Recitation 10: Extra Credit Overview

COP3402 FALL 2015 — ARYA POURTABATABAIE

FROM EURIPIDES MONTAGNE, FALL 2014

PL/0 Program

```
var f, n;←
                          Global variables used to send and receive
procedure fact;
                          data from procedure.
var ans1;
begin
       ans1:=n;
                                          Local variables needed to
       n := n-1;
                                          keep temporal calculations
       if n = 0 then f := 1;
                                          and values.
       if n > 0 then call fact;
       f:=f*ans1;
end;
begin
n := 3;
call fact;
write f;
end.
```

Extended PL/O Program

```
var num;
                                             Parameter Passing
procedure factorial(x); \leftarrow
begin
        if x = 0 then return := 1
        else return := x * call factorial(x-1);
 end;
begin
                                  Functional Value return
num_:= call factorial(3);
write num;
end.
                                                 Parameter Passing
                  Functional Value return
```

Decouples the output of a procedure.

Uses the Functional Value (FV) slot in the activation record to return a value.

When parsing the procedure declaration, we insert a variable called "return" with level+1 and address=0.

FV is accessed through "return".

o return := 5;

```
var x,y;
procedure foo;
var w;
begin
       w := x;
        w := w * w * w;
        return := w;
 end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

```
05 LOD 0 4
 06 LOD 0 4
 07 OPR 0 MUL
 08 LOD 0 4
 09 OPR 0 MUL
 10 STO 0 4
 11 LOD 0 4
→12 STO 0 0
 13 OPR 0 RET
 . . .
```

```
var x,y;
procedure foo;
var w;
begin
       w := x;
        w := w*w*w;
        return := w;
end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

| stack | |
|-------|-----|
| 0 | FV |
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | X |
| 5 | У |
| 6 | 25 |
| 7 | FV |
| 8 | SL |
| 9 | DL |
| 10 | RA |
| 11 | W |
| 12 | 125 |

```
var x,y;
procedure foo;
var w;
begin
       w := x;
        w := w*w*w;
        return := w;
end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

```
05 LOD 0 4
06 LOD 0 4
07 OPR 0 MUL
08 LOD 0 4
09 OPR 0 MUL
10 STO 0 4
11 LOD 0 4
12 STO 0 0
13 OPR 0 RET
```

| stack | |
|-------|-------|
| 0 | FV |
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | X |
| 5 | У |
| 6 | 25 |
| 7 | FV |
| 8 | SL |
| 9 | DL |
| 10 | RA |
| 11 | w=125 |

```
var x,y;
procedure foo;
var w;
begin
       w := x;
        w := w*w*w;
        return := w;
end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

| stack | |
|-------|-------|
| 0 | FV |
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | Х |
| 5 | У |
| 6 | 25 |
| 7 | FV |
| 8 | SL |
| 9 | DL |
| 10 | RA |
| 11 | w=125 |
| 12 | 125 |

```
var x,y;
procedure foo;
var w;
begin
       w := x;
        w := w*w*w;
        return := w;
end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

```
05 LOD 0 4
06 LOD 0 4
07 OPR 0 MUL
08 LOD 0 4
09 OPR 0 MUL
10 STO 0 4
11 LOD 0 4
12 STO 0 0
13 OPR 0 RET
```

| stack | |
|-------|--------|
| 0 | FV |
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | Х |
| 5 | У |
| 6 | 25 |
| 7 | FV=125 |
| 8 | SL |
| 9 | DL |
| 10 | RA |
| 11 | w=125 |

```
var x,y;
procedure foo;
var w;
begin
       w := x;
        w := w*w*w;
        return := w;
end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

| stack | |
|-------|----|
| 0 | FV |
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | Х |
| 5 | У |
| 6 | 25 |

The FV is still in the stack, but we must recover it before using it.

Two types of calls:

- Do not use FV (discard it).
- Use FV (recover it).

Recover FV

This happens when we call the function inside an expression.

The FV is one position over the SP.

Just increment the SP by 1 (INC 0 1).

```
var x,y;
procedure foo;
var w;
begin
        w := x;
        w := w*w*w;
        return := w;
end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

```
12 STO 0 0
   OPR 0 RET
16 LIT 0 25
  CAL 0 1
18 INC 0 1
19 OPR 0 SUM
20 LIT 0 10
21 OPR 0 DIV
22 STO 0 5
  LOD 0 5
24 SIO 0 1
25 OPR 0 RET
```

| stack | |
|-------|----|
| 0 | FV |
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | Х |
| 5 | У |
| 6 | 25 |

```
var x, y;
procedure foo;
var w;
begin
        w := x;
        w := w*w*w;
        return := w;
 end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

```
12 STO 0 0
13 OPR 0 RET
16 LIT 0 25
   CAL 0
  INC 0 1
19 OPR 0 SUM
21 OPR 0 DIV
22 STO 0 5
  LOD 0 5
24 SIO 0 1
25 OPR 0 RET
```

| stack | |
|-------|----|
| 0 | FV |
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | X |
| 5 | У |
| 6 | 25 |

```
var x,y;
procedure foo;
var w;
begin
       w := x;
        w := w*w*w;
        return := w;
end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

```
12 STO 0 0
13 OPR 0 RET
16 LIT 0 25
17 CAL 0 1
19 OPR 0 SUM
21 OPR 0 DIV
22 STO 0 5
  LOD 0 5
24 SIO 0 1
25 OPR 0 RET
```

stack 0 FV 1 SL 2 DL 3 RA 4 x 5 y 6 25 7 FV=125

```
var x, y;
procedure foo;
var w;
begin
        w := x;
        w := w*w*w;
        return := w;
 end;
begin
x := 5;
y := 25 + call foo / 10;
write y;
end.
```

```
12 STO 0 0
13 OPR 0 RET
16 LIT 0 25
17 CAL 0 1
18 INC 0 1
21 OPR 0 DIV
22 STO 0 5
  LOD 0 5
24 SIO 0 1
25 OPR 0 RET
```

| stack | |
|-------|-----|
| 0 | FV |
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | X |
| 5 | У |
| 6 | 150 |

Recover FV

To use function calls inside expressions, we must treat them as factors.

factor ::= ident | number | "(" expression ")" | "call" ident.

Differentiate call types

When used as factor, we must generate INC.

```
factor ::= ident | number | "(" expression ")" | "call" ident.
```

When used as statement, we do nothing else.

```
statement ::= [ ...
```

| "call" ident parameter-list

...

Parameter Passing

Decouples the input of a procedure.

We treat parameter exactly as variables, with one distinction:

• The value of a parameter is set before we call the procedure.

We must add parameter slots in the AR, before the variable slots.

| 0 | FV |
|---|--------|
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | params |
| X | ••• |
| X | vars |
| X | |

How Parameter Passing works?

Before the call, we must set the value of the parameters.

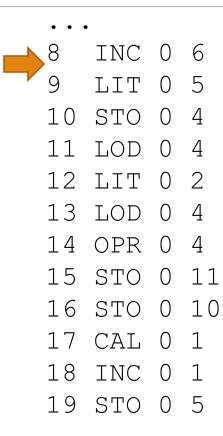
Each parameter is an expression, so first solve all expressions.

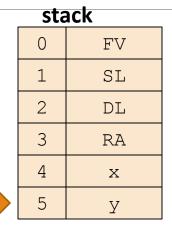
We must predict the position of parameters.

Copy expressions results to parameter slots.

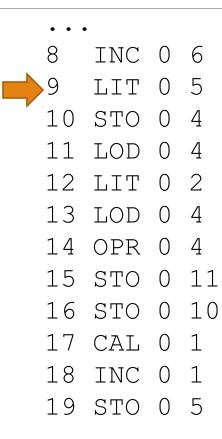
Then make the call.

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```



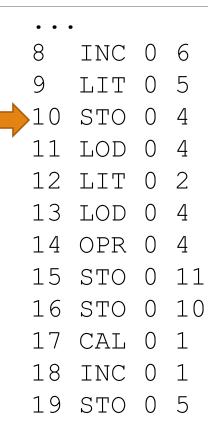


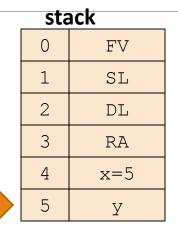
```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```



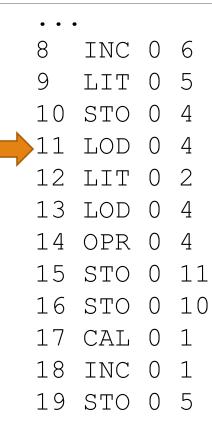


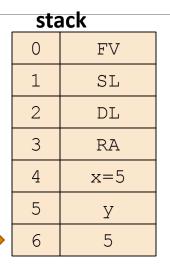
```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```





```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```





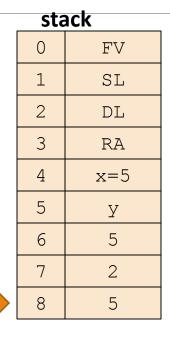
```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

```
INC 0 6
   LIT 0 5
10 STO 0 4
11 LOD 0 4
12 LIT 0 2
13 LOD 0 4
14 OPR 0 4
15 STO 0 11
16 STO 0 10
17 CAL 0 1
18 INC 0 1
19 STO 0 5
```

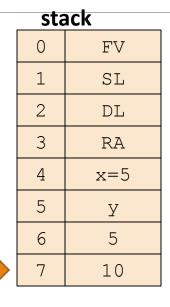


```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

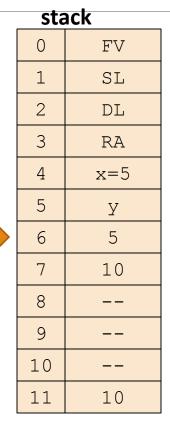
```
INC 0 6
   LIT 0 5
10 STO 0 4
11 LOD 0 4
12 LIT 0 2
13 LOD 0 4
14 OPR 0 4
15 STO 0 11
16 STO 0 10
17 CAL 0 1
18 INC 0 1
19 STO 0 5
```



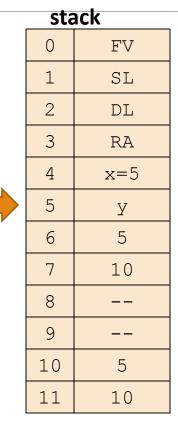
```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```



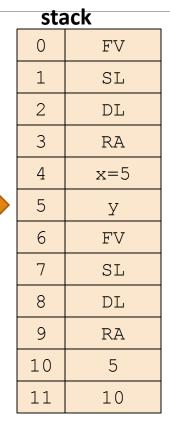
```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```



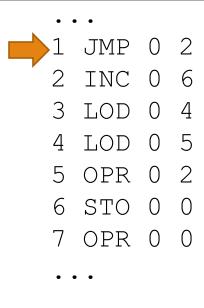
```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

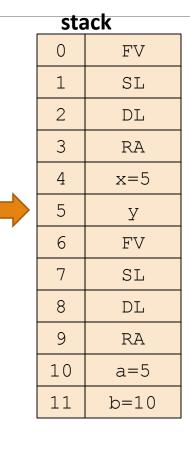


```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```



```
var x, y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```





```
var x, y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

stack FV SL DL 3 RA x=5FV SL DLRA 10 a=5b = 10

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

stack FV SL DL 3 RA x=5FV SL DLRA 10 a=5b = 105

```
var x, y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

stack FV SL DLRA x=5FV SL DLRA 10 a=5b = 1010 5 10

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

stack FV SL DL 3 RA x=5FV SL DLRA 10 a=5b = 1015 10

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

stack FV SL DL3 RA x=5FV=15 SL DLRA 10 a=5b = 1015 10

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

stack FV SL DLRA x=5FV=15 SL DLRA 10 a=5b = 1015 10 11 10

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

```
INC 0 6
   LIT 0 5
10 STO 0 4
  LOD 0 4
12 LIT 0 2
13 LOD 0 4
14 OPR 0 4
15 STO 0 11
16 STO 0 10
   CAL 0 1
   INC 0 1
19 STO 0 5
  LOD 0 5
  SIO 0 1
22 OPR 0 0
```

stack FV SL DLRA x=5FV=15 SL 8 DLRA 10 a=5b = 1011 10 15 10

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

| 0 | FV |
|----|-------|
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | x=5 |
| 5 | У |
| 6 | FV=15 |
| 7 | SL |
| 8 | DL |
| 9 | RA |
| 10 | a=5 |
| 11 | b=10 |
| 10 | 15 |
| 11 | 10 |

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

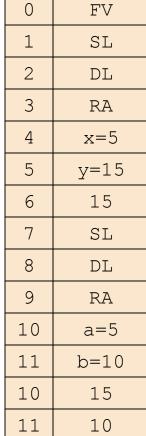
```
INC 0 6
   LIT 0 5
10 STO 0 4
  LOD 0 4
12 LIT 0 2
13 LOD 0 4
14 OPR 0 4
15 STO 0 11
16 STO 0 10
17 CAL 0 1
  INC 0 1
19 STO 0 5
20 LOD 0 5
21 SIO 0 1
22 OPR 0 0
```

stack FV SL DLRA x=5y = 15FV=15 SL 8 DLRA 10 a=5b = 1010 15 10

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

| 0 | FV |
|----|------|
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | x=5 |
| 5 | y=15 |
| 6 | 15 |
| 7 | SL |
| 8 | DL |
| 9 | RA |
| 10 | a=5 |
| 11 | b=10 |
| 10 | 15 |
| 11 | 10 |

```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```



```
var x,y;
procedure sum(a,b);
begin
      return := a+b;
end;
begin
x := 5;
y := call sum(x, 2*x);
write y;
end.
```

OPR 0 0



| 0 | FV |
|----|------|
| 1 | SL |
| 2 | DL |
| 3 | RA |
| 4 | x=5 |
| 5 | y=15 |
| 6 | 15 |
| 7 | SL |
| 8 | DL |
| 9 | RA |
| 10 | a=5 |
| 11 | b=10 |
| 10 | 15 |
| 11 | 10 |
| | |

PL/0 Extended grammar

(Not what you'll actually use, but close)

```
procedure-declaration ::= { "procedure" ident parameter-block ";" block ";" }
parameter-block ::= "(" [ ident { "," ident } ] ")".
parameter-list ::= " (" [ expression { "," expression } ] ")".
statement ::= [ ... | "call" ident parameter-list | ... ].
factor ::= ident | number | "(" expression ")" | "call" ident parameter-list.
```

Parsing the parameter block

In the procedure declaration, we found the parameter block.

For each ident found, insert a new symbol in the symbol table with level+1 and corresponding address.

The address of the first parameter is 4 (size of basic AR).

Each subsequent parameter gets the next address (5, 6, etc).

parameter-block()

```
parameter-block ::= "(" [ ident { "," ident } ] ")".
```

```
procedure parameter-block();
begin
addr := 4;
if(token <> "(") then ERROR("Procedure must have parameters");
get token();
 if(token = ident) then begin
          enter(ident, level+1, addr); addr++;
          get token();
          while (token = ",") begin
                    get token();
                    if(token <> ident) then ERROR();
                    enter(ident, level+1, addr); addr++;
                    get token();
          end;
 end;
if token <> ")" then ERROR("Bad procedure declaration");
get token();
end;
```

Address of variables

The address of the first variable is address of last parameter plus 1.

Each subsequent variable gets the next address.

We must somehow pass the address of the last parameter to function vardecl().

Don't forget to insert variable "return".

Parsing the parameter list

Each time we have a call, we found the parameter list.

We must generate code for solving each expression.

Then generate code to store results in parameter slots.

parameter-list()

parameter-list ::= " (" [expression { "," expression }] ")".

```
procedure parameter-list();
begin
 params = 0;
if(token <> "(") then ERROR("Missing parameter list at call");
get token();
 if(token <> ")") then begin
           expression();
           params++;
 end;
 while (token = ",") begin
           get token();
           expression();
           param++;
 end;
 while(param > 0) { // Save results into param slots
           gen(STO, 0, stack size+4-1);
           param--;
if (token <> ")") then ERROR("Bad calling formating.");
get token();
end;
```

What is stack_size?

In order to predict the position of the parameter slots, we need to know the size of the stack at the moment we do the call.

The stack only grows or shrinks for certain instructions.

We detect when those instructions are issued, and increment or decrement the stack size accordingly.

updateStackSize()

```
int gen(int op, int l, int m) {
        code[code size].op = op;
        code[code size].l = 1;
        code[code_size].m = m;
        updateStackSize(op, 1, m);
        return code_size++;
void updateStackSize(int op, int l, int m) {
if(op == LIT, LOD, READ) {
        stack size++;
}else if (op == STO, JPC, WRITE) {
        stack size--;
}else if (op == INC) {
        stack size += m;
}else if (op == OPR) {
        if(m == RET) {stack size = 0;}
        else if (m != NEG, ODD) {stack size--;}
```

Questions?