RSA Conference 2019

San Francisco | March 4 – 8 | Moscone Center

BETTER.

SESSION ID: HTA-W03

Guardians of the Port: Infinity War

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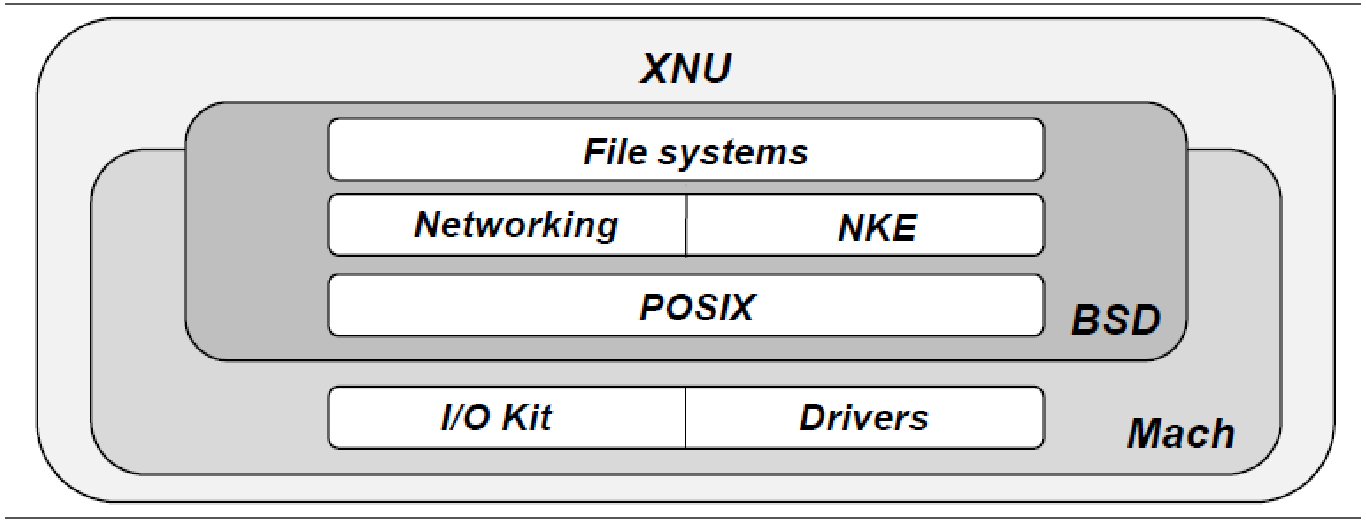
Apple Devices & Jailbreaking

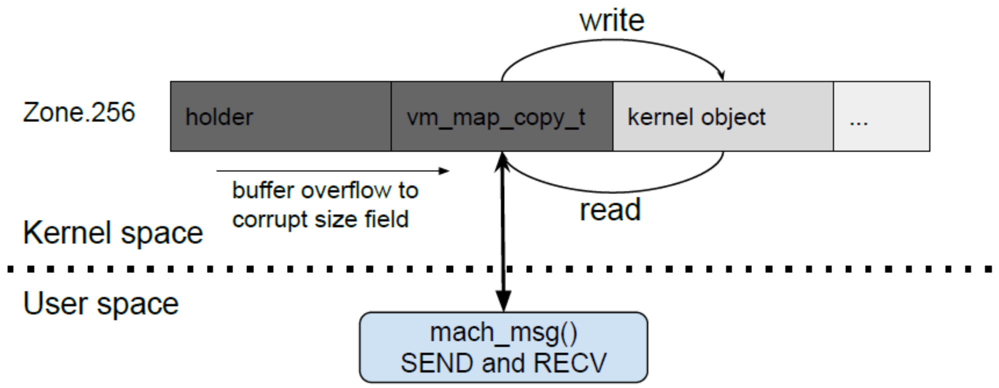
- **Jailbreaking** in general means breaking the device out of its "jail".
- Apple devices (e.g., iPhone, iPad) are most famous "**jail**" devices among the world.
- iOS, macOS, watchOS, and tvOS are **operating systems** developed by Apple Inc and used in Apple devices.

XNU



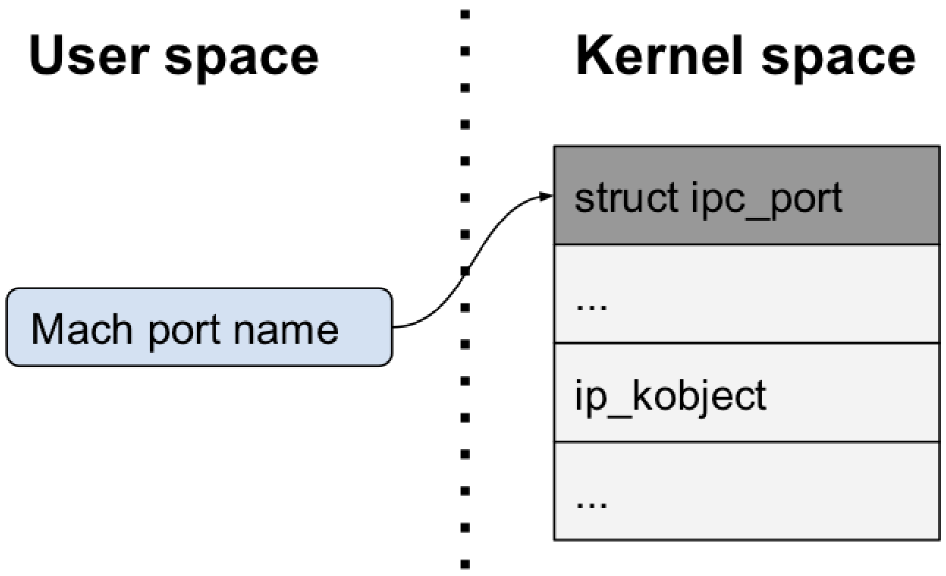
- All systems deploy a same hybrid kernel structure called **XNU**.
- There are cases that **kernel vulnerabilities** have been used to escalate the privileges of attackers and get full control of the system (hence jailbreak the device).
- Accordingly, Apple has deployed multiple **security mechanisms** that make the exploitation of the device harder.

Mitigations in History



- Apple deployed Data Execution Prevention (**DEP**) and Kernel Address Space Layout Randomization (**KASLR**) from iOS 6.
- Apple removed **kdata** field in vm_map_copy_t to prevent attackers abusing mach_msg() in iOS 9.
- Zone element **randomization** was introduced in iOS 9.2.
- To mitigate wrong zone free attacks, Apple added a new **zone_metadata_region** structure for each zone in iOS 10.

New Target - Mach Port in User Space



- A Mach port in XNU is a kernel controlled **communication channel**. It provides basic operations to pass messages between threads.

- **Ports** are used to represent resources, services, and facilities (e.g., hosts, tasks, threads, memory objects, and clocks) thus providing **object-style** access to these abstractions.

- In user space, Mach ports are **integer numbers** like handlers for kernel objects.

New Target – Struct ipc_port in Kernel Space

struct ipc_port	
io_bits	io_references
io_lock_data	
…	
struct ipc_space *receiver;	
ipc_kobject_t ip_kobject;	
…	
mach_vm_address_t ip_context;	
…	

- In the kernel, a Mach port is represented by a pointer to an **ipc_port** structure.
- There are **40** types of ipc_port objects in XNU and **io_bits** field defines the type of it. io_references field counts the reference number of the object. Locking related data is stored in the io_lock_data field.
- Receiver field is a pointer that points to receiver' s **IPC** space (e.g. ipc_space_kernel). **ip_kobject** field points to a kernel data structure according to the kernel object type.

(Mach) Port-oriented Programming (POP)

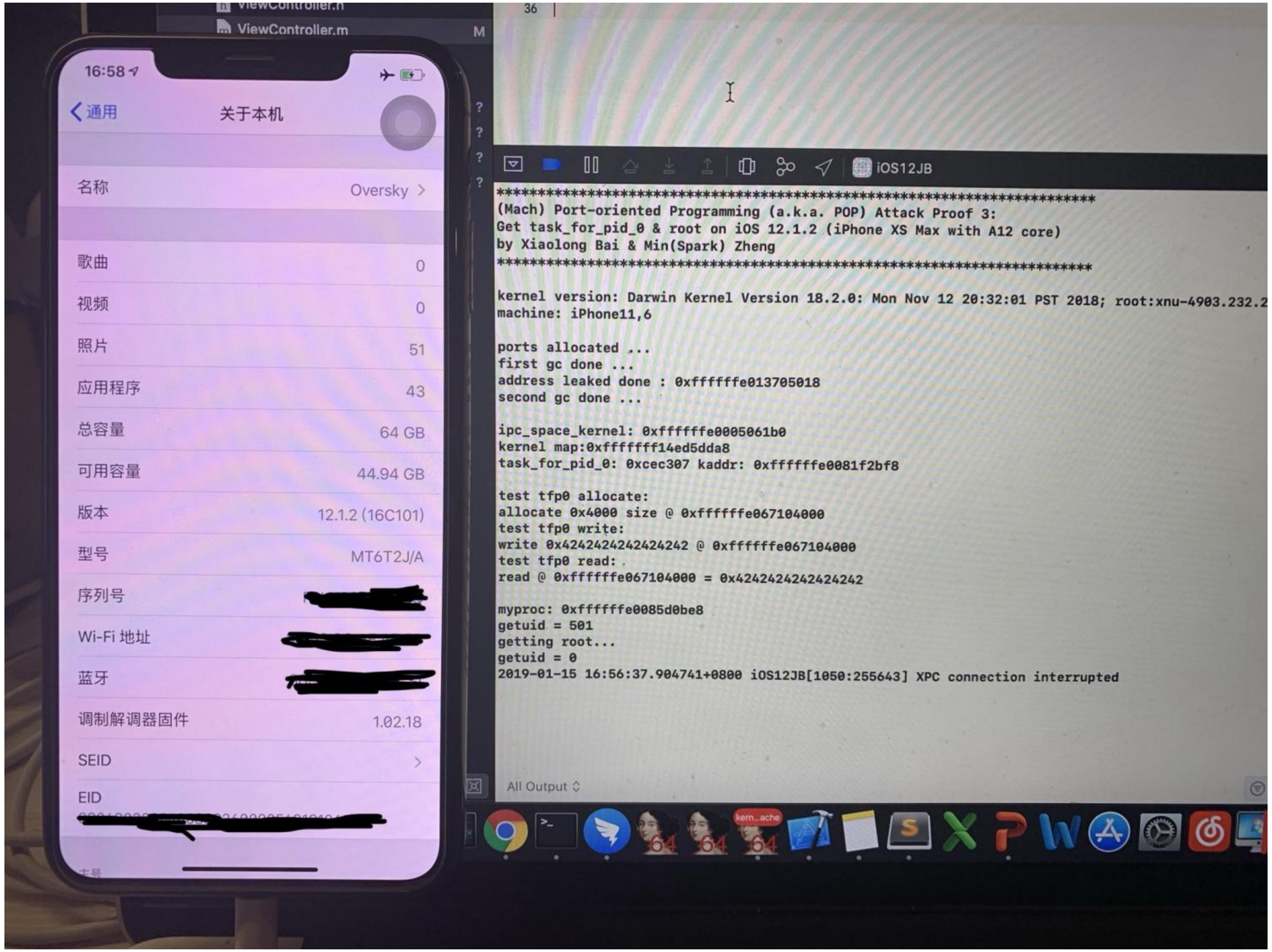
- The main goal is to obtain multiple **primitives** to read/write kernel memory and execute arbitrary kernel code, even in the case that multiple **mitigations** are deployed in the system.

(MACH) PORT-ORIENTED PROGRAMMING

- Attackers leverage a special kernel object, i.e., ipc_port, to obtain multiple primitives by issuing system calls in user mode. Since the proposed method is mainly based on the ipc_port kernel object, we call it **(Mach) Port-oriented Programming (POP)**.

- Note that POP technology was **not created** by **us**. We saw it in many **public** exploits and then **summarize** this code reuse attack technique for systematic study.

POP Attack Proofs (the bug was reported to apple)



How to find POP gadgets? - MIG in Source Code

```
const struct mig_subsystem *mig_e[] = {
        (const struct mig_subsystem *)&mach_vm_subsystem,
        (const struct mig_subsystem *)&mach_port_subsystem,
        (const struct mig_subsystem *)&mach_host_subsystem,
        (const struct mig_subsystem *)&host_priv_subsystem,
        (const struct mig_subsystem *)&host_security_subsystem,
        (const struct mig_subsystem *)&clock_subsystem,
        (const struct mig_subsystem *)&clock_priv_subsystem,
        (const struct mig_subsystem *)&processor_subsystem,
        (const struct mig_subsystem *)&processor_set_subsystem,
        (const struct mig_subsystem *)&is_iokit_subsystem,
    (const struct mig_subsystem *)&lock_set_subsystem,
    (const struct mig_subsystem *)&task_subsystem,
    (const struct mig_subsystem *)&thread_act_subsystem,
#ifdef VM32_SUPPORT
    (const struct mig_subsystem *)&vm32_map_subsystem,
#endif
    (const struct mig_subsystem *)&UNDReply_subsystem,
    (const struct mig_subsystem *)&mach_voucher_subsystem,
    (const struct mig_subsystem *)&mach_voucher_attr_control_subsystem,

#if     XK_PROXY
        (const struct mig_subsystem *)&do_uproxy_xk_uproxy_subsystem,
#endif /* XK_PROXY */
#if     MACH_MACHINE_ROUTINES
        (const struct mig_subsystem *)&MACHINE_SUBSYSTEM,
#endif  /* MACH_MACHINE_ROUTINES */
#if     MCMSG && iPSC860
    (const struct mig_subsystem *)&mcmsg_info_subsystem,
#endif  /* MCMSG && iPSC860 */
};
```

xnu-3248.60.10 〉 osfmk 〉 mach 〉 mach_host.defs 〉 No Selection

```
/*
 *  Return statistics from this host.
 */
routine host_statistics(
        host_priv   : host_t;
        flavor      : host_flavor_t;
    out host_info_out   : host_info_t, CountInOut);

routine host_request_notification(
        host        : host_t;
        notify_type : host_flavor_t;
        notify_port : mach_port_make_send_once_t);
```

xnu-3248.60.10 〉 osfmk 〉 kern 〉 host.c 〉 host_statistics

```
kern_return_t
host_statistics(host_t host, host_flavor_t flavor,
    host_info_t info, mach_msg_type_number_t * count)
{
    uint32_t i;

    if (host == HOST_NULL)
        return (KERN_INVALID_HOST);

    switch (flavor) {
    case HOST_LOAD_INFO: {
        host_load_info_t load_info;

        if (*count < HOST_LOAD_INFO_COUNT)
            return (KERN_FAILURE);
```

How to find POP gadgets? - MIG in Kernel Cache

```
__constdata:FFFFFF8000C5FD00                 public _host_priv_subsystem
__constdata:FFFFFF8000C5FD00 _host_priv_subsystem dq offset _host_priv_server_routine
__constdata:FFFFFF8000C5FD00                                         ; DATA XREF: host priv
__constdata:FFFFFF8000C5FD00                                         ;  host priv server+44↑o
__constdata:FFFFFF8000C5FD08                 db  90h
__constdata:FFFFFF8000C5FD09                 db    1
__constdata:FFFFFF8000C5FD0A                 db    0
__constdata:FFFFFF8000C5FD0B                 db    0
__constdata:FFFFFF8000C5FD0C                 db 0AAh
__constdata:FFFFFF8000C5FD0D                 db    1
__constdata:FFFFFF8000C5FD0E                 db    0
__constdata:FFFFFF8000C5FD0F                 db    0
__constdata:FFFFFF8000C5FD10                 db  34h ; 4
__constdata:FFFFFF8000C5FD11                 db  10h
__constdata:FFFFFF8000C5FD12                 db    0
__constdata:FFFFFF8000C5FD13                 db    0
__constdata:FFFFFF8000C5FD14                 db    0
__constdata:FFFFFF8000C5FD15                 db    0
__constdata:FFFFFF8000C5FD16                 db    0
__constdata:FFFFFF8000C5FD17                 db    0
__constdata:FFFFFF8000C5FD18                 db    0
__constdata:FFFFFF8000C5FD19                 db    0
__constdata:FFFFFF8000C5FD1A                 db    0
__constdata:FFFFFF8000C5FD1B                 db    0
__constdata:FFFFFF8000C5FD1C                 db    0
__constdata:FFFFFF8000C5FD1D                 db    0
__constdata:FFFFFF8000C5FD1E                 db    0
__constdata:FFFFFF8000C5FD1F                 db    0
__constdata:FFFFFF8000C5FD20                 db    0
__constdata:FFFFFF8000C5FD21                 db    0
__constdata:FFFFFF8000C5FD22                 db    0
__constdata:FFFFFF8000C5FD23                 db    0
__constdata:FFFFFF8000C5FD24                 db    0
__constdata:FFFFFF8000C5FD25                 db    0
__constdata:FFFFFF8000C5FD26                 db    0
__constdata:FFFFFF8000C5FD27                 db    0
__constdata:FFFFFF8000C5FD28                 dq offset sub_FFFFFF80002C0DA0
__constdata:FFFFFF8000C5FD30                 db    2
```

```
__int64 __fastcall host_priv_server_routine(__int64 a1)
{
  signed __int64 v1; // rcx
  __int64 result; // rax
  signed __int64 v3; // rcx

  v1 = *(signed int *)(a1 + 28);
  result = 0LL;
  if ( v1 >= 400 )
  {
    v3 = v1 - 400;
    if ( (signed int)v3 <= 25 )
      result = (__int64)*(&host_priv_subsystem + 5 * v3 + 5);
  }
  return result;
}
```

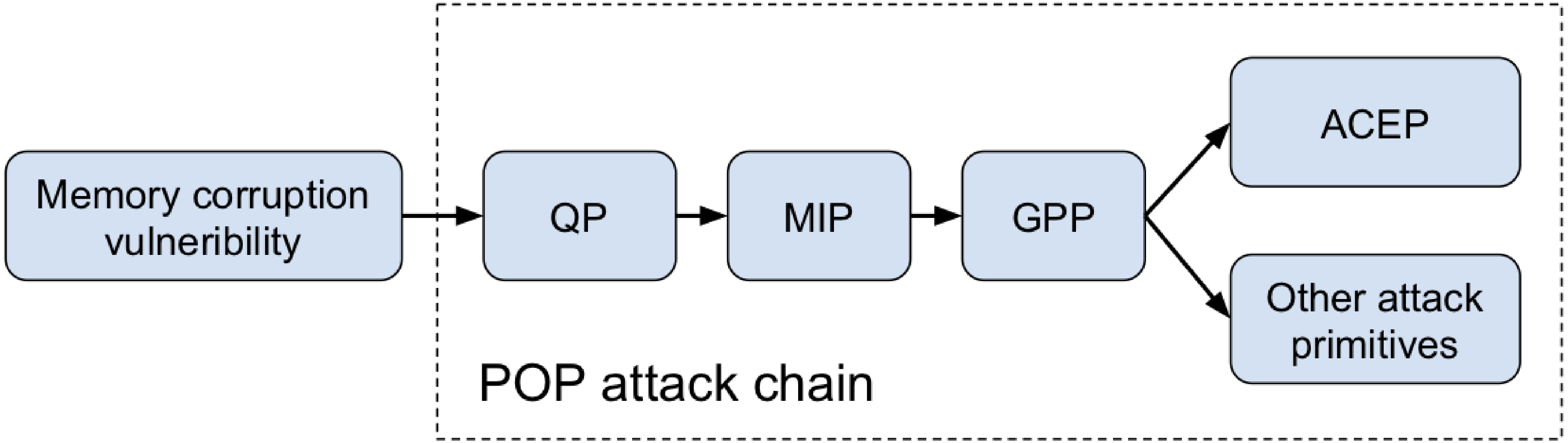
```
char __fastcall sub_FFFFFF80002C0FC0(mach_msg_header_t *a1, mach_msg_header_t *a2)
{
  mach_msg_id_t v2; // ecx
  host_t v3; // eax
  mach_msg_size_t v4; // eax
  mach_msg_size_t v5; // eax
  unsigned __int64 v6; // rax
  __int64 *v7; // rax
  unsigned __int64 v8; // rax
  __int64 *v9; // rax

  if ( kdebug_enable & 1 )
  {
    v8 = __readgsqword(8u);
    if ( v8 )
      v9 = *(__int64 **)(v8 + 976);
    else
      v9 = 0LL;
    sub_FFFFFF80006DFDC0(0LL, 0xFF000649, 0LL, 0LL, 0LL, 0LL, v9);
    if ( (a1->msgh_bits & 0x80000000) != 0 )
      goto LABEL_15;
  }
  else if ( (a1->msgh_bits & 0x80000000) != 0 )
  {
LABEL_15:
    a2[1].msgh_reserved = -304;
    goto LABEL_16;
  }
  if ( a1->msgh_size != 48 )
    goto LABEL_15;
  a2[1].msgh_id = 68;
  v2 = 68;
  if ( a1[1].msgh_id < 0x44u )
    v2 = a1[1].msgh_id;
  a2[1].msgh_id = v2;
  v3 = convert_port_to_host_priv(*(_QWORD *)&a1->msgh_remote_port);
  v4 = host_statistics(v3, a1[1].msgh_reserved, (host_info_t)&a2[2], (mach_msg_type_number_t *)&a2[1].msgh_id);
  a2[1].msgh_reserved = v4;
  if ( v4 )
  {
    a2[1].msgh_reserved = v4;
LABEL_16:
    LOBYTE(v5) = NDR_record.mig_vers;
    *(NDR_record_t *)&a2[1].msgh_remote_port = NDR_record;
    return v5;
  }
  *(NDR_record_t *)&a2[1].msgh_remote_port = NDR_record;
  v5 = 4 * a2[1].msgh_id + 48;
  a2->msgh_size = v5;
  if ( kdebug_enable & 1 )
  {
    v6 = __readgsqword(8u);
    if ( v6 )
      v7 = *(__int64 **)(v6 + 976);
    else
      v7 = 0LL;
    LOBYTE(v5) = sub_FFFFFF80006DFDC0(0LL, 0xFF00064A, 0LL, 0LL, 0LL, 0LL, v7);
  }
  return v5;
}
```

How to find POP gadgets? - Mach Traps

```
const mach_trap_t    mach_trap_table[MACH_TRAP_TABLE_COUNT] = {
/* 0 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 1 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 2 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 3 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 4 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 5 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 6 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 7 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 8 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 9 */      MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 10 */     MACH_TRAP(_kernelrpc_mach_vm_allocate_trap, 4, 5, munge_wwlw),
/* 11 */     MACH_TRAP(_kernelrpc_mach_vm_purgable_control_trap, 4, 5, munge_wlww),
/* 12 */     MACH_TRAP(_kernelrpc_mach_vm_deallocate_trap, 3, 5, munge_wll),
/* 13 */     MACH_TRAP(kern_invalid, 0, 0, NULL),
/* 14 */     MACH_TRAP(_kernelrpc_mach_vm_protect_trap, 5, 7, munge_wllww),
/* 15 */     MACH_TRAP(_kernelrpc_mach_vm_map_trap, 6, 8, munge_wwllww),
/* 16 */     MACH_TRAP(_kernelrpc_mach_port_allocate_trap, 3, 3, munge_www),
/* 17 */     MACH_TRAP(_kernelrpc_mach_port_destroy_trap, 2, 2, munge_ww),
/* 18 */     MACH_TRAP(_kernelrpc_mach_port_deallocate_trap, 2, 2, munge_ww),
/* 19 */     MACH_TRAP(_kernelrpc_mach_port_mod_refs_trap, 4, 4, munge_wwww),
/* 20 */     MACH_TRAP(_kernelrpc_mach_port_move_member_trap, 3, 3, munge_www),
/* 21 */     MACH_TRAP(_kernelrpc_mach_port_insert_right_trap, 4, 4, munge_wwww),
/* 22 */     MACH_TRAP(_kernelrpc_mach_port_insert_member_trap, 3, 3, munge_www),
/* 23 */     MACH_TRAP(_kernelrpc_mach_port_extract_member_trap, 3, 3, munge_www),
/* 24 */     MACH_TRAP(_kernelrpc_mach_port_construct_trap, 4, 5, munge_wwlw),
/* 25 */     MACH_TRAP(_kernelrpc_mach_port_destruct_trap, 4, 5, munge_wwwl),
```

- Most of the Mach traps are **fast-paths** for kernel APIs that are also exposed via the standard MIG kernel APIs.
- Mach traps provide a faster interface to these kernel functions by avoiding the **serialization** and **deserialization** overheads involved in calling kernel MIG APIs.
- **pid_for_task()** is a Mach trap which is heavily used in POP attacks.

POP Attack Chain



- **QP**: Querying primitives
- **MIP**: Memory interoperation primitives
- **GPP**: general purpose primitives
- **ACEP**: Arbitrary code execution primitives

Querying Primitives

```
kern_return_t
mach_port_kobject(
      ipc_space_t                 space,
      mach_port_name_t        name,
      natural_t                    *typep,
      mach_vm_address_t    *addrp)
{
 …

   *typep = (unsigned int) ip_kotype(port);
   *addrp = 0;
   return KERN_SUCCESS;
}
```

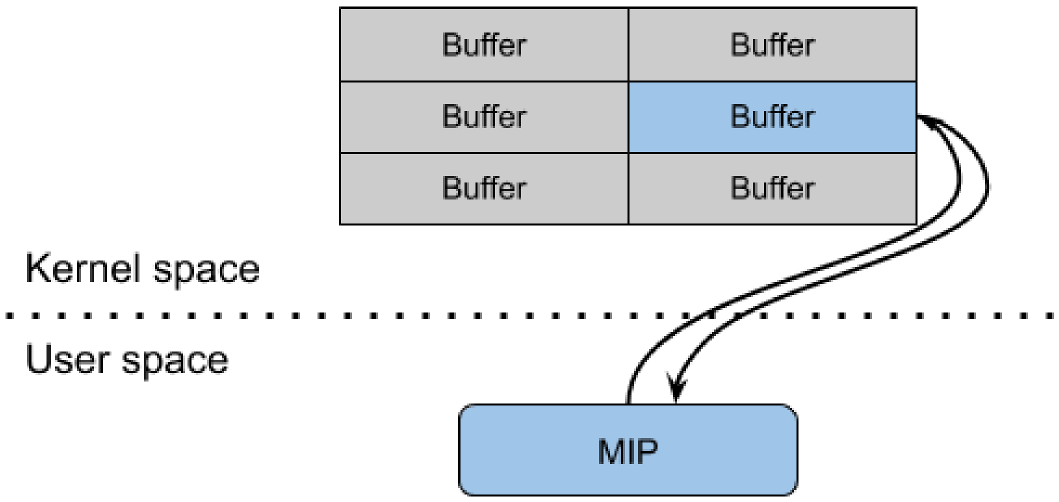
- QP is used to break the **randomization** based mitigations. It helps the attacker to find **corrupted** ports in kernel without **panicking** the system.
- It depends on **corrupted** port objects or **MIP** to change a port's attribute.
- The limitation is the attacker can only gain **limited** information from QP, thus a successful attack usually requires **several** attempts.

Querying Primitives

```
kern_return_t clock_sleep_trap(
struct clock_sleep_trap_args *args)
{
mach_port_name_t clock_name = args->clock_name;
...
if (clock_name == MACH_PORT_NULL)
        clock = &clock_list[SYSTEM_CLOCK];
else
        clock = port_name_to_clock(clock_name);
...
if (clock != &clock_list[SYSTEM_CLOCK])
        return (KERN_FAILURE);
...
return KERN_SUCCESS;
}
```

- clock_sleep_trap() is a system call expecting its first argument (if not NULL) to be a send right to the **global** system clock, and it will return **KERN_SUCCESS** if the port name is correct.
- If the attacker can manipulate an ipc_port kernel object and **change** its **ip_kobject** field, a side channel attack could be launched to break KASLR.

Memory Interoperation Primitives



- MIP can help the attacker to **read/write** the kernel memory in **limited/arbitrary** kernel addresses.
- It depends on corrupted port objects and system calls with **inadequate** type integrity checking.
- Specifically, some MIPs are not used for the **original intention** of the design.

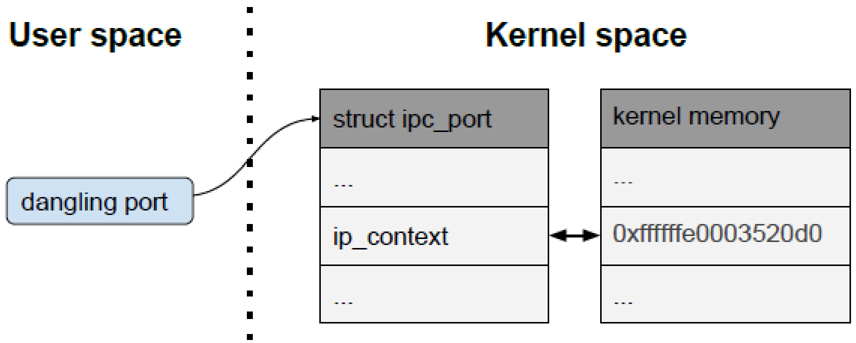
Memory Interoperation Primitives

```
kern_return_t pid_for_task(
struct pid_for_task_args *args)
{
mach_port_name_t t = args->t;
user_addr_t          pid_addr  = args->pid;
...
t1 = port_name_to_task_inspect(t);
...
p = get_bsdtask_info(t1);
if (p) {
        pid  = proc_pid(p);
        err = KERN_SUCCESS;
}
copyout(&pid, pid_addr, sizeof(int));
...
}
```

- **pid_for_task()** is such a system call which returns the **PID** number corresponding to a particular Mach **task**.
- However, the function does not check the validity of the task, and directly returns the value of **task -> bsd_info -> p_pid** to user space after calling get_bsdtask_info(), proc_pid() and copyout().

Memory Interoperation Primitives

```
static kern_return_t
mach_port_guard_locked(
    ipc_port_t       port,
    uint64_t         guard,
    boolean_t        strict)
{
    if (port->ip_context)
        return KERN_INVALID_ARGUMENT;

    port->ip_context = guard;
    port->ip_guarded = 1;
    port->ip_strict_guard = (strict)?1:0;
    return KERN_SUCCESS;
}
```



- A port referring to a freed ipc_port object is called a **dangling port**.
- System calls like mach_port_set/get_*(), mach_port_guard/unguard() are used to write and read the **member fields** of the ipc_port object.
- ip_context field in the ipc_port object is used to associate a userspace pointer with a port. By using mach_port_set/get_context() to a dangling port, the attacker can **retrieve** and **set** 64-bits value in the kernel space.

General Purpose Primitives

Category	Syscall number	Object types
RAW_PORT	36	IKOT_NONE
HOST	52	IKOT_HOST, IKOT_HOST_PRIV, IKOT_HOST_NOTIFY, IKOT_HOST_SEC
PROCESSOR	16	IKOT_PROCESSOR, IKOT_PSET, IKOT_PSET_NAME
TASK	163	IKOT_TASK, IKOT_TASK_NAME, IKOT_TASK_RESUME, IKOT_MEM_OBJ, IKOT_UPL, IKOT_MEM_OBJ_CONTROL, IKOT_NAMED_ENTRY
THREAD	28	IKOT_THREAD
DEVICE	86	IKOT_MASTER_DEVICE, IKOT_IOKIT_SPARE, IKOT_IOKIT_CONNECT
SYNC	29	IKOT_SEMAPHORE, IKOT_LOCK_SET
MACH_VOUCHER	7	IKOT_VOUCHER, IKOT_VOUCHER_ATTR_CONTROL
TIME	10	IKOT_TIMER, IKOT_CLOCK, IKOT_CLOCK_CTRL
MISC	18	IKOT_PAGING_REQUEST, IKOT_MIG, IKOT_XMM_PAGER, IKOT_XMM_KERNEL, IKOT_XMM_REPLY,IKOT_UND_REPLY,IKOT_LEDGER, IKOT_SUBSYSTEM, IKOT_IO_DONE_QUEUE, IKOT_AU_SESSIONPORT, IKOT_FILEPORT
Sum	445	

- GPP is used by the system designer’s **original** intention based on normal user space inputs.
- XNU provides powerful system calls for **privileged** ports (e.g., host_priv and kernel_task). If attackers can get a send right to privileged ports, they can enlarge the attack surface or even control the whole system.

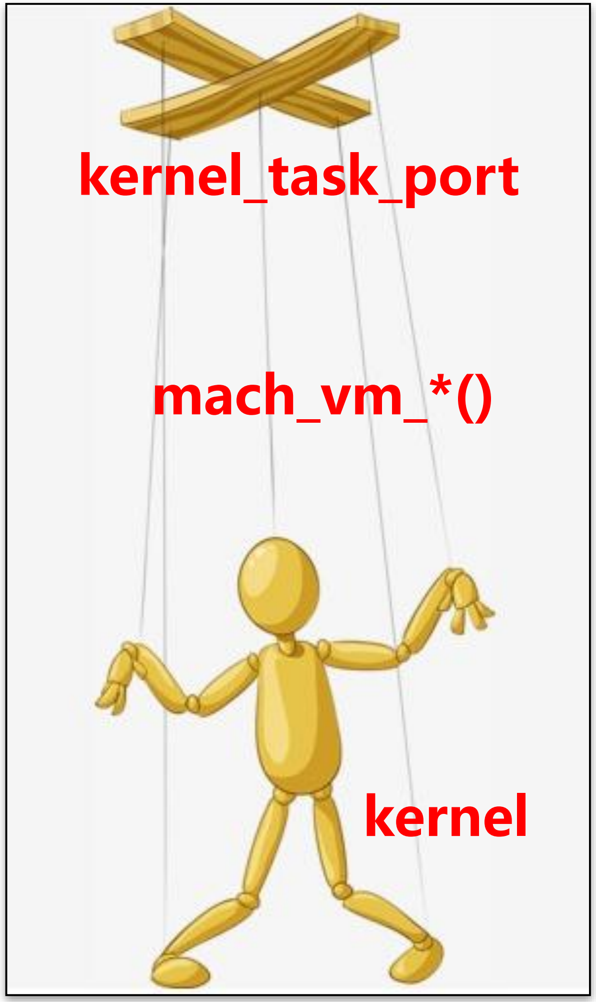
General Purpose Primitives for host_priv

- Mach represents the overall computer system as a **host** object.

- Through **host_*()** system calls, a userspace app can retrieve information (e.g., host_info()) or set properties (e.g.,host_set_multiuser_config_flags()) for a host.

- Moreover, with a send right to **host_priv** port (like **root** user) and related system calls like host_processor_set_priv(), an attacker can gain send rights to other **powerful ports** (e.g.,processor_set port).

General Purpose Primitives for kernel_task



- XNU provides a powerful set of routines, **mach_vm_*()** system calls, to userspace apps for manipulating task **memory** spaces.
- Before iOS 10.3, if the userspace app has a send right to kernel_task, it can manipulate the whole kernel.
- After iOS 10.3, with an information leak or MIP, the attacker could retrieve other tasks' map pointers and receivers. Then attackers can craft **fake** tasks to manage other memory spaces (especially for **kernel**'s memory space).

Arbitrary Code Execution Primitives

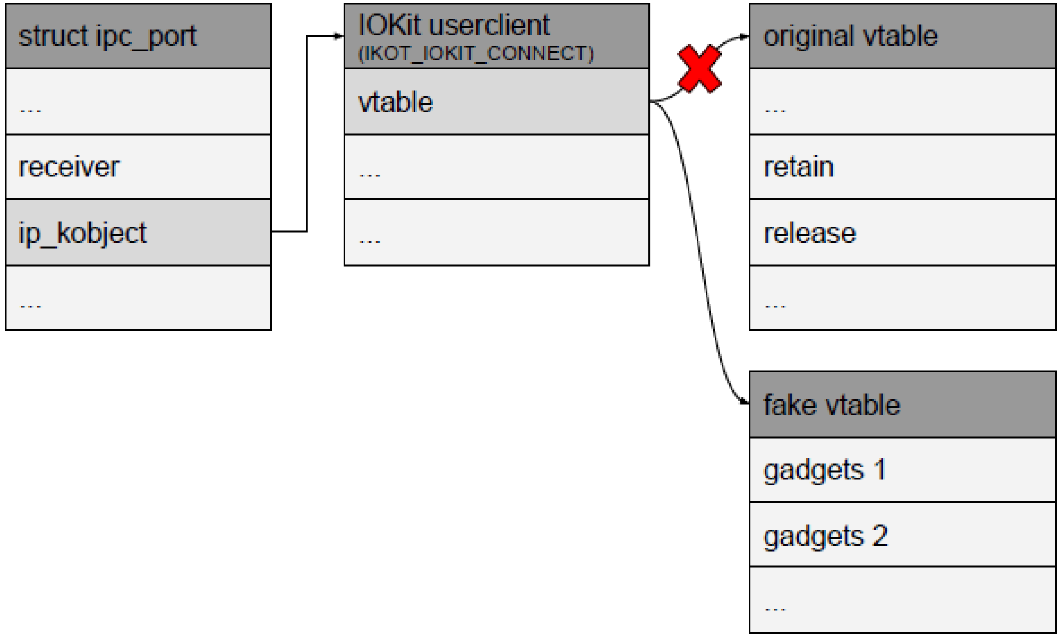
```
struct  clock   clock_list[] = {

    /* SYSTEM_CLOCK */
    { &sysclk_ops, 0, 0 },

    /* CALENDAR_CLOCK */
    { &calend_ops, 0, 0 }
};
```

```
struct clock_ops sysclk_ops = {
    NULL,
    rtclock_init,
    rtclock_gettime,
    rtclock_getattr,
};
```

```
/*
 * Get clock attributes.
 */
kern_return_t
clock_get_attributes(
    clock_t                 clock,
    clock_flavor_t          flavor,
    clock_attr_t            attr,       /* OUT */
    mach_msg_type_number_t  *count)     /* IN/OUT */
{
    if (clock == CLOCK_NULL)
        return (KERN_INVALID_ARGUMENT);
    if (clock->cl_ops->c_getattr)
        return (clock->cl_ops->c_getattr(flavor, attr, count));
    return (KERN_FAILURE);
}
```

- This type of primitives can be used to execute kernel code (e.g., a ROP chain or a kernel function) in **arbitrary** addresses.
- clock_get_attributes() is a system call to get attributes of target clock object. An attack can change the **global** function pointers or **fake** an object to hijack the control flow.
- This technique was used in the **Pegasus** APT attack in iOS 9.3.3.

Arbitrary Code Execution Primitives



- IOKit is an object-oriented device driver framework in XNU that uses a subset of **C++** as its language.

- If the attacker has the kernel write primitives, then he can change the **vtable** entry of an I/OKit userclient to hijack the control flow to the address of a ROP gadget to achieve a kernel code execution primitive.

Practical Case Study: Yalu Exp (fixed in iOS 10.2.1)

```
kern_return_t
mach_voucher_extract_attr_recipe_trap(
struct mach_voucher_..._args *args)
{
...
 mach_msg_type_number_t sz = 0;

 copyin(args->recipe_size, (void *)&sz, \
        sizeof(sz));
...
 uint8_t *krecipe = kalloc((vm_size_t)sz);
...
 //args->recipe_size should be sz
 copyin(args->recipe, (void *)krecipe, \
        args->recipe_size)
...
}
```

- CVE-2017-2370 is a heap **buffer overflow** in mach_voucher_extract_attr_recipe_trap().
- The function first copies 4 bytes from the user space pointer args->recipe_size to the **sz** variable. After that, it calls kalloc(sz).
- The function then calls copyin() to copy args->recipe_size sized data from the user space to the **krecipe** (should be **sz**) sized kernel heap buffer. Consequently, it will cause a buffer overflow.

Practical Case Study: Yalu Exp

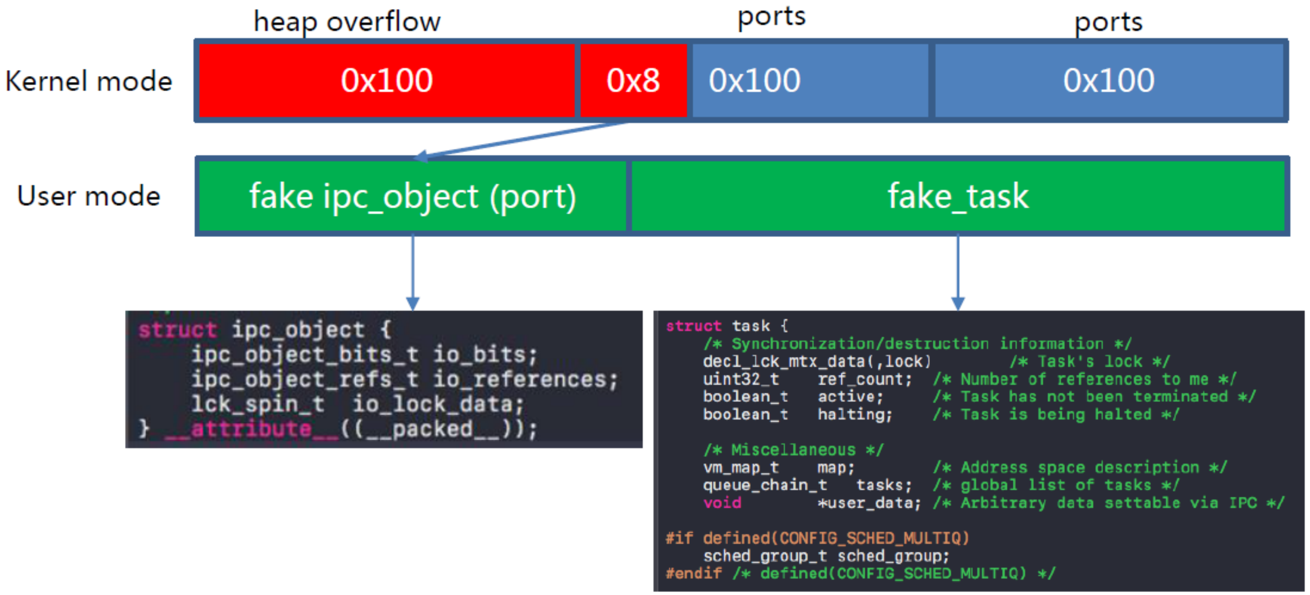
- **Before heap overflow**

```
(lldb) x/50x 0xffffff8029404c00
0xffffff8029404c00: 0xdeadbeefdeadbeef 0xffffffffffffffff
0xffffff8029404c10: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404c20: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404c30: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404c40: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404c50: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404c60: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404c70: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404c80: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404c90: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404ca0: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404cb0: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404cc0: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404cd0: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404ce0: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404cf0: 0xffffffffffffffff 0xdeadbeefdeadbeef
0xffffff8029404d00: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d10: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d20: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d30: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d40: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d50: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d60: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d70: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d80: 0xffffffffffffffff 0xffffffffffffffff
```

- **After heap overflow**

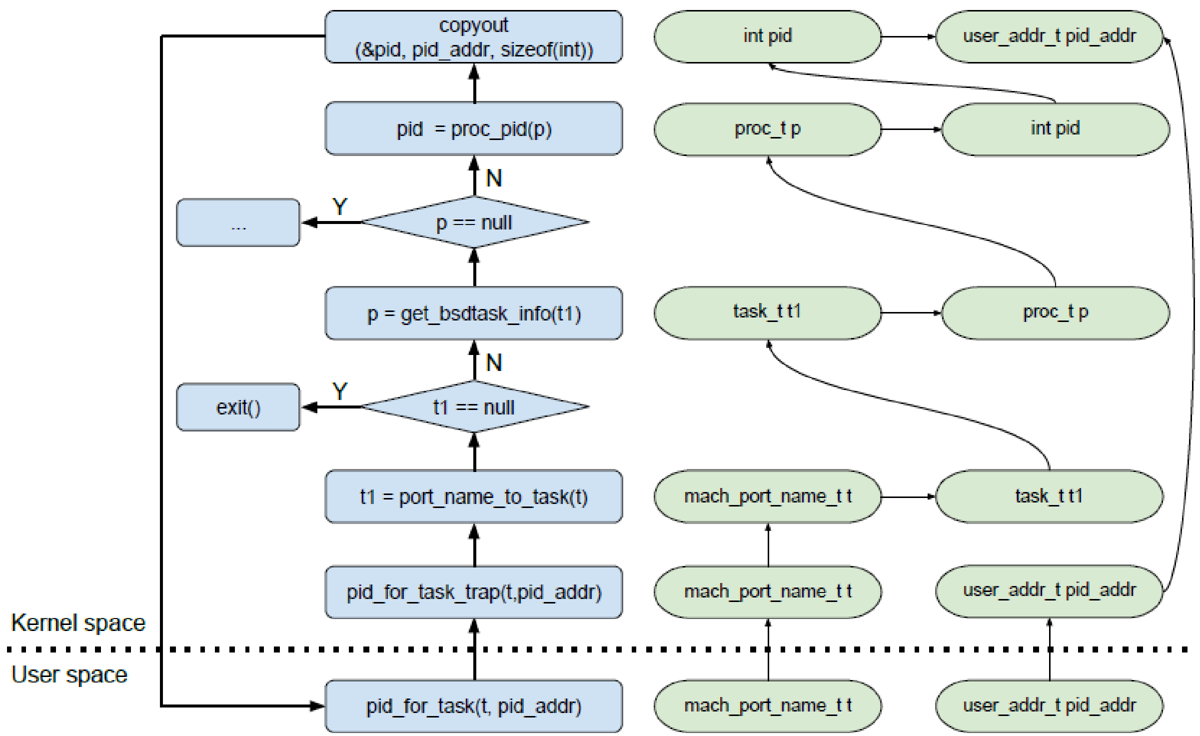
```
(lldb) x/50x 0xffffff8029404c00
0xffffff8029404c00: 0x4141414141414141 0x4141414141414141
0xffffff8029404c10: 0x4141414141414141 0x4141414141414141
0xffffff8029404c20: 0x4141414141414141 0x4141414141414141
0xffffff8029404c30: 0x4141414141414141 0x4141414141414141
0xffffff8029404c40: 0x4141414141414141 0x4141414141414141
0xffffff8029404c50: 0x4141414141414141 0x4141414141414141
0xffffff8029404c60: 0x4141414141414141 0x4141414141414141
0xffffff8029404c70: 0x4141414141414141 0x4141414141414141
0xffffff8029404c80: 0x4141414141414141 0x4141414141414141
0xffffff8029404c90: 0x4141414141414141 0x4141414141414141
0xffffff8029404ca0: 0x4141414141414141 0x4141414141414141
0xffffff8029404cb0: 0x4141414141414141 0x4141414141414141
0xffffff8029404cc0: 0x4141414141414141 0x4141414141414141
0xffffff8029404cd0: 0x4141414141414141 0x4141414141414141
0xffffff8029404ce0: 0x4141414141414141 0x4141414141414141
0xffffff8029404cf0: 0x4141414141414141 0x4141414141414141
0xffffff8029404d00: 0x4242424242424242 0x4242424242424242
0xffffff8029404d10: 0x4242424242424242 0x4242424242424242
0xffffff8029404d20: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d30: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d40: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d50: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d60: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d70: 0xffffffffffffffff 0xffffffffffffffff
0xffffff8029404d80: 0xffffffffffffffff 0xffffffffffffffff
```

Practical Case Study: Yalu Exp



- The exploit overflow those pointers and modify one ipc_object **pointer** to point to a **fake** ipc_object in user mode. The exploit creates a **fake** task in user mode for the fake port as well.
- After that, the exploit chain calls QP (e.g., clock_sleep_trap()) to **brute force** the address of the **global** system clock.

Practical Case Study: Yalu Exp



- The exploit sets io_bits of the fake ipc_object to IKOT_TASK and craft a **fake** task for the fake port. By setting the value at the faketask + bsdtask offset, an attacker could **read arbitrary** 32 bits kernel memory through pid_for_task() **without** break KASLR.

Practical Case Study: Yalu Exp

```
kern_return_t pid_for_task(struct pid_for_task_args *args)
{
    mach_port_name_t    t = args->t;
    user_addr_t      pid_addr  = args->pid;  //return value
    ...
    t1 = port_name_to_task(t);   //get faketask
    ...
    p = get_bsdtask_info(t1);    //get *(faketask + procoff)
    if (p) {
        pid  = proc_pid(p);      //get *(p + 0x10)
        err = KERN_SUCCESS;
    }
    ...
    //copy the value to pid_addr
    (void) copyout((char *) &pid, pid_addr, sizeof(int));
    return(err);
}
```

```
__int64 __fastcall get_bsdtask_info(__int64 a1)
{
  return *(_QWORD *)(a1 + 0x380);
}
```

```
signed __int64 __fastcall proc_pid(__int64 a1)
{
  signed __int64 result; // rax@1

  result = 0xFFFFFFFFLL;
  if ( a1 )
    result = *(_DWORD *)(a1 + 0x10);
  return result;
}
```

```
//copy the value to pid_addr
(void) copyout((char *) &pid, pid_addr, sizeof(int));
```

➡

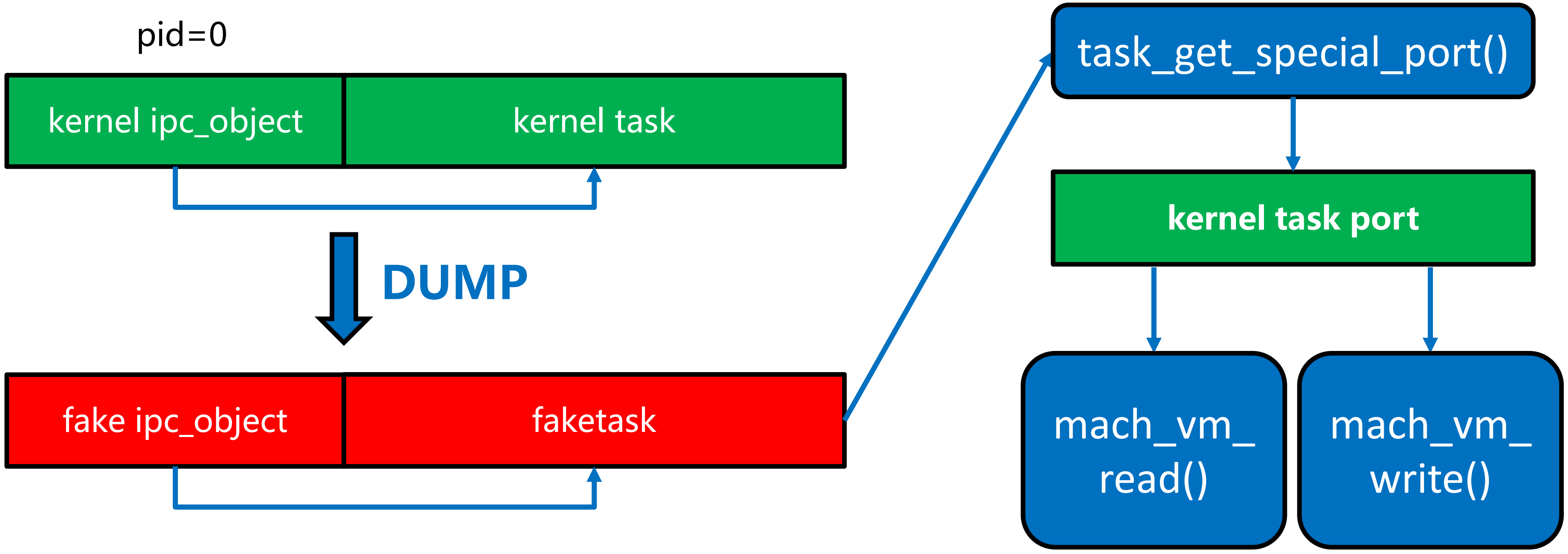
```
read 0xffffff800cc00000 : 0xfeedfacf
```

- As we mentioned before, the function doesn’t check the **validity** of the task, and just return the value of *(*(faketask + 0x380) + 0x10).

Practical Case Study: Yalu Exp

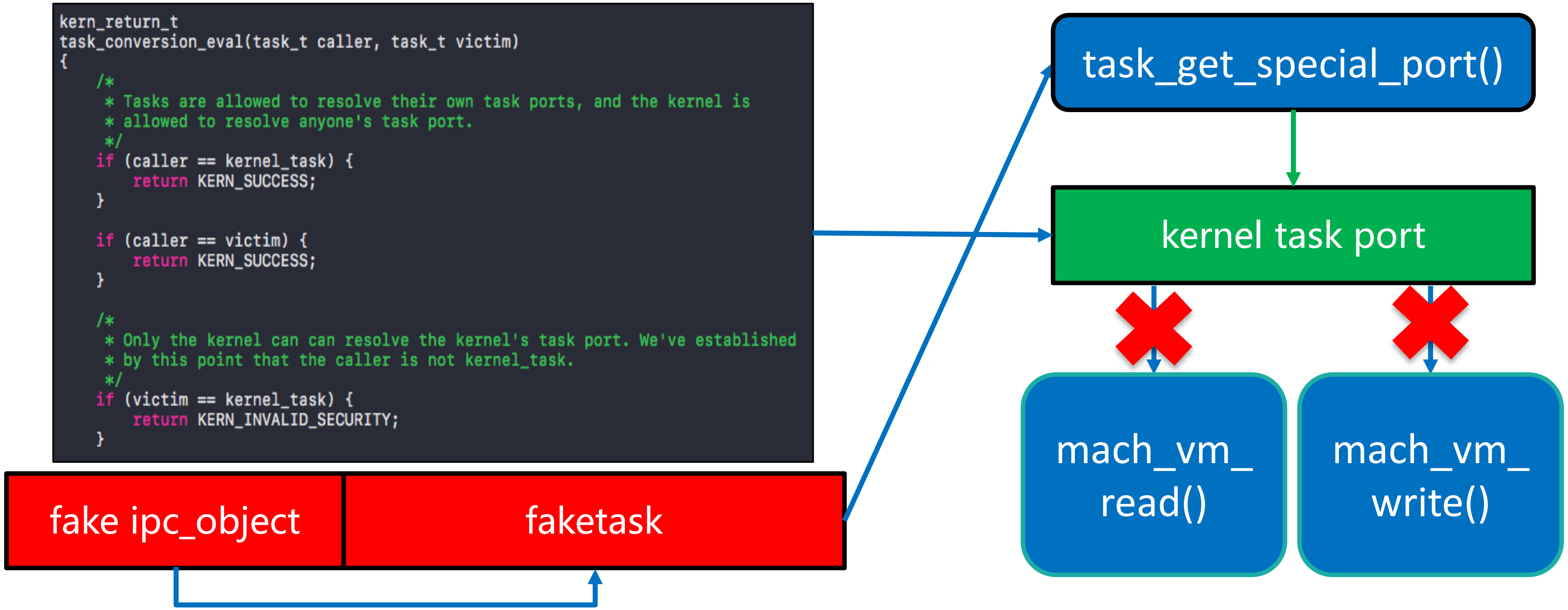
- The attacker **dumps** kernel ipc_object and kernel task to a fake ipc_object and a fake task. By using task_get_special_port() to the fake ipc_object and task, the attacker could get the **kernel task port**.

- Kernel task port can be used to do **arbitrary** kernel memory read and write.



iOS 11 Kernel Task Mitigation

- iOS 11 added a new **mitigation** that only the kernel can resolve the kernel’ s task port.
- We **cannot** use the task_get_special_port() trick on iOS 11.

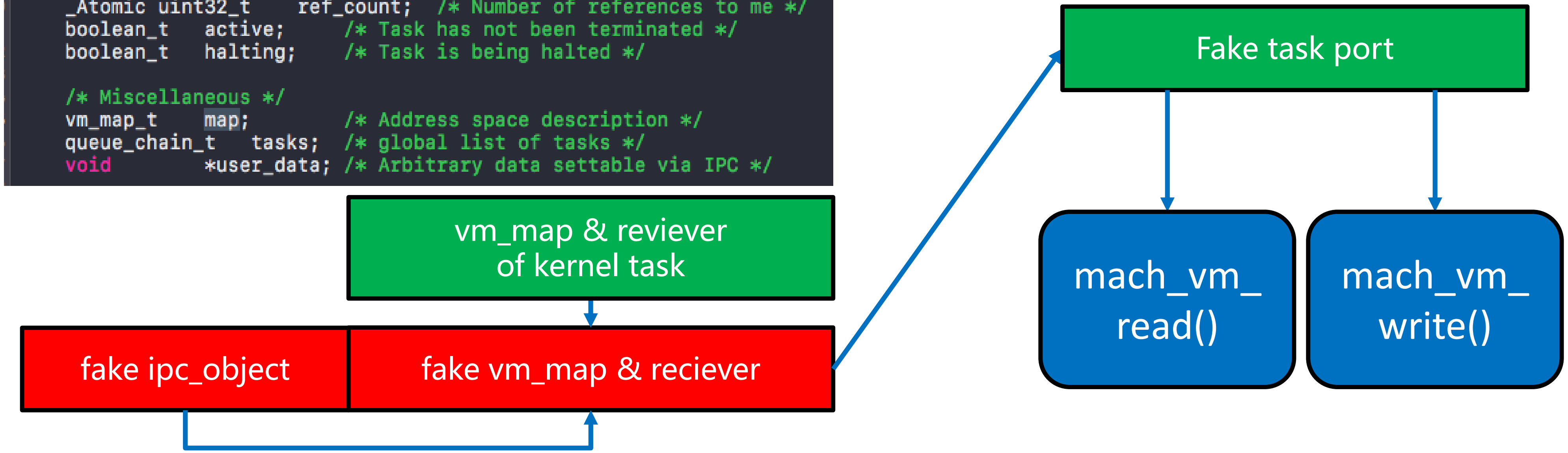



```
kern_return_t
task_conversion_eval(task_t caller, task_t victim)
{
    /*
     * Tasks are allowed to resolve their own task ports, and the kernel is
     * allowed to resolve anyone's task port.
     */
    if (caller == kernel_task) {
        return KERN_SUCCESS;
    }

    if (caller == victim) {
        return KERN_SUCCESS;
    }

    /*
     * Only the kernel can can resolve the kernel's task port. We've established
     * by this point that the caller is not kernel_task.
     */
    if (victim == kernel_task) {
        return KERN_INVALID_SECURITY;
    }
```

Mitigation bypass in Async_wake Exp

- The attacker cannot use a real kernel task port. But the attacker can copy reference pointer of kernel’ s **vm** to the fake task.
- Now the fake port has a same **address space** as the kernel task port. It’ s enough for the attacker to do arbitrary kernel read/write.

```
struct task {
    /* Synchronization/destruction information */
    decl_lck_mtx_data(,lock)          /* Task's lock */
    _Atomic uint32_t    ref_count;  /* Number of references to me */
    boolean_t    active;      /* Task has not been terminated */
    boolean_t    halting;     /* Task is being halted */

    /* Miscellaneous */
    vm_map_t     map;         /* Address space description */
    queue_chain_t    tasks;   /* global list of tasks */
    void         *user_data;  /* Arbitrary data settable via IPC */
```



Enterprise Computer Security

Pic from time.com

- Lots of companies (e.g., Alibaba Inc and Tencent) offer Macbooks as work computers to their employees.
- **Problems:**

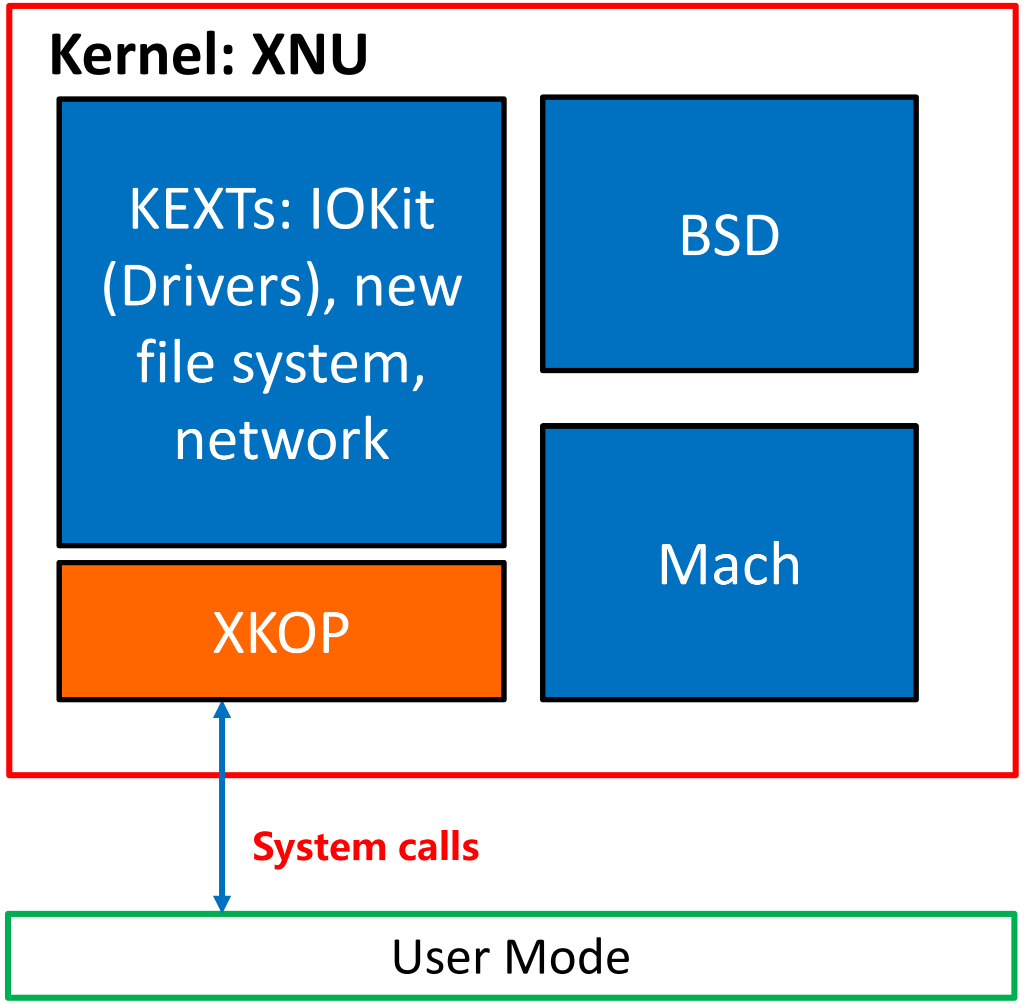
  1. macOS is not forced to upgrade like iOS.

  2. Less hardware based protections (e.g., AMCC and PAC) on Macbooks.

  3. Less secure sandbox rules than iOS.

- Hard to defend against advanced persistent threat (APT). Enterprise computers need a more **secure** system.

XNU Kernel Object Protector for macOS



- To mitigate the APT and POP attack, we propose a framework called ***XNU Kernel Object Protector*** (XKOP).
- **Basic idea**: a kernel extension to implement instrumentation for specific system calls and deploy integrity check for ipc_port kernel objects.
- In addition, XKOP could bring **new** mitigations to **old** macOS versions.

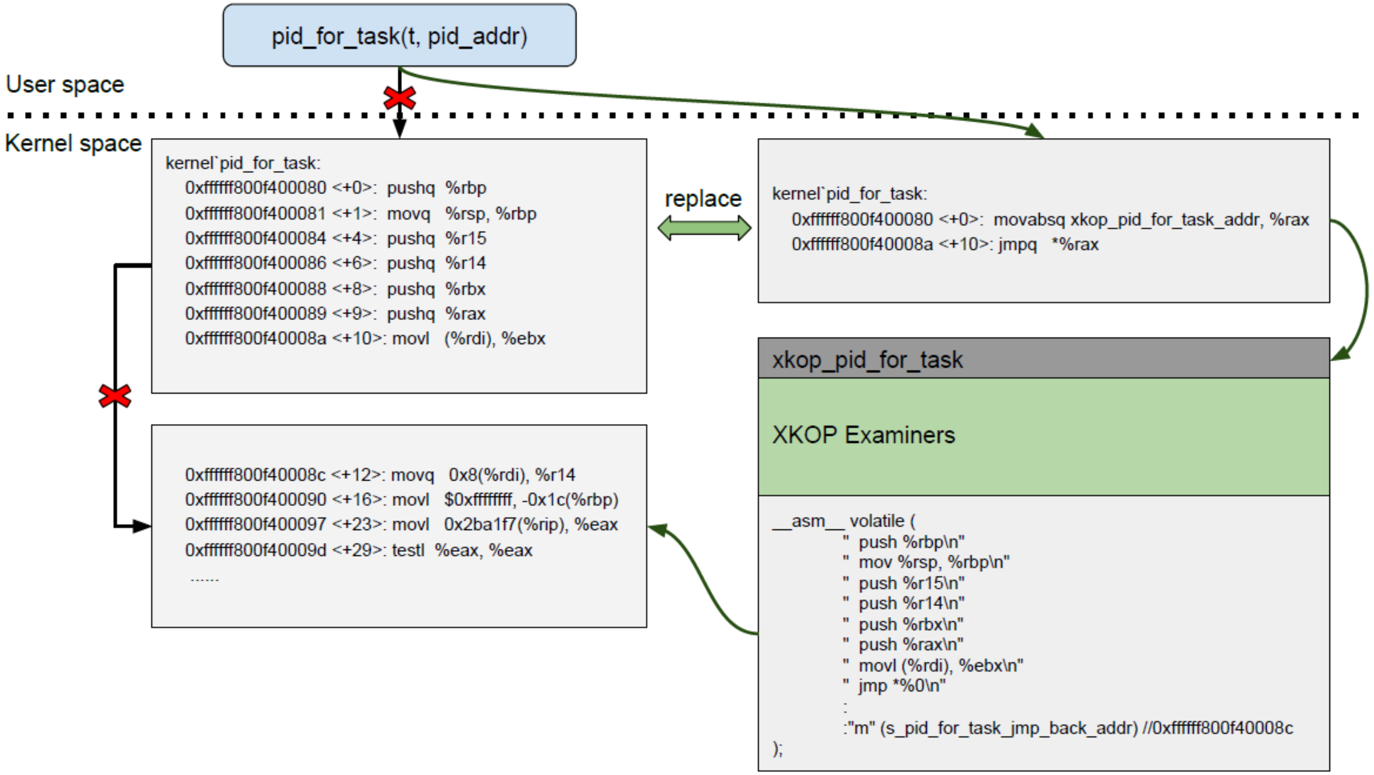
Instrumentation

xnu-4570.41.2 〉 security 〉 h mac_policy.h 〉 No Selection

```
/**
  @file mac_policy.h
  @brief Kernel Interfaces for MAC policy modules

  This header defines the list of operations that are defined by the
  TrustedBSD MAC Framwork on Darwin.  MAC Policy modules register
  with the framework to declare interest in a specific set of
  operations.  If interest in an entry point is not declared, then
  the policy will be ignored when the Framework evaluates that entry
  point.
*/

#ifndef _SECURITY_MAC_POLICY_H_
#define _SECURITY_MAC_POLICY_H_

#ifndef PRIVATE
#warning "MAC policy is not KPI, see Technical Q&A QA1574, this header
#endif

#include <security/_label.h>

struct attrlist;
struct auditinfo;
struct bpf_d;
struct cs_blob;
struct devnode;
struct exception_action;
struct fileglob;
struct ifnet;
struct inpcb;
struct ipq;
struct label;
struct mac_module_data;
```

- Our system needs to find reliable code points that the **examiners** could be executed.
- **KAuth** kernel subsystem exports a KPI that allows third-party developers to authorize actions within the kernel. However, the operation set is very limited.
- **MAC** framework is private and can only be used by Apple. In addition, the rules are hardcoded in the code of the XNU kernel.
- Finally, we choose **instrumentation**.

Instrumentation



- Based on the examiners, XKOP replaces the original code entry of the target system call into a **trampoline**. The trampoline jumps to the **examiner** stored in the XKOP kernel extension. Then, the examiner verifies the **integrity** of the target kernel object.

Examiners

```
kern_return_t pid_for_task(struct pid_for_task_args *args)
{
    mach_port_name_t    t = args->t;
    user_addr_t     pid_addr  = args->pid;  //return value
    ...
    t1 = port_name_to_task(t);  //get faketask
    ...
    p = get_bsdtask_info(t1);   //get *(faketask + procoff)
    if (p) {
        pid  = proc_pid(p);     //get *(p + 0x10)
        err = KERN_SUCCESS;
    }
    ...
    //copy the value to pid_addr
    (void) copyout((char *) &pid, pid_addr, sizeof(int));
    return(err);
}
```

Kernel object address checker: t1 should not be in the user space address. Must break KASLR first and put the payload into kernel. Just like a soft SMAP for old devices.

```
__int64 __fastcall get_bsdtask_info(__int64 a1)
{
  return *(_QWORD *)(a1 + 0x380);
}
```

Kernel object type examiner: a1 should be a real badtask_info structure with a valid pid number.

Examiners

```c
uint64_t textbase = 0xfffffff007004000;
while(1)
{
    k+=8;
    //guess the task of clock
    *(uint64_t*)(((uint64_t)fakeport) + 0x68) = textbase + k;
    *(uint64_t*)(((uint64_t)fakeport) + 0xa0) = 0xff;

    //fakeport->io_bits = IKOT_CLOCK | IO_BITS_ACTIVE ;
    kern_return_t kret = clock_sleep_trap(foundport, 0, 0, 0, 0);

    if (kret != KERN_FAILURE) {
        printf("task of clock = %llx\n",textbase + k);
        break;
        }
}
```

Through brute force attacks, clock_sleep_trap() can be used to guess the address of global clock object and break the KASLR.

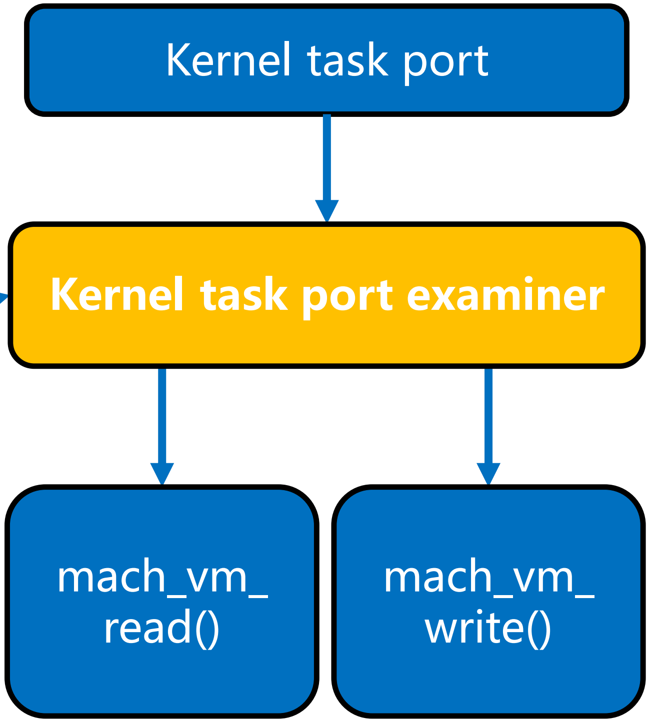
```c
clock_t
port_name_to_clock(
    mach_port_name_t clock_name)
{
    clock_t     clock = CLOCK_NULL;
    ipc_space_t space;
    ipc_port_t  port;

    if (clock_name == 0)
        return (clock);
    space = current_space();
    if (ipc_port_translate_send(space, clock_name, &port) != KERN_SUCCESS)
        return (clock);
    if (ip_active(port) && (ip_kotype(port) == IKOT_CLOCK))
        clock = (clock_t) port->ip_kobject;
    ip_unlock(port);
    return (clock);
}
```

Kernel object querying examiner: if the function returns too many errors, warning the user or panic according to the configuration.

Examiners

- **Kernel task port examiner**: firstly, bring task_conversion_eval(task_t caller, task_t victim) mitigation to old macOS system versions. Only the kernel can resolve the kernel's task port.

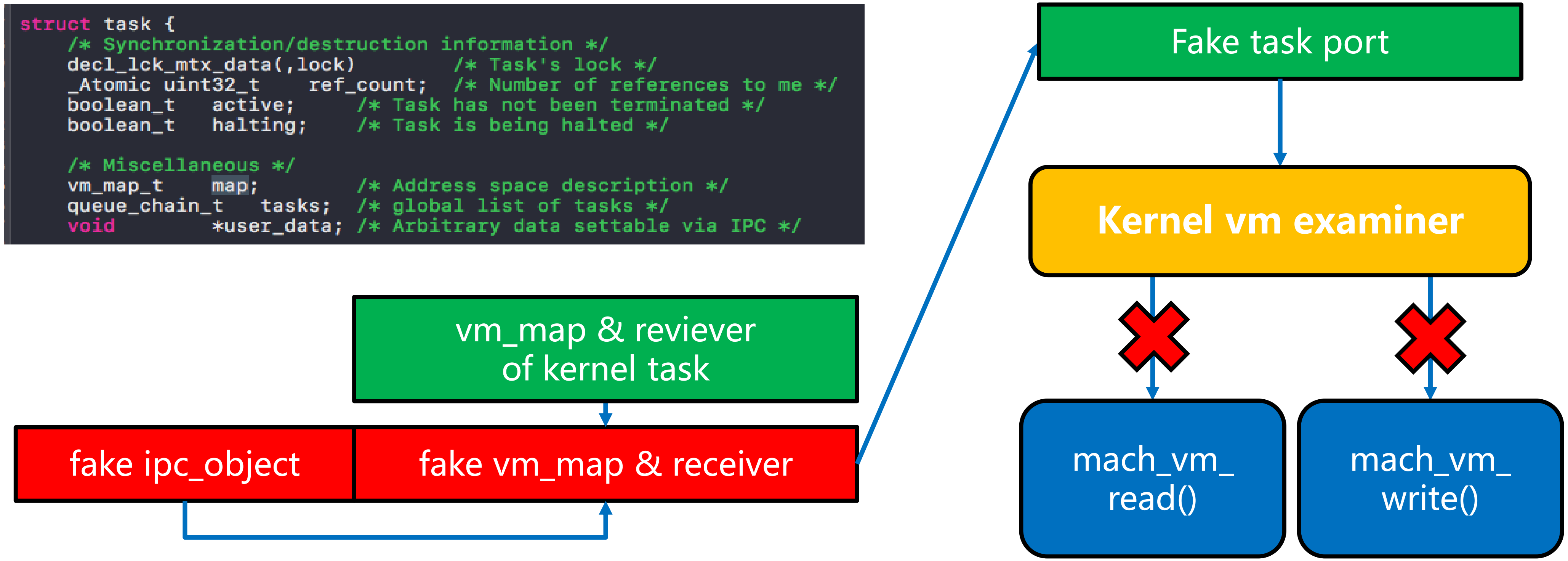
```
kern_return_t
task_conversion_eval(task_t caller, task_t victim)
{
    /*
     * Tasks are allowed to resolve their own task ports, and the kernel is
     * allowed to resolve anyone's task port.
     */
    if (caller == kernel_task) {
        return KERN_SUCCESS;
    }

    if (caller == victim) {
        return KERN_SUCCESS;
    }

    /*
     * Only the kernel can can resolve the kernel's task port. We've established
     * by this point that the caller is not kernel_task.
     */
    if (victim == kernel_task) {
        return KERN_INVALID_SECURITY;
    }
```



Examiners

- **Kernel vm examiner** for mach_vm_*(): if the caller process does not belong to kernel (pid == 0) and the target ipc_port object has the same map structure with the one of a kernel task, the examiner will trigger configured operations, e.g., error return or panic.

```
struct task {
    /* Synchronization/destruction information */
    decl_lck_mtx_data(,lock)          /* Task's lock */
    _Atomic uint32_t     ref_count;  /* Number of references to me */
    boolean_t    active;     /* Task has not been terminated */
    boolean_t    halting;    /* Task is being halted */

    /* Miscellaneous */
    vm_map_t     map;        /* Address space description */
    queue_chain_t   tasks;   /* global list of tasks */
    void         *user_data; /* Arbitrary data settable via IPC */
```



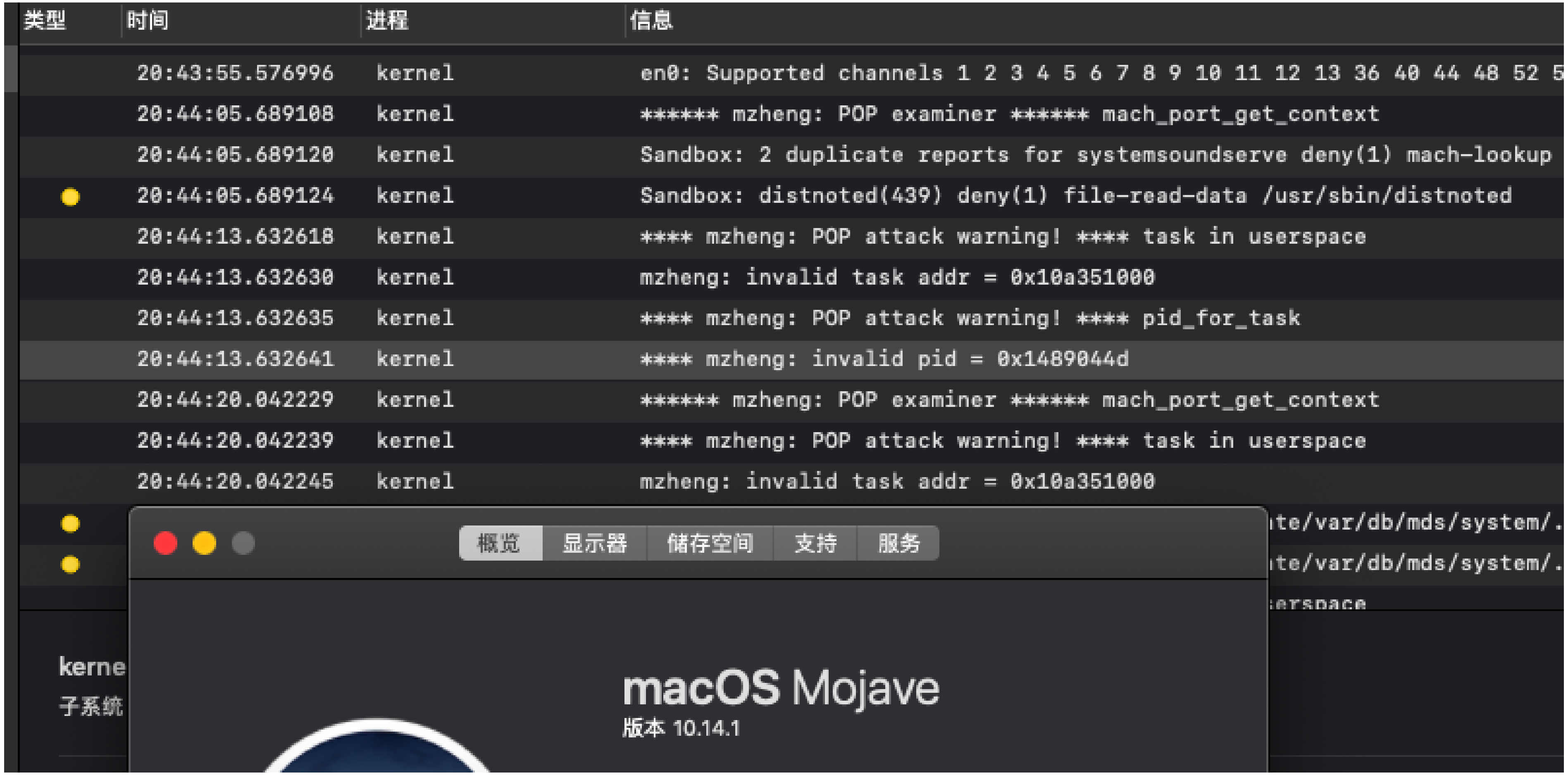
Evaluation

- We selected 5 kernel vulnerabilities (one or two exploitable bugs for each version of macOS) and state-of-the-art exploits to **evaluate** the effectiveness of our system.

macOS version	Vulnerability (CVE)	XKOP Protection
10.12	CVE-2016-4669	YES
	CVE-2017-2370	YES
10.13	CVE-2017-13861	YES
	CVE-2018-4241	YES
10.14	CVE-2019-6225	YES

- We first ensure that the **exploits** work on the corresponding macOS systems, and then we deploy the XKOP framework and run the exploits again to check whether our system detects and blocks the attack.

Evaluation



- The experiment (voucher_swap exp for macOS) shows that XKOP provides **deterministic** protection for every vulnerability and blocks each attempt to exploit the system.

Discussion

- Unfortunately, XKOP cannot mitigate **all** kinds of POP primitives:

  (1). Querying primitives use error return values to gain an extra source of information which is very similar to the **side-channel** attack.

  (2). No protection for arbitrary code execution primitives. Without hardware support, software-based CFI implementation can be very **expensive**. In addition, modern kernel could be patched by **pure data** which means kernel memory read and write primitives are enough for attackers to accomplish the aim.

- We may miss some potential vulnerabilities that can bypass XKOP protection. As an imperfect solution, XKOP supports **extensible** examiners to prevent new threats in the first place.

Conclusion

- We discuss the **mitigation** techniques in the XNU kernel, i.e., the kernel of iOS and macOS, and how these mitigations make the traditional exploitation technique ineffective.

- We summarize a new attack called **POP** that leverages multiple ipc_port kernel objects to bypass these mitigations.

- A defense mechanism called XNU Kernel Object Protector (**XKOP**) is proposed to protect the integrity of the kernel objects in the macOS kernel.

- We have discussed this issue with Apple's security team (follow-up id: 707542859). They appreciate our assistance in helping to maintain and improve the security of their products and plan to *deploy security enhancements against the Mach port issue for a future release of iOS and macOS*.

Reference

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- iOS 10 Kernel Heap Revisited, Stefan Esser
- Port(al) to the iOS Core, Stefan Esser
- mach voucher buffer overflow. https://bugs.chromium.org/p/projectzero/issues/detail?id=1004
- Mach portal: https://bugs.chromium.org/p/project-zero/issues/detail?id=965
- voucher_swap: https://googleprojectzero.blogspot.com/2019/01/voucherswap-exploiting-mig-reference.html
- PassiveFuzzFrameworkOSX: https://github.com/SilverMoonSecurity/PassiveFuzzFrameworkOSX

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Thanks