



西安交通大学
XI'AN JIAOTONG UNIVERSITY

SDN 实验报告

Lab2 自学习交换机与广播风暴

课程名称： 软件定义网络

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指导老师： 张鹏

2023 年 3 月 31 日

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实验名称： Lab2 自学习交换机与广播风暴 实验类型： 设计实验 同组学生姓名： _____

一、 实验任务一：自学习交换机

1. 背景介绍



图 1: ARPANET-1

1969 年的 ARPANET 非常简单，仅由四个结点组成。假设每个结点都对应一个交换机，每个交换机都具有一个直连主机，你的任务是实现不同主机之间的正常通信。前文给出的简单交换机洪泛数据包，虽然能初步实现主机间的通信，但会带来不必要的带宽消耗，并且会使通信内容泄露给第三者。因此，请在简单交换机的基础上实现二层自学习交换机，避免数据包的洪泛。

2. 自学习交换机原理

- (a) 控制器为每个交换机维护一个 **mac-port** 映射表。
- (b) 控制器收到 **packet_in** 消息后，解析其中携带的数据包。
- (c) 控制器学习 **src_mac - in_port** 映射。
- (d) 控制器查询 **dst_mac**，如果未学习，则洪泛数据包；如果已学习，则向指定端口转发数据包 (**packet_out**)，并向交换机下发流表项 (**flow_mod**)，指导交换机转发同类型的数据包。

3. 关键代码设计思路

(1) 未考虑 Buffer 代码

在构造函数中添加一个全局 mac 与 port 映射表, 相当于存储所有传统交换机的 mac 表。

```
1 self.mac_to_port = {}
```

将收到的包源 mac 地址与接收端口对应, 相当于传统交换机自学习 mac 表的过程。

```
1 self.mac_to_port[dpid][src] = in_port
```

将收到的包与映射表匹配, 如果能匹配到, 则输出端口为表中端口, 如果匹配不到, 则洪泛数据包。

```
1 #if match, then send to outport
2 if dst in self.mac_to_port[dpid]:
3     out_port = self.mac_to_port[dpid][dst]
4 #if not match, then flood
5 else:
6     out_port = ofp.OFPP_FLOOD
```

控制器对交换机作出指示, 如果输出端口不是洪泛, 添加流表项, 增加匹配次数, 完成数据包的转发, 如果洪泛, 则直接转发。

```
1 #pass information
2 actions = [parser.OFPACTIONOutput(out_port)]
3 if out_port != ofp.OFPP_FLOOD:
4     match = parser.OFPMATCH(in_port=in_port, eth_dst=dst, eth_src=src)
5     self.add_flow(dp, 1, match, actions)
6 data = msg.data
7 out = parser.OFPPACKETOut(datapath=dp,buffer_id=msg.buffer_id,in_port=in_port,
8     actions=actions, data=data)
9 dp.send_msg(out)
```

(2) 基于 Buffer 的代码改进

Buffer 是交换机中的概念, 意思就是说数据包发入交换机, 交换机之后与控制器交互, 此时可以不把整个数据包发给控制器, 可以通过 Bufferid 进行通信, 以下是对三种消息基于 Buffer 改进的介绍:

(a) Packetin 消息: 用于标记缓存在交换机中的数据报文 id, 如报文被 action 上送到控制器中 maxlen 字段或者 table_miss 消息限制长度, 而通过 bufferid 将报文缓存在交换机中, 以便被另外两种消息来调用;

(b) Packetout 消息: 用于控制器将原先 buffer 在交换机中的报文, 通过 Packetout 个形式从交换机的某个物理口送出去;

(c) Flowmod 消息: 如果 flowmod 中带有 bufferid, 那么说明这个 flowmod 需要做两件事情, 第一是正常下发一条 flow, 其次是把交换机中先前 buffer 的那个数据报文, Packetout 到 table 来匹配一次下的这条 flow; 注意以上两个指令都是通过这个带有 bufferid 的消息执行的, 不需要控制器另外下 packet_out 消息, 这种设计思路是非常巧妙的。

(d) 代码修改:

首先将 add_flow 函数修改成能够支持 Buffer 机制的函数, 即支持上文介绍的 Buffer 机制的 Flow-mod。

```

1 def add_flow(self, datapath, priority, match,
2   actions,buffer_id=None,idle_timeout=0,hard_timeout=0):
3     dp = datapath
4     ofp = dp.ofproto
5     parser = dp.ofproto_parser
6     inst = [parser.OFPInstructionActions(ofp.OFPIT_APPLY_ACTIONS, actions)]
7     if buffer_id:
8         mod = parser.OFPFlowMod(datapath=datapath,
9           priority=priority,buffer_id=buffer_id,idle_timeout=idle_timeout,
10          hard_timeout=hard_timeout,match=match,instructions=inst)
11     else:
12         mod = parser.OFPFlowMod(datapath=datapath,
13           priority=priority,idle_timeout=idle_timeout,
14          hard_timeout=hard_timeout,match=match, instructions=inst)
15     dp.send_msg(mod)

```

在代码中加入以下逻辑:

(1) 输出不是洪泛的情况下, 如果支持 Buffer 机制的话, 直接发出支持 Buffer 机制的 Flowmod 消息即可 (此时不需要 Packetout, 可以直接 return, 因为原报文重新发到 table 匹配)。

(2) 如果是支持 Buffer 机制且输出洪泛的话, 将数据 data 置为空, Packetout 只需要传输一个 buffer_id 即可转发消息。

```

1     actions = [parser.OFPActionOutput(out_port)]
2     if out_port != ofp.OFPP_FLOOD:
3         match = parser.OFPMatch(in_port=in_port, eth_dst=dst, eth_src=src)
4         if msg.buffer_id != ofp.OFP_NO_BUFFER:
5             self.add_flow(dp, 1, match, actions, msg.buffer_id)
6             return
7         else:
8             self.add_flow(dp, 1, match, actions)
9     data = None
10    if msg.buffer_id == ofp.OFP_NO_BUFFER:
11        data = msg.data
12    out = parser.OFPPacketOut(datapath=dp,buffer_id=msg.buffer_id,in_port=in_port,
13      actions=actions, data=data)
14    dp.send_msg(out)

```

4. 运行结果与分析

可以由图 2, 图 3 看出在 UCLA ping UTAH 的过程中, 数据包不再发给 UCSB, 实现了自学习交换机。

```

mininet> UCLA ping UTAH
PING 10.0.0.4 (10.0.0.4) 56(84) bytes of data.
64 bytes from 10.0.0.4: icmp_seq=1 ttl=64 time=340 ms
64 bytes from 10.0.0.4: icmp_seq=2 ttl=64 time=127 ms
64 bytes from 10.0.0.4: icmp_seq=3 ttl=64 time=128 ms
64 bytes from 10.0.0.4: icmp_seq=4 ttl=64 time=127 ms
64 bytes from 10.0.0.4: icmp_seq=5 ttl=64 time=126 ms
64 bytes from 10.0.0.4: icmp_seq=6 ttl=64 time=127 ms
64 bytes from 10.0.0.4: icmp_seq=7 ttl=64 time=131 ms
64 bytes from 10.0.0.4: icmp_seq=8 ttl=64 time=126 ms
64 bytes from 10.0.0.4: icmp_seq=9 ttl=64 time=126 ms
64 bytes from 10.0.0.4: icmp_seq=10 ttl=64 time=126 ms

```

图 2: UCLA ping UTAH

35	15.292901663	fe80::c462:66ff:feb1::f02::2	ICMPv6	70	Router Solicitation from c6:62:66:ba:89:71
36	19.141923857	06:57:5b:61:6d:8d	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
37	21.389717731	fe80::28b4:6dff:fe1::f02::2	ICMPv6	70	Router Solicitation from 2a:b4:6d:1a:74:ac
38	21.438180068	fe80::2852:48ff:fed::f02::2	ICMPv6	70	Router Solicitation from 2a:52:48:d1:74:bd
39	22.200959404	fe80::1473:b2ff:feb::f02::f0	MONS	107	Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipps._tcp.local, "QM" question
40	22.633843212	fe80::2852:48ff:fed::f02::f0	MONS	107	Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipps._tcp.local, "QM" question
41	22.779126039	fe80::e807:c3ff:fee::f02::f0	MONS	107	Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipps._tcp.local, "QM" question
42	23.871527963	fe80::08ba:56ff:fee::f02::f0	MONS	107	Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipps._tcp.local, "QM" question
43	25.437447610	fe80::3c74:59ff:fe7::f02::2	ICMPv6	70	Router Solicitation from 3e:74:59:75:85:b7
44	25.599948676	fe80::9457:5bff:fe6::f02::2	ICMPv6	70	Router Solicitation from 9e:57:5b:61:6d:8d
45	27.659393205	06:57:5b:61:6d:8d	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
46	29.535357775	fe80::08ba:56ff:fee::f02::2	ICMPv6	70	Router Solicitation from da:8a:56:ef:cc:52
47	31.637864625	fe80::e807:c3ff:fee::f02::2	ICMPv6	70	Router Solicitation from ea:d7:c3:e6:8d:2c
48	76.726274558	fe80::c462:66ff:feb1::f02::2	ICMPv6	70	Router Solicitation from c6:62:66:ba:89:71
49	84.868306279	fe80::28b4:6dff:fe1::f02::2	ICMPv6	70	Router Solicitation from 2a:b4:6d:1a:74:ac
50	84.920486755	fe80::2852:48ff:fed::f02::2	ICMPv6	70	Router Solicitation from 2a:52:48:d1:74:bd
51	86.202942376	fe80::1473:b2ff:feb::f02::f0	MONS	107	Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipps._tcp.local, "QM" question
52	86.630280933	fe80::2852:48ff:fed::f02::f0	MONS	107	Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipps._tcp.local, "QM" question
53	86.754598138	fe80::e807:c3ff:fee::f02::f0	MONS	107	Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipps._tcp.local, "QM" question
54	87.883047950	fe80::08ba:56ff:fee::f02::f0	MONS	107	Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipps._tcp.local, "QM" question
55	95.160342166	fe80::3c74:59ff:fe7::f02::2	ICMPv6	70	Router Solicitation from 3e:74:59:75:85:b7
56	97.192169997	fe80::9457:5bff:fe6::f02::2	ICMPv6	70	Router Solicitation from 9e:57:5b:61:6d:8d
57	101.212682281	fe80::08ba:56ff:fee::f02::2	ICMPv6	70	Router Solicitation from da:8a:56:ef:cc:52
58	101.248287819	fe80::1473:b2ff:feb::f02::2	ICMPv6	70	Router Solicitation from 16:73:b2:07:86:c7
59	109.457807089	fe80::e807:c3ff:fee::f02::2	ICMPv6	70	Router Solicitation from ea:d7:c3:e6:8d:2c

图 3: UCSB Capture

二、 实验任务二: 广播风暴

1. 背景介绍

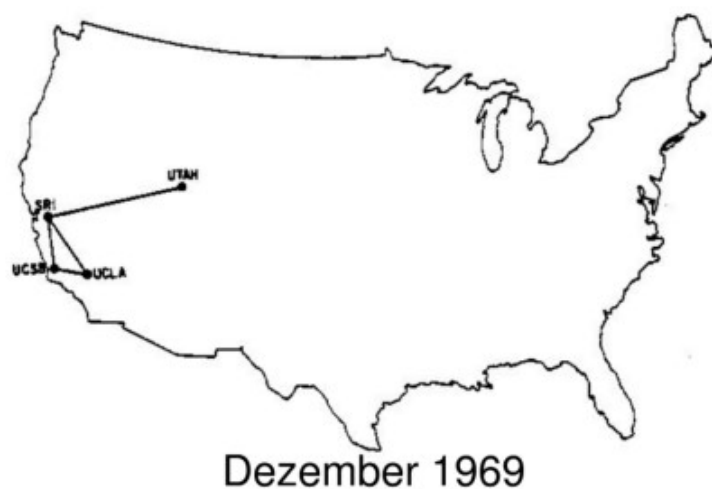


图 4: ARPANET-2

UCLA 和 UCSB 通信频繁, 两者间建立了一条直连链路。在新的拓扑 topo_1969_2.py 中运行自学习交换机, UCLA 和 UTAH 之间无法正常通信。分析流表发现, 源主机虽然只发了很少的几个

数据包, 但流表项却匹配了上千次; WireShark 也截取到了数目异常大的相同报文

2. 广播风暴及其解决思路

这实际上是 **ARP** 广播数据包在环状拓扑中洪泛导致的, 传统网络利用**生成树协议**解决这一问题。在 **SDN** 中, 不必局限于生成树协议, 可以通过多种新的策略解决这一问题。以下给出一种解决思路, 请在自学习交换机的基础上完善代码, 解决问题:

当序号为 **dpid** 的交换机从 **in_port** 第一次收到某个 **src_mac** 主机发出, 询问 **dst_ip** 的广播 **ARP Request** 数据包时, 控制器记录一个映射 (**dpid, src_mac, dst_ip**)->**in_port**。下一次该交换机收到同一 (**src_mac, dst_ip**) 但 **in_port** 不同的 **ARP Request** 数据包时直接丢弃, 否则洪泛。

3. 关键代码设计思路

在构造函数初始化一个表, 维护一个 (**dpid,src_mac,dst_ip**)->**in_port** 的映射。

```
1 self.sw = {}
```

即对于 arp 包来说, 问题在于如何分辨这个包是交换机转发而来还是主机发来的, 可以先存储一个表映射, arp 包一定是主机先发起, 所以先通过 sw 存储映射, 表示只能从这个端口进入 arp 包, 其他端口的 arp 包全部都要丢掉 (因为其他端口的 arp 数据包是环路广播的结果), 而对于其他交换机也是绑定端口和特定 arp 请求包的发送许可, 从而实现广播风暴的避免, 与生成树协议些许相似。对于符合映射的包正常转发, 而不符合映射则丢弃, 若没有对应映射, 则进行学习, 学习后正常转发。

```
1 if dst == ETHERNET_MULTICAST and ARP in header_list:
2     arp_dst_ip = header_list[ARP].dst_ip
3     if (dp.id, src, arp_dst_ip) in self.sw:
4         if self.sw[(dp.id, src, arp_dst_ip)] != in_port:
5             out = dp.ofproto_parser.OFPPacketOut(datapath=dp,
6             buffer_id=dp.ofproto.OFP_NO_BUFFER, in_port=in_port, actions=[], data=None)
7             dp.send_msg(out)
8             return
9         else:
10            self.sw[(dp.id, src, arp_dst_ip)] = in_port
```

4. 运行结果与分析

使用实验 1 的自学习交换机控制器来应用实验二, 发现明显的广播风暴现象, 流表匹配次数巨大。


```

mininet> dpctl dump-flows
*** s1 -----
cookie=0x0, duration=22.683s, table=0, n_packets=1094, n_bytes=45948, priority=1,in_port="s1-eth2",dl_src=de:36:b4:40:17:b9,dl_dst=92:14:36:df:63:48 actions=output:"s1-eth4"
cookie=0x0, duration=22.521s, table=0, n_packets=272853, n_bytes=11459826, priority=1,in_port="s1-eth3",dl_src=de:36:b4:40:17:b9,dl_dst=92:14:36:df:63:48 actions=output:"s1-eth4"
cookie=0x0, duration=99.352s, table=0, n_packets=20550, n_bytes=1777225, priority=0 actions=CONTROLLER:65535
*** s2 -----
cookie=0x0, duration=22.801s, table=0, n_packets=1094, n_bytes=45948, priority=1,in_port="s2-eth1",dl_src=de:36:b4:40:17:b9,dl_dst=92:14:36:df:63:48 actions=output:"s2-eth2"
cookie=0x0, duration=99.355s, table=0, n_packets=20522, n_bytes=1775546, priority=0 actions=CONTROLLER:65535
*** s3 -----
cookie=0x0, duration=22.594s, table=0, n_packets=273195, n_bytes=11474190, priority=1,in_port="s3-eth3",dl_src=de:36:b4:40:17:b9,dl_dst=92:14:36:df:63:48 actions=output:"s3-eth2"
cookie=0x0, duration=99.357s, table=0, n_packets=20564, n_bytes=1778389, priority=0 actions=CONTROLLER:65535
*** s4 -----
cookie=0x0, duration=22.629s, table=0, n_packets=273476, n_bytes=11485992, priority=1,in_port="s4-eth2",dl_src=de:36:b4:40:17:b9,dl_dst=92:14:36:df:63:48 actions=output:"s4-eth3"
cookie=0x0, duration=99.359s, table=0, n_packets=20554, n_bytes=1777523, priority=0 actions=CONTROLLER:65535

```

图 5: 解决前流表

15339	76.996508450	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15341	77.070429839	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15342	77.070543736	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15347	77.074058098	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15352	77.074148823	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15355	77.074206793	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15357	77.104899359	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15358	77.105023940	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15362	77.107829338	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15363	77.107888189	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15375	77.135029181	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15376	77.135146672	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15377	77.135267368	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15381	77.135725698	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15383	77.135953071	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15384	77.136064063	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15387	77.136412102	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15390	77.164599774	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15398	77.214043324	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15400	77.248398937	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15401	77.248513934	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15406	77.251441210	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15411	77.251533077	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15414	77.251585999	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15416	77.272868756	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15417	77.272997993	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15421	77.273512931	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15422	77.273638714	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15434	77.303889482	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15435	77.304007255	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15436	77.304123543	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
15440	77.304589135	92:14:36:df:63:48	Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2

图 6: 解决前 UCSB 抓包

解决 ARP 数据包在环状拓扑中的洪泛问题后, UCLA 和 UTAH 之间可以 ping 通, 并且流表项的匹配次数明显减少, UCSB 监听也没有了广播风暴现象。

```

mininet> dpctl dump-flows
*** s1 -----
cookie=0x0, duration=4.533s, table=0, n_packets=3, n_bytes=294, priority=1,in_port="s1-eth2",dl_src=5a:95:d9:35:74:c7,dl_dst=de:b9:60:4a:a9:61 actions=output:"s1-eth4"
cookie=0x0, duration=4.522s, table=0, n_packets=2, n_bytes=196, priority=1,in_port="s1-eth4",dl_src=de:b9:60:4a:a9:61,dl_dst=5a:95:d9:35:74:c7 actions=output:"s1-eth2"
cookie=0x0, duration=14.024s, table=0, n_packets=24, n_bytes=2347, priority=0 actions=CONTROLLER:65535
*** s2 -----
cookie=0x0, duration=4.542s, table=0, n_packets=3, n_bytes=294, priority=1,in_port="s2-eth1",dl_src=5a:95:d9:35:74:c7,dl_dst=de:b9:60:4a:a9:61 actions=output:"s2-eth2"
cookie=0x0, duration=4.524s, table=0, n_packets=2, n_bytes=196, priority=1,in_port="s2-eth2",dl_src=de:b9:60:4a:a9:61,dl_dst=5a:95:d9:35:74:c7 actions=output:"s2-eth1"
cookie=0x0, duration=14.022s, table=0, n_packets=11, n_bytes=983, priority=0 actions=CONTROLLER:65535
*** s3 -----
cookie=0x0, duration=14.024s, table=0, n_packets=16, n_bytes=1546, priority=0 actions=CONTROLLER:65535
*** s4 -----
cookie=0x0, duration=4.535s, table=0, n_packets=3, n_bytes=294, priority=1,in_port="s4-eth2",dl_src=5a:95:d9:35:74:c7,dl_dst=de:b9:60:4a:a9:61 actions=output:"s4-eth1"
cookie=0x0, duration=4.530s, table=0, n_packets=2, n_bytes=196, priority=1,in_port="s4-eth1",dl_src=de:b9:60:4a:a9:61,dl_dst=5a:95:d9:35:74:c7 actions=output:"s4-eth2"
cookie=0x0, duration=14.027s, table=0, n_packets=17, n_bytes=1644, priority=0 actions=CONTROLLER:65535

```

图 7: 解决后流表

1	0.00000000	fe80::fc57:ddff:fe3_ ff02::2	ICMPv6	70 Router Solicitation from fe:57:dd:34:3c:9e
2	2.048047670	fe80::4850:cbff:fe0_ ff02::2	ICMPv6	70 Router Solicitation from 4a:50:cb:0a:de:28
3	3.172501689	46:05:f6:0d:98:cb Broadcast	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2
4	3.304527659	fe80::fc57:ddff:fe3_ ff02::fb	MDNS	197 Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipp._tcp.local, "QM" question
5	15.615991905	fe80::fc57:ddff:fe3_ ff02::2	ICMPv6	70 Router Solicitation from fe:57:dd:34:3c:9e
6	19.302842640	fe80::fc57:ddff:fe3_ ff02::fb	MDNS	197 Standard query 0x0000 PTR _ipps._tcp.local, "QM" question PTR _ipp._tcp.local, "QM" question
7	20.480000513	fe80::4850:cbff:fe0_ ff02::2	ICMPv6	70 Router Solicitation from 4a:50:cb:0a:de:28

图 8: 解决后 UCSB 抓包

5. 方案优缺点分析

优点: 方法简单, 实现成本低。

缺点: 实际上还是堵塞了某端口, 占用了交换机链路带宽, 每次需要丢包。

三、 实验任务三: 附加方案

1. 新的解决思路

可以让控制器学习每个交换机与主机连接的 mac 地址和端口。直接让 arp 请求和响应报文不通过交换机链路, 直接通过其他交换机交付主机。这样就不会出现 arp 广播风暴, 示意图如图所示。

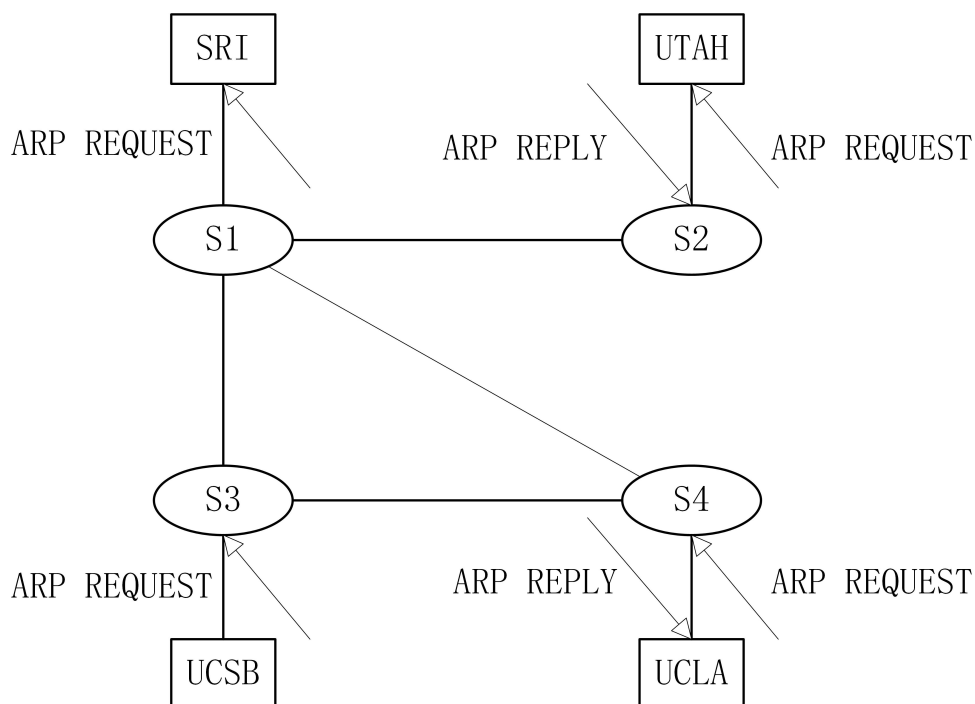


图 9: 解决后流表

2. 关键代码设计思路

在构造函数初始化一个表, 维护一个每个交换机主机和交换机端口的映射。

```
1 self.host_mac_port = {}
```

下面是一个实现上述机制的代码, 首先分析是否为 ARP 包:

(1) 如果是 ARP 包则查询 src_mac 地址是否在该交换机的表中, 如果在, 则说明是正常发送, 则其他交换机转发到连接的所有主机, 如果不在, 则说明不是正常发送, 这时候学习 src_mac 及其对应的交换机端口号, 并且不处理 arp 请求。

```

1  if ARP in header_list:
2      opcode = header_list[ARP].opcode
3      if opcode == arp.ARP_REQUEST:
4          self.host_mac_port.setdefault(dp, {})
5          if src in self.host_mac_port[dp]:
6              for switch in self.host_mac_port.keys():
7                  ofp = switch.ofproto
8                  parser = switch.ofproto_parser
9                  for out_port in self.host_mac_port[switch].values():
10                     actions = [parser.OFPActionOutput(out_port)]
11                     out = parser.OFPPacketOut(datapath=switch,buffer_id=ofp.OFP_NO_BUFFER,
12                     in_port=ofp.OFPP_CONTROLLER,actions=actions, data=msg.data)
13                     switch.send_msg(out)
14             return
15         else:
16             self.host_mac_port[dp][src] = in_port
17             out =
18                 dp.ofproto_parser.OFPPacketOut(datapath=dp,buffer_id=dp.ofproto.OFP_NO_BUFFER,
19                 ,in_port=in_port,actions=[], data=None)
20             dp.send_msg(out)
21             return

```

(2) 如果是 arp 响应报文, 则查询目的 mac 地址是否在表中, 如果在某个交换机的表中, 则某个交换机直接将 arp 响应报文发送给该主机, 实现响应。

```

1      elif opcode == arp.ARP_REPLY:
2          for switch in self.host_mac_port.keys():
3              if dst in self.host_mac_port[switch]:
4                  ofp = switch.ofproto
5                  parser = switch.ofproto_parser
6                  out_port = self.host_mac_port[switch][dst]
7                  actions = [parser.OFPActionOutput(out_port)]
8                  out = parser.OFPPacketOut(datapath=switch,buffer_id=ofp.OFP_NO_BUFFER,
9                  in_port=ofp.OFPP_CONTROLLER,actions=actions, data=msg.data)
10                  switch.send_msg(out)
11          return

```

3. 运行结果与分析

解决 ARP 数据包在环状拓扑中的洪泛问题后, UCLA 和 UTAH 之间可以 ping 通, 并且流表项的匹配次数明显减少, UCSB 监听也没有了广播风暴现象。

```

*** s1 -----
cookie=0x0, duration=26.789s, table=0, n_packets=0, n_bytes=0, priority=1,in_port="s1-eth4",dl_src=76:11:13:7b:3
9:99,dl_dst=46:29:8d:c8:29:25 actions=output:"s1-eth1"
cookie=0x0, duration=26.787s, table=0, n_packets=2, n_bytes=196, priority=1,in_port="s1-eth3",dl_src=76:11:13:7b
:39:99,dl_dst=46:29:8d:c8:29:25 actions=output:"s1-eth1"
cookie=0x0, duration=26.784s, table=0, n_packets=4, n_bytes=336, priority=1,in_port="s1-eth1",dl_src=46:29:8d:c8
:29:25,dl_dst=76:11:13:7b:39:99 actions=output:"s1-eth3"
cookie=0x0, duration=19.620s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s1-eth3",dl_src=9e:87:ec:54
:2c:eb,dl_dst=46:29:8d:c8:29:25 actions=output:"s1-eth1"
cookie=0x0, duration=19.619s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s1-eth1",dl_src=46:29:8d:c8
:29:25,dl_dst=9e:87:ec:54:2c:eb actions=output:"s1-eth3"
cookie=0x0, duration=15.493s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s1-eth2",dl_src=3a:ca:09:19
:b8:17,dl_dst=46:29:8d:c8:29:25 actions=output:"s1-eth1"
cookie=0x0, duration=15.492s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s1-eth1",dl_src=46:29:8d:c8
:29:25,dl_dst=3a:ca:09:19:b8:17 actions=output:"s1-eth2"
cookie=0x0, duration=15.481s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s1-eth2",dl_src=3a:ca:09:19
:b8:17,dl_dst=76:11:13:7b:39:99 actions=output:"s1-eth3"
cookie=0x0, duration=15.471s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s1-eth3",dl_src=76:11:13:7b
:39:99,dl_dst=3a:ca:09:19:b8:17 actions=output:"s1-eth2"
cookie=0x0, duration=15.459s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s1-eth2",dl_src=3a:ca:09:19
:b8:17,dl_dst=9e:87:ec:54:2c:eb actions=output:"s1-eth3"
cookie=0x0, duration=15.455s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s1-eth3",dl_src=9e:87:ec:54
:2c:eb,dl_dst=3a:ca:09:19:b8:17 actions=output:"s1-eth2"
cookie=0x0, duration=42.844s, table=0, n_packets=48, n_bytes=4092, priority=0 actions=CONTROLLER:65535
*** s2 -----
cookie=0x0, duration=15.494s, table=0, n_packets=4, n_bytes=336, priority=1,in_port="s2-eth2",dl_src=46:29:8d:c8
:29:25,dl_dst=3a:ca:09:19:b8:17 actions=output:"s2-eth1"
cookie=0x0, duration=15.473s, table=0, n_packets=3, n_bytes=238, priority=1,in_port="s2-eth2",dl_src=76:11:13:7b

```

图 10: 流表匹配次数减少

所有节点能够成功 ping 通。

```

*** Ping: testing ping reachability
SRI -> UCLA UCSB UTAH
UCLA -> SRI UCSB UTAH
UCSB -> SRI UCLA UTAH
UTAH -> SRI UCLA UCSB
*** Results: 0% dropped (12/12 received)

```

图 11: 能够成功 ping 通

在 UTAH 中抓包, 发现能够正常抓包。

9 4.075725649	10.0.0.2	10.0.0.2	ICMP	98 Echo (ping) request	id=0x13dd, seq=5/1280, ttl=64 (request in 7)
10 4.075737607	10.0.0.4	10.0.0.2	ICMP	98 Echo (ping) reply	id=0x13dd, seq=5/1280, ttl=64 (reply in 10)
11 5.099719240	3a:ca:09:19:b8:17	76:11:13:7b:39:99	ARP	42 Who has 10.0.0.2? Tell 10.0.0.4	
12 5.099815840	10.0.0.2	10.0.0.4	ICMP	98 Echo (ping) request	id=0x13dd, seq=6/1536, ttl=64 (request in 13)
13 5.099823247	10.0.0.4	10.0.0.2	ICMP	98 Echo (ping) reply	id=0x13dd, seq=6/1536, ttl=64 (request in 12)
14 5.100031612	76:11:13:7b:39:99	3a:ca:09:19:b8:17	ARP	42 Who has 10.0.0.4? Tell 10.0.0.2	
15 5.100032865	3a:ca:09:19:b8:17	76:11:13:7b:39:99	ARP	42 10.0.0.4 is at 3a:ca:09:19:b8:17	
16 5.100048642	76:11:13:7b:39:99	3a:ca:09:19:b8:17	ARP	42 10.0.0.2 is at 76:11:13:7b:39:99	
17 6.123791658	10.0.0.2	10.0.0.4	ICMP	98 Echo (ping) request	id=0x13dd, seq=7/1792, ttl=64 (reply in 18)
18 6.123810472	10.0.0.4	10.0.0.2	ICMP	98 Echo (ping) reply	id=0x13dd, seq=7/1792, ttl=64 (request in 17)
19 7.148000952	10.0.0.2	10.0.0.4	ICMP	98 Echo (ping) request	id=0x13dd, seq=8/2048, ttl=64 (reply in 20)
20 7.148011988	10.0.0.4	10.0.0.2	ICMP	98 Echo (ping) reply	id=0x13dd, seq=8/2048, ttl=64 (request in 19)
21 8.171744888	10.0.0.2	10.0.0.4	ICMP	98 Echo (ping) request	id=0x13dd, seq=9/2304, ttl=64 (reply in 22)

图 12: UTAH 抓包

可以在下图抓包发现, 链路上并没有广播 arp 包, 而是直接通过其他交换机广播发给主机, 很好地实现了结果。

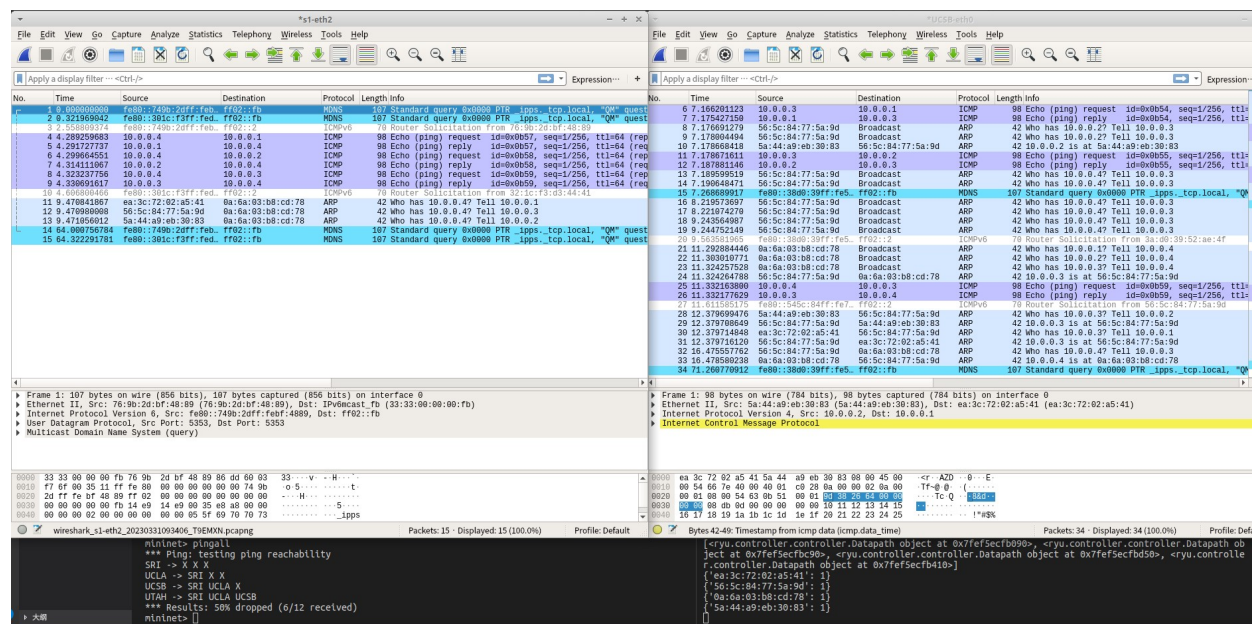


图 13: 交换机链路上并没有广播 arp 包

4. 方案优缺点分析

优点: 直接不占用交换机链路带宽, 可以充分利用交换机的带宽。

缺点: 方法复杂, 不够简洁, 而且需要先 ping 来获取主机与端口绑定信息。没有办法对后期带有 dst 的 arp 请求报文进行处理, 仍然会走交换机链路。

四、 附录

1. Learning_Switch.py

```

1 from ryu.base import app_manager
2 from ryu.controller import ofp_event
3 from ryu.controller.handler import MAIN_DISPATCHER, CONFIG_DISPATCHER
4 from ryu.controller.handler import set_ev_cls
5 from ryu.ofproto import ofproto_v1_3
6 from ryu.lib.packet import packet
7 from ryu.lib.packet import ethernet
8
9
10 class Switch(app_manager.RyuApp):
11     OFP_VERSIONS = [ofproto_v1_3.OFP_VERSION]
12
13
14     def __init__(self, *args, **kwargs):
15         super(Switch, self).__init__(*args, **kwargs)
16
17         # maybe you need a global data structure to save the mapping
18         self.mac_to_port = {}

```

```

18
19     def add_flow(self, datapath, priority, match, actions, idle_timeout=0, hard_timeout=0):
20         dp = datapath
21         ofp = dp.ofproto
22         parser = dp.ofproto_parser
23         inst = [parser.OFPInstructionActions(ofp.OFPIT_APPLY_ACTIONS, actions)]
24         mod = parser.OFPFlowMod(datapath=dp, priority=priority,
25                                 idle_timeout=idle_timeout,
26                                 hard_timeout=hard_timeout,
27                                 match=match, instructions=inst)
28         dp.send_msg(mod)
29
30     @set_ev_cls(ofp_event.EventOFPSwitchFeatures, CONFIG_DISPATCHER)
31     def switch_features_handler(self, ev):
32         msg = ev.msg
33         dp = msg.datapath
34         ofp = dp.ofproto
35         parser = dp.ofproto_parser
36         match = parser.OFPMatch()
37         actions = [parser.OFPActionOutput(ofp.OFPP_CONTROLLER, ofp.OFPCML_NO_BUFFER)]
38         self.add_flow(dp, 0, match, actions)
39
40     @set_ev_cls(ofp_event.EventOFPPacketIn, MAIN_DISPATCHER)
41     def packet_in_handler(self, ev):
42         msg = ev.msg
43         dp = msg.datapath
44         ofp = dp.ofproto
45         parser = dp.ofproto_parser
46
47         # the identity of switch
48         dpid = dp.id
49         self.mac_to_port.setdefault(dpid, {})
50         # the port that receive the packet
51         in_port = msg.match['in_port']
52         pkt = packet.Packet(msg.data)
53         eth_pkt = pkt.get_protocol(ethernet.ethernet)
54         # get the mac
55         dst = eth_pkt.dst
56         src = eth_pkt.src
57         # we can use the logger to print some useful information
58         self.logger.info('packet: %s %s %s %s', dpid, src, dst, in_port)
59         # you need to code here to avoid the direct flooding
60         # having fun
61         # :)
62         # learn src mac -> port
63         self.mac_to_port[dpid][src] = in_port
64         # if match, then send to output
65         if dst in self.mac_to_port[dpid]:
66             out_port = self.mac_to_port[dpid][dst]
67         # if not match, then flood

```

```

68     else:
69         out_port = ofp.OFPP_FLOOD
70         #pass information
71         actions = [parser.OFPACTIONOutput(out_port)]
72
73     if out_port != ofp.OFPP_FLOOD:
74         match = parser.OFPMATCH(in_port=in_port, eth_dst=dst, eth_src=src)
75         self.add_flow(dp, 1, match, actions)
76
77     data = msg.data
78     out = parser.OFPPACKETOut(datapath=dp,buffer_id=msg.buffer_id,in_port=in_port,
79                               actions=actions, data=data)
80     dp.send_msg(out)

```

2. Learning_Switch_Modified.py

```

1  from ryu.base import app_manager
2  from ryu.controller import ofp_event
3  from ryu.controller.handler import MAIN_DISPATCHER, CONFIG_DISPATCHER
4  from ryu.controller.handler import set_ev_cls
5  from ryu.ofproto import ofproto_v1_3
6  from ryu.lib.packet import packet
7  from ryu.lib.packet import ethernet
8
9
10 class Switch(app_manager.RyuApp):
11     OFP_VERSIONS = [ofproto_v1_3.OFP_VERSION]
12
13     def __init__(self, *args, **kwargs):
14         super(Switch, self).__init__(*args, **kwargs)
15         # maybe you need a global data structure to save the mapping
16         self.mac_to_port = {}
17
18     def add_flow(self, datapath, priority, match,
19                 actions,buffer_id=None,idle_timeout=0,hard_timeout=0):
20         dp = datapath
21         ofp = dp.ofproto
22         parser = dp.ofproto_parser
23         inst = [parser.OFPInstructionActions(ofp.OFPIT_APPLY_ACTIONS, actions)]
24         if buffer_id:
25             mod = parser.OFPFlowMod(datapath=datapath,
26                                     priority=priority,buffer_id=buffer_id,idle_timeout=idle_timeout,
27                                     hard_timeout=hard_timeout,match=match,instructions=inst)
28         else:
29             mod = parser.OFPFlowMod(datapath=datapath,
30                                     priority=priority,idle_timeout=idle_timeout,
31                                     hard_timeout=hard_timeout,match=match, instructions=inst)
32         dp.send_msg(mod)

```



```
30
31 @set_ev_cls(ofp_event.EventOFPSwitchFeatures, CONFIG_DISPATCHER)
32 def switch_features_handler(self, ev):
33     msg = ev.msg
34     dp = msg.datapath
35     ofp = dp.ofproto
36     parser = dp.ofproto_parser
37     match = parser.OFPMatch()
38     actions = [parser.OFPActionOutput(ofp.OFPP_CONTROLLER, ofp.OFPCML_NO_BUFFER)]
39     self.add_flow(dp, 0, match, actions)
40
41 @set_ev_cls(ofp_event.EventOFPPacketIn, MAIN_DISPATCHER)
42 def packet_in_handler(self, ev):
43     msg = ev.msg
44     dp = msg.datapath
45     ofp = dp.ofproto
46     parser = dp.ofproto_parser
47
48     # the identity of switch
49     dpid = dp.id
50     self.mac_to_port.setdefault(dpid, {})
51     # the port that receive the packet
52     in_port = msg.match['in_port']
53     pkt = packet.Packet(msg.data)
54     eth_pkt = pkt.get_protocol(ethernet.ethernet)
55     # get the mac
56     dst = eth_pkt.dst
57     src = eth_pkt.src
58     # we can use the logger to print some useful information
59     self.logger.info('packet: %s %s %s %s', dpid, src, dst, in_port)
60     # you need to code here to avoid the direct flooding
61     # having fun
62     # :)
63     # learn src mac -> port
64     self.mac_to_port[dpid][src] = in_port
65     #if match, then send to outport
66     if dst in self.mac_to_port[dpid]:
67         out_port = self.mac_to_port[dpid][dst]
68     #if not match, then flood
69     else:
70         out_port = ofp.OFPP_FLOOD
71     #pass information
72     actions = [parser.OFPActionOutput(out_port)]
73
74     if out_port != ofp.OFPP_FLOOD:
75         match = parser.OFPMatch(in_port=in_port, eth_dst=dst, eth_src=src)
76         if msg.buffer_id != ofp.OFP_NO_BUFFER:
77             self.add_flow(dp, 1, match, actions, msg.buffer_id)
78             return
79         else:
```

```

80         self.add_flow(dp, 1, match, actions)
81
82         data = None
83         if msg.buffer_id == ofp.OFP_NO_BUFFER:
84             data = msg.data
85
86         out = parser.OFPFpacketOut(datapath=dp,buffer_id=msg.buffer_id,in_port=in_port,
87                                     actions=actions, data=data)
87         dp.send_msg(out)

```

3. Broadcast_Loop.py

```

1  from ryu.base import app_manager
2  from ryu.controller import ofp_event
3  from ryu.controller.handler import MAIN_DISPATCHER, CONFIG_DISPATCHER
4  from ryu.controller.handler import set_ev_cls
5  from ryu.ofproto import ofproto_v1_3
6  from ryu.lib.packet import packet
7  from ryu.lib.packet import ethernet
8  from ryu.lib.packet import arp
9  from ryu.lib.packet import ether_types
10
11  ETHERNET = ethernet.ethernet.__name__
12  ETHERNET_MULTICAST = "ff:ff:ff:ff:ff:ff"
13  ARP = arp.arp.__name__
14
15
16  class Switch_Dict(app_manager.RyuApp):
17      OFP_VERSIONS = [ofproto_v1_3.OFP_VERSION]
18
19      def __init__(self, *args, **kwargs):
20          super(Switch_Dict, self).__init__(*args, **kwargs)
21          self.sw = {} #(dpid, src_mac, dst_ip)=>in_port, you may use it in mission 2
22          # maybe you need a global data structure to save the mapping
23          # just data structure in mission 1
24          self.mac_to_port = {}
25
26
27
28      def add_flow(self, datapath, priority, match, actions, idle_timeout=0, hard_timeout=0):
29          dp = datapath
30          ofp = dp.ofproto
31          parser = dp.ofproto_parser
32          inst = [parser.OFPInstructionActions(ofp.OFPIT_APPLY_ACTIONS, actions)]
33          mod = parser.OFPFlowMod(datapath=dp, priority=priority,
34                                  idle_timeout=idle_timeout,
35                                  hard_timeout=hard_timeout,
36                                  match=match, instructions=inst)

```

```

37     dp.send_msg(mod)
38
39     @set_ev_cls(ofp_event.EventOFPSwitchFeatures, CONFIG_DISPATCHER)
40     def switch_features_handler(self, ev):
41         msg = ev.msg
42         dp = msg.datapath
43         ofp = dp.ofproto
44         parser = dp.ofproto_parser
45         match = parser.OFPMatch()
46         actions = [parser.OFPActionOutput(ofp.OFPP_CONTROLLER, ofp.OFPCML_NO_BUFFER)]
47         self.add_flow(dp, 0, match, actions)
48
49     @set_ev_cls(ofp_event.EventOFPPacketIn, MAIN_DISPATCHER)
50     def packet_in_handler(self, ev):
51         msg = ev.msg
52         dp = msg.datapath
53         ofp = dp.ofproto
54         parser = dp.ofproto_parser
55
56         # the identity of switch
57         dpid = dp.id
58         self.mac_to_port.setdefault(dpid, {})
59         # the port that receive the packet
60         in_port = msg.match['in_port']
61         pkt = packet.Packet(msg.data)
62         eth_pkt = pkt.get_protocol(ethernet.ethernet)
63         if eth_pkt.ethertype == ether_types.ETH_TYPE_LLDP:
64             return
65         if eth_pkt.ethertype == ether_types.ETH_TYPE_IPV6:
66             return
67         # get the mac
68         dst = eth_pkt.dst
69         src = eth_pkt.src
70         # get protocols
71         header_list = dict((p.protocol_name, p) for p in pkt.protocols if type(p) != str)
72         if dst == ETHERNET_MULTICAST and ARP in header_list:
73             # you need to code here to avoid broadcast loop to finish mission 2
74             arp_dst_ip = header_list[ARP].dst_ip
75             if (dp.id, src, arp_dst_ip) in self.sw:
76                 if self.sw[(dp.id, src, arp_dst_ip)] != in_port:
77                     out = dp.ofproto_parser.OFPPacketOut(datapath=dp,
78                                                             buffer_id=dp.ofproto.OFP_NO_BUFFER,
79                                                             in_port=in_port, actions=[], data=None)
80                     dp.send_msg(out)
81                 return
82             else:
83                 self.sw[(dp.id, src, arp_dst_ip)] = in_port
84
85         # self-learning
86         # you need to code here to avoid the direct flooding

```

```

87     # having fun
88     # :)
89     # just code in mission 1
90     # learn src mac -> port
91     self.mac_to_port[dpid][src] = in_port
92     #if match, then send to outport
93     if dst in self.mac_to_port[dpid]:
94         out_port = self.mac_to_port[dpid][dst]
95     #if not match, then fllink_list ood
96     else:
97         out_port = ofp.OFPP_FLOOD
98     #pass information
99     actions = [parser.OFPActionOutput(out_port)]
100
101     if out_port != ofp.OFPP_FLOOD:
102         match = parser.OFPMatch(in_port=in_port, eth_dst=dst, eth_src=src)
103         self.add_flow(dp, 1, match, actions)
104
105     data = None
106     if msg.buffer_id == ofp.OFP_NO_BUFFER:
107         data = msg.data
108
109     out = parser.OFPPacketOut(datapath=dp,buffer_id=msg.buffer_id,in_port=in_port,
110                               actions=actions, data=data)
111     dp.send_msg(out)

```

4. Broadcast_Loop2.py

```

1  from ryu.base import app_manager
2  from ryu.controller import ofp_event
3  from ryu.controller.handler import MAIN_DISPATCHER, CONFIG_DISPATCHER
4  from ryu.controller.handler import set_ev_cls
5  from ryu.ofproto import ofproto_v1_3
6  from ryu.lib.packet import packet
7  from ryu.lib.packet import ethernet
8  from ryu.lib.packet import arp
9  from ryu.lib.packet import ether_types
10
11
12  ETHERNET = ethernet.ethernet.__name__
13  ETHERNET_MULTICAST = "ff:ff:ff:ff:ff:ff"
14  ARP = arp.arp.__name__
15
16
17  class Switch_Dict(app_manager.RyuApp):
18      OFP_VERSIONS = [ofproto_v1_3.OFP_VERSION]
19
20      def __init__(self, *args, **kwargs):

```

```
21     super(Switch_Dict, self).__init__(*args, **kwargs)
22     self.mac_to_port = {}
23     self.host_mac_port = {}
24
25
26     def add_flow(self, datapath, priority, match, actions, idle_timeout=0, hard_timeout=0):
27         dp = datapath
28         ofp = dp.ofproto
29         parser = dp.ofproto_parser
30         inst = [parser.OFPInstructionActions(ofp.OFPIT_APPLY_ACTIONS, actions)]
31         mod = parser.OFPFlowMod(datapath=dp, priority=priority,
32                                 idle_timeout=idle_timeout,
33                                 hard_timeout=hard_timeout,
34                                 match=match, instructions=inst)
35         dp.send_msg(mod)
36
37     @set_ev_cls(ofp_event.EventOFPSwitchFeatures, CONFIG_DISPATCHER)
38     def switch_features_handler(self, ev):
39         msg = ev.msg
40         dp = msg.datapath
41         ofp = dp.ofproto
42         parser = dp.ofproto_parser
43         match = parser.OFPMatch()
44         actions = [parser.OFPActionOutput(ofp.OFPP_CONTROLLER, ofp.OFPCML_NO_BUFFER)]
45         self.add_flow(dp, 0, match, actions)
46
47     @set_ev_cls(ofp_event.EventOFPPacketIn, MAIN_DISPATCHER)
48     def packet_in_handler(self, ev):
49         msg = ev.msg
50         dp = msg.datapath
51         ofp = dp.ofproto
52         parser = dp.ofproto_parser
53
54
55         # the identity of switch
56         dpid = dp.id
57         self.mac_to_port.setdefault(dp, {})
58         # the port that receive the packet
59         in_port = msg.match['in_port']
60         pkt = packet.Packet(msg.data)
61         eth_pkt = pkt.get_protocol(ethernet.ethernet)
62         if eth_pkt.ethertype == ether_types.ETH_TYPE_LLDP:
63             return
64         if eth_pkt.ethertype == ether_types.ETH_TYPE_IPV6:
65             return
66         # get the mac
67         dst = eth_pkt.dst
68         src = eth_pkt.src
69
70         # get protocols
```



```

71     header_list = dict((p.protocol_name, p) for p in pkt.protocols if type(p) != str)
72
73     self.mac_to_port[dp][src] = in_port
74     if dst in self.mac_to_port[dp]:
75         out_port = self.mac_to_port[dp][dst]
76     else:
77         out_port = ofp.OFPP_FLOOD
78     actions = [parser.OFPACTIONOutput(out_port)]
79     if out_port != ofp.OFPP_FLOOD:
80         match = parser.OFPMATCH(in_port=in_port, eth_dst=dst, eth_src=src)
81         self.add_flow(dp, 1, match, actions)
82
83     if ARP in header_list:
84         # you need to code here to avoid broadcast loop to finish mission 2
85         opcode = header_list[ARP].opcode
86         if opcode == arp.ARP_REQUEST:
87             self.host_mac_port.setdefault(dp, {})
88             if src in self.host_mac_port[dp]:
89                 for switch in self.host_mac_port.keys():
90                     ofp = switch.ofproto
91                     parser = switch.ofproto_parser
92                     for out_port in self.host_mac_port[switch].values():
93                         #print(out_port)
94                         actions = [parser.OFPACTIONOutput(out_port)]
95                         out =
96                             parser.OFPPACKETOut(datapath=switch,buffer_id=ofp.OFPP_NO_BUFFER,
97                             in_port=ofp.OFPP_CONTROLLER,actions=actions, data=msg.data)
98                         switch.send_msg(out)
99                     return
100             else:
101                 self.host_mac_port[dp][src] = in_port
102                 return
103
104         elif opcode == arp.ARP_REPLY:
105             for switch in self.host_mac_port.keys():
106                 if dst in self.host_mac_port[switch]:
107                     ofp = switch.ofproto
108                     parser = switch.ofproto_parser
109                     out_port = self.host_mac_port[switch][dst]
110                     actions = [parser.OFPACTIONOutput(out_port)]
111                     out = parser.OFPPACKETOut(datapath=switch,buffer_id=ofp.OFPP_NO_BUFFER,
112                     in_port=ofp.OFPP_CONTROLLER,actions=actions, data=msg.data)
113                     switch.send_msg(out)
114                 return
115
116     data = None
117     if msg.buffer_id == ofp.OFPP_NO_BUFFER:
118         data = msg.data
119     out = parser.OFPPACKETOut(datapath=dp,buffer_id=msg.buffer_id,in_port=in_port,
120         actions=actions, data=data)

```

119

```
dp.send_msg(out)
```