

Fork Consistency & Certificate Transparency

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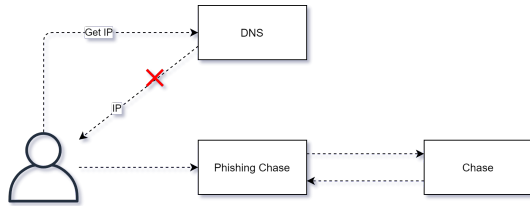
1. Why Certificate Transparency?

▼ HTTP

- Communication not encrypted
- [Man in the middle](#)
- Example

▼ HTTPS

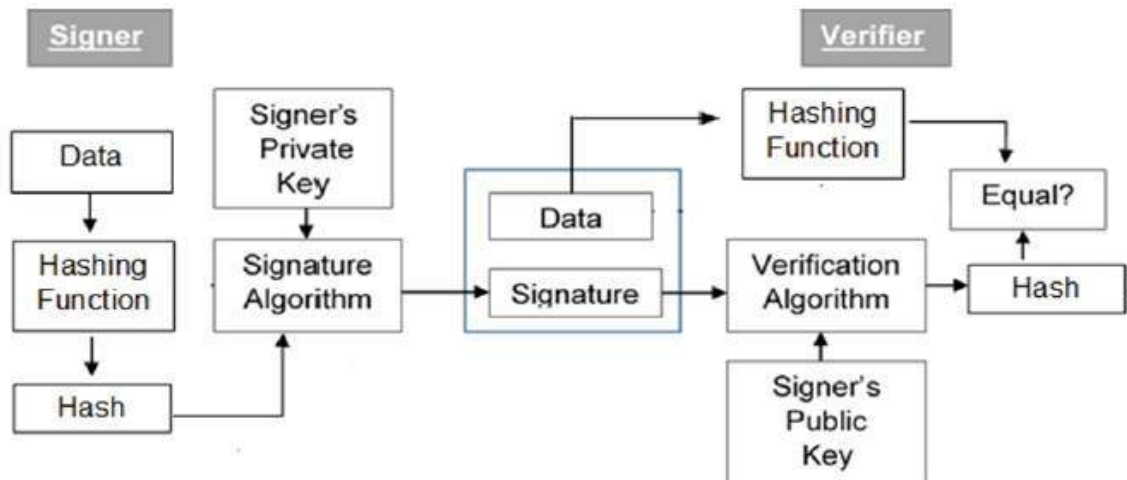
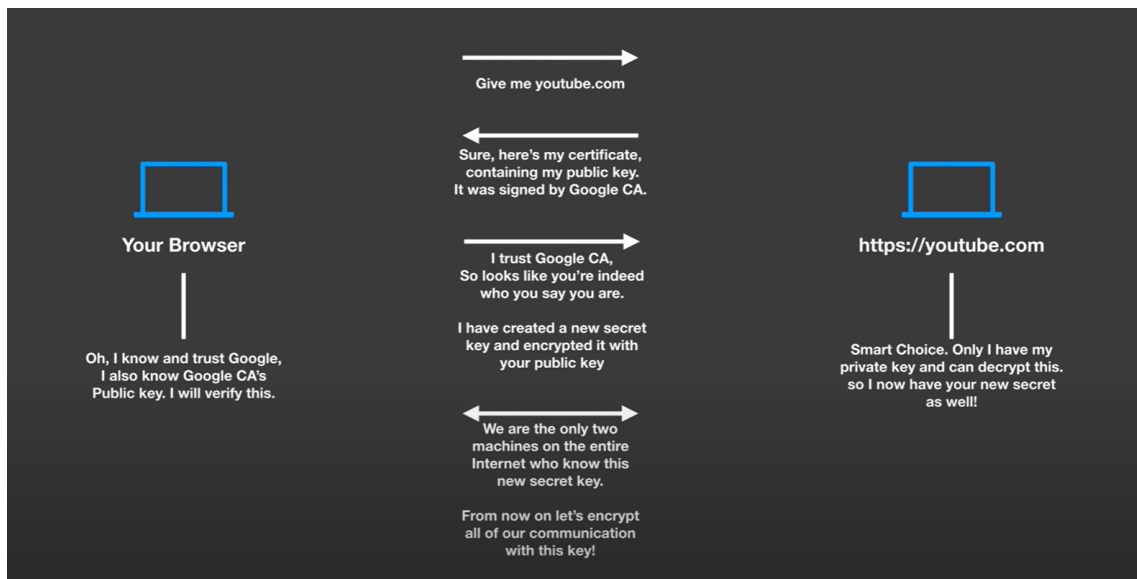
- Encrypted communication
- How it works?

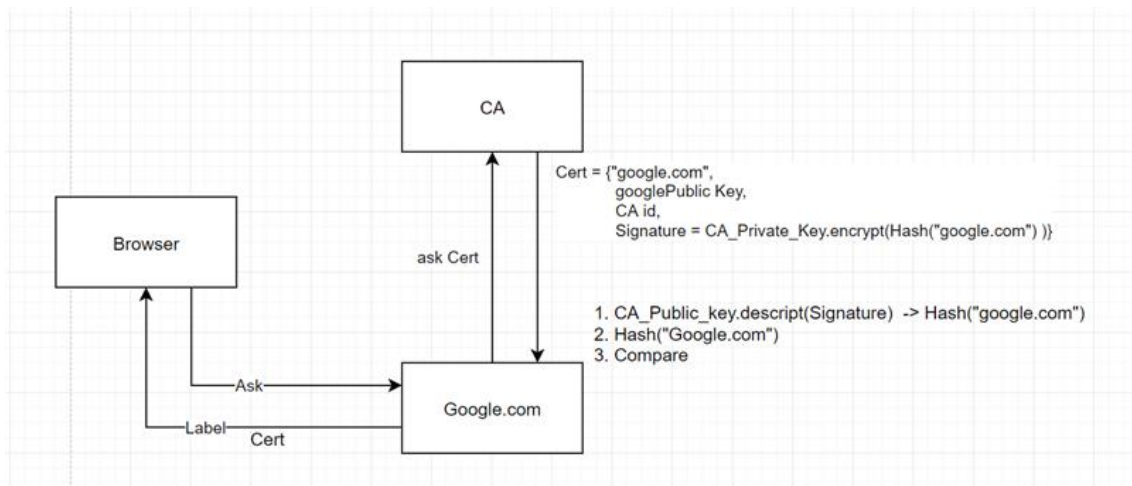


https://www.youtube.com/watch?v=T4Df5_cojAs

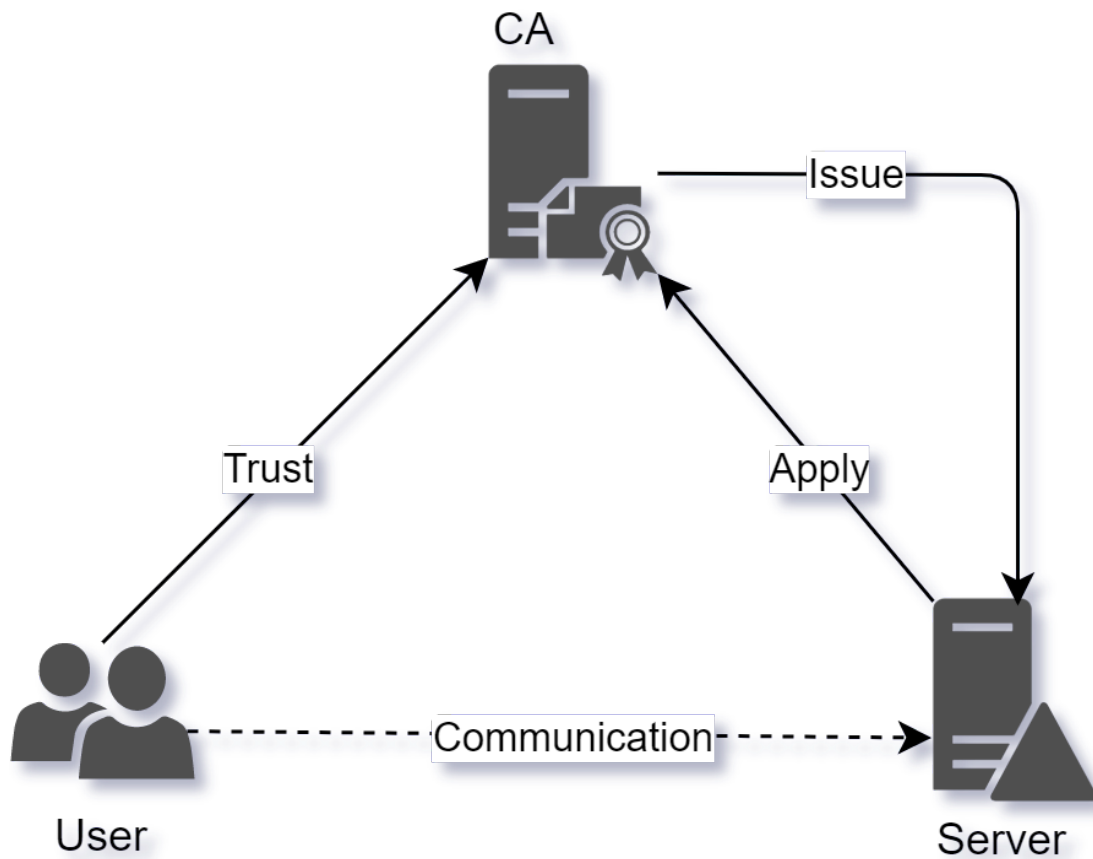
▼ Does HTTPS resolved all issues?

▼ Certificate Authority





▼ Components



▼ Problem



DigiNotar®

A VASCO COMPANY




Google

2. How CT Works?

How CT Works : Certificate Transparency

Certificate logs are append-only ledgers of certificates. Because they're distributed and independent, anyone can query them to see what certificates have been included and when. Because they're append-only, they are verifiable by Monitors. Organisations and individuals with the technical skills and capacity

 <https://certificate.transparency.dev/howctworks/>

3. Merkle Tree

▼ Basis: Cryptographic Hashes

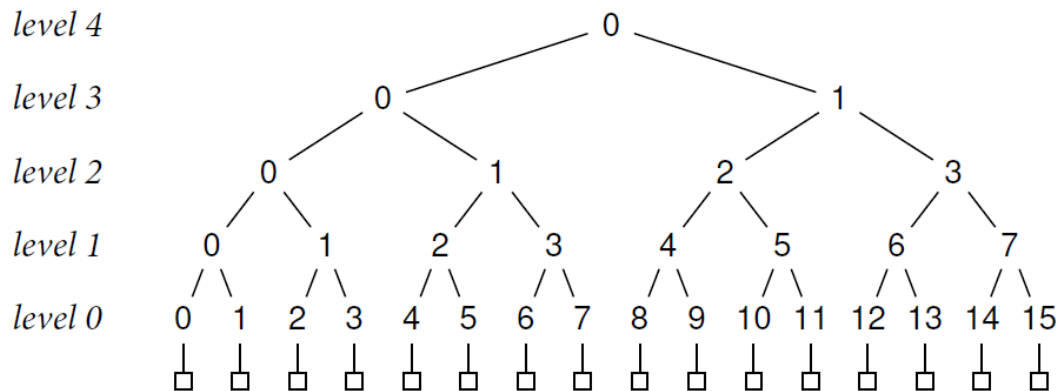
- A cryptographic hash function:
 - A deterministic function H that maps an arbitrary-size message M to a small fixed-size output $H(M)$. (Ex. SHA-256)

▼ Serving a Log

1. *Latest()* returns the **current log size** and top-level hash, cryptographically signed by the server for non-repudiation.
2. *RecordProof(R, N)* returns the proof that record R is

- It is infeasible in practice to produce a pair of distinct messages $M_1 \neq M_2$ with identical hashes $H(M_1) = H(M_2)$.
3. *TreeProof*(N, N') returns the proof that the tree of size N is a prefix of the tree of size N' .
 4. *Lookup*(K) returns the record index R matching lookup key K , if any.
 5. *Data*(R) returns the data associated with record R .

▼ How a Merkle tree got constructed?

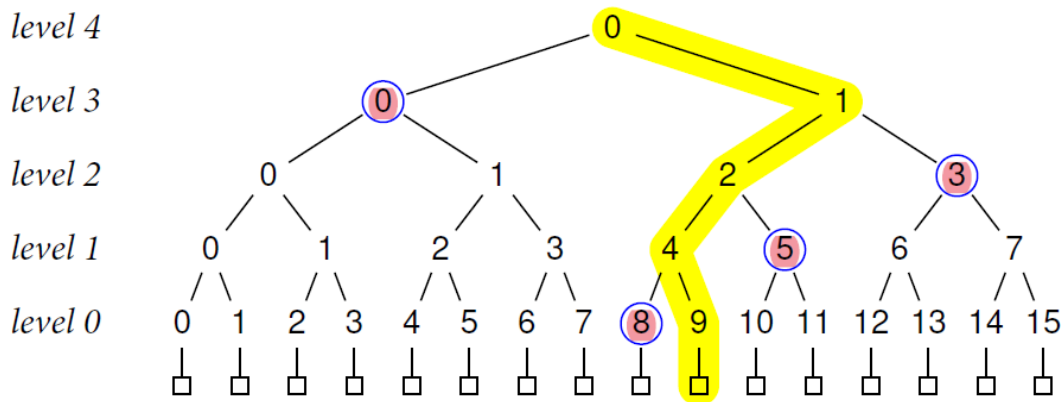


$$h(0, K) = H(\text{record}K)$$

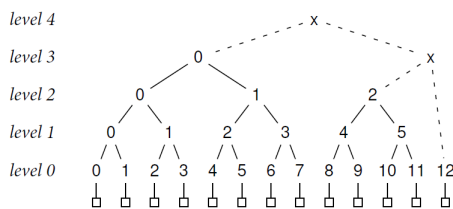
$$h(L + 1, K) = H(h(L, 2K), h(L, 2K + 1))$$

▼ How to verify a record is in the log?

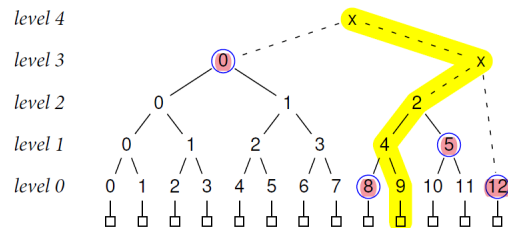
In general, the proof that a given is contained in the tree requires $\lg N$ hashes, one for each level below the root.



▼ Generalize the Merkle tree - Construction



▼ Generalize the Merkle tree - Verification



▼ Storing a log

- Log record data;
- Log index;
- Hashes;

▼ Optimization for storing

Maintaining lgN append-only files, each holding the sequence of hashes at one level of the tree.

▼ Verifying a Log

```
validate(bits B as record R):
    if R ≥ cached.N:
        N, T = server.Latest()
```

```
    if server.TreeProof(cached.N, N) cannot be verified:
        fail loudly
    cached.N, cached.T = N, T
    if server.RecordProof(R, cached.N) cannot be verified using B:
        fail loudly
    accept B as record R
```