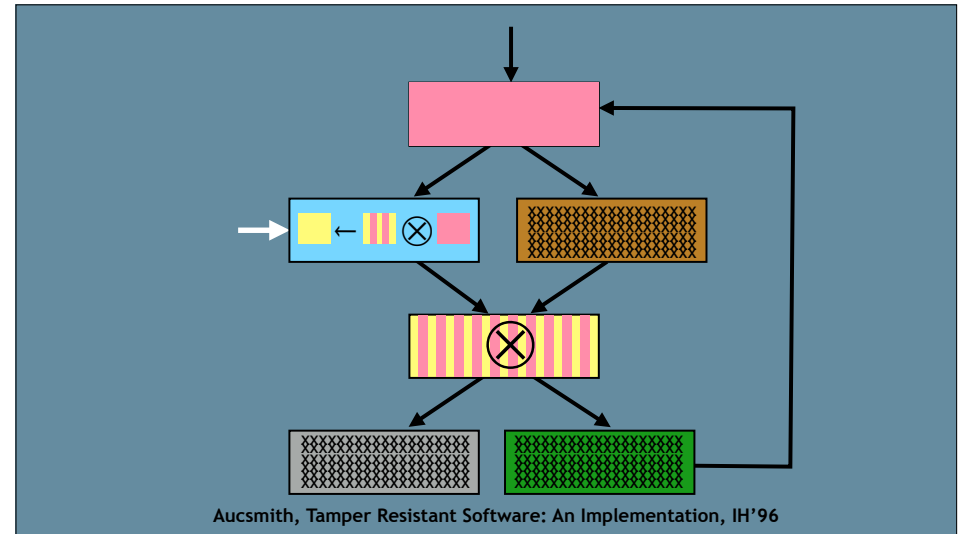
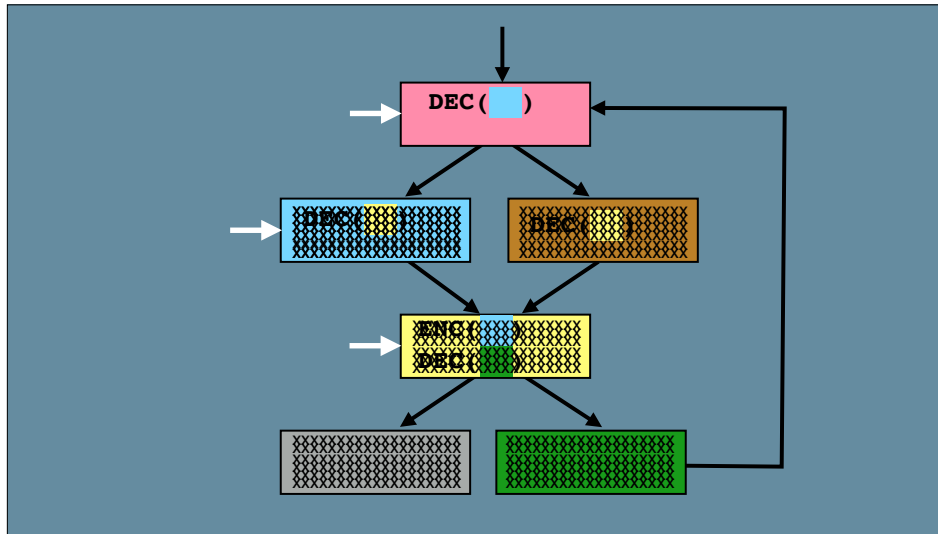
Flatten this CFG!
Work with your friends!

Dynamic Obfuscation

Dynamic Obfuscation

- Use self-modifying code
- Keep the code in constant flux
- The code should never exist in cleartext





Exercise!

```
*****
* 6) Dynamic Obfuscation
*****
tigress --Environment=x86_64:Linux:Gcc:4.6 \
--Transform=JitDynamic \
--Functions=fib \
--JitDynamicCodecs=xtea \
--JitDynamicDumpCFG=false \
--JitDynamicBlockFraction=%50 \
--out=fib6.c fib.c
```

- Please add this to fib.c:

```
#include "tigress_unstable/jitter-amd64.c"
```

```
// #include "apple.h"
#include <stdio.h>
#include <stdlib.h>
#include "tigress_unstable/jitter-amd64.c"

int fib(int n) {
    int a = 1; int b = 1; int i;
    for (i = 3; i <= n; i++) {
        int c = a + b; a = b; b = c;
    };
    return b;
}

int main(int argc, char** argv) {
    if (argc != 2) {
        printf("Give one argument!\n"); abort();
    };
    long n = strtol(argv[1], NULL, 10);
    int f = fib(n);
    printf("fib(%li)=%i\n", n, f);
}
```

<http://tigress.cs.arizona.edu/fib.c>

Anti Tamper

Tamperproofing has to do two things:

1. detect tampering
2. respond to tampering

Essentially:

```
if (tampering-detected())  
    respond-to-tampering()
```

but this is too unstealthy!

```
int foo () {  
    if (today > "Aug 17,2016"){  
        printf("License expired!");  
        abort;  
    }  
}
```

```
check(){  
    if (hash(foo)!=42)  
        abort()  
}
```



```
int foo() {  
    ... ..  
}
```



```
int foo_copy() {  
    ... ..  
}
```

Checker1

```
if (foo-has-changed-in-any-way())
```

copy

```
int foo_copy() {  
    ... ..  
}
```

```
int foo(){  
    ... ..  
}
```

Repair foo!!!

