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Q1 Signal :

a. Every process will be terminated, because in kill() function call, pid is negative means, it will sent to every child of that parent's process group. So, they all have pgid = 10. So, all will be kill.

b.
```
if ((pid = fork()) == 0)   // 2nd child
{
    printf("child 2");
    setpgid(getpid(), 12);
}
```

→ If same kill will execute as (a.) then parent & child 1 will be kill & not child 2 because child 2's pgid is changed.

c.
```
#include <signal.h>
#include <stdio.h>
#include <sys/wait.h>
#include <errno.h>
#include <sys/types.h>
#include <unistd.h>
```

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```
void alarmHandler ();
void (* oldhandler) ();
void disableCritical();

void disableCritical()
{
    printf (" Normal section");
    kill (getpgrp()*(-1), SIGUSR2);
    signal (SIGALRM, alarmHandler);
    alarm(5);
}

void alarmHander() {
    printf ("Critical section");
    kill (getpgrp()*(-1) SIGUSR1);
    signal (SIGALRM, disableCritical);
    alarm(3);
}

void siguserhandler (int signum)
{
    switch (signum)
    {
        case SIGUSR1:
            oldhandler = signal (SIGINT, SIG_IGN);
            break;
        case SIGUSR2:
            signal (SIGINT, oldhandler);
            break;
    }
}
```

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```
int main(void)
{
    int status;
    int pid = fork();
    if (pid != 0) {
        signal (SIGUSR1, sigusrhandler);
        wait (& status);
    }
    else
    {
        signal (SIGUSR1, sigusrhandler);
        signal (SIGUSR2, sigusrhandler);
        printf ("child");
        signal (SIGALRM, alarmHandler);
        alarm (5);
        while (1);
    }
}
```

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Q2 Concurrent Programming :

a. Thread 0
Thread 1

or Thread 1
Thread 0

→ Because, it is concurrent so sequence may differe.

b. Thread 2
Thread 2

or Thread 1
Thread 2

→ Because, it is concurrent so output may differ

c. Thread 2
Thread 2

d. Thread 0
Thread 1

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e.

Line 1 : sem_init (&mutex, 0, 1);
Line 2 : sem_init (&items, 0, 1);
Line 3 : sem_init (&slots, 0, 1);

Line 4 : sem_wait (&items);
Line 5 : sem_wait (&mutex);

Line 6 : sem_post (&items);
Line 7 : sem_post (&mutex);

Line 8 : sem_wait (&items);
Line 9 : sem_wait (&slots);

Line 10 : sem_post (&items);

Line 11 : sem_post (&slots);

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Q3 Network :

a.

Client	Server
Socket	Socket
	bind
	listen
Connect	accept

b. Socket : It will use to create socket file descriptor. It will create on both client & server side.
- It is end point of connection
- And socketfd is listenfd for server side & clientfd for client side

Bind : It will bind the socket address of server to file descriptor that

Bind : It will bind the listenfd (server side) to the server socket address.

listenfd : Listenfd is where the server is waiting for the client request & so as soon as the request is accepte accept() call will create a new fd.

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accept : As soon as request is accepted from client it will create a new file descriptor called connfd. Blocking call

connect : It will send / connection request to

connect : It will communicate between client & server with all details.

C.

```
main(){

    int clb = 0; int cur = 0;
    int port = 500;

    while (clb<4) {
        clb++;
        getaddrinfo (NULL, port, &hint,
                     &listf);
        listenfd = socket (listp->ai_family,
                           listp->ai_socktype,
                           listp->ai_protocol);
        bind (listenfd, 500);
        while(1) {
            connfd = accept (listenfd,
                     (struct sockaddr), &clientaddr,
                     &clientlen);
            cur++;
```

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```
pthread_create( &tid , NULL, Bundle_Client,
                                    commfd );
if ( cur == 500 ) {
    cli = 0;
    cli_cb ++ ;
    break; }
}
}
```

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Q4 Device Driver :

a. Major Number represents which driver will used to manage a particular device. It allows multiple devices to have the same major number if they are managed by the same driver.

Exc : there are 4 devices represented as minor numbers 0, 10, 11, 12 by they are manage by same driver number 4.

```
$ ls -ltr /dev
crw--w---1 root tl 4, 0 M tl 0
crw--w---1 root tl 4, 10 tl 10
crw--w---1 root tl 4, 12 tl 12
crw--w---1 root tl 4, 11 tl 11
```

b. kmalloc is used to dynamically allocate memory to buffers in kernal code.
Yes, we use this same function for user application by using ~~GFP_USR~~ GFP_USER.