

### CS 302 ASSIGNMENT 3

1(a) 1 digit:  $\begin{array}{l} \text{---} \\ - \text{---} \end{array} \} O(1)$

10 digits:  $\begin{array}{l} \text{---} \\ - \text{---} \\ \text{---} \end{array} \} O(1) \text{ } n \text{ times} = \boxed{O(n)}$   
 $\downarrow \downarrow \downarrow$   
 $O(1) \text{ } n \text{ times}$

1(b) Compute  $a^n$  where  $a$  is a 3-digit number.

$\begin{array}{l} \text{---} \\ \times \text{---} \\ \text{---} \end{array} \begin{array}{l} \text{3 digits} \\ \text{3 digits} \\ 3^2 \text{ multiplications} \end{array}$   
 $\begin{array}{l} \text{---} \times \\ \text{---} \times \end{array} \begin{array}{l} \text{each multiplication} = O(1) \\ 3^n \text{ multiplications} \end{array}$   
 Therefore  $\boxed{O(3^n)}$

1(c) Divide  $n$ -digit number with a 1-digit number  $> 1$

$\begin{array}{l} \text{---} \\ \text{---} \end{array} \leftarrow \begin{array}{l} \text{each digit in quotient must} \\ \text{try out digits 0-9 and is} \\ O(1) \text{ in runtime occurring} \\ n \text{ number of times} \end{array}$   
 $\uparrow$   
 1 digit  $\neq > 1$

Therefore  $\boxed{O(n)}$

2) Show for any constant  $K > 0$ ,  $O(\log^K n) = o(n)$

$$\lim_{n \rightarrow \infty} \frac{n}{\log^K(n)} = \lim_{n \rightarrow \infty} \frac{1}{K \cdot \log^{K-1}(n)}$$

$$= \lim_{n \rightarrow \infty} \frac{n}{K \cdot \log^{K-1}(n)}$$

$$= \frac{1}{K} \lim_{n \rightarrow \infty} \frac{n}{\log^{K-1}(n)} = \frac{\infty}{\text{undefined}} = \infty$$

$f(x) = n$  is faster in growth rate.

3) Prove or disprove  $2^n = \Theta(3^n)$

$$\lim_{n \rightarrow \infty} \frac{2^n}{3^n} = \lim_{n \rightarrow \infty} \frac{2^n \ln(2)}{3^n \ln(3)}$$

$$= \frac{\ln(2)}{\ln(3)} \cdot \lim_{n \rightarrow \infty} \left(\frac{2}{3}\right)^n = 0$$

$2^n$  and  $3^n$  do not have the same asymptotic growth rate.  $3^n$  has a faster growth rate.

$$4) \log_{10} n = 4 \log_2 n < \frac{n}{2^{64}} = 4000n + 22 < n \log_2 n = 2n^{1.5} - n \log n \\ < \frac{n^2}{2} + 20n - 4 < 1.0001^n = 2^{64} = n\sqrt{n}$$

5) INSERTION SORT 5, 3, 10, 2, 9, 1

3 5

5 ~~3~~ 10 2 9 1

$i \leftarrow 1$

key  $\leftarrow 3$

$j \leftarrow 0 - 1$

↓ (next iteration)

3 5 10 2 9 1

$i \leftarrow 2$

key  $\leftarrow 10$

$j \leftarrow 1$  (next iteration)

2 3 5 10

~~3~~ ~~5~~ 10 ~~2~~ 9 1

$i \leftarrow 3$

key  $\leftarrow 2$

$j \leftarrow 2 - 1 \rightarrow$  (next line iteration)

9 10

2 3 5 10 ~~9~~ 1

$i \leftarrow 4$

key  $\leftarrow 9$

$j \leftarrow 3$  (next iteration)

1 2 3 5 9 10

2 ~~3~~ ~~5~~ ~~9~~ 10 ~~1~~

$i \leftarrow 5$

key  $\leftarrow 1$

$j \leftarrow 4 - 1 \rightarrow$  (next iteration)

1 2 3 5 9 10

$i \leftarrow 6$  exit insertion sort.



# 6) MERGE SORT

1	30	29	31	11	25	32	21	35	23	13	2	4	6	5	10
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

1	30	29	31	11	25	32	21	35	23	13	2	4	6	5	10
---	----	----	----	----	----	----	----	----	----	----	---	---	---	---	----

1	30	29	31	11	25	32	21	35	23	13	2	4	6	5	10
---	----	----	----	----	----	----	----	----	----	----	---	---	---	---	----

1	30	29	31	11	25	32	21	35	23	13	2	4	6	5	10
---	----	----	----	----	----	----	----	----	----	----	---	---	---	---	----

1	29	30	31	32	21	23	35	2	13	4	5	6	10
---	----	----	----	----	----	----	----	---	----	---	---	---	----

21	32	23	35	2	13
----	----	----	----	---	----

11	21	25	32	2	13	25	35
----	----	----	----	---	----	----	----

1	11	21	25	29	30	31	32	2	4	5	6	10	13	25	35
---	----	----	----	----	----	----	----	---	---	---	---	----	----	----	----

1	2	4	5	6	10	11	13	21	25	25	29	30	31	32	35
---	---	---	---	---	----	----	----	----	----	----	----	----	----	----	----

# 7) QUICKSORT

P L R R R

10 23 7 6 2 5 13 8 19 11

P L R R R

10 8 7 6 2 5 13 23 19 11

P L R

5 8 7 6 2 10 13 23 19 11

P L R R

5 2 7 6 8

R L

P L R

2 5 7 6 8 10 13 23 19 11

R L

P R

2 5 6 7 8 10 13 23 19 11

P L R

2 5 6 7 8 10 13 23 19 11

R R

P L R

13 11 19 23

R R

P L

2 5 6 7 8 10 11 13 19 23

2 5 6 7 8 10 11 13 19 23