#### Generating Circumstances

#### Zeeshan Lakhani

1-9-2015 (Codemash)



#### What I Listened To



# Cod Mash 2015

08 JAN 2015

#### Is TDD Dead?

By DHH FOWLER

Test-first fundamentalism is like abstinence-only sex ed An unrealistic, ineffective morality campaign for self-loathing and shaming.

Long live testing.



tion Council, worldwide moose numbers are expected to grow markedly on last year due to the traditional moose strongholds of Canada and the United States, with the larger developing moose ecologies also poised to make gains. The largest percentagege increase in moose will likely come from China", says McRobson, The Chinese government has invested heavily in moose infrastructure over the past decade, and their committment to macrofauna is beginning to pay dividends". Since 2004 China has expanded moose pasture from 1.5% of arable land to nearly 3.648% and moose numbers are expected to rise to 60,000 making China a net moose dred million billion.

Europe's rise as an international moose power will slow slightly this year as a response to the European Union's move towards standardising the European moose. Stringent quality controls are holding back the development of the eastern european populations compared to last year when they contributed significantly to europe's strong growth figures. Norway, which is not an EU member but has observer status, strengthed in numbers relative to the Euro area with numbers of Norweigian moose. known locally as elk" expected to rise for the tenth consecutive year, particularly thanks to a strong showing in 1

<sup>1</sup>http://bit.ly/1ALmSGC

### In Papers

#### QuickCheck: A Lightweight Tool for Random Testing of Haskell Programs

cally tested on random input, but it is also possible to de-

case studies, in which the tool was successfully used, and also point out some pitfalls to avoid. Random testing is es-

a multi-core processor. In that congent residential programs because properties a hierarchical approach to testing in a property of the state of the s

a interactional approach to sensing engage-sensitivities of the first at miscous in ourse trees debugging of fault encountered. Separately is tested components, then mandom testing suffices The Lebusy programming languages/datastopte\_datastop

is a seal problem for Erlang descloper 1889 (Fried 1889), Moher ter time, or to test more thorover, these race conditions are often raised by decapturely, water, they make it easy to repeat

see meriodado expensivo o consideratador sea da medificacionada a recevam. In this passer se-

In this paper, we describe tools withard-skate/sped-familiading-ill suited to automatic testing

race conditions in Erlang code during any annual Geograph while pure functions are much easier

based on property-based texting toping Quick Quick Competing times, because one need not be

Haghes 2000), in a commercial spagion lightlylang descioped office and office execution. In an

Quying AB (Haghes 2007; Arts et al., 2006), do practi for away, and other pure

Quying All Congress court, Artis to the personner improvement or whose programs are described in section 3. We develop a personner to the proceedures from which

not share memory, and firlanc data structures are immutable

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races is difficult, so we developed a trace visualizer which

to an industrial case study, which is introduced in section

the kind of alass races which plaged unpiny FRODE CTION

#### Finding Born Conditions in Follow with Ontal Charle and Day of

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Keen Claessen Nicholas S	Michał Pałka mallbone	John Hughes Thor	Hans Svensson nas Arts	Koen Claessen Chaimers University of Techn Ulif Wiger koen@cs.chalmers.s	
Chalmers University of Technology, Gothenburg, Sweden		Chairners University of Technology and Qaviq AB		alf.uigsrborlung-consulting.con ABSTRACT	
keenGchalmers.se		rjsh@chalsers.se		QuickCheck is a tool which aids the Hashell pro	ene

#### Abstract

micematchalmers.se

We address the problem of testing and debugging concurrent, disinbuted Erlang applications. In concurrent programs, race conditions are a common class of hugs and are very hard to find in pra-, where due to the yest amount of code under test, it is often hand to diagnose the error resulting from race conditions. We present three tools (QuickCheck, PULSE, and a visualizer) that in combination can be used to test and debug concurrent programs in unit testing with a much better possibility of detecting race conditions

We evaluate our method on an industrial concurrent case study and Electrate how we find and analyze the race conditions. Categories and Subject Descriptors D.2.5 (Textics and Debar-

Keywords QuickCheck, Race Conditions, Erlang

Concurrent programming is notoriously difficult, because the nonlead software to work most of the time, but fail in rare and hardreproduce circumstances when an unfortunate order of events concerns on the property of th

components of a software system are tested together. Timing delays caused by other component Clarin Benne Carte Chris-Ake Fredhand tion 6. We evaluate our tools by applying them schedules of actions performed by the individual units. In the we Security of configuration (1) and the config

\*\*Common Country (Section 2014) Control Country (Section 2014) Count Poliseunica de Machid. 111. We begin by introducing the industrial case that we arely our Revwords: Software Testing, Java, Erlang.

Acknowledgements: The authors were supported by the DESAFIOS 10 project (TIN2009-14599-C03-O3) financed by the (former) Spanish Ministry of Education and Science, and the EU FP7 project PROWESS (317820)

#### 1 Introduction

This paper describes a methodology for testing Java code which uses the property-based random testing Ouvin OuickCheck tool [AHJW05], in the following often abbreviated as OuickCheck

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A testing tool must be able to determine whether a test matically checkable criterion of doing so. We have chopen introduction to use formal specifications for this purpose. We have to

signed a simple domain-specific language of totable specifi-outions which the tenter uses to define exacted properties of the nice things about purely functional languages is that functions of the functions under test. QuickCheck then checks that walking simple properties, and enjoy simple algebraic relationships. Indeed, properties hold in a large number of cases. The specificanctions of an API satisfy elegant laws, that in itself is a sign of a good desition lammare is embedded in Baskell using the class systems laws not only indicate conceptual simplicity, but are useful in practic Properties are normally written in the same module as difficilitieng programs that use the API, by equational reasoning or otherwi-Properties are normany written in the same and the same supportant to the same state of the same state A testing tool must also be able to generate test cases all

tomatically. We have chosen the simplest method, random reverse (xs++ys) == reverse xs++reverse ys testing [11], which competes surprisingly favourably with systematic methods in practice. However, it is meaninglethre reverse is the list reversal function, and ++ is list around to this over cardier feeling willows decreasing the modelline review of the last reversal fractions, and ++ is not appear.

It is producted to the feeling willows decreasing the modelline report of the feeling willows described in the feeling of the feeling of the feeling them in correctly. A stated law which is untrue is a the feeling them term this is not possible, since the distribution of schools to eshould try out the law in a few cases—just to avoid stupid mist data in all subsequent ropes is not known. A uniform Mis can ease that task a little bit by defining a function to test the law ; tribution is often used instead, but for data drawn features for its free variables:

infinite sets this is not even meaningful. how would one choose a resolven closed & term with a resilven distributionDEGO TAVASO X8 V8 \* for example? We have chosen to put distribution under the reverse (xx++yx) -- reverse xx++reverse yx increase (also embedded in Haskell), and a way to observe we can test the law just by applying prop\_revApp to suitable pairs of

generator, the tester can not only control the distribution for no have a tool that would perform that task for us? Then we could si generator, the tener can not only the control of test cases, but also ensure that they satisfy arbitragiftle laws in our programs and automatically check that they are reason mainly used from the Hugs interpreter. We have also writ-ten a small script to invoke it, which needs to know very quickCheck prop\_revApp

little about Haskell systex, and consequently supports the Falsifiable, after 2 tests: full language and its extensions. It is not dependent on audit. -1] particular Haskell system. A cost that comes with this deci{0}

method is in detecting faults. However, we have used Quigit: filebookt inspectifig the pulsperfyringer, dikelyness-kimitestring and for in

#### the distribution of test cases. By programming a suitabliaventing such pairs of lists, and running the tests, is tedious, however, Wou complex invariants. An important design goal was that QuickCheck shows Apotheses, at least. In 1999, Keen Classes and I built just such a tor An important design goal was that QuickCheck shows Apotheses and Classes and Classes and Classes are the Apotheses and Classes and Classes and Classes are the Apotheses and Classes and Classes are the Apotheses and Classes and Classes and Classes are the Apotheses and Classes and Classes and Classes are the Apotheses and Classes and Classes are the Apotheses and Classes and Classes and Classes are the Apotheses and Classes and Classes and Classes are the Apotheses and Classes are the Apotheses and Classes are the Apotheses are the Apotheses and Classes are the Apotheses are the Apot

Hashell'98 module of about 300 lines, which is in near pass prop\_revilep to quickCheck to test the property in 100 random cases:

Testing Monadic Code with QuickCheck Doing so exposes at office-that the agencerty is not taked Whe-baloes printed It is nonciously difficult to say how effective a toutsenter-example to the claim. [1,-1] being the value of xx, and [0] the val-

QuickCheck Testing for Fun and Profit

John Horbes Citalmers University of Technology, S-41296 Gothenburg. Sweden

M. Harris (Ed.), PADL 2007, LNCN 4354, pp. 1-32, 2007 M. Harrisan, Victor Berlin, Heidelberg 2007

#### Computing Laboratory, University of Kent, UK Abstract (H.Li, S.J.Thompson) Goest.ac.uk

at the University of Kent to support Erlang program refactoring; the for testing the implementation of Wrangler, QuickCirck is a specification writing properties in a restricted logic, and using the tool these properties are tested in randomly generated test cases This paper first gives overviews of Wrangler and Queiq QuickCheck,

QuickCheck is a previously published random testing tool cation language is sufficiently nowerful to reservest common make a link between program testing and the notion of oband nostconditions. We introd ionguage, which simplifies these

#### Zeeshan Lakhani

#### In the Year 2000...

We have designed a simple domain-specific language of testable specifications which the tester uses to define expected properties of the functions under test.

We have chosen to put distribution under the human tester's control, by defining a test data generation language. . .

We have taken two relatively old ideas, namely specifications as oracles and random testing, and found ways to make them easily available . . .

### Don't Write Tests<sup>2</sup>

- One Feature O(n)
- Pairs of Features O(n2) quadratic
- Triples of Features O(n3) cubic

# Thinking in Specifications<sup>3</sup>

- A "roundtrip", e.g encode/decode?
- An existing implementation with similar behavior
- A relationship between inputs/outputs?
- A set of client interactions with a platform

# Sample Properties<sup>3</sup>

- Reversing a list twice should equal the original list.
- Reversing and sorting a list should preserve its length.
- Popping an element from a queue should reduce its size by one

# Defining Truth?<sup>4</sup>

Invariant

### ?SOMETIMES(N,Prop)

A property which tests Prop repeatedly N times, failing only if all of the tests fail. In other words, the property passes if Prop sometimes passes. This is used in situations where test outcomes are non-deterministic, to search for test cases that consistently fail...

### fails(Prop::property()) -> property()

A property which succeeds when its argument fails. Sometimes it is useful to write down properties which do not hold (even though one might expect them to). This can help prevent misconceptions.

<sup>&</sup>lt;sup>4</sup>Notes on Erlang QC - http://bit.ly/14CmDBF

# The Only Sure Thing in Computer Science

- Everything is a tradeoff.<sup>5</sup>
  - QC Tradeoffs: Duration | Assertion Power

# A Personal Story

- ClojureWest 2014
- John Hugues / Reid Draper



### #BEZERKER

```
:: Release-Base and Release-Link are at the top of this file. This is
;; so Nedia-Link can correctly depend on Release-Link.
(def Release-New
   ;; TODO: This should probably be optional, but I really want cayno to send it
    ;; so leaving it as required just for now.
   :slug c/Slug
:name c/NonEmptyString
    ;; TODO: This should probably be optional, but I really want cayno to send it
    (s/optional-key :copyright) (s/maybe c/Localized-Map)
    (s/optional-key :images) c/Images
    (g/optional-key : created-with) c/Simple-Account-Link
    (s/optional-key :uploaded-with) c/Simple-Account-Link
    (s/optional-key :pro-id) s/Str
   (s/optional-key :pro-slug) s/Str))
(def Release-Existing
    :name c/NonEmptyString
    :label (s/maybe s/Account-Link)
    :copyright (s/maybe c/Localized-Map)
    media [(s/either Track-Link Video-Link)]
    ;; :total-plays s/Int
    :uploaded-with (s/maybe a/Account-Link)
    :pro-slug (s/maybe s/Str)
    ;; :listened-to c/Simple-Link
(def Playlist-Base
   (s/optional-key :tags) c/Tags
   (s/optional-key :duration-seconds) (s/maybe s/Int)))
(def Playlist-Link
    :url c/URL
    ;; :total-hearts s/Int
```

;; :total-plays s/Int

```
(def Mix-Existing
  (assoc Mix-Base
    :slug c/Slug
    :prior-slugs [c/Slug]
    :created-at sc/ISO-Date-Time
    :created-by a/Account-Link
    :artists [a/Account-Link]
    :release (s/maybe Release-Link)
    :label (s/maybe a/Account-Link)
    :images c/Images
    :hearted s/Bool
    :copyright (s/maybe c/Localized-Map)
    ;; : license c/Simple-Link
    :description c/Localized-Map
    :source-tracks [{:track Track-Link
                     :start-time-seconds s/Int}l
    :sharing c/Simple-Link
    :total-hearts s/Int
    ;; :total-plays s/Int
    :purchase (s/maybe c/Simple-Link)
    :created-with (s/maybe a/Account-Link)
    :uploaded-with (s/maybe a/Account-Link)
    :recorded-date (s/maybe sc/ISO-Date-Time)
    :hearted-by c/Simple-Link
    ;; : listened-to c/Simple-Link
   ))
```

### quickcheck in the wild

- test.check (clojure)
- Quickcheck Haskell
- Erlang Quickcheck (from QuviQ)
- ScalaCheck
- JSVerify
- FsCheck (for .NET)
- . . .

# Specification - The Transpose of a Transposed Matrix is the Original Matrix

```
(A^T)^T = A
(def transpose-of-transpose-prop
  (prop/for-all [m matrix-gen]
  (= m (transpose (transpose m)))))
(quick-check 50 transpose-of-transpose-prop)
;; Results:
{:result true, :num-tests 50, :seed 1405444353915}
```

```
(def matrix
  [[1 2]
     [3 4]])
(transpose matrix)
;; Results:
[[1 3]
  [2 4]]
```

```
(def matrix-gen
  (gen/such-that
  not-empty
   (gen/vector
    (gen/tuple gen/int gen/int gen/int))))
(gen/sample matrix-gen 10)
;; Results:
([[0 \ 0 \ 1]]
 [[0 -1 0]]
 [[0 2 1] [0 1 -2]]
 [[1 0 1] [0 1 -3] [0 4 2] [2 -3 4]]
 [[-3 5 -4] [3 -3 -2] [-2 3 -2] [-3 -5 -3] [0 -3 -1]]
 [[2 -5 -3] [-4 -3 5] [-4 -3 -4] [-4 4 3]]
 [[-4 4 5] [-4 2 0] [5 -6 0] [2 -3 5] [-6 -3 -5]]
 [[-5 2 -3] [-2 -2 5]]
 [[-1 3 -5] [5 -3 -2] [7 4 -7] [7 -3 -2] [1 -3 8]]
 [[-5 -3 -3] [2 7 -3]])
```

### gen/fmap

gen/fmap allows us to create a new generator by applying a function to the values generated by another generator<sup>6</sup>

```
(def gen1
  (gen/fmap (fn [n] (* n 2)) gen/nat))
(gen/sample gen1)
;; Results
(0 2 4 2 0 10 10 6 8 8)
```

<sup>6</sup>test.check docs - http://bit.ly/1xWxQ9R

### gen/bind

gen/bind allows us to create a new generator based on the value of a previously created generator<sup>6</sup>

• Generator a -> (a -> Generator b) -> Generator b

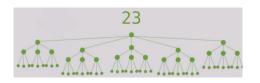
```
(def gen2
  (gen/bind gen1 (fn [v]
                 (gen/hash-map :codemash
                                (gen/return v)))
(gen/sample gen2)
;; Results
({:codemash 0} {:codemash 2} {:codemash 0}
{:codemash 4} {:codemash 6} {:codemash 10}
{:codemash 2} {:codemash 14} {:codemash 12}
{:codemash 16})
```

#### seed

#### Running Against the Same Set of Test Cases

# Shrinking<sup>7</sup>

Shrink trees are lazily generated for each generated result value



```
Remember our property (> 100 \text{ v})?
```

<sup>&</sup>lt;sup>7</sup>video - http://bit.ly/1zYwPwa

# schema->gen<sup>10</sup> for reasons

- Turn types, generics, schemas into generated data.
- My use-case: Prismatic Schema<sup>8, 9</sup>
- Work with Properties Testing API Workflow
- Regression Testing Out of the Box

<sup>8</sup>http://bit.ly/1BTaPW4

<sup>9</sup>another example - Herbert - http://bit.ly/1IxJs5V

<sup>&</sup>lt;sup>10</sup>http://bit.ly/1tSakMP

```
(def s-vector
    [(s/one s/Bool "first")
     (s/one s/Num "second")
     (s/one #"[a-z0-9]" "third")
     (s/optional s/Keyword "maybe")
     s/Int])
  ;; (true 3.0 "r"
  ;; :_1:r98l:Y!:npG-*:ZLyx4*+?:+I7:y08577B5D:392_:!1-+2:8-aMu7
  :: 1 7 3 -4)
  (def s-hashmap-with-hashmap
   {:foo s/Int
     :baz s/Str
     :bar {:foo s/Int}
     :far {(s/optional-key :bah) s/Bool}
     s/Keyword s/Num})
;; {:foo 0,
;; :baz "%?I\"".
;; :bar {:foo 9},
;; :far {:bah false},
:: 1.0,
   :0*37L:?K7+43:M?!?U_DIl*:GS90Ky**11 2.0}
  (s/check s-hashmap-with-hashmap datum)
```

#### A more realistic shrink

• Trying to Model Recursive Data Types<sup>11</sup>

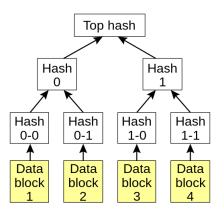
<sup>11</sup>http://bit.ly/1sc221b

```
Multiple Dispatch
(defmethod schema->gen* schema.core.One
  [e]
  (schema->gen (:schema e)))
(defmethod schema->gen* schema.core.RequiredKey
  [e]
  (gen/return (:k e)))
(defmethod schema->gen* schema.core.OptionalKey
  [e]
  (gen/return (:k e)))
(defmethod schema->gen* schema.core.Maybe
  [e]
  (gen/one-of
   [(gen/return nil)
    (schema->gen (:schema e))]))
```

#### Composing Generators

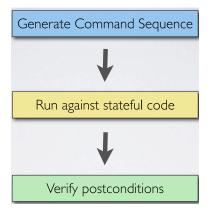
```
(defmethod schema->gen* clojure.lang.Sequential
  [e]
  (let [[ones [repeated]]
        (split-with #(instance? schema.core.One %) e)
        [required optional]
        (split-with (comp not :optional?) ones)]
    (g/apply-by
     (partial apply concat)
     (g/one-of
      (apply gen/tuple (map schema->gen required))
      (g/apply-by
       (partial apply concat)
       (apply gen/tuple
              (map schema->gen
                   (concat required optional)))
       (if repeated
         (gen/vector (schema->gen repeated))
         (gen/return [])))))))
```

# Changing Gears



#### eqc statem

Model State Transitions (as a FSM) -> Assert Against Implementation 12



<sup>12@</sup>jtuple - http://bit.ly/1t0fzaR

#### Side Effects

- Requires knowledge about the context and its possible histories<sup>13</sup>
- Symbolic values are generated during test generation and dynamic values are computed during test execution
  - dynamic state is computed at runtime
- next\_state callback operates during both test generation and test execution

<sup>13</sup>http://bit.ly/14CyorP

```
XX -----
%% Notes
YY -----
%% [1] Earle, Clara Benac, and Lars-Ake Fredlund. "Testing Java with QuickCheck."
%% This is a very basic eac statem test that I've updated a bit, dealing with
%% adding/cons'ing to a list and making sure those added values are members of a
%% list.
YY -----
%% Code
YY -----
-module(stateful sm).
-compile(export_all).
-include_lib("eqc/include/eqc.hrl").
-include_lib("eqc/include/eqc_statem.hrl").
setup() ->
  io:format("Setup Components If Need Be.~n"),
  ok.
cleanup() ->
  io:format("TearDown Components if Need Be.~n").
  ok.
test() ->
  test(100).
test(N) ->
  setup().
  try eqc:quickcheck(numtests(N, prop_codemash()))
  after
     cleanup()
   end.
```

```
%% Initialize State
initial_state() -> [].
%% ----- Grouped operator: add
%% @doc add_command - Command generator
-spec add_command(S :: egc_statem:symbolic_state()) ->
        eqc_gen:gen(eqc_statem:call()).
add command(S) ->
   {call, ?MODULE, add, [S, nat()]}.
%% @doc add_pre - Precondition for add
-spec add pre(S :: egc statem:symbolic state().
              Args :: [term()]) -> boolean().
add_pre(S, _Args) ->
   S /= undefined.
%% @doc add_next - Next state function
-spec add_next(S :: eqc_statem:symbolic_state(),
              V :: eqc_statem:var(),
               Args :: [term()]) -> eqc_statem:symbolic_state().
add next(S. Value, [ , N]) ->
    [NIS].
%% @doc add_post - Postcondition for add
-spec add post(S :: egc statem:dynamic state().
               Args :: [term()], R :: term()) -> true | term().
add_post(S, [_, N], Res) ->
    [N|S] = := Res.
%% @doc - Perform add action
-spec add(list(), non_neg_integer()) -> list().
add(AList, N) ->
    [N|AList].
```

```
%% ----- Grouped operator: is member
%% @doc is_member_command - Command generator
-spec is_member_command(S :: eqc_statem:symbolic_state()) ->
        eqc_gen:gen(eqc_statem:call()).
is_member_command(S) ->
    {call, ?MODULE, is member, [S, nat()]}.
%% @doc is_member_pre - Precondition for is_member
-spec is_member_pre(S :: egc_statem:symbolic_state(),
                    Args :: [term()]) -> boolean().
is member pre(S. Args) ->
    S /= undefined
%% @doc is_member_next - Next state function
-spec is_member_next(S :: eqc_statem:symbolic_state(),
                     V :: egc statem:var().
                     Args :: [term()]) -> eqc_statem:symbolic_state().
is_member_next(S, _Value, _Args) ->
    S.
%% @doc is_member_post - Postcondition for is_member
-spec is member post(S :: egc statem:dvnamic state().
                     Args :: [term()], R :: term()) -> true | term().
is_member_post(S, [_, N], Res) ->
    Res == lists:member(N, S).
%% @doc - Perform is_member action
-spec is_member(list(), non_neg_integer()) -> boolean().
is_member(S, N) ->
    lists:member(N, S).
```

```
Invariants
-spec invariant(eqc_statem:dynamic_state()) -> boolean().
invariant(S) when length(S) >= 0 ->
   true;
invariant(S) when length(S) > 0 ->
   FirstNum = hd(S).
    is_number(FirstNum) andalso FirstNum >= 0;
invariant() ->
   false.
%% Property Test
prop_codemash() ->
   ?FORALL(Cmds, commands(?MODULE),
            aggregate(command_names(Cmds),
                      begin
                      {H, S, Res} = run_commands(?MODULE, Cmds),
                      pretty_commands(?MODULE,
                                       Cmds,
                                       {H, S, Res},
                                       Res =:= ok)
                      end)).
```

```
stateful sm:test().
%% Results:
%% Licence for Basho reserved until
%% {{2015,1,8},{16,15,57}}
%% OK, passed 100 tests
%% 51.6% {stateful_sm, is_member, 2}
%% 48.4% {stateful_sm,add,2}
%% true
```

# WIP: Modeling KV-YZ HashTree in Riak

```
%% ----- Grouped operator: start uz tree
%% @doc start_yz_tree_command - Command generator
-spec start_yz_tree_command(S :: eqc_statem:symbolic_state()) ->
        eqc_gen:gen(eqc_statem:call()).
start_vz_tree_command(_S) ->
    {call, ?MODULE, start_yz_tree, []}.
%% @doc start_yz_tree_pre - Precondition for generation
-spec start_yz_tree_pre(S :: eqc_statem:symbolic_state()) -> boolean().
start_yz_tree_pre(S) ->
   S#state.yz_idx_tree == undefined.
-spec insert_kv_tree(sync|async, obj(), {ok, tree()}) -> ok.
insert_kv_tree(Method, RObj, {ok, TreePid}) ->
    {Bucket, Key} = eqc_util:get_bkey_from_object(RObj),
   Items = [{void, {Bucket, Key}, RObj}],
    case Method of
        sync ->
            riak_kv_index_hashtree:insert(Items, [], TreePid);
        asvnc ->
            riak_kv_index_hashtree:async_insert(Items, [], TreePid)
    end.
```

### We Talking About Tests? Tests?

- Automating Selenium Actions<sup>14</sup>
- Generating Models
- Sample Data
- Investigation
- Assertions

# Going Further

- Formal Specifications Thinking for Programmers<sup>15</sup>
- Molly Peter Alvaro<sup>16</sup>
  - A system for automatically detecting errors in a program with a correctness spec, the program, and malevolent sentience as input

<sup>&</sup>lt;sup>15</sup>http://bit.ly/1w5emM2

<sup>16</sup>http://bit.ly/1obnZLJ