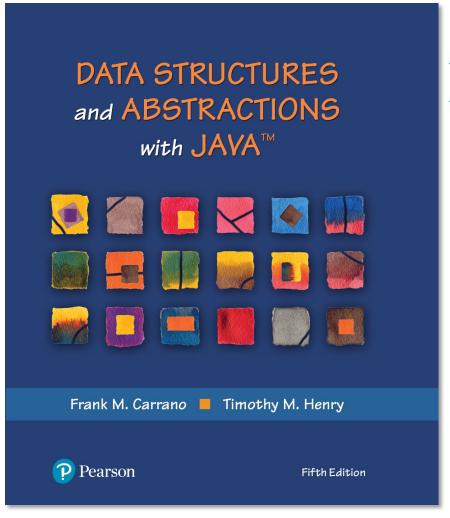
#### **Data Structures**

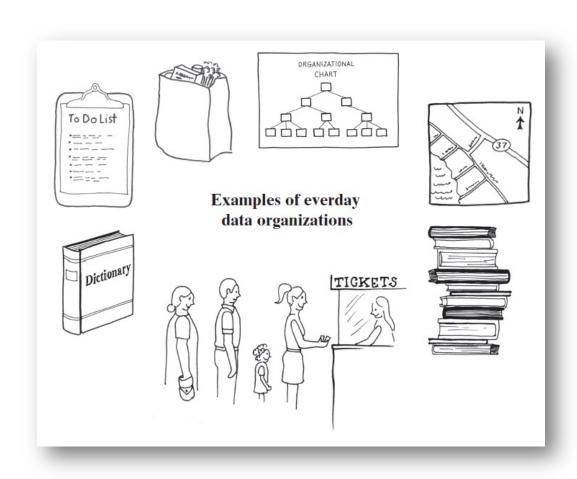


### Introduction



### Data Organization in Life

- Standing in a line
- Stack of books
- To-Do list
- Dictionary
- Folders, directories on your computer
- Road map



### Computer Data Organization

- Abstract Data Type: ADT
- Data Structure
- Collection
- Examples of containers
  - Bag
  - List
  - Stack
  - Queue

- Dictionary
- Tree
- Graph



### **Problems with Array Implementation**

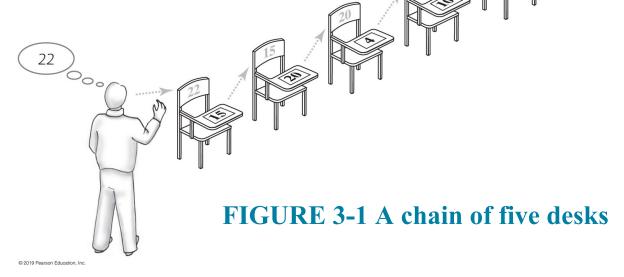
- Array has fixed size
- May become full
- Alternatively may have wasted space
- Resizing is possible but requires overhead of time



### **Analogy**

- Empty classroom
- Numbered desks stored in hallway
  - Number on back of desk is the "address"
- Number on desktop references another desk in chain of desks

Desks are linked by the numbers





### Forming a Chain by Adding to Its Beginning

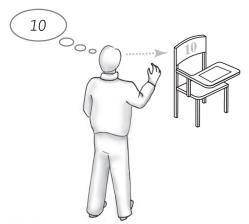


FIGURE 3-2
One desk in
the room

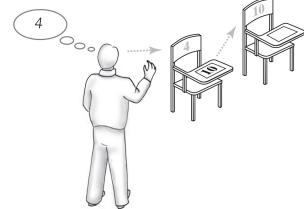


FIGURE 3-3
Two linked desks, with the newest desk first

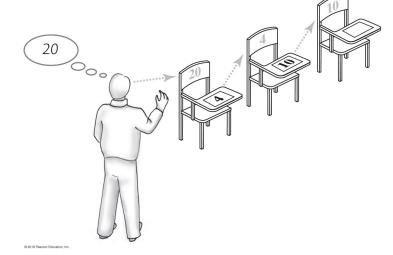


FIGURE 3-4
Three linked desks, with the newest desk first



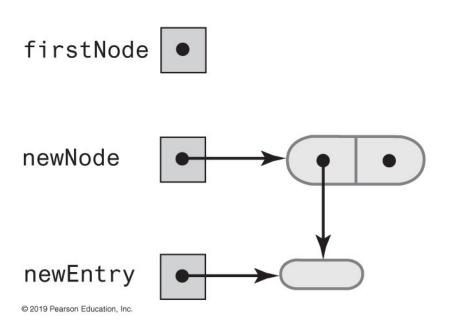
### What Is an Iterator?

- An object that traverses a collection of data
- During iteration, each data item is considered once
  - Possible to modify item as accessed
- Should implement as a distinct class that interacts with the ADT

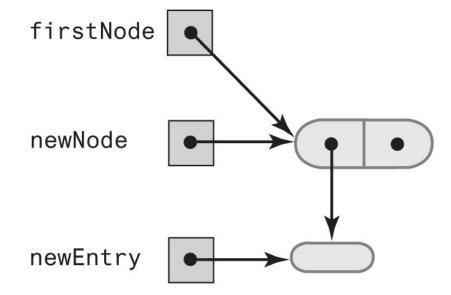


### Beginning a Chain of Nodes

(a) An empty chain and a new node



(b) After adding a new node to a chain that was empty

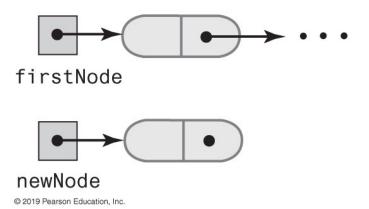


### FIGURE 3-6 Adding a new node to an empty chain

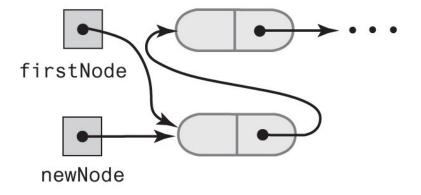


### Beginning a Chain of Nodes

(a) Before adding a node at the beginning



(b) After adding a node at the beginning



## FIGURE 3-7 A chain of nodes just before and just after adding a node at the beginning



#### Case 1:

Your desk is first in the chain of desks.

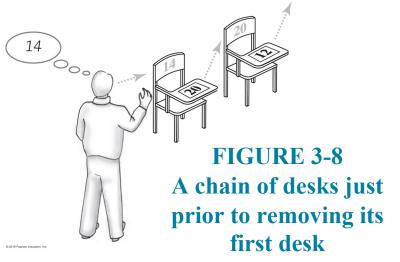
#### • Case 2:

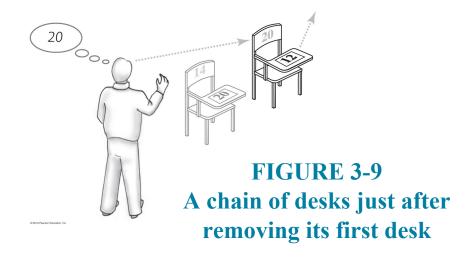
Your desk is not first in the chain of desks.



#### Case 1

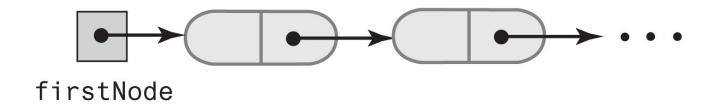
- Locate first desk by asking instructor for its address.
- Give address written on the first desk to instructor.
  - This is address of second desk in chain.
- Return first desk to hallway.







(a) A chain of linked nodes



(b) The chain after its first node is removed

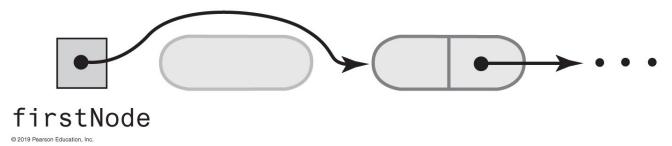


FIGURE 3-10 A chain of nodes just before and just after its first node is removed



#### Case 2

- Move the student in the first desk to your former desk.
- Remove the first desk using the steps described for Case 1.



### Pros of Using a Chain

- Bag can grow and shrink in size as necessary.
- Remove and recycle nodes that are no longer needed
- Adding new entry to end of array or to beginning of chain both relatively simple
- Similar for removal



### Cons of Using a Chain

- Removing specific entry requires search of array or chain
- Chain requires more memory than array of same length



### **Advantages of Linked Implementation**

- Uses memory only as needed
- When entry removed, unneeded memory returned to system
- Avoids moving data when adding or removing entries



### Adding a Node at Various Positions

- Possible cases:
  - Chain is empty
  - Adding node at chain's beginning
  - Adding node between adjacent nodes
  - Adding node to chain's end

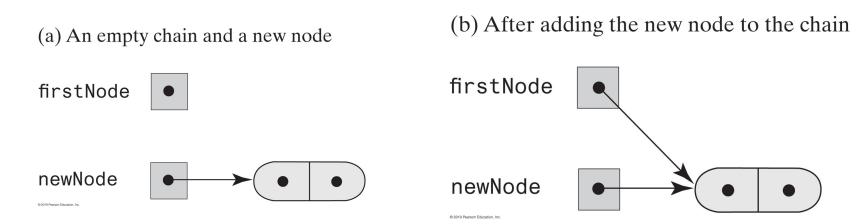


# Adding a Node This pseudocode establishes a new node for the given data

newNode references a new instance of Node

*Place* newEntry *in* newNode

firstNode = address of newNode



#### FIGURE 12-1 Adding a node to an empty chain



This pseudocode describes the steps needed to add a node to the beginning of a chain.

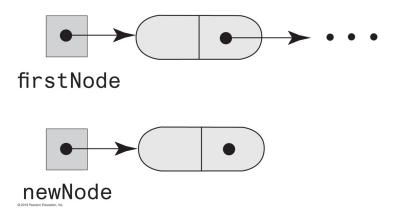
newNode references a new instance of Node

Place newEntry in newNode

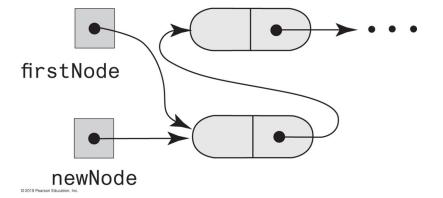
Set newNode's link to firstNode

Set firstNode to newNode

(a) A chain of nodes and a new node



(b) After adding the new node to the beginning of the chain



#### FIGURE 12-2 Adding a node to the beginning of a chain



Pseudocode to add a node to a chain between two existing, consecutive nodes

newNode references the new node

Place newEntry in newNode

Let nodeBefore reference the node that will be before the new node

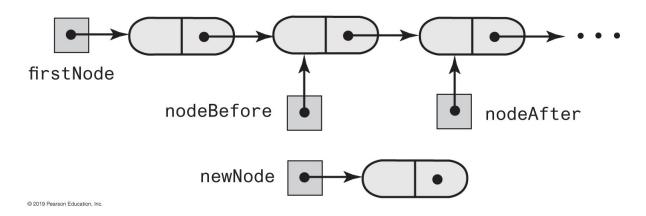
Set nodeAfter to nodeBefore's link

Set newNode's link to nodeAfter

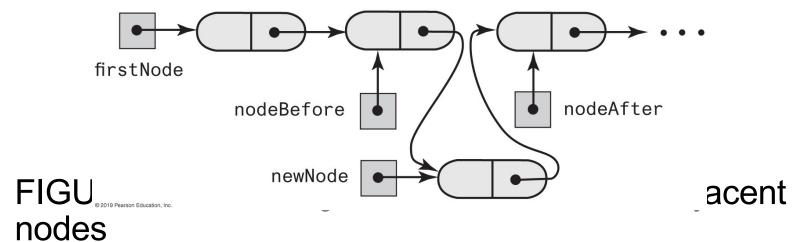
Set nodeBefore's link to newNode



(a) A chain of nodes and a new node



(b) After adding the new node between adjacent nodes





Steps to add a node at the end of a chain.

newNode references a new instance of Node

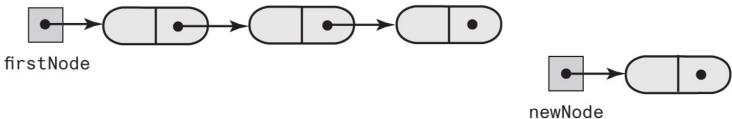
Place newEntry in newNode

Locate the last node in the chain

Place the address of newNode in this last node

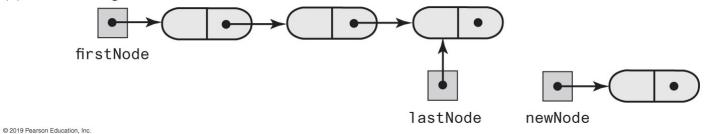


(a) A chain of nodes and a new node

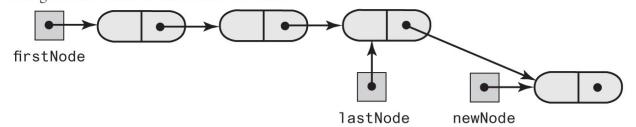


© 2019 Pearson Education, Inc.

(b) After locating the last node



(c) After adding the new node to the end of the chain



2019 Pearson Education, Inc.

#### FIGURE 12-4 Adding a node to the end of a chain



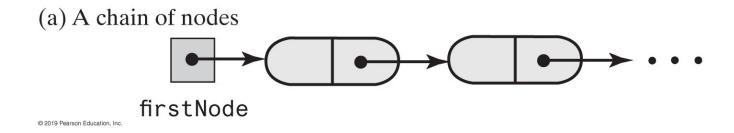
- Possible cases
  - Removing the first node
  - Removing a node other than first one



#### Steps for removing the first node.

Set firstNode to the link in the first node; firstNode now either references the second node or is null if the chain had only one node.

Since all references to the first node no longer exist, the system automatically recycles the first node's memory.



(b) After removing the first node

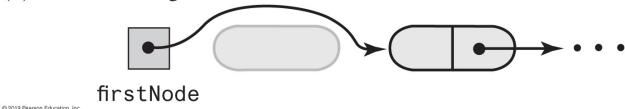


FIGURE 12-5 Removing the first node from a chain



### Removing a node other than the first one.

Let nodeBefore reference the node before the one to be removed.

Set nodeToRemove to nodeBefore's link; nodeToRemove now references the node to be removed.

Set nodeAfter to nodeToRemove's link; nodeAfter now either references the node after the one to be removed or is null.

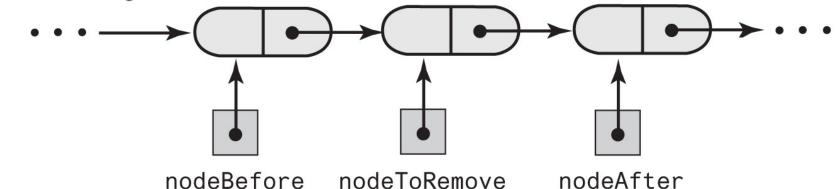
Set nodeBefore's link to nodeAfter. (nodeToRemove is now disconnected from the chain.)

Set nodeToRemove to null.

Since all references to the disconnected node no longer exist, the system automatically recycles the node's memory.

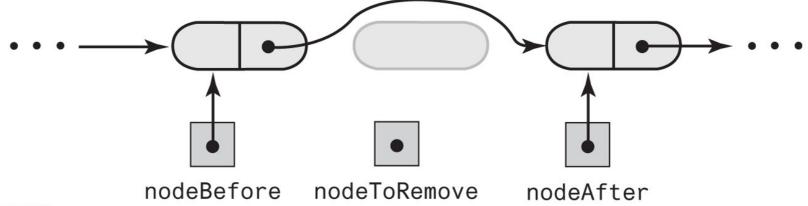


(a) After locating the node to remove



© 2019 Pearson Education, Inc.

(b) After removing the node



© 2019 Pearson Education, Inc.

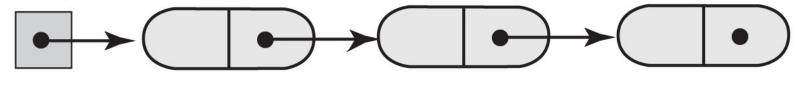
#### FIGURE 12-6 Removing an interior node from a chain



### Using a Tail Reference

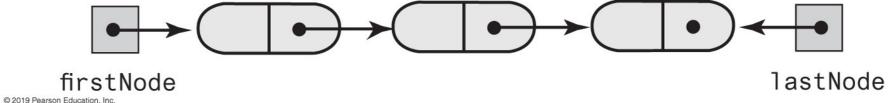
(a) With only a head reference

firstNode



© 2019 Pearson Education, Inc.

(b) With both a head reference and a tail reference



#### FIGURE 12-7 Two linked chains



### A Refined Linked Implementation

```
private Node firstNode;  // Head reference to first node
private Node lastNode;  // Tail reference to last node
private int numberOfEntries; // Number of entries in list
```

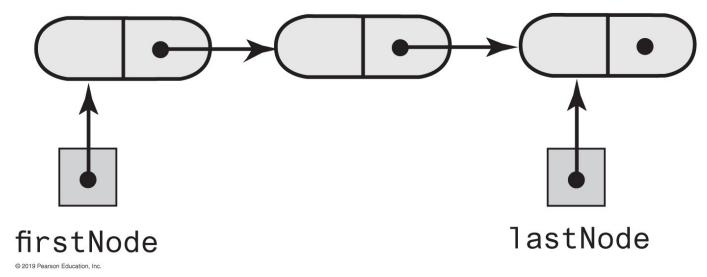
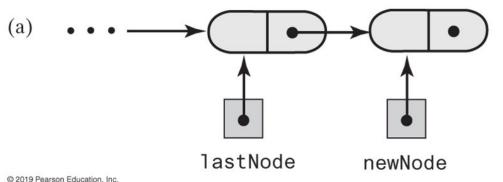


FIGURE 12-8 A linked chain with both a head reference and a tail reference



### A Refined Linked Implementation



After executing lastNode.setNextNode(newNode);

2019 Fearson Education, inc

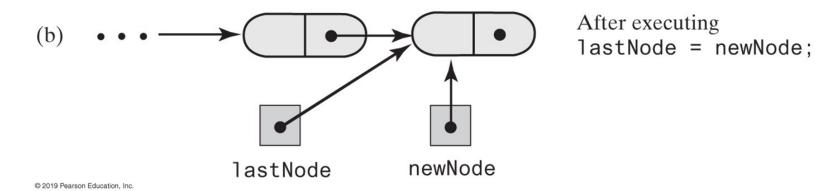


FIGURE 12-9 Adding a node to the end of a nonempty chain that has a tail reference



### A Refined Linked Implementation

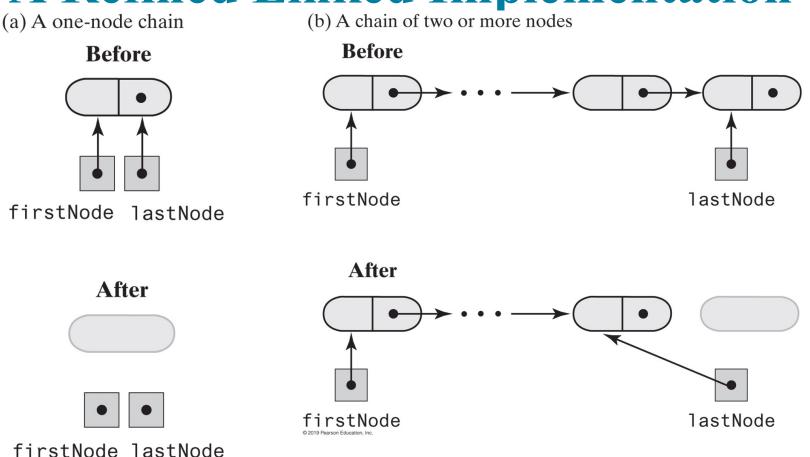


FIGURE 12-10 Before and after removing the last node from a chain that has both head and tail references and contains one or more nodes



### **Testing Core Methods**

#### **Program Output**

```
Create an empty list.
List should be empty; isEmpty returns true.

Testing add to end:
List should contain 15 25 35 45.
The list contains 4 entries, as follows:
15 25 35 45
List should not be empty; isEmpty() returns false.

Testing clear():
List should be empty; isEmpty returns true.
```

LISTING 12-2 A main method that tests part of the implementation of the ADT list

