

Programare declarativă

Monoid, Foldable

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Monoid

din nou **foldr**

```
foldr :: (a -> b -> b) -> b -> t a -> b
```

```
Prelude> foldr (+) 0 [1,2,3]
```

```
6
```

```
Prelude> foldr (*) 1 [1,2,3]
```

```
6
```

```
Prelude> foldr (++) [] ["1","2","3"]  
"123"
```

```
Prelude> foldr (||) False [True, False, True]  
True
```

```
Prelude> foldr (&&) True [True, False, True]  
False
```

Ce au in comun aceste operații?

Monoizi

(M, \circ, e) este **monoid** dacă

$\circ : M \times M \rightarrow M$ este asociativă

$m \circ e = e \circ m = m$ oricare $m \in M$

Monoizi

(M, \circ, e) este **monoid** dacă

- $\circ : M \times M \rightarrow M$ este asociativă
- $m \circ e = e \circ m = m$ oricare $m \in M$

Observații:

- $(\text{Int}, +, 0)$, $(\text{Int}, *, 1)$, $(\text{String}, ++, [])$, $(\{\text{True}, \text{False}\}, \&\&, \text{True})$ sunt monoizi

Monoizi

(M, \circ, e) este **monoid** dacă

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$m \circ e = e \circ m = m$ oricare $m \in M$

Observații:

- $(\text{Int}, +, 0)$, $(\text{Int}, *, 1)$, $(\text{String}, ++, [])$, $(\{\text{True}, \text{False}\}, \&\&, \text{True})$ sunt monoizi
- Operația de monoid poate fi generalizată pe liste:

sum = **foldr** (+) 0

product = **foldr** (*) 1

concat = **foldr** (++) []

all = **foldr** (&&) True

clasa Monoid

<https://en.wikibooks.org/wiki/Haskell/Monoids>//<https://hackage.haskell.org/package/base-4.10.0.0/docs/Data-Monoid.html>

Data.Monoid

```
class Monoid a where
    mempty  :: a                -- elementul neutru
    mappend :: a -> a -> a      -- operatia de monoid

    mconcat :: [a] -> a         -- generalizarea la liste
    mconcat = foldr mappend mempty
```

Observație: În loc de *mappend* se poate folosi (*<>*)

```
infixr 6 <>
(<>) :: Monoid m => m -> m -> m
(<>) = mappend          -- notatie infixă
```

clasa **Monoid**

Legile monoizilor

Instanțele clasei **Monoid** trebuie să satisfacă următoarele ecuații:

$$x \langle \rangle (y \langle \rangle z) == (x \langle \rangle y) \langle \rangle z$$

$$x \langle \rangle \text{mempty} == x$$

$$\text{mempty} \langle \rangle x == x$$

Atenție! Acest lucru este responsabilitatea programatorului!

clasa **Monoid**

Legile monoizilor

Instanțele clasei **Monoid** trebuie să satisfacă următoarele ecuații:

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$$x \langle \rangle \text{mempty} == x$$

$$\text{mempty} \langle \rangle x == x$$

Atenție! Acest lucru este responsabilitatea programatorului!

Listele ca instanță

```
instance Monoid [a] where
```

```
    mempty  = []
```

```
    mappend = (++)
```

```
Prelude> mempty :: [a]
```

```
[]
```

```
Prelude> mconcat [[1,2,3],[4,5],[6]]
```

```
[1,2,3,4,5,6]
```

clasa **Monoid**

$(\text{Int}, +, 0)$, $(\text{Int}, *, 1)$ sunt monoizi

$(\{\text{True}, \text{False}\}, \&\&, \text{True})$, $(\{\text{True}, \text{False}\}, ||, \text{False})$ sunt monoizi

Cum definim instante diferite pentru acelasi tip?

clasa **Monoid**

$(\text{Int}, +, 0)$, $(\text{Int}, *, 1)$ sunt monoizi

$(\{\text{True}, \text{False}\}, \&\&, \text{True})$, $(\{\text{True}, \text{False}\}, ||, \text{False})$ sunt monoizi

Cum definim instante diferite pentru acelasi tip?

- se crează o copie a tipului folosind **newtype**
- copia este definită ca instanță a tipului

newtype

newtype Nat = MkNat **Integer**

- **newtype** se folosește când un singur constructor este aplicat unui singur tip de date
- declarația cu **newtype** este mai eficientă decât cea cu **data**
- **type** redenumeste tipul; **newtype** face o copie și permite redefinirea operațiilor

clasa **Monoid**

- **Num a** ca monoid față de adunare

```
newtype Sum a = Sum { getSum :: a }
                  deriving (Eq, Read, Show)
```

```
instance Num a => Monoid (Sum a) where
    mempty = Sum 0
    Sum x 'mappend' Sum y = Sum (x + y)
```

- **Num a** ca monoid față de înmulțire

```
newtype Product a = Product { getProduct :: a }
                      deriving (Eq, Read, Show)
```

```
instance Num a => Monoid (Product a) where
    mempty = Product 1
    Product x 'mappend' Product y = Product (x * y)
```

clasa Monoid

Prelude> Sum 3

<interactive>:15:1: error:

Prelude> :m + Data.Monoid

Prelude Data.Monoid> Sum 3

Sum {getSum = 3}

Prelude Data.Monoid> Sum 3 <> Sum 4

Sum {getSum = 7}

Prelude Data.Monoid> Sum 3 + Sum 4

Sum {getSum = 7}

Prelude Data.Monoid> mconcat [Sum 3,Sum 4,Sum 5]

Sum {getSum = 12}

Prelude Data.Monoid> (getSum . mconcat) [Sum 3,Sum 4,Sum 5]
12

Prelude Data.Monoid> (getSum . mconcat) \$ **map** Sum [3,4,5]
12

Prelude Data.Monoid> getSum . mconcat . (**map** Sum) \$ [3,4,5]
12

Monoid Maybe

```
instance Monoid a => Monoid (Maybe a) where
  mempty                = Nothing
  Nothing 'mappend' m   = m
  m          'mappend' Nothing = m
  Just m1    'mappend' Just m2  = Just (m1 'mappend' m2)
```

Atentie! Monoid a => este o constrangere de tip

```
Prelude Data.Monoid> Nothing 'mappend' (Just 3)
<interactive>:35:1: error:
```

```
Prelude Data.Monoid> Nothing 'mappend' (Just (Sum 3))
Just (Sum {getSum = 3})
```

Funcții ca instanțe

(**a -> a**) ca instanța a clasei **Monoid**

```
newtype Endo a = Endo { appEndo :: a -> a }
```

```
instance Monoid (Endo a) where
```

```
    mempty                = Endo id
```

```
    Endo g 'mappend' Endo f = Endo (g . f)
```

Funcții ca instanțe

(**a** -> **a**) ca instanța a clasei **Monoid**

```
newtype Endo a = Endo { appEndo :: a -> a }
```

```
instance Monoid (Endo a) where
```

```
    mempty                = Endo id
```

```
    Endo g 'mappend' Endo f = Endo (g . f)
```

```
Prelude> :m + Data.Monoid
```

```
>let f = mconcat [Endo (+1), Endo (+2), Endo (+3)]
```

```
>:t f
```

```
f :: Num a => Endo a
```


Funcții ca instanțe

(**a** -> **a**) ca instanța a clasei **Monoid**

```
newtype Endo a = Endo { appEndo :: a -> a }
```

```
instance Monoid (Endo a) where
    mempty                = Endo id
    Endo g `mappend` Endo f = Endo (g . f)
```

```
Prelude> :m + Data.Monoid
>let f = mconcat [Endo (+1), Endo (+2), Endo (+3)]
>:t f
f :: Num a => Endo a

> (appEndo f) 0
6
> (appEndo . mconcat) [Endo (+1), Endo (+2), Endo (+3)] $ 0
6
```

Foldable

din nou **foldr**

foldr pe liste

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr f i [] = i
foldr f i (x:xs) = f x (foldr f i xs)
```

Problema: să generalizăm **foldr** la alte structuri recursive.

Exemplu: arbori binari

```
data BinaryTree a = Leaf a
                  | Node (BinaryTree a) (BinaryTree a)
deriving Show
```

din nou **foldr**

foldr pe liste

```
foldr :: (a -> b -> b) -> b -> [a] -> b
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Exemplu: arbori binari

```
data BinaryTree a = Leaf a
                  | Node (BinaryTree a) (BinaryTree a)
deriving Show
```

Cum definim **"foldr"** înlocuind listele cu date de tip **BinaryTree** ?

"foldr" folosind BinaryTree

```
data BinaryTree a = Leaf a
                  | Node (BinaryTree a) (BinaryTree a)
deriving Show
```

foldTree

```
foldTree :: (a -> b -> b) -> b -> BinaryTree a -> b
```

```
foldTree f i (Leaf x) = f x i
```

```
foldTree f i (Node l r) = foldTree f (foldTree f i r) l
```

foldTree

```
data BinaryTree a = Leaf a
                  | Node (BinaryTree a) (BinaryTree a)
deriving Show
```

```
foldTree :: (a -> b -> b) -> b -> BinaryTree a -> b
foldTree f i (Leaf x) = f x i
foldTree f i (Node l r) = foldTree f (foldTree f i r) l
```

```
myTree = Node (Node (Leaf 1)(Leaf 2))(Node (Leaf 3)(Leaf 4))
```

```
*Main> foldTree (+) 0 myTree
10
```

clasa **Foldable**

<https://en.wikibooks.org/wiki/Haskell/Foldable>

<https://hackage.haskell.org/package/base-4.10.0.0/docs/Data-Foldable.html>

Data.Foldable

```
class Foldable t where
    fold      :: Monoid m => t m -> m
    foldMap   :: Monoid m => (a -> m) -> t a -> m
    foldr     :: (a -> b -> b) -> b -> t a -> b

    fold = foldMap id
    ...
```

Observații:

- definiția minimală completă conține fie **foldMap**, fie **foldr**
- foldMap** și **foldr** pot fi definite una prin cealaltă
- pentru a crea o instanță este suficient să definim una dintre **foldMap** și **foldr**, cealaltă va fi automat accesibilă

Foldable cu foldr

```
instance Foldable BinaryTree where
  foldr = foldTree
```

```
tree1 = Node(Node(Leaf 1)(Leaf 2))(Node (Leaf 3)(Leaf 4))
treeS = Node (Node(Leaf "1")(Leaf "2"))
          (Node (Leaf "3")(Leaf "4"))
```

```
*Main> foldr (+) 0 tree1
10
```

```
*Main> foldr (++) [] treeS
"1234"
```


clasa **Foldable**

Data.Foldable

```
class Foldable t where
    fold      :: Monoid m => t m -> m
    foldMap   :: Monoid m => (a -> m) -> t a -> m
    foldr     :: (a -> b -> b) -> b -> t a -> b

    fold = foldMap id
    ...
```

```
instance Foldable BinaryTree where
    foldr = foldTree
```

Observație: în definiția clasei **Foldable**, variabila de tip **t** nu reprezintă un tip concret (`[a]`, `Sum a`) ci un **constructor de tip** (`BinaryTree`)

Foldable cu foldr

```
instance Foldable BinaryTree where
    foldr = foldTree
```

```
tree1 = Node(Node(Leaf 1)(Leaf 2))(Node (Leaf 3)(Leaf 4))
treeS = Node (Node(Leaf "1")(Leaf "2"))
           (Node (Leaf "3")(Leaf "4"))
```

Avem definite automat **foldMap** și alte funcții precum: **foldl**, **foldr'**, **foldr1**,...

```
*Main> foldl (++) [] treeS
"1234"
*Main> foldl (+) 0 tree1
10
```

Foldable cu foldr

```
instance Foldable BinaryTree where
  foldr = foldTree
```

```
tree1 = Node(Node(Leaf 1)(Leaf 2))(Node (Leaf 3)(Leaf 4))
treeS = Node (Node(Leaf "1")(Leaf "2"))
          (Node (Leaf "3")(Leaf "4"))
```

Avem definite automat **foldMap** și alte funcții precum: **foldl**, **foldr'**, **foldr1**,...

```
*Main> foldl (++) [] treeS
"1234"
*Main> foldl (+) 0 tree1
10
```

```
*Main Data.Monoid> foldMap Sum tree1
Sum {getSum = 10}
*Main Data.Monoid> foldMap id treeS
"1234"
```

foldMap

```
foldMap :: Monoid m => (a -> m) -> t a -> m
```

```
newtype Sum a = Sum { getSum :: a }
                  deriving (Eq, Read, Show)
```

```
instance Num a => Monoid (Sum a) where
    mempty = Sum 0
    Sum x 'mappend' Sum y = Sum (x + y)
```

```
tree1 = Node(Node(Leaf 1)(Leaf 2))(Node (Leaf 3)(Leaf 4))
```

```
*Main> foldMap Sum tree1    -- Sum :: a -> Sum a
Sum {getSum = 10}
```

Cum definim **foldMap** folosind **foldr**?

foldMap folosind foldr

<http://cmsc-16100.cs.uchicago.edu/2016/Lectures/13-monoid-foldable.php>

foldr :: (a -> b -> b) -> b -> t a -> b

foldMap :: Monoid m => (a -> m) -> t a -> m

```
foldMap f tr = foldr foo i tr      -- f :: a -> m
                where foo = ???    -- foo :: (a -> m -> m)
                          i = mempty
```

foldMap folosind foldr

<http://cmsc-16100.cs.uchicago.edu/2016/Lectures/13-monoid-foldable.php>

```
foldr    :: (a -> b -> b) -> b -> t a -> b
foldMap :: Monoid m => (a -> m) -> t a -> m
```

```
foldMap f tr = foldr foo i tr      -- f :: a -> m
                where foo  = ???    -- foo  :: (a -> m -> m)
                        i = mempty
```

```
foo = \x acc -> f x <> acc
     = \x acc -> (<>) (f x) acc
     = \x -> (<>) $ f x
     = \x -> ((<>) . f) x
     = (<>) . f
```

foldMap folosind foldr

<http://cmsc-16100.cs.uchicago.edu/2016/Lectures/13-monoid-foldable.php>

```
foldr    :: (a -> b -> b) -> b -> t a -> b
foldMap :: Monoid m => (a -> m) -> t a -> m
```

```
foldMap f tr = foldr foo i tr      -- f :: a -> m
               where foo  = ???    -- foo  :: (a -> m -> m)
                       i = mempty
```

```
foo = \x acc -> f x <> acc
    = \x acc -> (<>) (f x) acc
    = \x -> (<>) $ f x
    = \x -> ((<>) . f) x
    = (<>) . f
```

```
foldMap f = foldr (mappend . f) mempty
```

Foldable cu foldMap

```
instance Foldable BinaryTree where
```

```
  foldMap f (Leaf x)    = f x
```

```
  foldMap f (Node l r) = foldMap f l <> foldMap f r
```

```
tree1 = Node(Node(Leaf 1)(Leaf 2))(Node (Leaf 3)(Leaf 4))
```

```
treeS = Node (Node(Leaf "1")(Leaf "2"))
           (Node (Leaf "3")(Leaf "4"))
```

Avem definite automat **foldr** și alte funcții precum: **foldl**, **foldr'**, **foldr1**,...

```
*Main> foldr (++) [] treeS
"1234"
```

```
*Main> foldl (+) 0 tree1
10
```


Foldable cu foldMap

```
instance Foldable BinaryTree where
```

```
  foldMap f (Leaf x)    = f x
```

```
  foldMap f (Node l r) = foldMap f l <> foldMap f r
```

```
tree1 = Node(Node(Leaf 1)(Leaf 2))(Node (Leaf 3)(Leaf 4))
```

```
treeS = Node (Node(Leaf "1")(Leaf "2"))
           (Node (Leaf "3")(Leaf "4"))
```

Avem definite automat **foldr** și alte funcții precum: **foldl**, **foldr'**, **foldr1**,...

```
*Main> foldr (++) [] treeS
"1234"
```

```
*Main> foldl (+) 0 tree1
10
```

Cum definim **foldr** folosind **foldMap**?

foldr folosind foldMap

<https://en.wikibooks.org/wiki/Haskell/Foldable>

```
foldr    :: (a -> b -> b) -> b -> t a -> b  
foldMap :: Monoid m => (a -> m) -> t a -> m
```

foldr folosind foldMap

<https://en.wikibooks.org/wiki/Haskell/Foldable>

```
foldr    :: (a -> b -> b) -> b -> t a -> b
foldMap  :: Monoid m => (a -> m) -> t a -> m
```

Idee

```
foldr    :: (a -> (b -> b)) -> b -> t a -> b
```

- pentru fiecare element de tip **a** din **t a** se crează o funcție de tip **(b->b)**
*obținem, de exemplu, o lista de funcții sau
 un arbore care are ca frunze funcții*
- folosim faptul ca **(b->b)** este instanță a lui **Monoid** și aplicăm **foldMap**

foldr folosind foldMap

<https://en.wikibooks.org/wiki/Haskell/Foldable>

```
foldr    :: (a -> (b -> b)) -> b -> t a -> b
```

(b->b) instanță a lui **Monoid**

```
newtype Endo b = Endo { appEndo :: b -> b }
```

```
instance Monoid (Endo b) where
```

```
    mempty                = Endo id
```

```
    Endo g 'mappend' Endo f = Endo (g . f)
```

foldr folosind foldMap

<https://en.wikibooks.org/wiki/Haskell/Foldable>

```
foldr    :: (a -> (b -> b)) -> b -> t a -> b
```

(b->b) instanță a lui **Monoid**

```
newtype Endo b = Endo { appEndo :: b -> b }
```

```
instance Monoid (Endo b) where
```

```
    mempty                = Endo id
```

```
    Endo g 'mappend' Endo f = Endo (g . f)
```

Definim funcția ajutătoare

```
foldComposing :: (a -> (b -> b)) -> t a -> Endo b
```

astfel încât

```
foldr f i tr = appEndo (foldComposing f tr) $ i
```

foldr folosind foldMap

<https://en.wikibooks.org/wiki/Haskell/Foldable>

```
foldr    :: (a -> (b -> b)) -> b -> t a -> b  
foldComposing :: (a -> (b -> b)) -> t a -> Endo b
```

foldr folosind foldMap

<https://en.wikibooks.org/wiki/Haskell/Foldable>

```
foldr    :: (a -> (b -> b)) -> b -> t a -> b
```

```
foldComposing :: (a -> (b -> b)) -> t a -> Endo b
```

```
foldComposing f = foldMap (Endo . f)
```

foldr folosind foldMap

<https://en.wikibooks.org/wiki/Haskell/Foldable>

```
foldr    :: (a -> (b -> b)) -> b -> t a -> b
foldComposing :: (a -> (b -> b)) -> t a -> Endo b

foldComposing f = foldMap (Endo . f)
```

Exemplu:

```
foldComposing (+) [1, 2, 3]
foldMap (Endo . (+)) [1, 2, 3]
(Endo . (+)) 1 <> (Endo . (+)) 2 <> (Endo . (+)) 3
Endo (+1) <> Endo (+2) <> Endo (+3)
Endo ((+1) . (+2) . (+3))
Endo (+6)
```


foldr folosind foldMap

<https://en.wikibooks.org/wiki/Haskell/Foldable>

```
foldr    :: (a -> (b -> b)) -> b -> t a -> b
foldComposing :: (a -> (b -> b)) -> t a -> Endo b

foldComposing f = foldMap (Endo . f)
```

Exemplu:

```
foldComposing (+) [1, 2, 3]
foldMap (Endo . (+)) [1, 2, 3]
(Endo . (+)) 1 <> (Endo . (+)) 2 <> (Endo . (+)) 3
Endo (+1) <> Endo (+2) <> Endo (+3)
Endo ((+1) . (+2) . (+3))
Endo (+6)
```

```
foldr f i tr = appEndo (foldComposing f tr) $ i
```