

Lecture 17

CprE 308

February 19, 2013

Intro

Review

- Producer Consumer using Semaphores
- Condition Variable

Today's Topics

- Readers/Writer Problem

Review

Review: Producer Consumer using Semaphores

Shared Variables

- count (number of items in buffer)
- buffer
- N (maximum size of buffer)

Semaphores

- Empty - semaphore initialized to N (number of free slots in buffer)
- Full - semaphore initialized to zero (number of items in buffer)

Review: Producer Consumer using Semaphores (Example)

Producer

```
while(TRUE) {  
    item = produce();  
    down(Empty);  
    lock(mutex);  
    insert(item,buffer);  
    count++;  
    unlock(mutex);  
    up(Full);  
}
```

Consumer

```
while(TRUE) {  
    down(Full);  
    lock(mutex);  
    item = remove(buffer);  
    count--;  
    unlock(mutex);  
    up(Empty);  
    consume(item);  
}
```

Review: Condition Variables

- `pthread_cond_t` condition_variable
- `pthread_mutex_t` mutex;

Waiting Thread

```
pthread_mutex_lock(&mutex);  
while(!cond. satisfied) {  
    pthread_cond_wait(  
        &condition_variable,  
        &mutex);  
}  
pthread_mutex_unlock(  
    &mutex);
```

Signaling Thread

```
pthread_mutex_lock(&mutex);  
/* change variable value */  
if(cond. satisfied) {  
    pthread_cond_signal(  
        &condition_variable);  
}  
pthread_mutex_unlock(  
    &mutex);
```


Review: Solved using condition variables

Global

```
int thread1_done = 0;  
pthread_cond_t cv;  
pthread_mutex_t mutex;
```

Review: Solved using condition variables

Global

```
int thread1_done = 0;
pthread_cond_t cv;
pthread_mutex_t mutex;
```

Thread 1

```
printf("hello");
pthread_mutex_lock(&mutex);
thread1_done = 1;
pthread_cond_signal(&cv);
pthread_mutex_unlock(
    &mutex);
```

Review: Solved using condition variables

Global

```
int thread1_done = 0;
pthread_cond_t cv;
pthread_mutex_t mutex;
```

Thread 1

```
printf("hello");
pthread_mutex_lock(&mutex);
thread1_done = 1;
pthread_cond_signal(&cv);
pthread_mutex_unlock(
    &mutex);
```

Thread 2

```
pthread_mutex_lock(&mutex);
while(thread1_done == 0) {
    pthread_cond_wait(
        &cv, &mutex);
}
printf(" world\n");
pthread_mutex_unlock(
    &mutex);
```

Reader/Writers Problem

Readers/Writers Problem

- Multiple threads reading/writing a database
- Many threads can read simultaneously
- Only one can be writing at any time
 - When a writer is executing, nobody else can read or write

Readers-Writers Problem

- Database has multiple threads operating
 - Many threads can read simultaneously
 - Only one can be writing at a time
 - When a writer is executing, nobody else can read or write
- Example: Multithreaded web server cache
- Performance Problems with naive solution

Solution Idea

- Readers:
 - First reader locks the database
 - If a reader inside, other readers enter without locking again
 - Checking for readers occurs within a mutex
- Writer:
 - Always lock database before entering

Readers-Writers

Reader

```
while(1) {  
    down(&protector)  
    rc++;  
    if(rc == 1) // first  
        down(&database);  
    up(&protector);  
    read_data();  
    down(&protector);  
    rc--;  
    if(rc == 0) then // last  
        up(&database);  
    up(&protector);  
}
```

Initialized Variables

- Two semaphores:
 - Database (init 1)
 - Protector (init 1)
- $rc = 0$

Readers-Writers

Reader

```
while(1) {  
    down(&protector)  
    rc++;  
    if(rc == 1) // first  
        down(&database);  
    up(&protector);  
    read_data();  
    down(&protector);  
    rc--;  
    if(rc == 0) then // last  
        up(&database);  
    up(&protector);  
}
```

Initialized Variables

- Two semaphores:
 - Database (init 1)
 - Protector (init 1)
- $rc = 0$

Writer

```
while(1) {  
    generate_data();  
    down(&database);  
    write_data();  
    up(&database);  
}
```

Writer Starvation

- Readers might continuously enter while a writer waits
- Writer Priority Solution?
 - What does it mean to give writer priority?
- What is a fair solution?
 - Give a specification

Message Passing

Message Passing

- No shared variables
- Two primitives:
 - `send(destination,message)`
 - `receive(destination,message)`
 - Usually blocks till a message arrives
 - Non-blocking version also usually available

Issues

- Across different machines, message passing is the real thing
- Many issues to consider:
 - Marshaling data into messages
 - Provide reliable transmission across unreliable links?
 - Event-driven mode of programming
- Computer Networking: deals with sending messages across machines