LCS

Etapa 3 - relația de recurență Lmax[i][0]=0, i=1,...,n • Lmax[0][j]=0, j=1,...,m Lmax[i][j]=Lmax[i-1][j-1]+1, daca X[i]=Y[j] Lmax[i][j]=max{Lmax[i-1][j], Lmax[i][j-1], for $(i = 0; i \le m; i++)$ for $(j = 0; j \le n; j++)$ if((i == 0) | | (j == 0)) lcs[i][j] = 0;else if(s[i-1] == t[i-1]) |cs[i][i] = 1 + |cs[i-1][i-1];else lcs[i][j] = lcs[i-1][j] > lcs[i][j-1] ? lcs[i-1][j] : lcs[i][j-1];printf("Lungimea maxima a subsirului comun: %d\n", lcs[n][m]); i = m; j = n;printf("Subsirul maximal comun:\n"); while(lcs[i][i] != 0) if (s[i-1] == t[j-1]) { printf("%c " , s[i-1]); i--; j--; } else if (lcs[i][j] == lcs[i-1][j]) i--; else i--;

Primul sir: program

Al doilea sir: roman