01. Introduction to Convex Optimization –最优化导论

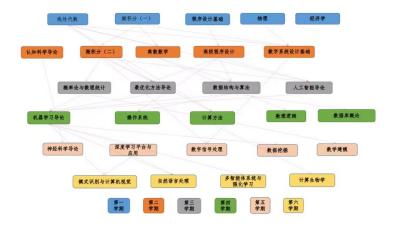
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0 About this course



0 About this course

1 助教

李紫微、张浩然、陈浩旭、(TBD)

2 课程时间

周一 3-4 节 (1-17 周); 周四 3-4 节 (双数周)

3 课程考核

- □ 期末考试(闭卷)60%
- □ 作业(两周一次)、课堂或期中考核 30%
- □ 课程项目(TBD, 文献阅读等)10%

4 参考书

"Convex Optimization" Stephen Boyd and Lieven Vandenberghe

1 Mathematical Optimization

- "Optimization" as a tool in engineering:
- "在给定的决策集中(通过某种方式)寻找最优解"

minimize
$$f_0(x)$$

subject to $f_i(x) \leq b_i, i = 1, 2, ..., m$

- \Box $x = (x_1, x_2, \dots, x_n)$: optimization variables
- \square $f_0: \mathbb{R}^n \to \mathbb{R}$: Objective function (objective)
- \square $f_i: \mathbb{R}^n \to \mathbb{R}, i = 1, ..., m$: Constraint functions (constraints)

Optimal solution x^* has smallest value of f_0 among all vectors that satisfy the constraints

Norm: a function $\|\cdot\|$ that satisfies:

- (1) $||x|| \ge 0$; ||x|| = 0 if and only if x = 0
- (2) $||tx|| = t||x||, t \in \mathbf{R}$
- (3) $||x+y|| \le ||x|| + ||y||$ (triangle inequality)

Notation: $\|\cdot\|$ is general (unspecified) norm; $\|\cdot\|_{\text{symb}}$ is particular norm

Distance: between vectors x and y as the length of their difference, *i.e.*,

$$\mathtt{dist}(x,y) = \|x - y\|$$

Some common norms on \mathbb{R}^n :

1. Sum-absolute-value, or l_1 -norm:

$$||x||_1 = |x_1| + |x_2| + \dots + |x_n|, \ x \in \mathbf{R}^n$$

2. l_2 -norm

$$||x||_2 = (x_1^2 + x_2^2 + \dots + x_n^2)^{\frac{1}{2}}, \ x \in \mathbf{R}^n$$

3. l_p -norm:

$$||x||_p = (|x_1|^p + |x_2|^p + \dots + |x_n|^p)^{\frac{1}{p}}, \ x \in \mathbf{R}^n$$

4. l_{∞} -norm (Chebyshev norm):

$$||x||_{\infty} = \lim_{p \to \infty} ||x||_p = \max\{|x_1|, \dots, |x_n|\}$$

5. For $P \in \mathbf{S}_{++}^n$, P-quadratic norm is

$$||x||_P = (x^T P x)^{\frac{1}{2}} = ||P^{\frac{1}{2}} x||_2$$



proof.

$$||x||_{p} = (|x_{1}|^{p} + |x_{2}|^{p} + \dots + |x_{n}|^{p})^{\frac{1}{p}} \ge |x_{\max}|$$

$$||x||_{p} \le (n \cdot |x_{\max}|^{p})^{\frac{1}{p}} = n^{\frac{1}{p}}|x_{\max}|$$

$$n^{\frac{1}{p}}|x_{\max}| \to |x_{\max}| \quad \text{when} \quad p \to \infty$$

. . .

Equivalence between norms:

1. Suppose that $\|\cdot\|_a$ and $\|\cdot\|_b$ are norms on \mathbb{R}^n , there exist positive constants α and β , for all $x \in \mathbb{R}^n$:

$$\alpha \|x\|_a \le \|x\|_b \le \beta \|x\|_a$$

2. If $\|\cdot\|$ is any norm on \mathbb{R}^n , then there exists a quadratic norm $\|\cdot\|_P$ for which

$$||x||_P \le ||x|| \le \sqrt{n} ||x||_P$$

holds for all x



proof. of first part:

It is sufficient to prove $\alpha ||x||_1 \le ||x||_b \le \beta ||x||_1$ Or $\alpha \le ||u||_b \le \beta$, where $u = x/||x||_1$ has norm $||u||_1 = 1$ Let

$$\alpha = \min_{\|u\|_1 = 1} \|u\|_b$$
$$\beta = \max_{\|u\|_1 = 1} \|u\|_b$$

 α , β exist for given bProof completes

2 Examples: least squares

Least squares formulation:

minimize
$$f_0(x) = ||Ax - b||_2^2 = \sum_{i=1}^k (a_i^T x - b_i)^2$$

- □ $A \in \mathbf{R}^{k \times n} (k \ge n), a_i^T$ 是矩阵 A 的行向量
- □ Solving the above problem is equivalent to solving $(A^T A)x = A^T b$. Hence, $x = (A^T A)^{-1} A^T b$
- \square Computation time: proportional to n^2k , less if structured

Properties of Least-squares problems:

- □ Easy to recognize and standard techniques
- □ Increase flexibility through weights, regularization terms

2 Examples: linear programming

Linear programming formulation:

minimize
$$f_0(x) = c^T x$$

subject to $a_i^T x \leq b_i, i = 1, \dots, m$.

- □ 向量 $c, a_1, \ldots, a_m \in \mathbf{R}^n, b_1, \ldots, b_m \in \mathbf{R}$ 为问题参数(parameter)
- □ No analytical solution (closed-form equation)
- \square Computation time: proportional to n^2m if $m \ge n$, less if structured

Properties of Least-squares problems:

- □ not as easy to recognize as least-squares problems
- □ A few standard tricks used to convert problems into linear programs (piecewise-linear functions)

2 Examples: convex optimization

Convex optimization formulation:

minimize
$$f_0(x)$$

subject to $f_i(x) \leq b_i, i = 1, ..., m$.

Object and constraint functions are both convex

$$f_i(\alpha x + \beta y) \le \alpha f_i(x) + \beta f_i(y),$$

if
$$\alpha + \beta = 1$$
, $\alpha \ge 0$, $\beta \ge 0$

□ Includes least-squares problems and linear programming as special cases

2 Examples: solving CO problems

- □ No analytical solution
- Reliable and efficient algorithms
- □ Computation time(roughly) proportional to $\max\{n^3, n^2m, F\}$, where F is cost of evaluating f_i 's and their first and second derivatives

When using convex optimization (why is it popular?)

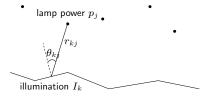
- Often difficult to recognize
- $\hfill\Box$ Many tricks for transforming problems into convex form
- $\hfill \square$ Surprisingly many problems can be solved via convex optimization

"In fact the great watershed in optimization isn't between linearity and nonlinearity, but convexity and nonconvexity."

——Rockafellar, 1993

2 Examples: lamps illuminating patches

m lamps illuminating n (small, flat) patches



 \square intensity I_k at patch k depends linearly on lamp powers p_j :

$$I_k = \sum_{j=1}^{m} a_{k,j} p_j, \quad a_{kj} = r_{kj}^{-2} \max\{\cos \theta_{kj}, 0\}$$

 \square Achieve desired illumination I_{des} with bounded lamp powers

minimize
$$\max_k |\log I_k - \log I_{\text{des}}|$$

subject to $0 \le p_i \le p_{\text{max}}, j = 1, \dots, m$.

2 Examples: lamps illuminating patches

How to solve?

- \square Use uniform power $p_i = p$, vary p
- □ Use least-squares

minimize
$$\sum_{k=1}^n (I_k - I_{\text{des}})^2$$

round p_j if $p_j < 0$ or $p_j > p_{\text{max}}$

□ Use weighted least-squares

minimize
$$\sum_{k=1}^{n} (I_k - I_{\text{des}})^2 + \sum_{j=1}^{m} w_j (p_j - p_{\text{max}}/2)^2$$

iteratively adjust weights w_j until $0 \le p_j \le p_{\max}$

□ Use linear programming

minimize
$$\max_{k=1,2,...,n} |I_k - I_{\text{des}}|$$

subject to $0 \le p_i \le p_{\max}, j = 1,..., m$



2 Examples: lamps illuminating patches

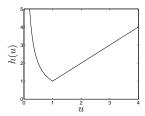
How to solve?

□ Use convex optimization: problem is equivalent to

minimize
$$f_0(p) = \max_{k=1,...,n} h(I_k/I_{\text{des}})$$

subject to $0 \le p_j \le p_{\max}, \ j=1,\ldots,m$

with $h(u) = \max\{u, 1/u\}$



 f_0 is convex because maximum of convex functions is convex

3 Examples in applications

Portfolio optimization

- □ Variables: amounts invested in different assets
- □ Constraints: budget, max./min. investment per asset, minimum return
- □ Objective: overall risk or return variance

Device sizing in electronic circuits

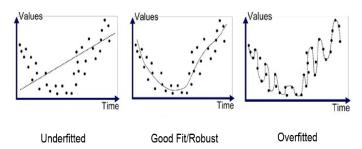
- □ Variables: device widths and lengths
- □ Constraints: manufacturing limits, timing requirements, maximum area
- □ Objective: power consumption

Data fitting

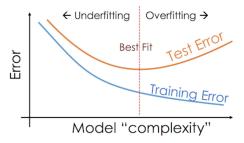
- □ Variables: model parameters
- □ Constraints: prior information, parameter limits
- □ Objective: measure of misfit or prediction error

3 Examples in applications

Overfitting/Underfitting



3 Examples in applications



Overfitting/Underfitting with different levels of complexity/training

4 Brief history of convex optimization

Theory(convex analysis): ca1900–1970

Algorithms

- □ 1947: simplex algorithm for linear programming (Dantzig)
- □ 1960s: early interior-point methods (Fiacco and McCormick, Dikin, ...)
- $\hfill\Box$ 1970s: ellipsoid method and other subgradient methods
- □ 1980s: polynomial-time interior-point methods for linear programming (Karmarkar1984)
- □ late 1980s-now: polynomial-time interior-point methods for nonlinear convex optimization (Nesterov and Nemirovski1994)

Applications

- □ Before 1990: mostly in operations research; few in engineering
- □ Since 1990: many new applications in engineering (control, signal processing, communications, circuit design, ...); new problem classes (semidefinite and second-order cone programming, robust optimization)