

# C++ ONLINE

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TALK:

EXTERNAL POLYMORPHISM  
AND TYPE ERASURE

A VERY USEFUL DANCE OF  
DESIGN PATTERNS

2025

# Redux

# Eduardo Madrid



# Introduction a bit later



# Objective

The background of the slide is a dark blue or black color, overlaid with a grid of thin, bright green lines. The grid consists of both horizontal and vertical lines, creating a pattern of small squares across the entire slide.

**See How Useful  
Type Erasure is**

# This is rather unique to C++

- C++ is perhaps unique in allowing users to implement features intrinsic to programming languages with practical benefits:
  - Better performance
  - Higher modeling power
  - A library that rivals, **successfully**, the very inheritance + `virtual` override mechanism in the language!

# Leading Edge

- We will talk about capabilities we didn't know we have:
  - A reaction might be to believe the new capabilities are “over engineering”, “solutions in search of problems”
  - But we can use them because they don't introduce hard problems, just:
    - More awful compilation errors
    - Still demonstrably better performance and codegen
- Most of this talk is material I've not seen discussed elsewhere



# Substitutability

**If we get substitutability right,  
We've partially succeeded!**

# Substitutability: Telephone

- We want to place a call, but we don't have a telephone:
  - We can ask a friend for a smartphone.
  - We can install an app in our computer to simulate a physical phone: a virtual/app phone
  - We can use an old telephone plugged to a landline
    - Even a rotary!

# Substitutability: Telephone

- We have an **abstract concept** of a phone, and we have **concrete types** of phone (smart, virtual/app, landline) that themselves have their own **more specific** concrete characteristics
- We typically model this with **subclassing**:
  - A “base class” or “interface” with “`virtual`” functions
  - That are “overridden” by derived classes

# Substitutability: Subclassing

- Suitable for runtime only
- It is:
  - A pre-determined need
  - Intrusive: it “intrudes” into the writing of the software component



# Substitutability: Container Iteration

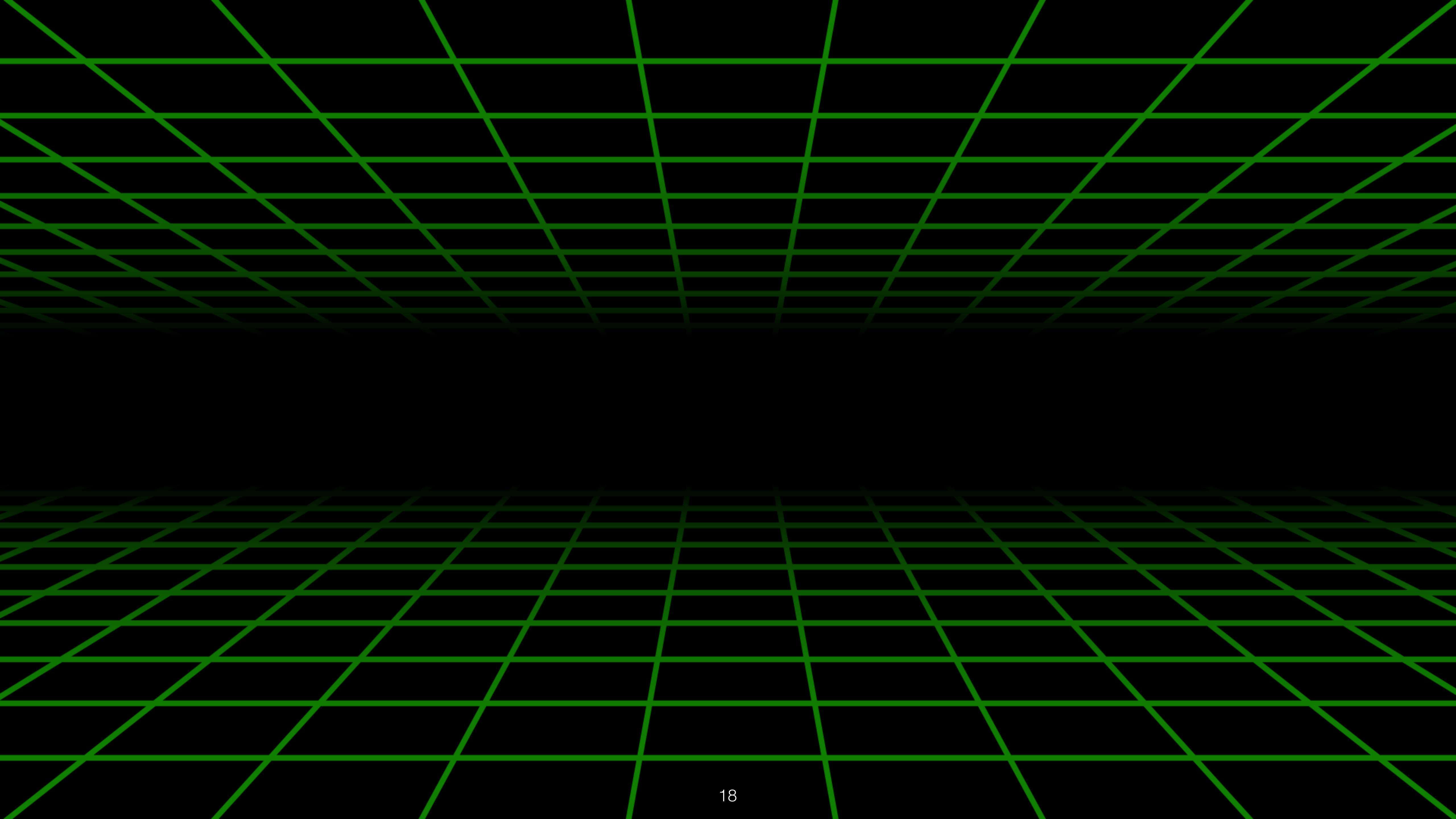
- Iterators:
  1. Equality (`==` and `!=`)
  2. Sentry (`end`)
  3. Dereferencing (`operators * and ->`)
  4. Increments (`++`)
  5. Semantics (not just syntax) of traversal
- This is incredibly general!
- Works great at program writing, a.k.a. “compilation time”

## Substitutability:

**We can substitute a concrete type where an abstract concept is required and things work.**

# That's the famous Liskov's Substitution Principle

This ability to substitute is called  
**polymorphism**





# Pains of Subclassing



# Pains:

Watch Sean Parent's  
"Inheritance Is the Base Class of Evil"

# Subclassing Pains

1. Difficulty to model many subtyping relations:  
Kevlin Henney's "Valued Conversions" [<https://citeseerx.ist.psu.edu/document?repid=rep1&type=pdf&doi=4610004b383e5c4f2dffbea0019c85847e18fff4>]:  
How would you like to pay for that?
  1. Money?
  2. Bartering?
2. Intrusive: we need to wrap perfectly good types to put them in a hierarchy. Busy work. Typical example: wrapping integers in `ISerialize` wrappers.
3. Take it or leave it: A feature of the language you can't finesse.

# Referential Semantics Pains: Their own hell

- One extra indirection
- Allocations: memory fragmentation, synchronization, may fail
- Shallow/deep copy?
- Disables local reasoning
- Incentivizes sharing, and this complicates lifetime management
- Performance hostile in many other ways: no “data” affordances, RTTI

# From Substitutability to C++ templates



We can use **Compile-Time** polymorphism  
for **runtime** polymorphism!

# Important note:

- In our industry, because of the poverty the other programming languages, **substitutability, polymorphism, runtime polymorphism, subclassing and the Liskov's Substitution Principle** are essentially the same thing.
- If this presentation helps you realize there are key differences, the partial success is much better.

# External Polymorphism

- Translate Compile Time Polymorphism to Runtime
  - Basically, a “virtual table”, that seems all you need, except very advanced new possibilities.
- Otherwise: It is a Decorator or Adapter Design Pattern.
- You can use many other Design Patterns, including Strategy

# External Polymorphism

- Example in the Compiler Explorer

# External Polymorphism Demo

- An adapter:
  - It refers to the object somehow (a pointer is good)
  - It gives the runtime polymorphism:
    - Via subclassing! (Nothing bad, the original object types are left undisturbed) see Sean Parent's presentation example.
    - Even better: via the virtual table mechanism.



**External Polymorphism:**  
It's essence is to give runtime polymorphism to types  
that don't have it

Not concerned with runtime  
polymorphism of the **ownership**.

# Type Erasure

- If you own the objects you're given External Polymorphism:
- External Polymorphism with **destruction, moving, and perhaps copying**

# `std::any`

- A container that needs `any_cast` to transform it into something usable

# `std::function`

- Canonical example
- Not intrusive!
- You can have local variables of type `std::function` (including function parameters) as well as members (not forced for them to be pointers or references)
- If it allocates, it is a fallback mechanism, this lessens the problems.

However, it is a **very**  
**bad** design

**Not the fault of the inventors, but  
the fault of our community to  
notice the problems and correct  
them more opportunely**



# `std::function`

- Performs type-erasure, like `std::any`, and a bunch of other things (does not follow the “single responsibility”) principle nor others:
- Without configurability:
  - No way to indicate the **size**, alignment of the local buffer
  - It is copyable, forcing the targets to need to be copyable: Hostile to move semantics, therefore it is hostile to the strong exception guarantee
- Throws an exception when misused: See John Lakos on `std::vector::at`

# std::function

- Supports **only one anonymous** call interface
  - Not a data member: Indirect function call penalty
- Annoyances like using RTTI and its inefficiencies
- **Not a function but a trampoline!**
  - Example const-call: the trampoline might be “const” and the target non-const, in the same way a pointer might be const and point to non-const
  - “Paternalistic” forwarding of arguments (discussion with colleagues)
- **I need to stop!**

# O'Dwyer's std::function design space

- At <https://quuxplusone.github.io/blog/2019/03/27/design-space-for-std-function/>
- Ownership
- Local Buffer Configuration
  - Disable heap allocation
- Fundamental Affordance set:
  - Move-Only? Not even movable?
- Is the user-specified affordance `const`, `noexcept`?

# *Legitimate* dimensions

- Most of the combinations of choices make sense and have good use cases
- The community drains discussing options that do not have a clear best
- Missing the point “how do we make these choices available to the user”, or perhaps, the implicit assumption is the **belief that it is impossible to make mechanisms that let the users choose.**

Consequences: People  
redesign, **poorly**, and  
implement even **worse**

**Because this is highly  
technical and  
misunderstood**

And it turns out the  
design space is much  
more vast!



**Rather than relying on the Standard Library to supply  
ever more species of type-erasure fishes, we should  
learn how to type-erasure fish ourselves!**

I thought it could not  
be done much better

To understand the  
roadblocks, I proceeded  
from first principles

**My discovery:**

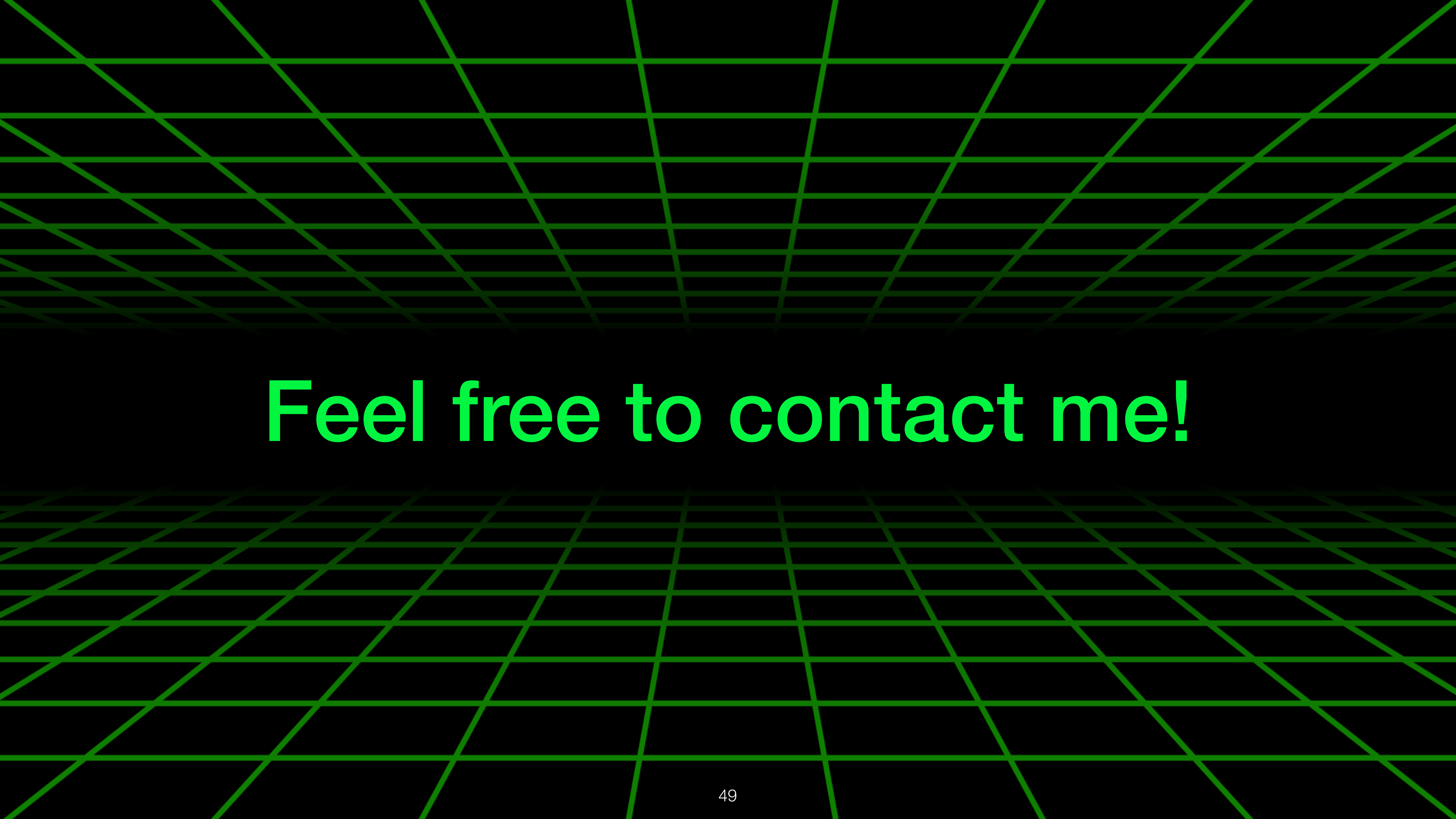
- 1. Subtle wrong assumptions**
- 2. Unclear thinking**

# Conflict of Interest

# Zoo Type Erasure

- Provably optimal performance solution, codegen.
- Partial proof of performance: Fedor Pikus presentation last year “Type Erasure Demystified” [[https://www.youtube.com/watch?v=p-qaf6OS\\_f4](https://www.youtube.com/watch?v=p-qaf6OS_f4)] that explains some of the significant performance improvements that apparently I first identified and articulated in Open Source code.
- I’ve shown the modeling powers of zoo’s type erasure are beyond any other framework
- Hence: **a solution exists!**
- I’d love if you try. Especially if you reject my framework and do it differently! I will try to help you as much as you want. Why? I know there is a lot to be learned, “tip of the iceberg” kind of thing.





**Feel free to contact me!**

# Who Am I?

- The author of the zoo libraries:
  - SWAR
  - Type Erasure
- Before, things like Financial Exchange Connectivity to exchanges such as the CME via “MDP3”, in production, for a Hedge Fund (my CPPCon 2016).
- Things in common of what I share with the public:
  - Production code, in very demanding scenarios
  - Out of the ordinary results.
  - Examples:
    - SWAR: beats, objectively, highly optimized code such as GLIBC’s.
    - Type Erasure: Used by Snap in Snapchat, at the critical places where `std::function` was used. Surfaced subtle errors in pre-existing code
    - Financial market data connectivity: Foundational work for other people who has also presented at these conferences.

**Type Erasure is  
“internal-external polymorphism”**

# Internal-External Polymorphism

- I tweaked the nomenclature to arrive to a deliberate contradiction:
  - External Polymorphism is an “stand-offish” way to give polymorphism to things that don’t have it.
  - Type Erasure does that and also takes complete control of the target:
    - **VALUE SEMANTICS!**
      - Even if underneath there might be a reference!
- Emergence of complexity: unpredictably interesting and useful behaviors of things in between contradicting design goals.

# Internal-External Polymorphism

- Example: a “Value Manager” that is neither a local buffer (also called “small buffer optimization, SBO”) nor a simple heap pointer but an *opt-in* value manager made by an *user* of the framework, so the pointer is a shared pointer.

# Internal-External Polymorphism

- Much more radical: The given objects to infuse them with runtime polymorphism are never stored, but rather, “scattered” as in “Data Orientation Scattering” into collections of homogeneous data types.
- We get rid of the types of given objects, and represent them *internally* in radically different ways—**while still preserving all of the runtime polymorphic interface!**
  - We have the cake and eat it too!
  - The process of abstraction inherent in the Liskov’s substitution principle reduces unnecessary details and this allows us to get more performance!
  - Negative performance cost abstraction!



**External Polymorphism and the Internal Polymorphism of  
controlling objects interact in very interesting and useful  
ways: The Dance.**

# Outgoing Remarks

- If you feel the advanced capabilities of Type Erasure are “over engineering”, I have good news: I think you got important parts of this presentation, because at least you know new things, those you are skeptical about. I just hope that by the time your imagination catches up to the capabilities, you still remember enough of them.



**END!**