

Practical Malware Analysis

CH 5: Advanced Dynamic Analysis

Previous session wrap-up

Dynamic analysis: Running malware in safe environment to study the behavior of the malware and how it works.

Real machines:

- No internet connection, some parts of the malware won't work.
- Difficult to remove malware → resetting machine is required.
- + Some malware detects VMs, so it'll run its native behaviour.

VMs:

- Sometimes it's detected by malware, thus it won't 100% work.
- + Most common method, easy to setup and reset.

Previous session wrap-up

why dynamic analysis?

You don't have to deal with:

- obfuscation
- packing
- advanced techniques.

Important tools:

- ProcMon
- Process explorer
- Process hacker
- RegShot
- INetSim
- wireShark

Dynamic analysis activities:

- Monitoring processes
- Tracking DLLs
- Detecting malicious documents.
- Registry before & after.
- Simulating a network and inspecting traffic

Debugging

Debuggers vs Disassemblers

- A disassembler like IDA Pro shows the state of the program just before execution begins
- Debuggers show
 - Every memory location
 - Registers
 - Argument to every function
 - At any point during processing

And it lets you change those values in real time



Two debuggers

ollydbg:

- Most popular
- user-mode debugging only.
- IDA Pro has a built-in debugger, but harder than ollyDBG

windbg:

- Supports kernel-mode debugging.

Two debuggers

ollydbg:

- Most popular
- user-mode debugging only.
- IDA Pro has a built-in debugger, but harder than ollyDBG

windbg:

- Supports kernel-mode debugging.

Source-level VS Assembly-level debuggers

Source-level:

- Built into development platform.
- Breakpoints are based on line of code.
- Can step through the program one line at a time.

Assembly-level:

- Operates on assembly code (instructions)
- Breakpoints are based on instructions instead of lines of code.
- Mostly used by malware analysts

User-mode vs Kernel-mode

User-mode

Debugger always runs on the same system as the code is analyzed on.

Separated from other executables by the OS

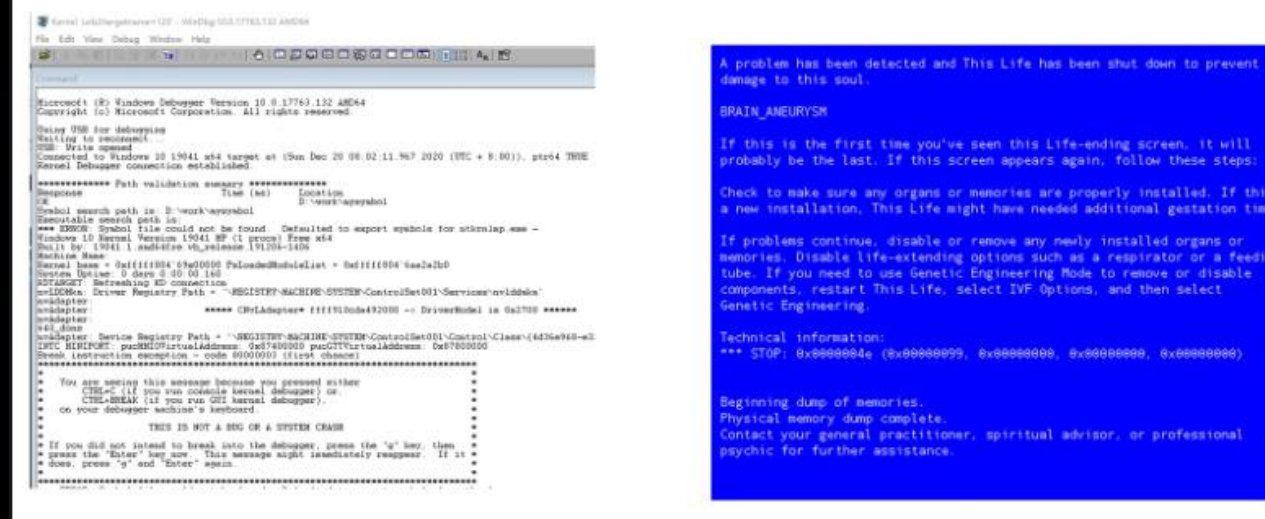
Debugs a single executable.

Kernel mode debugging (old way)

The two computers are connected, one runs debugger, the other runs the malware.

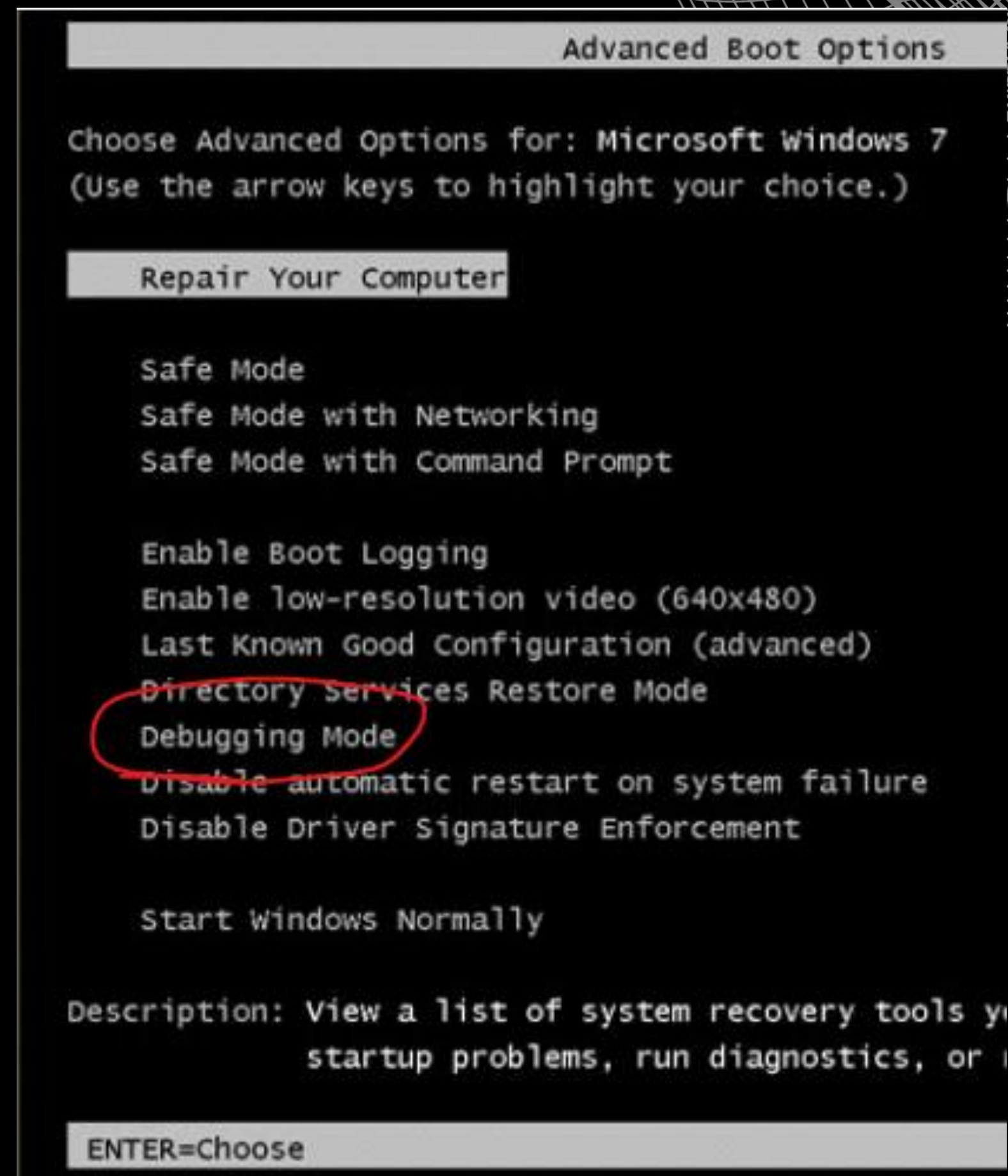
- Requires two computers
- Breakpoints stop the whole system, (not just the program).

windbg kernel-mode debugging



Kernel-mode debugging (new way)

- LiveKD tool No other computer is required. This tool has some limitations.
- On windows 7, you can enter debugging mode (kernel) when you press f8 during startup.



How to debug using a debugger ☺

Two ways

- Start malware with the debugger:
The debugger will stop the program after its entry point.
- Attach a debugger to a program that is already running:
Stop all the threads of the program and load it into the debugger. This is useful when you're debugging a process that is affected by malware.

Example

This code encrypts the string with XOR.

Single stepping is to stop after each Instruction.

```
mov     edi, DWORD_00406904
mov     ecx, 0x0d
LOC_040106B2
xor     [edi], 0x9C
inc     edi
loopw   LOC_040106B2
...
DWORD:00406904:  F8FDF3D01
```

```
D0F3FDF8 D0F5FEEE FDEEE5DD 9C (.....)
4CF3FDF8 D0F5FEEE FDEEE5DD 9C (L.....)
4C6FFDF8 D0F5FEEE FDEEE5DD 9C (Lo.....)
4C6F61F8 D0F5FEEE FDEEE5DD 9C (Loa.....)
. . . SNIP . . .
4C6F6164 4C696272 61727941 00 (LoadLibraryA.)
```

Step-over VS Step-into

- Stepping over is to finish the subroutine (function) as a single instruction, no pauses until the subroutine is finished.
- Stepping into is to stop at each instruction of the subroutine.

Breakpoint usage

We can't tell what this program does. Set a breakpoint at the call and see what's in eax.

This program encrypts data before sending it, pause the program before it encrypts the data, and see what data it's encrypting.

```
00401008  mov     ecx, [ebp+arg_0]
0040100B  mov     eax, [edx]
0040100D  call    eax
```

```
004010D0  sub     esp, 0CCh
004010D6  mov     eax, dword_403000
004010DB  xor     eax, esp
004010DD  mov     [esp+0CCh+var_4], eax
004010E4  lea     eax, [esp+0CCh+buf]
004010E7  call    GetData
004010EC  lea     eax, [esp+0CCh+buf]
004010EF  call    EncryptData
004010F4  mov     ecx, s
004010FA  push    0                ; flags
004010FC  push    0C8h            ; len
00401101  lea     eax, [esp+0D4h+buf]
00401105  push    eax              ; buf
00401106  push    ecx              ; s
00401107  call    ds:Send
```

Types of Breakpoints

Software execution

- default
- Overwrites the first byte of instruction with 0xcc (int 3 opcode) to raise exception to give control to the debugger.

Anti-debugging

- Sometimes anti-debugging malware checks this opcode (0xCC) in the memory and modifies it, and the breakpoint fails.
- Sometimes other code reads memory containing the breakpoint, and it will read 0xCC instead of original byte.
- Code that verifies integrity will notice the discrepancy

Hardware execution breakpoints

- Instead of opcode 0xCC, the DR (debug register) is used.
- This causes a breakpoint prior to any mov instruction that would change the contents of debug register.
- Does not detect other instructions, however.

Conditional breakpoints

- Same as software breakpoints but breaks when a condition is true.
- takes much longer than ordinary instructions.
- Sometimes so much that it never finishes.

For example: breakpoint on GetProcAddress function.

Exceptions

Used by debuggers to gain control of a running program.

- Generated by breakpoints.
- Also caused by invalid memory access.
- Other exceptions, such as division by 0, etc.

Common Exceptions

Examples of common exceptions:

- int 3 (opcode 0xCC, for software breakpoint.)
- Single-stepping: execute an instruction, generate exception.
- Memory-access violation
- violation of other privilege rules:
For example, executing a kernel-mode instruction while the code is in user-mode.

Other debugging activities (skipping a function)

while debugging, we change the current memory of the debugged code, by changing value of the instruction pointer, or the code itself.

You can skip a function by setting a breakpoint at the call, then changing the instruction pointer to the instruction after that function. (which isn't always safe for the program execution)

Other debugging activities (Testing a function)

You can run a function directly, without waiting for the main code to use it.

- You have to set the parameters
- This will destroy the program's attack.
- The program won't run properly when the function completes.

HTU.

OllyDbg



Loading malware

- You can load EXEs or DLLs into OllyDbg
- If the malware is already running, you can attach OllyDbg to it.

In case of opening the malware with OllyDbg, it'll stop at the entry point (WinMain, the main function if found), otherwise, it'll break at the entry point defined in PE header.

In the case of attaching OllyDbg to a process, it'll stop all the threads of the process.

OllyDbg Interface

OllyDbg - Lab09-01.exe - [CPU - main thread, module Lab09-01]

File View Debug Plugins Options Window Help

Registers (FPU)

Registers

Disassembler

Memory Dump

Stack

EBP=0019FF80

Lab09-01.<ModuleEntryPoint>

Address	Hex dump	ASCII
0040C000	00 00 00 00 00 00 00 00
0040C008	00 00 00 00 EF 42 40 00nB0.
0040C010	6C 5A 40 00 00 00 00 00	lZ0.....
0040C018	00 00 00 00 94 43 40 000C0.
0040C020	00 00 00 00 00 00 00 00
0040C028	00 00 00 00 00 00 00 00
0040C030	43 6F 6E 66 69 67 75 72	Configur
0040C038	61 74 69 6F 6E 00 00 00	ation...
0040C040	53 4F 46 54 57 41 52 45	SOFTWARE
0040C048	5C 4D 69 63 72 6F 73 6F	\Microso
0040C050	66 74 20 5C 58 50 53 00	ft \XPS.
0040C058	5C 68 65 72 6E 65 6C 33	\kernel3
0040C060	32 2E 64 6C 6C 00 00 00	2.dll...
0040C068	0D 0A 0D 0A 00 00 00 00
0040C070	20 48 54 54 50 2F 31 2E	HTTP/1.
0040C078	30 00 0A 00 0A 00 00 00	0.....
0040C080	47 45 54 20 00 00 00 00	GET
0040C088	27 60 27 60 27 60 27 60

0019FF74 76910179 RETURN to KERNEL32.76910179

0019FF78 002C5000

0019FF7C 76910160 KERNEL32.BaseThreadInitThunk

0019FF80 0019FFDC

0019FF84 77DB662D RETURN to ntdll.77DB662D

0019FF88 002C5000

0019FF8C F992FD2B

0019FF90 00000000

0019FF94 00000000

0019FF98 002C5000

0019FF9C 00000000

0019FFA0 00000000

0019FFA4 00000000

0019FFA8 00000000

0019FFAC 00000000

0019FFB0 00000000

0019FFB4 00000000

0019FFB8 00000000

0019FFBC 00000000



Memory dump

You can see which DLLs are identified, double click any row to see its memory.

You can also show the instructions in disassembler

OllyDbg - Lab09-01.exe - [Memory map]

File View Debug Plugins Options Window Help

LEMTW H C / K B R ... S

Address	Size	Owner	Section	Contains	Type	Access	Initial	Ma
755DE000	00005000	IMM32	755C0000	.rsrc	resources	Inag 01001002 R	RWE	
755E3000	00002000	IMM32	755C0000	.reloc	relocations	Inag 01001002 R	RWE	
755F0000	00001000	conbase	755F0000	(itself)	PE header	Inag 01001002 R	RWE	
755F1000	00230000	conbase	755F0000	.text	code,export	Inag 01001002 R	RWE	
75821000	00004000	conbase	755F0000	.proxy	code	Inag 01001002 R	RWE	
75825000	00004000	conbase	755F0000	.data	data	Inag 01001002 R	RWE	
75829000	00005000	conbase	755F0000	.idata	imports	Inag 01001002 R	RWE	
7582E000	00001000	conbase	755F0000	.didat		Inag 01001002 R	RWE	
7582F000	00014000	conbase	755F0000	.rsrc	resources	Inag 01001002 R	RWE	
75843000	00025000	conbase	755F0000	.reloc	relocations	Inag 01001002 R	RWE	
75920000	00001000	cryptsp	75920000	(itself)	PE header	Inag 01001002 R	RWE	
75921000	0000C000	cryptsp	75920000	.text	code,export	Inag 01001002 R	RWE	
7592D000	00001000	cryptsp	75920000	.data	data	Inag 01001002 R	RWE	
7592E000	00001000	cryptsp	75920000	.idata	imports	Inag 01001002 R	RWE	
7592F000	00001000	cryptsp	75920000	.didat		Inag 01001002 R	RWE	
75930000	00001000	cryptsp	75920000	.rsrc	resources	Inag 01001002 R	RWE	
75931000	00001000	cryptsp	75920000	.reloc	relocations	Inag 01001002 R	RWE	
75940000	00001000	KERNELBA	75940000	(itself)	PE header	Inag 01001002 R	RWE	
75941000	001C3000	KERNELBA	75940000	.text	code,export	Inag 01001002 R	RWE	
75B04000	00004000	KERNELBA	75940000	.data	data	Inag 01001002 R	RWE	
75B08000	00006000	KERNELBA	75940000	.idata	imports	Inag 01001002 R	RWE	
75B0E000	00001000	KERNELBA	75940000	.didat		Inag 01001002 R	RWE	
75B0F000	00001000	KERNELBA	75940000	.rsrc	resources	Inag 01001002 R	RWE	
75B10000	0002A000	KERNELBA	75940000	.reloc	relocations	Inag 01001002 R	RWE	
75B40000	00001000	USER32	75B40000	(itself)	PE header	Inag 01001002 R	RWE	
75B41000	000A0000	USER32	75B40000	.text	code,export	Inag 01001002 R	RWE	
75B4E000	00002000	USER32	75B40000	.data	data	Inag 01001002 R	RWE	
75B4F000	00003000	USER32	75B40000	.idata	imports	Inag 01001002 R	RWE	
75B4F000	000031000	USER32	75B40000	.didat		Inag 01001002 R	RWE	
75B4F000	0000E2000	USER32	75B40000	.rsrc	resources	Inag 01001002 R	RWE	
75CD2000	00007000	USER32	75B40000	.reloc	relocations	Inag 01001002 R	RWE	
75DE0000	00001000	windows_	75DE0000	(itself)	PE header	Inag 01001002 R	RWE	
75DE1000	00588000	windows_	75DE0000	.text	code,export	Inag 01001002 R	RWE	
76369000	00007000	windows_	75DE0000	.data	data	Inag 01001002 R	RWE	
76370000	00007000	windows_	75DE0000	.idata	imports	Inag 01001002 R	RWE	
76377000	00001000	windows_	75DE0000	.didat		Inag 01001002 R	RWE	
76378000	00004000	windows_	75DE0000	.rsrc	resources	Inag 01001002 R	RWE	
7637C000	00060000	windows_	75DE0000	.reloc	relocations	Inag 01001002 R	RWE	
765E0000	00001000	shlwapi	765E0000	(itself)	PE header	Inag 01001002 R	RWE	
765E1000	00039000	shlwapi	765E0000	.text	code,export	Inag 01001002 R	RWE	
7661A000	00001000	shlwapi	765E0000	.data	data	Inag 01001002 R	RWE	
7661B000	00004000	shlwapi	765E0000	.idata	imports	Inag 01001002 R	RWE	
7661F000	00001000	shlwapi	765E0000	.didat		Inag 01001002 R	RWE	
76620000	00001000	shlwapi	765E0000	.rsrc	resources	Inag 01001002 R	RWE	
76621000	00003000	shlwapi	765E0000	.reloc	relocations	Inag 01001002 R	RWE	
76630000	00001000	powrprof	76630000	(itself)	PE header	Inag 01001002 R	RWE	
76631000	00024000	powrprof	76630000	.text	code,export	Inag 01001002 R	RWE	
76655000	00001000	powrprof	76630000	.data	data	Inag 01001002 R	RWE	
76656000	00002000	powrprof	76630000	.idata	imports	Inag 01001002 R	RWE	
76658000	00001000	powrprof	76630000	.didat		Inag 01001002 R	RWE	
76659000	00029000	powrprof	76630000	.rsrc	resources	Inag 01001002 R	RWE	
76682000	00002000	powrprof	76630000	.reloc	relocations	Inag 01001002 R	RWE	
76830000	00001000	nsuort	76830000	(itself)	PE header	Inag 01001002 R	RWE	
76831000	000B2000	nsuort	76830000	.text	code,export	Inag 01001002 R	RWE	
768E3000	00006000	nsuort	76830000	.data	data	Inag 01001002 R	RWE	
768E9000	00002000	nsuort	76830000	.idata	imports	Inag 01001002 R	RWE	
768EB000	00001000	nsuort	76830000	.rsrc	resources	Inag 01001002 R	RWE	
768EC000	00004000	nsuort	76830000	.reloc	relocations	Inag 01001002 R	RWE	
768F0000	00001000	KERNEL32	768F0000	(itself)	PE header	Inag 01001002 R	RWE	
76900000	00064000	KERNEL32	768F0000	.text	code	Inag 01001020 R E	RWE	
769A0000	00001000	KERNEL32	768F0000	.rdata	imports,exp	Inag 01001002 R	RWE	
769B0000	00001000	KERNEL32	768F0000	.data	data	Inag 01001004 RW	RWE	
769C0000	00001000	KERNEL32	768F0000	.rsrc	resources	Inag 01001002 R	RWE	
769C0000	00005000	KERNEL32	768F0000	.reloc	relocations	Inag 01001002 R	RWE	
769C0000	00001000	nsuort	768F0000	(itself)	PE header	Inag 01001002 R	RWE	

Rebasing

- Rebasing is when a module isn't loaded at its preferred base address.
- PE's have a preferred base address. Usually, the files are loaded at that address, which often equals 0x00400000.
- DLLs are commonly relocated, because a single EXE can import many DLLs. The list of DLLs locations are in the .reloc section of the PE header.

Absolute vs Relative Address

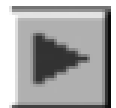
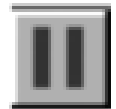

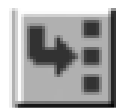
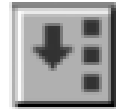
```
00401203      mov eax, [ebp+var_8]
00401206      cmp [ebp+var_4], 0
0040120a      jnz loc_0040120
0040120c      ❶mov eax, dword_40CF60
```

The first 3 instructions will run fine, since they use a relative address. → [ebp+var_8]

The last instruction won't always work, as it uses an absolute address to access the exact memory location.



Executing malware

Function	Menu	Hotkey	Button
Run/Play	Debug ▶ Run	F9	
Pause	Debug ▶ Pause	F12	
Run to selection	Breakpoint ▶ Run to Selection	F4	
Run until return	Debug ▶ Execute till Return	CTRL-F9	
Run until user code	Debug ▶ Execute till User Code	ALT-F9	
Single-step/step-into	Debug ▶ Step Into	F7	
Step-over	Debug ▶ Step Over	F8	



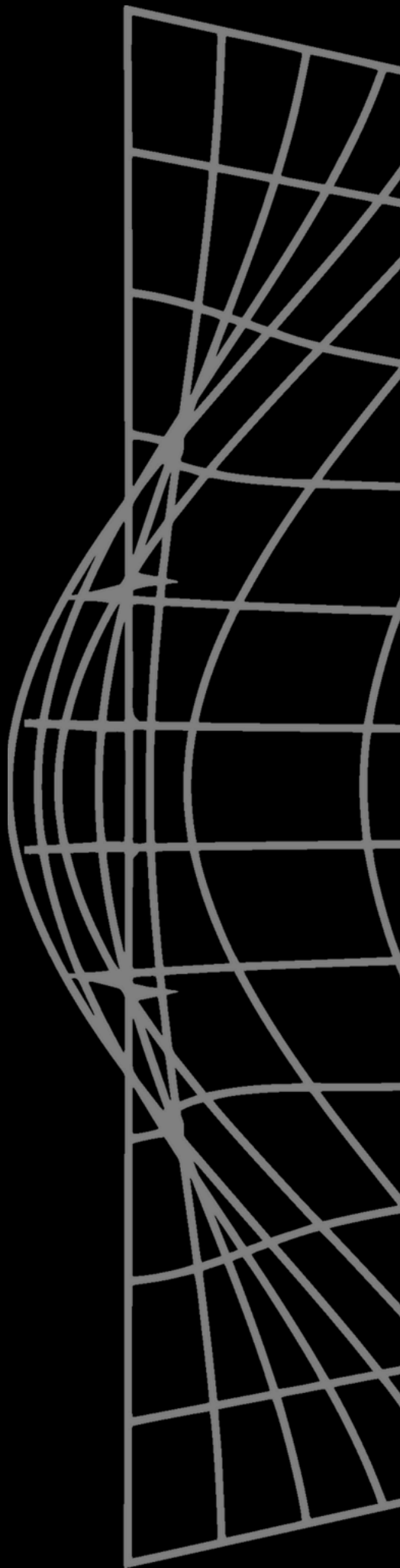
Notes on executing code

- Setting breakpoints is much better than just Run and Pause.
- You can execute till return
- Execute till User Code → will execute till it hits the compiled malware code
- Single-stepping: step-into or step-over
- Sometimes step-over will miss important functions, so the code will run forever.



Breakpoints

We've already explained breakpoints and their types, you can refer to [this pdf](#) to have a detailed explanation on how to set each type.



Loading DLLs

- DLLs aren't executed directly or alone.
- o1lyDbg uses a dummy loadDll.exe to load them.
- breaks at the DLL entry point DLLMain once DLL is loaded.



Tracing

Tracing is recording detailed execution information.

- Standard back trace
- Call stack trace
- Run trace



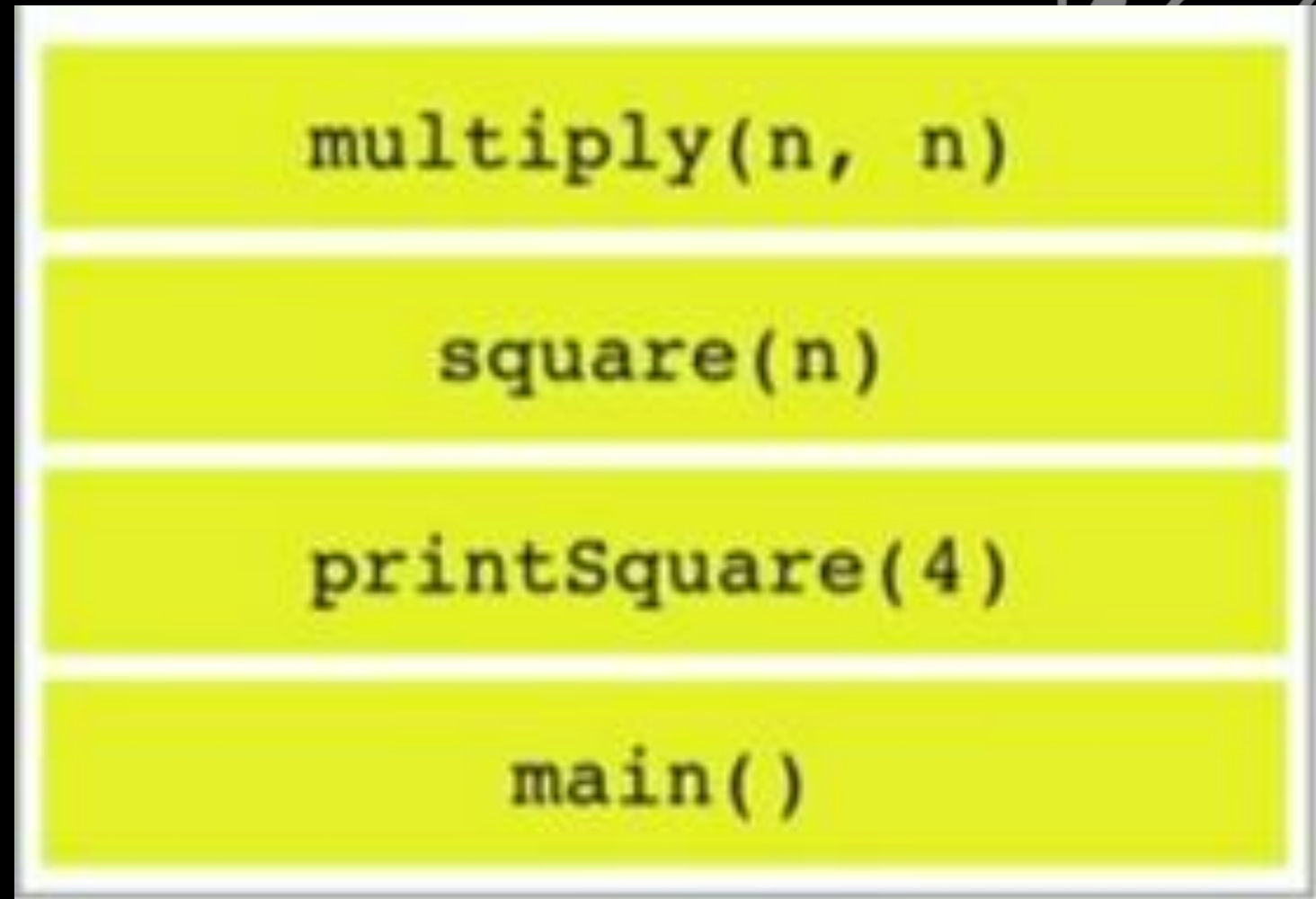
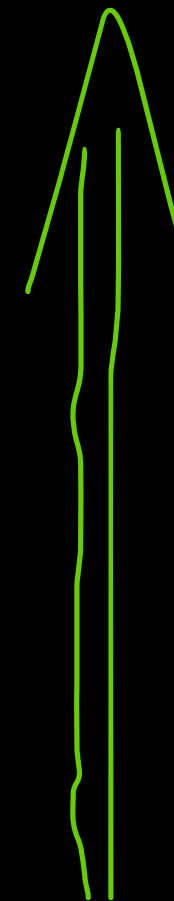
Standard back trace

- You move through the disassembler with the Step Into and Step Over buttons
- OllyDbg is recording your movement
- Use minus key on keyboard to see previous instructions, But you won't see previous register values
- If you used Step Over, you cannot go back and decide to step into



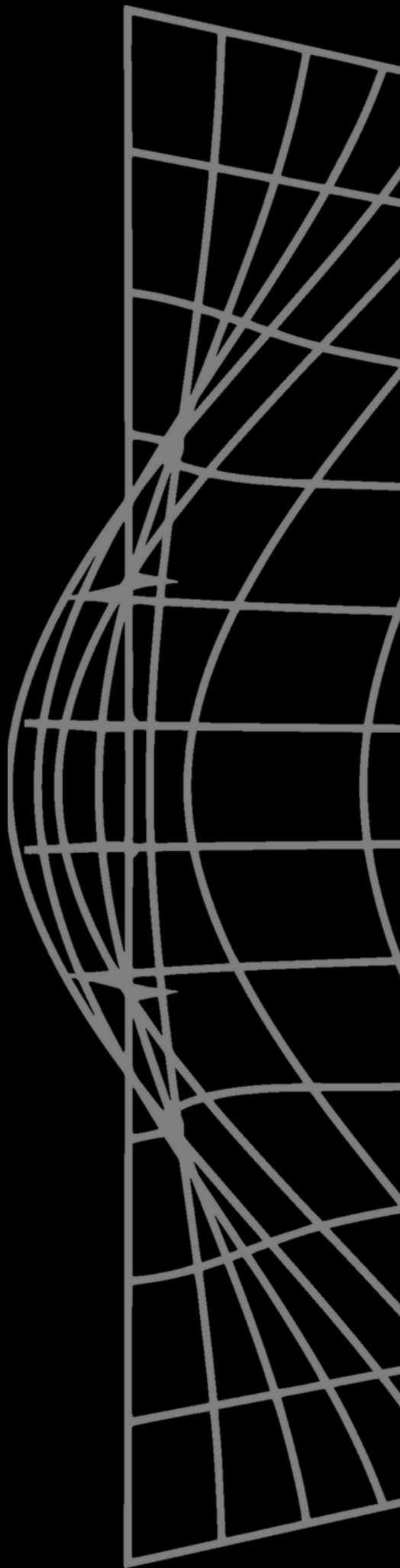
Call Stack Trace

- Views execution path to a given function.
- Displays the sequence of calls to reach your current location, which is the stack of calls (hence the name).



Run Trace

- Code runs, OllyDbg saves every executed instruction and all changes to Registers and flags.
- You can go back and forth, if you go back, you could see the previous values of registers.



Handling exceptions

- when an exception occurs, OllyDbg will stop, and you have to :
 - Step into exception
 - Step over exception (most common and practical)
 - Run exception handler



Patching

- A good way to skip some instructions in the binary.
- You edit the instructions, change them to 0x90, which is NOP instruction (No-Operation)
- This will skip instructions or force a branch



Drivers

- Drivers to hardware are like DLLs to EXEs. They make hardware communicate with the kernel indirectly.
- Example: USB Flash Drive. When it's plugged in, a device object is created.
- Application → OS → Device Driver → hardware.
- Multiple devices can use the same Driver (just like multiple EXEs can use the same DLL.)



Loading drivers

- Drivers are loaded in the kernel.
- DriverEntry procedure is called (like DLLMain)
- Callback objects and functions can be used by OS (like DLL Exports)
- Difference is that DLLs Exports are for user-mode programs, while Driver Callback functions will be a part of kernel-mode programs which interact with the kernel.



Kernel and user-mode calls

You can see how the program is using a malicious driver, even if it doesn't include malicious DLL.

Ntoskrnl.exe has core OS functions, so malware will often import functions from one it so it can manipulate the kernel.

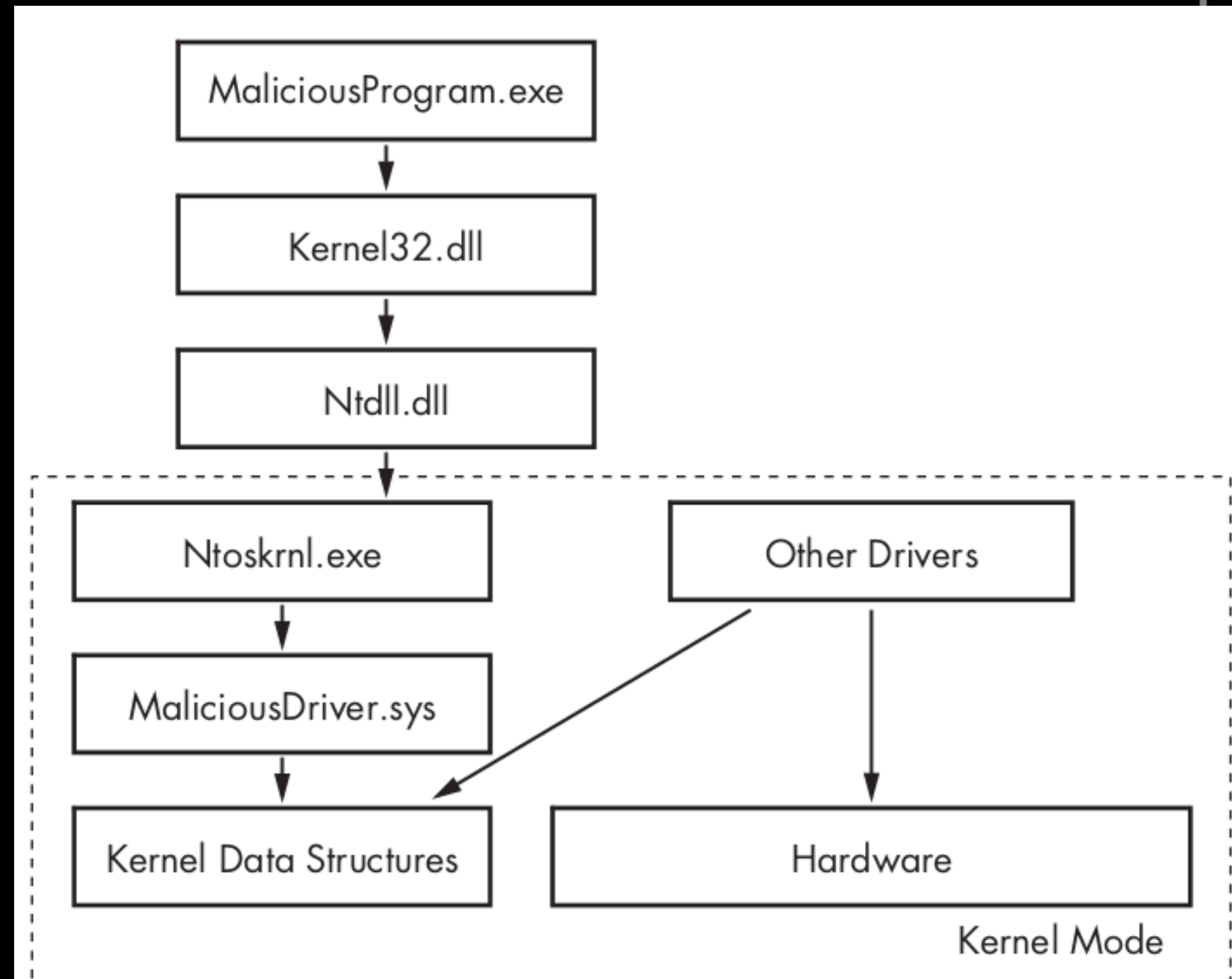


Figure 10-1: How user-mode calls are handled by the kernel

Practical Malware Analysis

CH 6: Malware behavior

Downloaders and launchers

They form the first stages of malware.

- Downloader often uses `URLDownloadToFileA`, followed by `WinExec`.
- Launcher (loader) will install the malware and installs malware in unexpected places like .rsrc section of PE.



Backdoor

- A backdoor will provide a persistent way to remotely access a victim machine.
- Most common.
- Communication is usually through http/s port or other common ports to prevent detection.
- Often manipulates registry, creates directories, searches files, etc.



Reverse shell

- Infected machine calls out to attacker, asking for commands to execute.
- On windows, this is done by calling `CreateProcess` and manipulating `STARTUPINFO` structure. Then creating a socket to a remote machine and bind the socket to `stdin/stdout/stderr` for `cmd.exe`.



RATs

(Remote Administration Tools)

- Carried on remotely managing a group of compromised computers.
- Used in targeted attacks. The structure is client-server.

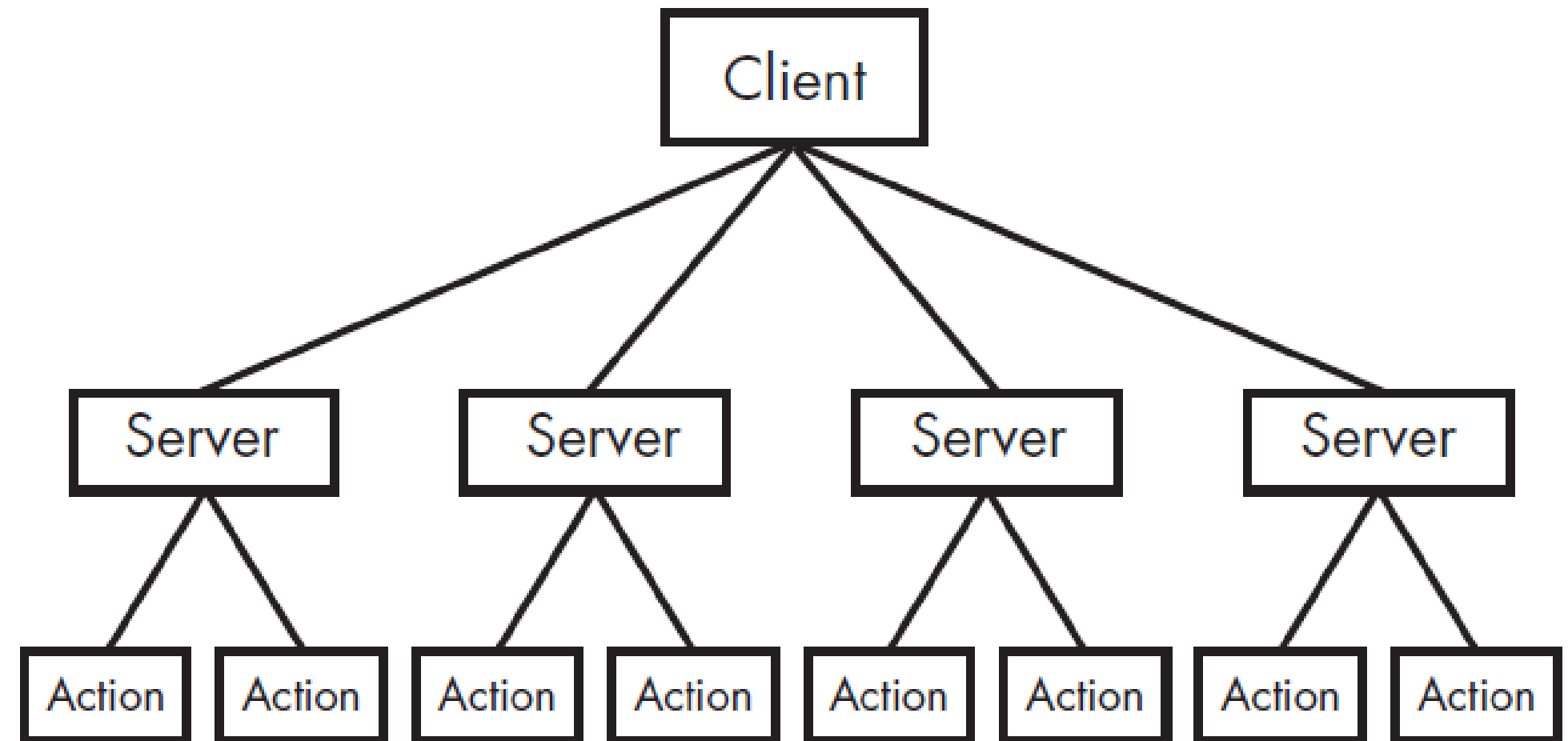


Figure 11-1: RAT network structure

Botnets

- A botnet is like a RAT, but it's for mass attacks.
- It has much bigger number of compromised machines than RATs. Compromised machines are called bots/zombies.
- All bots are controlled at once, while in RAT the victims are controlled one by one.



Credential stealers

- Dump stored data, such as password hashes, cached cookies, etc.
- Keyloggers.



Keyloggers (kernel-based)

- Built as keyboard driver.
- Bypasses user-space protection.
- Difficult to detect using user-space programs.



Keyloggers (user-mode)

- Uses windows API along with hooking.
- Uses SetWindowsHookEx function to notify the malware each time a key is pressed.
- When running strings on a keylogger, you might see terms like Num Lock, Up, Down.

```
[Up]  
[Num Lock]  
[Down]  
[Right]  
[UP]  
[Left]  
[PageDown]
```

Downloaders and launchers

They form the first stages of malware.

- Downloader often uses `URLDownloadToFileA`, followed by `WinExec`.
- Launcher (loader) will install the malware and installs malware in unexpected places like .rsrc section of PE.



Persistence

- Making your malware stay in the victim machine for a long time.
- Done through modifying registry, Trojanizing binaries, or DLL Load-Order hijacking.



Registry modification

- A very common registry that malware often modifies is `HKEY_LOCAL_MACHINE\SOFTWARE\Microsoft\Windows\CurrentVersion\Run`
- This registry is set to run on each boot.
- Registry actions can be detected using ProcMon when running the malware.



Trojanizing binaries

- Malware patches bytes of a system binary so it'll execute the malware next time that binary is loaded
- DLLs are common targets
- Modifies entry function so it jumps to the malware inserted in empty portion of the binary. Then it continues executing the binary normally.



Trojanized binaries

Original code	Trojanized code
<pre>DllEntryPoint(HINSTANCE hinstDLL, DWORD fdwReason, LPVOID lpReserved) mov edi, edi push ebp mov ebp, esp push ebx mov ebx, [ebp+8] push esi mov esi, [ebp+0Ch]</pre>	<pre>DllEntryPoint(HINSTANCE hinstDLL, DWORD fdwReason, LPVOID lpReserved) jmp DllEntryPoint_0</pre>

DLL Load-Order Hijacking

- The registry key KnownDLLs has a list of specific known DLL locations.
- Overrides the search order for listed DLLs, makes them load faster. And prevents load-order hijacking.
- Load-order hijacking can be used on binaries in directories other than system32 that load DLLs in system32 and aren't protected by KnownDLLs.



DLL Load-Order Hijacking

For example, explorer.exe. This binary is in \windows. Loads ntshrui.dll from System32, but ntshrui.dd isn't a known DLLs.

Default search is performed, so a malicious ntshrui.dll can be loaded instead.



Privilege escalation

- Metasploit has many privilege escalation exploits.
- privilege escalation can also be done using DLL order-load hijacking.



Covert Malware launching

Malware used to be visible in windows task manager. So malware authors try to blend their malware into normal windows landscape.

As we discussed before, a malware launcher sets itself or another piece of code for future execution.



Malware launching

The malicious piece of code is often compressed.

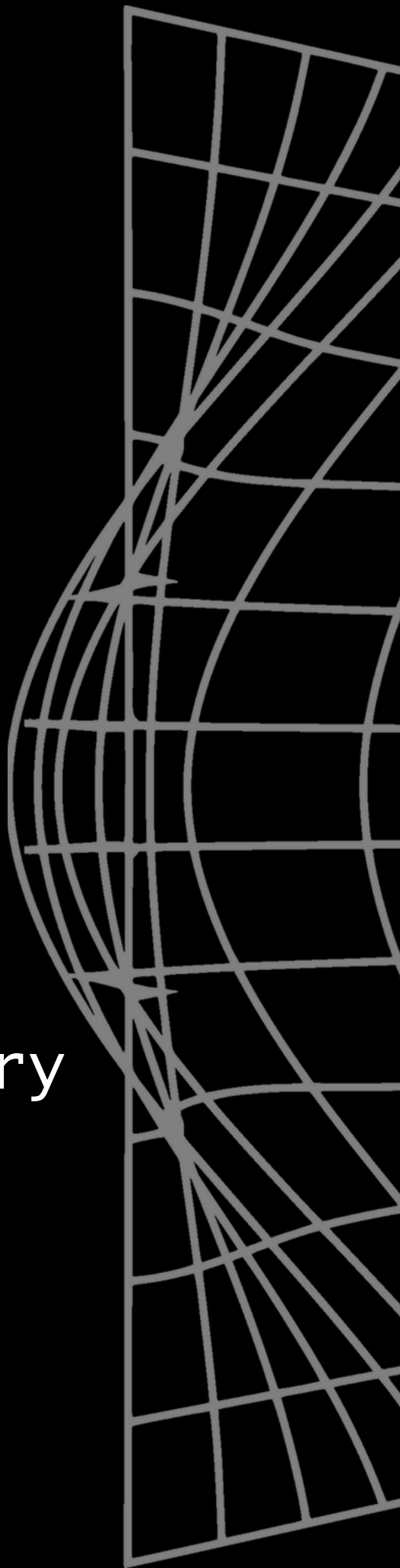
Resource extraction uses API functions such as **FindResource**, **LoadResource**, and **SizeofResource**.

Common techniques are process injection and DLL injection.



Process injection

- Injecting code to an already running process.
- Can be DLL injection or direct injection.
- Common API calls:
 - `VirtualAllocEx` to allocate space in remote process's memory
 - `WriteProcessMemory` to write to it.



Process injection

For example, launcher wants internet access to download code, but process-specific firewall won't let it access the internet.

We can inject that loader to a process (e.g. IExplorer) that already has internet access.



DLL injection

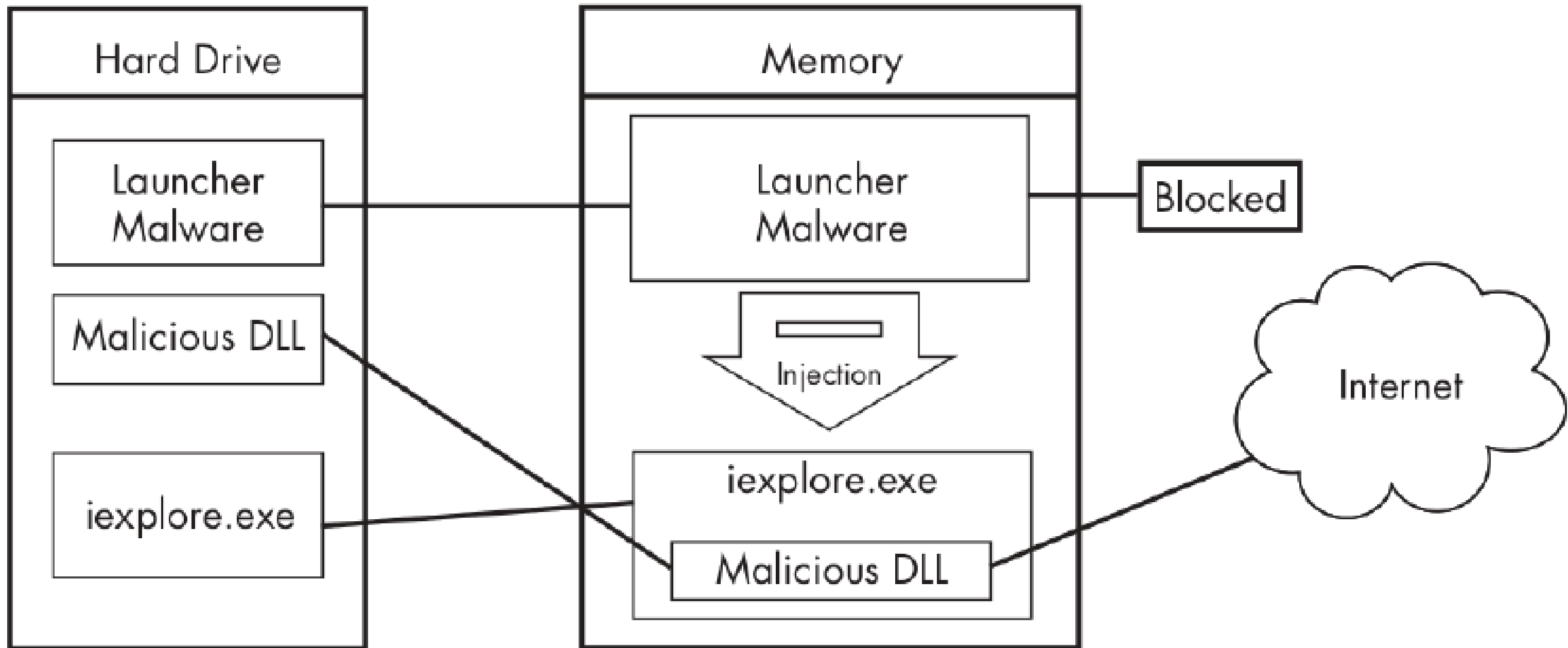


Figure 12-1: DLL injection—the launcher malware cannot access the Internet until it injects into iexplore.exe.

DLL Injection pseudocode

```
hVictimProcess = OpenProcess(PROCESS_ALL_ACCESS, 0, victimProcessID ❶);  
  
pNameInVictimProcess = VirtualAllocEx(hVictimProcess,...,sizeof(maliciousLibraryName),...,...);  
WriteProcessMemory(hVictimProcess,...,maliciousLibraryName, sizeof(maliciousLibraryName),...);  
GetModuleHandle("Kernel32.dll");  
GetProcAddress(...,"LoadLibraryA");  
❷ CreateRemoteThread(hVictimProcess,...,...,LoadLibraryAddress,pNameInVictimProcess,...,...);
```

Listing 12-1: C Pseudocode for DLL injection

Direct injection

- Injects code directly into remote process.
- No DLLs, difficult to write.
- Requires a lot of testing to make sure it doesn't mess with the flow of the native process.



Process replacement

- Instead of writing a process inside a remote process memory, we replace the whole memory of that process.
- Avoids the risk of crashing a process with process injection.
- Commonly replaces `svchost.exe`

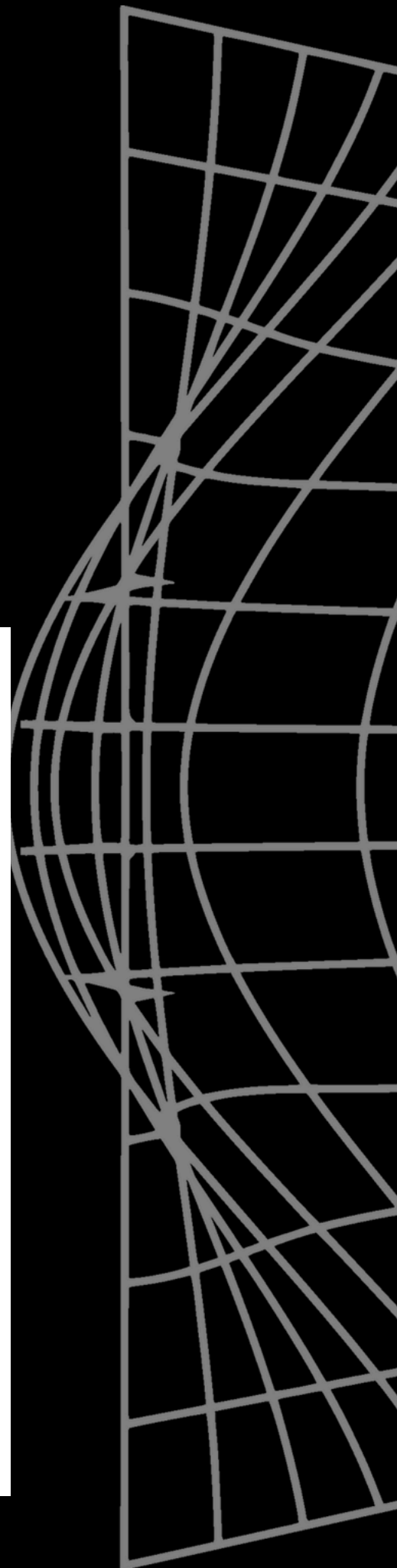


Suspended state

In a suspended state, the process is loaded into memory, but primary thread is suspended (so malware can overwrite it before the code runs)

```
CreateProcess(..., "svchost.exe", ..., CREATE_SUSPEND, ...);
ZwUnmapViewOfSection(...);
VirtualAllocEx(..., ImageBase, SizeOfImage, ...);
WriteProcessMemory(..., headers, ...);
for (i=0; i < NumberOfSections; i++) {
    ❶ WriteProcessMemory(..., section, ...);
}
SetThreadContext();
...
ResumeThread();
```

Listing 12-3: C pseudocode for process replacement



Direct injection

- Injects code directly into remote process.
- No DLLs, difficult to write.
- Requires a lot of testing to make sure it doesn't mess with the flow of the native process.



Hook Injection

- windows hooks intercept messages destined for applications.
- Malicious hooks:
 - Ensure that a malicious code will run whenever a message is intercepted.
 - Ensure that a DLL will be loaded in a victim process's memory.



Hook injection

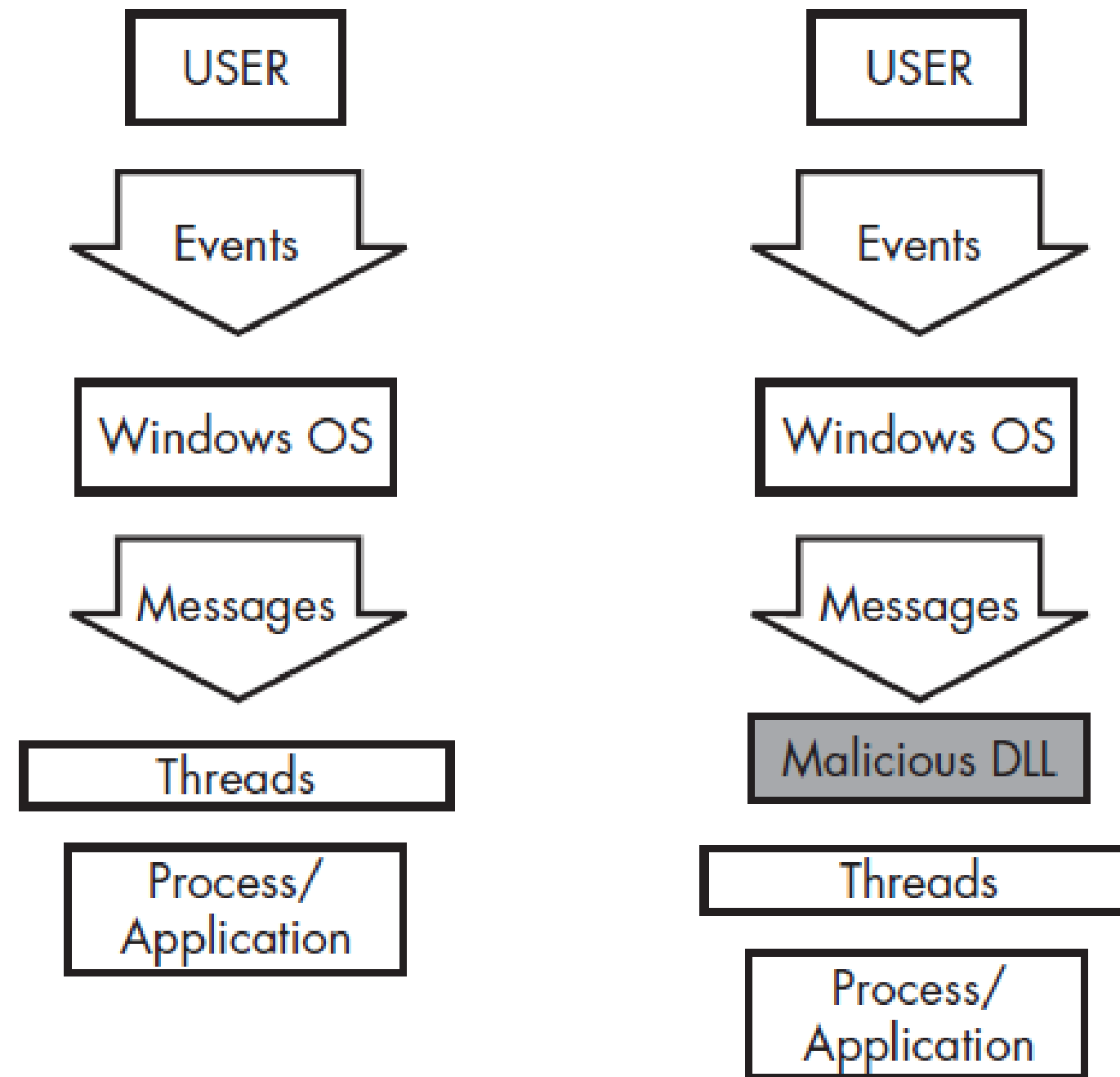


Figure 12-3: Event and message flow in Windows with and without hook injection

Keyloggers using hooks

In keyloggers, `WH_KEYBOARD` and `WH_KEYBOARD_DLL` procedures are used.

We can use `SetWindowsHookEx` for remote windows hooking.

Parameters:

- `idHook` → Type of hook procedure.
- `lpfn` → pointer to hook procedure
- `hMod` → Handle to DLL or local module containing hook procedure.
- `dwThreadId` → Thread to associate hook with.



Encoding

Malware uses encoding for some reasons:

- Hide some information (e.g. C2 domains)
- Save information for staging file
- Store strings needed by malware.
- Hide suspicious strings.

Encryption is also used, and it's much harder to decrypt than decoding encoded strings.



Common encodings

- Caesar cipher: Moving each letter by a certain number.
For example, CMD.EXE → FPG.HAH.
- XOR: Bitwise operation. Uses a key to encrypt data.
For example, malware ^ 0xa → gkf}kxo
NOTE: XOR loops can be found in IDA Pro.
- Base64: Includes [A-Za-z0-9\+\ \/]. Refer to it [here](#)



Encryption

- Much harder to break than encoding. Depends on the encryption algorithm
- For example, AES, SSL, etc.
- Cryptographic libraries required.
- Sometimes the code becomes less portable
- Symmetric encryption requires a way to hide the key, otherwise it's easy to decrypt.



Entropy

Entropy is the randomness of an executable.

High entropy probably means that there's encrypted content.

Can be found using IDA Pro Entropy Plugin.

Note that sometimes multiple methods of encryption can be applied in the same malware.



Entropy

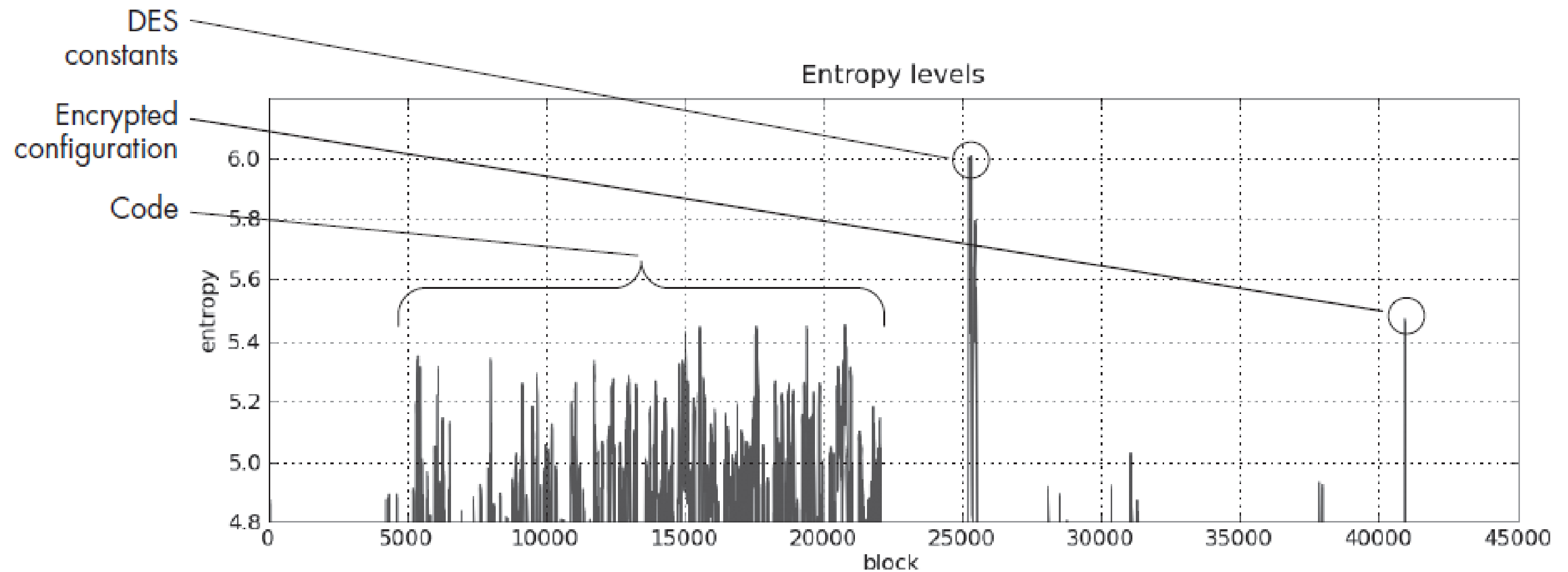
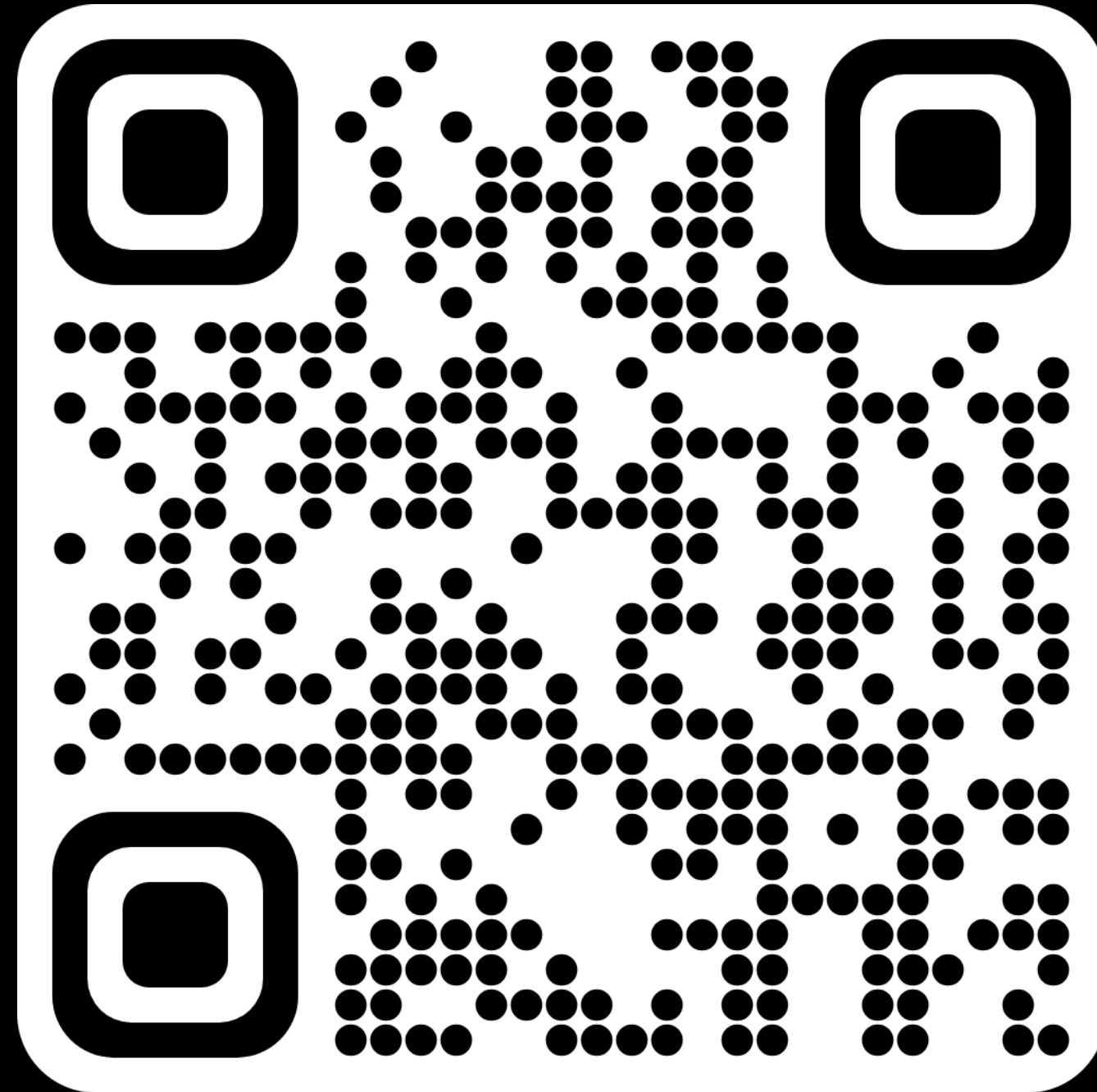


Figure 13-13: Entropy graph for a malicious executable



HTU.

Quiz



THANK YOU

