

# Cross interpolation

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## Definition

For a given matrix  $\mathbf{A} \in \mathbb{R}^{m \times n}$ , the cross interpolation is defined as:

$$\mathbf{A} \approx \mathbf{\tilde{A}} = \mathbf{A}(\mathbb{I}, J) \mathbf{A}(I, J)^{-1} \mathbf{A}(I, \mathbb{J}),$$

where  $\mathbb{I} = \{1, 2, \dots, m\}$  and  $\mathbb{J} = \{1, 2, \dots, n\}$ , while  $I \subseteq \mathbb{I}$  and  $J \subseteq \mathbb{J}$  with  $|I| = |J| = r$ .

## Main properties

- **Interpolation**

$$\begin{aligned} \mathbf{A}(\mathbb{I}, J) &= \mathbf{\tilde{A}}(\mathbb{I}, J), \quad \mathbf{A}(I, \mathbb{J}) \\ &= \mathbf{\tilde{A}}(I, \mathbb{J}). \end{aligned}$$

- **Low-rank decomposition**

$$\mathbf{A} = \mathbf{\tilde{A}} \quad \text{if} \quad \text{rank}(\mathbf{A}) \leq r.$$

## Algorithm

First, define an error function (matrix) as:

$$\begin{aligned} \mathcal{E}(i, j) &= |\mathbf{A} - \mathbf{\tilde{A}}|(i, j) = |\mathbf{A}(i, j) - \mathbf{\tilde{A}}(i, j)| \\ &= |\mathbf{A}(i, j) - \mathbf{A}(i, J) \mathbf{A}(I, J)^{-1} \mathbf{A}(I, j)| \end{aligned}$$

where  $i \in \mathbb{I} \setminus I$  and  $j \in \mathbb{J} \setminus J$ .

## Full Search

Here, it first implements a **full search** algorithm to find the optimal sets  $I$  and  $J$ .

Then the algorithm goes as follows:

1. Find a point  $(i, j)$  that maximizes  $\mathcal{E}(i, j)$ :

$$(i^\wedge, j^\wedge) = \arg\max_{i \in \mathbb{I}, j \in \mathbb{J}} \mathcal{E}(i, j).$$

2. Add the point  $(i^\wedge, j^\wedge)$  to the sets  $I$  and  $J$ :

$$I \leftarrow I \cup \{i^\wedge\}, \quad J \leftarrow J \cup \{j^\wedge\}.$$

3. Repeat steps 1 and 2 until the interpolation error satisfies a predefined threshold:

$$\|\mathbf{\tilde{A}} - \mathbf{A}\|_{\|\cdot\|} < \epsilon,$$

where  $\|\cdot\|$  is a norm (e.g., Frobenius norm) and  $\epsilon$  is a predefined threshold.

## Rook Search

Then, implements the **rook search** algorithm goes as:

1. Randomly initialize a new pivot  $(i^*, j^*)$ .
2. Column-wise movement:

$$i^* \leftarrow \arg\max_{i \in \mathbb{I}} \mathcal{E}(i, j^*).$$

3. Row-wise movement:

$$j^* \leftarrow \arg\max_{j \in \mathbb{J}} \mathcal{E}(i^*, j).$$

4. Repeat steps 2 and 3 until the limit of iterations is reached or the following condition (rook condition) is met:

$$\begin{aligned} i^* &= \arg\max_{i \in \mathbb{I}} \mathcal{E}(i, j^*), \quad j^* = \arg\max_{j \in \mathbb{J}} \mathcal{E}(i^*, j). \end{aligned}$$

5. Add the point  $(i^*, j^*)$  to the sets  $I$  and  $J$ :

$$I \leftarrow I \cup \{i^*\}, \quad J \leftarrow J \cup \{j^*\}.$$

6. Repeat steps 1 to 5 until the interpolation error satisfies a predefined threshold:

$$\|\mathbf{\mathcal{E}}\|_{\infty} < \epsilon.$$

## Implementation tips

### Schur Complement

Given a invertible matrix  $\mathbf{U} \in \mathbb{R}^{k \times k}$  and the inverse  $\mathbf{U}^{-1}$ , and consider the following block matrix

$$\mathbf{M} = \begin{bmatrix} \mathbf{U} & \mathbf{c} \\ \mathbf{r}^T & p \end{bmatrix} \in \mathbb{R}^{(k+1) \times (k+1)}$$

where

$$\begin{aligned} \mathbf{c} &\in \mathbb{R}^{k \times 1}, \quad \mathbf{r}^T \in \mathbb{R}^{1 \times k}, \quad p \\ &\in \mathbb{R}. \end{aligned}$$

Then the inverse of  $\mathbf{M}$  can be computed as

$$\begin{aligned} \mathbf{M}^{-1} &= \begin{bmatrix} \mathbf{U}^{-1} + \mathbf{U}^{-1} \mathbf{c} \mathbf{s}^{-1} \mathbf{r}^T \mathbf{U}^{-1} & -\mathbf{U}^{-1} \mathbf{c} \mathbf{s}^{-1} \\ \mathbf{r}^T \mathbf{U}^{-1} & \mathbf{s}^{-1} \end{bmatrix} \in \mathbb{R}^{(k+1) \times (k+1)} \\ &= \begin{bmatrix} \mathbf{U}^{-1} + \mathbf{U}^{-1} \mathbf{h} \mathbf{l}^T \mathbf{s}^{-1} & -\mathbf{U}^{-1} \mathbf{h} \mathbf{s}^{-1} \\ \mathbf{r}^T \mathbf{U}^{-1} & \mathbf{s}^{-1} \end{bmatrix} \end{aligned}$$

where

$$\begin{aligned} \mathbf{l} &= \mathbf{U}^{-1} \mathbf{c} \in \mathbb{R}^{k \times 1}, \quad \mathbf{h} = \mathbf{r}^T \mathbf{U}^{-1} \in \mathbb{R}^{1 \times k}, \quad \mathbf{s} = p - \mathbf{r}^T \mathbf{U}^{-1} \mathbf{c} \in \end{aligned}$$

- In the above equations,  $\mathbf{S}$  is actually the Schur complement of  $\mathbf{U}$  in  $\mathbf{M}$ , and the matrix  $\mathbf{M}$  is singular when  $s = 0$ .
- The above equations only hold when  $\mathbf{U}$  and  $\mathbf{M}$  are both invertible.

At step  $(t-1)$ , one have

Now, one somehow acquire a new pivot at  $(i^{\wedge}, j^{\wedge})$ , so that one have

$$\end{aligned} \quad \square$$

[illegible]

$$- \mathbf{\tilde{A}}_{t-1}(:, j^*) \mathbf{A}(i^*, :) + \mathbf{A}(:, j^*) \mathbf{A}(i^*, :) \setminus \mathbf{\tilde{A}}_{t-1} + \mathbf{s}^{t-1} (\mathbf{A}(:, j^*) - \mathbf{\tilde{A}}_{t-1}(:, j^*)) (\mathbf{A}(i^*, :) - \mathbf{\tilde{A}}_{t-1}(i^*, :)). \end{aligned} \quad \$\$$$

Note that

$$\begin{aligned} \mathbf{s} &= \mathbf{p} - \mathbf{r} \mathbf{U}_{t-1}^T \mathbf{c} = \mathbf{A}(i^*, j^*) - \mathbf{\tilde{A}}_{t-1}(i^*, j^*) = \mathcal{E}_{t-1}(i^*, j^*), \setminus \mathbf{A}(:, j^*) - \mathbf{\tilde{A}}_{t-1}(:, j^*) = \mathcal{E}_{t-1}(:, j^*), \setminus \mathbf{A}(i^*, :) - \mathbf{\tilde{A}}_{t-1}(i^*, :) = \mathcal{E}_{t-1}(i^*, :), \end{aligned} \quad \$\$$$

therefore

$$\mathbf{\tilde{A}}_t = \mathbf{\tilde{A}}_{t-1} + \frac{1}{\mathcal{E}_{t-1}(i^*, j^*)} \mathcal{E}_{t-1}(:, j^*) \mathcal{E}_{t-1}(i^*, :). \quad \$\$$$

With the above equation, one can establish the following algorithm:

1. Initialize  $\mathbf{\tilde{A}}_0 = \mathbf{A}$ ,  $\mathcal{E}_0 = \mathbf{A}$ .
2. Pick a new pivot  $(i^*, j^*)$ .
3. Update the matrices:
  - $\mathbf{D}_{t-1} = \frac{1}{\mathcal{E}_{t-1}(i^*, j^*)} \mathcal{E}_{t-1}(:, j^*) \mathcal{E}_{t-1}(i^*, :)$ ;
  - $\mathbf{\tilde{A}}_t = \mathbf{\tilde{A}}_{t-1} + \mathbf{D}_{t-1}$ ;
  - $\mathcal{E}_t = \mathcal{E}_{t-1} - \mathbf{D}_{t-1}$ .

## Efficient ACA

In some scenarios, one should assume that the original matrix  $\mathbf{A}$  is too large to fit in memory and expensive to evaluate. In such cases, one should follow a more memory-efficient implementation.

Let

$$\begin{aligned} \mathbf{p}_t &= \mathcal{E}_t(l_k, j_k), \setminus \mathbf{c}_t = \mathcal{E}_t(:, j_k), \setminus \mathbf{r}_t = \mathcal{E}_t(l_k, :), \end{aligned} \quad \$\$$$

then one has

$$\begin{aligned} \mathcal{E}_t &= \mathbf{A} - \mathbf{\tilde{A}}_t = \mathbf{A} - \sum_{k=0}^{t-1} \frac{1}{p_k} \mathbf{c}_k \mathbf{r}_k^T, \setminus \mathcal{E}_t(:, j) = \mathbf{A}(:, j) - \mathbf{\tilde{A}}_t(:, j) = \mathbf{A}(:, j) - \sum_{k=0}^{t-1} \frac{1}{p_k} \mathbf{c}_k \mathbf{r}_k^T(j), \setminus \mathcal{E}_t(i, :) = \mathbf{A}(i, :) - \mathbf{\tilde{A}}_t(i, :) = \mathbf{A}(i, :) - \sum_{k=0}^{t-1} \frac{1}{p_k} \mathbf{c}_k \mathbf{r}_k^T(i). \end{aligned} \quad \$\$$$

Therefore, one can follow

1. Find a new pivot  $(i^*, j^*)$ .
2. Find new vectors:
  - $\mathbf{c}_t = \mathcal{E}_t(:, j^*)$ ;
  - $\mathbf{r}_t = \mathcal{E}_t(i^*, :)$ .