



**RV
UNIVERSITY**
Go, change the world
an initiative of RV EDUCATIONAL INSTITUTIONS

USN

School of Computer Science and Engineering

B.Tech (Hons.)

CP-1 Question Paper
Academic Year 2024-2025

Course: Compiler Design

Course Code: CS3704

Semester: VI

Time: 2.00PM -3.00PM

Max Marks: 15

Date :13/02/2025

Sl. No.	Questions	Marks	L1-L6	CO
1.	A regular expression for accepting strings with exactly one 1 more than 0's is A. 0^*1 B. $(0/1)^*1(0/1)^*$ C. $(0/1)^*1(0/1)^* 1(0/1)^*$ D. Cannot be framed	1	L3	CO1
2.	Which of the following are Lexemes? A. Identifiers B. Constants C. Keywords D. All of the mentioned	1	L2	CO1
3.	In a lex specification file "?" stands for -----zero or 1 occurrence--	1	L2	CO1
4.	The number of tokens in the following C statement is <code>printf("i = %d, &i = %x", i, &i);</code> A.3 B.11 C.10 D.21	1	L3	CO1
5.	Output file of Lex is the input file is Myfile A. Myfile.asm C. Myfile.yy.c B. Myfile.lex D. Myfile.obj	1	L2	CO1
6.	Identify which one of the following grammars is free from left recursion? Option B A. $S \rightarrow AB$ $A \rightarrow Aa b$ $B \rightarrow c$ B. $S \rightarrow Ab Bb c$ $A \rightarrow Bd \epsilon$ $B \rightarrow e$ C. $S \rightarrow Aa B$ $A \rightarrow Bb Sc \epsilon$ $B \rightarrow d$ D. $S \rightarrow Aa Bb c$ $A \rightarrow Bd \epsilon$ $B \rightarrow Ae \epsilon$	1	L3	CO2

7.	Type checking is normally done during -Semantic Analysis----- phase of compiler.	1	L2	CO1																										
8.	<p>Eliminate left recursion in the productions given below:</p> <p>$S \rightarrow Ba \mid b$</p> <p>$B \rightarrow Bc \mid Sd \mid \epsilon$</p> <p>$S \rightarrow Ba \mid b$</p> <p>$B \rightarrow Bc \mid Bad \mid bd \mid \epsilon$</p> <p>$B \rightarrow bdB' \mid B'$</p> <p>$B' \rightarrow cB' \mid adB' \mid \epsilon$</p>	2	L3	CO2																										
9.	<p>Which one of the following statements is FALSE?</p> <p>A. Context-free grammar can be used to specify both lexical and syntax rules.</p> <p>B. Type checking is done before parsing.</p> <p>C. High-level language programs can be translated to different Intermediate Representations.</p> <p>D. Arguments to a function can be passed using the program stack.</p>	1	L2	CO2																										
10.	<p>Identify tokens generated by the scanner for the following statement and give the total count?</p> <p>$x = x * (a + b) - 5;$</p> <table> <tr> <th>token no</th> <th>character</th> </tr> <tr><td>1</td><td>x</td></tr> <tr><td>2</td><td>=</td></tr> <tr><td>3</td><td>x</td></tr> <tr><td>4</td><td>*</td></tr> <tr><td>5</td><td>(</td></tr> <tr><td>6</td><td>a</td></tr> <tr><td>7</td><td>+</td></tr> <tr><td>8</td><td>b</td></tr> <tr><td>9</td><td>)</td></tr> <tr><td>10</td><td>-</td></tr> <tr><td>11</td><td>5</td></tr> <tr><td>12</td><td>;</td></tr> </table>	token no	character	1	x	2	=	3	x	4	*	5	(6	a	7	+	8	b	9)	10	-	11	5	12	;	1	L3	CO1
token no	character																													
1	x																													
2	=																													
3	x																													
4	*																													
5	(
6	a																													
7	+																													
8	b																													
9)																													
10	-																													
11	5																													
12	;																													

11.	<p>Match all items in Group 1 with those given in Group 2.</p> <table> <tr> <td>Group 1</td> <td>Group 2</td> </tr> <tr> <td>A. Regular expression</td> <td>1. Syntax analysis</td> </tr> <tr> <td>B. Pushdown automata</td> <td>2. Code generation</td> </tr> <tr> <td>C. Dataflow analysis</td> <td>3. Lexical analysis</td> </tr> <tr> <td>D. Register allocation</td> <td>4. Code optimization</td> </tr> </table> <p>A—3 B—1 C—4 D—2</p>	Group 1	Group 2	A. Regular expression	1. Syntax analysis	B. Pushdown automata	2. Code generation	C. Dataflow analysis	3. Lexical analysis	D. Register allocation	4. Code optimization	2	L3	CO1
Group 1	Group 2													
A. Regular expression	1. Syntax analysis													
B. Pushdown automata	2. Code generation													
C. Dataflow analysis	3. Lexical analysis													
D. Register allocation	4. Code optimization													
12.	<p>A CFG G is given with the following productions where S is the start symbol, A is a non-terminal and a and b are terminals.</p> <p>$S \rightarrow aS \mid A$</p> <p>$A \rightarrow aAb \mid bAa \mid \epsilon$</p> <p>For the string "aabbaab" examine how many steps required to derive the string and how many parse trees are there?</p> $ \begin{array}{l} S \xrightarrow{1} aS \\ \xrightarrow{2} aA \\ \xrightarrow{3} aaAb \\ \xrightarrow{4} aabAab \\ \xrightarrow{5} aabbAaab \\ \xrightarrow{6} aabbaab \end{array} $ <p>Thus 6 steps are needed and only one way to derive the string so only one parse tree.</p>	2	L3	CO2										

Course Outcomes

1. Develop skills to devise, select, and apply appropriate tools and techniques for effective compiler design.
2. Apply context-free grammars (CFG) to develop language specifications.

Marks Distribution

L1	L2	L3	L4	L5	L6	CO1	CO2	CO3	CO4
	5	10				9	6		