# Integer Discrete Flows and Lossless Compression

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- Map every input to unique output such that probable inputs map to shorter codes and improbable inputs are mapped to longer codes.
- Minimum code length for a symbol x is close to  $-\log \mathcal{D}(x)$

Minimum expected code length:

$$\mathbb{E}_{x \sim \mathcal{D}}\left[|c(x)|\right] \approx \mathbb{E}_{x \sim \mathcal{D}}\left[-\log p_X(x)\right] \geq \mathcal{H}(\mathcal{D})$$

#### Normalizing flows for integer-valued data

Problem formulation: Define invertible  $f_{ heta}: \mathbb{Z}^d \mapsto \mathbb{Z}^d$ 

Integer Discrete Flows (IDFs): Remove scaling in coupling layers (step 1).



$$z = \begin{bmatrix} z_1 \\ z_2 \end{bmatrix} \leftarrow \begin{bmatrix} x_1 \\ s^{\theta}(x_1) \odot x_2 + t^{\theta}(x_1) \end{bmatrix}$$



$$x = egin{bmatrix} x_1 \ x_2 \end{bmatrix} \leftarrow egin{bmatrix} z_1 \ (z_2 - t^{ heta}(z_1)) \oslash s^{ heta}(z_1) \end{bmatrix}$$



#### Normalizing flows for integer-valued data

Problem formulation: Define invertible  $f_{ heta}: \mathbb{Z}^d \mapsto \mathbb{Z}^d$ 

Integer Discrete Flows (IDFs): Constrain translations to be integer valued (step 2).



$$z = egin{bmatrix} z_1 \ z_2 \end{bmatrix} \leftarrow egin{bmatrix} x_1 \ x_2 + egin{bmatrix} t^{ heta}(x_1) \end{bmatrix}$$



Use straight-through estimator to backprop gradients



$$x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \leftarrow \begin{bmatrix} z_1 \\ z_2 - \lfloor t^{\theta}(z_1) \rfloor \end{bmatrix}$$



# Obtaining the density

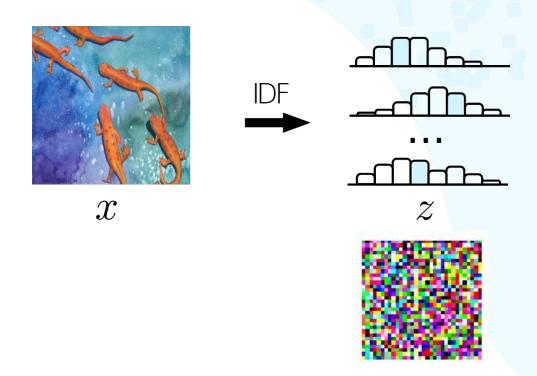
Continuous random variables:

$$p(x) = \int p(x|z)p(z) dz = \int \delta(x - f(z))p(z) dz = p(f^{-1}(x)) \left| \frac{\partial x}{\partial z} \right|^{-1}$$

Discrete random variables:

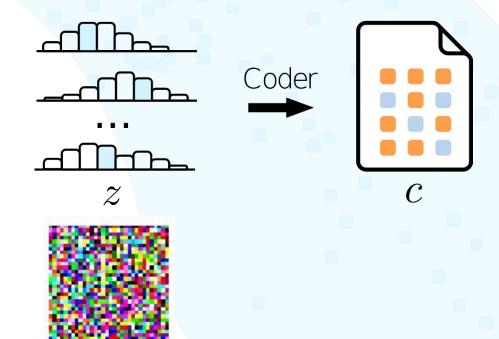
$$p(x) = \sum_{z} p(x|z)p(z) = \sum_{z} \delta_{z,f^{-1}(x)}p(z) = p(f^{-1}(x))$$

Step 1: transform input data to z-space using the IDF.

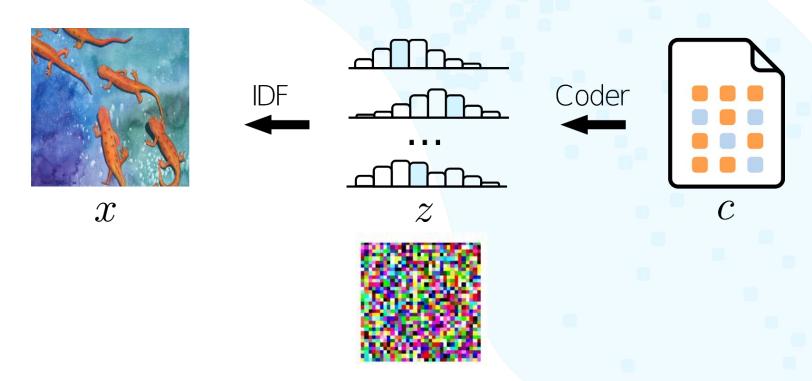


Step 2: encode *z* using off-the-shelve entropy encoder

 $\label{eq:code} \mbox{High-probability } z \rightarrow \mbox{Short code}$   $\mbox{Low-probability } z \rightarrow \mbox{Long code}$ 



Decompression works analogously in inverse order, using the inverse transformation: the entropy based decoder following by the inverse mapping defined by the IDF.



#### **Results**

Table 1: Compression performance of IDFs on CIFAR10, ImageNet32 and ImageNet64 in bits per dimension, and compression rate (shown in parentheses). The Bit-Swap results are retrieved from [23]. The column marked IDF<sup>†</sup> denotes an IDF trained on ImageNet32 and evaluated on the other datasets.

Dataset	IDF	$\mathrm{IDF}^\dagger$	Bit-Swap	FLIF [34]	PNG	JPEG2000
CIFAR10	3.34 (2.40×)	3.60 (2.22×)	$3.82(2.09\times)$	4.37 (1.83×)	5.89 (1.36×)	5.20 (1.54×)
ImageNet32	$4.18 (1.91 \times)$	$4.18 (1.91 \times)$	$4.50 (1.78 \times)$	$5.09(1.57\times)$	$6.42 (1.25 \times)$	$6.48 (1.23 \times)$
ImageNet64	$3.90 (2.05 \times)$	$3.94 (2.03 \times)$	_	$4.55 (1.76 \times)$	$5.74 (1.39 \times)$	$5.10 (1.56 \times)$

Table 3: Generative modeling performance of IDFs and comparable flow-based methods in bits per dimension (negative log<sub>2</sub>-likelihood).

Dataset	IDF	Continuous	RealNVP	Glow	Flow++
CIFAR10	3.32	3.31	3.49	3.35	3.08
ImageNet32	4.16	4.13	4.28	4.09	3.86
ImageNet64	3.90	3.85	3.98	3.81	3.69

#### Medical data: Histology dataset

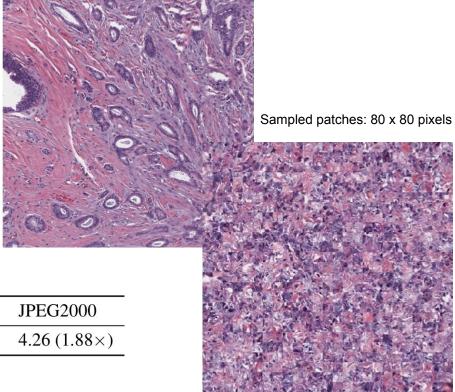
IDF trained on 80 x 80 px patches

Compression is done patch-wise (each patch is considered independent)

 Dataset
 IDF
 JP2-WSI
 FLIF [34]
 JPEG2000

 Histology
 2.42 (3.19×)
 3.04 (2.63×)
 4.00 (2.00×)
 4.26 (1.88×)

Resolution: 2000 x 2000 pixels



## Progressive image rendering

To partially render an image using IDFs, first the received variables are decoded. Next, using the hierarchical structure of the prior and ancestral sampling, the remaining dimensions are obtained. Below, the decoded images using 15, 30, 60, and 100% of the encoded data was user to decode the images.

