Submission

February 4, 2021

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[7]: # Craig Fox
print('Name: Craig Fox')
# TUID: 915781095
print('TUID: 915781095')

from math import comb
import math
```

Name: Craig Fox TUID: 915781095

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[29]: # Problem 1
      print('Problem 1\n')
      # Erdős-Rényi graph: N=4000, p=0.001
      n = 4000
      p = 0.001
      # What is the average degree of a node in this graph?
      print('Erdős-Rényi graphs have an average degree = p(n-1)')
      k = p*(n-1)
      print('This graph would have an average degree = ' + str(k) + '\n')
      # What is the variance in the degrees of the nodes?
      print('Erdős-Rényi graphs have a variance = p(1-p)(n-1)')
      print('This graph would have a variance = ' + str(k*(1-p)) + '\n')
      # What is the expected number of nodes with a degree which is at least twice_
      → larger that the average degree?
      print('The Probability of a node having k degrees = (n-1 pick,)
       \hookrightarrow k) (p^k) (1-p)^(n-1-k)'
      at_least_twice_degree = []
      for x in range(math.ceil(2*k), 100):
          #Note: The range should go from math.ceil(2*k) to n-1, but since n-1 is
       \rightarrow3999 the factorial in the combination becomes incredibly large. 100 is
       →already high enough for a high accuracy
          at_least_twice_degree.append(comb(n-1, x)*(p**x)*(1-p)**(n-1-x))
```

Problem 1

Erdős-Rényi graphs have an average degree = p(n-1)This graph would have an average degree = 3.999

Erdős-Rényi graphs have a variance = p(1-p)(n-1)This graph would have a variance = 3.9950010000000002

The Probability of a node having k degrees = $(n-1 \text{ pick k})(p^k)(1-p)^(n-1-k)$ So the number of nodes greater with greater than or equal to 2k degree = 203.93917397722828

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[32]: # Problem 2
      print('Problem 2\n')
      # Consider Gn,p, an Erdös-Rényi random graph with n nodes, m edges, and mean ⊔
       \rightarrow degree c
      # a) Compute the probability p of creating an edge in Gn,p.
      print('The maximum number of edge is n(n-1)/2. A graph has c*n/2 edges. ⊔
       \Rightarrow \text{Therefore p = } (c*n/2)/(n(n-1)/2) = c/(n-1) \setminus n')
      # b) Show that in the limit (large n) the expected number of triangles in Gn, p_{\sqcup}
       → is (1/6)*c^3
      print('For a triangle to exist, there has to be three codes all sharing edges. ⊔
       \hookrightarrowSo there are n Pick 3 possible triangles which equals n!/(3!(n-3)!). Each,
       \rightarrowedge in the triangle has a c/(n-1) chance of exisiting (as shown in the \sqcup
       ⇒previous question). So for a specific triangle to appear there is a (c/
       _{\hookrightarrow}(n-1))^3 chance. So the total number of triangles = the probability of a_{\sqcup}
       \negtriangle times the number of potential triangles = n!/(3!(n-3)!)*(c/(n-1))^3.
        \rightarrow As n approaches infinity this simplies to (c^3)/6')
```

Problem 2

The maximum number of edge is n(n-1)/2. A graph has c*n/2 edges. Therefore p = (c*n/2)/(n(n-1)/2) = c/(n-1)

For a triangle to exist, there has to be three codes all sharing edges. So there are n Pick 3 possible triangles which equals n!/(3!(n-3)!). Each edge in the triangle has a c/(n-1) chance of exisiting (as shown in the previous question). So for a specific triangle to appear there is a $(c/(n-1))^3$ chance. So the total number of triangles = the probability of a triangle times the number of potential triangles = $n!/(3!(n-3)!)*(c/(n-1))^3$. As n approaches infinity this

simplies to $(c^3)/6$

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