Matching Theory

Definition

A subset M of E is called a matching in G if its elements are links and no two are adjacent in G: A matching M saturates a vertex v, and v is said to be M-saturated, if some edge of M is traident with v.

• If every vertex of CT is

M-saturated, the matching M

is (perfect.)

Definition

An M-matching path in Gris a path whose edges are alternately in EIM and M.

An M-augmenting path is an M-alternating path whose origin and terminus are M-unsaturated.

Theorem (Berge, 1957)

A matching M in G is a maximum matching if and only of Grantains no M-augmenting path.

Proof:

→> ∨

Esuppose Mis not maximum

Let M' in G IMI>IMI

Let H= CT[M&M'] which is a

edge-induced subgraph.

Every ventex of H has degree

either 1 or 2 in H. Thus each

component of H is either an

even cycle with edges

alternately in M and M', or else a path with edges alternately in M and M'.

"! [M'| > [M], therefore some path component P of H must start and end with edges of M'.

Thus I an M-augmenting path.

Perfect Matching

Definition

A component of a graph is odd or even according as it has an odd or even number of vertices. We denote by O(G) the number of odd components.

Theorem (Tutte, 1947) Ghas a perfect matching if and only if $o(G-S) \leq |S|$ for all $S \subset V$.

Corollary C Peterson, 1891)

Every 3-regular without cut edges has a perfect matching.

Matching in Bipartite

Theorem CHall, 1935)

Let G be a bipartite graph with bipartite (X, Y). Then G contains a matching that saturates every vertex in X if and only if [N(S)] > |S| For all SEX

Definition

A R-factor of a graph is a spanning k-regular subgraph,

- a R-factorization partitions the edges of the graph into disjoint R-factors.
- A graph G is suid to be k-factorable of it admits a kfactorization.

Algorithm

To find maximum matching in Bipartite graph.

· Hungarian Algorithm

To find maximum Weights matching in weighted Sipartite Graph.

· Kuhn - Munkers Algorithm