

Question 1 (1.1 points)

Given the following main code of a program, where we assume that n is a constant defined with an integer value:

```
float a;  
float A[n][n], X[n][n], Y[n][n], Z[n][n];  
T1( A, X, Y, Z );  
a = T2( A );  
T3( a, X );  
T4( a, Y );  
T5( A, A, Z );  
T6( X, Y, Z );
```

and given the code of the following functions:

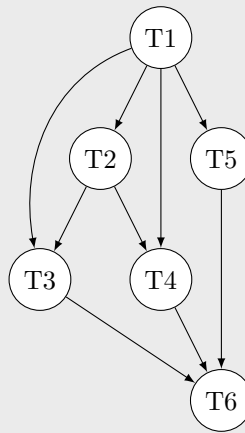
```
float T2( float A[n][n] ) {  
    float a = 0.0;  
    for( int i=0; i<n; i++ )  
        for( int j=i; j<n; j++ )  
            for( int k=i; k<n; k++ )  
                a = a + A[i][k];  
    return a;  
}  
  
void T5( float X[n][n], float Y[n][n], float Z[n][n] ) {  
    float aux;  
    for( int i=0; i<n; i++ )  
        for( int j=0; j<n; j++ ) {  
            aux=0;  
            for ( int k=0; k<n; k++ )  
                aux += X[i][k]*Y[k][j];  
            Z[i][j]=aux;  
        }  
}
```

where we know that the function T1 modifies all its arguments and has a cost of n^3 flops, while the functions T3, T4 and T6 modify only their last argument and also have a cost of n^3 flops. The function T2 has a cost of $\frac{n^3}{3}$ and the function T5 has a cost of $2n^3$.

0.3 p.

(a) Draw the graph of data dependencies between tasks.

Solution:



0.6 p.

- (b) Implement a parallel version with OpenMP using a single parallel region.

Solution:

```

float a;
float A[n][n], X[n][n], Y[n][n], Z[n][n];
T1( A, X, Y, Z );
#pragma omp parallel
{
    #pragma omp sections
    {
        #pragma omp section
        a = T2( A );
        #pragma omp section
        T5( A, A, Z );
    }
    #pragma omp sections
    {
        #pragma omp section
        T3( a, X );
        #pragma omp section
        T4( a, Y );
    }
} /* End of parallel */
T6( X, Y, Z );

```

0.2 p.

- (c) Get the speedup of the parallel version of the previous section for 2 processors.

Solution:

Sequential execution time:

$$t(n) = n^3 + \frac{n^3}{3} + n^3 + n^3 + 2n^3 + n^3 = \left(6 + \frac{1}{3}\right)n^3 = \frac{19}{3}n^3 \text{ flops}$$

Parallel execution time for $p = 2$:

$$t(n, p) = n^3 + \max\left(\frac{n^3}{3}, 2n^3\right) + \max(n^3, n^3) + n^3 = 5n^3 \text{ flops}$$

Speedup:

$$S = \frac{\frac{19}{3}n^3}{5n^3} = \frac{19}{15} = 1.27$$

Question 2 (1.2 points)

We want to manage a contest of `nP` participants, by a jury composed of `nM` members, numbered from position 0 onwards. Each jury casts two votes, one of 10 points for one participant and another of 5 for another. To do so, the function `manage_votes` receives the vector `votes_jury`, of `nM*2` components, with the two votes cast by each member of the jury and completes the vector `pts_participants`, of `nP` elements, with the score achieved by each participant. In addition to calculating which participant has won the contest (maximum score), the function completes the vectors `m10pts` and `m5pts` with the identifiers of the members of the jury who have awarded 10 points or 5 points, respectively, to the participant indicated as argument of the function. At the end, it obtains in the variable `nPZeroPts` the number of participants without any vote.

```
void manage_votes(int votes_jury[], int participant) {
    int i, p10pts, p5pts, pts_participants[nP];
    int m10pts[nM], m5pts[nM];
    int nM10pts=0, nM5pts=0, nPZeroPts=0;
    int pts_max=0, winner;
    ...
    for (i=0;i<nM;i++) {
        // Accumulate points to participants
        p10pts=votes_jury[i*2];
        p5pts=votes_jury[i*2+1];
        pts_participants[p10pts]+=10;
        pts_participants[p5pts]+=5;
        // Obtain the members that have voted the indicated participant
        if (p10pts==participant) {
            m10pts[nM10pts]=i;
            nM10pts++;
        }
        else if (p5pts==participant) {
            m5pts[nM5pts]=i;
            nM5pts++;
        }
    }
    // Calculate the winner and the number of participants with 0 points
    for (i=0;i<nP;i++) {
        if (pts_participants[i]>pts_max) {
            pts_max=pts_participants[i];
            winner=i;
        }
        else if (pts_participants[i]==0)
            nPZeroPts++;
    }
    ...
}
```

Parallelize the voting management in the most efficient way possible, using a single parallel region.

Solution:

```
void manage_votes(int votes_jury[], int participant) {
    int i, p10pts, p5pts, pts_participants[nP];
    int m10pts[nM], m5pts[nM];
    int nM10pts=0, nM5pts=0, nPZeroPts=0;
    int pts_max=0, winner;
    ...
    #pragma omp parallel
    {
        #pragma omp for private(p10pts, p5pts)
```

```

for (i=0;i<nM;i++) {
    // Accumulate points to participants
    p10pts=votes_jury[i*2];
    p5pts=votes_jury[i*2+1];
    #pragma omp atomic
    pts_participants[p10pts]+=10;
    #pragma omp atomic
    pts_participants[p5pts]+=5;
    // Obtain the members that have voted the indicated participant
    if (p10pts==participant) {
        #pragma omp critical (pts10)
        {
            m10pts[nM10pts]=i;
            nM10pts++;
        }
    }
    else if (p5pts==participant) {
        #pragma omp critical (pts5)
        {
            m5pts[nM5pts]=i;
            nM5pts++;
        }
    }
}
// Calculate the winner and the number of participants with 0 points
#pragma omp for reduction(+:nPZeroPts)
for (i=0;i<nP;i++) {
    if (pts_participants[i]>pts_max) {
        #pragma omp critical
        if (pts_participants[i]>pts_max) {
            pts_max=pts_participants[i];
            winner=i;
        }
    }
    else if (pts_participants[i]==0)
        nPZeroPts++;
}
}
...
}

```

Question 3 (1.2 points)

Given the following function:

```

void normalize_mat(double A[N][N]){
    int i, j;
    double s, norm1=0, norm2=0;
    for (i = 0; i < N; i++){
        s = 0;
        for (j=0; j<N; j++){
            norm1+=A[i][j];
            s+= fabs(A[i][j]);
        }
        if (s>norm2)
            norm2= s;
    }
    norm1=sqrt(norm1)*norm2;
}

```

```

    for (i = 0; i < N; i++)
        for (j = 0; j < N; j++)
            A[i][j]*=norm1;
}

```

0.9 p.

- (a) Parallelize the previous function with OpenMP, using a single parallel region.

Solution:

```

#pragma omp parallel
{
    #pragma omp for private(j,s) reduction(+:norm1) reduction(max:norm2)
    for (i = 0; i < N; i++){
        ...
    }
    #pragma omp single
    norm1=sqrt(norm1)*norm2;
    #pragma omp for private(j)
    for (i = 0; i < N; i++)
        for (j = 0; j < N; j++)
            A[i][j]*=norm1;
}

```

0.2 p.

- (b) Calculate the sequential and parallel cost, assuming that N is divisible by the number of threads p and that the cost of the functions `fabs` and `sqrt` is 1 flop. Indicate the intermediate computations done to reach the solution.

Solution: Sequential cost:

$$t(N) = \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} 3 + 2 + \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} 1 = 3N^2 + 2 + N^2 \approx 4N^2 \text{ flops.}$$

Parallel cost:

$$t(N, p) = \sum_{i=0}^{N/p-1} \sum_{j=0}^{N-1} 3 + 2 + \sum_{i=0}^{N/p-1} \sum_{j=0}^{N-1} 1 = \frac{3N^2}{p} + 2 + \frac{N^2}{p} \approx \frac{4N^2}{p} \text{ flops.}$$

0.1 p.

- (c) Calculate the speedup and the efficiency.

Solution:

$$S(n, p) = \frac{4N^2}{\frac{4N^2}{p}} = p.$$

$$E(n, p) = \frac{S(n, p)}{p} = 1.$$