CS170– Spring 2022— Homework 5

CurMack

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Knapsack with repetition

Define:

K(w) = maximum value achievable with a knapsack of capacity w.

express this in terms of smaller subproblems:

$$k(w) = \max_{i:w_i \le w} \{K(w - w_i) + v_i\}$$

where as usual our convention is that the maximum over an empty set is 0. So the algorithm:

```
1 K(0) \leftarrow 0

2 for w = 1 : W do

3 \lfloor K(w) \leftarrow \max\{K(w - w_i) + v_i : w_i \leq w\}

4 return K(W)
```

Runtime: $\mathcal{O}(nW)$

Knapsack without repetition

Definr:

 $K(w,j) = \text{maximum value achievable using a knapsack of capacity } w \text{ and items } 1, \dots, j.$

The answer we seek is K(W, n).

express a subproblem K(w, j) in terms of smaller subproblems:

$$K(w,j) = \max\{K(w-w_j, j-1) + v_j, K(w, j-1)\}.$$

So the algorithm:

```
1 Initialize all K(0,j) \leftarrow 0 and all K(w,0) \leftarrow 0

2 for j=1:n do

3 | for w=1:W do

4 | if w_j > w then

5 | K(w,j) \leftarrow K(w,j-1)

6 | else

7 | K(w,j) \leftarrow \max\{K(w,j-1),K(w-w_j,j-1)+v_j\}

8 return K(W,n)
```

Runtime: $\mathcal{O}(nW)$

2 Copper Pipes

Main Idea: Knapsack with repetition:

```
Algorithm 1: cutPipe(price, n)
```

```
1 val(0) \leftarrow 0

2 for w = 1 : n do

3 max\_val \leftarrow 0

4 for i = 1 : w do

5 max\_val \leftarrow max\{val(w - i) + price(i), max\_val\}

6 val(i) \leftarrow max\_val

7 return val(n)
```

Runtime: $\mathcal{O}(n^2)$

3 Egg Drop

- (a) f(1,k) = 1: drop the egg from the single floor.
 - f(0,k) = 0: there is only one possible value for l.
 - f(n,1) = n: drop it from floor 1 and going up, until it breaks.
 - $f(n,0) = \infty$: the problem is unsolvable if we have no eggs to drop.
- (b) the recurrence relation is:

$$f(n,k) = \min_{x \in \{1...n\}} \max\{f(x-1,k-1), f(n-x,k)\}$$

Suppose we drop at floor x:

- breaks: we know $l \in [0, x-1]$, and we have k-1 eggs $\Rightarrow f(x-1, k-1)$
- doesn't break: now $l \in [x, n], k \text{ eggs} \Rightarrow f(n x, k)$

And we pick the beat x