



π -CIPHER V2.0

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About π - Cipher

- Nonce-based authenticated encryption cipher with associated data
- Sponge based
 - key-less permutation function based on ARX operations
 - supports 16, 32 and 64-bit words
- Security in the range of 96 to 256 bits
- Uses secret message number (SMN)





What is new!?

• Padding rule

- Gaëtan Leurent and Thomas Fuhr Observation on picipher. Message on the cryptocompetitions mailing list, Nov, 2014



• The rule is now simple:

- "Append 1 in any case, and fill the rest of the block with 0s"

\mathbf{M}_1	M_2	• • •	$M_{\rm m}$	10*	1
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What is changed?

• The number of rounds R

Now R = 3
 (previously it was R = 4)







What is changed in explanation?

- From v2.0 π -Cipher supports the concept of "open authorship"
 - it gives opportunity to all people that contribute anyhow in the development of π -Cipher:
 - a tweak is introduced due to an analysis of the cipher,
 - a new mode of operation is proposed,
 - a new significantly different and improved implementation is given
 - if they want, they can be added to the list of designers for new versions or variants of π -Cipher.





What is changed in explanation?

- New parts in the documentation of π -Cipher:
 - The security proof of π -Cipher
 - Explanation of how to use tweakable parameter N for wide blocks
 - Explanation of how to securely use incremental property of π Cipher
 - Rational why we consider π -Cipher to be STREAM OAE2+ design





• F. Abed, C. Forler and S. Lucks, "General Overview of the Authenticated Schemes for the First Round of the CAESAR Competition", Cryptology ePrint Archive, Report 2014/792





Construction Candidate	Design	Primitive	Features Security
			Parallelizable Enc/Dec Online Inverse-Free Incremental AD/AE Fixed AD reuse Intermediate Tags Security proof Nonce-MR Decryption-MR
Sponge-based π -cipher [57]	ARX,Duplex	n.n.	· -/





Functional characteristics

1. Parallelizable

– π -Cipher is parallelizable in both encryption and decryption phases

2. Online

- Encryption of the *i*-th input message block M_i depends only on the common state CIS, *i* and M_i .

3. Inverse free

- π -Cipher does not use π^{-1} of underlying permutation

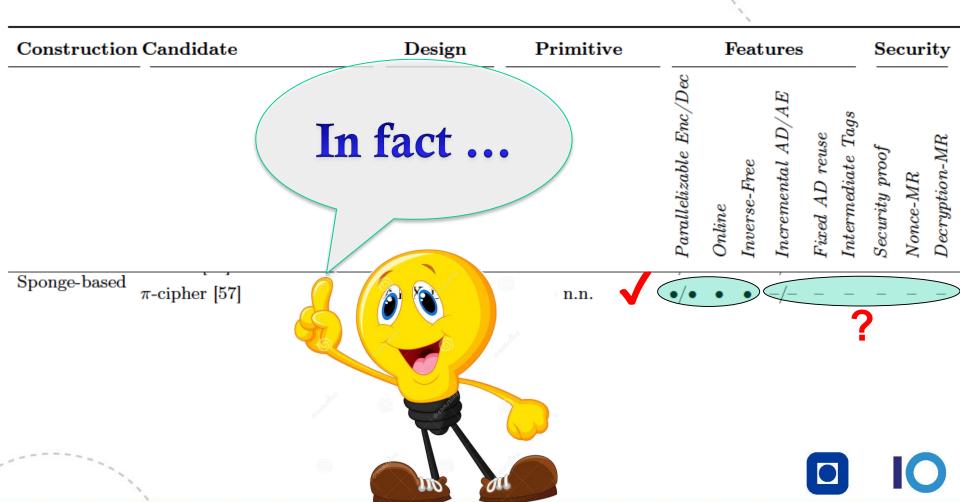


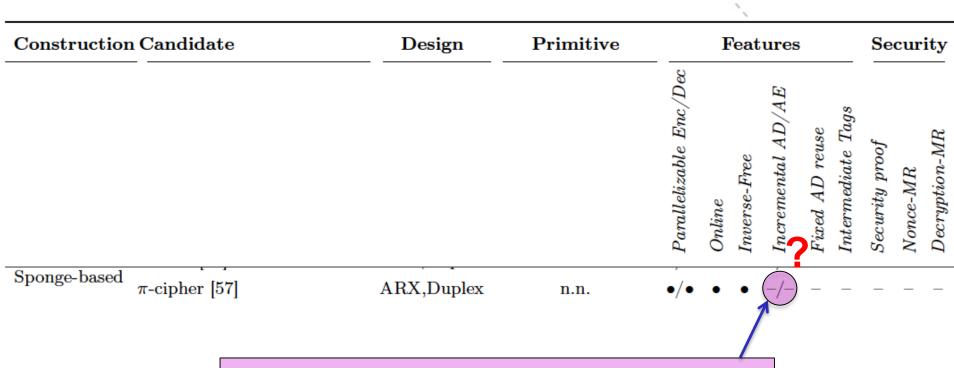


		*			
Design	Primitive	Features	Security		
		Parallelizable Enc/Dec Online Inverse-Free Incremental AD/AE Fixed AD reuse Intermediate Tags	Security proof Nonce-MR Decryption-MR		
ARX,Duplex	n.n.	•/• • • •			
	,		Parallelizable Enc/Dec Online Inverse-Free Incremental AD/AE Fixed AD reuse Intermediate Tags		









Yes, with additional metadata for the plaintext (overhead), in which case it is secure even with complete

NONCE = (PMN, SMN) REUSE





Incremental feature of π - Cipher

- Incremental schemes have advantage over standard one when longer messages are used (ex. encrypting data in rest)
- In π Cipher incrementality and NMR are achieved with additional **metadata** overhead of 64 bits per block
 - Update counter *UpdCtr* that records the history of updates for every data block





Incremental feature of π - Cipher

- Adding 64 bits of metadata to existing data blocks of π -Cipher (128, 256 and 512 bits) is unacceptable big overhead
- We need bigger blocks!
- How to do that?
 - Change the length of the state
 - In our case it is doable by changing the parameter N
 - Make π -Cipher a wide block cipher







π - Cipher as a wide block cipher

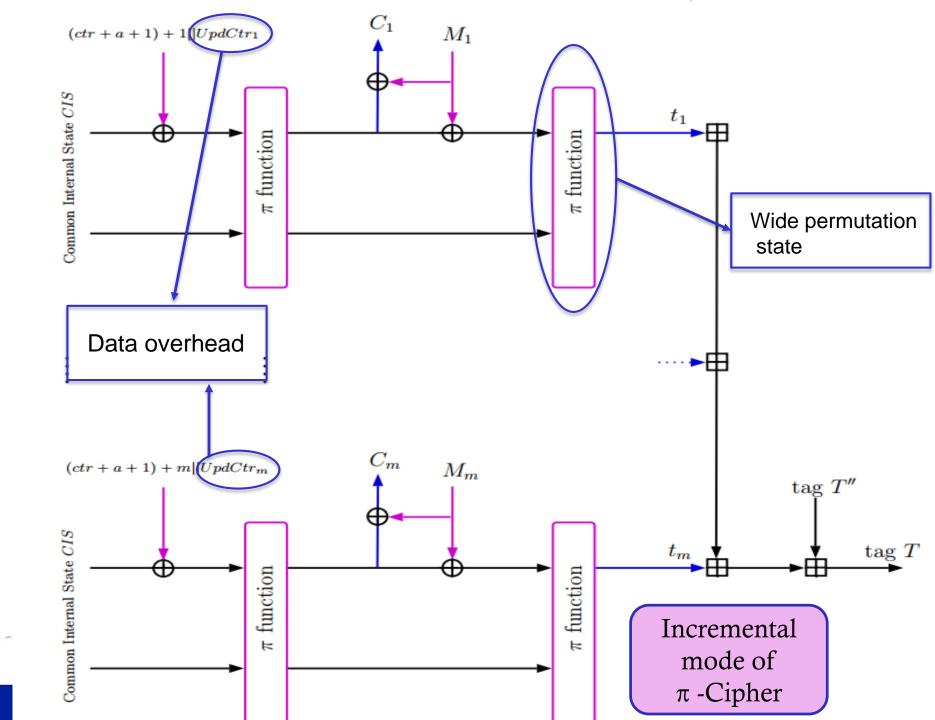
- Permutation state can be from 512B to 16KB
- Keeps the same security level even with 2 rounds

Table 4.1: Wide block characteristics of π 64-Cipher256

	klen (in bits)	PMN (in bits)	SMN (in bits)	Rate in Bytes	N	Tag T (in bits)	R
wide block of 512B	256	512	0	512	32	256	2
wide block of 2KB	256	512	0	2048	128	256	2
wide block of 4KB	256	512	0	4096	256	256	2
wide block of 8KB	256	512	0	8192	512	256	2
wide block of 16KB	256	512	0	16384	1024	256	2







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Sponge-based π -cipher [57]	ARX,Duplex	n.n.	•/• • • -/-
		s, when PM but SMN is	IN is reused different





Functional characteristics ...

4. Fixed Associated Data Reuse

- It is possible in the case where PMN is the same and SMN is different
- Allows considerable speed-up (Initialization phase and Processing the AD are skipped)
 - A typical use-case scenario would be a secure communication between devices in Internet Of Things. They run the initial setup procedure once where AD is used, and then they send only short encrypted messages.





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Sponge-based π -cipher [57]	ARX,Duplex	n.n.	•/• • -/

Yes, π- Cipher always computes intermediate tags for every block. It is just a matter of a mode of operation to use them. Additionally, with the wide-block feature, the relative overhead of having intermediate tags goes to zero.





Functional characteristics ...

5. By default, π - Cipher has no ciphertext expansion

- The length of the ciphertext is the same as the length of the message before padding + the length of the SMN
- But, as a mode of operation, it is possible to output intermediate tags for every block. Security of the cipher is not affected by publishing these intermediate tags.
- In order to reduce the relative overhead of having intermediate tags, the wide-block feature of π Cipher should be used.





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Security proof of π - Cipher

- Ensuring both privacy and authenticity for encrypted messages at the same time
 - Data privacy (IND-CPA)
 - Ciphertext integrity against forgery (INT-CTXT)
- π Cipher security proof is based on the proof for the sponge based authenticated ciphers given by P. Jovanovic, A. Luykx, B. Mennink in the ASIACRYPT 2014 paper "Beyond $2^{c/2}$ security in sponge based authenticated encryption modes"





IND-CPA

3.1.1 Privacy of π -Cipher

Theorem 2. Let $\Pi = (\mathcal{E}, \mathcal{D})$ be the proposed authenticated encryption scheme with an ideal permutation π which operates on b bits. Then,

$$\begin{aligned} \boldsymbol{A} \boldsymbol{d} \boldsymbol{v}_{II}^{priv}(q_p, q_{\varepsilon}, \lambda_{\varepsilon}) \leqslant & \frac{(q_p + \sigma_{\varepsilon} + \sigma_{\mathcal{D}})^2}{2^b} + \frac{q_{\mathcal{D}}}{2^r} + \frac{q_p + \sigma_{\varepsilon} + \sigma_{\mathcal{D}}}{2^k} + \frac{q_p r}{2^c} + \\ & \frac{q_{\varepsilon} a + q_{\mathcal{D}} a}{2^r} + \sqrt{\frac{8e\sigma_{\varepsilon} q_p}{2^b}} + \frac{\sigma_{\mathcal{D}}(q_p + \sigma_{\varepsilon} + \sigma_{\mathcal{D}}/2)}{2^c}, \end{aligned}$$

where σ_{ε} is defined in (3.1).





INT-CTXT

3.1.2 Authenticity of π -Cipher

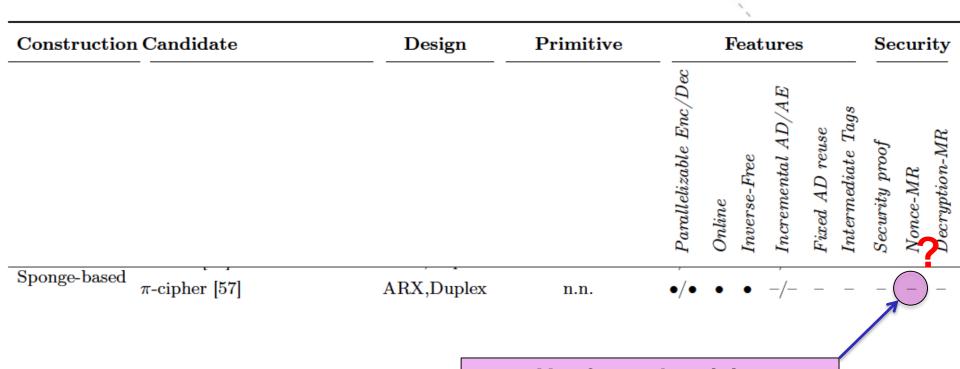
Theorem 3. Let $\Pi = (\mathcal{E}, \mathcal{D})$ be the proposed authenticated encryption scheme with an ideal permutation π which operates on b bits. Then,

$$\mathbf{Adv}_{II}^{auth}(q_p, q_{\varepsilon}, \lambda_{\varepsilon}, q_{\mathcal{D}}, \lambda_{\mathcal{D}}) \leqslant \frac{(q_p + \sigma_{\varepsilon} + \sigma_{\mathcal{D}})^2}{2^b} + \frac{q_{\mathcal{D}}}{2^{\tau}} + \frac{q_p + \sigma_{\varepsilon} + \sigma_{\mathcal{D}}}{2^k} + \frac{q_p r}{2^c} + \frac{q_p$$

where σ_{ε} and $\sigma_{\mathcal{D}}$ are defined in (3.1).







Yes for authenticity, Yes (conditional) for privacy (when SMN is not repeated)



Nonce Misuse Resistance

- Nonce = PMN (27 candidates)
- Nonce = (PMN, SMN) (2 candidates: *π* -Cipher and ICEPOLE-128)
- An intermediate level of nonce-misuse resistance is manifested when legitimate key holder reuses K, PMN and AD, but SMN is different



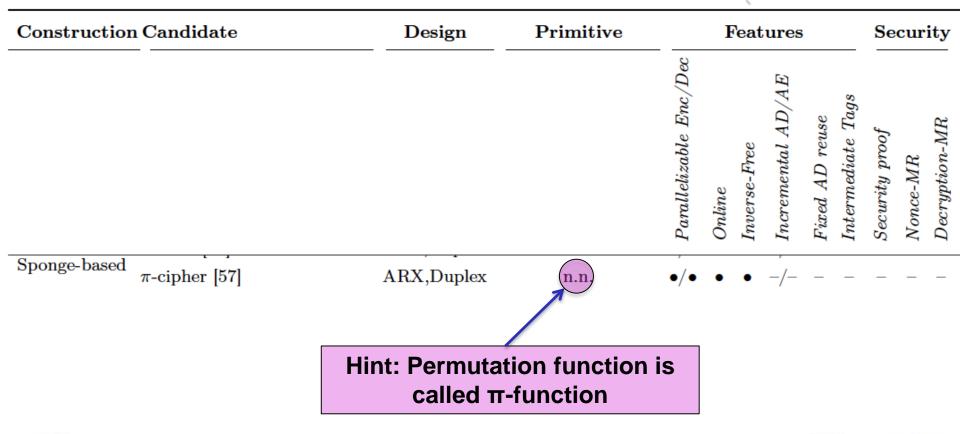


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Sponge-based π -cipher [57]	ARX,Duplex	n.n.	•/• • • -/

Yes, it is automatically achieved if it is implemented with intermediate tags, but still we need security proof (work in progress)







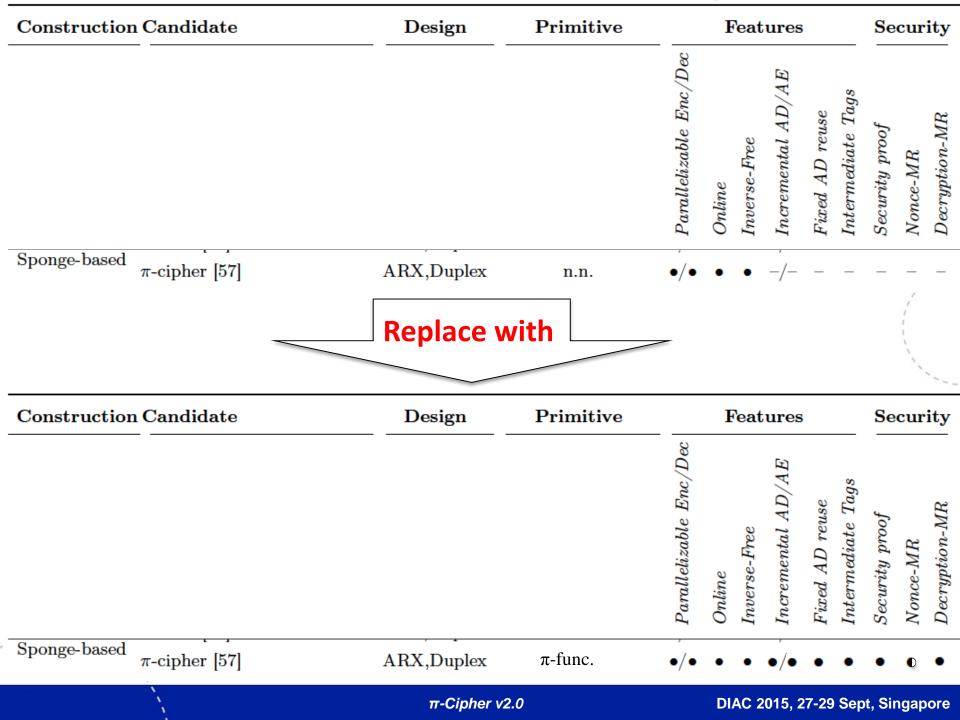


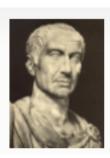


Construction Candidate	Design	Primitive	Features			Security				
			$Parallelizable\ Enc/Dec$	$Online \ Inverse-Free$	$Incremental\ AD/AE$	d AD reu		curity 1	once-MR	$Decryption ext{-}MR$
Sponge-based π -cipher [57]	ARX,Duplex	n.n.	•/•	• •	-/-	_	_	_	_	_







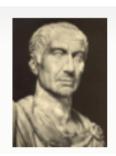


Authenticated Encryption Zoo

Name	Туре	Primitive	Parallel E/D	Online	Inverse- free	Security proof	Nonce- MR	Status
π-Cipher	Sponge	ARX	+/+	+	+	-	NONE	





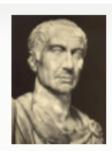


Authenticated Encryption Zoo

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π-Cipher	Sponge	ARX	+/+	+	+	- (NONE	



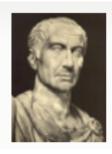




Authenticated Encryption Zoo

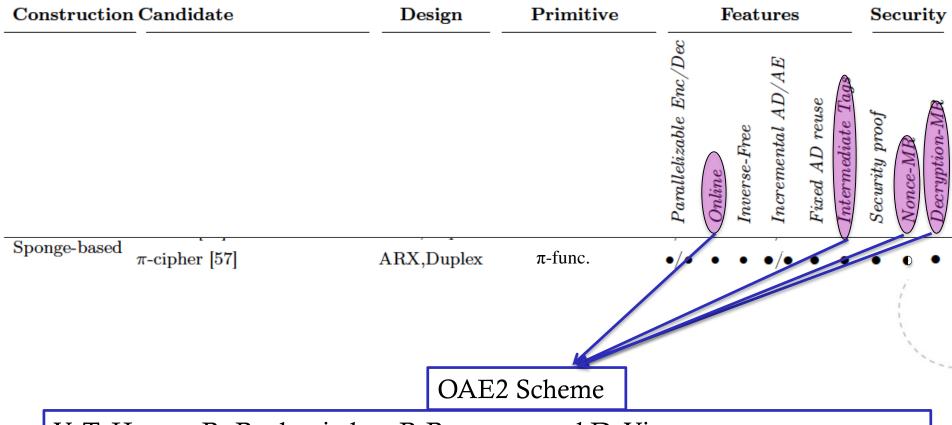
Name	Туре	Primitive	Parallel E/D	Online	Inverse- free	Security proof	Nonce- MR	Status
π-Cipher	Sponge	ARX	+/+	+	+	-	NONE	

Replace with



Authenticated Encryption Zoo

Name	Туре	Primitive	Parallel E/D	Online	Inverse- free	Security proof	Nonce- MR	Status
π-Cipher	Sponge	ARX	+/+	+	+	+	ON-SOME	



V. T. Hoang, R. Reyhanitabar, P. Rogaway, and D. Vizr. "Online Authenticated-Encryption and its Nonce-Reuse Misuse-Resistance", CRYPTO 2015. There, thay say: "Sponge duplex construction of Bertoni et al., resembles OAE2."





Construction Candidate	Design	Design Primitive	Features Security 2nd-roun	2nd-round
			Parallelizable Enc/Dec Online Inverse-Free Incremental AD/AE Fixed AD reuse Intermediate Tags Security proof Nonce-MR Decryption-MR	
Sponge-based π -cipher [57]	ARX,Duplex	n.n.		

 π - Cipher is based on sponge duplex construction of Bertoni et al., with additional cryptographic mechanisms that strengthen its robustness such as the features:

- tag second preimage resistance
- wide block tweakability
- incrementability
- use of SMN that guarantees confidentiality and integrity even when the K, AD and PMN are reused





Construction Candidate	Design	Primitive	Features Security 2nd-round
			Parallelizable Enc/Dec Online Inverse-Free Incremental AD/AE Fixed AD reuse Intermediate Tags Security proof Nonce-MR Decryption-MR
Sponge-based π -cipher [57]	ARX,Duplex	π -func.	

 π - Cipher is based on sponge duplex construction of Bertoni et al., with additional cryptographic mechanisms that strengthen its robustness su

- tag second
- wide block
- incrementa
- π Cipher is STREAM
 - OAE2+ cipher
- use of SMN that guarantees confidentiality and integrity even when the K, AD and PMN are reused





Efficiency

- Software speed of non SSE implementation of π 64-Cipher in v1.0 was around 11 cpb on Sandy Bridge. We expect v2.0 to be faster.
- Still we want to emphasize the incrementality feature of π -cipher by which it can outperform the speed of any non-incremental cipher even with 0.01 cpb





Efficiency

- Recent lightweight hardware implementation of π16-Cipher on Xilinx Virtex-7 platform XC7VX485T-2FFG1761 is:
 - 266 slices for the pi-function
 - 1114 slices for encryption engine without AD and SMN running at 347MHz
- Another lightweight implementation of $\pi 16$ -Cipher for AVR 8-bit MCU
 - 1.9 KB code size for encryption-authentication/decryption-verification part





Acknowledgements

- Gäetan Leurent and Thomas Fuhr
 - thanks for your detaild observation on the π -Cipher and pointing out the problem with padding
- Bart Mennink
 - thanks for your valuable and excellent advices in the process of proving the security of π -Cipher





Thank you for listening!

