

Introduction to Databases

Lecture 4:

Structured Query Language

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SQL

- How to retrieve data from the database?
 - Lingua franca = SQL
- SQL is the standard
 - Declarative language
 - Supported by all relational database systems

50 year of research on how to optimize

- Provably hard problem
- We will cover a representative subset of SQL
 - See also <http://www.w3schools.com/sql/>

SQL Syntax

- Declarative: say WHAT you want, not how to compute it
- Basic query form:

SELECT [distinct] [list of attributes L]

FROM [tables T]

WHERE

[conditions C]

(With distinct) equal to:
 $\pi_L \sigma_C (T_1 \times T_2 \times \dots \times T_k)$
Result is returned ordered.
If order by is absent, the
order is undefined.

ORDER BY [list of attributes] (**ASCENDING/DESCENDING**);

- Make any combination of tuples in T. For those that satisfy C, return L.
 - **Caveat: keeps duplicates unless **distinct** is used!**

Examples

- Works(name,salary)

Name	salary
Jan	1000
An	2000
George	2200

SELECT name
FROM works
WHERE salary > 1000;

name
An
George

SELECT **distinct** e1.name **AS** name
FROM works e1, works e2
WHERE e1.salary < e2.salary
ORDER BY name;

e1.name	e1.salary	e2.name	e2.salary
Jan	1000	Jan	1000
Jan	1000	An	2000
Jan	1000	George	2200
An	2000	Jan	1000
An	2000	An	2000
An	2000	George	2200
George	2200	Jan	1000
George	2200	An	2000
George	2200	George	2200

name
An
Jan
Jan

distinct
→

name
An
Jan

Examples

- Works(name,salary)
- Manages(m_name,e_name) // *m_name manages e_name*
- Ask for employees earning more than their manager

```
SELECT distinct e.name  
FROM works e, works m, manages man  
WHERE  
    e.name=man.e_name AND  
    m.name=man.m_name AND  
    e.salary > m.salary;
```

Syntax

- **SELECT** [list of attributes/expressions] or **SELECT ***
 - Attribute = *[Relation name].[Attribute name]* or *[Attribute name]* if there is no ambiguity
 - Expressions: A+B, 2C, A **AS** B
- **FROM** [list of tables]
 - Table = *[Table name]* or *[Table name] AS [alternative name]* if disambiguation is needed
 - E.g., FROM works e FROM works as e
 - Alternative name becomes the relation name of the copy
- **WHERE** [condition]
 - Condition = Boolean (AND, OR, NOT) combination of atoms
 - Atom = [Attribute or constant] OPER [Attribute or constant]
 - OPER can be =, <>, <, <=, >=, >

Exercises

Employee(ename, ecity)

Works(name,company)

Manages(mname,ename)

Located(company,city)

- Give all employees who live in their company's city.
- Give all employees who live in the same city as their managers.
- Give all pairs of companies that have employees living in the same city.
- Give all employees of a company in Brussels that are managed by a manager who lives in Antwerp.
- Give all employees who work for at least two companies.

Employee(ename, ecity)

Works(name,company)

Manages(mname,ename)

Located(company,city):

- Give all employees who live in their company's city.

```
SELECT E.ename
FROM Employee E, Works W, Located L
WHERE E.ename=W.name AND
        W.company=L.company AND
        E.ecity=L.city;
```


- Give all employees who live in the same city as their managers.

Employee(ename, ecity)
Works(name, company)
Manages(mname, ename)
located(company, city)

Need two
Employee tables!

```
SELECT E1.ename  
FROM Employee E1, Employee E2, Manages M  
WHERE E1.city=E2.city AND  
        E1.ename=M.ename AND  
        E2.ename=M.mname;
```

- Give all pairs of companies that have employees living in the same city

Employee(ename, ecity)

Works(name,company)

Manages(mname,ename)

Located(company,city)

```
SELECT W1.company, W2.company
FROM Works W1, Works W2, Employee E1, Employee E2
WHERE W1.name=E1.name AND
        W2.name=E2.name AND
        E1.ecity=E2.ecity;
```

- Give all employees of a company in Brussels that are managed by a manager who lives in Antwerp

Employee(ename, ecity)
Works(name,company)
Manages(mname,ename)
Located(company,city)

```
SELECT E1.ename  
FROM Employee E1, Employee E2, Manages M, Works W, Located L  
WHERE E1.ename=W.name AND W.company=L.company AND L.city="Brussels"  
        AND E1.ename=M.ename AND M.mname=E2.ename AND E2.ecity="Antwerp"
```

- Give all employees who work for at least two companies

Employee(ename, ecity)
Works(name,company)
Manages(mname,ename)
Located(company,city)

```
SELECT W1.name  
FROM Works W1, Works W2  
WHERE W1.name=W2.name AND W1.company <> W2.company
```

Other Basic Constructions: Set Operations

- Result of a query is a relation (with duplicates)
- If two queries are compatible (same schema, same order of attributes), then we can use set operations to combine them
 - **UNION**
 - **INTERSECT**
 - **EXCEPT** (MINUS in many systems)
- These operations remove duplicates. If you want duplicates:
 - **UNION ALL**
 - **INTERSECT ALL**
 - **EXCEPT ALL**

Examples: Set Operations

**Select city from members
where age>40;**

city
Antwerp
Antwerp
Gent
Gent

INTERSECT

=

city
Antwerp
Gent

**Select city from members
where salary<60;**

city
Antwerp
Antwerp
Brussels
Gent

name	city	salary	age
John	Antwerp	30	43
Mike	Antwerp	50	44
Ella	Brussels	40	23
Mia	Gent	20	42
Mike	Gent	60	60

Examples: Set Operations

**Select city from members
where age>40;**

city
Antwerp
Antwerp
Gent
Gent

INTERSECT ALL

=

city
Antwerp
Antwerp
Gent

**Select city from members
where salary<60;**

city
Antwerp
Antwerp
Brussels
Gent

name	city	salary	age
John	Antwerp	30	43
Mike	Antwerp	50	44
Ella	Brussels	40	23
Mia	Gent	20	42
Mike	Gent	60	60

Examples: Set Operations

**Select city from members
where age>40;**

city
Antwerp
Antwerp
Gent
Gent

UNION

=

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Brussels
Gent

**Select city from members
where salary<60;**

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Antwerp
Antwerp
Brussels
Gent

name	city	salary	age
John	Antwerp	30	43
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Mike	Gent	60	60

Examples: Set Operations

**Select city from members
where age>40;**

city
Antwerp
Antwerp
Gent
Gent

UNION ALL

=

**Select city from members
where salary<60;**

city
Antwerp
Antwerp
Brussels
Gent

city
Antwerp
Antwerp
Antwerp
Antwerp
Brussels
Gent
Gent
Gent

name	city	salary	age
John	Antwerp	30	43
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Mike	Gent	60	60

Examples: Set Operations

**Select city from members
where age>40;**

city
Antwerp
Antwerp
Gent
Gent

EXCEPT

=

city

**Select city from members
where salary<60;**

city
Antwerp
Antwerp
Brussels
Gent

name	city	salary	age
John	Antwerp	30	43
Mike	Antwerp	50	44
Ella	Brussels	40	23
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Mike	Gent	60	60

Examples: Set Operations

**Select city from members
where age>40;**

city
Antwerp
Antwerp
Gent
Gent

EXCEPT ALL

=

city
Gent

**Select city from members
where salary<60;**

city
Antwerp
Antwerp
Brussels
Gent

name	city	salary	age
John	Antwerp	30	43
Mike	Antwerp	50	44
Ella	Brussels	40	23
Mia	Gent	20	42
Mike	Gent	60	60

Other Basic Constructions: Joins

- In FROM:
 - R **JOIN** S:
 - table over the attributes of $R \cup S$, without duplicates, made out of all “compatible” pairs of tuples
 - If attribute names are different, can use:
R **JOIN** S on [condition]

```
SELECT distinct e1.name AS name  
FROM works e1 JOIN works e2 ON e1.name=e2.name  
WHERE e1.company <> e2.company;
```

Subqueries

- A subquery is a query inside another query
 - Everywhere a relation is expected, we can use a subquery instead
 - Everywhere a value is expected, we can use a subquery producing one result
- We can also use subqueries in combination with operators EXISTS, ANY, and ALL
- Subqueries can use attributes from the outer query
 - Called *correlated subquery*

Other Basic Constructions: Comparing with sets

- In Where:
 - **EXISTS** (query): true if query produces an answer
 - [Attribute or constant] OPER **ANY** (query Q): true if there is an answer A in Q such that [Attribute or constant] OPER A holds
E.g.: SELECT e.name FROM works e WHERE e.salary > **ANY** (SELECT salary FROM works)
 - [Attribute or constant] OPER **ALL** (query Q): true if for all answers A in Q [Attribute or constant] OPER A holds
E.g.: SELECT e.name FROM works e WHERE e.salary > **ALL** (SELECT salary FROM works)
 - **IN**: same as “= ANY”

NOT EXISTS

- Until now: all queries of the type:
 - If A and B and C and D, then: report E
- What if we want to output E if something is missing?
 - Employee NOT working for “UAntwerpen”
 - Employee NOT earning more than any of his/her managers?

NOT EXISTS

- Employee NOT working for “UAntwerpen”
- What’s wrong with:

```
SELECT e.name  
FROM works e  
WHERE e.company <> “UAntwerpen”;
```

Employee(name, city)

Works(name,company)

Manages(m_name,e_name)

Located(company,city)

NOT EXISTS

Employee(name, city)

Works(name,company)

Manages(m_name,e_name)

Located(company,city)

- Manager NOT earning less than any of his/her employees?
- What's wrong with:

SELECT m.name

FROM works e, works m, manages man

WHERE

e.name=man.e_name **AND**

m.name=man.m_name **AND**

e.salary <= m.salary;

NOT EXISTS

- We need “NOT EXISTS” in **WHERE** clause
- Syntax:
 - (NOT) EXISTS (SQL Query Q)
 - Q can reuse attribute names of the outer query
 - Is “TRUE” if Q does (not) return any answer

Correct Answers: NOT EXISTS

- Employee NOT working for “UAntwerpen”

Employee(name, city)

Works(name,company)

Manages(m_name,e_name)

Located(company,city)

SELECT e.name

FROM works e

WHERE

NOT EXISTS (

SELECT *

FROM works W

WHERE W.name=e.name AND W.company =“UAntwerpen”

);

Correct Answers: NOT EXIST

Employee(name, city)

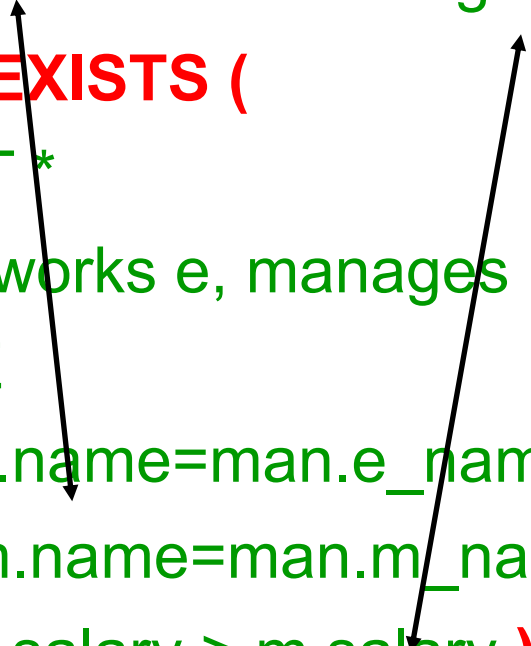
Works(name,company)

Manages(m_name,e_name)

Located(company,city)

- Manager NOT earning less than any of his/her employees?

```
SELECT m.name FROM Manages m
WHERE NOT EXISTS (
    SELECT *
    FROM works e, manages man
    WHERE
        e.name=man.e_name AND
        m.name=man.m_name AND
        e.salary > m.salary );
```



Examples: Subqueries

Serves(bar, beer) Assumption: All the bars in Serves are also bars in Visits and vice versa.
Visits(drinker, bar)

Select all drinkers that visit some bar where only Duvel is served:
(same as saying: select all drinkers that visit some bar where Duvel
is served and *no other* beer is served):

```
SELECT bar
FROM Serves
WHERE beer="Duvel" AND
      bar NOT IN
      (SELECT bar FROM Serves
       WHERE beer <> "Duvel");
```

```
SELECT bar
FROM Serves S1
WHERE beer="Duvel" AND
      NOT EXISTS
      (SELECT bar FROM Serves S2
       WHERE S1.bar=S2.bar AND
       S2.beer<> "Duvel");
```

Examples: Subqueries

`Serves(bar, beer)` Assumption: All the bars in Serves are also bars in Visits and vice versa.
`Visits(drinker, bar)`

Select all drinkers that visit some bar where only Duvel is served:
(same as saying: select all drinkers that visit some bar where Duvel
is served and *no other* beer is served):

Examples: Subqueries

Serves(bar, beer) Assumption: all the bars in Serves are also bars in Visits and vice versa.
Visits(drinker, bar)

Select all drinkers that visit some bar where only Duvel is served:
(same as saying: select all drinkers that visit some bar where Duvel
is served and *no other* beer is served):

```
SELECT DISTINCT Visits.drinker  
FROM Visits, Serves  
WHERE Visits.bar = Serves.bar AND Serves.beer = 'Duvel'  
      AND Serves.bar NOT IN  
        (SELECT Bar1.bar  
         FROM Serves Bar1  
         WHERE Bar1.beer <> 'Duvel');
```

```
SELECT DISTINCT drinker  
FROM Visits  
WHERE bar IN  
      (SELECT bar  
       FROM Serves  
       WHERE beer = 'Duvel')  
AND  
bar NOT IN  
  (SELECT bar  
   FROM Serves  
   WHERE beer <> 'Duvel');
```

Can also be done using **EXCEPT**

Views

- CREATE VIEW: Make a virtual table

```
CREATE VIEW temp AS
```

```
    SELECT bar
```

```
    FROM Serves
```

```
    WHERE beer <> 'Duvel';
```

```
SELECT DISTINCT Visits.drinker
```

```
FROM Visits, Serves
```

```
WHERE Visits.bar = Serves.bar
```

```
    AND Serves.beer = 'Duvel'
```

```
    AND Serves.bar NOT IN temp;
```

```
SELECT DISTINCT Visits.drinker
```

```
FROM Visits, Serves
```

```
WHERE Visits.bar = Serves.bar
```

```
    AND Serves.beer = 'Duvel'
```

```
    AND Serves.bar NOT IN
```

```
        (SELECT Bar1.bar
```

```
        FROM Serves Bar1
```

```
        WHERE Bar1.beer <> 'Duvel');
```


Band_member(bname,mname)

Concert(bname,date,location)

Exercises

- Give all bands that have never concertcd in “Wembley Stadium”
- Give all bands that have concertcd in any of the places in which “Queen” concertcd.
- Give all bands that have never concertcd on the same day in the same location as “Ramstein”
- Give all bands that have at least one member in common with the band “Abba”
- Give all bands of which “Mick Jagger” is a member and have no members in common with a band of which “John Lennon” is a member

Exercises

Band_member(bname,mname)

Concert(bname,date,location)

- Give all bands that have never concertted in “Wembley Stadium”

```
SELECT bname
FROM Band_member B
WHERE NOT EXISTS (
    SELECT * FROM Concert C
    WHERE C.bname=B.bname AND C.location=“WS”);
```

Exercises

Band_member(bname,mname)

Concert(bname,date,location)

- Give all bands that have concerted in any of the places in which “Queen” concerted.

```
SELECT B.bname  
FROM Band_member B  
WHERE EXISTS (  
    SELECT * FROM Concert C1, Concert C2  
    WHERE C1.bname=B.bname AND C2.bname=“Queen” AND C1.location=C2.location  
);
```

Exercises

Band_member(bname,mname)

Concert(bname,date,location)

- Give all bands that have never concertated on the same day in the same location as “Ramstein”

```
SELECT B.bname FROM Band_member B
WHERE NOT EXISTS (
    SELECT * FROM Concert C1, Concert C2
    WHERE C1.bname=B.bname AND C2.bname=“Ramstein” AND
    C2.date=C1.date AND C2.location=C1.location);
```

Band_member(bname,mname)

Concert(bname,date,location)

Exercises

- Give all bands that have at least one member in common with the band “Abba”

```
SELECT B1.bname FROM Band_member B1
```

```
WHERE EXISTS (
```

```
    SELECT * FROM Band_member B2
```

```
    WHERE B2.bname=“Abba” AND B2.mname=B1.mname);
```

```
SELECT B1.bname FROM Band_member B1, Band_member B2
```

```
WHERE B2.bname=“Abba” AND B2.mname=B1.mname);
```

Band_member(bname,mname)

Concert(bname,date,location)

Exercises

- Give all bands of which “Mick Jagger” is a member and have no members in common with a band of which “John Lennon” is a member

```
SELECT B.bname FROM Band_member B
WHERE B.mname="Mick Jagger"
AND NOT EXISTS (
    SELECT * FROM Band_member B1, Band_member B2, Band_member B3
    WHERE B1.bname=B.bname AND B2.mname=B1.mname AND
    B3.bname=B2.bname AND B3.mname="John Lennon");
```

Advanced Functionalities: Aggregation

- Aggregation (avg(A), min(A), max(A), sum(A), count(A), count(distinct A), ...)
 - group by / having (optional)
 - **All non-aggregated attributes in SELECT must appear in GROUP BY**

```
SELECT w.company as company, avg(salary) as salary
FROM works w
GROUP BY Company
HAVING min(salary) > 1000
ORDER BY salary;
```

Advanced Functionalities: Aggregation

SELECT w.company **as** company, **avg**(salary) **as** salary

FROM works w

GROUP BY Company

HAVING min(salary) > 1000

ORDER BY salary;

SELECT w.company **as** company, salary
FROM works w

Divide into groups
by company attribute

For each group, compute avg(salary)
and min(salary)

Throw out groups not satisfying
HAVING

Output one tuple (Company,avg(salary))
per group, ordered by Salary

Manages(m_name,e_name)

Employee(name, salary)

Aggregation Examples

- Give per manager the number of people they manage:

```
SELECT m_name, count(e_name)
```

```
FROM Manages
```

```
GROUP BY m_name ;
```

Manages(m_name, e_name)

Employee(name, salary)

Aggregation Examples

- Give all managers that manage more than 3 employees with a salary exceeding 3000 EUR.

SELECT m_name

FROM Manages **JOIN** Employee **ON** e_name=name

WHERE salary>3000

GROUP BY m_name

HAVING count(e_name) > 3;

Manages(m_name, e_name)

Employee(name, salary)

Aggregation Examples

- Give per manager the number of people they manage with a salary exceeding 3000 EUR.

```
SELECT m_name, count(e_name)
FROM Manages JOIN (select * from Employee
                     WHERE salary>3000)
                     ON e_name=name
GROUP BY m_name ;
```

NOT CORRECT for managers with zero such employees!

More joins

(R NATURAL FULL OUTER JOIN S)

(R NATURAL LEFT OUTER JOIN S)

(R NATURAL RIGHT OUTER JOIN S)

(R FULL OUTER JOIN S ON R.A > S.C)

(R LEFT OUTER JOIN S ON R.A > S.C)

(R RIGHT OUTER JOIN S ON R.A > S.C)

A	B	C
a	b	c
d	e	null
null	f	g

A	B	C
a	b	c
d	e	null

A	B	C
a	b	c
null	f	g

A	B ₁	B ₂	C
d	e	b	c
a	b	null	null
null	null	f	g

A	B ₁	B ₂	C
d	e	b	c
a	b	null	null

A	B ₁	B ₂	C
d	e	b	c
null	null	f	g

OUTER JOIN examples

R	A	B	S	B	C
	a	b		b	c
	d	e		f	g

Manages(m_name, e_name)

Employee(name, salary)

Aggregation Examples

- Give per manager the number of people they manage with a salary exceeding 3000 EUR.

```
SELECT m_name, count(e_name)
FROM Manages LEFT OUTER JOIN (select * from Employee
                                WHERE salary>3000)
                                ON e_name=name
GROUP BY m_name ;
```

What About Conditions With NULL ?

- Recall NULL being a special value indicating something is missing or unknown
 - Arithmetic operations (x, +, -, ...) with NULL return NULL.
 - Comparisons (>, <, =, ...) with NULL return UNKNOWN.
 - NULL=NULL → UNKNOWN ; NULL <> NULL → UNKNOWN
 - Use IS to test on NULL



```
SELECT distinct e_name  
FROM manages  
WHERE m_name=NULL;
```

Empty. Always. Because X=NULL is always UNKNOWN even if X=NULL

```
SELECT distinct e_name  
FROM manages  
WHERE m_name IS NULL;
```

All employees for which manager is missing/unknown

Three-Valued Logic

- Being able to evaluate complex Boolean expressions with NULL values requires an extension of the usual binary logic to a three-valued logic

X	Y	X AND Y	X OR Y	NOT X
TRUE	TRUE	TRUE	TRUE	FALSE
TRUE	UNK.	UNK.	TRUE	FALSE
TRUE	FALSE	FALSE	TRUE	FALSE
UNK.	TRUE	UNK.	TRUE	UNK.
UNK.	UNK.	UNK.	UNK.	UNK.
UNK.	FALSE	FALSE	UNK.	UNK.
FALSE	TRUE	FALSE	TRUE	TRUE
FALSE	UNK.	FALSE	UNK.	TRUE
FALSE	FALSE	FALSE	FALSE	TRUE

- Only tuples for which the condition in the **WHERE** clause evaluates to TRUE are returned.

Exercises

- Band_member(bname,mname)
- Concert(bname,date,location)
- Give the band(s) with fewest members.
- Give all bands that played more often in “Wembley stadium” than “Queen” did.

Exercises

Band_member(bname,mname)
Concert(bname,date,location)

Give the band(s) with fewest members.

```
CREATE VIEW nMembers AS (SELECT bname, count(mname) AS number  
                        FROM Band_member GROUP BY bname);
```

```
CREATE VIEW minimum AS (SELECT min(number) FROM nMembers);
```

```
SELECT B.bname from nMembers B  
WHERE EXISTS (SELECT * FROM minimum M  
              WHERE M.number=B.number);
```

Exercises

Band_member(bname,mname)
Concert(bname,date,location)

Give all bands that played more often in “Wembley stadium” than “Queen” did.

```
CREATE VIEW numWS AS (SELECT bname, count(date) AS num  
                        FROM Concert  
                        WHERE location=“WS” GROUP BY bname);
```

```
SELECT B.bname  
FROM numWS B, numWS Q  
WHERE Q.bname=“Queen” AND Q.num< B.num);
```

Summary

- The standard of all relational database query languages:
The Structured Query Language (SQL)
 - Declarative language: say what you want, not how to compute it
 - Allows for aggregation (sum, min, max, avg, ...)
 - Fully “composable”: output can be input of a SQL query
 - Has “bag” semantics
- Optimizer translates SQL to (an extended version of) the more procedural RA, and creates a physical query plan for RA query
- Expressive power: SQL without aggregation and without duplicates has same expressive power as the relational algebra