



## Andreas Karrenbauer

**Summer 2024** 

## Assignments for Optimization

Assignment Sheet 1 Due: Wednesday, 24 April 2024

The following students have contributed to the solution of this sheet:

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The cut-off time for submitting this homework 13:00 CEST (using Overleaf).

Exercise	1	2	3	4	$\sum$
Points					
(self-assessment)					
Check					
(at random)					

**Definition 1 (Convex set)** Let S be a set of points in an Euclidean space. S is convex if for any two points  $x, y \in S$  and any  $t \in (0,1)$  it holds that  $tx + (1-t)y \in S$ .

**Definition 2 (Convex hull)** Let  $S := \{x_1, \ldots, x_n\}$  be a set of points in an Euclidean space. We define  $conv(S) := \{\sum_{i=1}^n \lambda_i x_i : \sum_{i=1}^n \lambda_i = 1, \lambda_i \geq 0 \ \forall i \in [n]\}$  as the convex hull of S.

Exercise 1 (10 points) Prove the following statements:

- a) A polyhedron is convex. (2 points)
- b) Let  $x_1, \ldots, x_n \in \mathbb{R}^n$ . Show that  $conv(\{x_1, \ldots, x_n\})$  is convex. (2 points)
- c) The intersection of a finite number of convex sets is convex. (3 points)
- d) Let  $P := \{x \in \mathbb{R}^n : Ax \leq b\}$  be a polyhedron and  $(\pi, \pi_0) \in \mathbb{Z}^n \times \mathbb{Z}$ . We define the polyhedra

$$P_L := P \cap \{x \in \mathbb{R}^n : \pi^T x \le \pi_0\}, P_R := P \cap \{x \in \mathbb{R}^n : \pi^T x \ge \pi_0 + 1\}$$

Show that  $P \cap \mathbb{Z}^n = (P_L \cup P_R) \cap \mathbb{Z}^n$ . (3 points)

**Exercise 2** (10 points) Let  $A' \in \mathbb{Z}^{m' \times n}$ ,  $b' \in \mathbb{Z}^{m'}$ ,  $P := \{x \in \mathbb{R}^n : A'x \leq b'\}$ . The Integer Feasibility problem is to decide if  $P \cap \mathbb{Z}^n \neq \emptyset$ . Show that Integer Feasibility is NP-hard using the following scheme.

- a) Think of a reduction from any NP-hard problem that you prefer (e.g., 3-SAT, Vertex Cover, Independent Set). (6 points)
- b) Show that this reduction is polynomial. (4 points)

Exercise 3 (10 points) Consider the following problems and model them as ILPs.

- a) Given a set of elements  $U := [m] := \{1, \ldots, m\}$ , a set  $S = \{S_1, \ldots, S_n\}$  of n subsets of U (i.e.  $S_i \subseteq U$  for  $i \in [n]$ ), and a positive integer k. Find a subset  $\{S_{i_1}, \ldots, S_{i_j}\} \subset S$  with at most k members (that is,  $j \leq k$ ) such that the union of the selected subsets has maximum size (that is,  $\bigcup_{k \in [j]} S_{i_k}$  is maximized). (5 points)
- b) Given an undirected graph G = (V, E), find the largest subset  $Q \subseteq V$  such that  $\{u, v\} \in E$  for all distinct  $u, v \in Q$ . (5 points)

**Exercise 4** (10 points) Consider the following optimization problem in  $\mathbb{R}^2$ .

$$\begin{array}{lll} \text{maximize} & x_1 + x_2 \\ \text{subject to} & -x_1 + 2x_2 & \leq 3 \\ & 3x_1 + 4x_2 & \leq 22 \\ & -x_1 - 3x_2 & \leq -3 \\ & x_1, x_2 & \geq 0 \\ & x_1, x_2 & \text{integer} \end{array}$$

- a) Draw the polyhedron of the integer linear program. (4 points)
- b) Solve the problem with the minimum number of calls to an integer feasibility oracle. (6 points)