

Glenn Smith
Crypto Hw 2.a

1- Prove that

a) $a \equiv b \pmod{n}$ implies $b \equiv a \pmod{n}$

b) prove that $a \equiv b \pmod{n}$ and $b \equiv c \pmod{n}$ imply $a \equiv c \pmod{n}$

a.

Definition of $a \equiv b \pmod{n}$:

$$a = xn + A$$

$$b = yn + B$$

$$A = B$$

Since $(A = B) \rightarrow (B = A)$ then we could flip the two sides and say $b \equiv a \pmod{n}$

b.

$$a \equiv b \pmod{n}$$

$$a = xn + A$$

$$b = yn + B$$

$$A = B$$

$$b \equiv c \pmod{n}$$

$$c = zn + C$$

$$B = C$$

$$A = B, B = C \rightarrow A = C$$

$$A = C \rightarrow a \equiv c \pmod{n}$$

2- Using extended Euclidean algorithm find the multiplicative inverse of

a) 1234 mod 4321

$$n = 4321, m = 1234, v_n = (1, 0), v_m = (0, 1)$$

$$q = 3, n' = 619, v_{n'} = (1, -3)$$

$$q = 1, m' = 615, v_{m'} = (-1, 4)$$

$$q = 1, n' = 4, v_{n'} = (2, -7)$$

$$q = 153, m' = 3, v_{m'} = (-307, 1075)$$

$$q = 1, n' = 1, v_{n'} = (309, -1082),$$

$$q = 3, m' = 0, v_{m'} = (1234, -4321)$$

Inverse is $-1082 \equiv 3239 \pmod{4321}$

b) 24140 mod 40902

These are not coprime, so they don't have a multiplicative inverse

c) 550 mod 1769

$$n = 1769, m = 550, v_n = (1, 0), v_m = (0, 1)$$

$$q = 3, n' = 113, v_{n'} = (1, -3)$$

$$q = 4, m' = 98, v_{m'} = (-4, 13)$$

$q = 1, n' = 15, vn' = (5, -16)$
 $q = 6, m' = 8, vm' = (-34, 109)$
 $q = 1, n' = 7, vn' = (39, -125)$
 $q = 1, m' = 1, vm' = (-73, 234)$
 $q = 7, n' = 0, vn' = (550, -1763)$
 550 is its own inverse

3- Determine which of the following are reducible over $GF(2)$

a) $x^3 + 1$
 $(x+1)(x^2 + x + 1) = x^3 + 2x^2 + 2x + 1 \equiv x^3 + 1$
 Reducible

b) $x^3 + x^2 + 1$
 Irreducible

c) $x^4 + 1$
 $(x^2 + 1)(x^2 + 1) = x^4 + 2x^2 + 1 \equiv x^4 + 1 \pmod{2}$
 Reducible

4- Determine the GCD of following pair of polynomials:

a) $x^3 - x + 1$ and $x^2 + 1$ over $GF(2)$

$x^3 - x + 1$
 $-(x^3 + x)$
 $= -2x + 1 \equiv 1 \pmod{2}$

So the GCD is 1

b) $x^5 + x^4 + x^3 - x^2 - x + 1$ and $x^3 + x^2 + x + 1$ over $GF(3)$

$x^5 + x^4 + x^3 - x^2 - x + 1$
 $+2x^5 + 2x^4 + 2x^3 + 2x^2$
 $3x^5 + 3x^4 + 3x^3 + x^2 - x + 1 \equiv x^2 - x + 1 \pmod{3}$
 $x^3 + x^2 + x + 1 - x^3 + x^2 - x = 2x^2 + 1$
 $2x^2 + 1 - 2x^2 + 2x - 2 = 2x - 1$
 $x^2 - x + 1 - 4x^2 + 2x \equiv x + 1 \pmod{3}$
 $2x - 1 - 2x - 2 \equiv 0 \pmod{3}$
 So the GCD is $x + 1$

5- For a cryptosystem $\{P, K, C, E, D\}$ where

$P = \{a, b, c\}$ with

$PP(a) = 1/4$

$PP(b) = 1/4$

$PP(c) = 1/2$

$K = (k_1, k_2, k_3)$ with

$PK(k_1) = 1/2$

$PK(k_2) = 1/4$

$PK(k_3) = 1/4$

$C = \{1, 2, 3, 4\}$

Encryption table

$ek(P)$	a	b	c
k_1	1	2	1
k_2	2	3	1
k_3	3	2	4
k_4	3	4	4

$$PrC(1) = 1/8 + 1/4 + 1/8 = 1/2$$

$$PrC(2) = 1/16 + 1/8 + 1/16 = 1/4$$

$$PrC(3) = 1/16 + 1/16 + 0 = 1/8$$

$$PrC(4) = 1/8 + 0 + 0 = 1/8$$

$$H(K|C) = -\sum Pr(c) Pr(k|c) \log_2(Pr(k|c))$$

$$Pr(k|c) = Pr(c|k)Pr(k) / Pr(c)$$

$Pr(k_4) = 0$ so just ignore it

$$Pr(1|k_1) = Pr(a) + Pr(c) = 3/4$$

$$Pr(2|k_1) = Pr(b) = 1/4$$

$$Pr(3|k_1) = 0$$

$$Pr(4|k_1) = 0$$

$$Pr(1|k_2) = Pr(c) = 1/2$$

$$Pr(2|k_2) = Pr(a) = 1/4$$

$$Pr(3|k_2) = Pr(b) = 1/4$$

$$Pr(4|k_2) = 0$$

$$Pr(1|k_3) = 0$$

$$Pr(2|k_3) = Pr(b) = 1/4$$

$$Pr(3|k_3) = Pr(a) = 1/4$$

$$Pr(4|k_3) = Pr(c) = 1/2$$

$$Pr(k_1|1) = Pr(1|k_1) Pr(k_1) / Pr(1) = (3/4) (1/2) / (1/2) = 3/4$$

$$Pr(k_1|2) = Pr(2|k_1) Pr(k_1) / Pr(2) = (1/4) (1/2) / (1/4) = 1/2$$

$$\begin{aligned}
\Pr(k_1|3) &= \Pr(3|k_1) \Pr(k_1) / \Pr(3) = (0) (1/2) / (1/8) = 0 \\
\Pr(k_1|4) &= \Pr(4|k_1) \Pr(k_1) / \Pr(4) = (0) (1/2) / (1/8) = 0 \\
\Pr(k_2|1) &= \Pr(1|k_2) \Pr(k_2) / \Pr(1) = (1/2) (1/4) / (1/2) = 1/4 \\
\Pr(k_2|2) &= \Pr(2|k_2) \Pr(k_2) / \Pr(2) = (1/4) (1/4) / (1/4) = 1/4 \\
\Pr(k_2|3) &= \Pr(3|k_2) \Pr(k_2) / \Pr(3) = (1/4) (1/4) / (1/8) = 1/2 \\
\Pr(k_2|4) &= \Pr(4|k_2) \Pr(k_2) / \Pr(4) = (0) (1/4) / (1/8) = 0 \\
\Pr(k_3|1) &= \Pr(1|k_3) \Pr(k_3) / \Pr(1) = (0) (1/4) / (1/2) = 0 \\
\Pr(k_3|2) &= \Pr(2|k_3) \Pr(k_3) / \Pr(2) = (1/4) (1/4) / (1/4) = 1/4 \\
\Pr(k_3|3) &= \Pr(3|k_3) \Pr(k_3) / \Pr(3) = (1/4) (1/4) / (1/8) = 1/2 \\
\Pr(k_3|4) &= \Pr(4|k_3) \Pr(k_3) / \Pr(4) = (1/2) (1/4) / (1/8) = 1
\end{aligned}$$

$$\begin{aligned}
H(K|C) &= -\sum \Pr(c) \Pr(k|c) \log_2(\Pr(k|c)) \\
&= -(\Pr(1)(\Pr(k_1|1)\log_2 \Pr(k_1|1) + \Pr(k_2|1)\log_2 \Pr(k_2|1) + \Pr(k_3|1)\log_2 \Pr(k_3|1)) \\
&\quad + \Pr(2)(\Pr(k_1|2)\log_2 \Pr(k_1|2) + \Pr(k_2|2)\log_2 \Pr(k_2|2) + \Pr(k_3|2)\log_2 \Pr(k_3|2)) \\
&\quad + \Pr(3)(\Pr(k_1|3)\log_2 \Pr(k_1|3) + \Pr(k_2|3)\log_2 \Pr(k_2|3) + \Pr(k_3|3)\log_2 \Pr(k_3|3)) \\
&\quad + \Pr(4)(\Pr(k_1|4)\log_2 \Pr(k_1|4) + \Pr(k_2|4)\log_2 \Pr(k_2|4) + \Pr(k_3|4)\log_2 \Pr(k_3|4)))
\end{aligned}$$

$$\begin{aligned}
&= -((1/2)((3/4)\log_2 (3/4) + (1/4)\log_2 (1/4)) \\
&\quad + (1/4)((1/2)\log_2 (1/2) + (1/4)\log_2 (1/4) + (1/4)\log_2 (1/4)) \\
&\quad + (1/8)((1/2)\log_2 (1/2) + (1/2)\log_2 (1/2)) \\
&\quad + (1/8)((1)\log_2 (1)))
\end{aligned}$$

$$= 0.906$$

