

# CS150A Database

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Today:

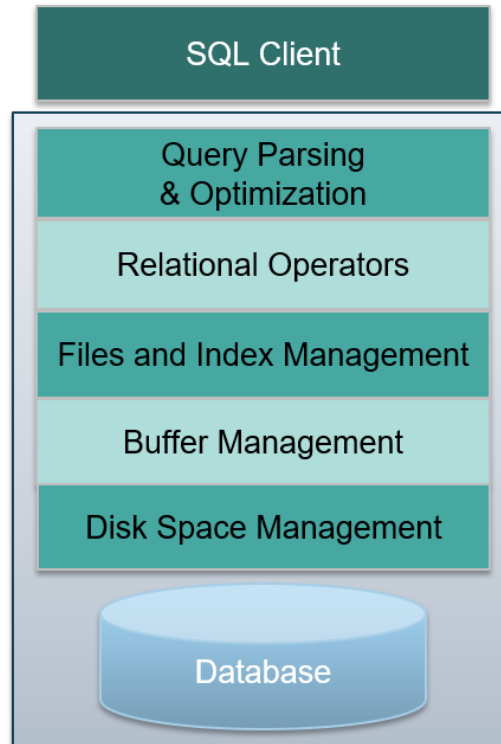
- Database Design I:
  - Entity-Relation Model

Readings:

- Database Management Systems (DBMS), Chapter 2

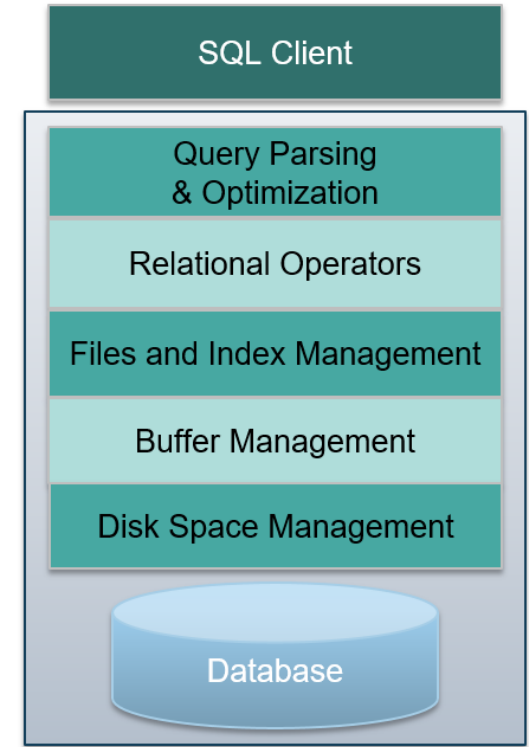
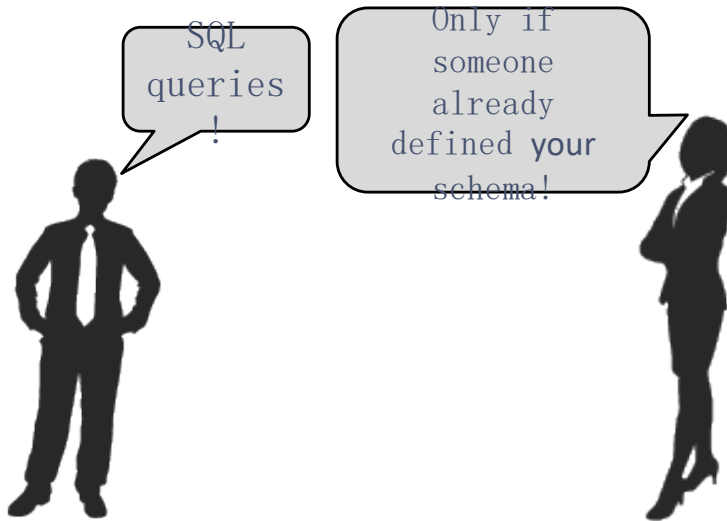
# Architecture of a DBMS

- Gives us a good sense of how to build a DBMS
- How about using one?



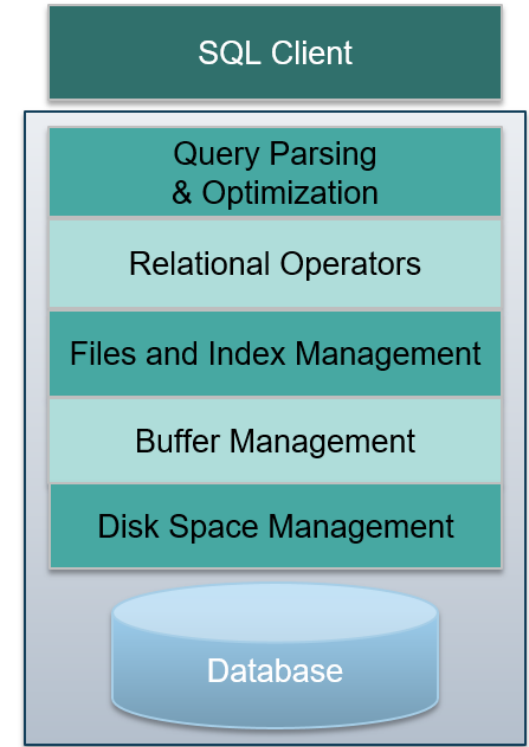
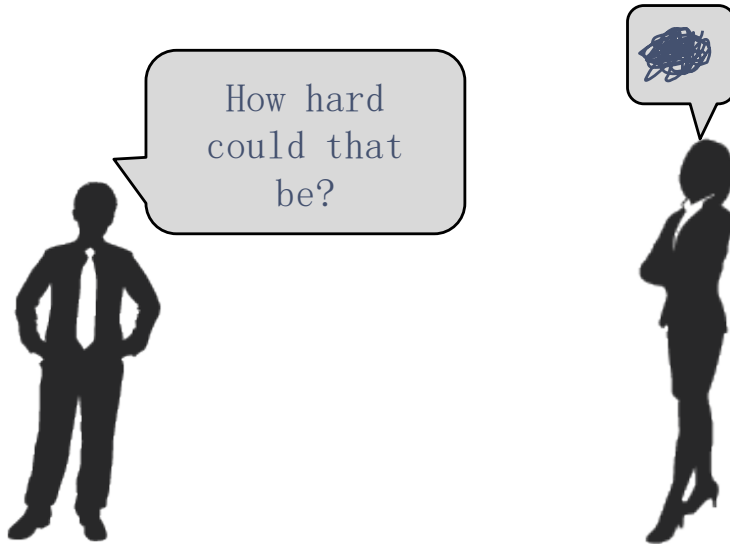
# Architecture of a DBMS, Pt 2

- Gives us a good sense of how to build a DBMS
- How about using one?



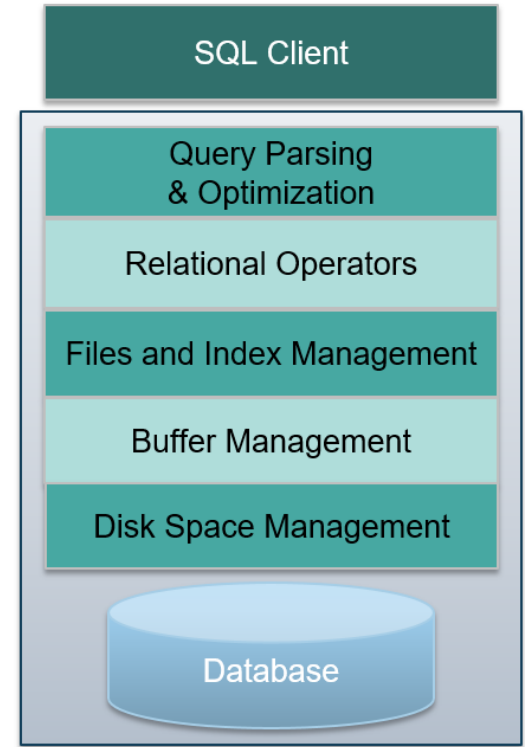
# Architecture of a DBMS, Pt 3

- Gives us a good sense of how to build a DBMS
- How about using one?



# Design of a Database

- Gives us a good sense of how to build a DBMS
- How about using one?
- Today let's talk about how to design a database
  - Not a database system



# Steps in Database Design

- **Requirements Analysis**
  - user **needs**; what must database do?
- **Conceptual Design**
  - *high level description (often done w/ER model)*
  - Object-Relational Mappings (ORMs: Hibernate, Rails, Django, etc)  
encourage you to program here
- **Logical Design**
  - translate ER into DBMS data model
  - ORMs often require you to help here too
- **Schema Refinement**
  - consistency, normalization
- **Physical Design** – indexes, disk layout
- **Security Design** – who accesses what, and how



You are here

# Describing Data: Data Models

- Data model: collection of concepts for describing data.
  - Relational model, hierarchical model, network model, ...
- Schema: description of a particular collection of data, using a given data model.
- Relational model of data
  - Main concept: relation (table), rows and columns
  - Every relation has a schema
    - describes the columns
    - column names and domains

# Levels of Abstraction

(UX)

**Users**

(programs connected to DB)

Views describe how users see the data.



Subsets / aggregation

log

Conceptual schema defines logical structure

View 1

View 2

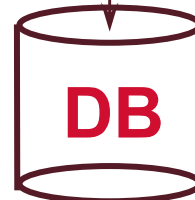
View 3

Conceptual Schema

Physical Schema

Physical schema describes the files and indexes used.

eg, B+





# Example: University Database



- Conceptual schema:

- Students(sid text, name text, login text, age integer, gpa float)
- Courses(cid text, cname text, credits integer)
- Enrolled(sid text, cid text, grade text)

- Physical schema:

- Relations stored as unordered files.
- Index on first column of Students.

B+

- External Schema (View):

- Course\_info(cid text, enrollment integer)

# Data Independence

- Insulate apps from structure of data  
*eg, add rows / cols*
- Logical data independence: *eg student + col* *eg, Student T*
  - Maintain views when logical structure changes *→ Student T1 + Student T2*
- Physical data independence:
  - Maintain logical structure when physical structure changes

# Levels of Abstraction, cont

**Users**



View 1

View 2

View 3

Logical data independence



Physical data independence



# Data Independence, cont

- Insulate apps from structure of data
- **Logical data independence:** *schema*
  - Maintain views when logical structure changes
- **Physical data independence:**
  - Maintain logical structure when physical structure changes *schema* *Bt, index..*
- Q: Why particularly important for DBMS?
  - Because databases and their associated applications persist

# Hellerstein's Inequality

eg. android app

$$\frac{dapp}{dt} \ll \frac{denv}{dt}$$

app : month update

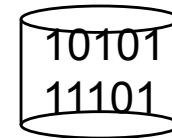
env : second update.  
(date)

Data independence is most important when the rate of change of your environment exceeds the rate of change of your applications.

# Data Models

- Connect concepts to bits!
- Many models exist
- We will ground ourselves in the Relational model
  - clean and common
  - generalization of key/value
- Entity-Relationship model also handy for design
  - Translates down to Relational


**Student** (*sid: string, name: string, login: string, age: integer, gpa: real*)



# Entity-Relationship Model

- Relational model is a great formalism
  - But a bit detailed for design time
  - Too fussy for brainstorming
  - Hard to communicate to “customers”
- Entity-Relationship model: a graph-based model
  - can be viewed as a graph, or a veneer over relations
    - “feels” more flexible, less structured
  - corresponds well to “Object-Relational Mapping”
    - (ORM) SW packages
    - Ruby-on-Rails, Django, Hibernate, Sequelize, etc.

# Steps in Database Design, again

- Requirements Analysis
  - user needs; what must database do?
- Conceptual Design
  - *high level description (often done w/ER model)*  You are here
  - ORM encourages you to program here
- Logical Design
  - translate ER into DBMS data model
  - ORMs often require you to help here too
- Schema Refinement
  - consistency, normalization
- Physical Design – indexes, disk layout
- Security Design – who accesses what, and how



# Conceptual Design

- What are the entities and relationships?
  - And what info about E' s & R' s should be in DB?
- What integrity constraints ( “business rules” ) hold?
- ER diagram is the “schema”
- Can map an ER diagram into a relational schema.
- Conceptual design is where the data engineering begins
  - If you' re familiar with the jargon, these are the “models” of the MVC pattern in ORMs

# ER Model Basics: Entities

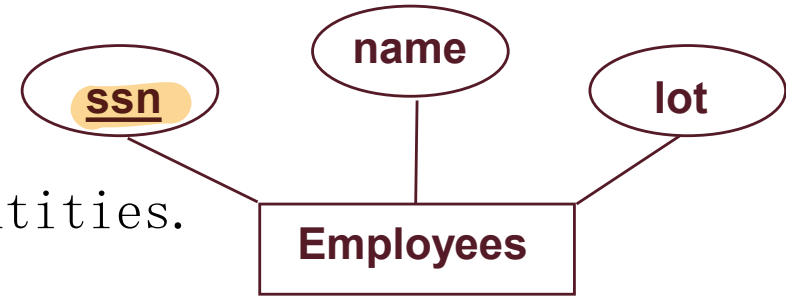


- Entity:

- A real-world object described by a set of attribute values.

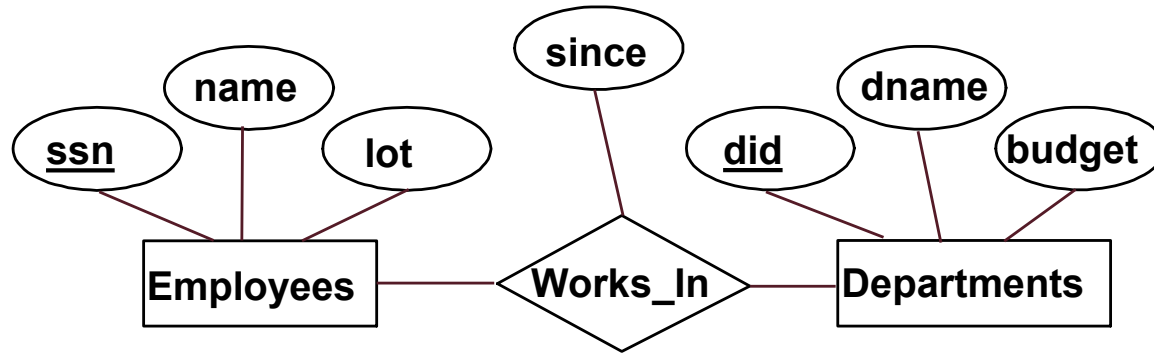
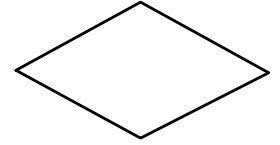
- Entity Set: A collection of similar entities.

- E.g., all employees.
- All entities in an entity set have the same attributes.
- Each entity set has a key (underlined)
- Each attribute has a domain



domain and Type

# ER Model Basics: Relationships



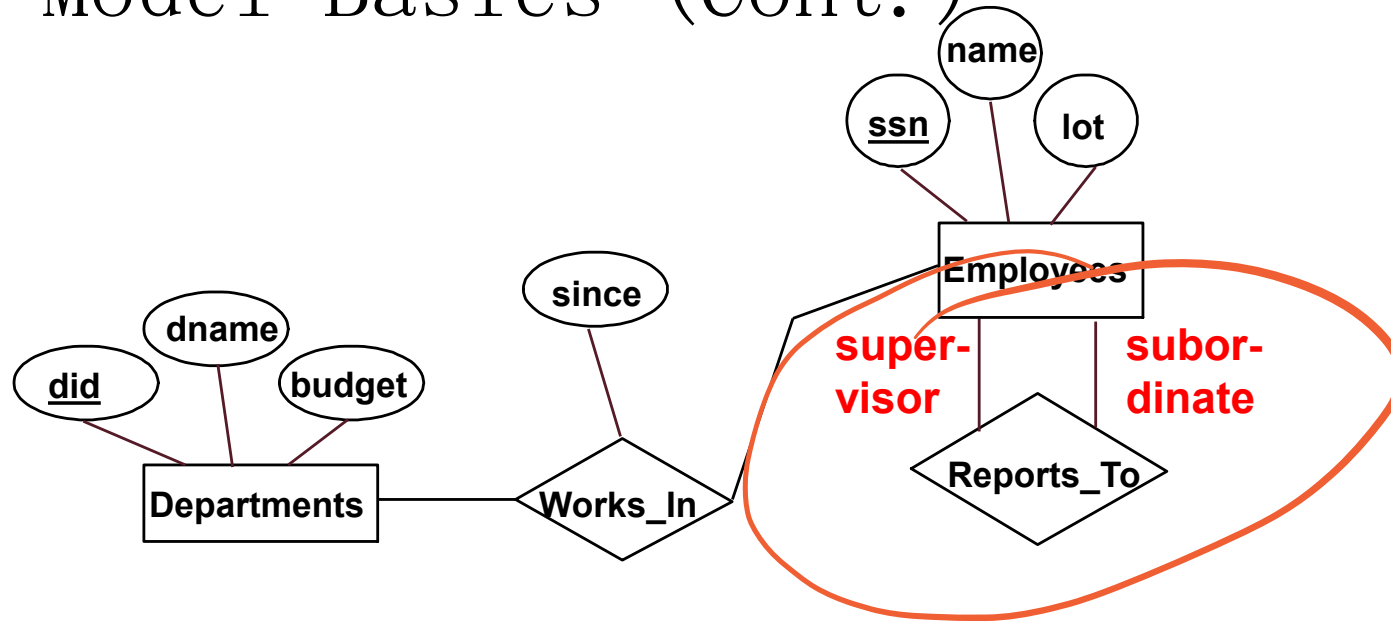
**Relationship:** Association among two or more entities.

- E.g., Attishoo works in Pharmacy department.
- Relationships can have their own attributes.

**Relationship Set:** Collection of similar relationships.

- An n-ary relationship set  $R$  relates  $n$  entity sets  $E_1 \dots E_n$

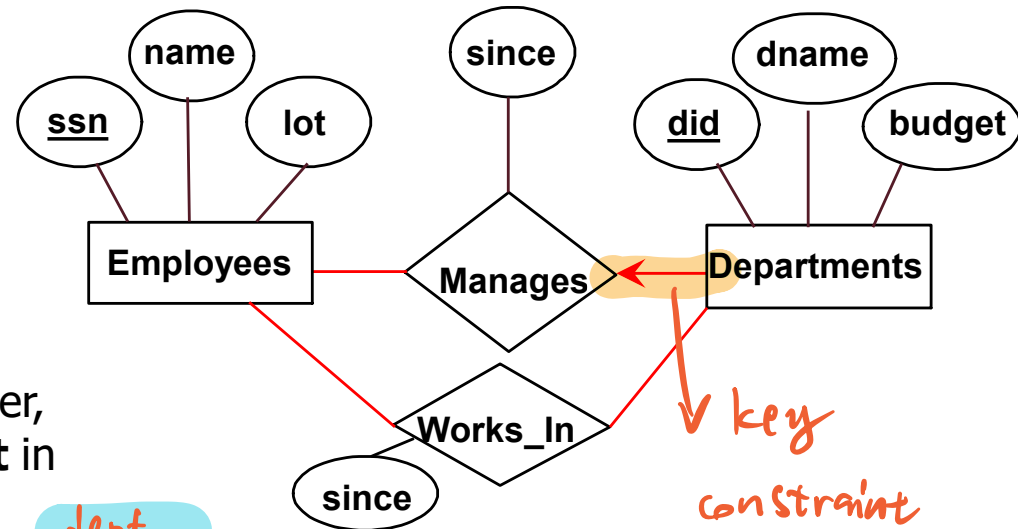
# ER Model Basics (Cont.)



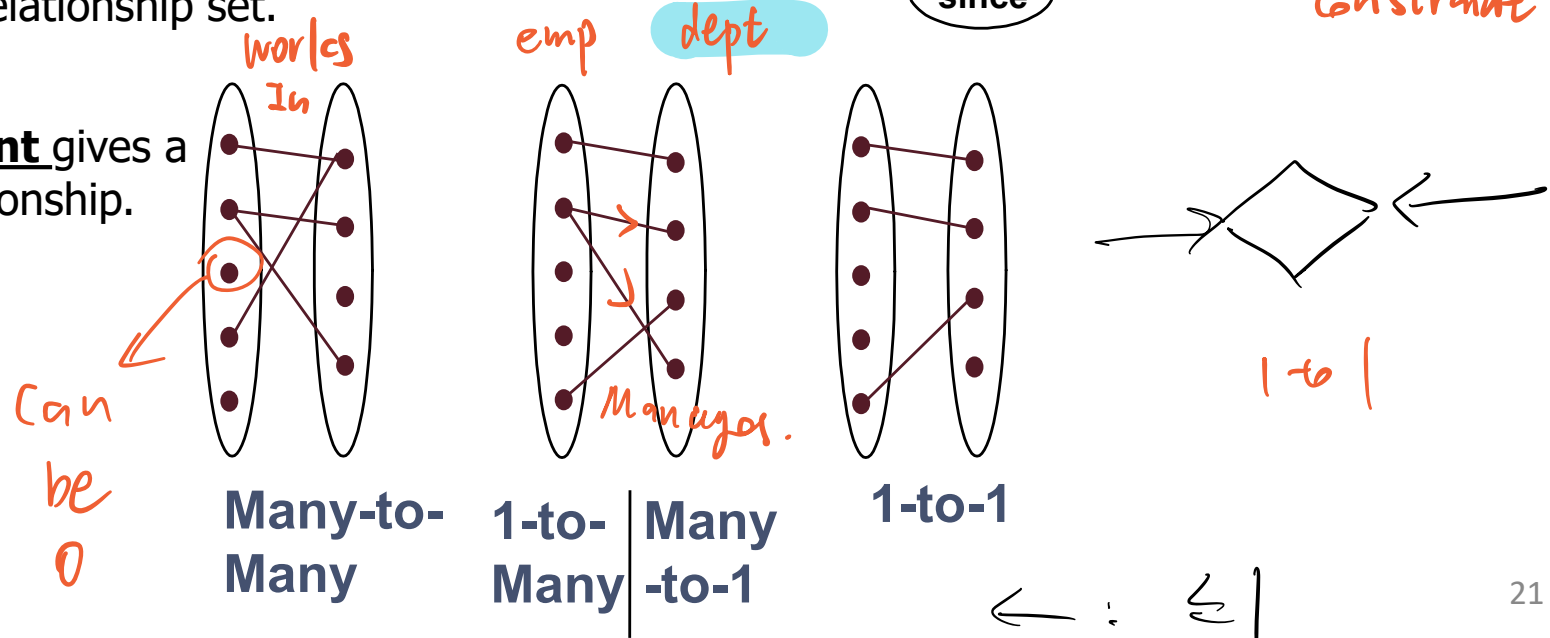
Same entity set can participate in different relationship sets, or in different “roles” in the same relationship set.

# Key Constraints

- An employee can work in **many** departments; a dept can have **many** employees.
- In contrast, each dept has **at most one** manager, according to the key constraint on **Department** in the **Manages** relationship set.



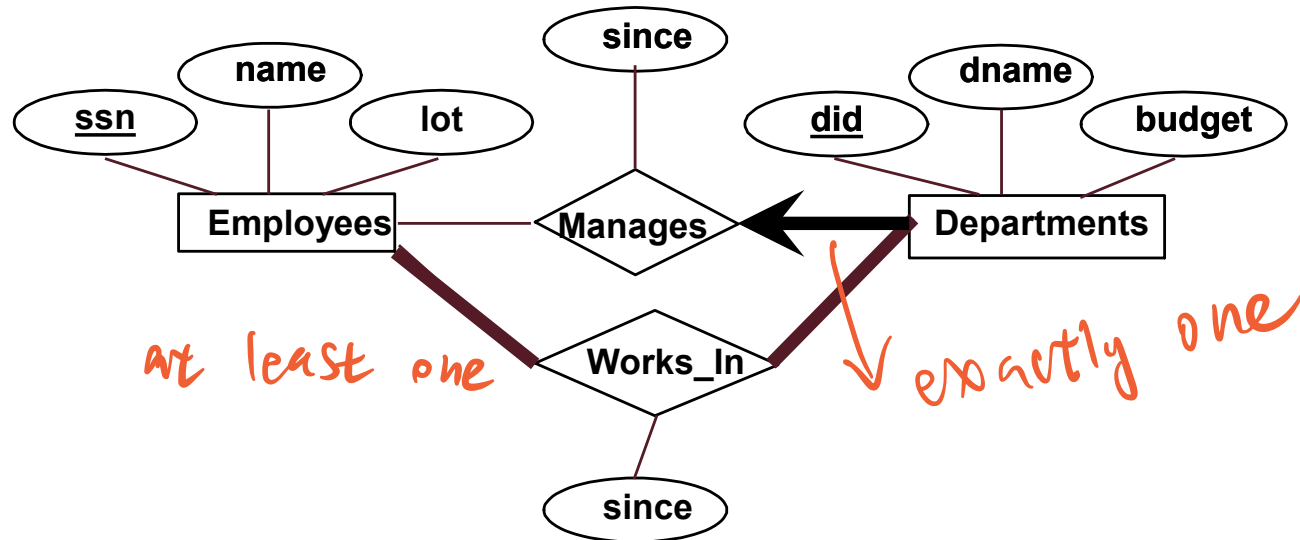
- A key constraint gives a **1-to-many** relationship.



# Participation Constraints

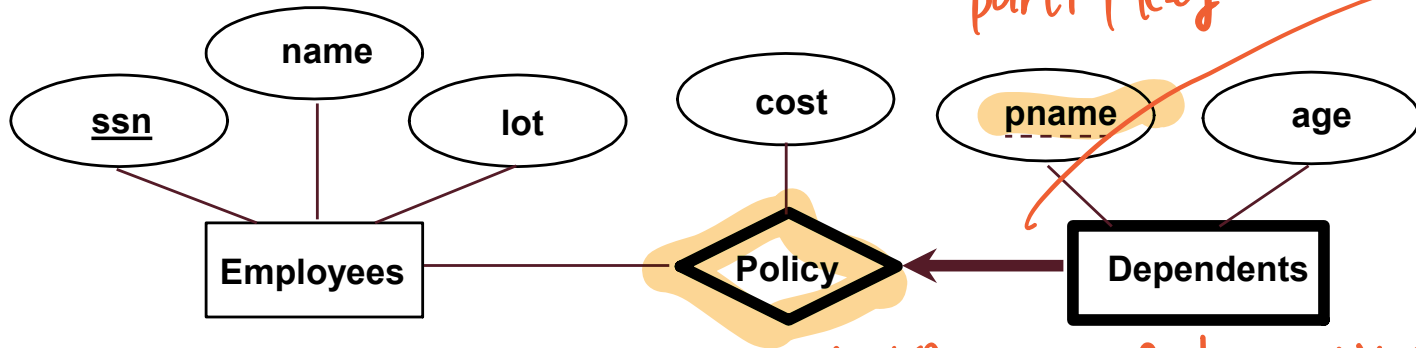
*no 0.*

- Does every employee work in a department?
- If so: a participation constraint
  - participation of Employees in Works\_In is total (vs. partial)
  - What if every department has an employee working in it?
- Basically means at least one.



# Weak Entities

- A **weak entity** can be identified uniquely only by considering the primary key of another (owner) entity.
  - Owner entity set and weak entity set must participate in a **one-to-many** relationship set (one owner, many weak entities).
  - Weak entity set must have **total participation** in this identifying relationship set.



partial key

① 1-to-many  
② total participation  
(exactly one)

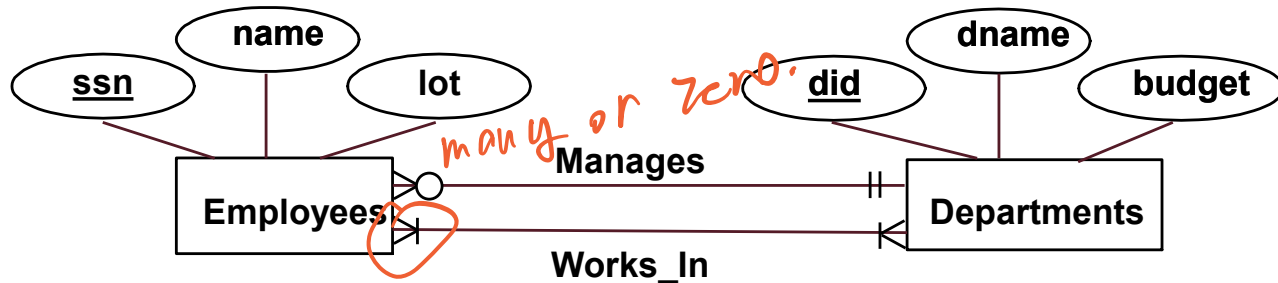
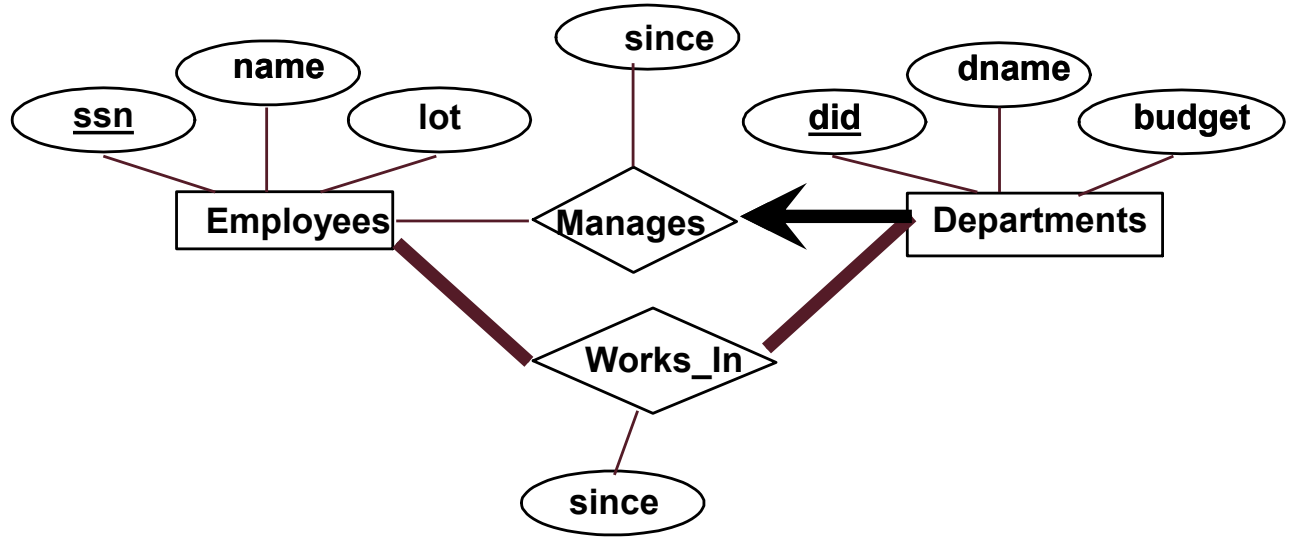
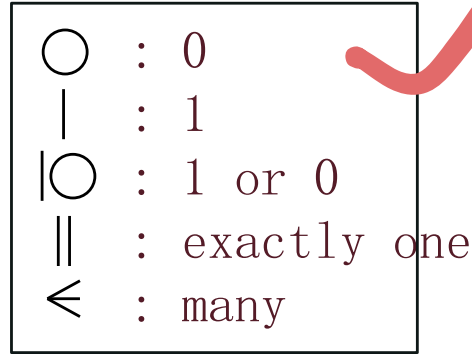
total

weak entity

- Weak entities have only a “partial key” (dashed underline) → follow policy to

# FYI: Crow's Foot Notation

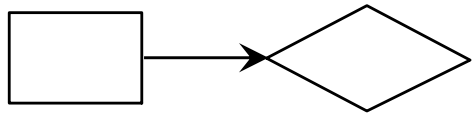
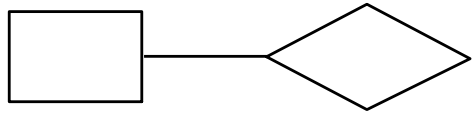
rent key  
(SSN here)



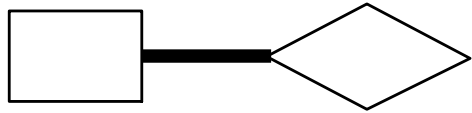
at least one 1 or many



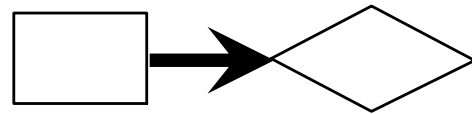
# Translating constraints across notations



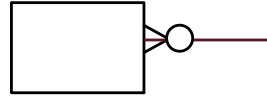
max 1



least 1



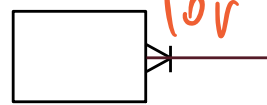
exact 1



0 to many



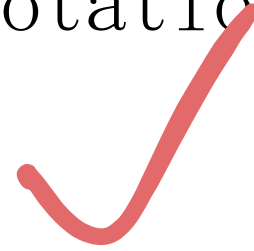
1 or zero



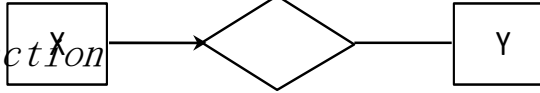
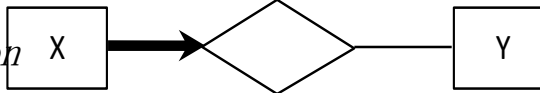
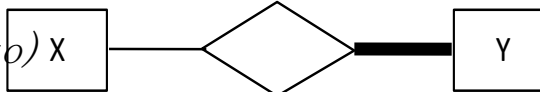
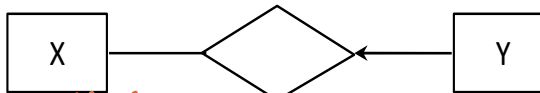
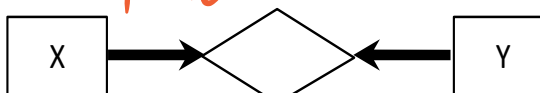
1 or many



exact 1



# Translation to Math Terminology on Relations

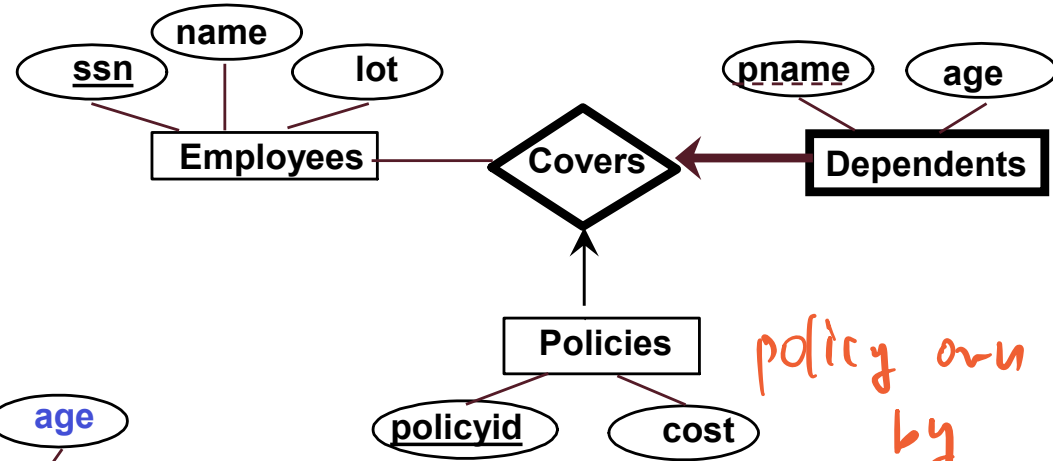
- Relation  $R(X, Y)$  is a *(partial) function* 
- Relation  $R(X, Y)$  is a *total function* 
- Relation  $R(X, Y)$  is *surjective (onto)* 
- Relation  $R(X, Y)$  is *injective (1-1)*    
 *not "func"*
- Relation  $R(X, Y)$  is a *bijection* 

# Binary vs. Ternary Relationships

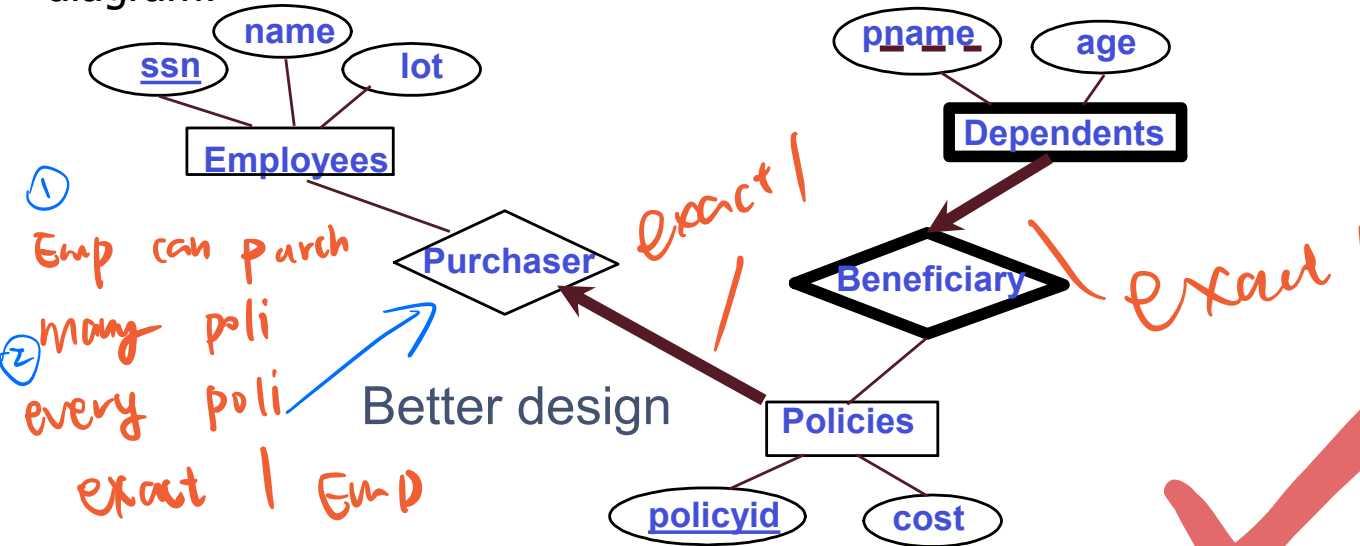
If each policy is owned by just 1 employee:

**Key constraint on Policies would mean policy can only cover 1 dependent!**

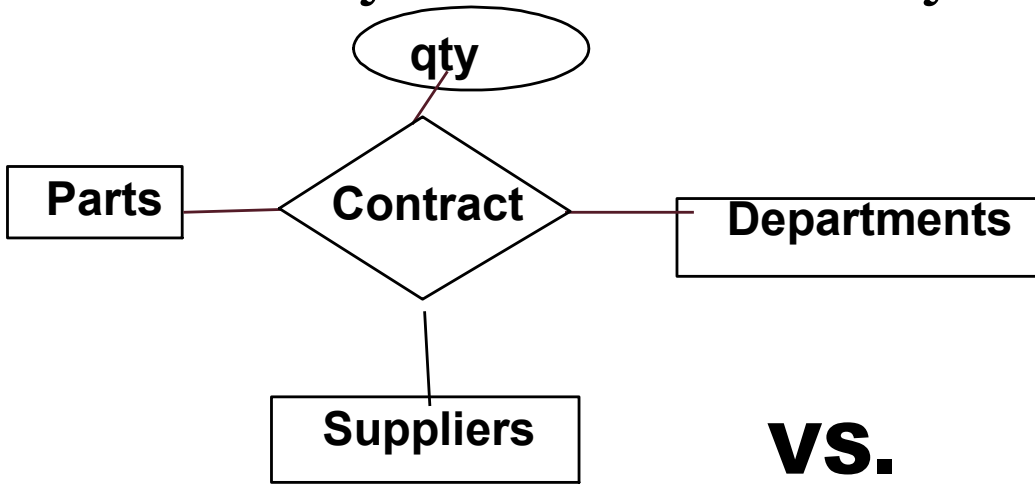
Think through **all** the constraints in the 2nd diagram!



policy own  
by  
1 emp



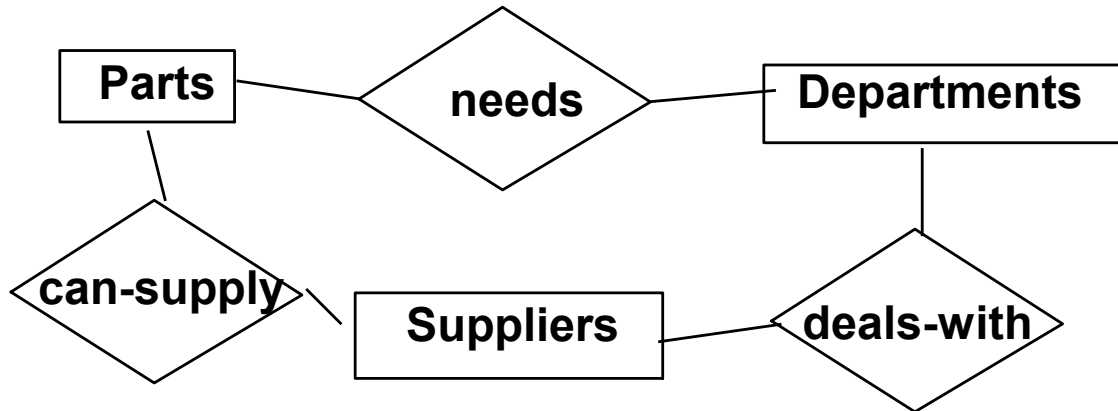
# Binary and Ternary Relationship (cont)



S “can-supply” P, D “needs” P, and D “deals-with” S does not imply that D has agreed to buy P from S.

How do we record *qty*?

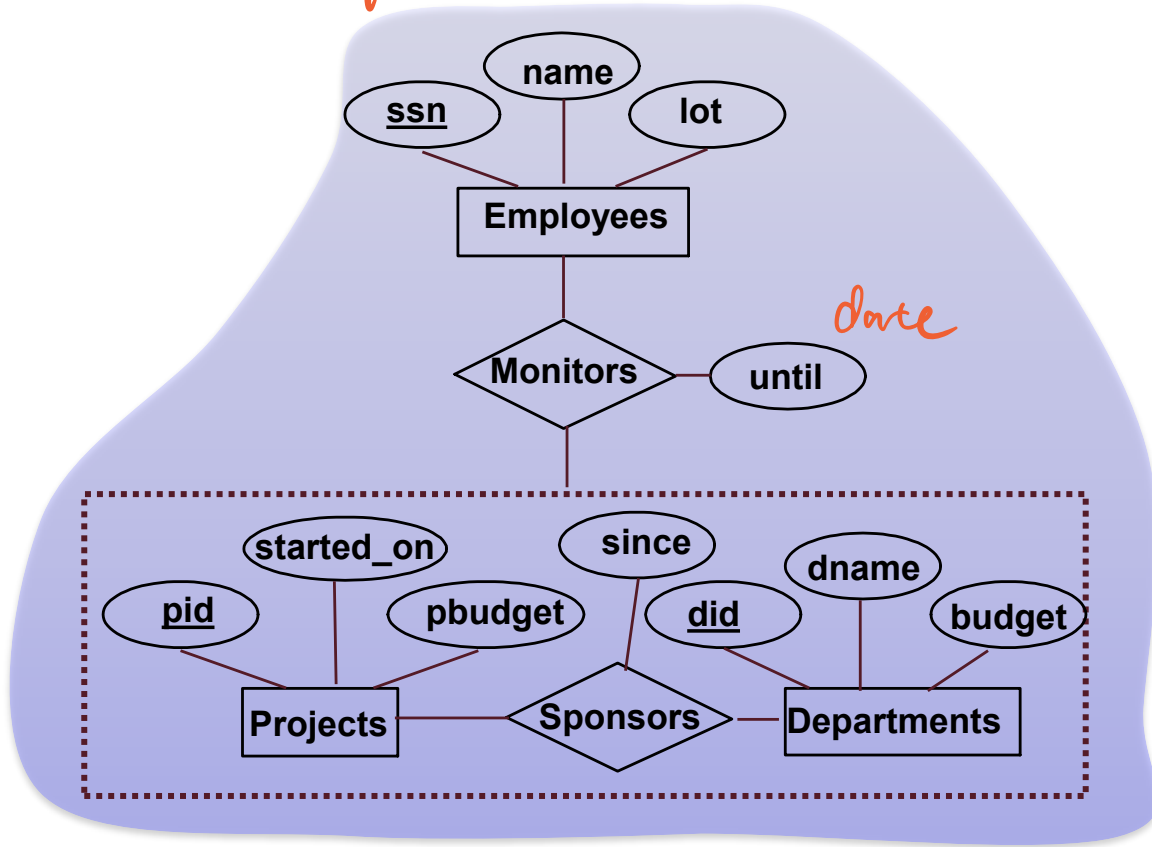
**VS.**



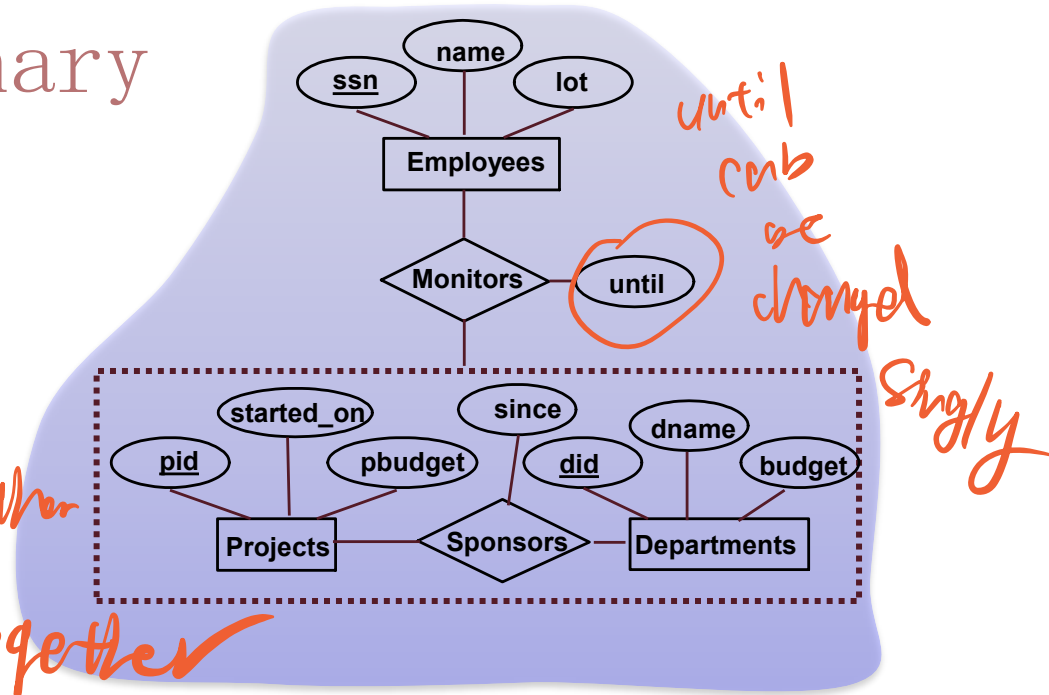
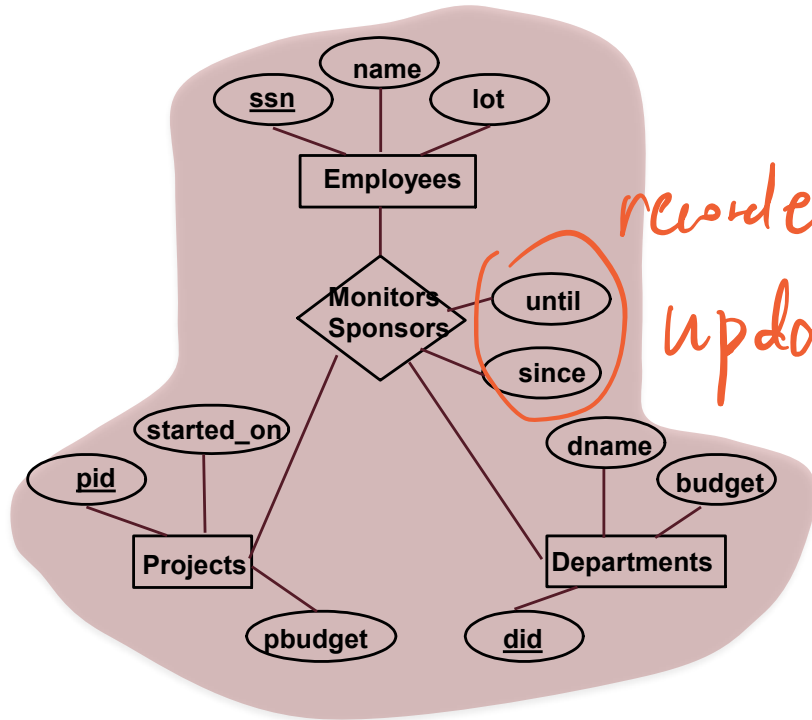
# Aggregation

Not sql aggr

Allows relationships to have relationships.



# Aggregation vs. Ternary




Rec'd: How information  
is changing?

# Conceptual Design Using the ER Model

- ER modeling can get tricky!
- Design choices:
  - **Entity** or **attribute**?
  - **Entity** or **relationship**?
  - Relationships: **Binary** or **ternary**? **Aggregation**?
- ER Model goals and limitations:
  - Lots of semantics can (and should) be captured.
  - Some constraints cannot be captured in ER.
    - We'll refine things in our logical (relational) design

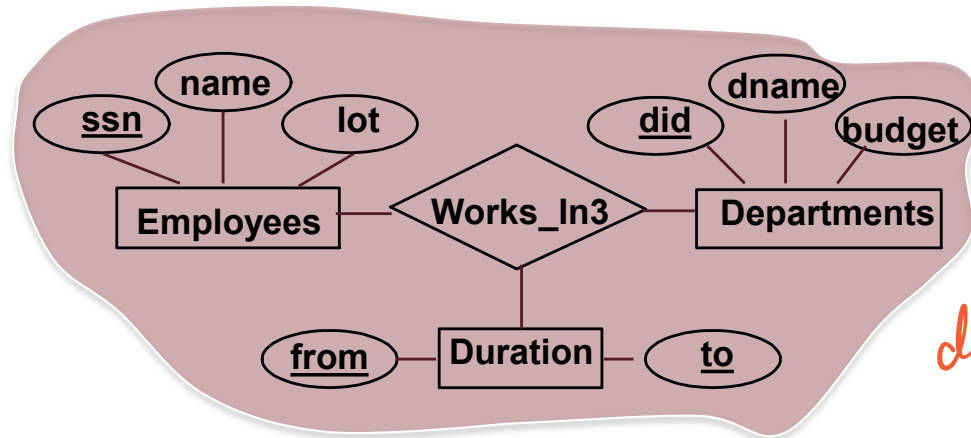
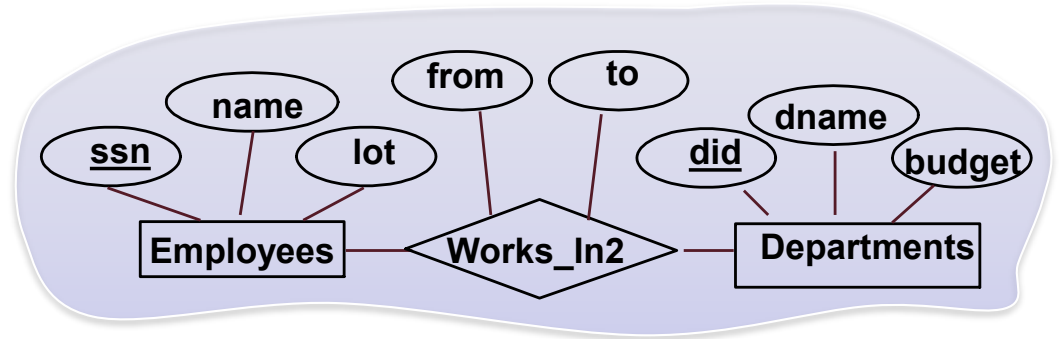
# Entity vs. Attribute

- “Address”:
  - attribute of Employees?
  - Entity of its own?
- It depends! Semantics and usage. 
  - Several addresses per employee?
    - must be an entity
    - atomic attribute types (no set-valued attributes!)
  - Care about structure? (city, street, etc.)
    - must be an entity!
    - atomic attribute types (no tuple-valued attributes!)



# Entity vs. Attribute (Cont.)

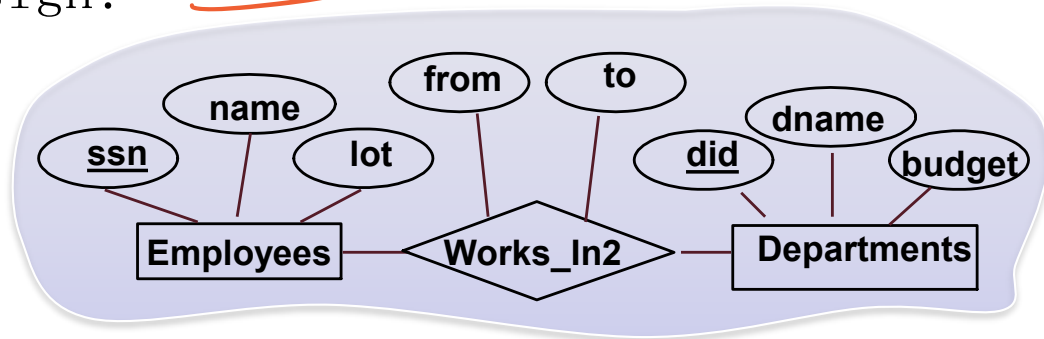
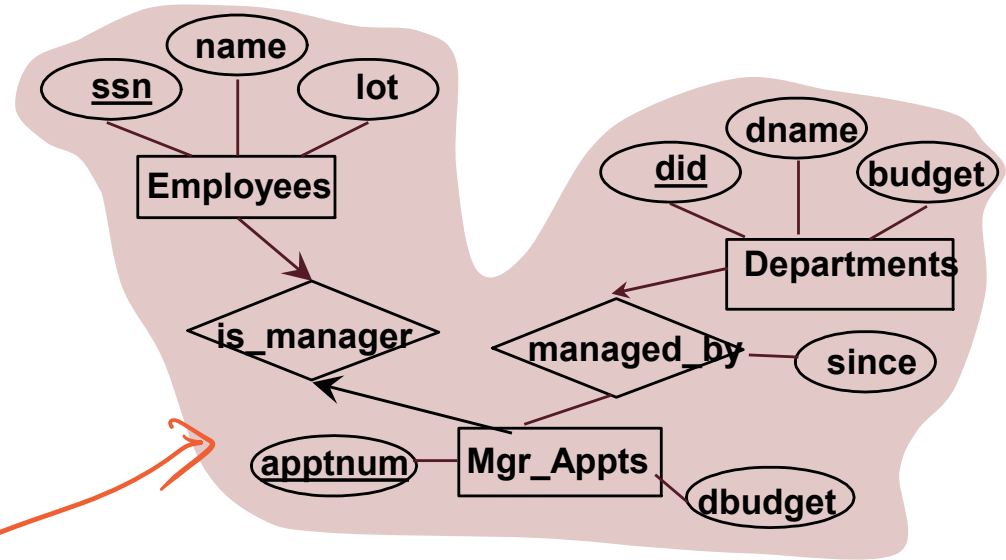
- Works\_In2: employee cannot work **in a** department for >1 period.
- Like multiple addresses per employee!



*different durations.*

# Entity vs. Relationship

- Separate discretionary budget (dbudget) for each dept.
- What if manager's dbudget covers all managed depts
  - Could repeat value
  - But redundancy = problems
- Better design:





# E-R Diagram as Wallpaper

- Very common for them to be wall-sized



# Steps in Database Design, Part 4

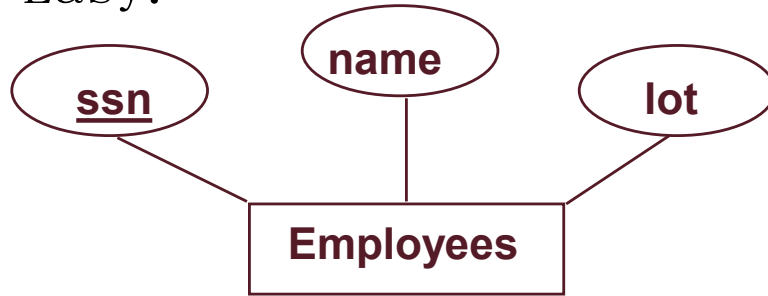
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  - *high level description (often done w/ER model)*  Completed
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- Logical Design *Schema*
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- Schema Refinement
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# Converting ER to Relational

- Fairly analogous structure
- But many simple concepts in ER are subtle to specify in relations

# Logical DB Design: ER to Relational

- Entity sets to tables. Easy.



ssn	name	lot
123-22-3666	Attishoo	48
231-31-5368	Smiley	22
131-24-3650	Smethurst	35

```
CREATE TABLE Employees
(ssn CHAR(11),
 name CHAR(20),
 lot INTEGER,
PRIMARY KEY (ssn))
```



# Relationship Sets to Tables

In translating a **many-to-many** relationship set to a relation, attributes of the relation must include:

- 1) Keys for each **participating entity set** (as **foreign keys**).

This set of attributes forms a **superkey** for the relation.

- 2) All descriptive ~~attributes~~

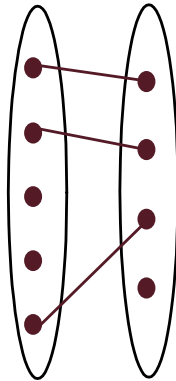
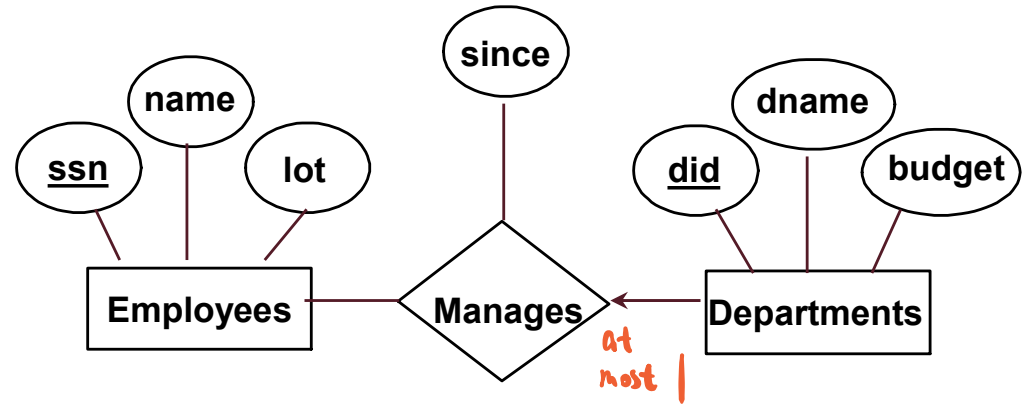
→ of the relation itself

ssn	did	since
123-22-3666	51	1/1/91
123-22-3666	56	3/3/93
231-31-5368	51	2/2/92

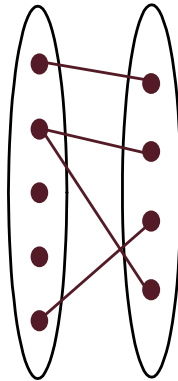
```
CREATE TABLE Works_In(  
    ssn CHAR(1),  
    did INTEGER,  
    since DATE,  
    PRIMARY KEY (ssn, did),  
    FOREIGN KEY (ssn)  
        REFERENCES Employees,  
    FOREIGN KEY (did)  
        REFERENCES Departments)
```

# Review: Key Constraints

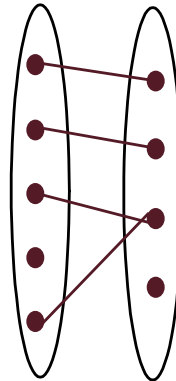
Each dept has at most one manager, according to the **key constraint** on Manages.



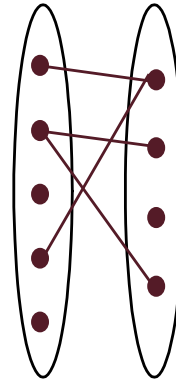
1-to-1



1-to Many



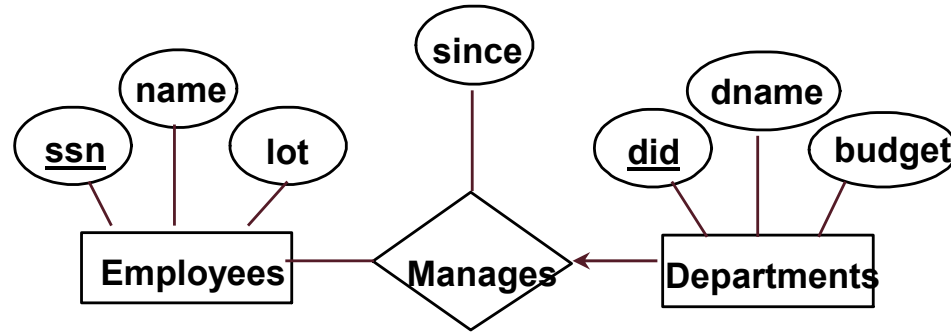
Many-to-1



Many-to-Many

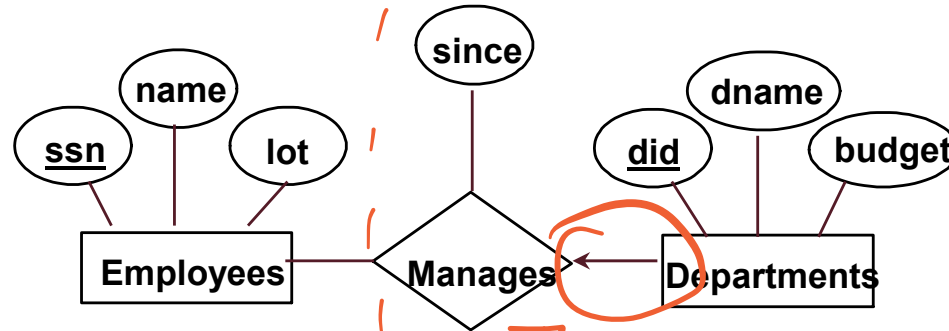


# Translating ER with Key Constraints



```
CREATE TABLE Manages(  
  ssn CHAR(11),  
  did INTEGER,  
  since DATE,  
  PRIMARY KEY (did),  
  FOREIGN KEY (ssn)  
    REFERENCES Employees,  
  FOREIGN KEY (did)  
    REFERENCES Departments)
```

# Translating ER with Key Constraints, cont



Since each department has a unique manager, we could instead combine Manages and Departments.

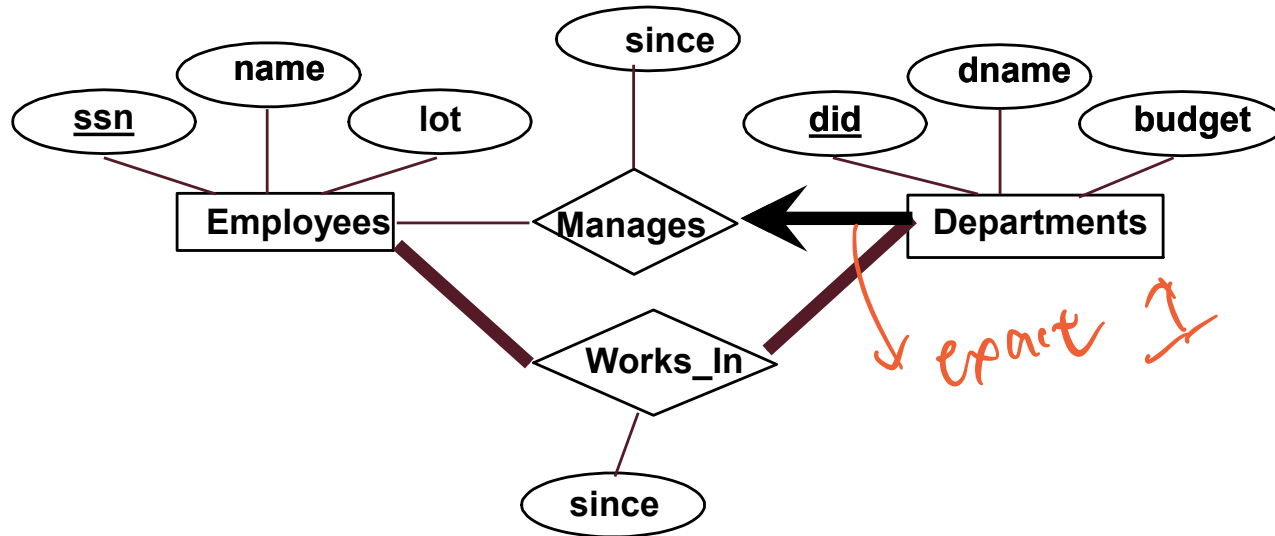
```
CREATE TABLE Manages(
  ssn CHAR(11),
  did INTEGER,
  since DATE,
  PRIMARY KEY (did),
  FOREIGN KEY (ssn)
    REFERENCES Employees,
  FOREIGN KEY (did)
    REFERENCES Departments)
```

Vs.

```
CREATE TABLE Dept_Mgr(
  did INTEGER,
  dname CHAR(20),
  budget REAL,
  ssn CHAR(11),
  since DATE,
  PRIMARY KEY (did),
  FOREIGN KEY (ssn)
    REFERENCES Employees)
```

# Review: Key+Participation Constraints

- Every department has one manager.
  - Every did value in Departments table must appear in a row of the Manages table (with a non-null ssn value!)



# Participation Constraints in SQL

- We can capture participation constraints involving one entity set in a binary relationship, but little else (without resorting to CHECK constraints which we'll learn later).

```
CREATE TABLE Dept_Mgr(  
  did INTEGER,  
  dname CHAR(20),  
  budget REAL,  
  ssn CHAR(11) NOT NULL, -- total participation!  
  since DATE,  
  PRIMARY KEY (did),  
  FOREIGN KEY (ssn) REFERENCES Employees  
  ON DELETE NO ACTION)
```

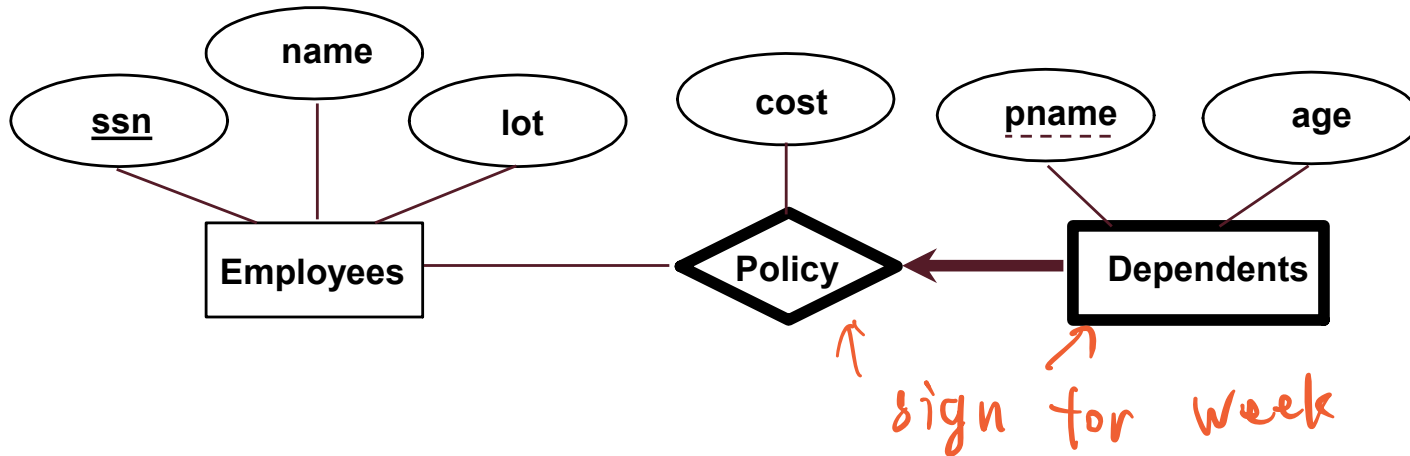
*exact 1*

*above deletion*



# Review: Weak Entities

- A **weak entity** can be identified uniquely only by considering the primary key of another (owner) entity.
  - Owner entity set and weak entity set **must participate in a one-to-many relationship set** (1 owner, many weak entities).
  - Weak entity set must have total participation in this **identifying** relationship set.



# Translating Weak Entity Sets

- Weak entity set and identifying relationship set are translated into a single table.
  - When the owner entity is deleted, all owned weak entities must also be deleted.

```
CREATE TABLE Dep_Policy (  
  pname CHAR(20),  
  age INTEGER,  
  cost REAL,  
  ssn CHAR(11) NOT NULL,  
  PRIMARY KEY (pname, ssn),  
  FOREIGN KEY (ssn) REFERENCES Employees  
  ON DELETE CASCADE)
```

on delete no action:

if employee has dependent  
can not remove it from  
e table.

delete thing be referred to,  
must delete things refer to it

# Summary of Conceptual Design

- **Conceptual design** follows requirements analysis
  - Yields a high-level description of data to be stored
- ER model popular for conceptual design
  - Constructs are expressive, close to the way we think about applications.
  - Note: There are many variations on ER model
    - Both graphically and conceptually
- Basic constructs: **entities**, **relationships**, and **attributes** (of entities and relationships).
- Some additional constructs: **weak entities**, ISA hierarchies (see text if you're curious), and aggregation.

# Summary of ER (Cont.)

- Basic integrity constraints
  - **key constraints**
  - **participation constraints**
- Some **foreign key** constraints are also implicit in the definition of a relationship set.
- Many other constraints (notably, **functional dependencies**) cannot be expressed.
- Constraints play an important role in determining the best database design for an enterprise.



# Summary of ER (Cont...)

- ER design is **subjective**. Many ways to model a given scenario!
- Analyzing alternatives can be tricky! Common choices include:
  - Entity vs. attribute, entity vs. relationship, binary or n-ary relationship, whether or not to use aggregation
- For good DB design: resulting relational schema should be analyzed and refined further.
  - Functional Dependency information  
+ normalization coming in subsequent lecture.