

BRAID VARIETIES (IV)

2026/01/20

1. FRAMED VERSION

The material in this section is from [1].

1.1. Motivation. Recall that any two weaves $\mathfrak{W}_1, \mathfrak{W}_2 : \beta_1 \rightarrow \beta_2$ consisting only of tetravalent and hexavalent vertices (i.e. braid relations) are related by **two-braid relations**. What we proved last time can be summarized as

the z -labeling is invariant under two-braid relations.

In this subsection, we will introduce another labeling invariant under two-braid relations. Recall that $y_i : \mathbb{C} \rightarrow G$ is defined by

$$y_i(t) = \text{image of } \begin{bmatrix} 1 \\ t & 1 \end{bmatrix} \in G.$$

We can replace B_i by y_i with the expanse of parametrizing only most of points.

Lemma. For a flag g_0B ,

$$\{xB : g_0B \xrightarrow{s_i} xB\} = g_0Bs_iB/B \approx \{g_0y_i(z)B/B : z \in \mathbb{C}\} \cong \mathbb{C}$$

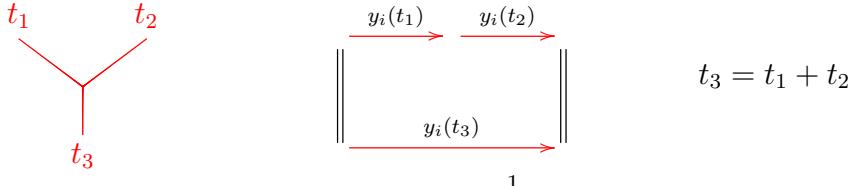
where \approx should be read as “birational to”. That is, once a representative of g_0B is chosen, there is a canonical choice of representative for **almost all** xB by $x = g_0y_i(z)$.

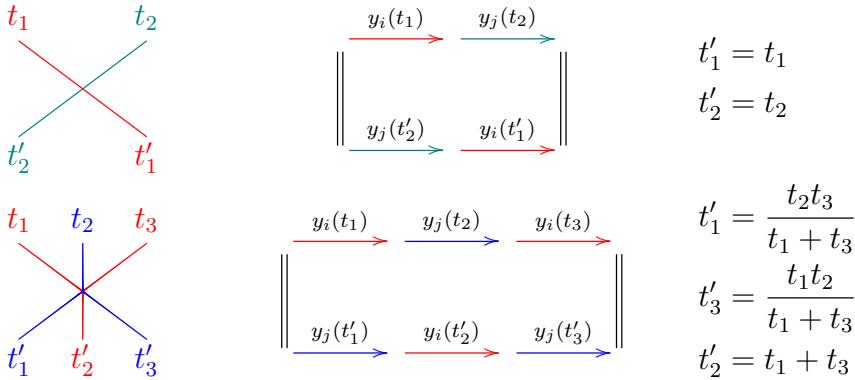
As a result, we can parametrize the **open Bott–Samelson variety**

$$BS(\beta) = \{(g_kB) : B \xrightarrow{s_{i_1}} g_1B \xrightarrow{s_{i_2}} \cdots g_{\ell-1}B \xrightarrow{s_{i_\ell}} g_\ell B\}.$$

(We met this variety in the proof in previous talks, but we did not name it). We can similar define $BS(\mathfrak{W})$, i.e. the same as $X(\mathfrak{W})$ but no restriction on the rightmost region.

There is a similar rule as the z -labeling.





Note that the situation is much easier, since no dashed edges need to be introduced.

Corollary. The t -labeling above satisfies two-braid relations.

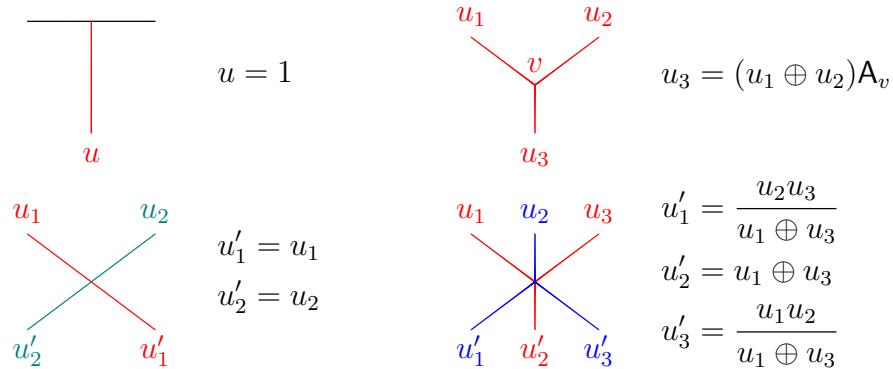
Note that $\frac{t_2 t_3}{t_1 + t_3}$ is not a polynomial function, some values (e.g. $t_1 = 1, t_3 = -1$) will lead to a problem. However, Lusztig [2, Proposition 2.5] first noticed that all above functions restrict to well-defined function between $\mathbb{R}_{>0} \rightarrow \mathbb{R}_{>0}$, since all functions are subtraction-free. In particular, it is a well-defined function between any **semifield**. Recall

$$\begin{aligned} \text{field} &= \text{a set with } +, -, \times, \div, \\ \text{semifield} &= \text{a set with } +, \quad \times, \div. \end{aligned}$$

1.2. The u -labeling. Let us introduce a variable A_v for each trivalent vertex $v \in \mathfrak{W}_3$. There is a semifield structure over the set of monomials in $\{A_v\}$

$$A^a \oplus A^b = A^{\min(a,b)} \quad \min = \text{entrywise minimum.}$$

We define the u -labeling as follows



Corollary. The u -labeling above satisfies two-braid relations.

We are going to define another parametrization of $X(\mathfrak{W})$ using u -labeling. Let us define $\chi_i : \mathbb{C}^\times \rightarrow G$ by

$$\chi_i(u) = \text{image of } \begin{bmatrix} u & \\ & u^{-1} \end{bmatrix} \in G.$$

We define

$$B_i(\tilde{z}, u) = B_i(\tilde{z})\chi_i(u) = \text{image of } \begin{bmatrix} uz & -u^{-1} \\ u & \end{bmatrix} \in G.$$

Let us define the following **framed parametrization** of weave variety:

$$\tilde{X}(\mathfrak{W}^=) = \left\{ \begin{array}{c} (\mathsf{A}_v)_{v \in \mathfrak{W}_3} \\ \times \\ (\tilde{\mathsf{g}}_r)_{r \in \text{Region}(\mathfrak{W}^=)} \end{array} \middle| \begin{array}{l} \mathsf{A}_v \in \mathbb{C}^\times, \tilde{\mathsf{g}}_r \in G \\ a \stackrel{i}{\mid} b \Rightarrow \tilde{\mathsf{g}}_b = \tilde{\mathsf{g}}_a B_i(\tilde{z}_e, u_e) \text{ for some } \tilde{z}_e \\ a \stackrel{a}{\dots} b \Rightarrow \tilde{\mathsf{g}}_a = \tilde{\mathsf{g}}_b \tilde{U}_e \text{ for some } \tilde{U}_e \in \text{Rad}(B) \\ \tilde{\mathsf{g}}_{\text{leftmost}} = 1, \tilde{\mathsf{g}}_{\text{rightmost}} \in w_0 B \end{array} \right\}.$$

Since each u_e is a monomial in $\{\mathsf{A}_v\}$, once A_v 's are valued, u_e 's are determined. Notice that difference — this time we require $\tilde{U}_e \in \text{Rad}(B)$ the unipotent subgroup. Similar as we did last time:

- there is an obvious map

$$\tilde{X}(\mathfrak{W}^=) \rightarrow X(\mathfrak{W}), \quad (\mathsf{A}_v) \times (\tilde{\mathsf{g}}_r) \mapsto (\tilde{\mathsf{g}}_r B);$$

- $\tilde{X}(\mathfrak{W}^=)$ also admits another description via

$$(\mathsf{A}_v)_{v \in \mathfrak{W}_3} \times (\tilde{z}_e)_{e \in \text{Edge}(\mathfrak{W}^=)} \times (\tilde{U}_e)_{e \in \text{Dashed}(\mathfrak{W}^=)}.$$

Theorem A. For a weave from $\beta \rightarrow \sigma_{w_0}$, the projection

$$\tilde{X}(\mathfrak{W}^=) \rightarrow (\mathbb{C}^\times)^{\mathfrak{W}_3}, \quad (\mathsf{A}_v) \times (\tilde{\mathsf{g}}_r) \mapsto (\mathsf{A}_v)$$

is an isomorphism.

Proof. We want to show $\{\tilde{\mathsf{g}}_r\}$ is a polynomial in $\{\mathsf{A}_v\}$. Equivalently, we need to solve (\tilde{z}_e) and (\tilde{U}_e) from the u -labeling. The variables \tilde{z}_e and \tilde{U}_e can be solved from bottom to top. We need to show

- (1) On the bottom of \mathfrak{W} , the \tilde{z} -labeling vanishes;
- (2) At each vertex, the bottom/left labeling determine the top/right labeling.

The claim (1) follows from a similar argument as last time. Assume $\sigma_{w_0} = \sigma_{i_1} \cdots \sigma_{i_l}$, then

$$\mathbb{C}^{\ell(w_0)} \xrightarrow{\sim} Bw_0 B/B, \quad (\tilde{z}_k) \mapsto B_{i_1}(\tilde{z}_1, u_1) \cdots B_{i_l}(\tilde{z}_l, u_l) B$$

for fixed $u_1, \dots, u_l \in \mathbb{C}^\times$. In particular, only vanishing (z_k) could be mapped to $w_0 B$. The claim (2) follows from explicit computation, see Appendix. \square

Theorem B. For a weave from $\beta \rightarrow \sigma_{w_0}$, the projection

$$\tilde{X}(\mathfrak{W}^=) \rightarrow X(\mathfrak{W}), \quad (\mathsf{A}_v) \times (\tilde{\mathsf{g}}_r) \mapsto (\tilde{\mathsf{g}}_r B)$$

is an isomorphism.

Proof. We can construct an inverse map. By above Theorem A, it suffices to solve $\{\mathsf{A}_v\}$ from $(g_r) \in X(\mathfrak{W}^=)$. For any coroot γ , denote

$$\mathbb{C}^\times \rightarrow T \subset G, \quad u \mapsto u^\gamma.$$

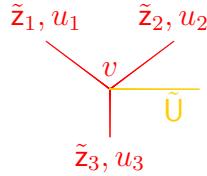
Note that $u^{\alpha_i^\vee} = \chi_i(u)$ and

$$u^\gamma B_j(\tilde{z}) = B_j(u^{\langle \gamma, \alpha_j \rangle} \tilde{z}) u^{s_j \gamma}.$$

By icking a region walk (without crossing the dashed edges), it is easy to see by induction that

$$\tilde{g}_r = g_r \cdot (\text{some element in } T), \quad \tilde{z}_r = z_r \cdot (\text{some nonzero constant})$$

where the brackets can be explicitly described in a Laurent monomial in $\{u_e\}$, thus in $\{\mathsf{A}_v\}$ (see below). Let us look at



Assume all $\mathsf{A}_{v'}$ are solved for v' above, and we are going to solve A_v . Then in particular,

$$u_1, u_2, \tilde{z}_2 \text{ are known,} \quad u = (u_1 \oplus u_2)\mathsf{A}_v \text{ is unknown.}$$

Since $\tilde{z}_2 = \frac{u}{u_1 u_2}$ (see Appendix), we can solve A_v as a regular function in $\mathsf{A}_{v'}$ and \tilde{z}_e for vertices v' and edges e above the vertex v . \square

Let us be precise for the claim in the proof. Assume the region walk is given by

$$\text{leftmost} = r_0 \underset{\tilde{z}_1, u_1}{\overset{i_1}{|}} r_1 \underset{\tilde{z}_2, u_2}{\overset{i_2}{|}} r_2 \underset{\tilde{z}_3, u_3}{\overset{i_3}{|}} \cdots \underset{\tilde{z}_{l-1}, u_{l-1}}{\overset{i_{l-1}}{|}} r_{l-1} \underset{\tilde{z}_l, u_l}{\overset{i_l}{|}} r$$

we have

$$\begin{aligned} \tilde{g}_r &= B_{i_1}(\tilde{z}_1, u_1) B_{i_2}(\tilde{z}_2, u_2) \cdots B_{i_l}(\tilde{z}_l, u_l) \\ &= B_{i_1}(\tilde{z}_1) \chi_{i_1}(u_1) B_{i_2}(\tilde{z}_2) \chi_{i_2}(u_2) \cdots B_{i_l}(\tilde{z}_l) \chi_{i_l}(u_l) \\ &= B_{i_1}(z'_1) B_{i_2}(z'_2) \cdots B_{i_l}(z'_l) \cdot u_1^{s_{i_1} \cdots s_{i_2} \alpha_1^\vee} u_2^{s_{i_2} \cdots s_{i_3} \alpha_2^\vee} \cdots u_l^{\alpha_l^\vee} \end{aligned}$$

with

$$z'_k = \tilde{z}_k u_1^{\langle s_{i_{k-1}} \cdots s_{i_2} \alpha_1^\vee, \alpha_{i_k} \rangle} u_2^{\langle s_{i_{k-1}} \cdots s_{i_3} \alpha_2^\vee, \alpha_{i_k} \rangle} \cdots u_k^{\langle \alpha_{k-1}^\vee, \alpha_{i_k} \rangle}.$$

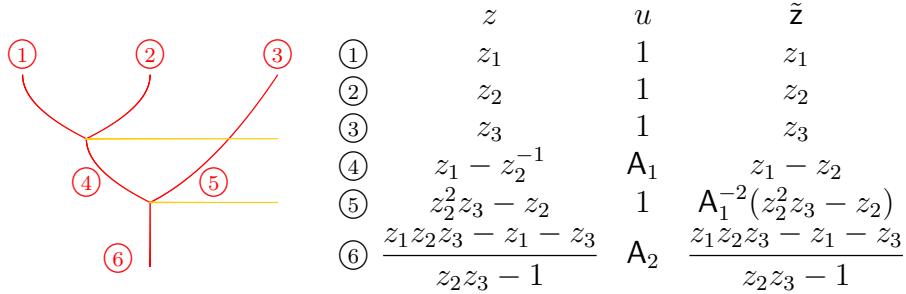
Note that we have to have $z_k = z'_k$ since (z'_k) and (z_k) parametrizes the same element in $X(\mathfrak{W})$.

1.3. Cluster chart. We will call the isomorphism

$$(\mathbb{C}^\times)^{\mathfrak{M}_3} \xrightarrow{\sim} \tilde{X}(\mathfrak{W}^=) \xleftarrow{\sim} X(\mathfrak{W})$$

the **cluster chart**.

Example. Let us consider the example



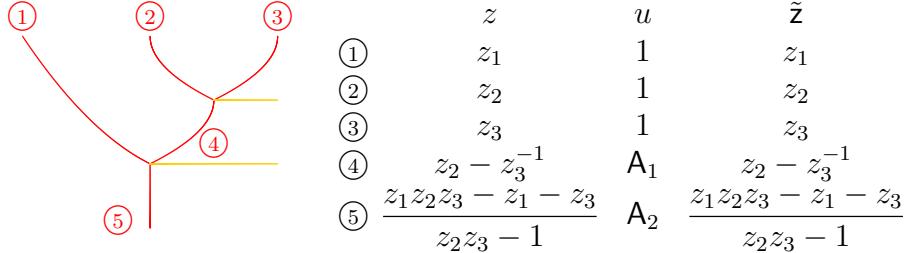
At vertices 1, 2, we have

$$z_2 = \tilde{z}_2 = \frac{u_{\textcircled{4}}}{u_{\textcircled{1}} u_{\textcircled{2}}} = A_1, \quad A_1^{-2}(z_2^2 z_3 - z_2) = \tilde{z}_5 = \frac{u_{\textcircled{6}}}{u_{\textcircled{4}} u_{\textcircled{5}}} = A_2 / A_1.$$

So

$$A_1 = z_2, \quad A_2 = z_2 z_3 - 1.$$

Let us consider another example



At vertices 1, 2, we have

$$z_3 = \tilde{z}_3 = \frac{u_{\textcircled{4}}}{u_{\textcircled{2}} u_{\textcircled{3}}} = A_1, \quad z_2 - z_3^{-1} = \tilde{z}_{\textcircled{4}} = \frac{u_{\textcircled{5}}}{u_{\textcircled{1}} u_{\textcircled{4}}} = A_2 / A_1.$$

So

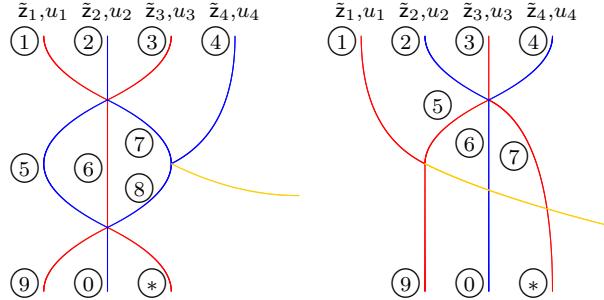
$$A_1 = z_3, \quad A_2 = z_2 z_3 - 1.$$

Theorem. For two weaves $\mathfrak{W}, \mathfrak{M}$ related by two-braid relations and 1212 relations, we have a natural bijection $\mathfrak{W}_3 \cong \mathfrak{M}_3$, the cluster charts

$$(\mathbb{C}^\times)^{\mathfrak{W}_3} \longrightarrow X(\beta)$$

are the same.

Proof. The proof is similar to the the proof we explained last time. The only computation we need to do is 1212 relations. Consider



We know $u_{(9)}, u_{(0)}, u_{(*)}$ agree. By an argument similar to z -labeling, we know $\tilde{z}_{(9)}, \tilde{z}_{(0)}, \tilde{z}_{(*)}$ agree. From the way we construct the inverse map in Theorem B, it suffices to show A agrees.

For the left diagram, we have

$$\tilde{z}_4 = \tilde{z}_{(4)} = A \frac{u_{(4)} \oplus u_{(7)}}{u_{(4)} u_{(7)}} = A \frac{\frac{u_1 u_2}{u_1 \oplus u_3} \oplus u_4}{\frac{u_1 u_2}{u_1 \oplus u_3} u_4} = A \frac{u_1 u_2 \oplus u_1 u_4 \oplus u_3 u_4}{u_1 u_2 u_4}.$$

For the right diagram, we have

$$\tilde{z}_{(5)} = A \frac{u_{(1)} \oplus u_{(5)}}{u_{(1)} u_{(5)}} = A \frac{u_1 \oplus \frac{u_3 u_4}{u_2 \oplus u_4}}{u_1 \frac{u_3 u_4}{u_2 \oplus u_4}} = A \frac{u_1 u_2 \oplus u_1 u_4 \oplus u_3 u_4}{u_1 u_3 u_4}.$$

By **(Comm)**, we have

$$\tilde{z}_4 = \tilde{z}_{(4)} = \frac{u_{(3)} \tilde{z}_{(5)}}{u_{(2)}} = \frac{u_3}{u_2} \tilde{z}_{(5)}.$$

We get the same equation for A . □

REFERENCES

- [1] Roger Casals, Eugene Gorsky, Mikhail Gorsky, Ian Le, Linhui Shen, José Simental. Cluster structures on braid varieties. *Journal of the American Mathematical Society*, 2025. [arXiv:2207.11607](#). 1
- [2] George Lusztig. Total Positivity in Reductive Groups 1.1

APPENDIX A. EXPLICIT SOLUTIONS

We can solve the equations explicitly.

	$\left\ \begin{array}{c} B_i(\tilde{z}_1, u_1) \rightarrow B_i(\tilde{z}_2, u_2) \\ B_i(\tilde{z}_3, u_3) \end{array} \right\ \quad \tilde{U}$	$\tilde{z}_2 = \frac{u_3}{u_1 u_2}, \quad \tilde{z}_1 = \tilde{z}_3 + \frac{u_2}{u_1 u_3},$ $\tilde{U} = \text{image of } \begin{bmatrix} 1 & -\frac{u_1}{u_2 u_3} \\ 0 & 1 \end{bmatrix}$
	$\left\ \begin{array}{c} B_i(\tilde{z}_1, u_1) \rightarrow B_j(\tilde{z}_2, u_2) \\ B_j(\tilde{z}'_2, u'_2) \end{array} \right\ \quad \tilde{U}$	$\tilde{z}_1 = \tilde{z}'_1$ $\tilde{z}_2 = \tilde{z}'_2$ <div style="border: 1px solid black; padding: 5px; display: inline-block;"> $u'_1 = u_1$ $u'_2 = u_2$ </div>
	$\left\ \begin{array}{c} B_i(\tilde{z}_1, u_1) \rightarrow B_j(\tilde{z}_2, u_2) \rightarrow B_i(\tilde{z}_3, u_3) \\ B_j(\tilde{z}'_1, u'_1) \rightarrow B_i(\tilde{z}'_2, u'_2) \rightarrow B_j(\tilde{z}'_3, u'_3) \end{array} \right\ \quad \tilde{U}$	$\tilde{z}_1 = \frac{u'_1 \tilde{z}'_3}{u'_2},$ $\tilde{z}_3 = \frac{u_2 \tilde{z}'_1}{u_1},$ $\tilde{z}_2 = \frac{u_1 u'_1 \tilde{z}'_1 \tilde{z}'_3 - u_1 \tilde{z}'_2}{u'_2 u'_1} = \frac{u_3 u'_3 \tilde{z}'_1 \tilde{z}'_3 - u_1 \tilde{z}'_2}{u'_2 u'_1}$ <div style="border: 1px solid black; padding: 5px; display: inline-block;"> $u'_2 u'_3 = u_1 u_2$ $u'_1 u'_2 = u_2 u_3$ </div>
	$\left\ \begin{array}{c} B_i(\tilde{z}, u) \rightarrow \\ B_i(\tilde{z}', u') \end{array} \right\ \quad \tilde{U} \quad \tilde{U}'$	$\tilde{z} = \tilde{z}' - \xi_i(U),$ $\tilde{U}' = B_i(\tilde{z}', u')^{-1} \tilde{U} B_i(\tilde{z}, u) \in \text{Rad}(B)$ <div style="border: 1px solid black; padding: 5px; display: inline-block;"> $u' = u$ </div>

Recall $\xi_i(U) = \text{coefficient of } E_i \text{ in } U$. Notice that $\xi_i(U') = 0$. This fact will be used later.

All the relations follow from the computation in $SL_2 \subset GL_2$, $SL_2 \times SL_2 \subset GL_4$, $SL_3 \subset GL_3$ and $SL_2 \subset GL_2$. The following is the code.

```
R.<u1,u2,u3p,z3p,z1,z2,z3,a> = QQ[];  
B = lambda z,u: matrix([[u*z,-u^(-1)],[u,0]])  
U = matrix([[1,a],[0,1]]);  
Rel = (B(z1,u1)*B(z2,u2) - B(z3,u3p)*U).change_ring(SR)  
solve([Rel[i][j]==0 for i in range(2) for j in range(2)], SR.var("u1p,z1,z2,a"))
```

```
[[u1p == r1,  
 z1 == (u1*u3p*z3 + u2)/(u1*u3p),  
 z2 == u3p/(u1*u2),  
 a == -u1/(u2*u3p)]]
```

```
R.<u1,u2,u1p,u2p,z1p,z2p,z1,z2> = QQ[]  
B1 = lambda z,u: matrix([[u*z,-u^(-1),0,0],[u,0,0,0],[0,0,1,0],[0,0,0,1]])  
B3 = lambda z,u: matrix([[1,0,0,0],[0,1,0,0],[0,0,u*z,-u^(-1)],[0,0,u,0]])  
Rel = (B1(z1,u1)*B3(z2,u2)-B3(z2p,u2p)*B1(z1p,u1p)).change_ring(SR)
```

```
solve([Rel[i][j]==0 for i in reversed(range(4)) for j in range(4)], SR.var("u1p,u2p,z1,z2"))
```

```
[[u1p == u1, u2p == u2, z1 == z1p, z2 == z2p]]
```

```
R.<u1,u2,u3,u1p,u2p,u3p,z1p,z2p,z3p,z1,z2,z3> = QQ[]
B1 = lambda z,u: matrix([[u*z,-u^(-1),0],[u,0,0],[0,0,1]])
B2 = lambda z,u: matrix([[1,0,0],[0,u*z,-u^(-1)],[0,u,0]])
Rel = (B1(z1,u1)*B2(z2,u2)*B1(z3,u3)-B2(z1p,u1p)*B1(z2p,u2p)*B2(z3p,u3p)).change_ring(SR)
solve([Rel[i][j]==0 for i in reversed(range(2)) for j in range(2)], SR.var("u1p,u2p,u3p,z1,z2,z3"))
```

```
[[u1p == r1*r2/z3p,
u2p == r1,
u3p == r1*r2*u1/(u3*z3p),
z1 == r2,
z2 == (r1^2*r2^2*u1*z1p - r1*u1*z2p*z3p)/(u2*u3*z3p),
z3 == r1^2*r2*z1p/(u1*u3*z3p)]]
```

```
R.<u,up,z,zp,a,ap> = QQ[]
B = lambda z,u: matrix([[u*z,-u^(-1)],[u,0]])
U = matrix([[1,a],[0,1]]); Up = matrix([[1,ap],[0,1]])
Rel = (U*B(z,u)-B(zp,up)*Up).change_ring(SR)
solve([Rel[i][j]==0 for i in range(2) for j in range(2)], SR.var("up,z,a,ap"))
```

```
[[up == u, z == r1, a == -r1 + zp, ap == 0]]
```

To convince the readers, we write down the matrix equations for the trivalent and hexavalent vertices

$$\begin{bmatrix} u_1 u_2 \tilde{z}_1 \tilde{z}_2 - u_1^{-1} u_2 & -u_1 u_2^{-1} \tilde{z}_1 \\ u_1 u_2 \tilde{z}_2 & -u_1 u_2^{-1} \end{bmatrix} = \begin{bmatrix} u_1 \tilde{z}_1 & -u_1^{-1} \\ u_1 & 0 \end{bmatrix} \begin{bmatrix} u_2 \tilde{z}_2 & -u_2^{-1} \\ u_2 & 0 \end{bmatrix}$$

$$\begin{bmatrix} u \tilde{z} & u \tilde{z} a - u^{-1} \\ u & u a \end{bmatrix} = \begin{bmatrix} u \tilde{z} & -u^{-1} \\ u & 0 \end{bmatrix} \begin{bmatrix} 1 & a \\ 0 & 1 \end{bmatrix}.$$

$$\begin{bmatrix} u_1 u_3 \tilde{z}_1 \tilde{z}_3 - u_1^{-1} u_2 u_3 \tilde{z}_2 & -u_1 u_3^{-1} \tilde{z}_1 & u_1^{-1} u_2^{-1} \\ u_1 u_3 \tilde{z}_3 & -u_1 u_3^{-1} & 0 \\ u_2 u_3 & 0 & 0 \end{bmatrix} = \begin{bmatrix} u_1 \tilde{z}_1 & -u_1^{-1} & 0 \\ u_1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & u_2 \tilde{z}_2 & -u_2^{-1} \\ 0 & u_2 & 0 \end{bmatrix} \begin{bmatrix} u_3 \tilde{z}_3 & -u_3^{-1} & 0 \\ u_3 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} u_2' \tilde{z}_{2'} & -u_{2'}^{-1} u_3' \tilde{z}_{3'} & u_{2'}^{-1} u_{3'}^{-1} \\ u_1' u_2' \tilde{z}_{1'} & -u_{1'}^{-1} u_{3'} & 0 \\ u_1' u_2' & 0 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & u_1' \tilde{z}_{1'} & -u_{1'}^{-1} \\ 0 & u_{1'} & 0 \end{bmatrix} \begin{bmatrix} u_2' \tilde{z}_{2'} & -u_{2'}^{-1} & 0 \\ u_{2'} & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & u_3' \tilde{z}_{3'} & -u_{3'}^{-1} \\ 0 & u_{3'} & 0 \end{bmatrix}.$$