

# Heuristic functions: admissibility, consistency, dominance

Albert Sanchis
Alfons Juan
Jorge Civera

DSIIC

Departamento de Sistemas Informáticos y Computación

### Learning objectives

- Describe the concept of heuristic function.
- Propose admissible (lower bound) heuristic functions as a relaxation of the original problem
- Evaluate the admissibility and consistency of a heuristic function
- Compare heuristic functions in terms of dominance



## **Contents**

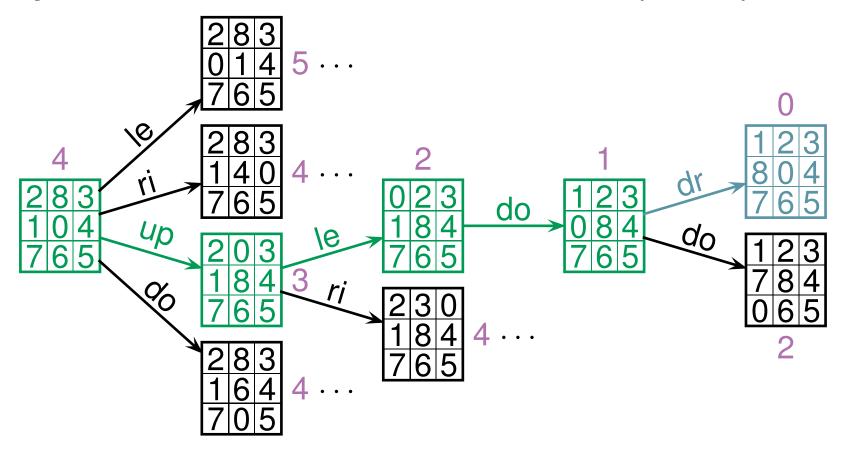
1	Concept of heuristic function	3
2	Admissibility	4
3	Consistency or monotonicity	5
4	Admissibility and consistency	6
5	Dominance	7
6	Conclusions	8



### 1 Concept of heuristic function

Provided a search problem on a state graph, a *heuristic function* h is a function that (efficiently) approximates the minimum cost  $h^*$  of reaching the goal state from a state n:

**Example:** sum of Manhattan distances in the 8-puzzle problem





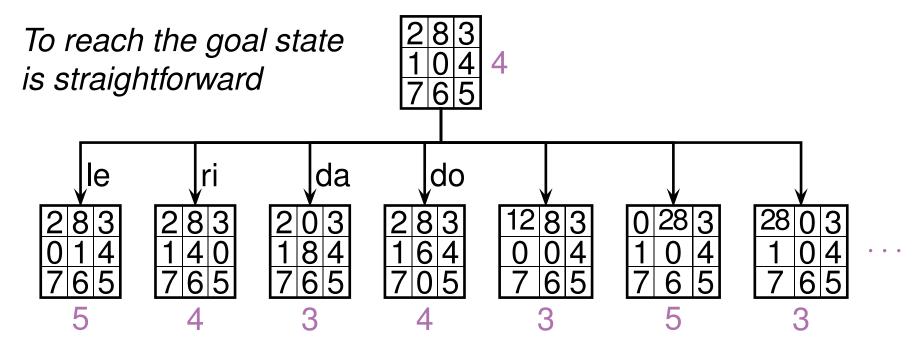
### 2 Admissibility

h is admissible (lower bound) if  $h(n) \le h^*(n) \forall$  node n.

It is usually obtained as a *relaxation of constraints* in the original problem, that is, smoothing or eliminating constraints to ease reaching the goal state.

Example: sum of Manhattan distances in the 8-puzzle problem

A can be moved to B if: B is adjacent to A and B is a blank





### 3 Consistency or monotonicity

h is *consistent* if,  $\forall$  node n [1, pp82–83]:

$$h(n) \le k(n, n') + h(n') \quad \forall \text{ node } n'$$

where k(n, n') is the minimum cost to go from n to n'. Therefore,

$$g(n) + h(n) \le g(n) + k(n, n') + h(n') \ \forall \text{ node } n'$$
  
 $f(n) \le f(n') \ \forall \text{ node } n'$ 

Equivalently, h is **monotone** if,  $\forall n$  [1, pp82-83]:

$$h(n) \le c(n, n') + h(n') \quad \forall n' \text{ adjacent to } n.$$

Consistency and monotonicity are equivalent properties.

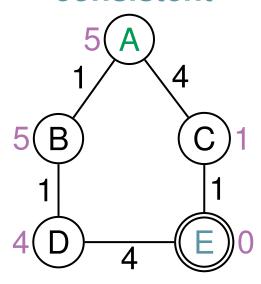


# 4 Admissibility and consistency

Consistency ⇒ Admissibility [1, pp83]

**Admissibility** ⇒ **Consistency** is not always true:

# Admissible and consistent



$$h(A) \leq w(A, B) + h(B)$$

$$h(A) \leq w(A, C) + h(C)$$

$$h(B) \leq w(B, A) + h(A)$$

$$h(B) \leq w(B, D) + h(D)$$

$$h(C) \leq w(C, A) + h(A)$$

$$h(C) \leq w(C, E) + h(E)$$

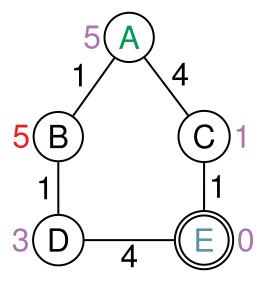
$$h(D) \leq w(D, B) + h(B)$$

$$h(D) \leq w(D, E) + h(E)$$

$$h(E) \leq w(E, C) + h(C)$$

$$h(E) \leq w(E, D) + h(D)$$

# Admissible and not consistent



$$h(B) \not\leq w(B,D) + h(D)$$



### 5 Dominance

**Dominance:**  $h_2(n)$  dominates  $h_1(n)$  if:

$$h_1(n) \le h_2(n) \le h^*(n) \ \forall n$$

**Example:** Manhattan dominates Misplaced tiles in 8-puzzle

**28**3 *Misplaced tiles:* 
$$1+1+1=3$$
 *7*65 *Manhattan:*  $1+1+2=4$ 

- ▶ A\* search with  $h_2(n)$  expands fewer nodes than  $h_1(n)$  [3].
- ▶ The closer h(n) approximates  $h^*(n)$ , fewer nodes are expanded

Non-admissible heuristics cannot be compared with admissible heuristics in terms of number of expanded nodes



### 6 Conclusions

#### We have studied:

- ► The concept of heuristic function
- Admissible heuristic functions as a relaxation of constraints
- Admissibility and consistency of a heuristic function
- Compare heuristic functions in terms of dominance



### References

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