

Int 201: Decision Computation and Language

Tutorial 7

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Question 1. Draw PDA for language $\{ww^R | w \in \{0,1\}^*\}$, where w^R is reversed of w , and show the language and $L(\text{PDA})$ (strings being accepted by PDA) are equivalent. ¹

Question 2. Can you draw PDA for language $\{ww | w \in \{0,1\}^*\}$? Can a PDA with two stacks recognize the language $\{ww | w \in \{0,1\}^*\}$? If the answer is yes for any of the questions, draw the PDA, no need for a proof. If the answer is no, give some intuitive explanation. (proof is better, but not necessary).

Question 3. Draw PDA P for language $L = \{a^i b^j c^k | i, j, k \geq 0, i = j \text{ or } j = k\}$, and show the $L = L(P)$. ²

Question 4. Complete the proof for the Kleene closure property of CFL. In the sense, that the Kleen closure of language and the language being accepted by the constructed CFG grammar is the same set.

¹Sets A and B being equivalent means $\forall x, x \in A \implies x \in B$ and $\forall x, x \in B \implies x \in A$

²Note that this is different from the pda in the lecture

Question 5 (Optional). Given the CFG $G = (V, \Sigma, R, S)$:

- $V = \{S, NP, VP, Det, Nominal, Noun, PP, Preposition, Verb\}$
- $\Sigma = \text{The, spy, saw, cop, with, a, telescope}$
- Rules

$S \rightarrow NP VP$

$NP \rightarrow Det Nominal$

$Nominal \rightarrow Noun \parallel Nominal PP$

$VP \rightarrow VP PP \parallel Verb NP$

$PP \rightarrow Preposition NP$

$Det \rightarrow \text{The} \parallel \text{a}$

$Noun \rightarrow \text{spy} \parallel \text{cop} \parallel \text{telescope}$

$Verb \rightarrow \text{saw}$

$Preposition \rightarrow \text{with}$

Is there a third derivation other than what we found in the lecture for The spy saw a cop with a telescope ?

Question 6 (Optional). Have fun with the notebook. <https://github.com/ND-CSE-30151/spring-2024/blob/main/notes/12-pdas.ipynb>