

北京邮电大学 2006——2007 学年第 2 学期

《数据库系统原理》期末考试试题（B）

考试 注意 事项	一、学生参加考试须带学生证或学院证明，未带者不准进入考场。学生必须按照监考教师指定座位就坐。 二、书本、参考资料、书包等物品一律放到考场指定位置。 三、学生不得另行携带、使用稿纸，要遵守《北京邮电大学考场规则》，有考场违纪或作弊行为者，按相应规定严肃处理。 四、学生必须将答题内容做在试题答卷上，做在试题及草稿纸上一律无效。 五、填空题用英文答，中文答对得一半分。										
考试 课程				考试时间		2007 年 6 月 25 日					
题号	一	二	三	四	五	六	七	八	九	十	总分
满分	9	11	6	10	20	15	6	6	8	9	
得分											
阅卷 教师											

1. Fill in blanks. (1×9 points)

- (1) Application programs are said to exhibit ___physical independency_____, if they do not depend on the physical schema, and thus need not be rewritten if the physical schema changes.
- (2) DBS design can be divided into three stages, at the ___conceptual design_____ stage, the E-R model is used to describe the data objects in the world and the associations among the objects.
- (3) An_____entity set_____ is a collection of entities of the same type.
- (4) With respect to integrity constraints in DBS, ___referential integrity, or foreign key_____ constraints ensures that a value appearing in a relation for a given set of attributes also appears for a certain set of attributes in another relation.
- (5) The commonly-used schemes of organization of records in files are heap file organization, ___sequential file organization_____, and hashing file organization.
- (6) With respect to ACID properties of the transaction, ___isolation_____guarantees that, even though multiple transactions may execute concurrently, the system guarantee that all transactions seem to execute serially, so that the consistent states of DB can be preserved.
- (7) The schedule S and S' are ___conflict equivalent_____if S can be transformed into S' by a series of swaps of non-conflicting instructions.

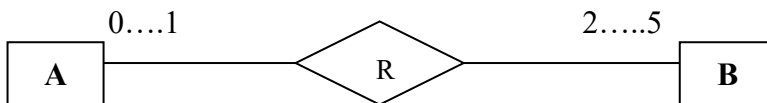
- (8) With respect to handling of transaction deadlock, two commonly used schemes are deadlock prevention and deadlock____detection and recovery_____
- (9).The several concurrent control strategies are lock-based protocols,____timestamp protocols, validation-based protocols and multiversion schemes.

答案:

- (1) physical independency
- (2) conceptual design
- (3) entity set
- (4) referential integrity, *or* foreign key
- (5) sequential file organization
- (6) isolation
- (7) conflict equivalent
- (8) detection and recovery
- (9) timestamp

2. Choice (1×11 points)

- (1) For the E-R diagram given below, the mapping cardinality from **A** to **B** is C



- A. one-to-many B. one-to-one C. many-to-one D. many-to-many

- (2) In SQL Select statement, the *where* subclause corresponds to the relational algebra operation D

- A. Π B. ρ C. \Join D. σ

- (3) Given the schema $R(A, B, C, D, E, F, G)$ and the functional dependencies $F = \{AE \rightarrow D, CDFG \rightarrow AE, D \rightarrow BF, CF \rightarrow AE\}$ holding on it, B is a trivial functional dependency.(平凡函数依赖)

- A. $AE \rightarrow D$ B. $CDFG \rightarrow DG$ C. $D \rightarrow BF$ D. $AE \rightarrow BF$

- (4) Given a relation $r(R)$, which one of the following functional dependencies is **not** satisfied by r

D

- A. $D \rightarrow C$ B. $BD \rightarrow A$ C. $C \rightarrow D$ D. $AB \rightarrow D$

A	B	C	D
4	4	1	1
6	5	4	9
3	4	2	3
6	5	5	2
9	1	0	5

(5) Given a schema $R(A, B, C, D)$ and $F=\{A \rightarrow B, B \rightarrow C\}$, R is decomposed into $R_1(A, B, D)$ and $R_2(A, C, D)$, this decomposition is A.

A. lossless B. lossy C. dependency preserving

(6) In SQL language, the statement that can be used for security control is C

A. commit B. rollback C. grant D. select

(7) With respect to the relation **account**(*account_number*, *branch_name*, *balance*), all information except D belong to meta-data and are stored in the data dictionary.

A. the name and type of the attribute *branch_name*
 B. the numbers of all the tuples in **account**
 C. the type of the index for **account**, defined on *branch_name*
 D. a tuple (A-101, downtown, 500) in **account**

(8) Considering the measures of query cost, A is the predominant cost

A. disk accesses, i.e. accessing data from disk
 B. CPU times to execute a query
 C. communication costs, for example, in distributed or parallel DBS

(9) As for the following equivalence rules for transformation of relational expressions, which one is not right? B

A. $\sigma_{\theta_1 \wedge \theta_2}(E) = \sigma_{\theta_1}(\sigma_{\theta_2}(E))$

B. $\sigma_{\theta}(E_1 \times E_2) = E_1 \div E_2$

C. $E_1 \cup E_2 = E_2 \cup E_1$

D. $\Pi_L(E_1 \cup E_2) = (\Pi_L(E_1)) \cup (\Pi_L(E_2))$

(10) With respect to data access in DBS, the operation C is used to transfer the block in the disk buffer to the disk, and replaces the appropriate physical block here.

A. input(B_X) B. read(X) C. output(B_X) D. write(X)

(11) Consider the following concurrent schedules on transaction T1 and T2. Assuming TS (T1) < TS(T2), which one is **not** in accordance with timestamp protocol ? A.

A. S1 B. S2 C. S3

S1

T1	T2...
...	...
	read(X)
	write(X)
	read(Z)
write(Y)	
	write(Y)
read(Y)	
read(Z)	

S2

T1	T2
...	...
	read(X)
write(Y)	
	write(X)
	read(Z)
read(Y)	
	write(Y)
read(Z)	

S3

T1	T2
...	...
write(Y)	
	read(X)
read(Y)	
read(Z)	
	write(X)
	read(Z)
	write(Y)

答案:

- (1) C
- (2) D
- (3) B
- (4) D
- (5) A
- (6) C
- (7) D
- (8) A
- (9) B
- (10) C
- (11) A

3. (6 points) Given relational schemas $Supplier(\underline{S\#}, SNAME, CITY)$, $Part(\underline{P\#}, PNAME, COLOR, WEIGHT)$, $Job(\underline{J\#}, JNAME, CITY)$, $SPJ(\underline{S\#}, \underline{P\#}, \underline{J\#}, QTY)$. In the relation $Supplier$, $S\#$ is the supplier identifier, $SNAME$ is the supplier name, and $CITY$ is the city where the supplier lives; In the relation $Part$, $P\#$ is the part code, $PNAME$ is the part name, $COLOR$ is the part color, and $WEIGHT$ is the part weight; In the engineering relation Job , $J\#$ is the engineering identifier, $JNAME$ is the engineering job name, $CITY$ is the city where the engineering is located; A *part* can be supplied by some *suppliers*, and be used for some engineering *jobs*. The supplying relation SPJ describes the association among the objects *part*, *supplier* and *job*, QTY is the total number of the *parts* supplied.

Give relational algebra expressions for the following queries:

- (1) Find the *names* of the suppliers who provide *red parts* for the engineering jobs located either in Beijing or Shanghai (3points)
- (2) Find the *code*, *name* and *weight* of the part provided by the suppliers living in the city, at which the engineering job is also located and makes use of this part. (3points)

答案:

(1) $\Pi_{SNAME}(\sigma_{(Job.CITY='Beijing' \vee Job.CITY='Shanghai')} \text{ and } Part.COLOR='red' (Supplier \bowtie Part \bowtie Job \bowtie SPJ))$ (3 points)

(2) $\Pi_{P\#, PNAME, WEIGHT} (\sigma_{supplier.CITY=Job.CITY} (Supplier \bowtie Part \bowtie Job \bowtie SPJ))$ (3 points)

4. (10 points) Considering the schema $R=(athlete-id, game-event, grade, category \text{ of games, manger of games})$ that describes the sports meeting. It is assumed that

- for each athlete, if he takes part in a game event, he will achieve one and only one grade

- each game event belongs to one and only one category
- each category is managed by one and only one manager

这两个条件就是，虽然语句是被动句，但是前面的决定了后面的。

- (1) According to the descriptions mentioned above, list the functional dependency set F that holds on R (3 points)
- (2) List all the candidate keys of R. (2 points)
- (3) Give a lossless and dependency-preserving decomposition of R into 3NF. (5 points)

答案：

函数依赖如下：

athlete-id, game-event → grade (1 point)

game-event → category of games (1 point)

category of games → manger of games (1 point)

候选键是：

athlete-id, game-event (2 points)

分解的 3NF 是：

(athlete-id, game-event, grade) (2 points)

(game-event, category of games) (2 points)

(category of games, manger of games) (1 point)

5. (20 points) A pharmacies database needs to store the information about the pharmaceutical companies (制药公司) and others. The relevant information is as follows:

- Each pharmaceutical company is identified by its name and has a phone number.
- Each pharmacy (药房) is identified by its name and has an address, and several phone numbers.
- For each drug, the trade name and formula must be recorded, and the trade name identifies a drug uniquely.
- Each drug is developed by only one given pharmaceutical company, and one pharmaceutical company can develop several drugs and perhaps develops no drug now.
- Each pharmacy sells several drugs and has a price for each. A drug could be sold by several pharmacies, and the price may vary from one pharmacy to another.
- Pharmaceutical companies make contracts with pharmacies, but some pharmaceutical companies may have no contract. A pharmaceutical company can sign contracts with several pharmacies, and a pharmacy can also enter into contracts with several pharmaceutical companies. Each contract is described by a start date, an end date and the content of the contract.

- (1) Design the E/R diagram for hospital database on basis of the information mentioned above. (10 points)

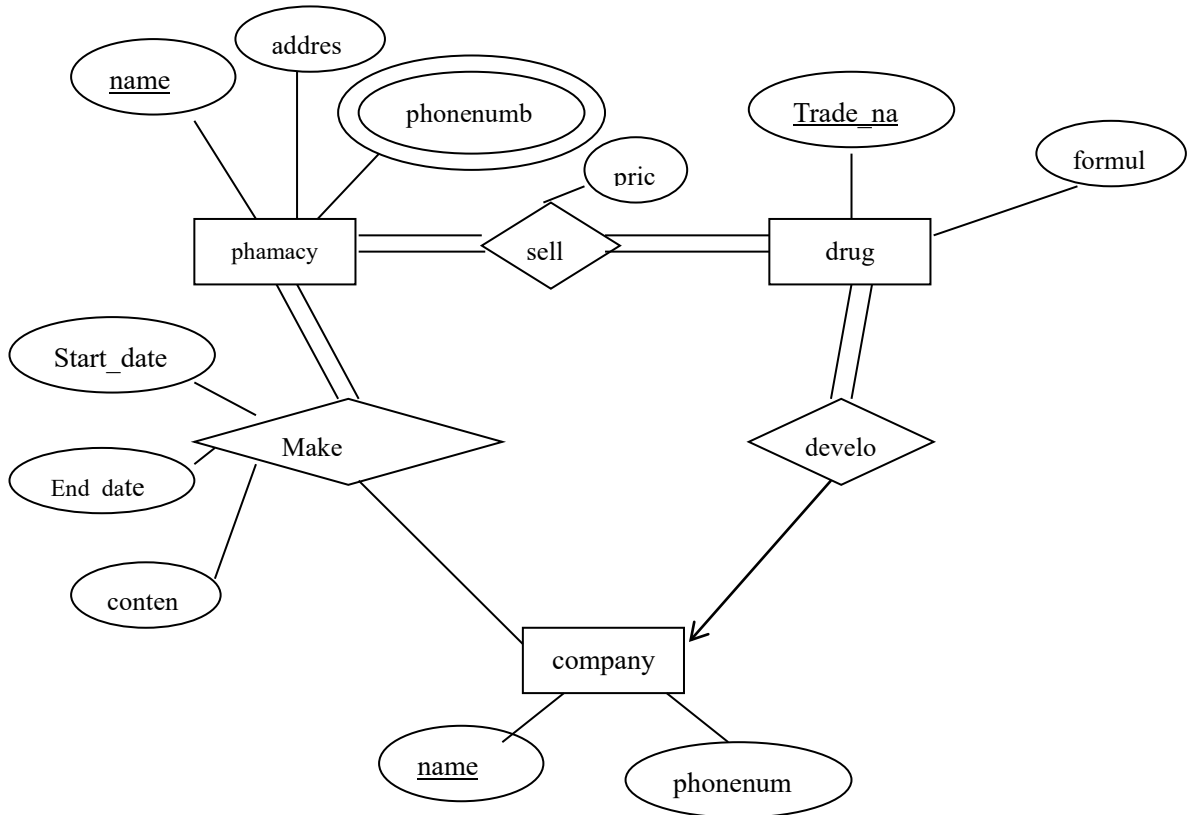
Note: the mapping cardinality of each relationship and participation of each entity to the

relationship should be described in the diagram.

(2) Convert the E-R diagram to the proper relational schemas, and give the primary key of each relation schema by underlines. (10 points)

答案:

(1) E-R 图



(2) 转换后的关系模式

Pharmacy(name, address)

Pharmacy_phone(name, phonenumber)

Drug(Trade_name, formulary, company_name)

Company(name, phonenumber)

Sell(pharmacy_name, trade_name, price)

Contract(pharmacy_name, company_name, start_date, end_date, content)

其中 drug、pharmacy 实体的参与联系答案不唯一。

评分标准:

ER 图有实体、属性、联系、主键、参与度、对应约束等得分点, 错一个小点扣 0.5 分。

表分基本表、属性、主键等几个点, 每一个小点错扣 0.5 分。

6. (15 points) There are three relations in the database, describing the employees and the companies that the employees work at,

Employee(employee-id, name, age, sex, e-city)

Company(company-id, company-name, company-city)

Works(employee-id, company-id, salary)

- (1) The salaries of the employees, who live in Beijing and work at the company named SHOUGANG, increase by \$100. Give a SQL statement to add this information into the database. (3 points)
- (2) Give an SQL statement to find out the *name* and *id* of all the *employees*, who are male, work at the company named SHOUGANG, and whose *salaries* are below \$2000. (3 points)
- (3) Translate the SQL statement in (2) into an initial query tree, and give an optimized query tree for it, by means of heuristic query optimization. (9 points)

答案:

```
(1) update  Works
      set    salary = salary + 100
      where  employee-id in ( select  employee-id
                              from    Employee
                              where e-city='Beijing'
                              and    company-id in ( select  company-id
                                                    from    Company
                                                    where  company-name='SHOUGANG' )
```

或:

```
update  Works
set     salary = salary + 100
where   employee-id in ( select  employee-id
                        from     employee as E, company as C, works as W
                        where    E.employee-id = W.employee-id   and
                        C.company-id = W. company-id   and
                        e-city = 'Beijing' and company-name='SHOUGANG')
```

或:

```
update  Works
set     salary = salary + 100
where   exists ( select  employee-id
                from    employee as E, company as C, works as W
                where   E.employee-id = W.employee-id   and
                C.company-id = W. company-id   and
                e-city = 'Beijing' and company-name='SHOUGANG'
```

或:

```

with employee-B-S (employee-id) as
  select employee-id
  from employee as E, company as C, works as W
  where E.employee-id = W.employee-id and
        C.company-id = W. company-id and
        e-city='Beijing' and company-name='SHOUGANG'

update Works
set salary = salary + 100
where Works.employee-id in employee-B-S

```

.....要好好研究下 with 语句，是怎么回事，
 评分标准：正确写出 Update 语句结构可得 1.5 分

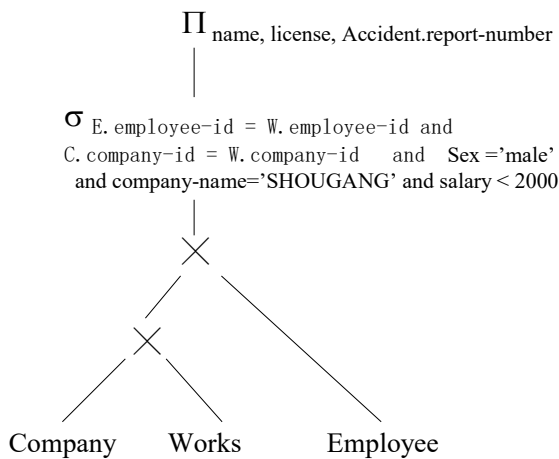
```

(2) select name, employee.employee-id
      from employee as E, company as C, works as W
      where E.employee-id = W.employee-id and
            C.company-id = W. company-id and
            Sex ='male' and company-name='SHOUGANG' and salary < 2000

```

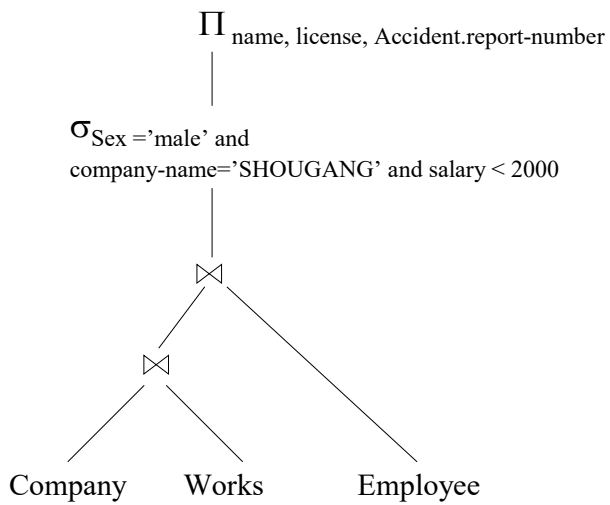
评分标准：漏掉连接条件扣 1 分

(3) 初始查询树



(a) initial query tree

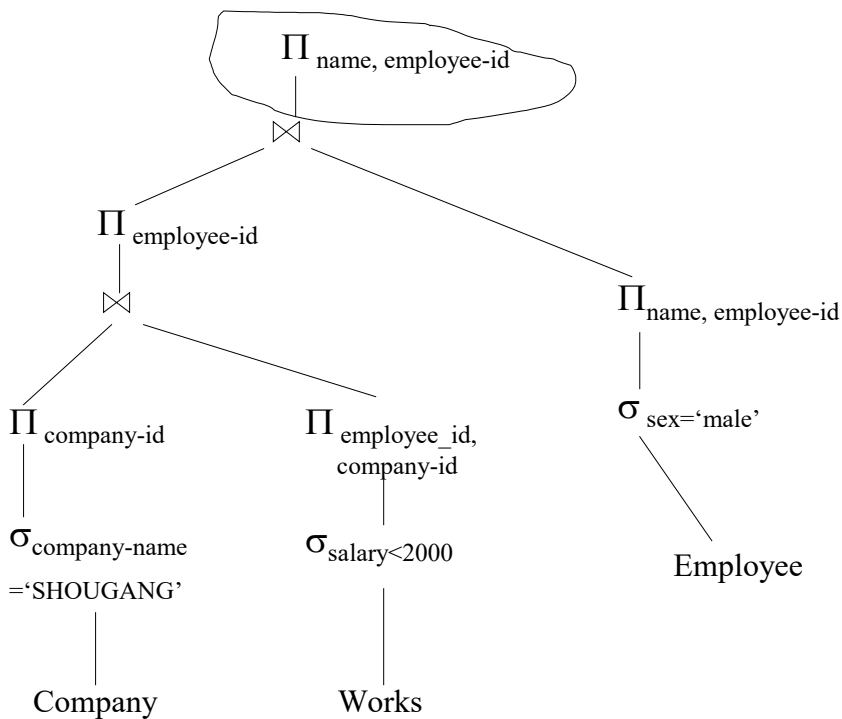
或：



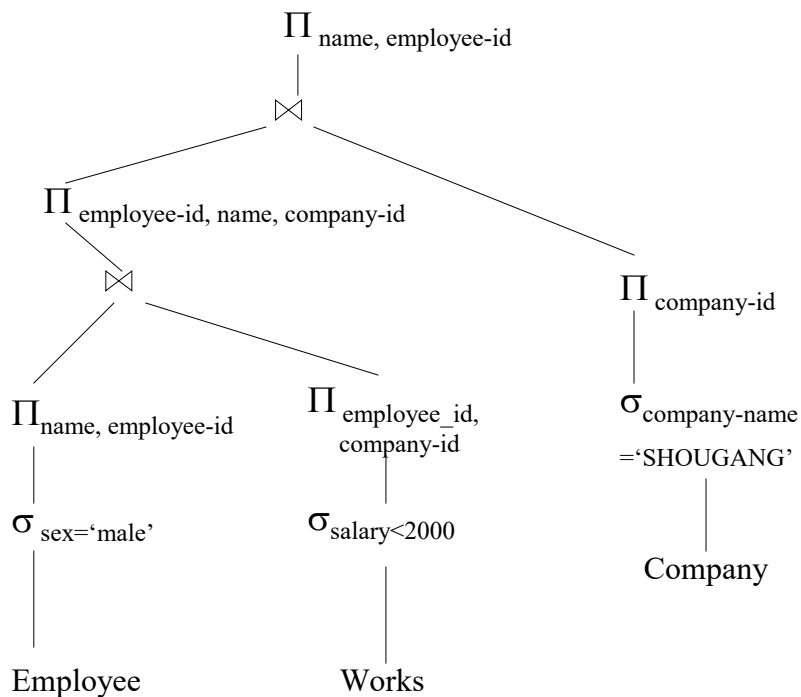
(a) initial query tree

评分标准：漏掉连接条件扣 1 分
叶结点关系表示错误扣一分

优化后的查询树：



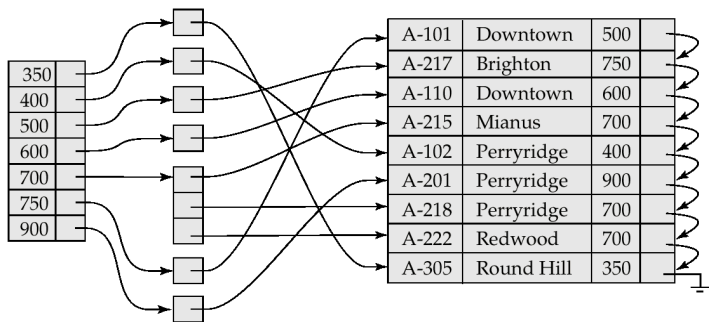
或：



评分标准：每个选择、投影操作 0.5 分

7. (6 points) Given the data file *account(account_number, branch_name, balance)* and the index defined on *balance*, as shown in the following figure,

- (1) Is the index a dense or sparse index, why? (3 points)
- (2) Is the index a clustering or non- clustering index, why? (3 points)



- 答案：
- (1) 稠密索引。因为对于被索引文件的 *balance* 属性上的每个值（共有 7 个），在索引中都有一个对应的索引项。
 - (2) non- clustering index。因为被索引文件中记录的排列顺序与索引项指定的查询顺序不一致。

评分标准：对各小题，正确回答是否得 1.5 分，正确阐述原因得 1.5 分。

8. (6 points) With respect to the concurrent schedule S on the transaction set {T1, T2, T3},

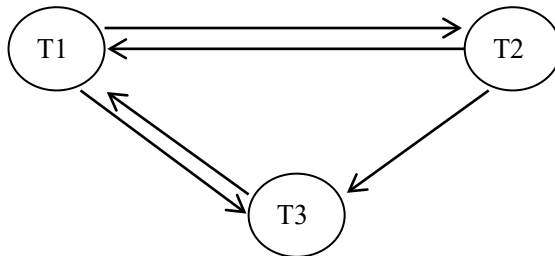
(1) Give the precedence graph G(S) for S (2 points)

(2) Is S conflict serializable? and why? (4 points)

S		
T1	T2	T3
read(P)		
	read(P)	
write(Q)		
	write(R)	
	write(Q)	
		write(Q)
		read(R)
write(R)		

答案:

(1)



评分标准: 缺一条边扣 0.5 分

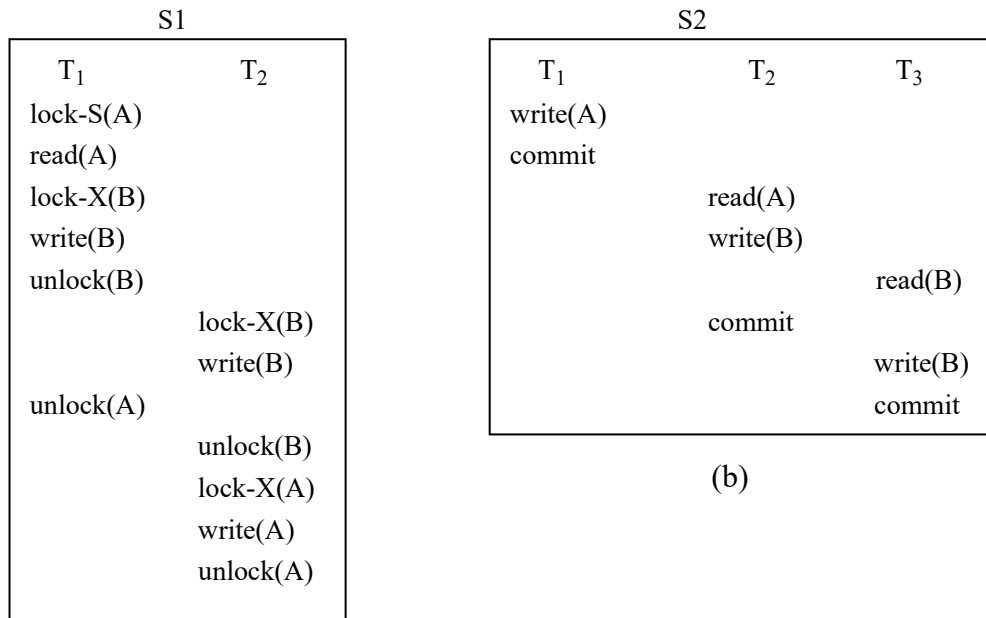
(2) 非冲突可串行的。因为前驱图中存在环路。

评分标准: 正确回答是否冲突可串行得 2 分, 正确阐述原因得 2 分。

9. (8 points) Considering the schedules in the following figures

(1) Is the concurrent schedule S1 shown in Fig.(a) under the two-phase locking protocol? and why ? (4 points)

(2) Is the concurrent schedule S2 shown in Fig.(b) a cascadeless schedule? and why ? (4 points)



(a)

答案:

(1) 不是符合两阶段锁协议的并发调度。因为事务 T₂ 在 unlock(B)后又 lock(A),不符合协议规定的锁增长和锁缩减的两个阶段。

(2) 不是无级联回滚的调度。因为事务 T₃ 读了 T₂ 修改过但没提交的数据 B, T₂ 若在提交前失败回滚, 将引起 T₃ 的级联回滚。

评分标准:

回答对第一问得两分。

10. (9 points) Considering the concurrent transactions T₁, T₂, T₃, T₄ and T₅, and the data items

A, B, C and D modified by these transactions. It is assumed that

- the initial values of these data items are A=10, B=20, C=0, D=40.
- immediate database modification and checkpoint techniques are employed

With respect to the log in the following figure that describes the concurrent executing of T₁, T₂, T₃, T₄ and T₅, when a failure occurs, the log-based recovery scheme consults the log to determine the recovery operations (i.e. **redo**, **undo**, **ignore**) done on T₁, T₂, T₃, T₄ and T₅.

(1) Which transaction needs to be undone, which transaction needs to be redone? And which transaction can be ignored ? (5 points)

- (2) After recovery operations on T₁, T₂, T₃, T₄ and T₅ are completed, what are the values of the data items A, B, C and D in the database? (4 points)

```
<T1 start>
<T1, A, 10, 20>
<T2 start>
<T3 start>
<T2, B, 20, 10>
<T1 commit>
<T3, C, 0, 10>
<T3, C, 10, 30>
<checkpoint {T2, T3}>
<T4 start>
<T4, A, 20, 40>
<T2, B, 10, 35>
<T5 start>
<T5, D, 40, 30>
<T4, A, 40, 50>
<T5, commit>
<T4, commit>

**Crash**
```

答案:

(1) redo: T₄, T₅ (2 points)

undo: T₂, T₃ (2 points)

ignored: T₁ (1 point)

(2) A = 50 B = 20 C = 0 D = 30 (4 points)