

Q 1

$E = \{ w \in \Sigma^* \mid w \text{ begins with } b \text{ \& ends with } a \}$

let the initial state = q_0

the dead state = q_d

define a DFA $M = (Q, \Sigma, \delta, q_0, F)$

where : $Q = \{ q_0, q_d, q_b, q_{ba} \}$

$\Sigma = \{ a, b \}$

$\delta: Q \times \Sigma \rightarrow Q$

Input state	a	b
q_0	q_d	q_b
q_d	q_d	q_d
q_b	q_{ba}	q_b
q_{ba}	q_{ba}	q_b

equally : $\delta(q_0, a) = q_d, \delta(q_0, b) = q_b$

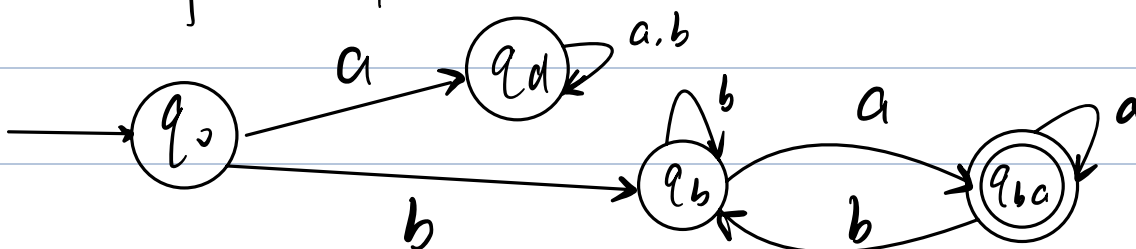
$\delta(q_d, a) = \delta(q_d, b) = q_d$

$\delta(q_b, a) = q_{ba}, \delta(q_b, b) = q_b$

$\delta(q_{ba}, a) = q_{ba}, \delta(q_{ba}, b) = q_b$

$F = \{ q_{ba} \}$

Drawing Description =



Q2

state \ input	0	1
q_1	$\{q_1\}$	$\{q_2\}$
q_2	$\{q_2, q_3\}$	$\{q_2\}$
q_3	$\{q_3\}$	$\{q_2, q_3\}$

by subset construction :

state \ input	0	1
$\{q_1\}$	$\{q_1\}$	$\{q_2\}$
$\{q_2\}$	$\{q_2, q_3\}$	$\{q_2\}$
$\{q_2, q_3\}$	$\{q_2, q_3\}$	$\{q_2, q_3\}$

so the new initial state : $A = \{q_1\}$

$B = \{q_2\}$

the new accept state : $C = \{q_2, q_3\}$

the new DFA : $M = \{Q, \Sigma, \delta, q_0, F\}$

$Q = \{A, B, C\}$

$\Sigma = \{0, 1\}$

$\delta = \delta(A, 0) = A, \delta(A, 1) = B$

$\delta(B, 0) = C, \delta(B, 1) = B$

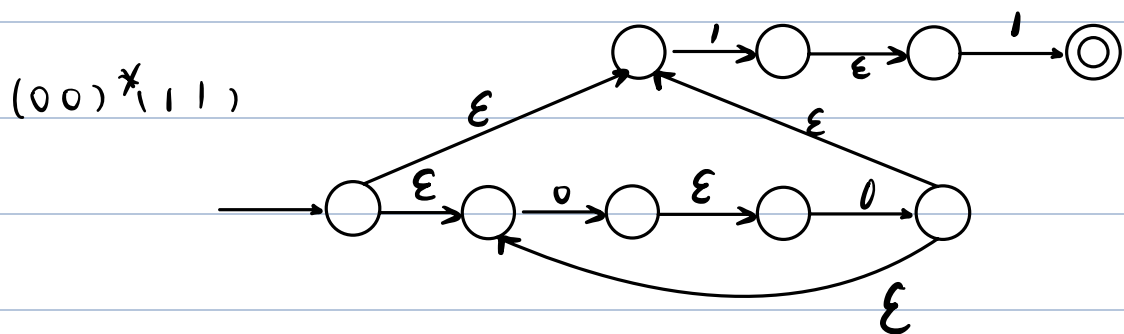
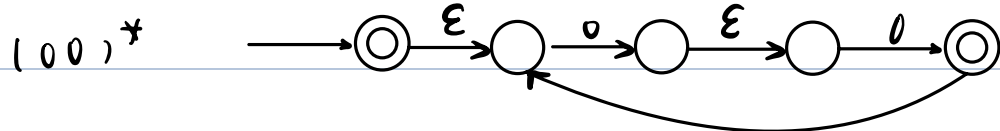
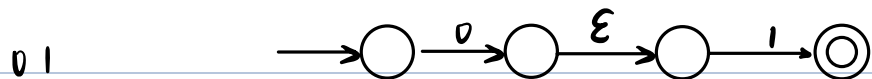
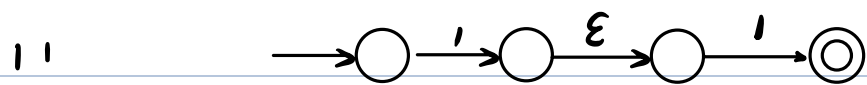
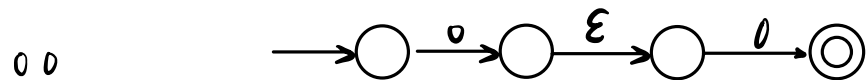
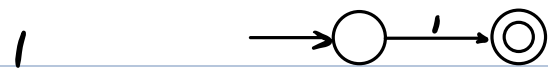
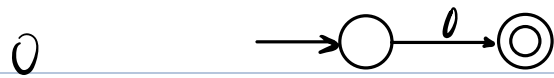
$\delta(C, 0) = C, \delta(C, 1) = C$

$q_0 = A$

$F = C$

The language accepted by the NFA consists of strings Containing
0 1

Q 3



Q4

Assume A_1 is regular

By pumping Lemma:

there's a pumping length $p \geq 1$ such that every string $s \in A_1$ with $|s| \geq p$ can be written $s = xyz$

$$|xy| \leq p$$

$$|y| > 0$$

for all $i \geq 0$ $xy^iz \in A_1$

$$\text{Let } w = a^p b^p$$

$$s = www = (a^p b^p)(a^p b^p)(a^p b^p)$$

$$\text{So } |s| = 6p \geq p \text{ and } s \in A_1$$

$$\therefore s = (a^p b^p)(a^p b^p)(a^p b^p)$$

\therefore the first p symbols of s are all "a"

$$\therefore |xy| \leq p$$

\therefore x and y are both contained in first a^p

$$\text{So let } y = a^t \quad (t > 0)$$

$$x = a^{p-t}$$

according to pumping Lemma:

for all $i \geq 0$ $xy^iz \in A_1$

Let $i = 0$

$$\text{So } xz = a^{p-t} b^p a^p b^p a^p b^p$$

We suppose $xz \in A$

$$|x^i y^j| = |s| - |y| = 6p - t$$

if $x^i y^j \in A$

So $|x^i y^j|$ should be the multiple of 3
 1° t isn't the multiple of 3:

$|x^i y^j|$ can't be divided into 3 same parts

So it contradicts to for all $i \geq 0$, $x y^i \in A$

2° t is the multiple of 3

$|x^i y^j|$ can be divided into 3 parts

in the first block $a^{p-t} b^t$

the first b appears after $p-t$ "a"

in the second block $a^p b$

the first b appears after p "a"

because $t > 0$

So $p-t \neq p$

they're different strings

Hence: $x^i y^j$ can not be uuu

This contradicted with pumping lemma which predicted
 $x y^i \in A$, so A is not regular