

实验报告

【实验目的】

1. 掌握椭圆曲线上的运算和常见的椭圆曲线密码算法;
2. 了解基于ECC的伪随机数生成算法和基于椭圆曲线的商用密码算法。

【实验环境】

1. 语言: C
2. 平台: clion 2021.2 版本

【实验内容】

一、ECC - 四则运算

1. 算法流程

1. ECC加法减法

对于点P, Q, 如果两点横坐标相同但纵坐标相反, 那么相加结果为无穷远点;
其中一个为无穷远点时则结果等于另一个点;
否则就有公式如下:

$$\delta = \begin{cases} (3Px^2 + a)(\text{mod } p) \times \text{invmod}(2Py, p)(\text{mod } p) \\ (Py - Qy)(\text{mod } p) \times \text{invmod}((Px - Qx), p)(\text{mod } p) \end{cases}$$
$$Rx = \delta^2 - Px - Qx(\text{mod } p)$$
$$Ry = \delta \times (Px - Qx) - Py(\text{mod } p)$$

R即为结果。

对于减法则只需要将点Q点坐标改为负数即可。

2. ECC乘法除法

即整数k与点坐标的乘法, 意为倍点。这里需要解决的是快速求倍点的问题。继续沿用第一次实验中的快速模

幂算法, 将整数k化为二进制, 将模乘运算变为点加即可, 伪代码如下:

Algorithm 1 ECC的数乘运算

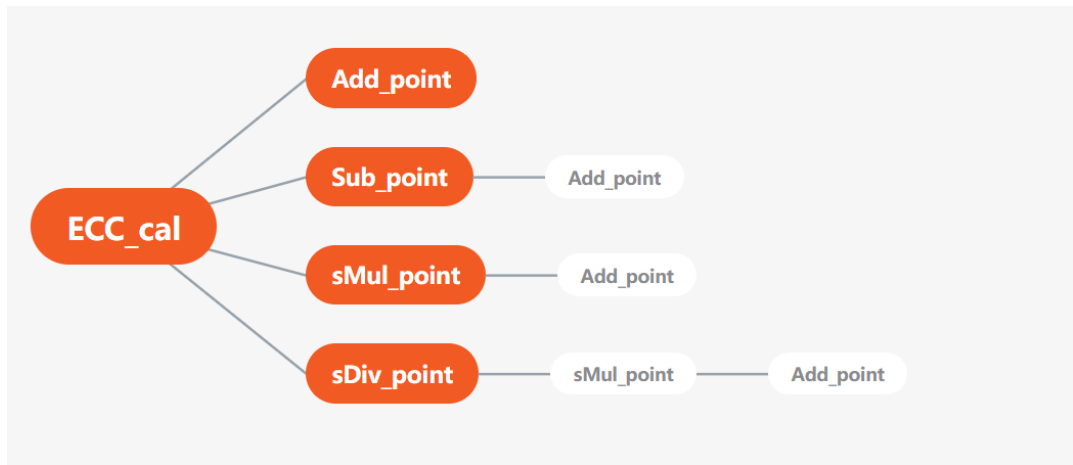
Input: *point p, mpz_t k, ECC_curve C*

Output: *point ans*

```
1: if k%2 == 1 then
2:   ans = p
3: end if
4: k = k/2
5: for k != 0 do
6:   p = p + p
7:   if k%2 == 1 then
8:     ans+ = p
9:   end if
10:  k = k/2
11: end for
```

对于除法, 情况也类似。即求出k的关于模数p的逆元, 将其作为整数参数传入乘法函数即可。

○ 函数调用图：



2. 测试样例及结果截图：

本地测试样例的运行结果如下所示：

```
"E:\_drive\clicon\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_cal.exe"
118792089210356240756420345214020892766250353991926191454421193933289684991999
118792089210356240756420345214020892766250353991926191454421193933289684991999
185089190222381800113072901827995639221458448578012075254897346196103069175443
664901889980129086613443884082074452575157116031103742365434905633820928013192
0465341252842791972320613308995459491213526647800598970633596412071601640913
664901889980129086613443884082074452575157116031103742365434905633820928013192
0465341252842791972320613308995459491213526647800598970633596412071601640913
109099916468336434299323842074559377173634036498855097587235085319472956
91829719240076595600910287219737299259627413891073174690491219092963035830325 314748222768498591041231146460709769749214
01394140157637420547181522913249875
0 0
57373321148051506091723093715148369162664271829435462559258806045759894672711 478105239341029923072192254309559256213954
44436399395870291687870429993066087
65262872771489155421945763814501787279403072416537416049831455048807745490812 852554443924726005842158206461884433642146
32091202947623642966545271132050758
Process finished with exit code 0

"E:\_drive\clicon\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_cal.exe"
118792089210356240756420345214020892766250353991926191454421193933289684991999
118792089210356240756420345214020892766250353991926191454421193933289684991999
185089190222381800113072901827995639221458448578012075254897346196103069175443
3047588008780961444347000912770690330667593770626312123619394819408538269140
90023601239643517308939790510900951707001876773407990710202404894089212360694
144887749067412590661957190874423408258661779066480012808012316926101212902384
5120871017130140832316096967145897741690421850918580077708230138429775187679
6171367966392211550327354279193125208614726020731250960120457053519700
76418670396892355653243957783665229567941867955910264199035186565505072408164 740130061550006943653380894990213374683302
5504950958270819328838872254906572
78895673407062940374715955825638547545005250291461468259382893293918199494264 220431590064127472402681065650530532965350
81007178340513578111610711191864756
8782182299491035655407366522213245608339654612217932552302421142270846681994 7814446842864845172049965856717415889780340
3506995461546832258623034756259383
133589242127889012189855145261361973853304174397400758918867054805888167102912 153808475768360990726523002710278200075379
1334843337231645846731885387715563
Process finished with exit code 0

"E:\_drive\clicon\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_cal.exe"
118792089210356240756420345214020892766250353991926191454421193933289684991999
118792089210356240756420345214020892766250353991926191454421193933289684991999
185089190222381800113072901827995639221458448578012075254897346196103069175443
614000108663701755347267550265267036422106036884610188438874385931303189573
400010350501720013995763627790702656151173077906702162748249716709420042009
88676605730604584948906802571952219104053026319105192479396867314569702342062
260140879533100808652265530537331909295338660822977861066649406404312102114
145169512400758286200212677930791136911712126778061663620991432494863976849
1513938498219190048759198378835370844090831636605751713241967093574427994743 1121938196345632583614385494949414167041864
937689326132933650053264392471602
112980429052981945792806805361750017649838533159287688989234058026821704764321 19896334991812450652080092926702658620188
688821228463595414901333673743909614
12637770518718272361069030959675909011950070361104001676990341902418348821166 409162228750944407497055609713911886363555
20322606034728049992680090883970559
72519289070001971517117073049519824288518639052252472087812461978663312943158 977821225842917535958913934187500212184061
45461157584281766145451975450555990
Process finished with exit code 0
```

```
"E:\E_drive\clion\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_cal.exe"
118792089210306240756420740214020892766250353991924191454421193933289606991999
118792089210306240756420740214020892766250353991924191454421193933289606991999
10509919022201000113072901027995639221450440970012070254007346196103069175443
1355215697908753363935906304612651321612502271075937256096304646962700530197
27103113344768520834601013461646548772036099432320515037653070109124730604412
090190125479308701301707376313209708049637371004031044919012473274996490499969
113625264091997307711342929607010016579774033982218109720700200914506087097272
740099011770021203953029989001301077104605466201306716706630296021974303
46416434919624558033056716946399905725184054194477621641981818870586970614985 870905641694899167862170143314564693880074
30460863965606250437495188833946443
104673815949190848064823792929504530617300375265159409007922482354341802926247 49720414426250787946558522178489712574956
91090986485116773778132342334607845
12155434101822315229785474237673275927256428374157073820978014707268450290347 592607920940730025477492132760973509069057
09657685638377740494081362239475981
49594380326687913930629211348076681153717837731434921282465624553827847405191 102275878924654171059450634351933573693661
31540118020353760541918655258792286
Process finished with exit code 0
```

二、ECC - 公钥加解密

1. 算法流程

下面是利用椭圆曲线进行加密通信的过程：

- 1、用户A选定一条椭圆曲线 $E_p(a,b)$ ，并取椭圆曲线上一点，作为基点 G 。
- 2、用户A选择一个私有密钥 k ，并生成公开密钥 $K=kG$ 。
- 3、用户A将 $E_p(a,b)$ 和点 K, G 传给用户B。
- 4、用户B接到信息后，将待传输的明文编码到 $E_p(a,b)$ 上一点 M （编码方法很多，这里不作讨论），并产生一个随机整数 r ($r < n$)。
- 5、用户B计算 $C_2 = M + rK, C_1 = rG$ 。
- 6、用户B将 (C_1, C_2) 传输给A。
- 7、A在接收到B的密文信息后，计算 $C_2 - kC_1 = M$ 得到明文，解码即可得到消息。

算法流程图如下：

○ 算法流程图：

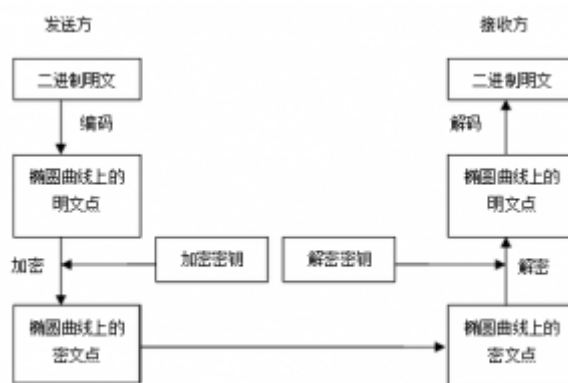


图 1 椭圆曲线密码算法的加密解密流程图

○ 伪代码如下：

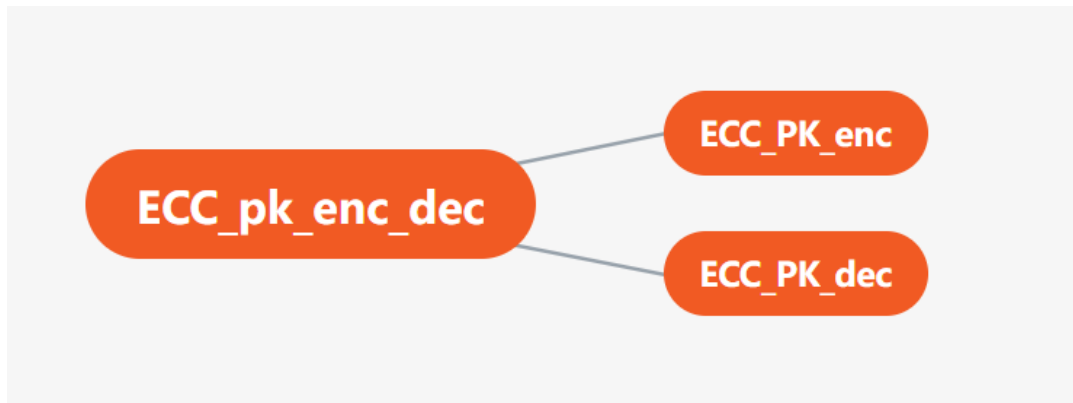
算法 1 ECC加解密

输入: $p, a, b, G, op, (Pm, k, Pb)/(C1, C2, nb)$

输出: $(C1, C2)/Pm$

```
1: function ECC( $p, a, b, G, op, (Pm, k, Pb)/(C1, C2, nb)$ )
2:   if  $op == 1$  then
3:      $C1 = kG$ 
4:      $C2 = Pm + kPb$ 
5:     return  $(C1, C2)$ 
6:   else
7:      $M = C2 - nbC1$ 
8:     return  $M$ 
9:   end if
10: end function
```

- 函数调用图如下:



2. 测试样例及结果截图:

本地测试样例的运行结果如下所示:

```
"E:\_drive\clion\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_pk_enc_dec.exe"
115792009210356248756420340214020092766250353991924191454421191933209604991996
115792009210356248756420340214020092766250353991924191454421191933209604991996
18509919022201800113072981827955639221456448570012075254857346196103069175443
2296314604723705050947953136250074570002567295341616970375194840004139615431
85132369209820560025618990617112496413000300631904050003203536607500077201560
1
6222030967324995506943059594303043316990201421105076004209933343560009233015
9063699496346492210878763452203053616231930003241827298170711426179690619737
01419392794900372206772030160930400931071740967362193779666947296900000
00690242911250790190307097601000251691466959665004540024100127307663566707
46820310853200066427314921969913203667136064079000112361092669654799097610074
89583431844721923129778721014009509006972090847660911645111898104009749186163 269307114890302101141102411704530064878973
95422517206355196890099778264510599
103440095822431431181554656770186818550669543322647226082461390490749218093950 20457019858049175615069165493796980004756
588511726052018196771904984755957123
Process finished with exit code 0

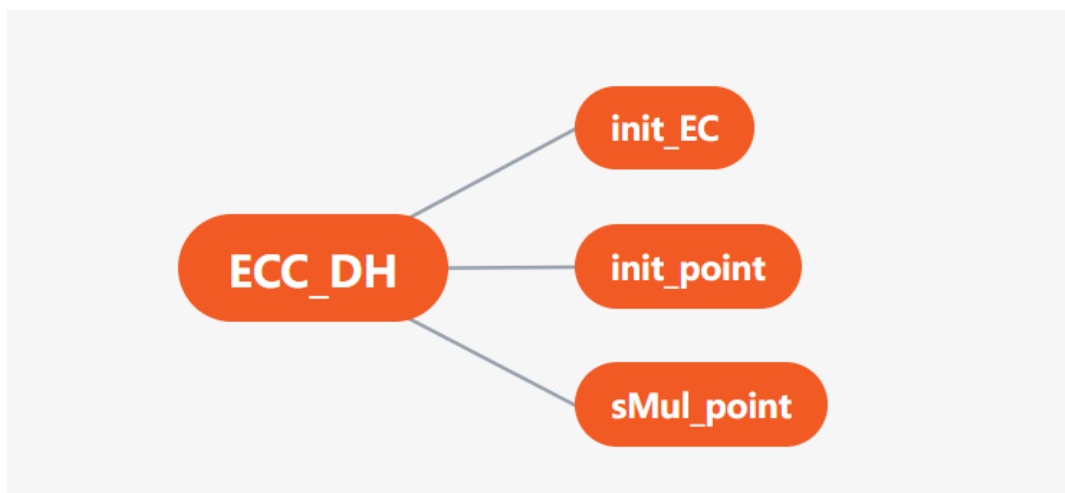
"E:\_drive\clion\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_pk_enc_dec.exe"
115792009210356248756420340214020092766250353991924191454421191933209604991996
115792009210356248756420340214020092766250353991924191454421191933209604991996
18509919022201800113072981827955639221456448570012075254857346196103069175443
2296314604723705050947953136250074570002567295341616970375194840004139615431
85132369209820560025618990617112496413000300631904050003203536607500077201560
0
6348961704231444103192743400167176760735397105250750730010361305127672364000
41730001692067197137950635040045231615660996567203100053910791740253018095163
07476900951670009791273002979092960516640912016001637796755117261903304423
10212109000102010003000639494300141231209726291312230500749410026645266390070
6385094248020599064256736496506600660061641279362962129000709550436131
9351157407290143374074872232898911884008298332267475567954737644846874112389 855600490203970259106461156246452983425447
129230744639232161292330016598209
Process finished with exit code 0
```

三、ECC - DH密钥交换协议

1. 算法流程

过程与Diffie-Hellman协议一致，这里一方将自己产生的随机数与基点相乘后传输给另一方，另一方再将该点与自己的随机数相乘。双方的密钥就都是 $K = X_a * X_b * G$ 。

- 函数调用图：



2. 测试样例及结果截图：

本地测试样例的运行结果如下所示：

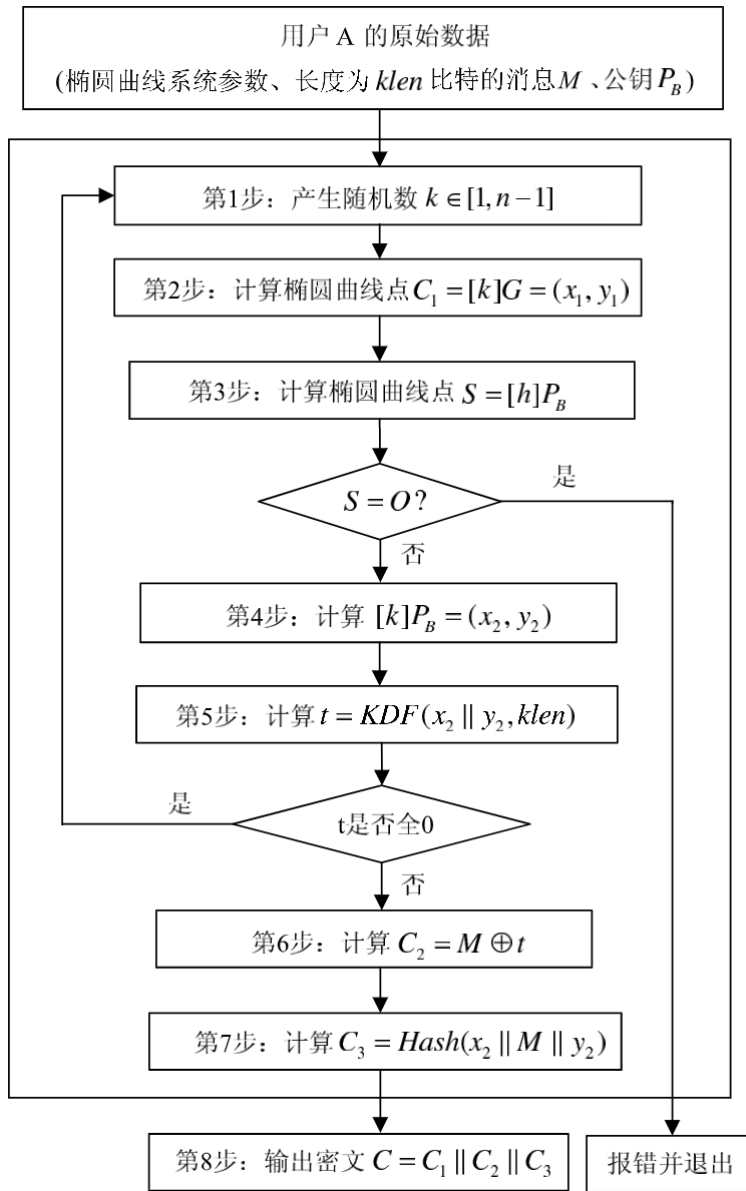
```
"E:\E_drive\clion\Project_List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_DH.exe"
115792089210356240756420345214020892766250351991924191454421193913209684991999
115792089210356240756420345214020892766250351991924191454421193913209684991999
18580919022201000113072901027955639221450440570012075254057346196103069175443
22963146547237080859479531362550074570802567295341616970375194840604139615431
05132369209020560025610990617112496413000300631904505003203536607500077201560
01100104610913722570604425355170616100410192111930555045711000030425000149
109900095266094733630946095666100030209722025233241660101030323414352037044396
396103300930611412395111936421564164706419275136927269416250005434604831729
17563276920628041528213370371379123587681909252989227496921943682543883264878 162907300695687139410749046539969910089816
85386084447436564101776231176534603
Process finished with exit code 0

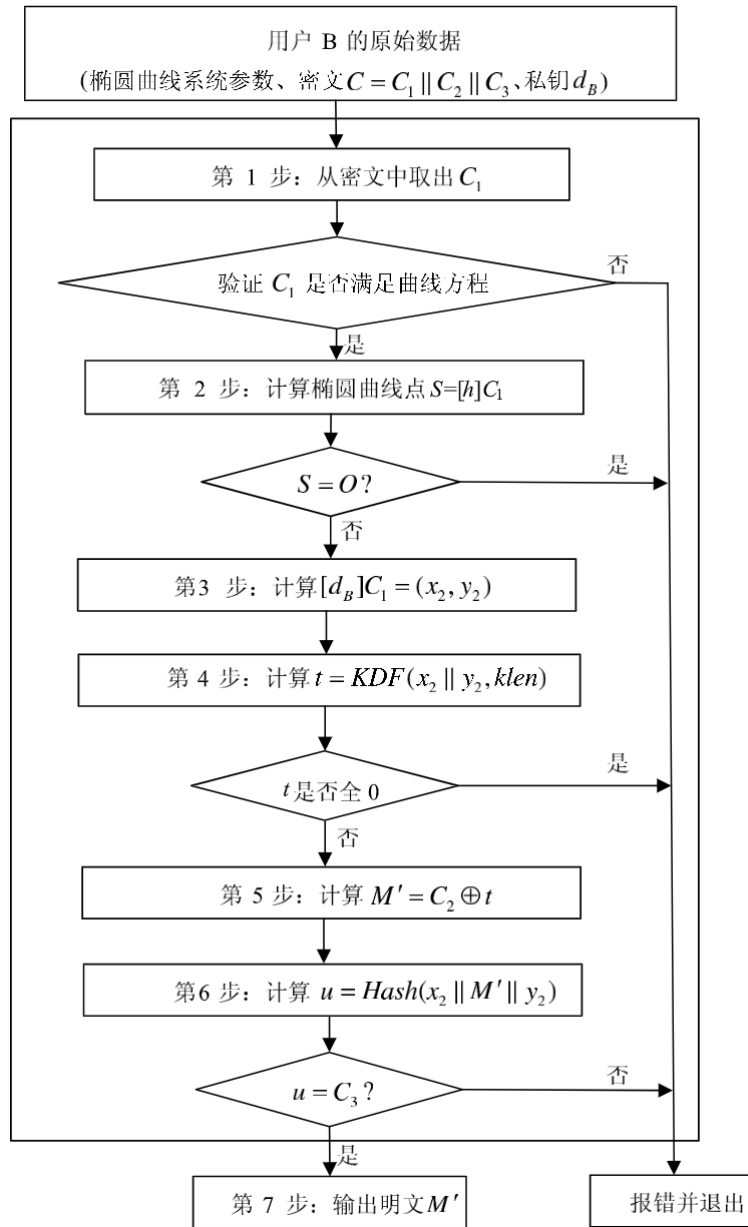
"E:\E_drive\clion\Project_List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_DH.exe"
115792089210356240756420345214020892766250351991924191454421193913209684991999
115792089210356240756420345214020892766250351991924191454421193913209684991999
18580919022201000113072901027955639221450440570012075254057346196103069175443
22963146547237080859479531362550074570802567295341616970375194840604139615431
05132369209020560025610990617112496413000300631904505003203536607500077201560
20165770951009370481409000442007347350569267400136481696605440041696184829
32440691193999261301314400034434061037205304746039122510920973565070356020367
44304050000701313197990303399207915644003330009170910907161710276202052040944
12512346806273286597779399744481781487246936530763666061969314356007739153319 965226313079202299373022269442812352712248
04951260748197377933029315574510741
Process finished with exit code 0
```

四、SM2 - 公钥加密

1. 算法流程

- 原理：





伪代码：

1. 伪代码：

算法 2 SM2

输入： $p, a, b, G, op, M/C, (Pb, k)/db$

输出： M/C

```

1: function SM2( $p, a, b, G, op, M/C, (Pb, k)/db$ )
2:   if  $op == 1$  then
3:      $C1 = kG$ 
4:      $tmp(x_2, y_2) = kPb$ 
5:      $t = KDF(x_2 \parallel y_2, klen)$ 
6:      $C2 = M \oplus t$ 
7:      $C3 = hash(x_2 \parallel M \parallel y_2)$ 
8:     return  $C1 \parallel C2 \parallel C3$ 
9:   else
10:     $C1 = M[2 * l + 1]$ 
11:     $tmp(x_2, y_2) = dbC1$ 
12:     $t = KDF(x_2 \parallel y_2, klen)$ 
13:     $C2 = M[2 * l + 1 : 2 * l + 1 + (klen // 8)]$ 
14:     $M = C2 \oplus t$ 
15:    return  $M$ 
16:   end if
17: end function
  
```

15: **end function**

```
graph LR; ECC_pk_code --> ECC_encode; ECC_pk_code --> ECC_decode; ECC_encode --> init_point; ECC_encode --> sMul_point; ECC_encode --> KDF_SHA256; ECC_encode --> SHA256; ECC_decode --> init_point; ECC_decode --> sMul_point; ECC_decode --> KDF_SHA256; ECC_decode --> SHA256;
```

The diagram illustrates the structure of the ECC_pk_code. It is a hierarchical tree where the root node, ECC_pk_code, branches into two main categories: ECC_encode and ECC_decode. Each of these categories further branches into four sub-categories: init_point, sMul_point, KDF_SHA256, and SHA256.

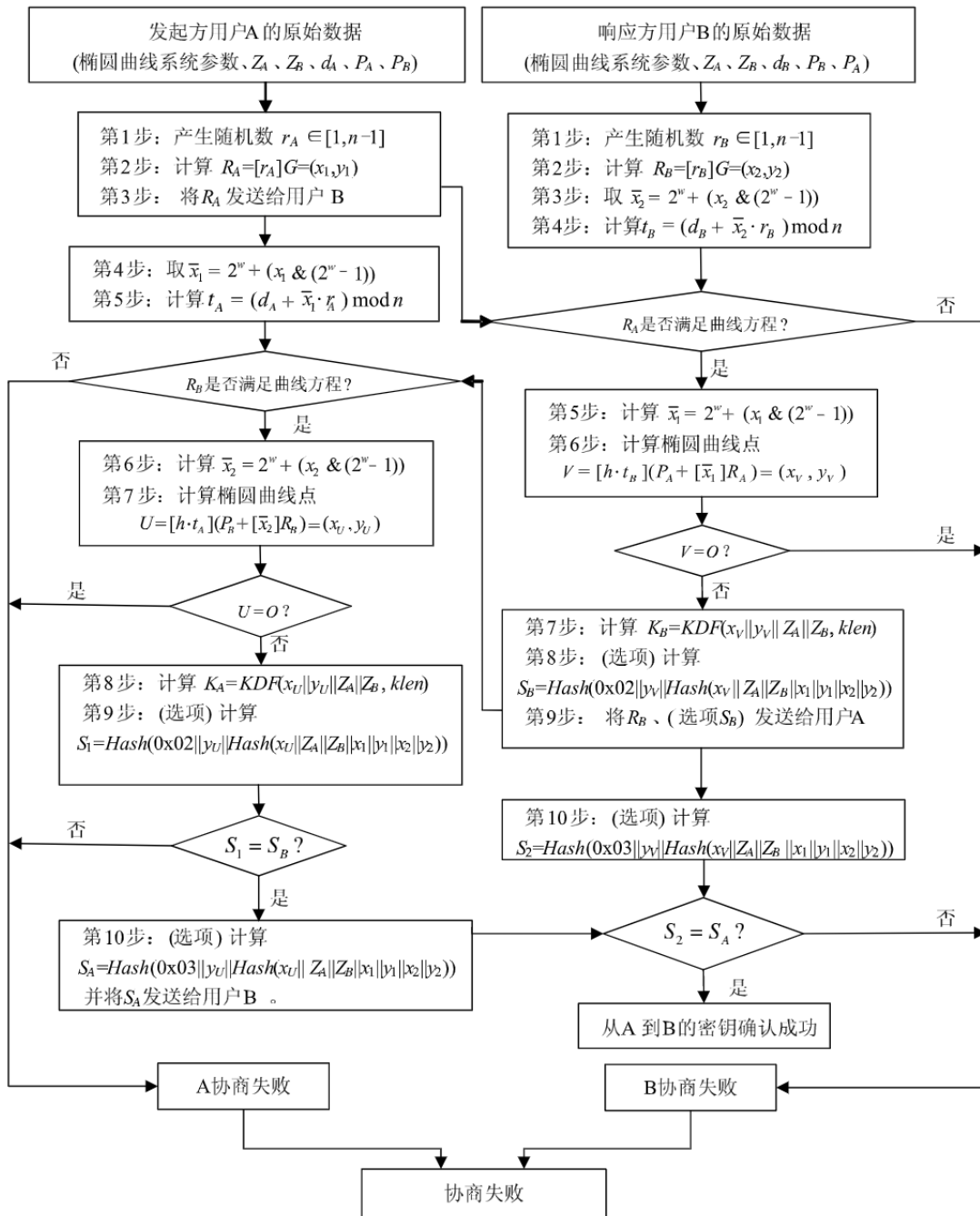
本地测试样例的运行结果如下所示:

```
"E:\_drive\c\ion\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_pk_code.exe"
a027870200924209a185a0017151521989a41a297121499a02209958370876a38005c19a1187
5a4a205298589947a080a43858a298570274016718417263133626858979785a5916325972248
431831851916001346a1a2550a9850180123187613551910337a09a391853873370470090074
2990851425a0703a123aa18a0000a7770023a3a3a499a0a2916a71209002818329800180225085
799a8593717970041915790390447892157254200a770830401260a12300519640063234001314
25a
1
0a09a0a372797074a9af0a00737a61a0a417244
394a6a142001137288a070019005200012003243716182190958300239831a0839a0243019794
533123a347099200212197984a4a0a031a120007141879a0291924009a71035137a9a1810274
3a08052a9920a43109977037519a18321907097a285780a4301827001209a09a0a01152a290255
0x04245c26f68b1dddb12c4b6bf9f2b6d5fe0a383b0d18d1c144abf17f6252e776cb9264c2a7e88e52b19903fdc47378f605e36811f5c07423a2
4b84400f01b82296e9aace2bb92cad649fe2c035689785da33be89139d07853100efa763f60cbe30099ea3df7f8f364f9d10a5e988e3c5aafc
Process finished with exit code 0

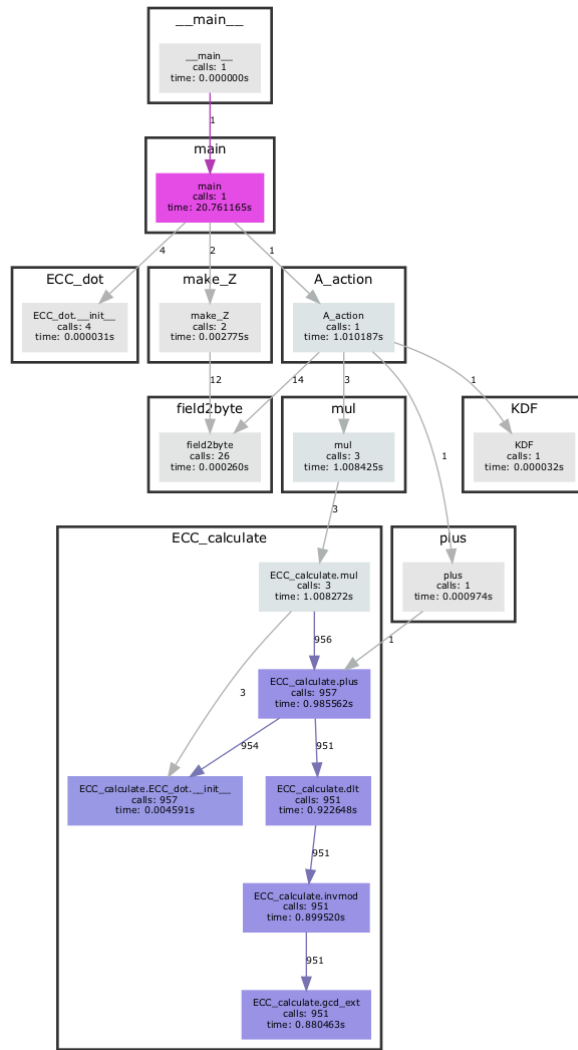
"E:\_drive\c\ion\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_pk_code.exe"
a017901151719a717a4a1a27a007071a3a0a17185070a3a2120a099
41920a71002077a0a9009070005240207180090023049173099a117
50a08025a08930a72a01829a421a4910a91212081782209899173a07a9
18349a0a7100a0703a51a019102a90a421a49199a0071a75721952 a0972a03a0005182a01770120370a4a1770a01972200120a09170
192
0
0a0a319a00b12174a9f071a0a7a0a3180031a0a01f0a0a0cf57071f1f100970a4777773001a01a82a411a0f0a0a1a1f0a10a0a13a091c409910a38a0a0f0a1a07051a0a42a0
4fe1e1a741a0a0f770a4220100a0a20a7009a050a0f1c412e0a0f0c
21700919900100a4219a1070122019170a000123117972174a08273113
0x1145141919810ab19260017cd947866efedcba
Process finished with exit code 0
```

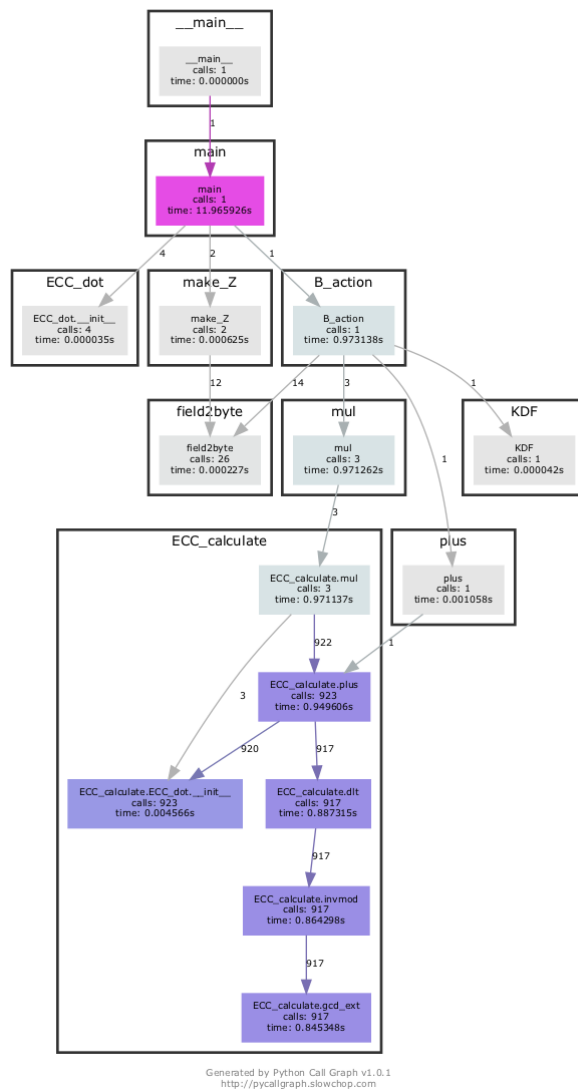

[illegible][illegible]

1. 算法流程



。函数调用图：





2. 测试样例及结果截图：

本地测试样例的运行结果如下所示：



✓ 您已通过本题！

查看历史提交

评测编号 ↓	提交时间	提交状态	代码语言	最大运行时间
22342	2022-05-18 21:48:18	Accepted	Python	109ms

六、【选做】基于ECC的公钥加密快速实现

1. 算法流程

在本题中我尝试了蒙哥马利算法。

。原理：

1. 蒙哥马利算法：

蒙哥马利算法（MontgomeryAlgorithm）蒙哥马利约简、模乘、模幂

Montgomery Algorithm(蒙哥马利算法)

蒙哥马利算法分为3种，蒙哥马利模乘，蒙哥马利约简，蒙哥马利模幂

1、从蒙哥马利模乘说起

模乘是为了计算 $ab \pmod N$ 。普通算法中，在计算模N时，利用的是带余除法，除法运算需要太多次乘法，计算复杂度较高，蒙哥马利算法的思想就是利用进制表示简化除法运算，转化成位运算。

- 蒙哥马利形式：为了计算 $ab \pmod N$ ，找一个 R ，然后使得 $a' \equiv aR \pmod N, b' \equiv bR \pmod N$ ，这就是蒙哥马利形式。
- 这个 R 不是随便一个数，需要满足两个条件：1: $R = 2^k > N$,其中 k 是满足条件的最小的正数，这个取法能够保证除以 R 就相当于右移 k 位，避免除法运算；2: $\gcd(R, N) = 1$ ，这使得能够求出下面的 m 。
- 蒙哥马利形式有什么用？往下看

前面只是铺垫，下面才真正开始。为了计算一开始的 $ab \pmod N$ ，需要用到上面的蒙哥马利形式。令 $X = a'b'$ ，我们可以设计一个函数来计算 $XR^{-1} \pmod N$ ，简单计算发现这个函数的计算结果为 $X_1 = XR^{-1} = a'b'R^{-1} = abR \pmod N$ ，这样在调用一遍函数计算 $X_1R^{-1} \pmod N$ 就得到我们最终需要的结果 $ab \pmod N$ 了。我们称这个算法叫蒙哥马利约简算法。所以说，蒙哥马利约简的产生是为了蒙哥马利模乘计算服务的。

2、蒙哥马利约简

根据上述分析可知蒙哥马利算法的核心在于蒙哥马利约简。而且前面提到，蒙哥马利算法的主要思想是把取模运算变得简单，到底是怎么做到的呢？

主要思想：蒙哥马利约简是计算 $XR^{-1} \pmod N$ ，这相当于 $\frac{X}{R} \pmod N$ ，如何才能避免除法？从前面的 R 的定义中我们知道， $R = 2^k$ ，所以 $\frac{X}{R} = X \gg k$ （ X 右移 k 位）。但是新的问题出现了，右移 k 位可能会抹掉 X 的低位中的一些1，如 $7 \div 4 = 0b111 \gg 2 = 0b1 = 1$ ，这个不是精确计算，而是向下取整的除法。当且仅当 X 是 R 的整数倍时， $X \div R$ 严格等于 $X \gg k$ 。所以我们实际上是找一个 m ，使得 $X + mN$ 是 R 的倍数，这样计算 $\frac{X+mN}{R} \pmod N$ 就可以了。

m 如何找

根据 R 的定义， $\gcd(R, N) = 1$ ，根据扩展的欧几里得算法，有 $RR' - NN' = 1$ 并且有 $0 < N' < R, 0 < R' < N < R_0$ 。

$$\begin{aligned} X + mN &\equiv 0 \pmod R \\ XN' + mNN' &\equiv 0 \pmod R \\ XN' + m(RR' - 1) &\equiv 0 \pmod R \\ XN' &\equiv m \pmod R \end{aligned}$$

这样就求出了 m 。

约简算法总流程

1. 计算 $N' = N^{-1} \pmod R$, 计算 $m = XN' \pmod R$;

2. 计算 $y = \frac{X+mN}{R}$:将 $X + mN$ 右移 k 位;

3. 若 $y > N$, 则 $y = y - N$;

这时的 y 满足: $0 < y < 2N$ 。因为

$$X < N^2, m < R, N < R$$

所以

$$\frac{X + mN}{R} < \frac{N^2 + R \cdot N}{R} < \frac{RN + R \cdot N}{R} = 2N$$

4. 返回 y 。

3、蒙哥马利模乘流程

Montgomery Multiply(a, b, N):

1. 计算 $a' = aR \pmod N, b' = bR \pmod N, X = a'b'$;

2. 调用蒙哥马利约简算法: $X_1 = \text{Montgomery reduction}(X, R, N) = X/R = abR \pmod N$;

3. 再调用蒙哥马利约简算法: $y = \text{Montgomery reduction}(X_1, R, N) = X_1/R = ab \pmod N$;

4. 返回 y 。

4、复杂度分析

蒙哥马利算法是为了简化模 N 的复杂度, 当 N 是比较大的数时, 模 N 需要多次的加、减、乘运算。从上述过程看来, 蒙哥马利约简的复杂度确实降低了, 因为只有模 R 的移位运算。但是看蒙哥马利模乘的流程, 在第一步中进行蒙哥马利表示时就计算了两次模 N , $(a' = aR \pmod N), b' = bR \pmod N, X = a'b'$, 总的来看复杂度也没有降低啊? 实际上, 第一步可以看成是蒙哥马利的预先计算。在硬件实现中, 先把预先计算的算好, 在后面运行就会快很多。尤其是当出现大量的模乘运算时, 可以通过并行运算进行预计算, 这会大量节省运行时间。

5、蒙哥马利模幂

当进行模幂运算时, 也可以利用蒙哥马利算法。如计算 $a^e \pmod N$ 。根据, 模幂运算中就有许多步的模乘运算, 这恰恰可以发挥蒙哥马利模乘的优势。具体的, 在重复平方方法的每一次模乘中都利用蒙哥马利模乘进行计算, 把需要的参数提前计算好就行。

算法的本质在于预处理计算, 进而简化过程, 加快运算速度。

2. 分析:

在模拟实现后发现并没有速度上的明显提升, 原因就在于gmp大数库中的模乘运算已经采用了这样的优化方式。

2. 测试样例及结果截图:

本地测试样例的运行结果如下所示:

```
"E:\E_drive\clion\Project_List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_Fast_enc_dec.exe"
11579208921035240756420340214020092766200353991924191654421191913209604991999
11579208921035240756420340214020092766200353991924191654421191913209604991999
10800919022201800113072901027950639221630440570012070294057346196103069179443
22963146547237050099479931362500074970001967299341616970379194040004139619431
08132169209020560025610990617112696413008300631904000083203536607500877201568
1
6065269490002100915615070009000039020551140764000003050070946647960945092102
972400066623091406333947300494196120167700904032023719329772202320742149407074
75620043076150007000370100296300916712317203964371566779374000036170990
96813097921121266207110203191003976449903069076196969700912196430013315470096
10720292036202791001000290060992459927703519692540074660076619311130191933901
99
2836196863401915602599022753213785933891666039351133884663098349326957490724 2742666222963389924701676552899180233370008
0038105369561847065403315807331346
84654655179098796912662543007901138302884985227327354714303500052960379059857 906878510114102502111975388065406507874686
46750805872850704138724948523318145
Process finished with exit code 0
```

```
"E:\E_drive\clion\Project List\cryptography\Crypto_Experiment\cmake-build-debug\Exp8--ECC_Fast_enc_dec.exe"
115792009210306248796420340214020092766230333991924191456421193933209604991999
1157920092103062487964203445214020092766250351991924191456421193933209604991996
1050891902220100011307290102798563922145044057001207325405736196103069175443
2796316647237000559479531362550076570002567295341616970375194060604139615431
05132369209020560029610990617112496413080300631904000003203536607500077201560
0
2036196063401910602099022753213700913091666039351133004663090349326957690724
2742666222963309924701676002099100213370000003010536961067065603115007331346
0665465517909079691266254300790130302004905227327354714303500052960379059057
90687001011410250211197530006540600707460646700000072050704130724940523310140
75679514760352709663961092596921506647349330721274567052923221116904421
00
60652694590882909356250788599098839028553140764880503559070946642968945892102 972485666236914563339473884941961201677959
04532823719329772252328742149457874
Process finished with exit code 0
```

七：讨论与思考

ECC算法抗攻击性强，CPU占用少，内容使用少，网络消耗低，加密速度快，更高的扩展性。而RSA的数学原理简单，在工程应用中易于实现，且已经较为成熟。