

# BRAID VARIETIES (IV)

2026/01/20

## 1. FRAMED VERSION

**1.1. Motivation.** Recall that any two waves  $\mathfrak{W}_1, \mathfrak{W}_2 : \beta_1 \rightarrow \beta_2$  consisting only of tetravalent and hexavalent vertices (i.e. braid relations) are related by **two-braid relations**. What we proved last time can be summarized as

the  $z$ -labeling is invariant under two-braid relations.

In this subsection, we will introduce another labeling invariant under two-braid relations. Recall that  $y_i : \mathbb{C} \rightarrow G$  is defined by

$$y_i(t) = \text{image of } \begin{bmatrix} 1 & \\ t & 1 \end{bmatrix} \in G.$$

We can replace  $B_i$  by  $y_i$  with the expense of parametrizing only most of points.

**Lemma.** For a flag  $g_0B$ ,

$$\{xB : g_0B \xrightarrow{s_i} xB\} = g_0Bs_iB/B \approx \{g_0y_i(z)B/B : z \in \mathbb{C}\} \cong \mathbb{C}$$

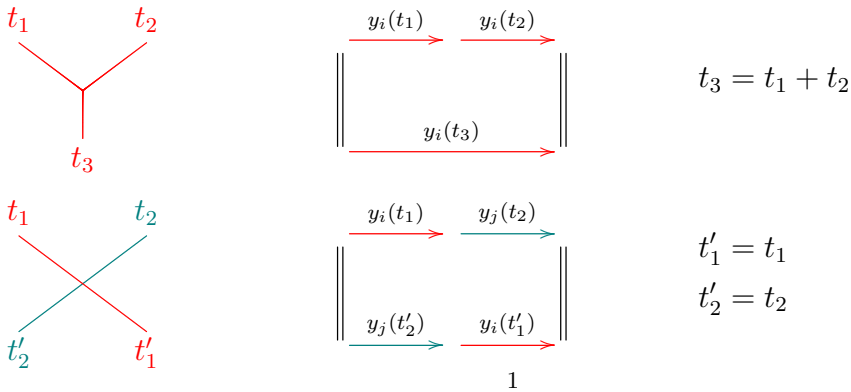
where  $\approx$  should be read as “birational to”. That is, once a representative of  $g_0B$  is chosen, there is a canonical choice of representative for **almost all**  $xB$  by  $x = g_0y_i(z)$ .

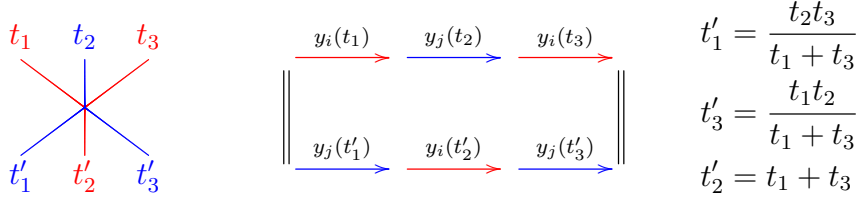
As a result, we can parametrize the **open Bott–Samelson variety**

$$BS(\beta) = \{(g_kB) : B \xrightarrow{s_{i_1}} g_1B \xrightarrow{s_{i_2}} \cdots g_{\ell-1}B \xrightarrow{s_{i_\ell}} g_\ell B\}.$$

(We met this variety in the proof in previous talks, but we did not name it). We can similarly define  $BS(\mathfrak{W})$ , i.e. the same as  $X(\mathfrak{W})$  but no restriction on the rightmost region.

There is a similar rule as the  $z$ -labeling.





Note that the situation is much easier, since no dashed edges need to be introduced.

**Corollary.** The  $t$ -labeling above satisfies two-braid relations.

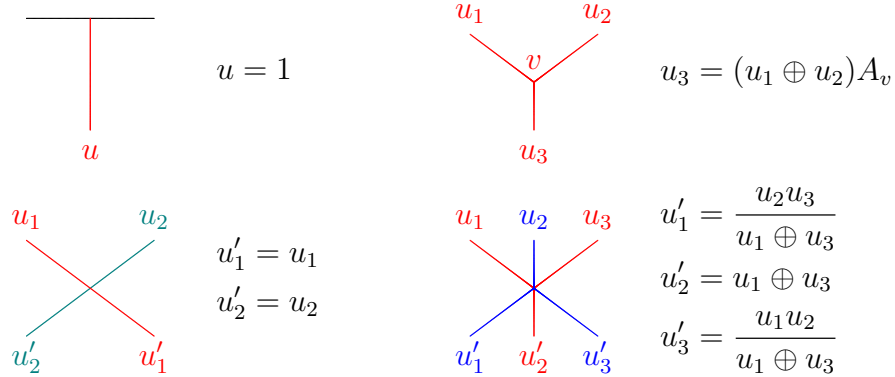
Note that  $\frac{t_2 t_3}{t_1 + t_3}$  is not a polynomial function, some values (e.g.  $t_1 = 1, t_3 = -1$ ) will lead to a problem. However, Lusztig noticed that all above functions restrict to well-defined function between  $\mathbb{R}_{>0} \rightarrow \mathbb{R}_{>0}$ , since all functions are subtraction-free. In particular, it is a well-defined function between any **semifield**. Recall

field = a set with  $+, -, \times, \div$ ,  
semifield = a set with  $+, \times, \div$ .

**1.2. The  $u$ -labeling.** Let us introduce a variable  $A_v$  for each trivalent vertex  $v \in \mathfrak{W}_3$ . There is a semifield structure over the set of monomials in  $\{A_v\}$

$$A^{\mathbf{a}} \oplus A^{\mathbf{b}} = A^{\min(\mathbf{a}, \mathbf{b})} \quad \text{min = entrywise minimum.}$$

We define the  $u$ -labeling as follows



**Corollary.** The  $u$ -labeling above satisfies two-braid relations.

We are going to define another parametrization of  $X(\mathfrak{W})$  using  $u$ -labeling. Let us define  $\chi_i : \mathbb{C}^\times \rightarrow G$  by

$$\chi_i(u) = \text{image of } \begin{bmatrix} u & \\ & u^{-1} \end{bmatrix} \in G.$$

We define

$$B_i(\tilde{z}, u) = B_i(\tilde{z})\chi_i(u) = \text{image of } \begin{bmatrix} uz & -u^{-1} \\ u & \end{bmatrix} \in G.$$

Let us define the following **framed parametrization** of wave variety:

$$\tilde{X}(\mathfrak{W}^=) = \left\{ \begin{array}{c} (A_v)_{v \in \mathfrak{W}_3} \\ \times \\ (\tilde{\mathbf{g}}_r)_{r \in \text{Region}(\mathfrak{W}^=)} \end{array} \left| \begin{array}{l} A_v \in \mathbb{C}^\times, \tilde{\mathbf{g}}_r \in G \\ a \mid b \Rightarrow \tilde{\mathbf{g}}_b = \tilde{\mathbf{g}}_a B_i(\tilde{\mathbf{z}}_e, u_e) \text{ for some } \tilde{\mathbf{z}}_e \\ \cdots \Rightarrow \tilde{\mathbf{g}}_a = \tilde{\mathbf{g}}_b \tilde{\mathbf{U}}_e \text{ for some } \tilde{\mathbf{U}}_e \in \text{Rad}(B) \\ \tilde{\mathbf{g}}_{\text{leftmost}} = 1, \tilde{\mathbf{g}}_{\text{rightmost}} \in w_0 B \end{array} \right. \right\}.$$

Since each  $u_e$  is a monomial in  $\{A_v\}$ , once  $A_v$ 's are valued,  $u_e$ 's are determined. Notice that difference — this time we require  $\tilde{\mathbf{U}}_e \in \text{Rad}(B)$  the unipotent subgroup. Similar as we did last time:

- there is an obvious map

$$\tilde{X}(\mathfrak{W}^=) \rightarrow X(\mathfrak{W}), \quad (A_v) \times (\tilde{\mathbf{g}}_r) \mapsto (\tilde{\mathbf{g}}_r B);$$

- $\tilde{X}(\mathfrak{W}^=)$  also admits another description via

$$(A_v)_{v \in \mathfrak{W}_3} \times (\tilde{\mathbf{z}}_e)_{e \in \text{Edge}(\mathfrak{W}^=)} \times (\tilde{\mathbf{U}}_e)_{e \in \text{Dashed}(\mathfrak{W}^=)}.$$

**Theorem A.** For a wave from  $\beta \rightarrow \sigma_{w_0}$ , the projection

$$\tilde{X}(\mathfrak{W}^=) \rightarrow (\mathbb{C}^\times)^{\mathfrak{W}_3}, \quad (A_v) \times (\tilde{\mathbf{g}}_r) \mapsto (A_v)$$

is an isomorphism.

*Proof.* We want to show  $\{\tilde{\mathbf{g}}_r\}$  is a polynomial in  $\{A_v\}$ . Equivalently, we need to solve  $(\tilde{\mathbf{z}}_e)$  and  $(\tilde{\mathbf{U}}_e)$  from the  $u$ -labeling. The variables  $\tilde{\mathbf{z}}_e$  and  $\tilde{\mathbf{U}}_e$  can be solved from bottom to top. We need to show

- (1) On the bottom of  $\mathfrak{W}$ , the  $\tilde{\mathbf{z}}$ -labeling vanishes;
- (2) At each vertex, the bottom/left labeling determine the top/right labeling.

The claim (1) follows from a similar argument as last time. Assume  $\sigma_{w_0} = \sigma_{i_1} \cdots \sigma_{i_l}$ , then

$$\mathbb{C}^{\ell(w_0)} \xrightarrow{\sim} Bw_0B/B, \quad (\tilde{\mathbf{z}}_k) \mapsto B_{i_1}(\tilde{\mathbf{z}}_1, u_1) \cdots B_{i_l}(\tilde{\mathbf{z}}_l, u_l)B$$

for fixed  $u_1, \dots, u_l \in \mathbb{C}^\times$ . In particular, only vanishing  $(z_k)$  could be mapped to  $w_0B$ . The claim (2) follows from explicit computation, see Appendix.  $\square$

**Theorem B.** For a wave from  $\beta \rightarrow \sigma_{w_0}$ , the projection

$$\tilde{X}(\mathfrak{W}^=) \rightarrow X(\mathfrak{W}), \quad (A_v) \times (\tilde{\mathbf{g}}_r) \mapsto (\tilde{\mathbf{g}}_r B)$$

is an isomorphism.

*Proof.* We can construct an inverse map. By above Theorem A, it suffices to solve  $\{A_v\}$  from  $(g_r) \in X(\mathfrak{W}^=)$ . For any coroot  $\gamma$ , denote

$$\mathbb{C}^\times \rightarrow T \subset G, \quad u \mapsto u^\gamma.$$

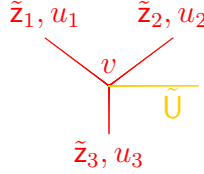
Note that  $u^{\alpha_i^\vee} = \chi_i(u)$  and

$$u^\gamma B_j(\tilde{\mathbf{z}}) = B_j(u^{\langle \gamma, \alpha_j \rangle} \tilde{\mathbf{z}}) u^{s_j \gamma}.$$

By icking a region walk (without crossing the dashed edges), it is easy to see by induction that

$$\tilde{g}_r = g_r \cdot (\text{some element in } T), \quad \tilde{z}_r = z_r \cdot (\text{some nonzero constant})$$

where the brackets can be explicitly described in a Laurent monomial in  $\{u_e\}$ , thus in  $\{A_v\}$  (see below). Let us look at



Assume all  $A_{v'}$  are solved for  $v'$  above, and we are going to solve  $A_v$ . Then in particular,

$$u_1, u_2, \tilde{z}_2 \text{ are known} \quad u = (u_1 \oplus u_2)A_v \text{ is unknown.}$$

Since  $\tilde{z}_2 = \frac{u}{u_1 u_2}$  (see Appendix), we can solve  $A_v$  as a regular function in  $A_{v'}$  and  $\tilde{z}_e$  for vertices  $v'$  and edges  $e$  above the vertex  $v$ .  $\square$

Let us be precise for the claim in the proof. Assume the region walk is given by

$$\text{leftmost} = r_0 \begin{array}{c} i_1 \\ | \\ \tilde{z}_1, u_1 \end{array} r_1 \begin{array}{c} i_2 \\ | \\ \tilde{z}_2, u_2 \end{array} r_2 \begin{array}{c} i_3 \\ | \\ \tilde{z}_3, u_3 \end{array} \cdots \begin{array}{c} i_{l-1} \\ | \\ \tilde{z}_{l-1}, u_{l-1} \end{array} r_{l-1} \begin{array}{c} i_l \\ | \\ \tilde{z}_l, u_l \end{array} r$$

we have

$$\begin{aligned} \tilde{g}_r &= B_{i_1}(\tilde{z}_1, u_1) B_{i_2}(\tilde{z}_2, u_2) \cdots B_{i_l}(\tilde{z}_l, u_l) \\ &= B_{i_1}(\tilde{z}_1) \chi_{i_1}(u_1) B_{i_2}(\tilde{z}_2) \chi_{i_2}(u_2) \cdots B_{i_l}(\tilde{z}_l) \chi_{i_l}(u_l) \\ &= B_{i_1}(z'_1) B_{i_2}(z'_2) \cdots B_{i_l}(z'_l) \cdot u_1^{s_{i_l} \cdots s_{i_2} \alpha_1^\vee} u_2^{s_{i_l} \cdots s_{i_3} \alpha_2^\vee} \cdots u_l^{\alpha_l^\vee} \end{aligned}$$

with

$$z'_k = \tilde{z}_k u_1^{\langle s_{i_{k-1}} \cdots s_{i_2} \alpha_1^\vee, \alpha_{i_k} \rangle} u_2^{\langle s_{i_{k-1}} \cdots s_{i_3} \alpha_2^\vee, \alpha_{i_k} \rangle} \cdots u_k^{\langle \alpha_{k-1}^\vee, \alpha_{i_k} \rangle}.$$

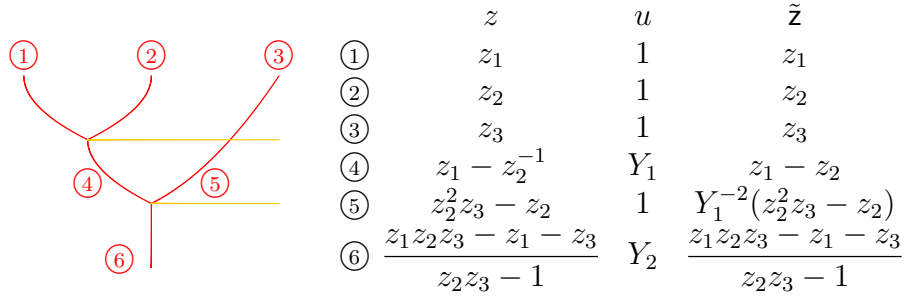
Note that we have to have  $z_k = z'_k$  since  $(z'_k)$  and  $(z_k)$  parametrizes the same element in  $X(\mathfrak{W})$ .

**1.3. Cluster chart.** We will call the isomorphism

$$(\mathbb{C}^\times)^{\mathfrak{M}_3} \xrightarrow{\sim} \tilde{X}(\mathfrak{W}^-) \xleftarrow{\sim} X(\mathfrak{W})$$

the **cluster chart**.

**Example.** Let us consider the example



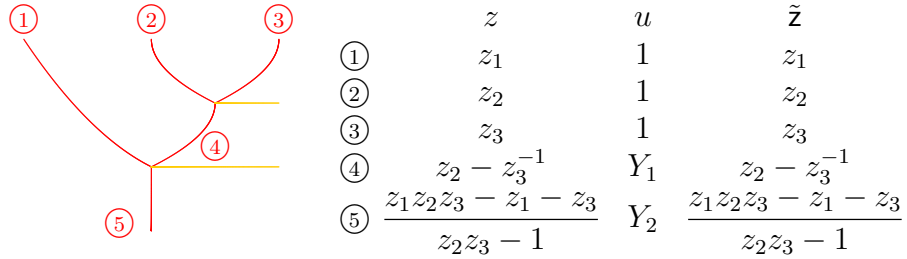
At vertices 1, 2, we have

$$z_2 = \tilde{z}_2 = \frac{u_4}{u_1 u_2} = Y_1, \quad Y_1^{-2}(z_2^2 z_3 - z_2) = \tilde{z}_5 = \frac{u_6}{u_4 u_5} = Y_2 / Y_1.$$

So

$$Y_1 = z_2, \quad Y_2 = z_2 z_3 - 1.$$

Let us consider another example



At vertices 1, 2, we have

$$z_3 = \tilde{z}_3 = \frac{u_4}{u_2 u_3} = Y_1, \quad z_2 - z_3^{-1} = \tilde{z}_4 = \frac{u_5}{u_1 u_4} = Y_2 / Y_1.$$

So

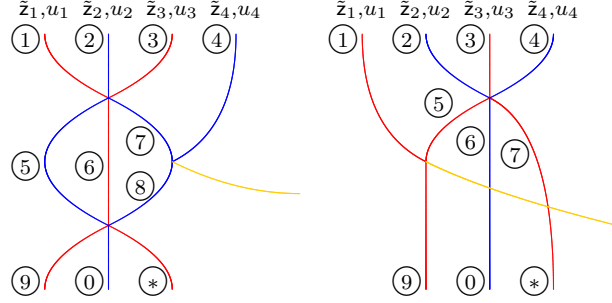
$$Y_1 = z_3, \quad Y_2 = z_2 z_3 - 1.$$

**Theorem.** For two waves  $\mathfrak{W}, \mathfrak{M}$  related by two-braid relations and 1212 relations, we have a natural bijection  $\mathfrak{W}_3 \cong \mathfrak{M}_3$ , the cluster charts

$$(\mathbb{C}^\times)^{\mathfrak{W}_3} \longrightarrow X(\beta)$$

are the same.

*Proof.* The proof is similar to the the proof we explained last time. The only computation we need to do is 1212 relations. Consider



We know  $u_{\textcircled{9}}, u_{\textcircled{0}}, u_{\textcircled{*}}$  agree. By an argument similar to  $z$ -labeling, we know  $\tilde{z}_{\textcircled{9}}, \tilde{z}_{\textcircled{0}}, \tilde{z}_{\textcircled{*}}$  agree. From the way we construct the inverse map in Theorem B, it suffices to show  $Y$  agrees.

For the left diagram, we have

$$\tilde{z}_4 = \tilde{z}_{\textcircled{4}} = Y \frac{u_{\textcircled{4}} \oplus u_{\textcircled{7}}}{u_{\textcircled{4}} u_{\textcircled{7}}} = Y \frac{\frac{u_1 u_2}{u_1 \oplus u_3} \oplus u_4}{\frac{u_1 u_2}{u_1 \oplus u_3} u_4} = Y \frac{u_1 u_2 \oplus u_1 u_4 \oplus u_3 u_4}{u_1 u_2 u_4}.$$

For the right diagram, we have

$$\tilde{z}_{\textcircled{5}} = Y \frac{u_{\textcircled{1}} \oplus u_{\textcircled{5}}}{u_{\textcircled{1}} u_{\textcircled{5}}} = Y \frac{u_1 \oplus \frac{u_3 u_4}{u_2 \oplus u_4}}{u_1 \frac{u_3 u_4}{u_2 \oplus u_4}} = Y \frac{u_1 u_2 \oplus u_1 u_4 \oplus u_3 u_4}{u_1 u_3 u_4}.$$

By **(Comm)**, we have

$$\tilde{z}_4 = \tilde{z}_{\textcircled{4}} = \frac{u_{\textcircled{3}}}{u_{\textcircled{2}}} \tilde{z}_{\textcircled{5}} = \frac{u_3}{u_2} \tilde{z}_{\textcircled{5}}.$$

We get the same equation for  $Y$ . □

## APPENDIX A. EXPLICIT SOLUTIONS

We can solve the equations explicitly.

$$\begin{aligned}
 & \begin{array}{ccc} \tilde{z}_1, u_1 & \tilde{z}_2, u_2 & \\ & \tilde{U} & \\ & \tilde{z}_3, u_3 & \end{array} & \begin{array}{ccc} B_i(\tilde{z}_1, u_1) & B_i(\tilde{z}_2, u_2) & \\ \parallel & & \uparrow \tilde{U} \\ B_i(\tilde{z}_3, u_3) & & \end{array} & \begin{aligned} \tilde{z}_2 &= \frac{u}{u_1 u_2}, & \tilde{z}_1 &= \tilde{z}_3 + \frac{u_2}{u_1 u}, \\ \tilde{U} &= \text{image of } \begin{bmatrix} 1 & -\frac{u_1}{u_2 u} \\ & 1 \end{bmatrix} \end{aligned} \\
 & \begin{array}{ccc} \tilde{z}_1, u_1 & \tilde{z}_2, u_2 & \\ & \tilde{z}_2', u_2' & \\ \tilde{z}_2', u_2' & \tilde{z}_1', u_1' & \end{array} & \begin{array}{ccc} B_i(\tilde{z}_1, u_1) & B_j(\tilde{z}_2, u_2) & \\ \parallel & & \parallel \\ B_j(\tilde{z}_2', u_2') & B_i(\tilde{z}_1', u_1') & \end{array} & \begin{aligned} \tilde{z}_1 &= \tilde{z}_1' \\ \tilde{z}_2 &= \tilde{z}_2' \end{aligned} \quad \boxed{\begin{array}{l} u_1' = u_1 \\ u_2' = u_2 \end{array}} \\
 & \begin{array}{ccc} \tilde{z}_1, u_1 & \tilde{z}_2, u_2 & \tilde{z}_3, u_3 \\ & \tilde{z}_1', u_1' & \tilde{z}_2', u_2' \\ \tilde{z}_1', u_1' & \tilde{z}_2', u_2' & \tilde{z}_3', u_3' \end{array} & \begin{array}{ccc} B_i(\tilde{z}_1, u_1) & B_j(\tilde{z}_2, u_2) & B_i(\tilde{z}_3, u_3) \\ \parallel & & \parallel \\ B_j(\tilde{z}_1', u_1') & B_i(\tilde{z}_2', u_2') & B_j(\tilde{z}_3', u_3') \end{array} & \begin{aligned} \tilde{z}_1 &= \frac{u_1'}{u_2'} \tilde{z}_3', \\ \tilde{z}_3 &= \frac{u_2}{u_1} \tilde{z}_1', \\ \tilde{z}_2 &= \frac{u_1 u_3'}{u_2'} \tilde{z}_1' \tilde{z}_3' - \frac{u_1}{u_1'} \tilde{z}_2' \end{aligned} \quad \boxed{\begin{array}{l} u_2' u_3' = u_1 u_2 \\ u_1' u_2' = u_2 u_3 \end{array}} \\
 & \begin{array}{ccc} \tilde{z}, u & & \\ \tilde{U} & & \tilde{U}' \\ & \tilde{z}', u' & \end{array} & \begin{array}{ccc} & B_i(\tilde{z}, u) & \\ \uparrow \tilde{U} & & \uparrow \tilde{U}' \\ & B_i(\tilde{z}', u') & \end{array} & \begin{aligned} \tilde{z} &= \tilde{z}' - \xi_i(U), \\ \tilde{U}' &= B_i(\tilde{z}', u')^{-1} \tilde{U} B_i(\tilde{z}, u) \in \text{Rad}(B) \end{aligned} \quad \boxed{u' = u}
 \end{aligned}$$

Recall  $\xi_i(U)$  = coefficient of  $E_i$  in  $U$ . Notice that  $\xi_i(U') = 0$ . This fact will be used later.

All the relations follow from the computation in  $SL_2 \subset GL_2$ ,  $SL_2 \times SL_2 \subset GL_4$ ,  $SL_3 \subset GL_3$  and  $SL_2 \subset GL_2$ . The following is the code.

```

R.<u1,u2,u3p,z3p,z1,z2,z3,a> = QQ[];
B = lambda z,u: matrix([[u*z,-u^(-1)], [u,0]])
U = matrix([[1,a], [0,1]]);
Rel = (B(z1,u1)*B(z2,u2) - B(z3,u3p)*U).change_ring(SR)
solve([Rel[i][j]==0 for i in range(2) for j in range(2)], SR.var("u1p,z1,z2,a"))

```

```

[[u1p == r1,
 z1 == (u1*u3p*z3 + u2)/(u1*u3p),
 z2 == u3p/(u1*u2),
 a == -u1/(u2*u3p)]

```

```

R.<u1,u2,u1p,u2p,z1p,z2p,z1,z2> = QQ[]
B1 = lambda z,u: matrix([[u*z,-u^(-1),0,0], [u,0,0,0], [0,0,1,0], [0,0,0,1]])
B3 = lambda z,u: matrix([[1,0,0,0], [0,1,0,0], [0,0,u*z,-u^(-1)], [0,0,u,0]])
Rel = (B1(z1,u1)*B3(z2,u2)-B3(z2p,u2p)*B1(z1p,u1p)).change_ring(SR)

```

```
solve([Rel[i][j]==0 for i in reversed(range(4)) for j in range(4)], SR.var("u1p,u2p,z1,
z2"))
```

```
[[u1p == u1, u2p == u2, z1 == z1p, z2 == z2p]]
```

```
R.<u1,u2,u3,u1p,u2p,u3p,z1p,z2p,z3p,z1,z2,z3> = QQ[]
B1 = lambda z,u: matrix([[u*z,-u^(-1)], [u,0,0], [0,0,1]])
B2 = lambda z,u: matrix([[1,0,0], [0,u*z,-u^(-1)], [0,u,0]])
Rel = (B1(z1,u1)*B2(z2,u2)*B1(z3,u3)-B2(z1p,u1p)*B1(z2p,u2p)*B2(z3p,u3p)).change_ring(
SR)
solve([Rel[i][j]==0 for i in reversed(range(2)) for j in range(2)], SR.var("u1p,u2p,u3p
,z1,z2,z3"))
```

```
[[u1p == r1*r2/z3p,
u2p == r1,
u3p == r1*r2*u1/(u3*z3p),
z1 == r2,
z2 == (r1^2*r2^2*u1*z1p - r1*u1*z2p*z3p)/(u2*u3*z3p),
z3 == r1^2*r2*z1p/(u1*u3*z3p)]]
```

```
R.<u,up,z,zp,a,ap> = QQ[]
B = lambda z,u: matrix([[u*z,-u^(-1)], [u,0]])
U = matrix([[1,a], [0,1]]); Up = matrix([[1,ap], [0,1]])
Rel = (U*B(z,u)-B(zp,up)*Up).change_ring(SR)
solve([Rel[i][j]==0 for i in range(2) for j in range(2)], SR.var("up,z,a,ap"))
```

```
[[up == u, z == r1, a == -r1 + zp, ap == 0]]
```

To convince the readers, we write down the matrix equations for the trivalent and hexavalent vertices

$$\begin{bmatrix} u_1 u_2 \tilde{z}_1 \tilde{z}_2 - u_1^{-1} u_2 & -u_1 u_2^{-1} \tilde{z}_1 \\ u_1 u_2 \tilde{z}_2 & -u_1 u_2^{-1} \end{bmatrix} = \begin{bmatrix} u_1 \tilde{z}_1 & -u_1^{-1} \\ u_1 & 0 \end{bmatrix} \begin{bmatrix} u_2 \tilde{z}_2 & -u_2^{-1} \\ u_2 & 0 \end{bmatrix}$$

$$\begin{bmatrix} u \tilde{z} & u \tilde{z} a - u^{-1} \\ u & u a \end{bmatrix} = \begin{bmatrix} u \tilde{z} & -u^{-1} \\ u & 0 \end{bmatrix} \begin{bmatrix} 1 & a \\ 0 & 1 \end{bmatrix}.$$

$$\begin{bmatrix} u_1 u_3 \tilde{z}_1 \tilde{z}_3 - u_1^{-1} u_2 u_3 \tilde{z}_2 & -u_1 u_3^{-1} \tilde{z}_1 & u_1^{-1} u_2^{-1} \\ u_1 u_3 \tilde{z}_3 & -u_1 u_3^{-1} & 0 \\ u_2 u_3 & 0 & 0 \end{bmatrix} = \begin{bmatrix} u_1 \tilde{z}_1 & -u_1^{-1} & 0 \\ u_1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & u_2 \tilde{z}_2 & -u_2^{-1} \\ 0 & u_2 & 0 \end{bmatrix} \begin{bmatrix} u_3 \tilde{z}_3 & -u_3^{-1} & 0 \\ u_3 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} u_{2'} \tilde{z}_{2'} & -u_{2'}^{-1} u_{3'} \tilde{z}_{3'} & u_{2'}^{-1} u_{3'}^{-1} \\ u_{1'} u_{2'} \tilde{z}_{1'} & -u_{1'}^{-1} u_{3'} & 0 \\ u_{1'} u_{2'} & 0 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & u_{1'} \tilde{z}_{1'} & -u_{1'}^{-1} \\ 0 & u_{1'} & 0 \end{bmatrix} \begin{bmatrix} u_{2'} \tilde{z}_{2'} & -u_{2'}^{-1} & 0 \\ u_{2'} & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & u_{3'} \tilde{z}_{3'} & -u_{3'}^{-1} \\ 0 & u_{3'} & 0 \end{bmatrix}.$$