OPERATIONS RESEARCH DOI 10.1287/opre.1080.0550ec pp. ec1-ec5



e - c o m p a n i o n ONLY AVAILABLE IN ELECTRONIC FORM

Electronic Companion—"Simulation-Based Optimization of Virtual Nesting Controls for Network Revenue Management" by Garrett van Ryzin and Gustavo Vulcano, *Operations Research*, DOI 10.1287/opre.1080.0550.

Simulation-Based Optimization of Virtual Nesting Controls for Network Revenue Management

ONLINE APPENDIX

Garrett van Ryzin¹

Gustavo Vulcano²

A1 Derivatives of the revenue function

From equation (5), the deduction of the derivative of the revenue function with respect to y_{ic} proceeds as follows:

$$\frac{\partial}{\partial y_{ic}} R_t(x(t), y, \omega) = r_{j_t} \frac{\partial}{\partial y_{ic}} u_{j_t}(x(t) - \zeta_t, y, q_t) + \sum_{k=1}^m \frac{\partial}{\partial x_k(t+1)} R_{t+1}(x(t+1), y, \omega) \frac{\partial}{\partial y_{ic}} x_k(t+1) + \frac{\partial}{\partial y_{ic}} R_{t+1}(x(t+1), y, \omega)$$

Take the column A_{j_t} in equation (6). It has a 1 in position k if and only if $k \in A_{j_t}$. Then,

$$\frac{\partial}{\partial y_{ic}} R_t(x(t), y, \omega) = r_{jt} \frac{\partial}{\partial y_{ic}} u_{jt}(x(t) - \zeta_t, y, q_t) - \sum_{k \in A_{j_t}} \frac{\partial}{\partial x_k(t+1)} R_{t+1}(x(t+1), y, \omega) \frac{\partial}{\partial y_{ic}} u_{j_t}(x(t) - \zeta_t, y, q_t)
+ \frac{\partial}{\partial y_{ic}} R_{t+1}(x(t+1), y, \omega)
= \left(r_{j_t} - \sum_{k \in A_{j_t}} \frac{\partial}{\partial x_k(t+1)} R_{t+1}(x(t+1), y, \omega) \right) \frac{\partial}{\partial y_{ic}} u_{j_t}(x(t) - \zeta_t, y, q_t)
+ \frac{\partial}{\partial y_{ic}} R_{t+1}(x(t+1), y, \omega).$$

Now we have to solve for the derivative with respect to capacity, involved in the second term. Taking equation (5), we have that:

$$\frac{\partial}{\partial x_i(t)} R_t(x(t), y, \omega) = r_{j_t} \frac{\partial}{\partial x_i(t)} u_{j_t}(x(t) - \zeta_t, y, q_t)
+ \sum_{k=1}^m \frac{\partial}{\partial x_k(t+1)} R_{t+1}(x(t+1), y, \omega) \frac{\partial}{\partial x_i(t+1)} x(t+1)$$

Now, consider again equation (6). According to the resources used by product j_t , we can rewrite the equation above as:

$$\frac{\partial}{\partial x_i(t)} R_t(x(t), y, \omega) = r_{jt} \frac{\partial}{\partial x_i(t)} u_{jt}(x(t) - \zeta_t, y, q_t)$$

¹Graduate School of Business, Columbia University, 412 Uris Hall, New York, NY 10027, gjv1@columbia.edu.

²Stern School of Business, New York University, 44 West Fourth St., Suite 8-76, New York, NY 10012, gvulcano@stern.nyu.edu.

$$-\sum_{\substack{k \in A_{j_t} \\ k \neq i}} \frac{\partial}{\partial x_k(t+1)} R_{t+1}(x(t+1), y, \omega) \frac{\partial}{\partial x_i(t)} u_{j_t}(x(t) - \zeta_t, y, q_t)$$

$$+ \frac{\partial}{\partial x_i(t+1)} R_{t+1}(x(t+1), y, \omega) \left(1 - \frac{\partial}{\partial x_i(t)} u_{j_t}(x(t) - \zeta_t, y, q_t) \mathbf{I}\{i \in A_{j_t}\}\right)$$

Finally, regrouping terms, we have:

$$\frac{\partial}{\partial x_{i}(t)} R_{t}(x(t), y, \omega) = \left(r_{j_{t}} - \sum_{k \in A_{j_{t}}} \frac{\partial}{\partial x_{k}(t+1)} R_{t+1}(x(t+1), y, \omega) \right) \frac{\partial}{\partial x_{i}(t)} u_{j_{t}}(x(t) - \zeta_{t}, y, q_{t}) + \frac{\partial}{\partial x_{i}(t+1)} R_{t+1}(x(t+1), y, \omega)$$

A2 Displacement adjusted virtual nesting (DAVN)

We implemented displacement adjusted virtual nesting (DAVN) to obtain the indexing and the initial protection levels y in our numerical experiments. The method first uses a deterministic linear program (DLP) to compute a set of static bid prices $\overline{\pi} = (\overline{\pi}_1, \dots, \overline{\pi}_m)$ as follows: Define D_j to be the aggregated demand for product j over the whole time horizon (i.e. $D_j = \sum_{t=1}^T \mathbf{I}\{J_t = j\}$). The DLP is then

$$\max r^{\mathsf{T}} z$$
 s.t. $Az \le x, 0 \le z \le \mathrm{E}[D]$

Then, for all products j that use i, a displacement adjusted revenue \overline{r}_{ij} is computed using

$$\overline{r}_{ij} = r_j - \sum_{k \in A_j, k \neq i} \overline{\pi}_k$$

Roughly, \overline{r}_{ij} is an estimate of the net benefit of accepting product j over leg i, in which the revenue of j is reduced by the static bid price values of the other resources used by j. Note that \overline{r}_{ij} can be negative.

We then proceed to cluster or index the different products on a resource i into a fixed number of virtual classes or buckets. A particular virtual class is denoted $v \in \{1, ..., \bar{c} + 1\}$. A variety of techniques can be used in this step. Ours is a rather simple clustering procedure that consists of evenly partitioning the range of displacement adjusted revenue values on each leg i – one for each virtual class – and assigning product j on leg i to virtual class $c_i(j)$ if it falls in the corresponding interval. More precisely, define:

$$\overline{r}_i^{\min} = \min\left\{\overline{r}_{ij}: j \in A^i\right\} \quad \text{and} \quad \overline{r}_i^{\max} = \max\left\{\overline{r}_{ij}: j \in A^i\right\} + \epsilon, \ \epsilon > 0$$

and the "width" of a virtual class on resource i to be $w_i = (\overline{r}_i^{\text{max}} - \overline{r}_i^{\text{min}})/(\overline{c} + 1)$. The indexing is then described by the mapping:

$$c_i(j) = v$$
 if $\overline{r}_{ij} \in [\overline{r}_i^{\min} + (v-1)w_i, \overline{r}_i^{\min} + vw_i)$

Let $c_i^{-1}(v) = \{j \in A^i : c_i(j) = v\}$ be the set of products on resource i that map to virtual class v. For all virtual classes v such that $c_i^{-1}(v) \neq \emptyset$, a representative revenue is computed by taking the revenue given by the mean demand weighted average:

$$R_i(v) = \left(\frac{\sum_{j \in c_i^{-1}(v)} \overline{r}_{ij} \operatorname{E}[D_j]}{\sum_{j \in c_i^{-1}(v)} \operatorname{E}[D_j]}\right)^+$$

By taking the positive part, we are guaranteed not to get a negative representative revenue for virtual class v over leg i, which could happen due to the \overline{r}_{ij} 's.

Define $\bar{c}_i = |\{v : c_i^{-1}(v) \neq \emptyset\}| - 1$ (i.e. \bar{c}_i is the significant number of virtual classes over leg i, minus one). Virtual classes are then relabeled consecutively so that $R_i(1) \geq R_i(2) \geq \cdots \geq R_i(\bar{c}_i + 1)$. Then, the distribution of total demand for a virtual class is computed by adding the means and variances of the demand of the involved products:

$$\mathrm{E}[D_i(v)] = \sum_{j \in c_i^{-1}(v)} \mathrm{E}[D_j] \quad \text{and} \quad Var[D_i(v)] = \sum_{j \in c_i^{-1}(v)} Var[D_j]$$

Next, we solve a multiclass, single-resource problem based on this data, using the EMSR-b heuristic of Belobaba [3]. This procedure yields a set of protection levels for the virtual classes at each resource i. The heuristic is based on the assumption that demand arrives from lowest to highest revenue in the order of virtual classes (i.e. demand for virtual class $\bar{c}_i + 1$ arrives first, followed by demand for virtual class \bar{c}_i , and so on, until finally demand for virtual class 1 is realized). In particular, in stage k + 1 in which demand for virtual class k + 1 arrives and we want to set protection level y_{ik} for classes k and higher (in terms of revenue), DAVN computes the future aggregated demand to come

$$S_{ik} = \sum_{v=1}^{k} D_i(v)$$

Let the weighted average revenue from classes $1, \ldots, k$ over leg i, denoted \overline{R}_{ik} , be defined by

$$\overline{R}_{ik} := \frac{\sum_{v=1}^{k} R_i(v) \operatorname{E}[D_i(v)]}{\sum_{v=1}^{k} \operatorname{E}[D_i(v)]}$$

The EMSR-b protection level over leg i for virtual classes k and higher, y_{ik} , is chosen so that

$$P(S_{ik} > y_{ik}) = \frac{R_i(k+1)}{\overline{R}_{ik}}$$

In particular, if demand for each virtual class v is independent and normally distributed, then

$$y_{ik} = \mu_{S_{ik}} + z_{\alpha} \sigma_{S_{ik}}$$

where $\mu_{S_{ik}} = \sum_{v=1}^k \mathrm{E}[D_i(v)]$, $\sigma_{S_{ik}}^2 = \sum_{v=1}^k Var[D_i(v)]$ and $z_{\alpha} = \Phi^{-1}(1 - R_i(k+1)/\overline{R}_{ik})$. In case $y_{ic} > x_i$, we set $y_{ic} = x_i$. Finally, we round the vector y to end up with integer protection levels.

A3 Pseudo-code for sample path gradient calculation

```
procedure getgrads(x,j,q,r,T,y);
/* input definitions
x[i] = remaining capacity on resource i (modified by program)
j[t] = product requested by customer t on sample path
q[t] = quantity requested by customer t on sample path
r[j] = revenue generated per unit sold of product j
T = total number of customer requests
y[i,c] = protection level for class c on resource i. We assume y[i,0]=0
/*
/* local variables
min_space[t] = minimum remaining capacity for customer t
u[t] = quantity accepted for customer t
gradX[i] = sample path subgradient with respect to x
gradY[i,c] = sample path subgradient with respect to y
/* functions used /*
GextNextRes(j,n); /* returns the n-th resource used by product j;
                        or -1 if less than n resources are used by j /*
Min(a,b);
                  /* returns minimum of a and b /*
/* initialize gradients to zero /*
for (i:=1 to m)
 {gradX[i]:=0.0;
  for (c:=1 to cbar)
       gradY[i,c]:=0.0;
/* forward pass /*
for (t:=1 to T)
 {MIN:=INFTY;
  j:=j[t];
  q:=q[t];
  n:=1;
              /* compute minimum remaining capacity for j /*
  while (i:=GetNextRes(j,n) NEQ -1)
    {if (x[i]-y[i,c[i,j]-1] < MIN)
        {if (x[i]-y[i,c[i,j]-1] >= 0)
           MIN:=x[i]-y[i,c[i,j]-1]; /* save the new minimum remaining capacity for j /*
         else
           MIN:=0; /* protection level for resource i is beyond the remaining capacity /*
       }
```

```
n:=n+1;
    }
   min_space[t]:=MIN;  /* save minimum remaining capacity /*
   u[t] := min(q,MIN); /* accepted amount is minimum of q and MIN /*
  n := 1;
  while (i:=GetNextRes(j,n) NEQ -1) /* update remaining capacity /*
    {x[i] := x[i] - u[t];}
     n := n+1;
    }
 }
/* backward pass /*
for (t:=T down to 1)
 {j:=j[t];
  q:=q[t];
  MIN:=min_space[t]; /* retrieve minimum remaining capacity for j /*
  SUM:=0.0;
  n:=1; /* go through resources used by product j /*
  while (i:=GetNextRes(j,n) NEQ -1)
    {x[i]:=x[i]+u[t]; /* reconstruct capacity used by t on resource i /*}
     SUM:=SUM + gradX[i]; /* add up partial with respect to in x on path of j */
     n:=n+1;
    }
   if ( 0 < MIN <= q)
    {n:=1};
     while (i:=GetNextRes(j,n) NEQ -1)
        /* gradients affected only if resource i is binding /*
        \{if (x[i]-y[i,c[i,j]-1] == MIN)
           {gradX[i]:=gradX[i] + r[j] - SUM; /* update the gradX for next iteration /*}
           for (c:=1 to c[i,j]-1) /* search for protection levels for classes higher than c[i,j] /*
               {if (x[i]-y[i,c] == MIN) /* update gradY if y[i,c] is binding /*
                   gradY[i,c]:=gradY[i,c] - r[j] + SUM;
              }
          }
        n:=n+1;
       }
    }
 }
```