### Monads

A4M36TPJ, 2013/2014

#### Introduction

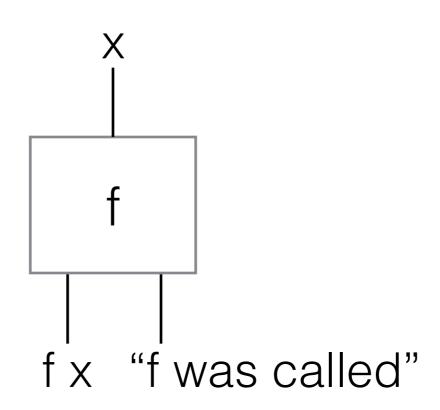
- In pure-functional languages no side-effects are allowed.
- Functions in pure-functional languages depend only on input arguments.
- Monads can be used to simulate (not only) sideeffects.

• We have functions **f** and **g** that both map floats to floats, but we'd like to modify these functions to also output strings for **debugging purposes**.

f,g: Float -> Float

- How can we modify the types of f and g to admit side effects?
- The only possible way is for these strings to be returned alongside the floating point numbers.

f',g' :: Float -> (Float,String)

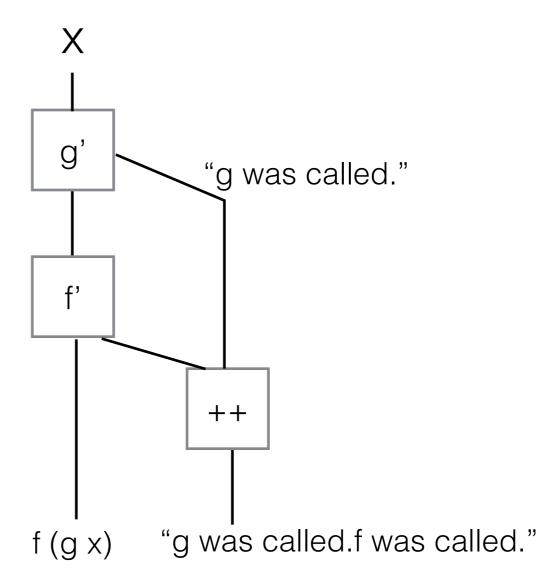


What about function composition?

f' . g'

- These functions cannot be composed straightforward.
- Return type of **g**' is not same as input type of **f**'.

 We would like to compose functions f' and g' same way as f and g.



• To implement previous diagram you can do:

let 
$$(y,s) = g' x$$
  
 $(z,t) = f' y in (z,s++t)$ 

 But you have to do it every time you want to compose functions f' and g'.

- How can we do it easier programmatically?
- We need to find higher-order function which will do this plumbing for us.
- As the problem is that the output from g' can't simply be plugged into the input to f', we need to 'upgrade' f'.

 We introduce new function bind with the following type:

```
bind f' :: (Float, String) -> (Float, String)
```

```
bind :: (Float -> (Float,String)) -> ((Float,String) -> (Float,String))
```

- bind must serve two purposes:
  - It must apply f' to the correct part of g' x.
  - Concatenate the string returned by g' with the string returned by f'.

bind f'(gx,gs) = let(fx,fs) = f'gx in(fx,gs++fs)

- Given a pair of debuggable functions, f' and g', we can now compose them together to make a new debuggable function bind f'.g'.
- We will write this composition as f' \* g'.

- Even though the output of **g'** is incompatible with the input of **f'** we still have a nice easy way to concatenate their operations.
- And this suggests another question: Is there an 'identity' debuggable function?

Identity have the following properties:

$$f.id = f$$
 and  $id.f = f$ 

According that we are looking for the function unit:

 The function unit does not change the output of the function f.

unit 
$$x = (x, "")$$

 The unit allows us to 'lift' any function into a debuggable one.

lift 
$$f x = (f x,"")$$

# Debuggable Functions Summary

 The functions, bind and unit, allow us to compose debuggable functions in a straightforward way, and compose ordinary functions with debuggable functions in a natural way.

Exercise: Show that **lift f \* lift g = lift (f.g)** 

- Consider functions sqrt and cbrt that compute the square root and cube root, respectively, of a real number. These are straightforward functions of type
   Double -> Double.
- Consider a version of these functions that works with complex numbers.
- Every complex number, besides zero, has two square roots. Similarly, every non-zero complex number has three cube roots.

sqrt',cbrt' :: Complex Double -> [Complex Double]

- Suppose we want to find the sixth root of a real number. We can just concatenate the cube root and square root functions. In other words we can define sixthroot x = sqrt (cbrt x).
- How do we define a function that finds all six sixth roots of a complex number using sqrt' and cbrt'?

- We face the similar problem like in Debuggable Functions. The return type (list) is not compatible with the input type (complex).
- We declare higher-order function bind with the following type:

bind :: (Complex Double -> [Complex Double]) -> ([Complex Double] -> [Complex Double])

```
bind :: (Complex Double -> [Complex Double])
-> ([Complex Double] -> [Complex Double])
bind f x = concat (map f x)
unit x = [x]
```

```
f * g = bind f . g
lift f = unit . f
```

#### Random Numbers

#### random :: StdGen -> (a,StdGen)

- To generate a random number you need a seed, and after you've generated the number you need to update the seed to a new value.
- A function that is conceptually a randomised function a -> b can be written as a function a -> StdGen -> (b,StdGen) where StdGen is the type of the seed.

#### Random Numbers

```
bind :: (a -> StdGen -> (b,StdGen)) ->
    (StdGen ->(a,StdGen)) -> (StdGen -> (b,StdGen))
bind f x seed = let (x',seed') = x seed in f x' seed'
unit :: a -> (StdGen -> (a,StdGen))
unit x g = (x,g)
```

## Random Numbers Complete Example in Haskell

import Random

```
bind :: (a -> StdGen -> (b,StdGen)) -> (StdGen -> (a,StdGen)) -> (StdGen -> (b,StdGen))
```

bind f x seed = let (x', seed') = x seed in f x' seed'

```
unit x g = (x,g)
```

lift f = unit . f

## Random Numbers Complete Example in Haskell

```
addDigit n g =
```

```
let (a,g') = random g in (n + a \mod 10,g')
```

```
shift = lift (*10)
```

test :: Integer -> StdGen -> (Integer,StdGen)

test = bind addDigit . bind shift . addDigit

g = mkStdGen 123

main = print \$ test 0 g

## Summary

```
type Debuggable a = (a,String)
```

type Multivalued a = [a]

type Randomised a = StdGen -> (a,StdGen)

m ∈ {Debuggable, Multivalued, Randomised}

- We're given a function a -> m b but we need to somehow apply this function to an object of type m a instead of one of type a.
- In each case we do so by defining a function called bind of type
   (a -> m b) -> (m a -> m b) and introducing a kind of identity
   function unit :: a -> m a.

## Summary

 The triple of objects (m,unit,bind) is the monad, and to be a monad they must satisfy the Monad laws such as unit \* f = f \* unit = f, ...

#### Monads in Haskell

- Haskell is a lazy evaluated pure-functional language.
- Monads are there used for I/O operations, State and other standard side-effects.
- In Haskell we write bind as infix operator >>=. So bind f x is written as x >>= f.
- unit function is called return.
- From previous examples Debuggable is the Writer monad, Multivalued is the List monad and Randomised is the State monad.

#### Monads in Haskell

```
return 7 >>= (\x -> Writer (x+1,"inc."))

>>= (\x -> Writer (2*x,"double."))

>>= (\x -> Writer (x-1,"dec."))
```

## Haskell Syntax

do x <- y

more code

$$y >>= (\x -> do$$

more code).

## Haskell Syntax

```
do
   let x = 7
   y \leftarrow Writer(x+1,"inc\n")
   z \leftarrow Writer (2*y, "double\n")
   Writer (z-1, "dec\n")
```

#### References

- http://www.haskell.org/haskellwiki/Monad
- http://blog.sigfpe.com/2006/08/you-could-haveinvented-monads-and.html?m=1