

# University of Trieste Data Management for Big Data Course

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# Data Warehouse case study

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# Contents

1	Intr	roduction	2
	1.1	TPC-H benchmark database	2
		1.1.1 Database statistics	2
		1.1.2 Database SQL definition	3
	1.2	Organization of work	
2	Set	of queries	3
	2.1	Export/import revenue value	3
	2.2	Late delivery	5
	2.3	Returned item loss	
	2.4	Execution times	7
3	Mat	terialization	7
	3.1	Execution times	G
4	Ind	exes design	9
5	Cor	nclusions	g

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### 1 Introduction

The aim of this project is to study an efficient implementation of a suite of business oriented ad-hoc queries over the public TPC-H benchmark, which can be considered as a Big Data database, that has been implemented in Postgres.

#### 1.1 TPC-H benchmark database

The TPC-H benchmark is a decision support benchmark that can be downloaded from the TPC official website. The data generator lets the user specify a *scale factor* in order to control the size of the resulted database. Our choices was to use a *scale factor* of 10, meaning that the overall database size is approximately 13 GB.

#### 1.1.1 Database statistics

The benchmark is composed by eight tables:

- CUSTOMER, with 16 columns and 1500000 tuples (312 MB);
- LINEITEM, with 32 columns and 59 986 052 tuples (11 GB); the main attributes that are going to be used are:
  - 1\_extendedprice (1 351 462 distinct values, i.e. there is an average of 44 tuples with the same value, that range from 900.91 to 104 949.50),
  - 1\_discount (11 distinct values, i.e. there is an average of 5 453 277 tuples with the same value, that range from 0.00 to 0.10),
  - o l\_returnflag (which can assume values A→accepted, R→returned, N→not yet delivered; the percentage of tuples for A and R are almost 25%, while the percentage of tuples where l\_returnflag is N is about 50%),
  - 1\_commitdate (2466 distinct values, i.e. there is an average of 24325 tuples with the same value, that range from 1992-01-31 to 1998-10-31),
  - 1\_receiptdate (2555 distinct values, i.e. there is an average of 23478 tuples with the same value, that range from 1992-01-03 to 1998-12-31);
- NATION, with 8 columns and 25 tuples (24 kB);
- ORDERS, with 18 columns and 1500000 tuples (2481 kB); the main attributes that are going to be used are:
  - o o\_orderdate (2406 distinct values, i.e. there is an average of 6234 tuples with the same value, that range from 1992-01-01 to 1998-08-02);
- PART, with 18 columns and 2000000 tuples (363 MB); the main attributes that are going to be used are:
  - p\_type (150 distinct values, i.e. there is an average of 13 333 tuples with the same value);

- PARTSUPP, 10 columns and with 8 000 000 tuples (1535 MB);
- REGION, 6 columns and with 5 tuples (24 kB);
- SUPPLIER, 14 columns and with 100 000 tuples (20 MB).

Other attributes have been used, but statistics about them have been omitted for lack of usefulness (e.g., keys of the tables, for which the cardinality is exactly the cardinality of the corresponding table).

#### 1.1.2 Database SQL definition

The SQL definition of the tables can be found on the official benchmark download.

#### 1.2 Organization of work

Each of the group component implemented a query (specifically, Della Rovere designed the first query, Capone on the second one, and Stefanel the last one). After regular meetings between the group members comparing the solo works, common substructures between queries have been detected, and the following work proceed in a coral way.

All the execution times reported in the present report are the results of runs on the same machine, with roughly the same external factors. In particular, a computer with an Apple M1 processor, 8 GB of RAM, and macOS operating system has been used.

# 2 Set of queries

## 2.1 Export/import revenue value

It is asked to return the *export/import revenue* between two different nations (E, I) where E is the nation of the lineitem supplier and I the nation of the lineitem customer, and where the *revenue* is defined as

```
SUM(l_extendedprice*(1-l_discount))
```

. The aggregation should be performed with the Month  $\to$  Quarter  $\to$  Year, (Part) Type and Nation  $\to$  Region roll-ups.

```
9
       FROM lineitem JOIN orders ON (l_orderkey = o_orderkey)
10
   ), customer_location AS (
11
       SELECT
12
           c_custkey,
13
           c_name,
14
           n_nationkey AS c_nationkey,
15
           n_name AS c_nationname,
16
           r_regionkey AS c_regionkey,
17
           r_name AS c_regionname
       FROM customer
18
19
           JOIN nation ON (c_nationkey = n_nationkey)
           JOIN region ON (n_regionkey = r_regionkey)
20
   ), supplier_location AS (
21
22
       SELECT
23
           s_suppkey,
24
           s_name,
25
           n_nationkey AS s_nationkey,
26
           n_name AS s_nationname,
27
           r_regionkey AS s_regionkey,
28
           r_name AS s_regionname
29
       FROM supplier
30
           JOIN nation ON (s_nationkey = n_nationkey)
31
           JOIN region ON (n_regionkey = r_regionkey)
32
   ), query1 AS (
       SELECT
33
34
           EXTRACT (YEAR FROM o_orderdate) AS _year,
35
           EXTRACT (QUARTER FROM o_orderdate) AS _quarter,
36
           EXTRACT (MONTH FROM o_orderdate) AS _month,
37
           c_regionname,
38
           c_nationname,
39
           c_name,
40
           s_regionname,
41
           s_nationname,
42
           s_name,
43
           p_type,
           SUM(l_extendedprice * (1 - l_discount)) AS revenue
44
45
       FROM lineitem_orders
46
           JOIN part ON l_partkey = p_partkey
           JOIN supplier_location ON (s_suppkey = l_suppkey)
47
48
           JOIN customer_location ON (c_custkey = o_custkey)
       WHERE s_nationkey <> c_nationkey
49
       GROUP BY
50
51
           _year,
```

```
52
            _quarter,
53
            _month,
54
            c_regionkey,
55
            c_regionname,
56
            c_nationkey,
57
            c_nationname,
58
            c_custkey,
59
            c_name,
60
            s_regionkey,
            s_regionname,
61
62
            s_nationkey,
            s_nationname,
63
64
            s_suppkey,
65
            s_name,
66
           p_type
67
68
   SELECT * FROM query1;
```

# 2.2 Late delivery

It is asked to retrieve the number of orders where at least one "lineitem" has been received later than the committed date. The aggregation should be performed with the Month  $\rightarrow$  Year roll-up, and the (Customer's) Nation  $\rightarrow$  Region roll-up.

```
WITH lineitem_orders AS (
2
       SELECT
3
           o_orderkey,
           1_partkey,
4
5
           1_suppkey,
6
           o_orderdate,
7
           o_custkey,
8
           l_commitdate,
9
           l_receiptdate
10
       FROM lineitem JOIN orders ON (l_orderkey = o_orderkey)
11
   ), customer_location AS (
12
       SELECT
13
           c_custkey,
14
           n_nationkey AS c_nationkey,
           n_name AS c_nationname,
15
16
           r_regionkey AS c_regionkey,
17
           r_name AS c_regionname
18
       FROM customer
19
           JOIN nation ON (c_nationkey = n_nationkey)
```

```
20
           JOIN region ON (n_regionkey = r_regionkey)
   ), query2 AS (
21
22
   SELECT
       EXTRACT (YEAR FROM o_orderdate) AS _year,
23
24
       EXTRACT (MONTH FROM o_orderdate) AS _month,
25
       c_regionname,
26
       c_nationname,
       COUNT(DISTINCT(o_orderkey)) AS orders_no
27
28
   FROM lineitem_orders
29
       JOIN part ON l_partkey = p_partkey
30
       JOIN customer_location ON (c_custkey = o_custkey)
31
   WHERE
32
       l_receiptdate > l_commitdate
       -- AND _month = 1
33
34
       -- AND p_type = 'PROMO BURNISHED COPPER'
   GROUP BY
35
       _year,
36
37
       _month,
38
       c_regionkey,
39
       c_regionname,
       c_nationkey,
40
41
       c_nationname
42 )
43 SELECT * FROM query2;
```

#### 2.3 Returned item loss

It is asked to retrieve the *revenue loss* for customers who might be having problems with the parts that are shipped to them, where a *revenue loss* is defined as

```
SUM(l_extendedprice*(1-l_discount))
```

for all qualifying lineitems.

```
WITH lineitem_orders AS (
1
2
       SELECT
3
           o_orderkey,
4
           o_orderdate,
           o_custkey,
5
           l_extendedprice,
6
7
           l_discount,
           1_returnflag
8
9
       FROM lineitem JOIN orders ON (l_orderkey=o_orderkey)
10),
```

```
query3 AS (
   SELECT
12
13
       EXTRACT(YEAR FROM o_orderdate) AS _year,
       EXTRACT(QUARTER FROM o_orderdate) AS _quarter,
14
15
       EXTRACT(MONTH FROM o_orderdate) AS _month,
16
       SUM(l_extendedprice*(1-l_discount)) AS returnloss
17
18
   FROM
19
       lineitem_orders
       JOIN customer ON (o_custkey=c_custkey)
20
21
   WHERE
22
       l_returnflag='R'
23
       -- AND c_name='Customer#000129976'
       -- AND EXTRACT(QUARTER FROM o_orderdate) = 1
24
   GROUP BY
25
26
       _year,
27
       _quarter,
28
       _month,
29
       c_custkey,
30
       c_name
31
   )
32
   SELECT * FROM query3;
```

#### 2.4 Execution times

The query timings have been measured as previously discussed in subsection 1.2. Furthermore, no consecutive runs have been performed on the same query, in order to reduce external biases (following a Round-Robin schema).

Query	Run 1	Run 2	Run 3	Run 4	Run 5	$\mu$	$\sigma$
1	643	685	724	698	813	713	63
2	201	203	205	203	203	203	2
3	98	102	101	102	101	101	2

Table 1: Naïve query timings, in seconds.

### 3 Materialization

After a study on the given set of queries, some common intermediate results have been detected between the three. In order to try lowering the average execution cost, materialized views have been defined starting from the said results.

```
-- used by Q1, Q2, Q3
1
2
   CREATE MATERIALIZED VIEW lineitem_orders_mv AS
       SELECT
3
4
           o_orderkey,
5
           1_partkey,
6
           1_suppkey,
7
           o_orderdate,
8
           o_custkey,
9
           l_extendedprice,
10
           l_discount,
11
           1_returnflag,
12
           l_commitdate,
13
           1_receiptdate
       FROM lineitem JOIN orders ON (l_orderkey = o_orderkey);
14
15
16
   -- used by Q1 and Q2
   CREATE MATERIALIZED VIEW customer_location_mv AS
17
18
       SELECT
19
           c_custkey,
20
           c_name,
           n_nationkey AS c_nationkey,
21
22
           n_name AS c_nationname,
23
           r_regionkey AS c_regionkey,
24
           r_name AS c_regionname
25
       FROM customer
26
           JOIN nation ON (c_nationkey = n_nationkey)
27
           JOIN region ON (n_regionkey = r_regionkey);
28
29
   -- used by Q1 and Q2
30
   CREATE MATERIALIZED VIEW supplier_location_mv AS
31
       SELECT
32
           s_suppkey,
33
           s_name,
34
           n_nationkey AS s_nationkey,
35
           n_name AS s_nationname,
36
           r_regionkey AS s_regionkey,
37
           r_name AS s_regionname
38
       FROM supplier
39
           JOIN nation ON (s_nationkey = n_nationkey)
40
           JOIN region ON (n_regionkey = r_regionkey);
```

#### 3.1 Execution times

As for the base versions of queries, we collected statistics on execution timings and results can be observed on Table 2. Further considerations are reported in the section 5.

	Query	Run 1	Run 2	Run 3	Run 4	Run 5	$\mu$	$\sigma$
	1	643	685	724	698	813	713	63
	2	201	203	205	203	203	203	2
Ī	3	98	102	101	102	101	101	2

Table 2: Query timings with materialized views, in seconds.

# 4 Indexes design

. . .

# 5 Conclusions

Siccome nelle viste materializzate ci sono più attributi rispetto alle tableele intermedie con WITH, nella prima query, allora quella query è più pesante con le viste materializzate.

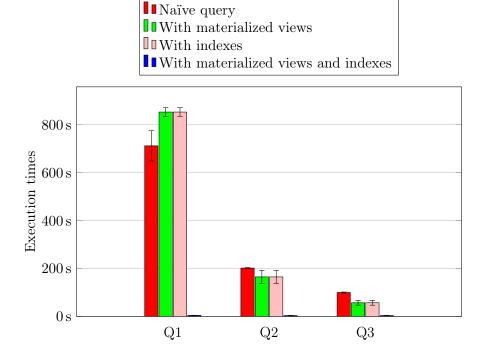


Figure 1: Query timings