

# CS150A Quiz09

## FD Properties

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I'd like some properties for my functional dependencies please.

### 1) Select all the FD's that follow from Armstrong's Axioms

Hint: there's at least one  
Check all that apply.

$XZ \rightarrow YZ$   
 $YZ \rightarrow YW$

- ☒ if  $X \rightarrow Y$  and  $Z \rightarrow W$ , then  $XZ \rightarrow YW$
- ☒ if  $\tilde{X} \rightarrow \tilde{Y}$  and  $WY \rightarrow Z$ , then  $WX \rightarrow Z$
- ☐ if  $XZ \rightarrow Y$ , then  $X \rightarrow Y$
- ☒ if  $X \rightarrow YZ$ , then  $X \rightarrow Y$
- ☒ if  $X \rightarrow Y$  and  $X \rightarrow Z$ , then  $X \rightarrow YZ$

## FD Example

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We have a relation  $R(A, B, C, D, E)$ . We are told that the set of functional dependencies is

$F = \{E \rightarrow BD, A \rightarrow BC, C \rightarrow \underline{DE}, D \rightarrow C\}$ .

Find the attribute closures for each of the attributes. If the attribute closure for  $X$  was  $WXZ$ , you would fill in "WXZ" without quotes in the answer box.

### 2) $A^+$ :

ABCDE

### 3) $B^+$ :

B

4) C+:

CDEB

5) D+:

DC

6) E+:

EBDC

7) Select the attribute set(s) that are keys for relation R

Hint: there's at least one

Check all that apply.

- ☐ E
- ☒ A
- ☒ AD
- ☐ BCE
- ☒ ABCDE

8) The attribute closure of (BC)+ is equivalent to the attribute closure of (BD)+.

By equivalent we mean the intersection is equivalent to the union of both closure sets.

Mark only one oval.

- ☐ True
- ☒ False

9) Is relation R already in Boyce-Codd Normal Form (BCNF)?

Mark only one oval.

- ☒ Yes
- ☐ No

A superkey

## Normalization

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BCNF stands for Boyce-Codd Normal Form. For this question, assume the decomposition is performed using the algorithm described in lecture.

10) Decomposing a relation into BCNF does not always guarantee a lossless decomposition.

Mark only one oval.

- ☐ True
- ☒ False

11) Decomposing a relation into BCNF will always guarantee a dependency preserving decomposition.

Mark only one oval.

☐ True

☒ False

12) Relation  $R(A, B, C, D, E)$  is decomposed into  $R(A, C, D, E)$  and  $R(A, B, C)$  with the set of functional dependencies  $F = \{BC \rightarrow A, C \rightarrow D\}$ . Is this decomposition lossless?

Note: the decomposition might not follow the BCNF algorithm discussed in class.

Mark only one oval.

☐ Yes

☒ No

$\downarrow$   $\downarrow$   
 $A \approx AC.$   
 $AC \rightarrow ACD. \nexists R_1$   
 $\nexists R_2$