

Q 1

$E = \{w \in \Sigma^* \mid w \text{ begins with } b \text{ & ends with } a\}$

let the initial state = q_0 .

the dead state = q_d

define a DFA $M = (Q, \Sigma, S, q_0, F)$

where: $Q = \{q_0, q_d, q_b, q_{ba}\}$

$\Sigma = \{a, b\}$

$S: Q \times \Sigma \rightarrow Q$

~~state~~
^{Input}

	a	b
q_0	q_d	q_b
q_d	q_d	q_d
q_b	q_{ba}	q_b
q_{ba}	q_{ba}	q_b

equally = $S(q_0, a) = q_d, S(q_0, b) = q_b$

$S(q_d, a) = S(q_d, b) = q_d$

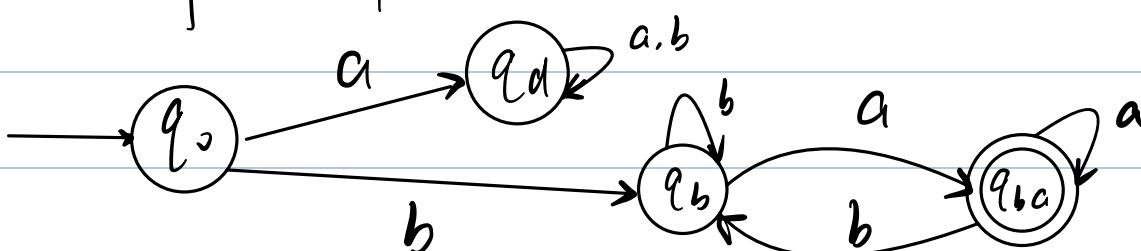
$S(q_b, a) = q_{ba}, S(q_b, b) = q_b$

$S(q_{ba}, a) = q_b, S(q_{ba}, b) = q_b$

$F = \{q_{ba}\}$

..

Drawing Description =



Q2

State Input	0	1
q_1	$\{q_1\}$	$\{q_2\}$
q_2	$\{q_2, q_3\}$	$\{q_2\}$
q_3	$\{q_3\}$	$\{q_2, q_3\}$

by subset construction:

State Input	0	1
$\{q_1\}$	$\{q_1\}$	$\{q_2\}$
$\{q_2\}$	$\{q_2, q_3\}$	$\{q_2\}$
$\{q_2, q_3\}$	$\{q_2, q_3\}$	$\{q_2, q_3\}$

so the new initial state: $A = \{q_1\}$

$$B = \{q_2\}$$

the new accept state: $C = \{q_2, q_3\}$

the new DFA: $M = \{Q, \Sigma, \delta, q_0, F\}$

$$Q = \{A, B, C\}$$

$$\Sigma = \{0, 1\}$$

$$\delta = \delta(A, 0) = A, \delta(A, 1) = B$$

$$\delta(B, 0) = C, \delta(B, 1) = B$$

$$\delta(C, 0) = C, \delta(C, 1) = C$$

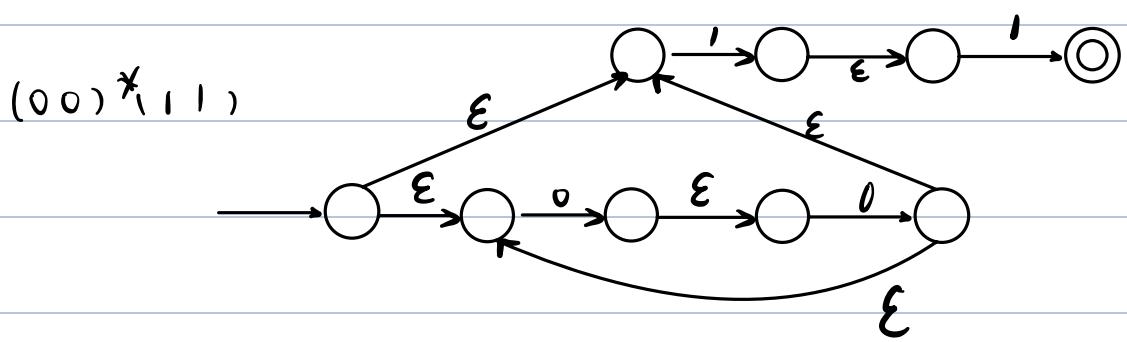
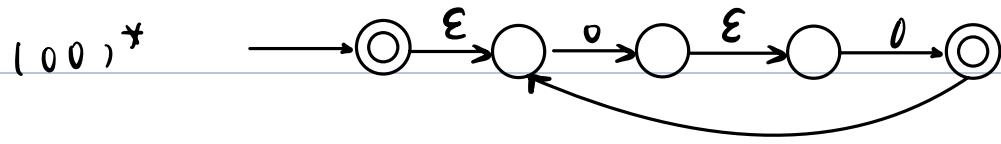
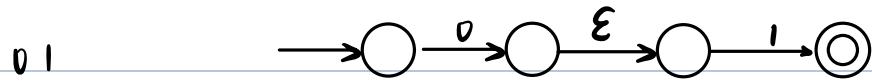
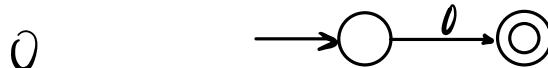
$$q_0 = A$$

$$F = C$$

contains 01

The language accepted by the NFA consists of strings 01

Q 3



Q4

Assume A_1 is regular

By pumping Lemma:

there's a pumping length $p \geq 1$ such that every string $s \in A_1$ with $|s| \geq p$ can be written $s = xy$

$$\begin{aligned} |xy| &\leq p \\ |y| &> 0 \end{aligned}$$

for all $i \geq 0$ $xy^i y \in A_1$

Let $w = a^p b^p$

$$s = www = (a^p b^p) (a^p b^p) (a^p b^p)$$

$$\therefore |s| = 6p \geq p \text{ and } s \in A_1$$

$$\therefore s = (a^p b^p) (a^p b^p) (a^p b^p)$$

\therefore the first p symbols of s are all "a"

$$\therefore |xy| \leq p$$

$\therefore x$ and y are both contained in first a^p

$$\therefore \text{Let } y = a^t \quad (t > 0) \\ x = a^{p-t}$$

according to pumping Lemma:

$$\text{for all } i \geq 0 \quad xy^i y \in A_1$$

$$\text{Let } i = 2$$

$$\therefore xy^2 = a^{p-t} b^p a^p b^p a^p b^p$$

we suppose $xy^2 \notin A_1$

$$|x^3| = |s| - |y| = 6p - t$$

if $x^3 \in A$

so $|x^3|$ should be the multiple of 3
 t is not the multiple of 3:

$|x^3|$ can't be divided into 3 same parts

so if contradict to for all $i \geq 0$, $xy^i z \in A$

t is the multiple of 3

$|x^3|$ can be divided into 3 parts
in the first block $a^{p-t} b^t$

the first b appears after $p-t$ "a"

in the second block $a^p b$

the first b appears after p "a"

because $t > 0$

so $p-t \neq p$

they're different strings

Hence x^3 can not be uuu

This contradicted with pumping Lemma which predicted
 $xy^3 \in A_1$, so A_1 is not regular