Lecture 42 - OS Security

CprE 308

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Security Countermeasures

What do we really need?

- From user perspective
- From process/thread perspective
- From file/directory/file system perspective
- From memory management and other I/O device perspective
- From service perspective
- From network perspective
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What we need in terms of security

- Authentication
 - Username/password
 - One-time password
 - Smartcards/Activebadge
 - Biometrics
- Access Control
 - User-based
 - Role-based
 - Location-based
 - Separation/Interaction, Multi-level Security
- Data Confidentiality & Integrity
 - Encrypted file
 - Encrypted file system
- Service/system availability/reliability

Access Control Fundamentals

- Lampson's Access Matrix
- Reference Monitor
- A secure OS is the one that satisfies:
 - Complete Mediation: TOCTTOU (Time-of-check-to-time-of-use)
 - Tamperproof
 - Verifiable
- Assessment Criteria

Verifiable Security Goals

- Information Flow
- IF Secrecy
 - Denning's Lattice Model
 - Bell-LaPadula Model
- IF Integrity
 - Biba Integrity Model
 - Low-water Mark Integrity
 - Clar-Wilson Integrity
- Covert Channels

History of Secure OSes

- Multics
- UNIX/Windows Security
- Security Kernels/TCB/SELinux
- Microkernels/MicroVM
- TPM
- System Assurance
 - Orange BOok
 - Common Criteria

Cryptography

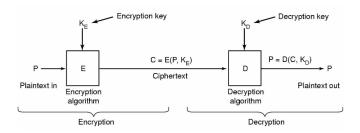
Crypto Building Blocks

- Data Confidentiality & Integrity
 - Encrypted file
 - Encrypted file system
- Hash Function

Cryptography

- Concerns the algorithms in security systems
 - Encryption, digital signatures, etc.
- Cryptography insecure system insecure
- Cryptography secure system might still be insecure!

Encryption and Decryption



Relationship between the plaintext and the ciphertext

Secret-key and Public-key

- Secret keys
 - Encryption and decryption keys are private
 - Need a way to transfer these keys securely one time
- Public Keys
 - Encryption key is common knowledge (public)
 - Decryption key is private

(Symmetric) Secret Key Cryptography

- Same (private) key used to encrypt and ecrypt
 - hence "symmetric"
- Encryption and Decryption functions public
- Requirement: Need a way to securely pass the key to the receiver

Secret Key Example: Substitution Cipher

```
A - X
B - F
```

C - G

••

..

Z - Y

Substitution Ciphers can be (easily) broken

- Plain Text: "Meet at four pm tomorrow"
- Cipher Text: "Tddi xi hosa ..."
- 26 char Encryption Key: "XFG...Y"
- If the encryption key is known, decryption is easy.

Secret Key Cryptography

- Substitution Ciphers are not secure enough, but
- Commonly used schemes:
 - Block Ciphers (work on blocks of data rather than individual characters)
 - Stream Ciphers (use the position of each character in the stream)
- Requirement: Need a way to securely transmit the key to receiver

(Asymmetric) Public Key Cryptography

- Different keys for encryption and decryption, hence "Asymmetric"
 - Encryption key is public
 - Decryption key is private
- If Alice wants to send a secure message to Bob:
 - Alice encrypts message using Bob's public key
 - Bob decrypts using his private key
 - An eavesdropper (Eve) can't make sense out of the message unless she knows Bob's private key
- Diffie and Hellman, 1976

Public Key Cryptography

- Keys are generated as a pair (P,S) for each user
 - P = public, S = secret key
 - Usually inverses of each other
 - \blacksquare S(P(message)) = message
 - ightharpoonup P(S(message)) = message
- No need for sender and receiver to agree on keys beforehand
- Public key dissemination through email, webpage, phonebooks(?)

RSA Cryptosystem

- Rivest, Shamir, and Adleman in 1977
- Rests upon the difficulty of factoring large numbers (which are the product of two primes)
- If factoring is easy, then RSA can be broken
- Converse is not proven!



Case Studies

- UNIX Password
- Unix/Linux Access Control
 - Users and groups
 - File systems controls
- (HW) Windows NT/XP Security Executive
 - Access Tokens
 - Security descriptors
 - ACLs
- (HW) Windows Vista
 - Security additions

Unix Reading Material

- Man pages
 - groups, newgroup
 - chmod, chown, ghgrp
- Unix and Security: The Influences of History
 - ftp://coast.cs.purdue.edu/pub/doc/misc/spaf-influences-ofhistory.ps.Z

Basic Unix Security Model

- User authenticated on logon
 - User ID associated with process
 - Default Group ID associated with process
 - Default Process listed in passwd file
- Groups defined in /etc/groups
 - Set of users listed with each group definition
 - User can be member of multiple groups

Passwords in UNIX

- Login: bpeck
- Password: cpre308
- How does the system check if the password is correct?
- One solution:
 - Password has (username, password) pairs
 - Store [bpeck, cpre308] in /etc/passwd
 - Password file readable only by privileged user
- Privileged users can get your password
 - Why is this a problem?

Solution: One-Way Functions

- f(x) is easy to compute
- $f^{-1}(x)$ is extremely difficult, if not impossible, to compute
- Unix password file contains image of each password
 - /etc/password contains bpeck:y
 - bpeck logs in, supplies x
 - if f(x) == y, then okay
- Password file can now be world-readable

Dictionary Attack (Morris and Thompson)

- For all words in dictionary, compute f(word)
- Find word such that f(word) == y
- Many users use simple passwords
- Systems that employ just one-way functions to protect their passwords are vulnerable to dictionary attacks

Counterattack

Salt

- for each password, create random "salt" value
- \blacksquare /etc/passwd contains f(append(word, salt)), salt)
- 12-bit salt values in Unix
- attacker must do dictionary attack 4096 times, for each salt value
 - done . . .
 - Feldmeier and Karn produced list of 732,000 most common passwords concatenated with each of 4096 salt values: covers "30% of all passwords

Shadow Files

- /etc/passwd and /etc/group must be readable by everyone
- Both files contain crypt'ed passwords
 - Access enable offline attacks
- Add shadow versions of each file
 - Password obscured in passwords and group
 - Stored in more restricted shadow version of these files

Authentication

- Username/password
- One-time password
- Smartcards/Activebadge
- Biomatrics

Access Control

Unix Access Control

- Three permission octets associated with each file and directory
 - Owner, group, and other
 - Read, write, execute
- For each file/directory
 - Can specify RWX permission sfor one owner, one group, and one other

Directory Permissions

- To list the contents of a directory (do an ls), do we need:
 - read permission?
 - execute permission?
 - both?
- To create/delete files?
- To change directory to /home/abc, what permission do we need on /home?
- To rename files mv /home/abc/f1 /home/abc/f2, what permissions do we need on:
 - /home/abc,
 - /home/abc/f1

Unix Access Check

- First test effective user ID against owner
 - If match, then use owner rights
- Then test all groups user is a member of against group
 - If match, then use group rights
- Otherwise, use other rights
- Can view as rwx, or a value from 0-7
 - E.g. rx = 5, and rw = 6

Constraining Control of New Objects

- Umask can be set to constrain allowed access on new objects created by user
- Expressed as a 3 octet mask
 - E.g. 0022
- Inverse of umask anded by requested access for new object
 - E.g. open requests 0666 (read and write for all)
 - 0666 & ~0022 = 0666 & 755 = 644

Other Bits

- Set UID and Set GUID bits
 - When set, the process created by executing file takes on user ID or group ID associated with file
- Sticky bit
 - On directories, prevents anyone but owner of file removing file in directory

File System Extensions

- Ext2 extra attributes
 - a: append only
 - c: compressed
 - s: secure deletion
 - u: undeletable
 - i: immutable

Access Control

- User-based
- Role-based
- Location-based
- Separation/Interaction, Multi-level Security

Unix Security Problems

- Created as a subset of more complete Multics model
 - Expedient at this time
 - Limits modern expressibility
- Security evolved over 30 years
 - Inconsistencies
- Early evolution occurred in open university environments
 - Encourages bad habits