算法竞赛个人模板

Cu_OH_2

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目录

1 常用

1.1 基础框架

```
#include<bits/stdc++.h>
using namespace std;
using ll = long long;

void solve()
{
    return;
}

int main()
{
    ios::sync_with_stdio(0);
    cin.tie(0);
    cout.tie(0);
    int T = 1;
    //cin >> T;
    while (T--) solve();
    return 0;
}
```

1.2 调试技巧

```
#define debug(x) cout << #x << " = " << x << endl
freopen("A.in", "r", stdin);
freopen("A.out", "w", stdout);
```

1.3 注意事项

```
//1.爆long long了吗？
//2.数组首尾边界初始化了吗？
//3.测试组间数据清空重置了吗？
//4.交互题用endl了吗？
//5.clear()重置数据了吗？
//6.size()参与减法溢出了吗？
//7.for(j)循环写成++i了吗？
```

2 动态规划

2.1 单调队列优化多重背包

```
/*****************************************************************
* 时间复杂度: O(nm)
* 说明: dp[j]只有可能从dp[j-k*w[i]]转移来
*****************************************************************/
const int N = 100005;
const int M = 40005;

ll n, m; //种数、容积
ll v[N], w[N], k[N]; //价值、体积、数量
ll dp[M]; //使用i容积的最大价值

struct Node
{
    ll key, id;
};

void solve()
{
    cin >> n >> m;
    for (int i = 1; i <= n; ++i) cin >> v[i] >> w[i] >> k[i];
    for (int i = 1; i <= n; ++i)
    {
        vector<deque<Node>> dq(w[i]);
        auto key = [&](int j) { return dp[j] - j / w[i] * v[i]; }; //
            dp[j]在比较基准下的指标
        auto join = [&](int j) //dp[j]入队
        {
            auto& q = dq[j % w[i]];
            while (q.size() && key(j) >= q.back().key) q.pop_back();
            q.push_back({ key(j),j });
            return;
        };
        for (int j = m; j >= max(0ll, m - k[i] * w[i]); --j) join(j);
        for (int j = m; j >= w[i]; --j)
        {
            auto& q = dq[j % w[i]];
            while (q.size() && q.front().id >= j) q.pop_front();
            if (j - k[i] * w[i] >= 0) join(j - k[i] * w[i]);
            dp[j] = max(dp[j], q.front().key + j / w[i] * v[i]);
        }
    }
    ll ans = 0;
    for (int i = 0; i <= m; ++i) ans = max(ans, dp[i]);
    cout << ans << '\n';
    return;
}
```

2.2 二进制分组优化多重背包

```
/*****************************************************************
* 时间复杂度: O(nmlogk)
* 说明: 二进制分组优化多重背包, 可bitset优化
*****************************************************************/
const int N = 100005;
const int M = 40005;

struct Item
{
    ll v, w; //价值、体积
};

ll n, m; //种数、容积
ll dp[M]; //使用i容积的最大价值

void solve()
{
    cin >> n >> m;
    vector<Item> items;
    ll x, y, z;
    for (int i = 1; i <= n; ++i)
    {
        ll b = 1;
        cin >> x >> y >> z;
        while (z > b)
        {
            z -= b;
            items.push_back({ x * b, y * b });
            b <<= 1;
        }
        items.push_back({ x * z, y * z });
    }
    for (auto e : items)
    {
        for (int i = m; i >= e.w; --i)
        {
            dp[i] = max(dp[i], dp[i - e.w] + e.v);
        }
    }
    ll ans = 0;
    for (int i = 0; i <= m; ++i) ans = max(ans, dp[i]);
    cout << ans << '\n';
    return;
}
```

2.3 动态 DP

```
/*****************************************************************
* 时间复杂度: O(qlogn)
* 说明: 将dp转移方程表示为矩阵乘法, 用线段树维护矩阵, 实现带修改dp。
*****************************************************************/
//CF1814E
const int N = 200005;
const ll INFLL = 0x3f3f3f3f3f3f3f3f;

struct SegTree
{
    struct Node
    {
        int lef, rig;
        array<array<ll, 2>, 2> mat;
    };
    vector<Node> tree;

    inline int ls(int src) { return src * 2; }
    inline int rs(int src) { return src * 2 + 1; }
```

```cpp
    inline Node& ln(int src) { return tree[ls(src)]; }
    inline Node& rn(int src) { return tree[rs(src)]; }

    inline void update(int src)
    {
        for (int i = 0; i < 2; ++i)
        {
            for (int j = 0; j < 2; ++j)
            {
                auto v1 = ln(src).mat[i][1] + rn(src).mat[1][j];
                auto v2 = ln(src).mat[i][0] + rn(src).mat[1][j];
                auto v3 = ln(src).mat[i][1] + rn(src).mat[0][j];
                tree[src].mat[i][j] = min({ v1, v2, v3 });
            }
        }
        return;
    }

    inline void calc(int src, ll val)
    {
        tree[src].mat[1][1] = val;
        tree[src].mat[0][0] = 0;
        tree[src].mat[0][1] = tree[src].mat[1][0] = INFLL;
        return;
    }

    SegTree(int x) { tree.resize(x * 4 + 1); }

    void build(int src, int lef, int rig, ll arr[])
    {
        tree[src].lef = lef;
        tree[src].rig = rig;
        if (lef == rig)
        {
            calc(src, arr[lef]);
            return;
        }
        int mid = lef + rig >> 1;
        build(ls(src), lef, mid, arr);
        build(rs(src), mid + 1, rig, arr);
        update(src);
        return;
    }

    void modify(int src, int pos, ll val)
    {
        if (tree[src].lef == tree[src].rig)
        {
            calc(src, val);
            return;
        }
        int mid = tree[src].lef + tree[src].rig >> 1;
        if (pos <= mid) modify(ls(src), pos, val);
        else modify(rs(src), pos, val);
        update(src);
        return;
    }

    ll query() { return tree[1].mat[1][1] * 2; }
};

int n, q, k;
ll a[N], x;

void solve()
{
    cin >> n;
    for (int i = 1; i <= n - 1; ++i) cin >> a[i];
    SegTree sgt(n - 1);
    sgt.build(1, 1, n - 1, a);
    cin >> q;
    for (int i = 1; i <= q; ++i)
    {
        cin >> k >> x;
        sgt.modify(1, k, x);
        cout << sgt.query() << '\n';
    }
    return;
}
```

3 字符串

3.1 KMP 算法

```cpp
/**********************************************************************
* 时间复杂度: O(n)
* 说明:
* 1.nxt[i]表示t[i](下标从0开始)失配时下一次匹配的位置
* 2.nxt[n]在匹配中无必要作用,但构成前缀数组
* 3.前缀数组pi[i]=nxt[i+1]+1,代表前缀t[0,i]的最长前后缀长度
**********************************************************************/
struct KMP
{
    string t;
    vector<int> nxt;

    KMP() {}
    KMP(const string& str) { init(str); }

    void init(const string& str)
    {
        t = str;
        nxt.resize(t.size() + 1);
        nxt[0] = -1;
        for (int i = 1; i <= t.size(); ++i)
        {
            int now = nxt[i - 1];
            while (now != -1 && t[i - 1] != t[now]) now = nxt[now];
            nxt[i] = now + 1;
        }
        return;
    }

    int first(const string& s)
    {
        int ps = 0, pt = 0;
        while (ps < s.size())
        {
            while (pt != -1 && s[ps] != t[pt]) pt = nxt[pt];
            ps++, pt++;
            if (pt == t.size()) return ps - t.size();
        }
        return -1;
    }

    vector<int> every(const string& s)
    {
        vector<int> v;
        int ps = 0, pt = 0;
        while (ps < s.size())
        {
            while (pt != -1 && s[ps] != t[pt]) pt = nxt[pt];
            ps++, pt++;
            if (pt == t.size())
            {
                v.push_back(ps - t.size());
                pt = nxt[pt];
            }
        }
        return v;
    }
};
```

3.2 扩展 KMP 算法

```cpp
/**********************************************************************
* 时间复杂度: O(n)
* 说明: Z函数代表后缀与母串的最长公共前缀
**********************************************************************/
struct ExKMP
{
    string t;
    vector<int> z;

    ExKMP(const string& str)
    {
        t = str;
        z.resize(t.size());
        z[0] = t.size();
        int l = 0, r = -1;
        for (int i = 1; i < t.size(); ++i)
        {
            if (i <= r && z[i - l] < r - i + 1) z[i] = z[i - l];
            else
            {
                z[i] = max(0, r - i + 1);
                while (i + z[i] < t.size() && t[z[i]] == t[i + z[i]]) z
                    [i]++;
            }
            if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
        }
```

```
    }

    vector<int> ext(const string& s)
    {
        vector<int> res(s.size());
        int l = 0, r = -1;
        for (int i = 0; i < s.size(); ++i)
        {
            if (i <= r && z[i - l] < r - i + 1) res[i] = z[i - l];
            else
            {
                res[i] = max(0, r - i + 1);
                while (i + res[i] < s.size() && res[i] < t.size() && t[
                    res[i]] == s[i + res[i]]) res[i]++;
            }
            if (i + res[i] - 1 > r) l = i, r = i + res[i] - 1;
        }
        return res;
    }
};
```

3.3 字典树

```
/*********************************************************************
* 时间复杂度: O(sigma(n))
* 说明:
* 1.字典树也即前缀树，每个结点代表一个前缀
* 2.字母表变化只需要修改映射函数F()
* 3.若需要遍历trie树可以用out数组记录出边降低复杂度
*********************************************************************/
struct Trie
{
    const int ALPSZ = 26;
    vector<vector<int>> trie;
    vector<int> tag;
    //vector<vector<int>> out;

    inline int F(char c) { return c - 'a'; }

    Trie() { init(); }

    void init()
    {
        create();
        return;
    }
    int create()
    {
        trie.push_back(vector<int>(ALPSZ));
        tag.push_back(0);
        //out.push_back(vector<int>());
        return trie.size() - 1;
    }
    void insert(const string& t)
    {
        int now = 0;
        for (auto e : t)
        {
            if (!trie[now][F(e)])
            {
                int newNode = create();
                //out[now].push_back(F(e));
                trie[now][F(e)] = newNode;
            }
            now = trie[now][F(e)];
            tag[now]++;
        }
        return;
    }
    int count(const string& pre)
    {
        int now = 0;
        for (auto e : pre)
        {
            now = trie[now][F(e)];
            if (now == 0) return 0;
        }
        return tag[now];
    }
};
```

3.4 AC 自动机

```
/*********************************************************************
* 时间复杂度: O(alpsz*sigma(len(t))+len(s))
* 说明:
* 1.本模板以小写英文字母为字母表举例，修改字母表可以通过修改F()函数完成。
* 2.Trie图优化: 建立fail指针时，fail指针指向的结点有可能依然失配，需要多
* 次跳转才能到达匹配结点。可以将所有结点的空指针补全，置为该结点的跳转
* 终点。此时根据BFS序，在计算tr[x][i]的fail指针时，fail[x]一定已遍历
* 过，且tr[fail[x]][i]一定存在，要么为fail[x]接收i的后继状态，要么为
* tr[x][i]的跳转终点。无论哪种情况，fail[tr[x][i]]都可以直接置为
* tr[fail[x]][i]。
* 3.last优化: 多模式匹配过程中，对于文本串的每个前缀s'，沿fail指针路径寻
* 找为s'后缀的模式串，途中可能经过无贡献的模式串真前缀结点; last优化使
* 得跳转时跳过真前缀结点直接到达上方第一个模式串结点。last数组可以完全
* 替代fail数组。
* 4.树上差分优化: 统计每种模式串出现次数时，每匹配到一个模式串都要向上跳
* 转一次，这个过程相当于区间加一，可以用更新差分数组代替，最后再计算前
* 缀和即可。
* 5.注意: 统计出现的模式串种类数时会将标记清空
*********************************************************************/
struct ACAM
{
    vector<vector<int>> trie; //trie树指针
    vector<int> tag; //标记数组
    vector<int> fail; //失配函数
    vector<int> last; //跳转路径上一个模式串结点
    vector<int> cnt; //计数器
    const int ALPSZ = 26; //字母表大小
    int ord; //结点个数

    inline int F(char c) { return c - 'a'; }

    ACAM() { init(); }

    void init()
    {
        ord = -1;
        newNode();
    }
    int newNode()
    {
        trie.push_back(vector<int>(ALPSZ));
        tag.push_back(0);
        return ++ord;
    }
    void addPat(const string& t)
    {
        int now = 0;
        for (auto e : t)
        {
            if (!trie[now][F(e)]) trie[now][F(e)] = newNode();
            now = trie[now][F(e)];
        }
        tag[now]++;
        return;
    }
    void buildAM()
    {
        fail.resize(ord + 1);
        last.resize(ord + 1);
        cnt.resize(ord + 1);
        queue<int> q;
        for (int i = 0; i < 26; ++i)
        {
            //第一层结点的fail指针都指向0，不需要处理
            if (trie[0][i]) q.push(trie[0][i]);
        }
        while (q.size())
        {
            int now = q.front();
            q.pop();
            for (int i = 0; i < 26; ++i)
            {
                int son = trie[now][i];
                if (son)
                {
                    fail[son] = trie[fail[now]][i];
                    if (tag[fail[son]]) last[son] = fail[son];
                    else last[son] = last[fail[son]];
                    q.push(trie[now][i]);
                }
                else trie[now][i] = trie[fail[now]][i];
            }
        }
        return;
    }
    int count(const string& s) //统计出现的模式串种类数
    {
        int now = 0, ans = 0;
        for (auto e : s)
```

```cpp
        {
            now = trie[now][F(e)];
            int p = now;
            while (p)
            {
                ans += tag[p];
                tag[p] = 0;
                p = last[p];
            }
        }
        return ans;
    }
};
```

3.5 后缀自动机

```cpp
/******************************************************************
* 时间复杂度：O(n*ALPSZ)
* 说明：字符集较大可以将next换成map<char,int>
******************************************************************/
struct SAM
{
    struct State
    {
        int maxlen; //结点代表的最长子串长度
        int link; //后缀链接，连向不在该点中的最长后缀
        vector<int> next;
        State(): maxlen(0), link(-1) { next.resize(26); }
    };
    vector<State> node;
    vector<ll> cnt; //子串出现次数（endpos集合大小）
    int now; //接收上一个字符到达的结点
    int size; //当前结点个数

    inline int F(char c) { return c - 'a'; }

    SAM(int x)
    {
        node.resize(x * 2 + 5);
        cnt.resize(x * 2 + 5);
        now = 0; //从根节点开始转移
        size = 1; //建立一个代表空串的根节点
    }

    void extend(char c)
    {
        int nid = size++;
        cnt[nid] = 1;
        node[nid].maxlen = node[now].maxlen + 1;
        int p = now;
        while (p != -1 && node[p].next[F(c)] == 0)
        {
            node[p].next[F(c)] = nid;
            p = node[p].link;
        }
        if (p == -1) node[nid].link = 0; //连向根结点
        else
        {
            int ori = node[p].next[F(c)];
            if (node[p].maxlen + 1 == node[ori].maxlen) node[nid].link
                = ori;
            else
            {
                //将ori结点的一部分拆出来分成新结点split
                int split = size++;
                node[split].maxlen = node[p].maxlen + 1;
                node[split].link = node[ori].link;
                node[split].next = node[ori].next;
                while (p != -1 && node[p].next[F(c)] == ori)
                {
                    node[p].next[F(c)] = split;
                    p = node[p].link;
                }
                node[ori].link = node[nid].link = split;
            }
        }
        now = nid;
        return;
    }

    void build(const string& s)
    {
        for (auto e : s) extend(e);
        return;
    }

    void DFS(int x, vector<vector<int>>& son)
    {
        for (auto e : son[x])
        {
            DFS(e, son);
            cnt[x] += cnt[e]; //link树上父节点endpos为所有子结点endpos之
                和
        }
        return;
    }

    void count() //计算endpos大小
    {
        //建立link树
        vector<vector<int>> son(size);
        for (int i = 1; i < size; ++i) son[node[i].link].push_back(i)
            ;

        //在link树上dfs
        DFS(0, son);
        return;
    }

    ll substr() //本质不同子串个数
    {
        ll res = 0;
        for (int i = 1; i < size; ++i)
        {
            res += node[i].maxlen - node[node[i].link].maxlen;
        }
        return res;
    }
};
```

3.6 回文自动机

```cpp
/******************************************************************
* 时间复杂度：O(n)
* 说明：
* 1.每个结点代表一个本质不同回文串。link链：多字串->单字符->偶根->奇根。
* 2.每个本质不同回文子串次数：最后由母串向子串传递。
    每个前缀的后缀回文子串个数：新建时由最长回文后缀向新串传递。
******************************************************************/
struct PAM
{
    struct State
    {
        int len; //长度
        int link; //最长回文后缀结点
        vector<int> next; //两边加上某字符时对应的结点
        State() { next.resize(26); }
        State(int x, int y): len(x), link(y) { next.resize(26); }
    };
    vector<State> node;
    vector<ll> cnt; //本质不同回文串出现次数
    int now; //接收上一个字符到达的结点
    int size; //当前结点个数

    inline int F(char c) { return c - 'a'; }

    PAM(int x)
    {
        node.resize(x + 3);
        node[0] = State(-1, 0); //奇根，link无意义
        node[1] = State(0, 0); //偶根，link指向奇根
        cnt.resize(x + 3);
        now = 0; //第一个字符由奇根转移
        size = 2;
    }

    void build(const string& s)
    {
        auto find = [&](int x, int p) //寻找x后缀中左方为s[p]的最长回文
            子串
        {
            while (p - node[x].len - 1 < 0 || s[p] != s[p - node[x].
                len - 1]) x = node[x].link;
            return x;
        };
        for (int i = 0; i < s.size(); ++i)
        {
            now = find(now, i);
            if (!node[now].next[F(s[i])]) //对应结点不存在则需要新建
            {
                int nid = size++;
                node[nid].len = node[now].len + 2; //新建状态结点
```

```
            node[nid].link = 1; //若now=0，对应结点为单字符，指向偶根
            if (now) node[nid].link = node[find(node[now].link, i)].next[F(s[i])]; //否则指向再前一个结点的扩展
            node[now].next[F(s[i])] = nid;
        }
        now = node[now].next[F(s[i])];
        cnt[now]++;
    }
    for (int i = size - 1; i >= 2; --i) cnt[node[i].link] += cnt[i]; //数量由母串向子串传递
    return;
  }
};
```

3.7 Manacher 算法

```
/*********************************************************************
* 时间复杂度：O(n)
* 说明：用n+1个分隔符将字符串分隔可以将奇偶回文子串过程统一处理
*********************************************************************/
struct Manacher
{
    vector<int> odd, even; //以[i]或[i,i+1]为中心的最长回文串半径
    void work(const string& s)
    {
        odd.resize(s.size());
        even.resize(s.size() - 1);
        int lef = 0, rig = -1, r;
        for (int i = 0; i < s.size(); ++i)
        {
            if (i > rig) r = 1;
            else r = min(odd[lef + rig - i], rig - i) + 1; //利用对称位置答案
            while (i - r >= 0 && i + r < s.size() && s[i - r] == s[i + r]) r++; //暴力扩展
            odd[i] = --r; //记录答案
            if (i + r > rig) lef = i - r, rig = i + r; //扩展lef,rig范围
        }
        lef = 0, rig = -1;
        for (int i = 0; i + 1 < s.size(); ++i)
        {
            if (i + 1 > rig) r = 1;
            else r = min(even[lef + rig - i - 1], rig - i) + 1;
            while (i + 1 - r >= 0 && i + r < s.size() && s[i + 1 - r] == s[i + r]) r++;
            even[i] = --r;
            if (i + r > rig) lef = i + 1 - r, rig = i + r;
        }
        return;
    }
};
```

3.8 最小表示法

```
/*********************************************************************
* 时间复杂度：O(n)
* 说明：求循环rotate得到的n种表示中字典序最小的一种
*********************************************************************/
const int N = 300005;

int n, a[N];

void solve()
{
    cin >> n;
    for (int i = 1; i <= n; ++i) cin >> a[i];
    auto norm = [](int x) { return (x - 1) % n + 1; };
    int p1 = 1, p2 = 2, len = 1;
    while (p1 <= n && p2 <= n & len <= n)
    {
        if (a[norm(p1 + len - 1)] == a[norm(p2 + len - 1)]) len++;
        else if (a[norm(p1 + len - 1)] < a[norm(p2 + len - 1)]) p2 += len, len = 1;
        else p1 += len, len = 1;
        if (p1 == p2) p1++;
    }
    int ans = min(p1, p2);
    return;
}
```

3.9 字符串哈希

```
/*********************************************************************
* 时间复杂度：O(n)
* 说明：字符串传入前必须处理为下标从1开始的模式！
*********************************************************************/
const int M1 = 998244389;
const int M2 = 998244391;
const int B1 = 31;
const int B2 = 29;
const int N = 1000005;

struct Base
{
    array<ll, N> pow{};
    Base(int base, int mod)
    {
        pow[0] = 1;
        for (int i = 1; i <= N - 1; ++i)
        {
            pow[i] = pow[i - 1] * base % mod;
        }
    }
    const ll operator[](int idx) const { return pow[idx]; }
} p1(B1, M1), p2(B2, M2);

struct Hash
{
    vector<ll> hash1, hash2;
    void build(const string& s)
    {
        int n = s.size() - 1;
        hash1.resize(n + 1);
        hash2.resize(n + 1);
        for (int i = 1; i <= n; ++i)
        {
            hash1[i] = (hash1[i - 1] * B1 % M1 + s[i] - 'a' + 1) % M1;
            hash2[i] = (hash2[i - 1] * B2 % M2 + s[i] - 'a' + 1) % M2;
        }
        return;
    }
    ll merge(ll x, ll y) { return x << 31 | y; }
    ll calc(int lef, int rig)
    {
        ll res1 = (hash1[rig] - hash1[lef - 1] * p1[rig - lef + 1] % M1 + M1) % M1;
        ll res2 = (hash2[rig] - hash2[lef - 1] * p2[rig - lef + 1] % M2 + M2) % M2;
        return merge(res1, res2);
    }
};
```

4 数学

4.1 快速幂

```
/*********************************************************************
* 时间复杂度：O(sqrt(n))
* 说明：
* 1.特殊情况下需要对res和a的初值进行取模，注意p不可取模
* 2.利用费马小定理求乘法逆元时注意仅当mod为质数时有效
* 3.若p较大且mod为质数可以将p对mod-1取模
*********************************************************************/
ll qpow(ll a, ll p, ll mod)
{
    ll res = 1;
    while (p)
    {
        if (p & 1) res = res * a % mod;
        a = a * a % mod;
        p >>= 1
    }
    return res;
}

ll inv(ll a, ll mod)
{
    return qpow(a, mod - 2, mod);
}
```

4.2 矩阵快速幂

```cpp
/*******************************************************************
* 时间复杂度：O(n^3logp)
* 说明：已知递推式可以表示为矩阵乘法，快速求数列第i项
*******************************************************************/
const int MOD = 1e9 + 7;

struct Square
{
   int n;
   vector<vector<ll>> a;
   Square(int n): n(n) { a.resize(n, vector<ll>(n)); }
   void unit()
   {
      for (int i = 0; i < n; ++i)
      {
         a[i][i] = 1;
      }
      return;
   }
};

Square mult(const Square& lhs, const Square& rhs)
{
   assert(lhs.n == rhs.n);
   int n = lhs.n;
   Square res(n);
   for (int i = 0; i < n; ++i)
   {
      for (int j = 0; j < n; ++j)
      {
         for (int k = 0; k < n; ++k)
         {
            res.a[i][j] += lhs.a[i][k] * rhs.a[k][j] % MOD;
            res.a[i][j] %= MOD;
         }
      }
   }
   return res;
}

Square qpow(Square a, ll p)
{
   int n = a.n;
   Square res(n);
   res.unit();
   while (p)
   {
      if (p & 1) res = mult(res, a);
      a = mult(a, a);
      p >>= 1;
   }
   return res;
}
```

4.3 排列奇偶性

```cpp
/*******************************************************************
* 时间复杂度：O(n)
* 说明：
* 1.顺序排列为偶排列
* 2.交换任意两个数，排列奇偶性改变
* 3.排列奇偶性等于逆序对数奇偶性
* 4.求环的个数可以O(n)求得排列奇偶性
*******************************************************************/
void solve()
{
   cin >> n;
   for (int i = 1; i <= n; ++i) cin >> a[i];
   bool inv = n & 1;
   vector<bool> vis(n + 1);
   for (int i = 1; i <= n; ++i)
   {
      if (vis[i]) continue;
      int cur = i;
      while (!vis[cur])
      {
         vis[cur] = 1;
         cur = a[cur];
      }
      inv ^= 1;
   }
   return;
}
```

4.4 组合数递推

```cpp
/*******************************************************************
* 时间复杂度：O(nm)
* 说明：递推预处理组合数
*******************************************************************/
struct Comb
{
   vector<vector<ll>> c;
   Comb(int x, int y, int mod)
   {
      c.resize(x + 1, vector<ll>(y + 1));
      for (int i = 0; i <= x; ++i) c[i][0] = 1;
      for (int i = 1; i <= x; ++i)
      {
         for (int j = 1; j <= i; ++j) c[i][j] = (c[i - 1][j - 1] +
              c[i - 1][j]) % mod;
      }
   }
   ll val(int x, int y)
   {
      if (x < 0 || y < 0) return 0;
      else return c[x][y];
   }
};
```

4.5 线性基

```cpp
/*******************************************************************
* 时间复杂度：插入O(b)/求最大异或和O(b)
* 说明：
* 1.可以求子序列最大异或和
* 2.v中非零元素表示一组线性基
* 3.线性基大小表征线性空间维数
* 4.线性基中没有异或和为0的子集
* 5.线性基中各数二进制最高位不同
*******************************************************************/
const int N = 55;
const int B = 50;

template<int bit>
struct LinearBasis
{
   vector<ll> v;
   LinearBasis() { v.resize(bit); }
   void insert(ll x)
   {
      for (int i = bit - 1; i >= 0; --i)
      {
         if (x >> i & 1ll)
         {
            if (v[i]) x ^= v[i];
            else
            {
               v[i] = x;
               break;
            }
         }
      }
      return;
   }
   ll qmax()
   {
      ll res = 0;
      for (int i = bit - 1; i >= 0; --i)
      {
         if ((res ^ v[i]) > res) res ^= v[i];
      }
      return res;
   }
   void merge(const LinearBasis<bit>& b)
   {
      for (auto e : b.v) insert(e);
      return;
   }
};
```

4.6 高精度

```cpp
/*******************************************************************
* 时间复杂度：O(n)/O(n^2)
* 说明：待完善，注意复杂度
```

```cpp
*******************************************************************/
const int N = 5005;

struct Large
{
   array<ll, N> ar{};
   int len = 0;
   Large() {}
   Large(ll x)
   {
      int p = 0;
      while (x)
      {
         ar[p++] = x % 10;
         x /= 10;
      }
      updateLen();
   }
   Large(const string& s)
   {
      for (int i = 0; i < s.size(); ++i)
      {
         ar[i] = s[s.size() - 1 - i] - '0';
      }
      updateLen();
   }

   void updateLen()
   {
      len = ar.size();
      for (int i = ar.size() - 1; i >= 0; --i)
      {
         if (ar[i]) break;
         len = i;
      }
      return;
   }

   Large& operator=(const Large& rhs)
   {
      for (int i = 0; i < ar.size(); ++i) ar[i] = rhs.ar[i];
      updateLen();
      return *this;
   }

   Large operator+(const Large& rhs) const
   {
      Large res;
      for (int i = 0; i < ar.size(); ++i) res.ar[i] = ar[i] + rhs.
          ar[i];
      for (int i = 0; i < ar.size() - 1; ++i)
      {
         res.ar[i + 1] += res.ar[i] / 10;
         res.ar[i] %= 10;
      }
      res.updateLen();
      return res;
   }

   Large& operator+=(const Large& rhs)
   {
      for (int i = 0; i < ar.size(); ++i) ar[i] += rhs.ar[i];
      for (int i = 0; i < ar.size() - 1; ++i)
      {
         ar[i + 1] += ar[i] / 10;
         ar[i] %= 10;
      }
      updateLen();
      return *this;
   }

   Large operator-(const Large& rhs) const
   {
      Large res;
      for (int i = 0; i < ar.size(); ++i) res.ar[i] = ar[i] - rhs.
          ar[i];
      for (int i = 0; i < ar.size() - 1; ++i)
      {
         if (res.ar[i] < 0)
         {
            res.ar[i] += 10;
            res.ar[i + 1]--;
         }
      }
      res.updateLen();
      return res;
   }

   Large operator*(const ll rhs) const
   {
      Large res;
      for (int i = 0; i < ar.size(); ++i) res.ar[i] = ar[i] * rhs;
      for (int i = 0; i < ar.size() - 1; ++i)
      {
         if (res.ar[i] > 9)
         {
            res.ar[i + 1] += res.ar[i] / 10;
            res.ar[i] %= 10;
         }
      }
      res.updateLen();
      return res;
   }

   Large& operator*=(const ll rhs)
   {
      for (int i = 0; i < ar.size(); ++i) ar[i] *= rhs;
      for (int i = 0; i < ar.size() - 1; ++i)
      {
         if (ar[i] > 9)
         {
            ar[i + 1] += ar[i] / 10;
            ar[i] %= 10;
         }
      }
      updateLen();
      return *this;
   }

   Large operator*(const Large& rhs) const
   {
      Large res;
      Large dup = *this;
      for (int i = 0; i < rhs.len; ++i)
      {
         res += dup * rhs.ar[i];
         dup *= 10;
      }
      return res;
   }

   Large& operator*=(const Large& rhs)
   {
      *this = *this * rhs;
      return *this;
   }
};

ostream& operator<<(ostream& out, const Large& large)
{
   if (large.len == 0)
   {
      cout << '0';
      return out;
   }
   for (int i = large.len - 1; i >= 0; --i) cout << large.ar[i];
   return out;
}
```

4.7 连续乘法逆元

```cpp
/*******************************************************************
* 时间复杂度: O(n)
* 说明:
* 1.线性时间复杂度求[1,n]的逆元
* 2.注意mod必须与1-n所有数互质，否则不存在逆元
*******************************************************************/
struct ConInv
{
   vector<ll> inv;
   ConInv(int sz, ll mod)
   {
      inv.resize(sz);
      inv[1] = 1;
      for (int i = 2; i <= sz; ++i)
      {
         inv[i] = (mod - mod / i) * inv[mod % i] % mod;
      }
   }
};
```

4.8 数论分块

```
/*****************************************************************
* 时间复杂度：O(sqrt(n))
* 说明：k%i=k-k/i*i => sigma(k%i)=k*n-sigma(k/i*i)
*****************************************************************/
ll n, k;

int main()
{
    //求sigma[i=1,n](k%i)
    ll ans = 0;
    cin >> n >> k;
    for (ll lef = 1, rig; lef <= n; lef = rig + 1) //分块
    {
        if (k >= lef)
        {
            rig = min(n, k / (k / lef));
        }
        else //该区间大于k（余数都为k）
        {
            rig = n;
        }
        ans += k * (rig - lef + 1) - (k / lef) * (lef + rig) * (rig -
            lef + 1) / 2;
    }
    cout << ans << '\n';
    return 0;
}
```

4.9 欧拉函数

```
/*****************************************************************
* 时间复杂度：O(sqrt(n))
* 说明：
* 1.欧拉函数的性质：
* I.phi(x)=x*П((p[i]-1)/p[i])，p[i]为x的第i个质因数；
* II.若x为质数：
* i%x==0 => phi(i*x)=x*phi(i)
* i%x!=0 => phi(i*x)=(x-1)*phi(i)
* 2.若求[l,r]内的欧拉函数，可以先筛出sqrt(r)以内的质数，用这些质数
* 贡献范围内的数，再特判所有数sqrt(r)以上的质因子即可，类似素数筛。
*****************************************************************/
//求n的欧拉函数，类似于质因数分解
int phi(int n)
{
    int res = n;
    for (int i = 2; i * i <= n; i++)
    {
        if (n % i == 0) res = res / i * (i - 1);
        while (n % i == 0) n /= i;
    }
    if (n > 1) res = res / n * (n - 1);
    return res;
}
```

4.10 线性素数筛

```
/*****************************************************************
* 时间复杂度：O(n)
* 说明：
* 1.筛出x以内所有质数
* 2.sieve[i]表征i是否为合数
* 3.prime中为升序排列的质数
*****************************************************************/
struct PrimeSieve
{
    vector<int> sieve;
    vector<ll> prime;

    void build(int x)
    {
        sieve.resize(x+1);
        sieve[1] = 1;
        for (int i = 2; i <= x; ++i)
        {
            if (sieve[i] == 0) prime.push_back(i);
            for (auto e : prime)
            {
                if (e > x / i) break;
                sieve[i * e] = 1;
                if (i % e == 0) break;
            }
        }
        return;
    }
};
```

4.11 欧几里得算法 + 扩展欧几里得算法

```
/*****************************************************************
* 时间复杂度：O(logn)
* 说明：
* 1.欧几里得算法：求最大公因数
* 2.扩展欧几里得算法：求解ax+by=gcd(a,b)
* 3.由扩展欧几里得算法求出一组解x1,y1后，可得解集：
    x=x1+b/gcd(a,b)*k;
    y=y1-a/gcd(a,b)*k;
    其中k为任意整数
* 4.ax+by=1有解=>1是gcd(a,b)倍数=>gcd(a,b)=1
* 5.扩展欧几里得还可以求乘法逆元
*****************************************************************/
ll gcd(ll a, ll b)
{
    return b == 0 ? a : gcd(b, a % b);
}

ll exgcd(ll a, ll b, ll& x, ll& y)
{
    if (b == 0) { x = 1, y = 0; return a; }
    ll d = exgcd(b, a % b, x, y);
    ll newx = y, newy = x - a / b * y;
    x = newx, y = newy;
    return d;
}

ll inv(ll a, ll mod)
{
    ll x, y;
    exgcd(a, mod, x, y);
    return x;
}

ll a, b, x, y, g;

void solve()
{
    cin >> a >> b;
    g = exgcd(a, b, x, y);
    auto M = [](ll x, ll m) {return (x % m + m) % m; };
    cout << M(x, b / g) << '\n';
    return;
}
```

4.12 哥德巴赫猜想

```
// 1. >=6 的整数可以写成三个质数之和
// 2. >=4 的偶数可以写成两个质数之和
// 3. >=7 的奇数可以写成三个奇质数之和
```

5 数据结构

5.1 哈希表

```
/*****************************************************************
* 时间复杂度：O(1)
* 说明：
* 1.自定义随机化哈希函数，降低碰撞概率
* 2.unordered_map采用开链法，gp_hash_table采用查探法
*****************************************************************/
#include<bits/stdc++.h>
#include<unordered_map>

#include<ext/pb_ds/assoc_container.hpp>
#include<ext/pb_ds/hash_policy.hpp>

using namespace std;
using ll = long long;

using namespace __gnu_pbds;

struct CustomHash
```

```cpp
{
    static uint64_t splitmix64(uint64_t x)
    {
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }

    size_t operator()(uint64_t x) const
    {
        static const uint64_t FIXED_RANDOM = chrono::steady_clock::
            now().time_since_epoch().count();
        return splitmix64(x + FIXED_RANDOM);
    }
};

unordered_map<ll, ll, CustomHash> mp;
gp_hash_table<ll, ll, CustomHash> ht;
```

5.2 ST 表

```cpp
/*******************************************************************
* 时间复杂度: 建表O(nlogn)/查询O(1)
* 说明: 可重复贡献问题[f(r,r)=r]的静态区间查询, 一般是最值/gcd
*******************************************************************/
struct ST
{
    int sz;
    vector<vector<ll>> st;

    ST(int x) { init(x); }
    void init(int x)
    {
        sz = x;
        st.resize(sz + 1, vector<ll>(32));
    }
    void build(ll arr[])
    {
        for (int i = 1; i <= sz; ++i) st[i][0] = arr[i];
        int lg = log2(sz);
        for (int i = 1; i <= lg; ++i)
        {
            for (int j = 1; j <= sz; ++j)
            {
                st[j][i] = st[j][i - 1];
                if (j + (1 << (i - 1)) <= sz)
                {
                    st[j][i] = max(st[j][i], st[j + (1 << (i - 1))][i -
                        1]);
                }
            }
        }
    }
    ll query(int lef, int rig)
    {
        int len = int(log2(rig - lef + 1));
        return max(st[lef][len], st[rig - (1 << len) + 1][len]);
    }
};
```

5.3 并查集

```cpp
/*******************************************************************
* 时间复杂度: 查找近似O(1)/合并近似O(1)
* 说明: 路径压缩/启发式合并保证时间复杂度
*******************************************************************/
struct DSU
{
    vector<int> f;
    vector<int> v; //集合大小
    DSU() {}
    DSU(int x) { init(x); }
    void init(int x)
    {
        f.resize(x + 1);
        v.resize(x + 1);
        for (int i = 1; i <= x; ++i) f[i] = i;
        for (int i = 1; i <= x; ++i) v[i] = 1;
        return;
    }
    int find(int id) { return f[id] == id ? id : f[id] = find(f[id]);
        }
    void merge(int x, int y)
    {
        int fx = find(x), fy = find(y);
        if (fx == fy) return;
        if (v[fx] > v[fy]) swap(fx, fy);
        f[fx] = fy;
        v[fy] += v[fx];
        return;
    }
};
```

5.4 树状数组

```cpp
/*******************************************************************
* 时间复杂度: 建立O(n)/修改O(logn)/查询O(logn)
* 说明:
* 1.动态维护满足区间减法的性质, 一般是求和
* 2.单点修改, 区间查询
* 3.时间戳优化可以替代暴力清空
* 4.将加法换成取最值就可以维护不可逆前缀最值
*******************************************************************/
struct Fenwick
{
    int sz;
    vector<ll> tree;
    //vector<int> tag;
    //int now;

    inline int lowbit(int x) { return x & -x; }

    Fenwick() {}
    Fenwick(int x) { init(x); }
    void init(int x)
    {
        sz = x;
        tree.resize(sz + 1);
        //tag.resize(sz + 1);
        //now = 0;
    }
    /*
    void clear()
    {
        now++;
        return;
    }
    */
    void add(int dst, ll v)
    {
        while (dst <= sz)
        {
            //if (tag[dst] != now) tree[dst] = 0;
            tree[dst] += v;
            //tag[dst] = now;
            dst += lowbit(dst);
        }
        return;
    }
    ll pre(int dst)
    {
        ll res = 0;
        while (dst)
        {
            /*
            if (tag[dst] == now) res += tree[dst];
            dst -= lowbit(dst);
            */
            res += tree[dst];
            dst -= lowbit(dst);
        }
        return res;
    }
    inline ll rsum(int lef, int rig) { return pre(rig) - pre(lef - 1)
        ; }
    void build(ll arr[])
    {
        for (int i = 1; i <= sz; ++i)
        {
            tree[i] += arr[i];
            int j = i + lowbit(i);
            if (j <= sz) tree[j] += tree[i];
        }
        return;
    }
};
```

5.5 二维树状数组

```
/*******************************************************************
* 时间复杂度：修改O(log^2n)/查询O(log^2n)
* 说明：普通树状数组的二维版，维护矩阵
*******************************************************************/
struct Fenwick2
{
   int sz;
   vector<vector<ll>> tree;

   inline int lowbit(int x) { return x & -x; }

   Fenwick2() {}
   Fenwick2(int x) { init(x); }

   void init(int x)
   {
      sz = x;
      tree.resize(sz + 1, vector<ll>(sz + 1));
      return;
   }

   void add(int x, int y, ll val)
   {
      for (int i = x; i <= sz; i += lowbit(i))
      {
         for (int j = y; j <= sz; j += lowbit(j))
         {
            tree[i][j] += val;
         }
      }
      return;
   }

   ll pre(int x, int y)
   {
      ll res = 0;
      for (int i = x; i >= 1; i -= lowbit(i))
      {
         for (int j = y; j >= 1; j -= lowbit(j))
         {
            res += tree[i][j];
         }
      }
      return res;
   }

   ll sum(int x1, int y1, int x2, int y2)
   {
      return pre(x2, y2) - pre(x1 - 1, y2) - pre(x2, y1 - 1) + pre(
          x1 - 1, y1 - 1);
   }
};
```

5.6 线段树

```
/*******************************************************************
* 时间复杂度：建立O(n)/询问O(logn)/修改O(logn)
* 说明：
* 1.维护区间性质，要求性质能由子区间性质得到。
* 2.区间修改，区间查询。若仅单点修改则不需要标记。
* 3.使线段树维护不同性质只需要改变上方函数和默认值。
* 4.modify也要push。
*******************************************************************/
struct SegTree
{
   struct Node
   {
      int lef, rig;
      ll val, tag;
   };
   vector<Node> tree;

   const ll VDEF = 0;
   const ll TDEF = 0;

   inline void update(int src) //由子节点及其标记更新父节点
   {
      ll lw = tree[src << 1].rig - tree[src << 1].lef + 1;
      ll rw = tree[src << 1 | 1].rig - tree[src << 1 | 1].lef + 1;
      ll lv = tree[src << 1].val + tree[src << 1].tag * lw;
      ll rv = tree[src << 1 | 1].val + tree[src << 1 | 1].tag * rw;
      tree[src].val = merge(lv, rv);
      return;
   }
   inline void push(int src) //下传标记并消耗
   {
      if (tree[src].lef < tree[src].rig)
      {
         tree[src << 1].tag += tree[src].tag;
         tree[src << 1 | 1].tag += tree[src].tag;
      }
      ll wid = tree[src].rig - tree[src].lef + 1;
      tree[src].val += tree[src].tag * wid;
      tree[src].tag = TDEF;
      return;
   }
   inline void mark(int src, ll val) //更新标记
   {
      tree[src].tag += val;
      return;
   }
   inline ll merge(ll x, ll y) //合并两个查询结果
   {
      return x + y;
   }

   SegTree() {}
   SegTree(int x) { init(x); }
   void init(int x) { tree.resize(x * 4 + 1); }

   void build(int src, int lef, int rig, ll arr[])
   {
      tree[src] = { lef, rig, VDEF, TDEF };
      if (lef == rig) tree[src].val = arr[lef];
      else
      {
         int mid = lef + (rig - lef) / 2;
         build(src << 1, lef, mid, arr);
         build(src << 1 | 1, mid + 1, rig, arr);
         update(src);
      }
      return;
   }
   void build(int src, int lef, int rig)
   {
      tree[src] = { lef, rig, VDEF, TDEF };
      if (lef == rig) return;
      int mid = lef + (rig - lef) / 2;
      build(src << 1, lef, mid);
      build(src << 1 | 1, mid + 1, rig);
      update(src);
      return;
   }
   void assign(int src, int pos, ll val)
   {
      push(src);
      if (tree[src].lef == tree[src].rig)
      {
         tree[src].val = val;
         return;
      }
      if (p <= tree[src << 1].rig) assign(src << 1, pos, val);
      else assign(src << 1 | 1, pos, val);
      update(src);
      return;
   }
   void modify(int src, int lef, int rig, ll val)
   {
      if (lef <= tree[src].lef && tree[src].rig <= rig)
      {
         mark(src, val);
         return;
      }
      push(src);
      if (lef <= tree[src << 1].rig) modify(src << 1, lef, rig, val
          );
      if (rig >= tree[src << 1 | 1].lef) modify(src << 1 | 1, lef,
          rig, val);
      update(src);
      return;
   }
   ll query(int src, int lef, int rig)
   {
      push(src);
      if (lef <= tree[src].lef && tree[src].rig <= rig) return tree
          [src].val;
      ll res = VDEF;
      if (lef <= tree[src << 1].rig) res = merge(res, query(src <<
          1, lef, rig));
      if (rig >= tree[src << 1 | 1].lef) res = merge(res, query(src
           << 1 | 1, lef, rig));
      return res;
   }
};
```

5.7 动态开点线段树

```
/*****************************************************************
* 时间复杂度：询问O(logn)/修改O(logn)
* 说明：注意空间大小
*****************************************************************/
struct SegTree
{
   struct Node
   {
      int ls = 0, rs = 0;
      ll val = 0, tag = 0;
   };
   vector<Node> tree;
   int ord;
   SegTree(int x)
   {
      tree.resize(x * 64 + 1);
      ord = 1;
   }
   void push(int src, int lef, int rig)
   {
      if (lef < rig)
      {
         if (!tree[src].ls) tree[src].ls = ++ord;
         if (!tree[src].rs) tree[src].rs = ++ord;
         tree[tree[src].ls].tag += tree[src].tag;
         tree[tree[src].rs].tag += tree[src].tag;
      }
      tree[src].val += tree[src].tag * (rig - lef + 1);
      tree[src].tag = 0;
      return;
   }
   void modify(int src, int lef, int rig, int l, int r, ll val)
   {
      if (lef >= l && rig <= r)
      {
         tree[src].tag += val;
         return;
      }
      int mid = lef + (rig - lef) / 2;
      if (l <= mid)
      {
         if (!tree[src].ls) tree[src].ls = ++ord;
         modify(tree[src].ls, lef, mid, l, r, val);
      }
      if (r >= mid + 1)
      {
         if (!tree[src].rs) tree[src].rs = ++ord;
         modify(tree[src].rs, mid + 1, rig, l, r, val);
      }
      tree[src].val += (min(rig, r) - max(lef, l) + 1) * val;
      return;
   }
   ll query(int src, int lef, int rig, int l, int r)
   {
      push(src, lef, rig);
      if (lef >= l && rig <= r) return tree[src].val;
      ll res = 0;
      int mid = lef + (rig - lef) / 2;
      if (l <= mid)
      {
         if (!tree[src].ls) tree[src].ls = ++ord;
         res += query(tree[src].ls, lef, mid, l, r);
      }
      if (r >= mid + 1)
      {
         if (!tree[src].rs) tree[src].rs = ++ord;
         res += query(tree[src].rs, mid + 1, rig, l, r);
      }
      return res;
   }
};
```

5.8 可持久化线段树

```
/*****************************************************************
* 时间复杂度：所有操作O(log(seglen))
* 说明：
* 1.建空根：可以不靠离散化维护大区间，但要谨慎考虑空间复杂度。
* 2.主席树维护区间值域上性质：用可持久化权值线段树维护值域，将序列元素逐
* 个插入，由前缀和性质，区间值域上性质蕴含在新树和旧树的差之中。
* 3.标记永久化：路过结点时标记不下放，也不通过子结点更新，而是直接改变其
* 值；向下搜索时记录累积标记值并在最后作用（因此assign()在维护最值时
* 无效）。
* 4.区间第k大也可以整体二分/划分树。
* 5.若维护区间超过int，记得把32换成64。
*****************************************************************/
struct PerSegTree
{
   struct Node
   {
      int ls, rs;
      ll val, tag;
      Node(): ls(0), rs(0), val(0), tag(0) {}
   };
   vector<Node> tree;
   vector<int> root;
   int size;
   ll L, R;

   int _build(ll l, ll r, ll a[])
   {
      int now = size++;
      if (l == r) tree[now].val = a[l];
      else
      {
         ll m = l + (r - l) / 2;
         tree[now].ls = _build(l, m, a);
         tree[now].rs = _build(m + 1, r, a);
         tree[now].val = tree[tree[now].ls].val + tree[tree[now].rs
             ].val;
      }
      return now;
   }
   void init(ll l, ll r, int cnt, ll a[]) //建初始树
   {
      size = 0;
      L = l, R = r;
      tree.resize(cnt * 32 + 5);
      root.push_back(_build(L, R, a));
      return;
   }
   void init(ll l, ll r, int cnt) //建一个空根
   {
      size = 1;
      L = l, R = r;
      tree.resize(cnt * 32 + 5);
      root.push_back(0);
      return;
   }
   void assign(int ver, ll pos, ll val) { root.push_back(_assign(
       root[ver], L, R, pos, val, 0)); }
   int _assign(int src, ll l, ll r, ll pos, ll val, ll tag)
   {
      int now = size++;
      tree[now] = tree[src];
      tag += tree[now].tag;
      if (l == r) tree[now].val = val - tag;
      else
      {
         ll m = l + (r - l) / 2;
         if (pos <= m) tree[now].ls = _assign(tree[now].ls, l, m,
             pos, val, tag);
         else tree[now].rs = _assign(tree[now].rs, m + 1, r, pos,
             val, tag);
      }
      return now;
   }
   void modify(int ver, ll lef, ll rig, ll val) { root.push_back(
       _modify(root[ver], L, R, lef, rig, val)); }
   int _modify(int src, ll l, ll r, ll lef, ll rig, ll val)
   {
      int now = size++;
      tree[now] = tree[src];
      if (lef <= l && r <= rig) tree[now].tag += val;
      else if (l <= rig && r >= lef)
      {
         tree[now].val += val * (min(rig, r) - max(lef, l) + 1);
         ll m = l + (r - l) / 2;
         if (lef <= m) tree[now].ls = _modify(tree[now].ls, l, m,
             lef, rig, val);
         if (rig > m) tree[now].rs = _modify(tree[now].rs, m + 1, r
             , lef, rig, val);
      }
      return now;
   }
   ll query(int ver, ll lef, ll rig) { return _query(root[ver], L, R
       , lef, rig, 0); }
   ll _query(int src, ll l, ll r, ll lef, ll rig, ll tag)
   {
```

```cpp
        tag += tree[src].tag;
        if (lef <= l && r <= rig) return tree[src].val + (r - l + 1)
            * tag;
        else if (l <= rig && r >= lef)
        {
            int m = l + (r - l) / 2;
            ll res = 0;
            if (lef <= m) res += _query(tree[src].ls, l, m, lef, rig,
                tag);
            if (rig > m) res += _query(tree[src].rs, m + 1, r, lef,
                rig, tag);
            return res;
        }
        else return 0;
    }
    ll kth(ll lef, ll rig, int k) { return _kth(root[lef - 1], root[
        rig], L, R, k); }
    ll _kth(int osrc, int nsrc, ll l, ll r, int k)
    {
        if (l == r) return l;
        int nsum = tree[tree[nsrc].ls].val + tree[tree[nsrc].ls].tag;
        int osum = tree[tree[osrc].ls].val + tree[tree[osrc].ls].tag;
        int dif = nsum - osum;
        int m = l + (r - l) / 2;
        if (dif >= k) return _kth(tree[osrc].ls, tree[nsrc].ls, l, m,
            k);
        else return _kth(tree[osrc].rs, tree[nsrc].rs, m + 1, r, k -
            dif);
    }
};
```

5.9 李超线段树

```cpp
/**********************************************************************
* 时间复杂度: 建立O(n)/修改O(log^2n)/查询O(logn)
* 说明:
* 1.谨慎使用, 注意浮点数精度和结点初始化问题
* 2.标记永久化, 整条链每一层的值都可能是答案
**********************************************************************/
const int N = 100005;
const double EPS = 1e-9;

struct Seg
{
    double k, b;
    int lef, rig;
    void init(int x0, int y0, int x1, int y1)
    {
        lef = x0, rig = x1;
        if (x0 == x1)
        {
            k = 0, b = max(y0, y1);
        }
        else
        {
            k = double(y1 - y0) / (x1 - x0);
            b = y0 - x0 * k;
        }
    }
    double at(int x) { return k * x + b; }
} seg[N];

struct LCSegTree
{
    struct Node
    {
        int lef, rig, id;
    };
    vector<Node> tree;

    LCSegTree(int x) { tree.resize(x * 4 + 1); }

    void build(int src, int lef, int rig)
    {
        tree[src] = { lef, rig, 0 };
        if (lef == rig) return;
        int mid = (lef + rig) / 2;
        build(src << 1, lef, mid);
        build(src << 1 | 1, mid + 1, rig);
        return;
    }

    void add(int src, int id)
    {
        if (seg[id].lef <= tree[src].lef && seg[id].rig >= tree[src].
            rig)
        {
            update(src, id);
            return;
        }
        if (seg[id].lef <= tree[src << 1].rig) add(src << 1, id);
        if (seg[id].rig >= tree[src << 1 | 1].lef) add(src << 1 | 1,
            id);
        return;
    }

    bool compare(int id1, int id2, int x)
    {
        if (id1 == 0) return 1;
        if (id2 == 0) return 0;
        double r1 = seg[id1].at(x);
        double r2 = seg[id2].at(x);
        if (fabs(r1 - r2) < EPS) return id2 < id1;
        else return r2 > r1 + EPS;
    }

    void update(int src, int id)
    {
        int mid = (tree[src].lef + tree[src].rig) / 2;
        if (compare(tree[src].id, id, mid)) swap(tree[src].id, id);
        if (tree[src].lef == tree[src].rig) return;
        if (compare(tree[src].id, id, tree[src].lef)) update(src <<
            1, id);
        if (compare(tree[src].id, id, tree[src].rig)) update(src << 1
            | 1, id);
        return;
    }

    int query(int src, int x)
    {
        if (tree[src].lef == tree[src].rig) return tree[src].id;
        if (x <= tree[src << 1].rig)
        {
            int r = query(src << 1, x);
            if (compare(r, tree[src].id, x)) return tree[src].id;
            else return r;
        }
        else
        {
            int r = query(src << 1 | 1, x);
            if (compare(r, tree[src].id, x)) return tree[src].id;
            else return r;
        }
    }
};
```

6 树论

6.1 LCA

```cpp
/**********************************************************************
* 时间复杂度: O(logM)
* 说明: 适用于有根树
**********************************************************************/
const int N = 500005;

vector<int> node[N];

struct LCA
{
    vector<int> d; //到根距离
    vector<vector<int>> st;

    void dfs(int x)
    {
        for (auto e : node[x])
        {
            if (e == st[x][0]) continue;
            d[e] = d[x] + 1;
            st[e][0] = x;
            dfs(e);
        }
        return;
    }

    void build(int sz)
    {
        int lg = int(log2(sz));
        for (int i = 1; i <= lg; ++i)
        {
            for (int j = 1; j <= sz; ++j)
```

```cpp
            {
                if (d[j] >= (1 << i))
                {
                    st[j][i] = st[st[j][i - 1]][i - 1];
                }
            }
        }
        return;
    }

    LCA() {}
    LCA(int x, int root) { init(x, root); }

    void init(int x, int root)
    {
        d.resize(x + 1);
        st.resize(x + 1, vector<int>(32));
        dfs(root);
        build(x);
        return;
    }

    int query(int a, int b)
    {
        if (d[a] < d[b]) swap(a, b);
        int dif = d[a] - d[b];
        for (int i = 0; dif >> i; ++i)
        {
            if (dif >> i & 1) a = st[a][i];
        }
        if (a == b) return a;
        else
        {
            for (int i = 31; i >= 0; --i)
            {
                while (st[a][i] != st[b][i])
                {
                    a = st[a][i];
                    b = st[b][i];
                }
            }
            return st[a][0];
        }
    }
};
```

6.2 树的直径

```cpp
/*********************************************************************
* 时间复杂度：O(N)
* 说明：
* 1.距离任一点最远的点一定是直径的一端
* 2.任一点距所有叶的最远距离对应的叶一定是直径端点
*********************************************************************/
const int N = 200005;

struct Edge { int to; ll v; };
vector<Edge> node[N];

pair<int, ll> farthest(int id, ll d, int pa)
{
    pair<int, ll> ret = { id,d };
    for (auto e : node[id])
    {
        pair<int, ll> res;
        if (e.to != pa) res = farthest(e.to, d + e.v, id);
        if (res.second > ret.second) ret = res;
    }
    return ret;
}

int n, m;

void solve()
{
    cin >> n >> m;
    int u, v;
    ll w;
    for (int i = 1; i <= m; ++i)
    {
        cin >> u >> v >> w;
        node[u].push_back({ v,w });
        node[v].push_back({ u,w });
    }
    int s = farthest(1, 0, -1).first;
    auto res = farthest(s, 0, -1);
    int t = res.first;
    ll d = res.second;
    return;
}
```

6.3 树哈希

```cpp
/*********************************************************************
* 时间复杂度：O(nlogn)
* 说明：判断有根树同构。无根树可通过找重心转换为有根树。
*********************************************************************/
struct TreeHash
{
    int n, root;
    vector<vector<int>> node;
    vector<int> hav;
    map<vector<int>, int> mp;
    int ord = 0;

    void getTree(vector<int>& p)
    {
        n = p.size() - 1;
        node.clear();
        node.resize(n + 1);
        hav.clear();
        hav.resize(n + 1);
        root = -1;
        for (int i = 1; i <= n; ++i)
        {
            if (p[i])
            {
                node[p[i]].push_back(i);
                node[i].push_back(p[i]);
            }
            else root = i;
        }
        return;
    }

    void getD(int id, int pa, vector<int>& sz, vector<int>& d)
    {
        sz[id] = 1;
        int res = 0;
        for (auto e : node[id])
        {
            if (e != pa)
            {
                getD(e, id, sz, d);
                sz[id] += sz[e];
                res = max(res, sz[e]);
            }
        }
        if (id == root) d[id] = res;
        else d[id] = max(res, n - sz[id]);
        return;
    }

    vector<int> center()
    {
        vector<int> res;
        vector<int> sz(n + 1), d(n + 1);
        int mnn = n;
        getD(root, -1, sz, d);
        for (int i = 1; i <= n; ++i) mnn = min(mnn, d[i]);
        for (int i = 1; i <= n; ++i) if (d[i] == mnn) res.push_back(i
            );
        return res;
    }

    vector<int> hash(vector<int>& p)
    {
        vector<int> res;
        getTree(p);
        auto v = center();
        for (auto e : v) dfs(e, -1), res.push_back(hav[e]);
        sort(res.begin(), res.end());
        return res;
    }

    int hash(vector<int>& p, int root)
    {
        getTree(p);
        dfs(root, -1);
        return hav[root];
    }
```

```
    void dfs(int id, int pa)
    {
        vector<int> v;
        for (auto e : node[id])
        {
            if (e != pa)
            {
                dfs(e, id);
                v.push_back(hav[e]);
            }
        }
        sort(v.begin(), v.end());
        if (mp.count(v) == 0) mp[v] = ++ord;
        hav[id] = mp[v];
        return;
    }
};
```

6.4 树链剖分

```
/*******************************************************************
* 时间复杂度: O(nlogn)
* 说明: 维护树上两点间路径相关性质, 也可求LCA。
*******************************************************************/
const int N = 100005;

vector<int> node[N];

struct HLD
{
    vector<int> pa, dep, sz, hson;
    vector<int> top, dfn, rnk;
    int ord = 0;

    HLD(int x, int root)
    {
        pa.resize(x + 1);
        dep.resize(x + 1);
        sz.resize(x + 1);
        hson.resize(x + 1);
        top.resize(x + 1);
        dfn.resize(x + 1);
        rnk.resize(x + 1);
        build(root);
        decom(root);
    }

    void build(int x)
    {
        sz[x] = 1;
        int mxsz = 0;
        for (auto e : node[x])
        {
            if (e != pa[x])
            {
                pa[e] = x;
                dep[e] = dep[x] + 1;
                build(e);
                sz[x] += sz[e];
                if (sz[e] > mxsz)
                {
                    mxsz = sz[e];
                    hson[x] = e;
                }
            }
        }
        return;
    }

    void decom(int x)
    {
        top[x] = x;
        dfn[x] = ++ord;
        rnk[ord] = x;
        if (hson[pa[x]] == x) top[x] = top[pa[x]];
        for (auto e : node[x]) if (e == hson[x]) decom(e);
        for (auto e : node[x]) if (e != pa[x] && e != hson[x]) decom(
            e);
        return;
    }

    int lcm(int u, int v)
    {
        while (top[u] != top[v])
        {
            if (dep[u] < dep[v]) v = pa[top[v]];
            else u = pa[top[u]];
        }
        if (dep[u] < dep[v]) return u;
        else return v;
    }
};
```

6.5 树上启发式合并

```
/*******************************************************************
* 时间复杂度: O(nlogn)(*状态更新复杂度)
* 说明:
* 1.维护一个用于得出答案的状态, 离线预处理每个子树的答案
* 2.用dfn序代替递归的贡献计算和清除可以优化常数
*******************************************************************/
const int N = 100005;

vector<int> node[N];

int n;
ll a[N];

struct DsuOnTree
{
    struct State
    {
        vector<int> cnt;
        map<int, ll> mp;
        State() { init(); }
        void init() { cnt.resize(1e5 + 1); }
        void add(ll val)
        {
            if (cnt[val]) mp[cnt[val]] -= val;
            if (mp[cnt[val]] == 0) mp.erase(cnt[val]);
            cnt[val]++;
            mp[cnt[val]] += val;
            return;
        }
        void del(ll val)
        {
            mp[cnt[val]] -= val;
            if (mp[cnt[val]] == 0) mp.erase(cnt[val]);
            cnt[val]--;
            if (cnt[val]) mp[cnt[val]] += val;
            return;
        }
        ll ans() { return mp.rbegin()->second; }
    } state;
    vector<int> big; //每个结点的重子
    vector<int> sz; //每个子树的大小
    vector<ll> ans; //每个子树的答案
    const int root = 1;

    DsuOnTree()
    {
        big.resize(n + 1);
        sz.resize(n + 1);
        ans.resize(n + 1);
    }
    void dfs0(int x, int p)
    {
        sz[x] = 1;
        for (auto e : node[x])
        {
            if (e == p) continue;
            dfs0(e, x);
            sz[x] += sz[e];
            if (sz[big[x]] < sz[e]) big[x] = e;
        }
        return;
    }
    void del(int x, int p) //删除子树贡献
    {
        state.del(a[x]);
        for (auto e : node[x])
        {
            if (e == p) continue;
            del(e, x);
        }
        return;
    }
    void add(int x, int p) //计算子树贡献
    {
        state.add(a[x]);
        for (auto e : node[x])
        {
```

```cpp
            if (e == p) continue;
            add(e, x);
        }
        return;
    }
    void dfs(int x, int p, bool keep)
    {
        for (auto e : node[x]) //计算轻子子树答案
        {
            if (e == big[x] || e == p) continue;
            dfs(e, x, 0);
        }
        if (big[x]) dfs(big[x], x, 1); //计算重子子树答案和贡献
        for (auto e : node[x]) //计算轻子子树贡献
        {
            if (e == big[x] || e == p) continue;
            add(e, x);
        }
        state.add(a[x]); //计算自己贡献
        ans[x] = state.ans(); //计算答案
        if (keep == 0) del(x, p); //删除子树贡献
        return;
    }
    void work()
    {
        dfs0(root, 0);
        dfs(root, 0, 0);
        return;
    }
};

void solve()
{
    cin >> n;
    for (int i = 1; i <= n; ++i) cin >> a[i];
    int u, v;
    for (int i = 1; i <= n - 1; ++i)
    {
        cin >> u >> v;
        node[u].push_back(v);
        node[v].push_back(u);
    }
    DsuOnTree dot;
    dot.work();
    for (int i = 1; i <= n; ++i) cout << dot.ans[i] << ' ';
    cout << endl;
    return;
}
```

6.6 点分治

```cpp
/**********************************************************************
* 时间复杂度：处理结点次数O(nlogn)
* 说明：用于树上路径计数问题：以重心为根分治子树，再计算经过重心的路径
**********************************************************************/
const int N = 100005;
const int D[3][2] = { -1, 0, 1, -1, 0, 1 };

int n, sz[N], maxd[N];
string s;
vector<int> node[N];
bool vis[N];
multiset<pair<int, int>> st;

void getRoot(int x, int fa, int sum, int& root)
{
    sz[x] = 1, maxd[x] = 0;
    for (auto e : node[x])
    {
        if (vis[e] || e == fa) continue;
        getRoot(e, x, sum, root);
        sz[x] += sz[e];
        maxd[x] = max(maxd[x], sz[e]);
    }
    maxd[x] = max(maxd[x], sum - sz[x]);
    if (maxd[x] < maxd[root]) root = x;
    return;
}

void dfs(int x, int fa, pair<int, int> p)
{
    p.first += D[s[x] - 'a'][0];
    p.second += D[s[x] - 'a'][1];
    st.insert(p);
    for (auto e : node[x])
    {
        if (vis[e] || e == fa) continue;
        dfs(e, x, p);
    }
    return;
}

ll work(int x)
{
    ll res = 0;
    multiset<pair<int, int>> ns;
    for (auto e : node[x])
    {
        if (vis[e]) continue;
        dfs(e, x, make_pair(0, 0));
        for (auto p : st)
        {
            pair<int, int> inv;
            inv.first = -(p.first + D[s[x] - 'a'][0]);
            inv.second = -(p.second + D[s[x] - 'a'][1]);
            if (inv == make_pair(0, 0)) res++;
            res += ns.count(inv);
        }
        for (auto p : st) ns.insert(p);
        st.clear();
    }
    return res;
}

ll divide(int x)
{
    ll res = 0;
    vis[x] = 1;
    res += work(x);
    for (auto e : node[x])
    {
        if (vis[e]) continue;
        int root = 0;
        getRoot(e, x, sz[e], root);
        res += divide(root);
    }
    return res;
}

void solve()
{
    cin >> n >> s;
    s = ' ' + s;
    for (int i = 1; i <= n - 1; ++i)
    {
        int u, v;
        cin >> u >> v;
        node[u].push_back(v);
        node[v].push_back(u);
    }
    maxd[0] = n + 1;
    int root = 0;
    getRoot(1, 0, n, root);
    cout << divide(root) << '\n';
    return;
}
```

7 图论

7.1 2-SAT

```cpp
/**********************************************************************
* 时间复杂度：O(N+M)
* 说明：按照推导关系建有向图，判断是否有两个矛盾点在同一强连通分量中
**********************************************************************/
const int N = 2000005;

ll n, m, x, y;
bool a, b;
vector<int> node[N];

struct Tarjan
{
    int sz, cnt, ord;
    stack<int> stk;
    vector<vector<int>> g; //新图
    vector<int> dfn, low, id, val;
    Tarjan(int x)
    {
        sz = x; //点数
        cnt = 0; //强连通分量个数
```

```cpp
        ord = 0; //时间戳
        dfn.resize(sz + 1); //dfs序
        low.resize(sz + 1); //能到达的最小dfn
        id.resize(sz + 1); //对应的强连通分量编号
        val.resize(sz + 1); //新图点权
    }
    void dfs(int x)
    {
        stk.push(x);
        dfn[x] = low[x] = ++ord;
        for (auto e : node[x])
        {
            if (dfn[e] == 0)
            {
                dfs(e);
                low[x] = min(low[x], low[e]);
            }
            else if (id[e] == 0)
            {
                low[x] = min(low[x], low[e]);
            }
        }
        if (dfn[x] == low[x]) //x为强连通分量的根
        {
            cnt++;
            while (dfn[stk.top()] != low[stk.top()])
            {
                id[stk.top()] = cnt;
                stk.pop();
            }
            id[stk.top()] = cnt;
            stk.pop();
        }
        return;
    }
    void shrink()
    {
        for (int i = 1; i <= sz; ++i)
        {
            if (id[i] == 0) dfs(i);
        }
        return;
    }
    void rebuild()
    {
        for (int i = 1; i <= sz; ++i)
        {
            for (auto e : node[i])
            {
                if (id[i] != id[e]) g[id[i]].push_back(id[e]);
            }
        }
        return;
    }
};

struct TwoSat
{
    int sz;
    vector<int> res;
    inline int negate(int x)
    {
        if (x > n) return x - n;
        else return x + n;
    }
    TwoSat(int x)
    {
        sz = x;
        res.resize(sz + 1);
    }
    bool work()
    {
        Tarjan tj(sz * 2);
        tj.shrink();
        for (int i = 1; i <= n; ++i)
        {
            if (tj.id[i] == tj.id[negate(i)]) return 0;
        }
        for (int i = 1; i <= n; ++i)
        {
            res[i] = tj.id[i] > tj.id[negate(i)];
        }
        return 1;
    }
};

void solve()
{
    cin >> n >> m;
    for (int i = 1; i <= m; ++i)
    {
        cin >> x >> a >> y >> b;
        node[x + (!a) * n].push_back(y + b * n);
        node[y + (!b) * n].push_back(x + a * n);
    }
    TwoSat ts(n);
    if (!ts.work()) cout << "IMPOSSIBLE\n";
    else
    {
        cout << "POSSIBLE\n";
        for (int i = 1; i <= n; ++i) cout << ts.res[i] << ' ';
    }
    return;
}
```

7.2 Bellman-Ford 算法

```cpp
/*********************************************************************
* 时间复杂度: O(NM)
* 说明:
* 1.适用于带负权边的单源最短路问题
* 2.可判断负环, negCir()要在work()后调用
*********************************************************************/
const int N = 1505;
const ll INFLL = 0x3f3f3f3f3f3f3f3f;

struct Edge {ll to, v;};

vector<Edge> node[N];

struct BellmanFord
{
    int sz;
    vector<ll> dis;

    BellmanFord(int x)
    {
        sz = x;
        dis.resize(sz + 1, INFLL);
    }

    void work(int s)
    {
        dis[s] = 0;
        for (int i = 1; i <= sz - 1; ++i)
        {
            for (int j = 1; j <= sz; ++j)
            {
                for (auto e : node[j])
                {
                    dis[e.to] = min(dis[e.to], dis[j] + e.v);
                }
            }
        }
        return;
    }

    bool negCir()
    {
        for (int i = 1; i <= sz; ++i)
        {
            for (auto e : node[i])
            {
                if (dis[e.to] > dis[i] + e.v) return 1;
            }
        }
        return 0;
    }
};
```

7.3 Dijkstra 算法

```cpp
/*********************************************************************
* 时间复杂度: 朴素O(N^2)/堆优化O(MlogM)
* 说明:
* 1.只适用于非负边权
* 2.稀疏图用堆优化, 稠密图用朴素
* 3.注意处理图不连通的情况 (dis==INFLL)
*********************************************************************/
const int N = 100005;
const ll INFLL = 0x3f3f3f3f3f3f3f3f;

struct Edge {int to, v;};
```

```cpp
vector<Edge> node[N];

struct Dijkstra
{
   struct NodeInfo
   {
      int id;
      ll d;
      bool operator < (const NodeInfo& p1) const
      {
         return d > p1.d;
      }
   };

   int sz;
   vector<int> vis;
   vector<ll> dis;

   Dijkstra(int x)
   {
      sz = x;
      vis.resize(sz + 1);
      dis.resize(sz + 1, INFLL);
   }

   void workO(int s)
   {
      priority_queue<NodeInfo> pq;
      dis[s] = 0;
      pq.push({ s,0 });
      while (pq.size())
      {
         int now = pq.top().id;
         pq.pop();
         if (vis[now] == 0)
         {
            vis[now] = 1; //被取出一定是最短路
            for (auto e : node[now])
            {
               if (vis[e.to] == 0 && dis[e.to] > dis[now] + e.v)
               {
                  dis[e.to] = dis[now] + e.v;
                  pq.push({ e.to,dis[e.to] });
               }
            }
         }
      }
      return;
   }

   void workS(int s)
   {
      auto take = [&](int x)
      {
         vis[x] = 1;
         for (auto e : node[x])
         {
            dis[e.to] = min(dis[e.to], dis[x] + e.v);
         }
         return;
      };
      dis[s] = 0;
      take(s);
      for (int i = 1; i <= sz - 1; ++i)
      {
         ll mnn = INFLL;
         int id = 0;
         for (int j = 1; j <= sz; ++j)
         {
            if (vis[j] == 0 && dis[j] < mnn)
            {
               mnn = dis[j];
               id = j;
            }
         }
         if (mnn == INFLL) return;
         take(id);
      }
      return;
   }
};
```

7.4 Dinic 算法

```cpp
/*********************************************************************
* 时间复杂度: 最差O(N^2*M)/二分图匹配O(sqrt(N)*M)
* 说明:
* 1.求有向网络最大流/最小割
* 2.也可以求二分图最大匹配
* 3.cap表示残量, cap为0的边满流
*********************************************************************/
const ll INFLL = 0x3f3f3f3f3f3f3f3f;
const int N = 3005;

struct Edge
{
   int to; //终点
   int rev; //反向边对其起点的编号
   ll cap; //残量
   Edge() {}
   Edge(int to, int rev, ll cap) :to(to), rev(rev), cap(cap) {}
};

vector<Edge> node[N];

void AddEdge(int from, int to, ll cap)
{
   int x = node[to].size();
   int y = node[from].size();
   node[from].push_back(Edge(to, x, cap));
   node[to].push_back(Edge(from, y, 0));
   return;
}

struct Dinic
{
   int sz;
   vector<int> dep; //每个点所属层深度
   vector<int> done; //每个点下一个要处理的邻接边
   queue<int> q;

   Dinic(int x)
   {
      sz = x;
      dep.resize(sz + 1);
      done.resize(sz + 1);
   }

   bool bfs(int s, int t) //建立分层图
   {
      for (int i = 1; i <= sz; ++i) dep[i] = 0;
      q.push(s);
      dep[s] = 1;
      done[s] = 0;
      bool f = 0;
      while (q.size())
      {
         int now = q.front();
         q.pop();
         if (now == t) f = 1; //到达终点说明存在增广路
         for (auto e : node[now])
         {
            if (e.cap && dep[e.to] == 0) //还有残量且未访问过
            {
               q.push(e.to);
               done[e.to] = 0; //有增广路, 需要重新处理
               dep[e.to] = dep[now] + 1;
            }
         }
      }
      return f;
   }

   ll dfs(int x, int t, ll flow) //统计增广路总流量
   {
      if (x == t || flow == 0) return flow; //找到汇点或断流
      ll rem = flow; //结点x当前剩余流量
      for (int i = done[x]; i < node[x].size() && rem; ++i)
      {
         done[x] = i; //前i-1条边已经搞定, 不会再有增广路
         auto& e = node[x][i];
         if (e.cap && dep[e.to] == dep[x] + 1)//还有残量且为下一层
         {
            ll inflow = dfs(e.to, t, min(rem, e.cap)); //计算流向e.
               to的最大流量
            if (inflow == 0) dep[e.to] = 0; //e.to无法流入, 本次增广
               不再考虑
            e.cap -= inflow; //更新残量
            node[e.to][e.rev].cap += inflow; //更新反向边
            rem -= inflow; //消耗流量
         }
      }
      return flow - rem;
   }
```

```
    ll work(int s, int t)
    {
        ll aug = 0, ans = 0;
        while (bfs(s, t))
        {
            while (aug = dfs(s, t, INFLL))
            {
                ans += aug;
            }
        }
        return ans;
    }
};
```

7.5 Floyd 算法

```
/**********************************************************************
* 时间复杂度: O(N^3)
* 说明:
* 1.求任意两点间最短路
* 2.最短路计数
**********************************************************************/
const int N = 105;
const ll INFLL = 0x3f3f3f3f3f3f3f3f;

int n, m;
ll dis[N][N], cnt[N][N];

void floyd()
{
    for (int i = 1; i <= n; ++i)
    {
        for (int j = 1; j <= n; ++j)
        {
            for (int k = 1; k <= n; ++k)
            {
                if (dis[j][k] > dis[j][i] + dis[i][k])
                {
                    dis[j][k] = dis[j][i] + dis[i][k];
                    cnt[j][k] = cnt[j][i] * cnt[i][k];
                }
                else if (dis[j][k] == dis[j][i] + dis[i][k])
                {
                    cnt[j][k] += cnt[j][i] * cnt[i][k];
                }
            }
        }
    }
    return;
}

void solve()
{
    cin >> n >> m;
    for (int i = 1; i <= n; ++i)
    {
        for (int j = 1; j <= n; ++j)
        {
            dis[i][j] = INFLL;
        }
    }
    ll u, v, w;
    for (int i = 1; i <= m; ++i)
    {
        cin >> u >> v >> w;
        dis[u][v] = dis[v][u] = w;
        cnt[u][v] = cnt[v][u] = 1;
    }
    floyd();
    return;
}
```

7.6 Kosaraju 算法

```
/**********************************************************************
* 时间复杂度: O(N+M)
* 说明: 有向图强连通分量
**********************************************************************/
const int N = 10005;

vector<int> node[N];

struct Kosaraju
{
    int sz, index = 0;
    vector<int> vis, ord;
    vector<vector<int>> rev;
    vector<int> id; //强连通分量编号
    Kosaraju(int x)
    {
        sz = x;
        vis.resize(sz + 1);
        id.resize(sz + 1);
        rev.resize(sz + 1);
        ord.resize(1);
        for (int i = 1; i <= sz; ++i)
        {
            for (auto e : node[i])
            {
                rev[e].push_back(i);
            }
        }
        for (int i = 1; i <= sz; ++i) if (vis[i] == 0) dfs1(i);
        for (int i = sz; i >= 1; --i) if (id[ord[i]] == 0) index++,
            dfs2(ord[i]);
    }

    void dfs1(int x)
    {
        vis[x] = 1;
        for (auto e : node[x])
        {
            if (vis[e] == 0) dfs1(e);
        }
        ord.push_back(x);
        return;
    }

    void dfs2(int x)
    {
        id[x] = index;
        for (auto e : rev[x])
        {
            if (id[e] == 0) dfs2(e);
        }
        return;
    }
};
```

7.7 Tarjan 算法

```
/**********************************************************************
* 时间复杂度: O(N+M)
* 说明: 求强连通分量, 也可求缩点后新图
**********************************************************************/
const int N = 10005;

int n, m;
int a[N]; //旧图点权
vector<int> node[N];

struct Tarjan
{
    int sz, cnt, ord;
    stack<int> stk;
    vector<vector<int>> g; //新图
    vector<int> dfn, low, id, val;
    Tarjan(int x)
    {
        sz = x; //点数
        cnt = 0; //强连通分量个数
        ord = 0; //时间戳
        dfn.resize(sz + 1); //dfs序
        low.resize(sz + 1); //能到达的最小dfn
        id.resize(sz + 1); //对应的强连通分量编号
        val.resize(sz + 1); //新图点权
    }
    void dfs(int x)
    {
        stk.push(x);
        dfn[x] = low[x] = ++ord;
        for (auto e : node[x])
        {
            if (dfn[e] == 0) //未访问
            {
                dfs(e);
                low[x] = min(low[x], low[e]);
            }
            else if (id[e] == 0) //在栈中
```

```cpp
        {
            low[x] = min(low[x], low[e]);
        }
    }
    if (dfn[x] == low[x]) //x为强连通分量的根
    {
        cnt++;
        while (dfn[stk.top()] != low[stk.top()]) //强连通分量中只有
            根dfn=low
        {
            val[cnt] += a[stk.top()];
            id[stk.top()] = cnt;
            stk.pop();
        }
        val[cnt] += a[stk.top()];
        id[stk.top()] = cnt;
        stk.pop();
    }
    return;
}
void shrink()
{
    for (int i = 1; i <= sz; ++i)
    {
        if (id[i] == 0) dfs(i);
    }
    return;
}
void rebuild()
{
    for (int i = 1; i <= sz; ++i)
    {
        for (auto e : node[i])
        {
            if (id[i] != id[e]) g[id[i]].push_back(id[e]);
        }
    }
    return;
}
};
```

7.8 K 短路

```cpp
/*******************************************************************
* 时间复杂度: O(NklogN)
* 说明: 利用A*算法。以估价函数值优先搜索，第k次访问某结点即k短路。
*******************************************************************/
const int N = 1005;
const ll INFLL = 0x3f3f3f3f3f3f3f3f;

struct E
{
    ll to, v;
};

struct V
{
    ll id, d;
    bool operator<(const V& v) const { return d > v.d; }
};

int n, m, k;
vector<E> node[N];

struct Dijkstra
{
    int sz;
    vector<ll> d;
    vector<int> vis;
    priority_queue<V> pq;
    vector<vector<E>> rev;

    void rebuild()
    {
        for (int i = 1; i <= sz; ++i)
        {
            for (auto e : node[i])
            {
                rev[e.to].push_back({ i,e.v });
            }
        }
        return;
    }
    Dijkstra(int x, int s)
    {
        sz = x;
        d.resize(sz + 1, INFLL);
        vis.resize(sz + 1);
        rev.resize(sz + 1);
        rebuild();
        d[1] = 0;
        pq.push({ 1,0 });
        while (pq.size())
        {
            auto now = pq.top();
            pq.pop();
            if (vis[now.id]) continue;
            vis[now.id] = 1;
            for (auto e : rev[now.id])
            {
                if (vis[e.to] == 0 && d[e.to] > d[now.id] + e.v)
                {
                    d[e.to] = d[now.id] + e.v;
                    pq.push({ e.to, d[e.to] });
                }
            }
        }
    }
};

void solve()
{
    cin >> n >> m >> k;
    int u, v, w;
    for (int i = 1; i <= m; ++i)
    {
        cin >> u >> v >> w;
        node[u].push_back({ v,w });
    }
    Dijkstra dj(n, n);
    priority_queue<V> pq;
    vector<int> vis(n + 1);
    pq.push({ n,dj.d[n] });
    vector<ll> ans(k, -1);
    while (pq.size())
    {
        auto now = pq.top();
        pq.pop();
        if (now.id == 1 && vis[now.id] < k) ans[vis[now.id]] = now.d;
        vis[now.id]++;
        for (auto e : node[now.id])
        {
            if (vis[e.to] >= k) continue;
            pq.push({ e.to,now.d - dj.d[now.id] + e.v + dj.d[e.to] });
        }
    }
    for (int i = 0; i < k; ++i) cout << ans[i] << '\n';
    return;
}
```

7.9 SSP 算法

```cpp
/*******************************************************************
* 时间复杂度: O(NMF) (伪多项式，与最大流有关)
* 说明:
* 1.求最小费用最大流
* 2.无法处理负环，需要提前排除
*******************************************************************/
const int N = 5005;
const ll INFLL = 0x3f3f3f3f3f3f3f3f;

struct Edge
{
    int to; //终点
    int rev; //反向边对其起点的编号
    ll cap; //残量
    ll cost; //单位流量费用
    Edge() {}
    Edge(int to, int rev, ll cap, ll cost) :to(to), rev(rev), cap(cap
        ), cost(cost) {}
};

vector<Edge> node[N];

void addEdge(int from, int to, ll cap, ll cost)
{
    int x = node[to].size();
    int y = node[from].size();
    node[from].push_back(Edge(to, x, cap, cost));
    node[to].push_back(Edge(from, y, 0, -cost));
    return;
}
```

```cpp
struct SSP
{
    int sz;
    vector<ll> dis; //源点到i的最小单位流量费用
    vector<int> vis;
    vector<int> done; //每个点下一个要处理的邻接边
    queue<int> q;
    ll minc, maxf;

    SSP(int x)
    {
        sz = x;
        dis.resize(sz + 1);
        vis.resize(sz + 1);
        done.resize(sz + 1);
        minc = maxf = 0;
    }

    bool spfa(int s, int t) //寻找单位流量费用最小的增广路
    {
        vis.assign(sz + 1, 0);
        done.assign(sz + 1, 0);
        dis.assign(sz + 1, INFLL);
        dis[s] = 0;
        q.push(s);
        vis[s] = 1;
        while (q.size())
        {
            int now = q.front();
            q.pop();
            vis[now] = 0;
            for (auto e : node[now])
            {
                if (e.cap && dis[e.to] > dis[now] + e.cost) //还有残量且
                    可松弛
                {
                    dis[e.to] = dis[now] + e.cost;
                    if (vis[e.to] == 0) q.push(e.to), vis[e.to] = 1;
                }
            }
        }
        return dis[t] != INFLL;
    }

    ll dfs(int x, int p, int t, ll flow) //沿增广路计算流量和费用
    {
        if (x == t || flow == 0) return flow; //找到汇点或断流
        vis[x] = 1; //防止零权环死循环
        ll rem = flow; //结点x当前剩余流量
        for (int i = done[x]; i < node[x].size() && rem; ++i)
        {
            done[x] = i; //前i-1条边已经搞定，不会再有增广路
            auto& e = node[x][i];
            if (e.to != p && vis[e.to] == 0 && e.cap && dis[e.to] ==
                dis[x] + e.cost)
            {
                ll inflow = dfs(e.to, x, t, min(rem, e.cap)); //计算流向
                    e.to的最大流量
                e.cap -= inflow; //更新残量
                node[e.to][e.rev].cap += inflow; //更新反向边
                rem -= inflow; //消耗流量
            }
        }
        vis[x] = 0; //出递归栈后可重新访问
        return flow - rem;
    }

    void work(int s, int t)
    {
        ll aug = 0;
        while (spfa(s, t))
        {
            while (aug = dfs(s, 0, t, INFLL))
            {
                maxf += aug;
                minc += dis[t] * aug;
            }
        }
        return;
    }
};
```

7.10 原始对偶算法

```cpp
/*********************************************************************
* 时间复杂度：O(MlogMF)（伪多项式，与最大流有关）
* 说明：
* 1.求最小费用最大流
* 2.无法处理负环，需要提前排除
*********************************************************************/
const int N = 5005;
const ll INFLL = 0x3f3f3f3f3f3f3f3f;

struct Edge
{
    int to; //终点
    int rev; //反向边对其起点的编号
    ll cap; //残量
    ll cost; //单位流量费用
    Edge() {}
    Edge(int to, int rev, ll cap, ll cost) :to(to), rev(rev), cap(cap
        ), cost(cost) {}
};

vector<Edge> node[N];

void addEdge(int from, int to, ll cap, ll cost)
{
    int x = node[to].size();
    int y = node[from].size();
    node[from].push_back(Edge(to, x, cap, cost));
    node[to].push_back(Edge(from, y, 0, -cost));
    return;
}

struct PrimalDual
{
    struct NodeInfo
    {
        int id;
        ll d;
        bool operator < (const NodeInfo& p1) const
        {
            return d > p1.d;
        }
    };

    int sz;
    vector<ll> h; //势能
    vector<int> vis;
    vector<int> done; //每个点下一个要处理的邻接边
    vector<ll> dis;
    queue<int> q;
    priority_queue<NodeInfo> pq;
    ll minc, maxf;

    PrimalDual(int x)
    {
        sz = x;
        h.resize(sz + 1, INFLL);
        vis.resize(sz + 1);
        done.resize(sz + 1);
        dis.resize(sz + 1);
        minc = maxf = 0;
    }

    void spfa(int s) //求初始势能
    {
        h[s] = 0;
        q.push(s);
        vis[s] = 1;
        while (q.size())
        {
            auto now = q.front();
            q.pop();
            vis[now] = 0;
            for (auto e : node[now])
            {
                if (e.cap && h[e.to] > h[now] + e.cost)
                {
                    h[e.to] = h[now] + e.cost;
                    if (vis[e.to] == 0) q.push(e.to), vis[e.to] = 1;
                }
            }
        }
        return;
    }

    bool dijkstra(int s, int t)
    {
        dis.assign(sz + 1, INFLL);
        vis.assign(sz + 1, 0);
        done.assign(sz + 1, 0);
        dis[s] = 0;
        pq.push({ s,0 });
```

```cpp
        while (pq.size())
        {
            int now = pq.top().id;
            pq.pop();
            if (vis[now] == 0)
            {
                vis[now] = 1; //被取出一定是最短路
                for (auto e : node[now])
                {
                    ll cost = e.cost + h[now] - h[e.to];
                    if (vis[e.to] == 0 && e.cap && dis[e.to] > dis[now]
                         + cost)
                    {
                        dis[e.to] = dis[now] + cost;
                        pq.push({ e.to,dis[e.to] });
                    }
                }
            }
        }
        vis.assign(sz + 1, 0); //还原vis
        return dis[t] != INFLL;
    }

    ll dfs(int x, int t, ll flow) //沿增广路计算流量和费用
    {
        if (x == t || flow == 0) return flow; //找到汇点或断流
        vis[x] = 1; //防止零权环死循环
        ll rem = flow; //结点x当前剩余流量
        for (int i = done[x]; i < node[x].size() && rem; ++i)
        {
            done[x] = i; //前i-1条边已经搞定，不会再有增广路
            auto& e = node[x][i];
            if (vis[e.to] == 0 && e.cap && e.cost == h[e.to] - h[x])
                 //势能差等于费用表明是最短路
            {
                ll inflow = dfs(e.to, t, min(rem, e.cap)); //计算流向e.
                     to的最大流量
                e.cap -= inflow; //更新残量
                node[e.to][e.rev].cap += inflow; //更新反向边
                rem -= inflow; //消耗流量
            }
        }
        vis[x] = 0; //出递归栈后可重新访问
        return flow - rem;
    }

    void work(int s, int t)
    {
        spfa(s);
        ll aug = 0;
        while (dijkstra(s, t))
        {
            for (int i = 1; i <= sz; ++i) h[i] += dis[i]; //更新势能
            while (aug = dfs(s, t, INFLL))
            {
                maxf += aug;
                minc += aug * h[t];
            }
        }
        return;
    }
};
```

7.11 Prim 算法

```cpp
/*********************************************************************
* 时间复杂度：O(N^2)
* 说明：
* 1.选点法最小生成树，适用于稠密图
* 2.注意考虑图不连通的情况
*********************************************************************/
const int N = 5005;
const int M = 200005;

const ll INFLL = 0x3f3f3f3f3f3f3f3f;

struct Edge {ll to, v;};

vector<Edge> node[N];
int n, m;

struct Prim
{
    int sz;
    vector<int> vis;
    vector<ll> dis;

    Prim(int x)
    {
        sz = x;
        vis.resize(sz + 1);
        dis.resize(sz + 1, INFLL);
    }

    ll work()
    {
        int now = 1;
        ll ans = 0;
        for (int i = 1; i <= sz - 1; ++i)
        {
            vis[now] = 1;
            for (auto e : node[now])
            {
                dis[e.to] = min(dis[e.to], e.v);
            }
            ll mnn = INFLL;
            for (int j = 1; j <= sz; ++j)
            {
                if (vis[j] == 0 && dis[j] < mnn)
                {
                    mnn = dis[j];
                    now = j;
                }
            }
            if (mnn == INFLL) return 0; //不连通
            ans += mnn;
        }
        return ans;
    }
};
```

7.12 Kruskal 算法

```cpp
/*********************************************************************
* 时间复杂度：O(MlogM)
* 说明：
* 1.选边法最小生成树，适用于稀疏图
* 2.注意考虑图不连通的情况
*********************************************************************/
const int N = 5005;
const int M = 200005;

struct Edge
{
    ll x, y, v;
    bool operator <(const Edge& e)
    {
        return v < e.v;
    }
};

Edge e[M];
int n, m;

ll kruskal()
{
    DSU dsu(n);
    ll ans = 0;
    sort(e + 1, e + 1 + m);
    for (int i = 1; i <= m; ++i)
    {
        if (dsu.find(e[i].x) != dsu.find(e[i].y))
        {
            ans += e[i].v;
            dsu.merge(e[i].x, e[i].y);
        }
    }
    return ans;
}
```

7.13 Kruskal 重构树

```cpp
/*********************************************************************
* 时间复杂度：建立O(N)/查询O(logN)
* 说明：
* 1.用于解决最小瓶颈路问题
* 2.考虑了初始图不连通的问题
* 3.注意n=1特殊情况（不用建树）
*********************************************************************/
```

```cpp
const int N = 100005;

struct DSU
{
    vector<int> f;
    void init(int x)
    {
        f.resize(x + 1);
        for (int i = 1; i <= x; ++i) f[i] = i;
        return;
    }
    int find(int id) { return f[id] == id ? id : f[id] = find(f[id]);
        }
    void attach(int x, int y) //将fx连向fy，不按秩合并
    {
        int fx = find(x), fy = find(y);
        f[fx] = fy;
        return;
    }
};

struct LCA
{
    vector<int> d;
    vector<vector<int>> st;
    void dfs(int x, vector<vector<int>>& son)
    {
        for (auto e : son[x])
        {
            d[e] = d[x] + 1;
            st[e][0] = x;
            dfs(e, son);
        }
        return;
    }
    void build(int x)
    {
        int lg = int(log2(x));
        for (int i = 1; i <= lg; ++i)
        {
            for (int j = 1; j <= x; ++j)
            {
                if (d[j] >= (1 << i))
                {
                    st[j][i] = st[st[j][i - 1]][i - 1];
                }
            }
        }
        return;
    }
    void init(int x)
    {
        d.resize(x + 1);
        st.resize(x + 1, vector<int>(32));
        return;
    }
    int query(int x, int y)
    {
        if (d[x] < d[y]) swap(x, y);
        int dif = d[x] - d[y];
        for (int i = 0; dif >> i; ++i)
        {
            if (dif >> i & 1) x = st[x][i];
        }
        if (x == y) return x;
        for (int i = 31; i >= 0; --i)
        {
            while (st[x][i] != st[y][i])
            {
                x = st[x][i];
                y = st[y][i];
            }
        }
        return st[x][0];
    }
};

struct Edge
{
    ll x, y, v;
    bool operator<(const Edge& rhs) const { return v < rhs.v; }
} edg[N];

struct KrsRebTree
{
    int size; //当前结点数，最多为n*2-1
    vector<vector<int>> son; //子结点
    vector<ll> val; //点权
    LCA lca;
    DSU dsu;
    void build(int n, int m)
    {
        son.resize(n * 2);
        val.resize(n * 2);
        dsu.init(n * 2 - 1);
        size = n;
        sort(edg + 1, edg + 1 + m);
        for (int i = 1; i <= m && size < n * 2 - 1; ++i)
        {
            int fx = dsu.find(edg[i].x);
            int fy = dsu.find(edg[i].y);
            if (fx == fy) continue;
            size++;
            dsu.attach(fx, size);
            dsu.attach(fy, size);
            son[size].push_back(fx);
            son[size].push_back(fy);
            val[size] = edg[i].v;
        }
        lca.init(size);
        for (int i = n + 1; i <= size; ++i)
        {
            if (dsu.find(i) == i) lca.dfs(i, son); //对所有树的根dfs
        }
        lca.build(size);
        return;
    }
    ll query(int x, int y)
    {
        if (dsu.find(x) == dsu.find(y)) return val[lca.query(x, y)];
        else return -1;
    }
};
```

8 计算几何

8.1 平面坐标旋转

```cpp
/*********************************************************************
* 时间复杂度: O(1)
* 说明: 二维平面上一点绕另一点逆时针旋转
*********************************************************************/
const double PI = acos(-1);

inline double deg_to_rad(int x) { return x * PI / 180; }

struct Point
{
    double x, y;

    void rotate(double rad)
    {
        double newx = x * cos(rad) - y * sin(rad);
        double newy = x * sin(rad) + y * cos(rad);
        x = newx;
        y = newy;
        return;
    }

    void rotate(Point p, double rad)
    {
        Point rela = { x - p.x,y - p.y };
        rela.rotate(rad);
        x = rela.x + p.x;
        y = rela.y + p.y;
        return;
    }
};
```

9 杂项算法

9.1 普通莫队算法

```cpp
/*********************************************************************
* 时间复杂度: O((n+m)sqrt(n))
* 说明: 线性序列普通莫队
*********************************************************************/
const int N = 50005;
const int M = 50005;
```

```cpp
ll n, m, k, a[N], BLOCK;
ll ans[M];

struct Q
{
    ll l, r, id;
    bool operator<(const Q& rhs) const
    {
        //奇偶化排序优化常数
        int lb = l / BLOCK, rb = rhs.l / BLOCK;
        if (lb == rb)
        {
            if (r == rhs.r) return 0;
            else return (r < rhs.r) ^ (lb & 1);
        }
        else return lb < rb;
    }
} q[M];

void solve()
{
    cin >> n >> m >> k;
    BLOCK = n / sqrt(m); //块大小
    for (int i = 1; i <= n; ++i) cin >> a[i];

    //离线处理询问
    for (int i = 1; i <= m; ++i) q[i].id = i, cin >> q[i].l >> q[i].r
        ;
    sort(q + 1, q + 1 + m);

    //计算首个询问答案
    vector<int> cnt(k + 1);
    for (int i = q[1].l; i <= q[1].r; ++i) cnt[a[i]]++;
    ll res = 0;
    for (int i = 1; i <= k; ++i) res += cnt[i] * cnt[i];
    ans[q[1].id] = res;

    //开始转移
    ll l = q[1].l, r = q[1].r;
    auto del = [&](int p)
    {
        res -= cnt[a[p]] * cnt[a[p]];
        cnt[a[p]]--;
        res += cnt[a[p]] * cnt[a[p]];
        return;
    };
    auto add = [&](int p)
    {
        res -= cnt[a[p]] * cnt[a[p]];
        cnt[a[p]]++;
        res += cnt[a[p]] * cnt[a[p]];
        return;
    };
    for (int i = 2; i <= m; ++i)
    {
        while (r < q[i].r) add(++r);
        while (r > q[i].r) del(r--);
        while (l < q[i].l) del(l++);
        while (l > q[i].l) add(--l);
        ans[q[i].id] = res;
    }
    for (int i = 1; i <= m; ++i) cout << ans[i] << '\n';
    return;
}
```

9.2 带修改莫队算法

```cpp
/**********************************************************************
* 时间复杂度：n,m,t同级时O(n^(5/3))
* 说明：
**********************************************************************/
const int N = 150005;
const int M = 150005;

ll BLOCK;

struct Q
{
    ll l, r, id, t;
    bool operator<(const Q& rhs) const
    {
        // 左右端点都分块
        if (l / BLOCK == rhs.l / BLOCK)
        {
            if (r / BLOCK == rhs.r / BLOCK) return t < rhs.t;
            else return r / BLOCK < rhs.r / BLOCK;
        }
        else return l / BLOCK < rhs.l / BLOCK;
    }
} q[M];

struct C
{
    ll p, o, v;
} c[M];

ll n, m, a[N], ans[N];

void solve()
{
    cin >> n >> m;
    BLOCK = pow(n, 2.0 / 3);
    for (int i = 1; i <= n; ++i) cin >> a[i];
    ll mxx = *max_element(a + 1, a + 1 + n);

    // 离线处理询问
    char op;
    ll t = 0, ord = 0, u, v;
    for (int i = 1; i <= m; ++i)
    {
        cin >> op >> u >> v;
        if (op == 'R') c[++t] = { u, a[u], v }, a[u] = v;
        else ord++, q[ord] = { u, v, ord, t };
    }
    sort(q + 1, q + 1 + ord);

    // 计算首个询问答案
    vector<ll> cnt(mxx + 1);
    ll res = 0, l = q[1].l, r = q[1].r, nowt = t;
    auto del = [&](int p)
    {
        cnt[a[p]]--;
        if (cnt[a[p]] == 0) res--;
        return;
    };
    auto add = [&](int p)
    {
        cnt[a[p]]++;
        if (cnt[a[p]] == 1) res++;
        return;
    };
    auto chg = [&](int p, ll v)
    {
        if (p >= l && p <= r) del(p);
        a[p] = v;
        if (p >= l && p <= r) add(p);
        return;
    };
    while (nowt > q[1].t) a[c[nowt].p] = c[nowt].o, nowt--;
    for (int i = l; i <= r; ++i) add(i);
    ans[q[1].id] = res;

    // 开始转移
    for (int i = 2; i <= ord; ++i)
    {
        for (int j = q[i - 1].t + 1; j <= q[i].t; ++j) chg(c[j].p, c[
            j].v);
        for (int j = q[i - 1].t; j > q[i].t; --j) chg(c[j].p, c[j].o)
            ;
        while (r < q[i].r) add(++r);
        while (r > q[i].r) del(r--);
        while (l < q[i].l) del(l++);
        while (l > q[i].l) add(--l);
        ans[q[i].id] = res;
    }
    for (int i = 1; i <= ord; ++i) cout << ans[i] << '\n';
    return;
}

int main()
{
    ios::sync_with_stdio(0);
    cin.tie(0);
    cout.tie(0);
    int T = 1;
    // cin >> T;
    while (T--) solve();
    return 0;
}
```

9.3 整体二分

```cpp
/*****************************************************************
* 时间复杂度: 框架O(qlogm)
* 说明:
* 1.对多个需要二分解决的询问同时二分
* 2.二分对象为答案值域，但也将询问序列分到两个值域区间中
* 3.对于区间[l,r)的check不能到达O(q)/O(m)，应只考虑[l,r)中的值或询问
* 4.注意分到右半区间的询问目标值要削减
* 5.注意值域区间和询问区间的开闭
* 6.注意必要时对元素值去重
*****************************************************************/
const int N = 300005;

struct Fenwick { /*带时间戳树状数组*/ }fen;
struct Discret { /*离散化*/ }D;

struct Q
{
    int l, r, k, id;
}q[N];

int n, m;
pair<int, int> a[N];
int ans[N];

void bis(int lef, int rig, int ql, int qr)
{
    if (lef == rig - 1)
    {
        for (int i = ql; i < qr; ++i) ans[q[i].id] = lef;
        return;
    }
    int mid = lef + rig >> 1;
    for (int i = lef; i < mid; ++i)
    {
        fen.add(a[i].second, 1);
    }
    queue<Q> q1, q2;
    for (int i = ql; i < qr; ++i)
    {
        int cnt = fen.rsum(q[i].l, q[i].r);
        if (cnt < q[i].k) q2.push({ q[i].l,q[i].r,q[i].k - cnt,q[i].
            id });
        else q1.push(q[i]);
    }
    int qm = ql + q1.size();
    for (int i = ql; i < qr; ++i)
    {
        if (q1.size()) q[i] = q1.front(), q1.pop();
        else q[i] = q2.front(), q2.pop();
    }
    fen.clear();
    bis(lef, mid, ql, qm);
    bis(mid, rig, qm, qr);
    return;
}

void solve()
{
    cin >> n >> m;
    fen.init(n);
    for (int i = 1; i <= n; ++i)
    {
        cin >> a[i].first;
        a[i].second = i;
        D.insert(a[i].first);
    }
    D.work();
    for (int i = 1; i <= n; ++i) a[i].first = D[a[i].first];
    sort(a + 1, a + 1 + n);
    for (int i = 1; i <= m; ++i)
    {
        cin >> q[i].l >> q[i].r >> q[i].k;
        q[i].id = i;
    }
    bis(1, n + 1, 1, m + 1);
    for (int i = 1; i <= m; ++i) cout << D.v[ans[i] - 1] << '\n';
    return;
}
```

9.4 离散化

```cpp
/*****************************************************************
* 时间复杂度: O(logn)
* 说明: 注意起始序号
*****************************************************************/
struct Discret
{
    vector<ll> v;
    void insert(ll val)
    {
        v.push_back(val);
        return;
    }
    void work()
    {
        sort(v.begin(), v.end());
        v.erase(unique(v.begin(), v.end()), v.end());
        return;
    }
    void clear()
    {
        v.clear();
        return;
    }
    ll operator[](ll val)
    {
        return lower_bound(v.begin(), v.end(), val) - v.begin();
    }
};
```

9.5 快速排序

```cpp
/*****************************************************************
* 时间复杂度: O(nlogn)
* 说明: 两倍常数，但跳过所有与基准相等的值
*****************************************************************/
const int N = 100005;

int n;
ll a[N];

int median(int x, int y, int z)
{
    if (a[x] > a[y] && a[z] > a[y]) return a[x] > a[z] ? z : x;
    else if (a[x] < a[y] && a[z] < a[y]) return a[x] < a[z] ? z : x;
    else return y;
}

void QuickSort(int lef, int rig)//[lef, rig]
{
    if (rig <= lef) return;
    int mid = lef + (rig - lef) / 2;
    int pivot = median(lef, mid, rig);
    swap(a[pivot], a[lef]);
    int lp = lef; //第一个等于基准的值
    for (int i = lef + 1; i <= rig; ++i)
    {
        if (a[i] < a[lef]) swap(a[i], a[++lp]);
    }
    swap(a[lef], a[lp]);
    int rp = lp; //最后一个等于基准的值
    for (int i = lp + 1; i <= rig; ++i)
    {
        if (a[i] == a[lp]) swap(a[i], a[++rp]);
    }
    QuickSort(lef, lp - 1);
    QuickSort(rp + 1, rig);
    return;
}
```

9.6 枚举集合

```cpp
/*****************************************************************
* 时间复杂度: O(枚举对象个数)
* 说明: 枚举子集、超集、固定大小集合
*****************************************************************/
struct EnumSet
{
    vector<int> subset(int x) // 枚举x的子集
    {
        vector<int> res;
        for (int i = x; i >= 1; i = (i - 1) & x) res.push_back(i);
        res.push_back(0);
        return res;
    }

    vector<int> kset(int b, int k) // 枚举b位大小为k的集合
    {
        vector<int> res;
```

```
        int now = (1 << k) - 1;
        while (now < (1 << b))
        {
            res.push_back(now);
            int lowbit = now & -now;
            int x = now + lowbit;
            int y = ((now & ~x) / lowbit) >> 1;
            now = x | y;
        }
        return res;
    }

    vector<int> superset(int x, int b) // 枚举x的b位超集
    {
        vector<int> res;
        for (int i = x; i < (1 << b); i = (i + 1) | x) res.push_back(
            i);
        return res;
    }
};
```

9.7 CDQ 分治 +CDQ 分治 = 多维偏序

```
/********************************************************************
* 时间复杂度：O(nlog^(d-1)n)
* 说明：
* 1.cdq注意事项详见[CDQ分治+数据结构=多维偏序]
* 2.n维偏序需要n-1层cdq
* 3.第i层cdq将整个区间按第i+1维归并排序，同时将第i维降为二进制，然后调用
    第i+1层cdq；第n-1层cdq递归将左右分别按第n维排序，再用双指针按照第n维
    大小归并，同时计算左部前n-2维全0元素对右部前n-2维全1元素的贡献
********************************************************************/
const int N = 100005;

struct Elem
{
    ll a, b, c;
    ll cnt, id;
    bool xtag;
    bool operator!=(const Elem& e) const
    {
        return a != e.a || b != e.b || c != e.c;
    }
}e[N], ee[N], eee[N];

int n, k, ans[N], res[N];

bool bya(const Elem& e1, const Elem& e2)
{
    if (e1.a == e2.a && e1.b == e2.b) return e1.c < e2.c;
    else if (e1.a == e2.a) return e1.b < e2.b;
    else return e1.a < e2.a;
}

void cdq2(int lef, int rig)
{
    if (lef == rig - 1) return;
    int mid = lef + rig >> 1;
    cdq2(lef, mid);
    cdq2(mid, rig);
    int p1 = lef, p2 = mid, now = lef;
    int sum = 0;
    while (now < rig)
    {
        //左半部分xtag为0的可以贡献右半部分xtag为1的
        if (p2 == rig || p1 < mid && ee[p1].c <= ee[p2].c)
        {
            eee[now] = ee[p1++];
            sum += eee[now].cnt * (eee[now].xtag == 0);
        }
        else
        {
            eee[now] = ee[p2++];
            res[eee[now].id] += sum * (eee[now].xtag == 1);
        }
        now++;
    }
    for (int i = lef; i < rig; ++i) ee[i] = eee[i];
    return;
}

void cdq1(int lef, int rig)
{
    if (lef == rig - 1) return;
    int mid = lef + rig >> 1;
    cdq1(lef, mid);
    cdq1(mid, rig);
    int p1 = lef, p2 = mid, now = lef;
    while (now < rig)
    {
        if (p2 == rig || p1 < mid && e[p1].b <= e[p2].b)
        {
            ee[now] = e[p1++];
            ee[now].xtag = 0;
        }
        else
        {
            ee[now] = e[p2++];
            ee[now].xtag = 1;
        }
        now++;
    }
    for (int i = lef; i < rig; ++i) e[i] = ee[i];
    cdq2(lef, rig);
    return;
}

void solve()
{
    cin >> n >> k;
    vector<Elem> ori(n);
    for (int i = 0; i < n; ++i)
    {
        cin >> ori[i].a >> ori[i].b >> ori[i].c;
        ori[i].cnt = 1;
    }
    sort(ori.begin(), ori.end(), bya);
    int cnt = 0;
    for (auto& x : ori)
    {
        if (cnt == 0 || e[cnt] != x) cnt++, e[cnt] = x, e[cnt].id =
            cnt;
        else e[cnt].cnt++;
    }
    cdq1(1, cnt + 1);
    for (int i = 1; i <= cnt; ++i)
    {
        res[e[i].id] += e[i].cnt - 1;
        ans[res[e[i].id]] += e[i].cnt;
    }
    for (int i = 0; i < n; ++i) cout << ans[i] << '\n';
    return;
}
```

9.8 CDQ 分治 + 数据结构 = 多维偏序

```
/********************************************************************
* 时间复杂度：O(nlog^(d-1)n)
* 说明：
* 1.每降一维需要乘O(logn)时间
* 2.适用于高维偏序等小元素对大元素有贡献的问题
* 3.元素需要提前去重
* 4.注意小于等于和小于做法不同，如分治顺序与排序复原/mid的移动
* 5.贡献有顺序要求如dp时，先左再合并再右
* 6.有时需要离散化才能利用数据结构
********************************************************************/
const int N = 100005;

struct Fenwick { /*带时间戳最大值树状数组*/ }fen;
struct Discret { /*离散化*/ }D;

struct Elem
{
    ll a, b, c;
    ll w, dp;
    bool operator!=(const Elem& e) const { return a != e.a || b != e.
        b || c != e.c; }
} e[N];

int n;

bool bya(const Elem& e1, const Elem& e2)
{
    if (e1.a == e2.a && e1.b == e2.b) return e1.c < e2.c;
    else if (e1.a == e2.a) return e1.b < e2.b;
    else return e1.a < e2.a;
}

bool byb(const Elem& e1, const Elem& e2)
{
    if (e1.b == e2.b) return e1.c < e2.c;
    else return e1.b < e2.b;
```

```
}

void cdq(int lef, int rig)
{
    if (e[lef].a == e[rig - 1].a) return;
    int mid = lef + (rig - lef) / 2;

    // 需要保证e[mid-1].a和e[mid].a不同
    if (e[lef].a == e[mid].a)
    {
        while (e[lef].a == e[mid].a) mid++;
    }
    else
    {
        while (e[mid - 1].a == e[mid].a) mid--;
    }

    // 解决左半
    cdq(lef, mid);

    // 解决合并
    sort(e + lef, e + mid, byb);
    sort(e + mid, e + rig, byb);
    int p1 = lef, p2 = mid;
    while (p2 < rig)
    {
        while (p1 < mid && e[p1].b < e[p2].b)
        {
            fen.add(D[e[p1].c], e[p1].dp);
            p1++;
        }
        e[p2].dp = max(e[p2].dp, e[p2].w + fen.pres(D[e[p2].c] - 1));
        p2++;
    }
    fen.clear();

    // 解决右半
    sort(e + mid, e + rig, bya); // 复原排序
    cdq(mid, rig);
    return;
}

void solve()
{
    cin >> n;
    vector<Elem> ori(n);
    for (int i = 0; i < n; ++i)
    {
        cin >> ori[i].a >> ori[i].b >> ori[i].c >> ori[i].w;
        ori[i].dp = ori[i].w;
        D.insert(ori[i].c);
    }
    D.work();
    fen.init(D.v.size());
    sort(ori.begin(), ori.end(), bya);
    int cnt = 0;
    for (auto& x : ori)
    {
        if (cnt == 0 || e[cnt] != x) e[++cnt] = x;
        else e[cnt].dp = e[cnt].w = max(e[cnt].w, x.w);
    }
    cdq(1, cnt + 1);
    ll ans = 0;
    for (int i = 1; i <= cnt; ++i) ans = max(ans, e[i].dp);
    cout << ans << '\n';
    return;
}
```

10 博弈论

10.1 Fibonacci 博弈

```
/**********************************************************************
* 时间复杂度: O(1)
* 说明:
* 1.有一堆石子，两人轮流取。先手第一次不能直接取完。每次至少取一个，
* 但最多取上一个人的两倍。取走最后一个石子的人获胜。
* 2.结论: 是斐波那契数则先手必败，否则先手必胜。
**********************************************************************/
bool Fibonacci(ll x) //返回先手是否必胜
{
    ll a = 1, b = 1;
    while (max(a, b) <= x)
    {
        if (a < b) a += b;
        else b += a;
        if (max(a, b) == x) return 0;
    }
    return 1;
}
```

10.2 Wythoff 博弈

```
/**********************************************************************
* 时间复杂度: O(1)
* 说明:
* 1.有两堆石子，两人轮流取。每次可以在一堆中取任意个石子或在两堆中取同
* 样多的任意个石子，取走最后一个石子的人获胜。
* 2.若x和y极大则需要注意精度问题。
**********************************************************************/
bool Wythoff(ll x, ll y) //返回先手是否必胜
{
    const double K = ((1.0 + sqrt(5.0)) / 2.0);
    ll res = abs(x - y) * K;
    return res != min(x, y);
}
```

10.3 Green Hackenbush 博弈

```
/**********************************************************************
* 时间复杂度: O(n)
* 说明:
* 1.有一棵有根树，两人轮流选择一个子树删除，删除根结点的人失败。
* 2.有一颗有根树，两人轮流删除一条边以及不与根相连的部分，无边可删
* 的人失败。
* 3.结论: 以边为对象: 叶结点父边sg值为1，中间结点父边sg值为所有邻接
* 边sg值异或和+1; 以点为对象: 叶结点sg值为0，其他结点sg值为所有邻接
* 点sg值+1的异或和。
**********************************************************************/
void dfs(int x, int fa)
{
    sg[x] = 0;
    for (auto e : node[x])
    {
        if (e == fa) continue;
        dfs(e, x);
        sg[x] ^= sg[e] + 1;
    }
    return;
}
```